## ------ASSIGNMENTS 1-----

#### Q1 ans-

- O(logn).
- Instead of using Nested loop, we can into a single loop and update a given array by using first loops which makes time complexity is O(n).

#### Q2 ans-

Given: T(n) = 3T(n-1) + 12n T(0) = 5

Let's calculate T(1) using the recurrence relation:

$$T(1) = 3T(1-1) + 12(1) T(1) = 3T(0) + 12 T(1) = 3(5) + 12 T(1) = 15 + 12 T(1) = 27$$

Now, let's calculate T(2) using the recurrence relation:

$$T(2) = 3T(2-1) + 12(2) T(2) = 3T(1) + 24 T(2) = 3(27) + 24 T(2) = 81 + 24 T(2) = 105$$

Therefore, the value of T(2) for the given recurrence relation is 105.

#### Q3 ans-

we'll assume a closed form solution of the form T(n) = a\*n + b, where a and b are constants to be determined.

Let's substitute this assumed solution into the recurrence relation:

$$T(n) = T(n - 1) + c$$

Substituting T(n) = an + b and T(n - 1) = a(n - 1) + b:

$$an + b = a(n - 1) + b + c$$

Simplifying the equation:

$$an + b = an - a + b + c$$

a\*n and b cancel out, resulting in:

$$0 = -a + c$$

Solving for a:

a = c

Now that we have found the value of a, we can substitute it back into the assumed solution to find b:

$$0 = -c + c$$

Since a = c, this equation simplifies to:

0 = 0

This equation doesn't provide any information about b. Therefore, we can choose any value for b.

Let's assume b = k, where k is a constant.

Finally, the solution to the recurrence relation T(n) = T(n - 1) + c is:

$$T(n) = c*n + k$$

where c is a constant, and k is an arbitrary constant.

#### Q4 ans-

To determine the time complexity of the recurrence relation  $T(n) = 16T(n/4) + n^2\log(n)$  using the master theorem, we need to compare the equation to the standard form of the master theorem:

$$T(n) = aT(n/b) + f(n)$$

In this case, a = 16, b = 4, and  $f(n) = n^2\log(n)$ .

Now, we can compare the values of f(n) to n^log\_b(a) and analyze the different cases:

- 1. If f(n) is polynomially smaller than  $n^{\log}b(a)$ , i.e.,  $f(n) = O(n^{c})$  for some constant  $c < \log_{b}(a)$ , then the time complexity is dominated by the recursive term T(n/b).
- 2. If f(n) and  $n^{\log}b(a)$  have the same growth rate, i.e.,  $f(n) = \Theta(n^{c} \log^{k}k(n))$ , where  $c = \log_{b}(a)$ , then the time complexity is affected by both the recursive term T(n/b) and the additional term f(n).
- 3. If f(n) is polynomially larger than  $n^{\log}b(a)$ , i.e.,  $f(n) = \Omega(n^{c})$ , where  $c > \log_{b}(a)$ , then the time complexity is dominated by the additional term f(n).

In our case, a = 16, b = 4, and  $f(n) = n^2\log(n)$ . Let's calculate  $n^\log_b(a)$ :

$$n^{\log_4(16)} = n^2$$

Comparing f(n) to n^log\_b(a), we have:

n^2log(n) vs. n^2

f(n) is in case 2 since they have the same growth rate.

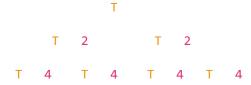
Therefore, using the master theorem, the time complexity of the recurrence relation  $T(n) = 16T(n/4) + n^2\log(n)$  is  $\Theta(n^2\log(n))$ .

#### Q5 ans-

At each level of the tree, we have the recurrence relation T(n) = 2T(n/2) + K, and we keep dividing n by 2 until we reach the base case T(1) or T(0).

Let's draw the recursion tree:

scssCopy code



We can observe that the tree has a total of log n levels (base 2) since we keep dividing n by 2 at each level until we reach 1.

Now, let's calculate the cost at each level:

- At the root level, the cost is K.
- At level 1, there are two recursive calls of size n/2 each, resulting in a cost of K for each subtree, totaling
   2K.
- At level 2, there are four recursive calls of size n/4 each, resulting in a cost of K for each subtree, totaling
   4K
- At level 3, there are eight recursive calls of size n/8 each, resulting in a cost of K for each subtree, totaling 8K.
- And so on...

We can see that at each level, the total cost is K. Since we have log n levels, the total cost of the recursion tree is:

Total cost = K \* log n = K log n

Therefore, using the recursion tree method, the time complexity of the recurrence relation T(n) = 2T(n/2) + K is  $\Theta(\log n)$ .

## ------ASSIGNMENT 2-----

### Q1 ans -

```
public class Assignment_Question {
  public static void main(String[] args) {
  int[] arr = {3,20,4,6,9};
  int sum = 0;
  for(int i = 0; i< arr.length;i++)
  {
  if(i%2==0)
  {
    sum+=arr[i];
  }
  }
  System.out.println(sum);
  }
}</pre>
```

### Q2 ans -

```
public class Assignment_Question {
  public static void main(String[] args) {
  int[] arr = {34, 21, 54, 65, 43};
  for (int i = 0; i < arr.length; i++) {
   if (arr[i] % 2 == 0) {
    System.out.print(arr[i] + " ");
  }
  }
}</pre>
```

### Q3 ans-

```
public class Assignment_Question
{
  public static void main(String[] args)
  {
   int[] arr = {34,21,54,65,43};
   int max = Integer.MIN_VALUE;
  for(int i = 0; i< arr.length;i++)
   {
   if(arr[i] > max)
   {
   max = arr[i];
   }
  }
  System.out.println(max);
  }
}
```

### Q4 ans-

```
public class Assignment_Question
{
  public static void main(String[] args)
  {
  int[] arr = {34,21,54,65,43};
  Arrays.sort(arr);
  System.out.println(arr[arr.length-2]);
```

```
}
}
```

### Q5 ans -

```
public class Assignment_Question
{
  public static void main(String[] args)
  {
  int[] arr = {34,21,54,65,43};
  int max = Integer.MIN_VALUE;
  for(int i = 0; i< arr.length;i++)
  {
  if(arr[i] > max)
  {
   max = arr[i];
  }
  }
  System.out.println(max);
  }
}
```

-----2D Arrays Assignment-----

## Q1 ans -

```
import java.util.*;

public class check {
  public static void main(String[] args) {
    Scanner sc = new Scanner(System.in);
    System.out.print("Enter the number of rows : ");
    int m = sc.nextInt();
    System.out.print("Enter the number of Columns : ");
    int n = sc.nextInt();
```

```
int[][] arr = new int[m][n];
int number_of_positive_numbers = 0;
int number_of_negative_numbers = 0;
int number_of_odd_numbers = 0;
int number_of_even_numbers = 0;
int number_of_0 = 0;
System.out.println("Enter the elements");
for(int i = 0; i<m;i++)
{
for(int j=0;j<n;j++)
arr[i][j] = sc.nextInt();
if(arr[i][j]>0)
number_of_positive_numbers++;
else if (arr[i][j] < 0)
number_of_negative_numbers++;
if (arr[i][j]%2== 1 || arr[i][j] %2 == -1)
number_of_odd_numbers++;
if(arr[i][j] %2 == 0 || arr[i][j] == 0)
number_of_even_numbers++;
if (arr[i][j] == 0)
number_of_0++;
}
}
}
System.out.println("number_of_positive_numbers: "+ number_of_positive_numbers);
System.out.println("number_of_negative_numbers: "+number_of_negative_numbers);
System.out.println("number_of_odd_numbers: "+number_of_odd_numbers);
System.out.println("number_of_even_numbers: "+number_of_even_numbers);
System.out.println( "number_of_0 : " +number_of_0);
```

```
}
}
```

## Q2 ans-

```
public class Diagnoal {
  public static void main(String[] args) {
  int[][] arr = {
    {1,2,3},
    {4,5,6},
    {7,8,9}
  };
  int n = arr.length;
  for(int i = 0; i<n;i++)
  {
    for(int j =0; j< n-i-1;j++)
    {
        System.out.print(arr[i][j]+" ");
    }
    }
}</pre>
```

# Q3 ans -

```
else if(j == n-i-1)
{
   System.out.print(arr[i][j]+" ");
}
}
}
```

## Q4 ans-

```
public class Largest_2D {
public static void main(String[] args) {
int[][] arr1 =
{1,2,4,0},
{2,5,7,-1},
{4,2,6,9}
};
int n = arr1.length;
int m = arr1[0].length;
int max = Integer.MIN_VALUE;
for(int i = 0; i<n; i++)
for(int j = 0; j < m; j++)
if(arr1[i][j] > max)
max = arr1[i][j];
}
System.out.println(max);
}
```

# Q5 ans-

```
public class Middle_2D {
public static void main(String[] args) {
int [][] arr2 = {
{1,2,3,4,5},
{3,4,5,6,7},
{7,6,5,4,3},
{8,7,6,5,4},
{1,2,37,8,0}
};
int n = arr2.length;
for(int i = 0; i<n;i++)
System.out.print(arr2[i][n/2]+" ");
for(int j = 0; j < n; j++)
if(j == n/2)continue;
System.out.print(arr2[n/2][j]+" ");
}
}
}
```