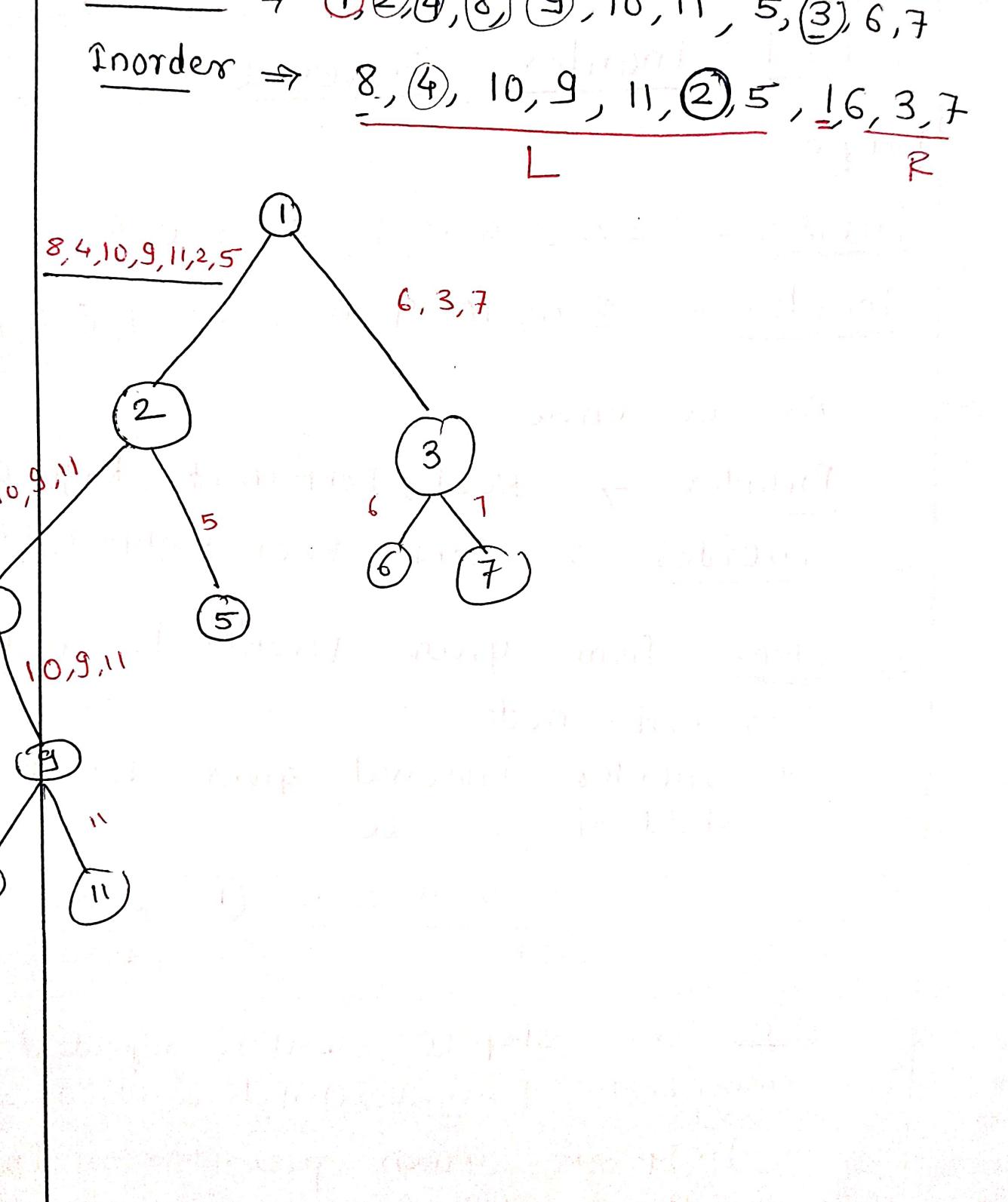


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Topic	Unit No.
<u>Preorder</u> \Rightarrow <u>1, 2, 4, 8, 9, 10, 11, 5, 3, 6, 7</u> <u>Inorder</u> \Rightarrow <u>8, 4, 10, 9, 11, 2, 5, 6, 3, 7</u> 	Root 1 Root 2 Root 3 Root 4 Root 5 L R
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Topic

Unit No.

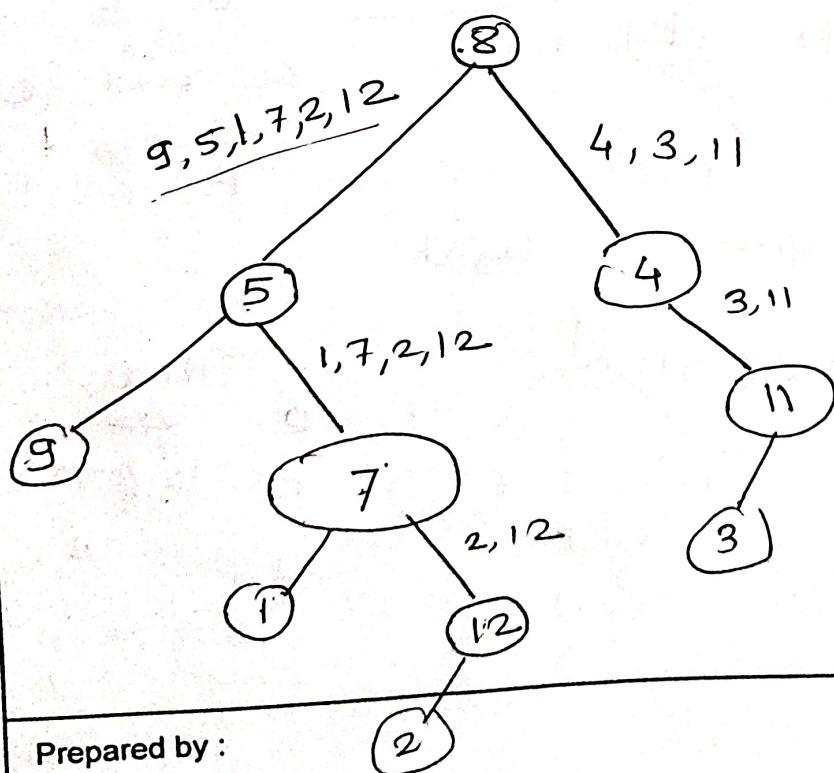
Construct Binary Tree from postorder and Inorder Traversal

Postorder - $(9, 1, 2, 12, 7, \underline{5}, 3, 11, 4) \xrightarrow{\text{R2}} 8$
 (Left, Right, Root)

Inorder - $9, \underline{5} \underline{1, 7, 2, 12, 8, 4, 3, 11} \xrightarrow{\text{R1}} (\text{Left, Root, Right})$

→ Step-1 → As we know in postorder traversal root node is last thus, "8" is root node of new tree

& $9, 5, 1, 7, 2, 12$ 8 $\frac{4, 3, 11}{\text{Root}}$ Left child Right child



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Topic

Unit No.

* Construct Binary Tree from Preorder and Postorder Traversal (Unique Full BT)
 (only with these two it is not possible to construct unique Binary Tree)

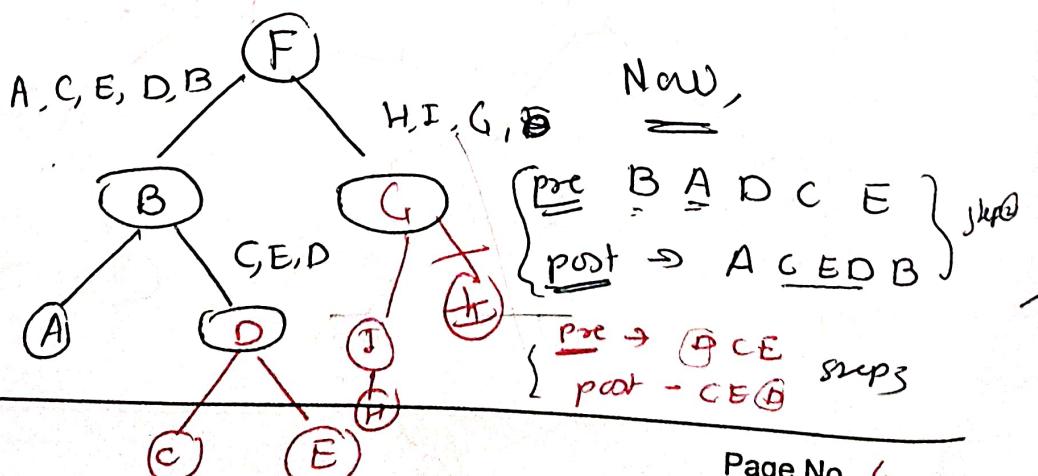
PreOrder -

F B A D C E G I H (Root, L, R)

PostOrder

A C E D B H I G F (L, R, Root)

As in therefore, Preorder First is root Acc. To postorder Last is root i.e. F.
 for given & check its successor (next) of F is preorder is B as its Left child so all elements in postorder of B is left of after B Right.



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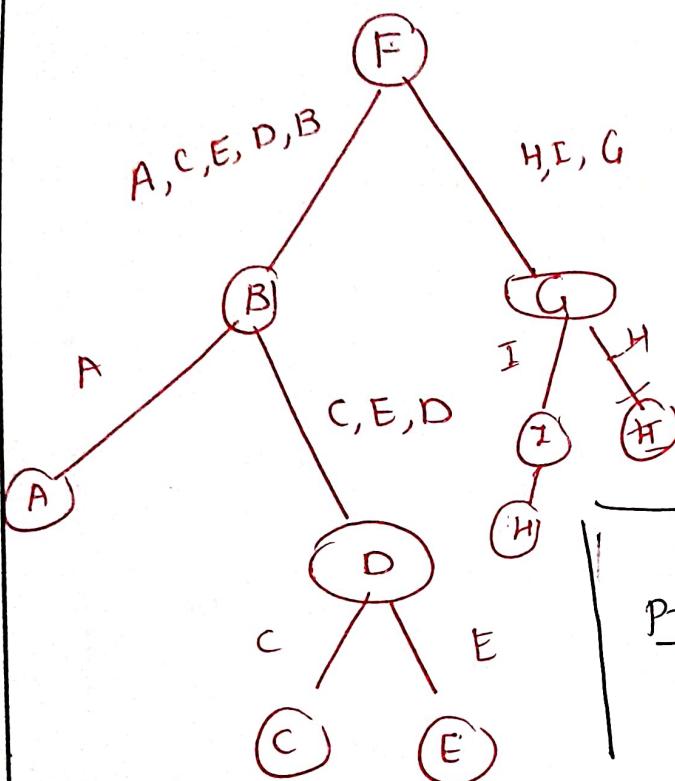
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Prof. G. I. H.
 P.D.T. - H. S. G.

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Topic

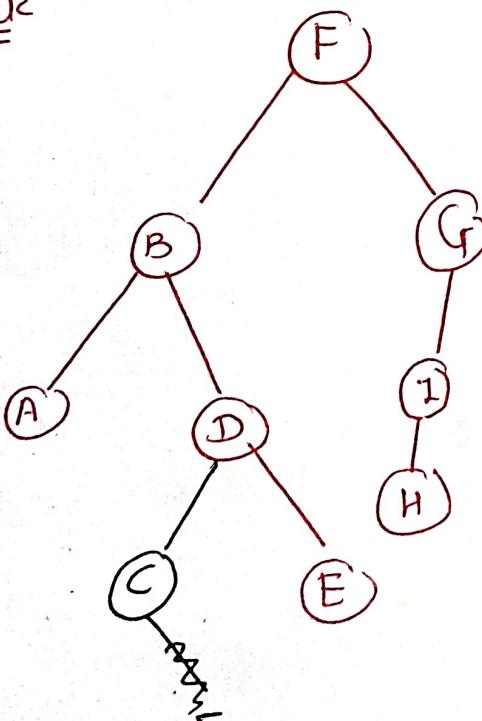
Unit No.



F B A D C E G H I

Given
Preorder Postorder
 $\overbrace{F} \overbrace{B} \overbrace{A} \overbrace{D} \overbrace{C} \overbrace{E} \overbrace{G} \overbrace{I} \overbrace{H}$
 $A \overbrace{C} \overbrace{E} \overbrace{D} \overbrace{\underbrace{B}_{\text{left}}} \overbrace{H} \overbrace{I} \overbrace{G} \overbrace{F}$
only left

Trick



preorder \rightarrow Right side e
 & postorder \rightarrow Left side
 of currently added
 or near)

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Topic

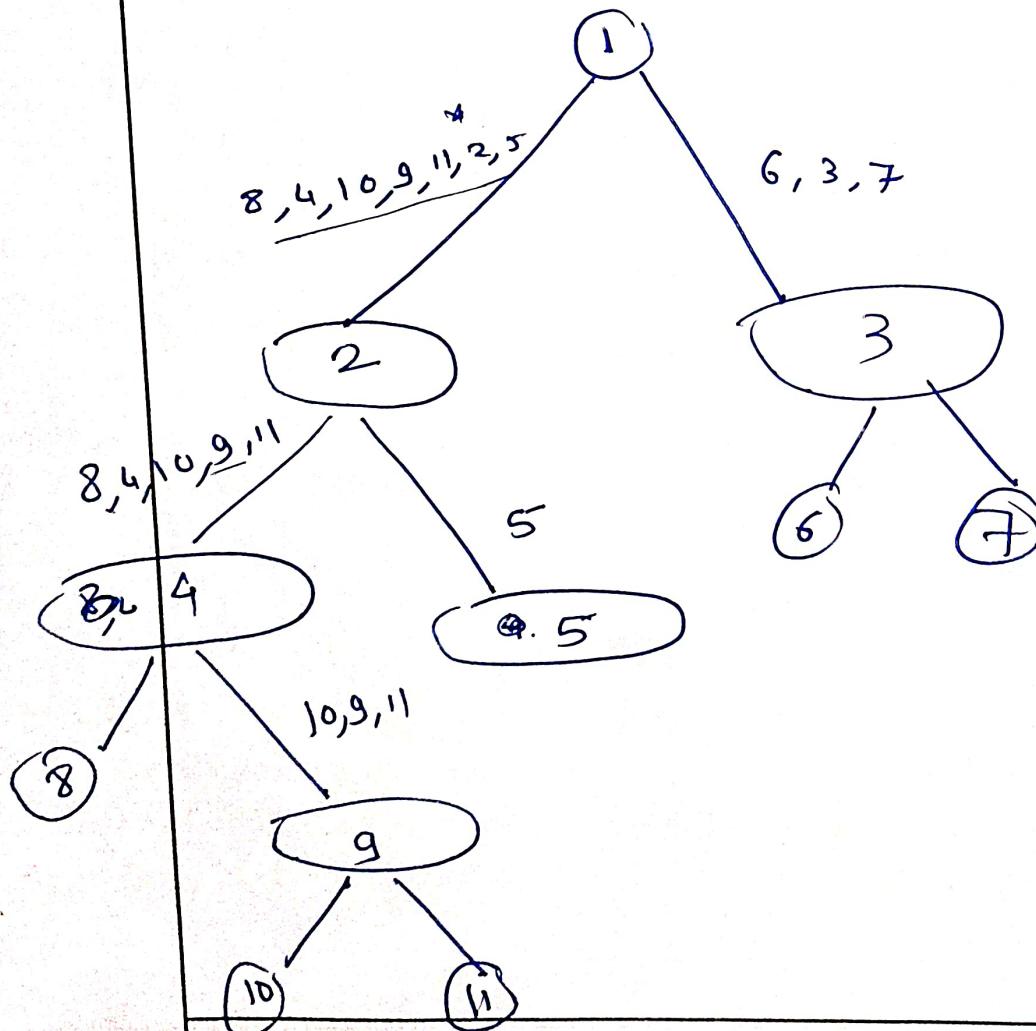
Unit No.

* Construct A Binary Tree from Preorder & Inorder

Preorder \rightarrow 1, 2, 4, 8, 9, 10, 11, 5, 3, 6, 7

Inorder - 8, 4, 10, 9, 11, 2, 5, 1, 6, 3, 7

as '1' is root because Preorder



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Topic Tree (BST)

Unit No. 02

Binary Search Tree

Definition - A BST is a BT which is either empty or in which each node contains a key that satisfies the following conditions:

- ① All keys (if any) in the left subtree of the root is less than the key in the root.
- ② The key in the root precedes all keys (if any) in its right subtree.
- ③ The right and left subtrees are again search trees.

In other words, BST is a BT

in which all the keys having values less than root will be in the left subtree and those greater than root lie in right subtree.

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Topic	Unit No.
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Algorithm creating of BST

- ① Read the key value & store it in node say 'p'
- ② if $(\text{root} == \text{NULL})$ i.e. if tree is currently empty then attach 'p' to root & goto step ④
 $\Rightarrow \boxed{\text{root} = p}$
- ③ $q = \text{root}$
 if $(P \rightarrow \text{info} < q \rightarrow \text{info})$ then
 if $(q \rightarrow \text{left} == \text{NULL})$
 then attach P to $q \rightarrow \text{left}$
 and goto step ④ i.e. $\boxed{q \rightarrow \text{left} = P}$
 else move to $q = q \rightarrow \text{left}$ and goto step ③
 (i.e. if new element to be inserted is less than root element move to left once left becomes NULL)

Topic

else

if ($q \rightarrow \text{right} == \text{NULL}$) thenattach p to $q \rightarrow \text{right}$ & goto step ④

i.e.
$$\boxed{q \rightarrow \text{right} = p}$$

else

$$q = q \rightarrow \text{right}$$

and goto step ③

(if new element is greater than root
value, move to right and connect
the node there)

Step ④ Repeat step 1-3 for all the
nodes

Step ⑤ stop

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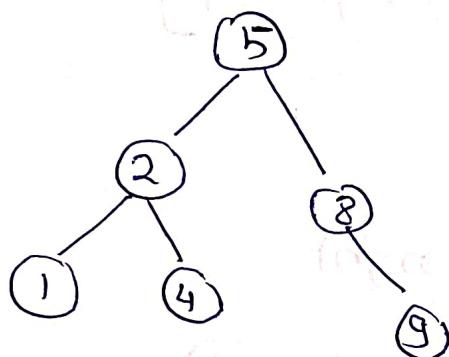
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Topic

Unit No.

Create a BST for the input elements:

5, 2, 8, 4, 1, 9



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Topic

Unit No.

Deletion In BST

For Deletion of Node in BST we must consider three cases,

① Leaf Node

These are the easiest, since what we do is delete the leaf and set the pointer of this node's parent to null.

② Nodes with a single child

These are also fairly simple, since we redirect pointer from nodes parent to the nodes child and then delete the node.

③ Nodes having both children

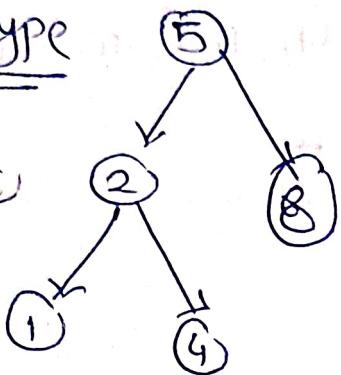
Here some arrangement is needed. what we can do is, replace the node being deleted by the rightmost node.

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Topic	Unit No.
in its left subtree of any node	1

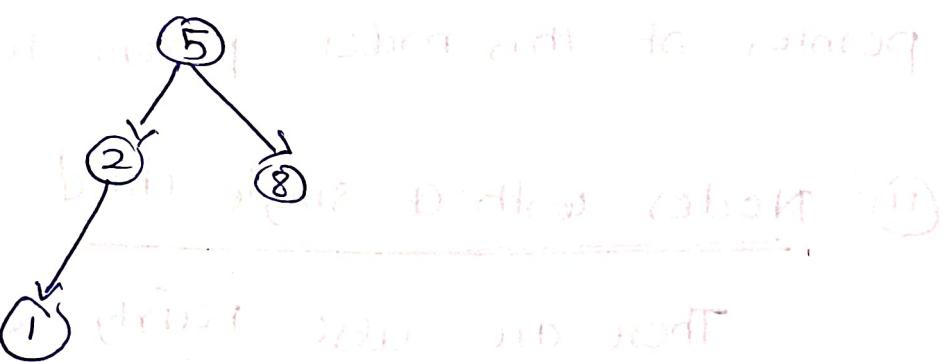
1st type in its left subtree of any node

Example ①

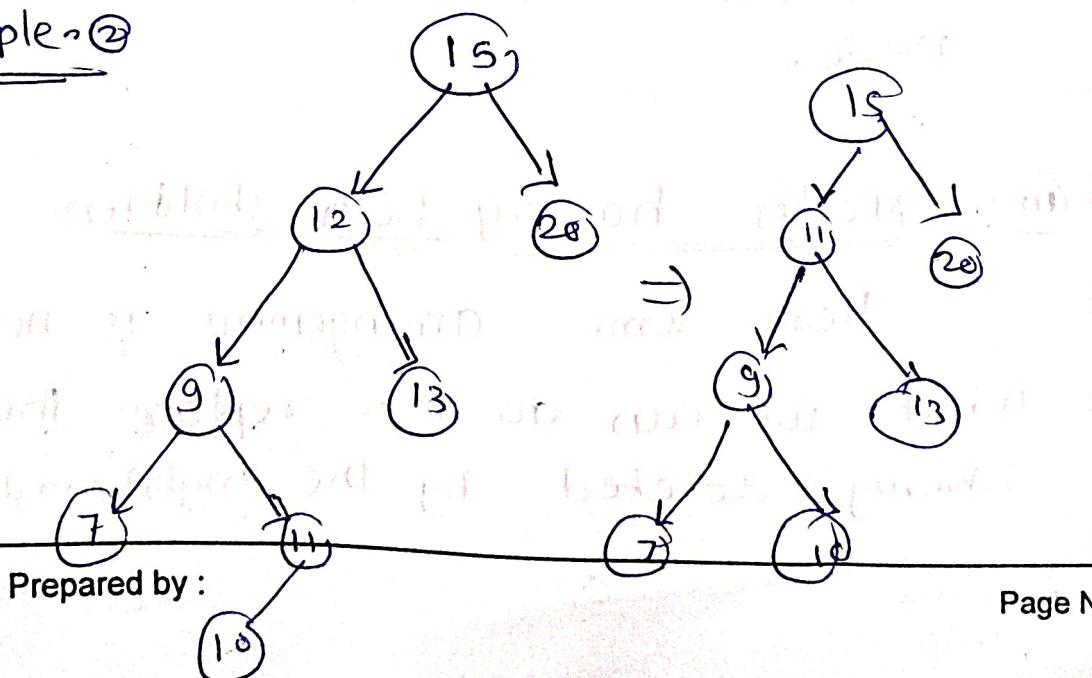


show find(1)

let us say node 4 is deleted, what will happen?



Example ②



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