

# Parallel Implementation of Graph Algorithms

Group 27

Bhagyashri Bhamare (181IT111)

Chinmayi C. Ramakrishna (181IT113)

K. Keerthana (181IT221)

Utkarsh Meshram (181IT250)

# Introduction

- ❑ The **Single-Source Shortest Path** (SSSP) problem consists of finding the shortest paths between a given vertex  $v$  and all other vertices in the graph.
- ❑ Four algorithms for parallel implementation of SSSP problem : Dijkstra's algorithm, Bellman Ford algorithm, Floyd Warshall algorithm and Prim's algorithm.
- ❑ Parallel implementation of these algorithms using OpenMP.
- ❑ Dijkstra's algorithm is a graph search algorithm that solves single-source shortest path for a graph with non-negative weights.
- ❑ The Floyd Warshall is a classic dynamic programming algorithm that solves the all-pairs shortest path (APSP) problem on directed weighted graphs. Floyd Warshall can be applied to graphs with negative weight edges to determine whether the graph has negative cycles or not.
- ❑ Prim's algorithm is a popular greedy algorithm that finds a minimum spanning tree for a weighted undirected graph.

# Problem Statement

Solving Single Source Shortest Path Problem using:

- ❑ Dijkstra's algorithm
- ❑ Bellman Ford algorithm
- ❑ Floyd Warshall algorithm
- ❑ Prim's algorithm

# Objectives

- ❑ To parallelise the Shortest Source Path Problem using four algorithms: Dijkstra's algorithm, Bellman Ford algorithm, Floyd Warshall algorithm and Prim's algorithm.
- ❑ To compare sequential and parallel execution of each algorithm.
- ❑ To compare the algorithms.
- ❑ To find the best algorithm for different size of graphs.

# Literature Survey

S. No	References	Work Done
1.	F. Busato and N. Bombieri An efficient implementation of the Bellman-Ford algorithm for Kepler GPU architectures.	Use of frontier data structure to implement Bellman-Ford algorithm.
2.	H. Ortega-Arranz et. al. A new GPU-based approach to the shortest path problem.	Parallelizing internal working of sequential Dijkstra algorithm.
3.	Vladimir Lonc̃ar, Srdjan Škrbic´ and Antun Balaz̃ Parallelization of Minimum Spanning Tree Algorithms Using Distributed Memory Architectures	Parallelizing Minimum Spanning Tree using Prim's algorithm and Kruskal algorithm.
4.	An Implementation of Parallel Floyd-Warshall Algorithm Based on Hybrid MPI and OpenMP	Parallelizing Floyd-Warshall Algorithm to exploit the parallelism inside a multi-core node computer.

# Serial Implementation

# Dijkstra's Algorithm

```
function Dijkstra(Graph, source):  
    for each vertex v in Graph:  
        dist[v] := infinity  
        previous[v] := undefined  
    dist[source] := 0  
    Q := the set of all nodes in Graph  
    while Q is not empty:  
        u := node in Q with smallest dist[ ]  
        remove u from Q  
        for each neighbor v of u:  
            alt := dist[u] + dist_between(u, v)  
            if alt < dist[v]  
                dist[v] := alt  
                previous[v] := u  
    return previous[ ]
```

# Bellman Ford Algorithm

```
function bellmanFord(G, S)
  for each vertex V in G
    distance[V] <- infinite
    previous[V] <- NULL
  distance[S] <- 0
  for each vertex V in G
    for each edge (U,V) in G
      tempDistance <- distance[U] + edge_weight(U, V)
      if tempDistance < distance[V]
        distance[V] <- tempDistance
        previous[V] <- U
  for each edge (U,V) in G
    If distance[U] + edge_weight(U, V) < distance[V]
      Error: Negative Cycle Exists
  return distance[], previous[]
```



# Floyd Warshall Algorithm

```
function FloydWarshall(Ak,n):  
    n = no of vertices  
    A = matrix of dimension n*n  
    for k = 1 to n  
        for i = 1 to n  
            for j = 1 to n  
                A[i, j] = min (Ak-1[i, j], Ak-1[i, k] + Ak-1[k, j])  
    return A
```

# Prim's Algorithm

```
function Prim:
    T =  $\emptyset$ ;
    U = { 1 }
    while (U  $\neq$  V)
        let (u, v) be the lowest cost edge such that u  $\in$  U and v  $\in$  V - U
            T = T  $\cup$  { (u, v) }
            U = U  $\cup$  {v}
```

# Parallel Implementation

# Dijkstra's Algorithm

```
#pragma omp parallel for schedule(runtime) private(i)
    for(i = 0; i < V; i++)
        if(vertices[i].visited == FALSE)
            int c = findEdge( u, vertices[i], edges, weights);
            len[vertices[i].title] = minimum(len[vertices[i].title],
len[u.title] + c);
```

# Bellman Ford Algorithm

```
function bellmanFord:
    omp_set_num_threads(p)
    while(!queue.empty() and no negative cycle):
        u <- queue.front()
        queue.pop()
        In_queue[u] = false
    #pragma omp parallel for
        for (int v = 0; v < n; v++):
            weight = mat[u * n + v];
            if (weight < INF):
                new_dist = weight + dist[u];
                if (new_dist < dist[v]):
                    dist[v] = new_dist;
                    enqueue_counter[v]++;
                    if (in_queue[v] == false)
                        in_queue[v] = true;
                    if (enqueue_counter[v] >= n)
                        *has_negative_cycle = true;
    #pragma omp critical
        queue.push(v)
```



# Prim's Algorithm

```
function minKey(key[], visited[]):  
    min = INT_MAX, index, i;  
    #pragma omp parallel  
        index_local = index;  
        min_local = min  
    #pragma omp for nowait  
        for (i to n)  
            if (visited[i] == false and key[i] < min_local)  
                min_local = key[i]  
                index_local = i  
    #pragma omp critical  
        if (min_local < min)  
            min = min_local  
            index = index_local  
    return index
```

# Individual Contribution

Bhagyashri Bhamare (181IT111)	Bellman Ford algorithm
K. Keerthana (181IT221)	Dijkstra's algorithm
Utkarsh Meshram (181IT250)	Floyd Warshall algorithm
Chinmayi C. Ramakrishna (181IT113)	Prim's algorithm