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Seymour

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Chapter 1

Code

1.1 TU Matrices

Definition 1 (TU matrix). `Matrix.TU` A rational matrix is *totally unimodular* (TU) if its every subdeterminant (i.e., determinant of every square submatrix) is 0 or ± 1 .

Lemma 2 (entries of a TU matrix). `def:codetumatrixMatrix.TU.applyIfAisTU,theneveryentryofAis0or ± 1 .`

Proof sketch. `def:codetumatrixEveryentryisasquaresubmatrixofsize1,andthereforehasdeterminant(andva`
□

Lemma 3 (any submatrix of a TU matrix is TU). `def:codetumatrixMatrix.TU.submatrixLetAbea rationalmat`

Proof sketch. `def:codetumatrixAnysquaresubmatrixofBisasubmatrixofA,soitsdeterminantis0or ± 1 .`
Thus, B is TU. □

Lemma 4 (transpose of TU is TU). `def:codetumatrixMatrix.TU.transposeLetAbeaTUMatrix.Then A^T`
is TU.

Proof sketch. `def:codetumatrixAsubmatrixTof A^T` is a transpose of a submatrix of A , so $\det T \in \{0, \pm 1\}$. □

Lemma 5 (appending zero vector to TU). `def:codetumatrixMatrix.TUadjoinrow0siffLetAbea matrix.Letabe`
 $= [A/a]$ is TU exactly when A is.

Proof sketch. `def:codetumatrix,lem : codesubmatrixoftuLetTbeasquaresubmatrixofC,and supposeAisTU.If`
0. Otherwise T is a submatrix of A , so $\det T \in \{0, \pm 1\}$. For the other direction, because A is a submatrix of C , we can apply lemma 3. □

Lemma 6 (appending unit vector to TU). `def:codetumatrixLetAbea matrix.Letabea unitrow.ThenC`
 $= [A/a]$ is TU exactly when A is.

Proof sketch. `def:codetumatrix,lem : codesubmatrixoftuLetTbeasquaresubmatrixofC,andsupposeAisTU.If`
entry of the unit row, then $\det T$ equals the determinant of some submatrix of A
times ± 1 , so $\det T \in \{0, \pm 1\}$. If T contains some entries of the unit row except
the ± 1 , then $\det T = 0$. Otherwise T is a submatrix of A , so $\det T \in \{0, \pm 1\}$.
For the other direction, simply note that A is a submatrix of C , and use lemma
3. \square

Lemma 7 (TUness with adjoint identity matrix). `def:codetumatrixMatrix.TUadjoinidbelowiff,Matrix.TUa`
has determinant ± 1 .

Proof sketch. `def:codetumatrixGaussianelimination.Basis`submatrix : its columns form a basis of all columns,

Lemma 8 (block-diagonal matrix with TU blocks is TU). `def:codetumatrixMatrix.fromBlocksTU`Let A be a matrix
and A_2 are both TU. Then A is also TU.

Proof sketch. `def:codetumatrixAny`squaresubmatrixTofAhas the form

T_1	0
0	T_2

 where T_1
and T_2 are submatrices of A_1 and A_2 , respectively.

- If T_1 is square, then T_2 is also square, and $\det T = \det T_1 \cdot \det T_2 \in \{0, \pm 1\}$.
- If T_1 has more rows than columns, then the rows of T containing T_1 are linearly dependent, so $\det T = 0$.
- Similar if T_1 has more columns than rows.

\square

Lemma 9 (appending parallel element to TU). `def:codetumatrixLetAbeaTUMatrix.Let`a be some row of A . Then
 $[A/a]$ is TU.

Proof sketch. `def:codetumatrixLetTbeasquaresubmatrixofC.IfT`contains the same row twice, then the rows are
dependent, so $\det T = 0$. Otherwise T is a submatrix of A , so $\det T \in \{0, \pm 1\}$. \square

Lemma 10 (appending rows to TU). `def:codetumatrixLetAbeaTUMatrix.LetBbea`matrix whose every row is a
 $[A/B]$ is TU.

Proof sketch. `def:codetumatrix,lem : codetuaddzerorow,lem : codetuaddunitrow,lem :`
`codetuaddcopyrow` Either repeatedly apply Lemmas 5, 6, and 9 or perform a similar case analysis directly. \square

Corollary 11 (appending columns to TU). `def:codetumatrix,lem : codetuaddzerorow,lem :`
`codetuaddunitrow,lem : codetuaddcopyrow` Let A be a TU matrix. Let B be a matrix whose every column is a column of
 $[A \mid B]$ is TU.

Proof sketch. `def:codetumatrix,lem : codetuaddzerorow,lem : codetuaddunitrow,lem :`
`codetuaddcopyrow,lem : codetutransposeCT` is TU by Lemma 10 and construction, so C is TU by Lemma 4. \square

Definition 12 (\mathcal{F} -pivot). Let A be a matrix over a field \mathcal{F} with row index set X and column index set Y . Let A_{xy} be a nonzero element. The result of a \mathcal{F} -pivot of A on the *pivot element* A_{xy} is the matrix A' over \mathcal{F} with row index set X' and column index set Y' defined as follows.

- For every $u \in X - x$ and $w \in Y - y$, let $A'_{uw} = A_{uw} + (A_{uy} \cdot A_{xw}) / (-A_{xy})$.
- Let $A'_{xy} = -A_{xy}$, $X' = X - x + y$, and $Y' = Y - y + x$.

Lemma 13 (pivoting preserves TUness). *def:code_tu_matrix, def:code_pivot Let A be a TU matrix and let A_{xy} be a nonzero element. Let A' be the matrix obtained by performing a real pivot in A on A_{xy} . Then A' is TU.*

Proof sketch. def:code_tu_matrix, def:code_pivot, lem:code_tu_adjoin_id

By Lemma 7 A is TU iff every basis matrix of $[I \mid A]$ has determinant ± 1 . The same holds for A' and $[I \mid A']$.

Determinants of the basis matrices are preserved under elementary row operations in $[I \mid A]$ corresponding to the pivot in A , under scaling by ± 1 factors, and under column exchange, all of which together convert $[I \mid A]$ to $[I \mid A']$.

□

Lemma 14 (pivoting preserves TUness). *def:code_tu_matrix, def:code_pivot Let A be a matrix and let A_{xy} be a nonzero element. Let A' be the matrix obtained by performing a real pivot in A on A_{xy} . If A' is TU, then A is TU.*

Proof sketch. def:code_tu_matrix, def:code_pivot, lem:code_pivot_tuReverse_the row operations, scaling, and column

1.1.1 Minimal Violation Matrices

Definition 15 (minimal violation matrix). *def:code_tu_matrix Let A be a rational $\{0, \pm 1\}$ matrix that is not TU but all of whose proper submatrices are TU. Then A is called a *minimal violation matrix of total unimodularity* (minimal violation matrix).*

Lemma 16 (simple properties of MVMs). *def:code_mvm Let A be a minimal violation matrix.*

A is square.

$\det A \notin \{0, \pm 1\}$.

If A is 2×2 , then A does not contain a 0.

Proof sketch. def:code_mvm

If A is not square, then since all its proper submatrices are TU, A is TU, contradiction.

If $\det A \in \{0, \pm 1\}$, then all subdeterminants of A are 0 or ± 1 , so A is TU, contradiction.

If A is 2×2 and it contains a 0, then $\det A \in \{\pm 1\}$, which contradicts the previous item.

□

Lemma 17 (pivoting in MVMs). *def:code_mvm, def:code_ppivot Let A be a minimal violation matrix. Suppose A has 3 rows. Suppose we perform a real pivot in A , then delete the pivot row and column. Then the resulting matrix A' is also a minimal violation matrix.*

Proof sketch. def:code_mvm, lem:code_ddiagonal_with_tu_blocks, lem:code_reverse_pivot_tu, lem:code_pivot_tu, lem:code_submatrix_of_tu

Let A'' denote matrix A after the pivot, but before the pivot row and column are deleted.

Since A is not TU, Lemma 14 implies that A'' is not TU. Thus A' is not TU by Lemma 8.

Let T' be a proper square submatrix of A' . Let T'' be the submatrix of A'' consisting of T' plus the pivot row and the pivot column, and let T be the corresponding submatrix of A (defined by the same row and column indices as T'').

T is TU as a proper submatrix of A . Then Lemma 13 implies that T'' is TU. Thus T' is TU by Lemma 3.

□

1.2 Matroid Definitions

Definition 18 (binary matroid). StandardRepresentation Let B be a binary matrix, let $A = [I \mid B]$, and let E denote the column index set of A . Let \mathcal{I} be all index subsets $Z \subseteq E$ such that the columns of A indexed by Z are independent over \mathbb{Z}_2 . Then $M = (E, \mathcal{I})$ is called a *binary matroid* and B is called its *(standard) representation matrix*.

Definition 19 (regular matroid). StandardRepresentation.IsRegular StandardRepresentation, def:code_tu_matri.

A is a signed version of B , i.e., $|A| = B$,

A is totally unimodular.

Then M is called a *regular matroid*.

Lemma 20 (regularity is ignostic of representation). StandardRepresentation_oMatroid_isRegular_i f f Standard add

1.3 k -Separation and k -Connectivity

Definition 21 (k -separation). *StandardRepresentation* Let M be a binary matroid generated by a standard representation matrix B . Suppose that B is

partitioned as $\begin{array}{cc|cc} & & Y_1 & Y_2 \\ X_1 & B_1 & D_2 \\ X_2 & D_1 & B_2 \end{array}$ where $X_1 \sqcup X_2$ is a partition of the rows of B and $Y_1 \sqcup Y_2$ is a partition of its columns. Let $k \in \mathbb{Z}_{\geq 1}$ and suppose that

- $|X_1 \cup Y_1| \geq k$ and $|X_2 \cup Y_2| \geq k$,
- $\mathbb{Z}_2\text{-rank } D_1 + \mathbb{Z}_2\text{-rank } D_2 \leq k - 1$.

Then $(X_1 \cup Y_1, X_2 \cup Y_2)$ is called a (*Tutte*) k -separation of B and M .

Definition 22 (exact k -separation). *def:code_{ksep}Ap* k -separation is called exact if the rank condition holds with equality.

Definition 23 (k -separability). *def:code_{ksep}W* We say that B and M are (exactly) (*Tutte*) k -separable if they have an exact (k -separation).

Definition 24 (k -connectivity). *def:code_{ksep}F* For $k \geq 2$, M and B are (*Tutte*) k -connected if they have no ℓ -separation for $1 \leq \ell < k$. When M and B are 2-connected, they are also called *connected*.

1.4 Sums

1.4.1 1-Sums

Definition 25 (1-sum of matrices). *Matrix_{1sum}Composition* Let B be a matrix that can be represented as $\begin{array}{cc|c} & & Y \\ X_1 & B_1 & B_2 \\ X_2 & 0 & 0 \end{array}$

and B_2 are the two *components* of a 1-sum decomposition of B .

Conversely, a 1-sum composition with components B_1 and B_2 is the matrix B above.

The expression $B = B_1 \oplus_1 B_2$ means either process.

Definition 26 (matroid 1-sum). *StandardRepresentation, Matrix_{1sum}Composition, StandardRepresentation* Let M be a binary matroid. Then the binary matroids M_1 and M_2 represented by B_1 and B_2 , respectively, are the two *components* of a 1-sum decomposition of M .

Conversely, a 1-sum composition with components M_1 and M_2 is the matroid M defined by the corresponding representation matrix B .

The expression $M = M_1 \oplus_1 M_2$ means either process.

Lemma 27 (1-sum is commutative). *StandardRepresentation*. Is 1 sum Of Standard Representation 1 sum commutative

Proof.

□

Theorem 28 (1-sum of regular matroids is regular). *StandardRepresentation.Is1sumOf.isRegular*
StandardRepresentation.Is1sumOf,StandardRepresentation.IsRegular Let M_1 and M_2 be regular matroids. Then $M = M_1 \oplus_1 M_2$ is a regular matroid.

Conversely, if a regular matroid M can be decomposed as a 1-sum $M = M_1 \oplus_1 M_2$, then M_1 and M_2 are both regular.

Proof sketch. *StandardRepresentation.Is1sumOf,StandardRepresentation.IsRegular,lem:code_ddiagonal_wwith_tub_l*
extractintolemmasaboutTUmatrices Let B, B_1 , and B_2 be the representation matrices of M, M_1 , and M_2 , respectively.

- Converse direction. Let B' be a TU signing of B . Let B'_1 and B'_2 be signings of B_1 and B_2 , respectively, obtained from B . By Lemma 3, B'_1 and B'_2 are both TU, so M_1 and M_2 are both regular.
- Forward direction. Let B'_1 and B'_2 be TU signings of B_1 and B_2 , respectively. Let B' be the corresponding signing of B . By Lemma 8, B' is TU, so M is regular.

□

Lemma 29 (left summand of regular 1-sum is regular). *StandardRepresentation.Is1sumOf.isRegular_{left}StandardRepresentation.IsRegular*
add

Proof. *StandardRepresentation.Is1sumOf,StandardRepresentation.IsRegular,lem:code_submatrix_of_tu* □

Lemma 30 (right summand of regular 1-sum is regular). *StandardRepresentation.Is1sumOf.isRegular_{right}StandardRepresentation.IsRegular*
add

Proof. *StandardRepresentation.Is1sumOf,StandardRepresentation.IsRegular,lem:code_submatrix_of_tu* □

1.4.2 2-Sums

Definition 31 (2-sum of matrices). *Matrix₂sumComposition* Let B be a matrix of the form $\begin{matrix} & Y_1 & Y_2 \\ X_1 & A_1 & 0 \\ X_2 & D & A_2 \end{matrix}$ Let B_1

be a matrix of the form $\begin{matrix} & Y_1 \\ X_1 & A_1 \\ \text{Unit} & x \end{matrix}$ Let B_2 be a matrix of the form $\begin{matrix} & Y_2 \\ X_2 & A_2 \\ \text{Unit} & y \end{matrix}$

Suppose that \mathbb{Z}_2 -rank $D = 1$, $x \neq 0$, $y \neq 0$, $D = y \cdot x$ (outer product).

Then we say that B_1 and B_2 are the two *components* of a *2-sum decomposition* of B .

Conversely, a *2-sum composition* with *components* B_1 and B_2 is the matrix B above.

The expression $B = B_1 \oplus_2 B_2$ means either process.

Definition 32 (matroid 2-sum). *StandardRepresentation.Is2sumOfStandardRepresentation,Matrix₂sumComposition*
 and B_2 satisfy the assumptions of Definition 31. Then the binary matroids M_1

and M_2 represented by B_1 and B_2 , respectively, are the two *components* of a 2-sum decomposition of M .

Conversely, a 2-sum composition with components M_1 and M_2 is the matroid M defined by the corresponding representation matrix B .

The expression $M = M_1 \oplus_2 M_2$ means either process.

Lemma 33 (2-sum of TU matrices is a TU matrix). *StandardRepresentation_{2sum}, sRegularMatrix_{2sumComp} and B_2 be TU matrices. Then $B = B_1 \oplus_2 B_2$ is a TU matrix.*

Proof sketch. Matrix_{2sumComposition}, def : code_{tu}matrix, lem : code_{tu}add_{okrows}, cor : code_{tu}add_{okcols}, lem : code_{vm}pivot, lem : code_{vm}props

Let B'_1 and B'_2 be TU signings of B_1 and B_2 , respectively. In particular, let A'_1 , x' , A'_2 , and y' be the signed versions of A_1 , x , A_2 , and y , respectively. Let B' be the signing of B where the blocks of A_1 and A_2 are signed as A'_1 and A'_2 , respectively, and the block of D is signed as $D' = y' \cdot x'$ (outer product).

Note that $[A'_1/D']$ is TU by Lemma 10, as every row of D' is either zero or a copy of x' . Similarly, $[D' | A'_2]$ is TU by Corollary 11, as every column of D' is either zero or a copy of y' . Additionally, $[A'_1 | 0]$ is TU by Corollary 11, and $[0/A'_2]$ is TU by Lemma 10.

todo: prove lemma below, separate into statement about TU matrices

Lemma: Let T be a square submatrix of B' . Then $\det T \in \{0, \pm 1\}$.

Proof: Induction on the size of T . *Base:* If T consists of only 1 element, then this element is 0 or ± 1 , so $\det T \in \{0, \pm 1\}$. *Step:* Let T have size t and suppose all square submatrices of B' of size $\leq t - 1$ are TU.

- Suppose T contains no rows of X_1 . Then T is a submatrix of $[D' | A'_2]$, so $\det T \in \{0, \pm 1\}$.
- Suppose T contains no rows of X_2 . Then T is a submatrix of $[A'_1 | 0]$, so $\det T \in \{0, \pm 1\}$.
- Suppose T contains no columns of Y_1 . Then T is a submatrix of $[0/A'_2]$, so $\det T \in \{0, \pm 1\}$.
- Suppose T contains no columns of Y_2 . Then T is a submatrix of $[A'_1/D']$, so $\det T \in \{0, \pm 1\}$.
- Remaining case: T contains rows of X_1 and X_2 and columns of Y_1 and Y_2 .
- If T is 2×2 , then T is TU. Indeed, all proper submatrices of T are of size ≤ 1 , which are $\{0, \pm 1\}$ entries of A' , and T contains a zero entry (in the row of X_2 and column of Y_2), so it cannot be a minimal violation matrix by Lemma 16. Thus, assume T has size ≥ 3 .
- . todo: complete proof, see last paragraph of Lemma 11.2.1 in Truemper

□

Theorem 34 (2-sum of regular matroids is a regular matroid). *StandardRepresentation.Is2sumOf.isRegular StandardRepresentation.Is2sumOf,StandardRepresentation.IsRegular*
Let M_1 and M_2 be regular matroids. Then $M = M_1 \oplus_2 M_2$ is a regular matroid.

Proof sketch. *StandardRepresentation.Is2sumOf,StandardRepresentation.IsRegular,Matrix2sumComposition*
and B_2 be the representation matrices of M , M_1 , and M_2 , respectively. Apply Lemma 33. \square

Lemma 35 (left summand of regular 2-sum is regular). *StandardRepresentation.Is2sumOf.isRegular_leftStandardRepresentation.IsRegular*
add

Lemma 36 (right summand of regular 2-sum is regular). *StandardRepresentation.Is2sumOf.isRegular_rightStandardRepresentation.IsRegular*
add

1.4.3 3-Sums

Definition 37 (3-sum of matrices). *Matrix3sumComposition* *todo* : *add*

Definition 38 (matroid 3-sum). *StandardRepresentation.Is3sumOf StandardRepresentation,Matrix3sumComposition*
add

Theorem 39 (3-sum of regular matroids is regular). *StandardRepresentation.Is3sumOf.isRegular*
StandardRepresentation.Is3sumOf,StandardRepresentation.IsRegular *todo*: *add*

Lemma 40 (left summand of regular 3-sum is regular). *StandardRepresentation.Is3sumOf.isRegular_leftStandardRepresentation.IsRegular*
add

Lemma 41 (right summand of regular 3-sum is regular). *StandardRepresentation.Is3sumOf.isRegular_rightStandardRepresentation.IsRegular*
add