# 1 ADD, ADC and EFLAGS

## 1.1 The ADD instruction

Consider the following program:

```
section .data

section .text

global _start

_start:

MOV eax, 1

MOV ebx, 2

ADD eax, ebx

INT 80h
```

The instruction ADD eax, ebx means the same as: eax = eax + ebx.

## 1.1.1 Debugging

Below are the values of the registers **eax** and **ebx** in before and after performing addition.

```
(gdb) info registers eax
eax 0x1 1
(gdb)
```

Figure 1: Value of eax after addition

```
(gdb) info registers ebx
ebx 0x2 2
(gdb)
```

Figure 2: Value of  $\mathbf{ebx}$  before addition

```
(gdb) info registers eax
eax 0x3 3
(gdb)
```

Figure 3: Value of eax after addition

### 1.2 EFLAGS

There is one more register called the **eflags** register. It is a 32-bit register used to indicate the **status/information** about the computations and control CPU operations. Each bit represents a specific flag.

```
(gdb) info registers eflags
eflags 0x206 [ PF IF ]
(gdb)
```

Figure 4: The **eflags** register.

Here **eflags** is giving information about the **ADD** operation and we see that two of its flags are set **PF** and **IF**.

**PF** is the **P**arity **F**lag. This flag is *set* if the value of an operation is **odd**. If the value of an operation is an even number then PF is set to 0. So, the **ADD** operation resulted into an odd value which is 3 and that's why PF was set.

**IF** is the Interrupt Flag. This flag is *set* to 1 when we allow interrupts to be done on our system.

The figure below shows the the **eflags** register:



Figure 5: The eflags register.[Image source: GeeksForGeeks]

<u>NOTE</u>: In the assembly program above, include the line: MOV eax, 1 before the interrupt statement otherwise the program won't exit and it will throw segmentation fault.

### 1.2.1 More flags of the EFLAGS register

Consider the following program:

```
section .data

section .text

MOV al, Ob11111111

MOV bl, Ob0001

ADD al, bl

MOV eax, 1

INT 80h
```

The notation **0b** indicates a binary number. Therefore **0b11111111** indicates a binary number. Basically, this program adds two numbers. One is -1(in **al**) and the other is 1(in **bl**). We know that the result is 0 but we want to see how this is operation performed and if any carry is generated then where that carry is stored.

```
(gdb) si
0x08049002 in _start ()
(gdb) info registers al
al 0xff -1
(gdb)
```

Figure 6: al stores -1

```
(gdb) si

0x08049004 in _start ()

(gdb) info registers bl

bl 0x1 1

(gdb)
```

Figure 7: bl stores 1

```
(gdb) si

0x08049006 in _start ()

(gdb) info registers al

al 0x0 0

(gdb)
```

Figure 8: al stores 0

```
(gdb) info registers eax
eax 0x0 0
(gdb)
```

Figure 9: eax also stores 0

```
(gdb) info registers eflags
eflags 0x257 [ CF PF AF ZF IF ]
(gdb)
```

Figure 10: eflags has some set bits

The **ADD** operation looks like this:



This operation was shown for **al** and **bl**. So when we check the register **al**(Figure 8), we see that it stores 0. But interestingly enough, when we check the **eax** register(Figure 9), we see that it also stores 0.

The reason is that, the **ADD** operation was performed for only 1 byte registers(**al** and **bl** in this case). So, when the extra 1(carry) resulted from the addition, it was put into the **CF** flag of the **eflags** register and that is the reason why we see 0 as the value of **eax** as well.

 ${f CF}$  is the Carry Flag. When there is a carry from the previous operation, then it's set to 1.

**ZF** is the **Z**ero **F**lag. It is set if the register that results from the operation is 0. In our case, **al** got set to 0.

#### 1.3 The ADC instruction

Consider this assembly program:

```
section .data
1
    section .text
3
             global _start
              _{	t start:}
5
                       MOV al, 0b11111111
6
                       MOV bl, 0b0001
                       ADD al, bl ; al = al + bl
8
                       ADC ah, 0
9
                       MOV eax, 1
10
                       INT 80h
11
12
```

ADC instruction is called add with carry. It works similar to the **ADD** operation. It takes a destination and adds a source to it. It then adds the CF to the destination.

In this program we are using the **ah**(higher 8 bits) to store the carry. The instruction **ADC ah**, **0** will basically add 0 to **ah** and then will also add the CF(the carry bit in CF) to it and store the result in **ah**.

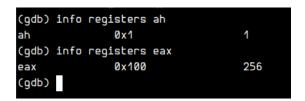


Figure 11: Value of ah and eax

We can see that **ah** stores 1(which is the CF) and **eax** stores 256.

Now the reason **eax** stores 256 is that it is a 32-bit register and **ah** is its higher 8-bits. So we know that the lower 8 bits(**al**) were all zeros(from addition) and the higher 8-bits were all 0 as well(from the previous addition). But now, we used the **ADC** instruction to put the carry into the higher 8 bits of **eax** which make it look like this:

 $00000000\ 00000000\ 00000001\ 00000000$ 

which is 256 in decimal.