CLAW language specification

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Conventions used in this document

Keywords and punctuation that are part of the actual specification will appear in typewriter font:

!\$claw

Italic font is used where a keyword or other name must be used:

!Sclaw directive-name

Optional syntax is enclosed in square brackets; an option that may be repeated more than once is followed by ellipses:

```
!$claw directive-name [clause [[,] clause]. . . ]
```

1 General information about the CLAW language

The directives are either local or global.

- Local directive: those directives have a limited impact on a local block of code (for example, only in a subroutine or only in a do statement block)
- Global directive: those directives can have an impact on the whole translation unit.

The CLAW transformations are separated in two sections:

- CLAW high-level abstraction transformations: perform transformations based on abstraction based on domain science.
- CLAW low-level directive transformations: perform basic transformations based on the code structure.

Loop transformation

Specific transformation

Accelerator abstraction/helper transformation

Utility transformation

1.1 Line continuation

CLAW directives can be defined on several lines using the continuation line symbol. The syntax is described as follows:

```
1 !$claw directive clause & 2 !$claw clause
```

1.2 Interpretation order of the CLAW directives

The CLAW transformations can be combined together. For example, loop-fusion and loop-interchange can be used together in a group of nested loops.

The default application order of the transformations is defined as follows:

- 1. ignore
- 2. remove
- 3. primitive
- 4. array-transform
- 5. loop-extract
- 6. loop-fusion
- 7. loop-hoist
- 8. loop-interchange
- 9. on-the-fly
- 10. kcache
- 11. parallelize
- 12. parallelize forward
- 13. formatting transformation (internal use transformation)

Users must be aware that transformations are applied sequentially and therefore, a transformation can be performed on code that has already been transformed.

1.3 Accelerator compile guard

As CLAW generates accelerator directives, it might be useful sometimes to protect the compilation of source code including accelerator directives as it might be wrong if the CLAW FORTRAN compiler didn't process it.

For this reason, the CLAW language has a compile guard that is removed by the CLAW FORTRAN compiler while it applies the other transformations. This compile guard is designed to make the compilation failed when transformations are not applied and accelerator directives are present in the source code.

The compile guard is shown in the listing below. It uses the target accelerator language with a specific claw directive. Example for OpenACC on line 2 and OpenMP on line 3.

This line is then removed once the transformations are applied.

2 CLAW high-level abstraction transformations

2.1 Column model abstraction

The parallelize directive is used to parallelize a column based algorithm. In weather prediction models, the physical parametrizations are column independent problems. This means that the algorithm to compute the different output fields, can be defined with as a single do statement iterating over the vertical dimension and there is no horizontal dependency. The parallelize directive is then used to parallelize efficiently this column algorithm over the domain (horizontal dimensions). The directive is local to a subroutine. Therefore, as a pre-condition, the column algorithm must be enclosed in a module subroutine.

```
!$claw define dimension dim_id(lower_bound:upper_bound) &
[!$claw define dimension dim_id(lower_bound:upper_bound) &] ...
!$claw parallelize [data(data(var_1[,var_2] ...)] [over (dim_id|:[,dim_id|:] ...)]
```

Options and details

• define dimension: Define a new dimension on which the column model will be parallelized. The lower bound and upper bound of the defined dimension can be either an integer constant or an integer variable. If a variable is given, it will be added in-order to the signature of the subroutine as an INTENT(IN) parameter.

- data (optional): Define a list of variables to be promoted and bypass the automatic promotion deduction.
- over (optional): Define the location of the new dimensions in promoted variables. The : symbol reflects the position of the current dimensions before the promotion.

If the data clause is omitted, the automatic promotion of variables is done as follows: Array variables declared with the INOUT or OUT intents are automatically promoted with the defined dimensions. In function of the target, intermediate scalar or array variables may also be promoted. Scalar/array variables placed on the left hand-side of an assignment statement will be promoted if the right hand-side references any variables already promoted. If the over clause is omitted, the dimensions are added on the left-side following the the definition order if there is multiple new dimensions. The over clause allows to change this behavior and to place the new dimensions where needed.

As the subroutine/function signature is updated by the parallelize directive with new dimensions, the call graph that leads to this subroutine/function must be updated as well. For this purpose, the clause forward allows to replicated the changes along the call graph. Each subroutine/function along the call graph must be decorated with this directive.

```
! Call of a parallelized subroutine
!$claw parallelize forward
CALL one_column_fct(x,y,z)
! Call of a parallelized function
!$claw parallelize forward
result = one_column_fct(x,y,z)
```

The different variables declared in the current function/subroutine might be promoted if the call requires it. In this case, the array references in the function/subroutine are updated accordingly. This can lead to extra promotion of local scalar variables or arrays when needed.

The first function call in the call graph might be an iteration over several column to reproduce the algorithm on the grid for testing purpose. If the directive with the forward clause is placed just before one or several nested do statements with a function call, the corresponding do statements will be removed and the function call updated accordingly.

```
1  !$claw parallelize forward
2  DO i=istart, iend
3  CALL one_columne_fct(x,y(i,:),z(i,:))
END DO
```

Code example

Simple example of a column model wrapped into a subroutine and parallelize with CLAW.

```
1 PROGRAM model
```

```
2
     USE mo_column, ONLY: compute_column
3
 4
     REAL, DIMENSION (20,60) :: q, t ! Fields as declared in the whole model
     INTEGER :: nproma, nz
 5
                                     ! Size of array fields
 6
     INTEGER :: p
                                      ! Loop index
7
 8
      nproma = 20
9
      nz = 60
10
     DO p = 1, nproma
11
12
       q(p,1) = 0.0
13
       t(p,1) = 0.0
     END DO
14
15
16
      !$claw parallelize forward
17
     DO p = 1, nproma
       CALL compute_column(nz, q(p,:), t(p,:))
18
19
     END DO
20
21
     PRINT*, SUM(q)
22
     PRINT*, SUM(t)
23
   END PROGRAM model
```

```
1
   MODULE mo column
2
     IMPLICIT NONE
3
   CONTAINS
4
     ! Compute only one column
5
     SUBROUTINE compute_column(nz, q, t)
6
       IMPLICIT NONE
7
8
       INTEGER, INTENT(IN) :: nz ! Size of the array field
       REAL, INTENT(INOUT) :: t(:) ! Field declared as one column only
9
       REAL, INTENT (INOUT)
10
                             :: q(:) ! Field declared as one column only
11
       INTEGER :: k
                                      ! Loop index
12
       REAL :: C
                                      ! Coefficient
13
       REAL :: d
                                      ! Intermediate variable
14
15
       ! CLAW definition
16
17
       ! Define one dimension that will be added to the variables defined in the
18
       ! data clause.
       ! Apply the parallelization transformation on this subroutine.
19
20
21
       !$claw define dimension proma(1:nproma) &
22
       !$claw parallelize
23
24
       c = 5.345
25
       DO k = 2, nz
26
         t(k) = c * k
         d = t(k) * *2
27
```

Only the code of the module is shown as it is the only one which changes.

CLAW generated code for GPU target with OpenACC accelerator directive language.

```
MODULE mo column
1
2
3
   CONTAINS
4
    SUBROUTINE compute_column ( nz , q , t , nproma )
5
     INTEGER , INTENT(IN) :: nz
6
     REAL , INTENT(INOUT) :: t ( : , : )
7
     REAL , INTENT(INOUT) :: q ( : , : )
     \textbf{INTEGER} \ :: \ k
8
9
     REAL :: C
10
     REAL :: d
11
     INTEGER , INTENT(IN) :: nproma
12
     INTEGER :: proma
13
   !$acc data present(t,q,nproma,nz)
14
15
   !$acc parallel private(k,proma,d,c)
16
   !$acc loop
17
     DO proma = 1 , nproma , 1
18
      c = 5.345
19
   !$acc loop seq
20
      DO k = 2 , nz , 1
21
       t (proma, k) = c * k
22
       d = t (proma, k) ** (2)
23
       q (proma, k) = q (proma, k-1) + t (proma, k) * c
24
      END DO
25
      q (proma, nz) = q (proma, nz) * c
26
     END DO
27
    !$acc end parallel
28
   !$acc end data
29
    END SUBROUTINE compute_column
30
31
   END MODULE mo_column
```

CLAW generated code for CPU target with OpenMP accelerator directive language.

```
MODULE mo_column

CONTAINS
SUBROUTINE compute_column ( nz , q , t , nproma )
```

```
5
     INTEGER , INTENT(IN) :: nz
6
     REAL , INTENT(INOUT) :: t ( : , : )
7
     REAL , INTENT (INOUT) :: q ( : , : )
8
     INTEGER :: k
9
     REAL :: C
     REAL :: d ( 1 : nproma )
10
     INTEGER , INTENT(IN) :: nproma
11
12
     INTEGER :: proma
13
   !$omp parallel
14
15
     c = 5.345
     DO k = 2 , nz , 1
16
17
   !$omp do
18
      DO proma = 1 , nproma , 1
19
      t (proma, k) = c * k
20
      END DO
21
   !$omp end do
22
   !$omp do
23
      DO proma = 1 , nproma , 1
24
       d (proma) = t (proma, k) ** (2)
25
      END DO
26
    !$omp end do
27
   !$omp do
28
      DO proma = 1 , nproma , 1
29
       q (proma, k) = q (proma, k-1) + t (proma, k) * c
30
      END DO
31
   !$omp end do
32
    END DO
33
   !$omp do
34
     DO proma = 1 , nproma , 1
35
      q (proma , nz) = q (proma , nz) * c
36
    END DO
37
   !$omp end do
38
   !$omp end parallel
39
    END SUBROUTINE compute_column
40
   END MODULE mo_column
```

Compilation with the CLAW FORTRAN Compiler

```
# For GPU target with OpenACC directive
clawfc --target=gpu --directive=openacc -o mo_column.gpu_acc.f90 mo_column.f90
clawfc -o main.claw.f90 main.f90

# For CPU target with OpenMP directive
clawfc --target=gpu --directive=openmp -o mo_column.cpu_omp.f90 mo_column.f90
clawfc -o main.claw.f90 main.f90
```

3 CLAW low-level directive transformations

3.1 Global clauses

The following clauses can be applied to any kind of low-level transformations.

3.1.1 Target clause

The target clause specify for which transformed target the transformation is active. If no target is specified, the transformation is active for any target.

```
1 | !$claw directive [target(specific_target)]
```

The specific_target is one of the value provided by the compiler.

3.2 Loop transformation

3.2.1 Loop interchange/reordering

```
1 [!$claw loop-interchange [(induction_var[, induction_var] ...)]
```

Loop reordering is a common transformation applied on do statements when adding parallelization. This transformation is mainly used to improve the data locality and avoid the need to transpose the arrays.

The **loop-interchange** directive must be placed just before the first do statement composing the nested do statements group. With a group of two nested do statements, the induction variable and iteration range of the first statement is swap with the information of the second statement (see example 3.2.1).

If the list of induction variables is specified, the do statements are reordered according to the order defined in the list (see example A.3).

Options and details

1. *induction_var*: induction variable of the do statement. The list gives the new order of the do statement after the transformation. For group of 2 nested do statements, this information is optional.

Behavior with other directives

When the do statements to be reordered are decorated with additional directives, those directives

stay in place during the code transformation. In other words, they are not reordered together with the do statements iteration ranges (see example A.4).

Limitations

Currently, the **loop-interchange** directive is limited to 3 level of nested do statements. More nested levels can be declared but the transformation is limited to the first 3 levels from the directive declaration.

Code example

Example with 2 levels of loop.

Original code

```
1 !$claw loop-interchange
2 DO i=1, iend
3 DO k=1, kend
4 ! loop body here
5 END DO
6 END DO
```

Transformed code

```
! CLAW transformation (loop-interchange i < -- > k)

DO k=1, kend

DO i=1, iend

! loop body here

END DO

END DO
```

More code examples in the appendix. Example with 3 nested statements (see example A.3). Example with others directives (see example A.4).

3.2.2 Loop jamming/fusion

```
1 [!$claw loop-fusion [group(group_id)] [collapse(n)]
```

Loop jamming or fusion is used to merge 2 or more do statements together. Sometimes, the work performed in a single do statement is too small to create significant impact on performance when it is parallelized. Merging some do statements together create bigger blocks (kernels) to be parallelized.

If the *group* clause is not specified, all the do statements decorated with the directive in the same structured block will be merged together as a single do statement (see code example in this section).

If the *group* clause is specified, the loops are merged in-order with do statements sharing the same group in the current structured block (see example A.5).

The *collapse* clause is used to specify how many tightly nested do statements are associated with the **loop-fusion** construct (see example A.7). The argument to the *collapse* clause is a constant positive integer. If the *collapse* clause is not specified, only the do statement that follows immediately the directive is associated with the **loop-fusion** construct.

Options and details

- group_id: An identifier that represents a group of do statement in a structured block (see example A.5).
- n: A constant positive integer.

Behavior with other directives

When the do statement to be merged are decorated with other directives, only the directives on the first do statement of the merge group are kept in the transformed code.

Limitations

All the do statement within a group must share the same iteration range. If the *collapse* clause is specified, the do statements must share the same iteration range at the corresponding depth (see example A.7).

Code example

Example without clauses.

Original code

```
DO k=1, iend
1
2
     !$claw loop-fusion
3
     DO i=1, iend
4
        ! loop #1 body here
5
     END DO
6
7
     !$claw loop-fusion
8
     DO i=1, iend
9
        ! loop #2 body here
10
     END DO
11
   END DO
```

```
DO k=1, iend
! CLAW transformation (loop-fusion same block group)

DO i=1, iend
! loop #1 body here
! loop #2 body here
END DO
END DO
```

More code examples in the appendix. Example with the group clause (see example A.5). Example with others directives (see example A.6). Example with the collapse clause (see example A.7).

3.2.3 Loop extraction

```
!$claw loop-extract range(range_info) [[map(var[,var]...:mapping[,mapping]...) [map(var[,var]...:mapping[,mapping]...)] ...] [fusion [group(group_id)]] [ parallel] [acc(directives)]
```

Loop extraction can be performed on a subroutine/function call. The do statement corresponding to the defined iteration range is extracted from the subroutine/function and is wrapped around the subroutine/function call. In the transformation, a copy of the subroutine/function is created with the corresponding transformation (demotion) for the parameters.

Options and details

- 1. range: Correspond to the iteration range of the loop to be extracted. Notation i = istart, iend, istep. Step value is optional.
- 2. map: Define the mapping of variable that are demoted during the loop extraction. As seen in the example 3.2.3, the two parameters (1 dimensional array) are mapped to a scalar with the induction variable *i*.
 - (a) The var can be defined as two parts variable (e.g. a/a1). The first part is the function call part and refers to the variable as it is defined in the function call. The second part is the function definition part and refers to the name of the variable to be mapped as it defined in the function declaration. If a var is defined as a single part variable, the same name is used for both the function call and function definition part.
 - (b) The mapping can be defined as two parts variable (e.g. i/il). The first part is the function call part and refers to the mapping variable as it is defined in the function call. The second part is the function definition part and refers to the name of the mapping variable as it defined in the function declaration. If a mapping is defined as a single part mapping variable, the same name is used for both the function call and function definition part.
- 3. fusion: Allow the extracted loop to be merged with other loops. Options are identical with the loop-fusion directive.
- 4. parallel: Wrap the extracted loop in a parallel region.
- 5. acc: Add the accelerator directives to the extracted loop.

If the directive **loop-extract** is used for more than one call to the same subroutine, the extraction can generate 1 to N dedicated subroutines.

Behavior with other directives

Accelerator directives are generated by the transformation. If the function call was decorated with accelerator directives prior to the transformation, those directives stay in place.

Code example

Example without additional clauses.

Original code

```
PROGRAM main
1
      !$claw loop-extract range(i=istart,iend) map(value1,value2:i)
     CALL xyz (value1, value2)
3
   END PROGRAM main
4
5
6
   SUBROUTINE xyz (value1, value2)
7
     REAL, INTENT (IN) :: value2(x:y), value2(x:y)
8
9
     DO i = istart, iend
        ! some computation with value1(i) and value2(i) here
10
     END DO
11
12
   END SUBROUTINE xyz
```

Transformed code

```
PROGRAM main
 1
 2
      !CLAW extracted loop
3
     DO i = istart, iend
 4
       CALL xyz_claw(value1(i), value2(i))
 5
     END DO
 6
   END PROGRAM main
8
   SUBROUTINE xyz (value1, value2)
9
     REAL, INTENT (IN) :: value2(x:y), value2(x:y)
10
11
     DO i = istart, iend
12
       ! some computation with value1(i) and value2(i) here
13
     END DO
   END SUBROUTINE xyz
14
15
16
   !CLAW extracted loop new subroutine
   SUBROUTINE xyz_claw(value1, value2)
17
18
     REAL, INTENT (IN) :: value1, value2
19
      ! some computation with value1 and value2 here
20
   END SUBROUTINE xyz_claw
```

More code examples in the appendix. Example with the fusion clause (see example A.8). Example with a more complicated mapping (see example A.9).

3.2.4 Loop hoisting

The **loop-hoist** directive allows nested loops in a defined structured block to be merged together and to hoist the beginning of those nested loop just after the directive declaration. Loops with different lower-bound indexes can also be merged with the addition of an IF statement. This feature works only when the lower-bound are integer constant.

Options and details

- 1. induction_var: List of induction variables of the do statements to be hoisted.
- 2. *interchange*: Allow the group of hoisted loops to be reordered. Options are identical with the **loop-interchange** directive.
- 3. reshape: Reshape arrays to scalar or to array with fewer dimensions. The original declaration is replaced with the corresponding demoted declaration in the current block (function/module).
 - (a) array_name: Name of the array type to be reshapped.
 - (b) target_dimension: Number of dimension after reshapaing. 0 indicates that the type is reshapped to a scalar.
 - (c) kept_dimension: comma separated list of integer that indicates which dimensions are kept when the array is not totally demoted to scalar. Dimension index starts at 1.
- 4. fusion: Allow the extracted loop to be merged with other loops. Options are identical with the loop-fusion directive.

Code example

Example without clauses.

```
!$acc parallel loop gang vector collapse(2)

DO jt=1, jtend
   !$claw loop-hoist(j,i) interchange

IF ( .TRUE. ) CYCLE
   ! outside loop statement

END IF

DO j=1, jend
```

```
8
        DO i=1, iend
9
          ! first nested loop body
10
11
      END DO
12
      DO j=2, jend
13
        DO i=1, iend
          ! second nested loop body
14
15
        END DO
16
      END DO
17
      !$claw end loop-hoist
18
    END DO
19
    !$acc end parallel
```

```
1
    !$acc parallel loop gang vector collapse(2)
2
    DO jt=1, jtend
3
      DO i=1, iend
        DO j=1, jend
4
5
          IF ( .TRUE. ) CYCLE
6
             ! outside loop statement
7
          END IF
8
           ! first nested loop body
9
          IF (\dot{\gamma} > 1) THEN
10
             ! second nested loop body
11
          END IF
12
        END DO
13
      END DO
   END DO
14
15
    !$acc end parallel
```

3.3 Specific transformation

3.3.1 K caching (column caching)

```
1 | !$claw kcache data(var_1[,var_2] ...) [offset(offset_1[,offset_2] ...)]] [init] [private]
```

In memory-bound problem, it might be useful to cache array values used several times during the current do statement iteration or for the next one.

The **kcache** directive is applied in a limited do statement block. It will cache the corresponding assigned value and update the array index in the block according to the given plus/minus offsets. If the array referenced in the data clause is assigned in the block, the caching will take its place (see example 3.3.1). Otherwise, an assignment statement is created to start the caching (see example A.1).

If the offset values are omitted, there are inferred from the dimension of the variable to be cached and set to 0.

Options and details

- data: The data clause specifies which array will be impacted by the column caching transformation. List of array identifiers.
- offset: Integer value separated by a comma that represents the offset at each dimension.
- *init*: If the *init* clause is specified, the cache variable will be initialized with the corresponding array value at the given offset during the 1st loop iteration.
- private: it declares that a copy of each item on the list will be created for each parallel gang on the accelerator. The list is the one specified on the data clause.

Code example

Where caching takes assignment place and init clause.

Original code

```
DO j1 = k1start, k1end
 1
 2
     array_u6(j1,k3end) = pc1(j1,kend) * p2(j1)
 3
    ENDDO
 4
5
   DO j3 = k3start + 1, k3end
 6
     DO j1 = k1start, k1end
 7
        ! More computation here with others variables
 8
9
        value1 = 1.0_{dp} / (1.0_{dp} - p2(j1) * (pc2(j1,j3) * array6(j1,j3-1))
10
                                           + pc2(j1, j3) * array8(j1, j3-1))
        value2 = p2(j1) * (pc2(j1, j3) * array6(j1, j3-1)
11
12
                           + pc2(j1, j3) * array8(j1, j3-1))
13
        value6 = p1(j1) * (pc2(j1,j3) * array6(j1,j3-1)
                                                                                 &
14
                            +pc2(j1, j3) * array8(j1, j3-1))
15
            !$claw kcache data(array6) offset(0 -1) init
16
        array6(j1, j3) = p2(j1) * pc1(j1, j3)
                         + value2 * array8(j1, j3) + value6 * array8(j1, j3)
17
18
            !$claw kcache data(array6, array8) offset(0 -1) init
19
        array8(j1, j3) = p2(j1) * pc1(j1, j3)
20
                         + value6 * array2(j1, j3) + value7 * array4(j1, j3)
21
     END DO
22
    END DO
```

```
DO j1 = k1start, k1end
array_u6(j1,k3end) = pc1(j1,kend) * p2(j1)
ENDDO
```

```
4
5
    DO j1 = k1start, k1end
6
     DO j3 = k3start + 1, k3end
 7
        IF (j3 = k3start + 1) THEN
8
          array6_k_m1 = array6(j1, j3-1)
9
          array8_k_m1 = array8(j1, j3-1)
10
        END IF
11
        ! More computation here with others variables
12
        value1 = 1.0_dp / (1.0_dp - p2(j1) * (pc2(j1,j3) * array6_k_m1)
13
                                                                              δ
14
                                           + pc2(j1,j3) * array8_k_m1))
        value2 = p2(j1) * (pc2(j1,j3) * array6_k_m1
15
                                                                              δ
                           + pc2(j1,j3) * array8_k_m1
16
17
        value6 = p1(j1) * (pc2(j1,j3) * array6_k_m1
                                                                              &
18
                            +pc2(j1, j3) * array8_k_m1
19
            array6_k_m1 = p2(j1) * pc1(j1, j3)
                        + value2 * array8(j1, j3) + value6 * array8(j1, j3)
20
21
        array6(j1,j3) = array6_k_m1
22
        array8_k_m1 = p2(j1) * pc1(j1, j3)
                                                                              &
23
                         + value6 * array2(j1, j3) + value7 * array4(j1, j3)
24
        array8(j1,j3) = array8_k_m1
25
     END DO
26
    END DO
```

More code examples in the appendix. Example with assignment created during the transformation (see example A.1). Example with init and private clause (see example A.2).

3.3.2 On-the-fly computation (array access to function call)

```
1 | !$claw call array_name=function_call(arg_list)
```

Sometimes, replacing access to pre-computed arrays with computation on-the-fly can increase the performance. It can reduce the memory access for memory-bound kernel and exploit some unused resources to execute the computation.

This transformation is local to the current do statement. Array access from the directive will be replaced.

Options and details

- array_name: Array identifier. References to this array will be replaced by the function call in the current structured block.
- function_call: Name of the function provided by the user to compute the value. Return value must be of the same type as accessed value.
- arg_list: Comma separated list of arguments to be passed to the function call.

3.4 Accelerator abstraction/helper transformation

3.4.1 Conditional primitive directives

```
! $claw primitive_prefix primitive_directives

! Example for OpenACC
! $claw acc routine seq

! Example for OpenMP
! $claw omp target enter data
```

Primitive compiler directive like OpenACC or OpenMP can be enabled conditionally during the transformation. Examples on line 4 and 7 are conditional directives. They will be transformed into standard OpenACC respectively OpenMP directives only if the primitive directive language is specified for the transformation.

3.4.2 Array notation to do statements

Computations using the array notation are not suitable to be parallelized with language like OpenACC. The **array-transform** directive allows to transform those notation with the corresponding do statements which are more suitable for parallelization.

The goal of this directive is to pass from an array notation assignment like this:

To a do statement statement like this:

```
1 DO i=1, n
2 A(i) = A(i+m) + B(i) * C(n+i)
END DO
```

If the directive is used as a block directive, the assignments are wrapped in a single do statement if their induction range match.

```
1 !$claw array-transform
```

To

```
1 DO i=1, n
2 A(i) = A(i+m) + B(i) * C(n+i)
3 B(i) = B(i) * 0.5
END DO
```

Options and details

- 1. induction: Allow to name the induction variable created for the do statement.
- 2. fusion: Allow the extracted loop to be merged with other loops. Options are identical with the **loop-fusion** directive
- 3. parallel: Wrap the extracted loop in a parallel region.
- 4. acc: Define accelerator clauses that will be applied to the generated loops.

Behavior with other directives

Directives declared before the array-transform directive will be kept in the generated code.

Code example

Example with array notation following the pragma statement.

Original code

```
SUBROUTINE vector_add
INTEGER :: i = 10
INTEGER, DIMENSION(0:9) :: vec1

!$claw array-transform
vec1(0:i) = vec1(0:i) + 10;
END SUBROUTINE vector_add
```

```
SUBROUTINE vector_add

INTEGER :: claw_i

INTEGER :: i = 10

INTEGER, DIMENSION(0:9) :: vec1

! CLAW transformation array notation to do loop

DO claw_i = 0, i

vec1(claw_i) = vec1(claw_i) + 10;
```

```
9 END DO 10 END SUBROUTINE vector_add
```

More code examples in the appendix. Example with the induction and acc (see example A.10). Example with fusion clause (see example A.11). Example with 2-dimensional arrays (see example A.12).

3.5 Utility transformation

3.5.1 Remove

The **remove** directive allows the user to remove section of code during the transformation process. This code section is lost in the transformed code.

Options and details

If the directive is directly followed by a structured block (IF or DO), the end directive is not mandatory (see example 3.5.1). In any other cases, the end directive is mandatory.

Code example

Original code

```
1
    DO k=1, kend
2
      DO i=1, iend
        ! loop #1 body here
3
 4
      END DO
 5
 6
      !$claw remove
 7
      IF (k > 1) THEN
 8
        PRINT*, k
9
      END IF
10
11
      DO i=1, iend
12
        ! loop #2 body here
13
      END DO
   END DO
14
```

```
1 DO k=1, kend
```

```
DO i=1, iend
   ! loop #1 body here

END DO

DO i=1, iend
   ! loop #2 body here

END DO

END DO

END DO
```

More code examples in the appendix. Example with block remove (see example A.13).

3.5.2 Ignore

```
! $claw ignore

! code block

! sclaw end ignore
```

The ignore directive allows the user to ignore section of code for the transformation process. This code section is retrieved in the transformed code. Any CLAW directives in an ignored code section is also ignored.

Code example

In the example below, the first remove block is ignored and the second is not.

```
PROGRAM testignore
 2
3
      !$claw ignore
      !$claw remove
 4
5
     PRINT*,'These lines'
6
     PRINT*,'are ignored'
 7
     PRINT*, 'by the CLAW compiler'
8
     PRINT*, 'but kept in the final transformed code'
9
     PRINT*,' with the remove directives.'
10
      !$claw end remove
      !$claw end ignore
11
12
13
      !$claw remove
14
      PRINT*,'These lines'
15
     PRINT*, 'are not ignored.'
16
      !$claw end remove
17
   END PROGRAM testignore
```

```
1
   PROGRAM testignore
2
3
      !$claw remove
4
     PRINT*,'These lines'
5
     PRINT*,'are ignored'
6
     PRINT*,'by the CLAW compiler'
7
     \mathtt{PRINT}\star, but kept in the final transformed code'
8
     PRINT*,'with the remove directives.'
9
      !$claw end remove
10
   END PROGRAM testignore
```

3.5.3 Verbatim

The verbatim directive allows the user to insert code after the transformation process. This code is not analyzed by the compiler.

Code example

In the example below, the first remove block is ignored and the second is not.

Original code

```
1
   PROGRAM testverbatim
2
3
     !$claw verbatim IF (.FALSE.) THEN
4
     PRINT*,'These lines'
5
     PRINT*,'are not printed'
6
     PRINT*,'if the the CLAW compiler has processed'
     PRINT*,'the file.'
7
8
     !$claw verbatim END IF
9
10
   END PROGRAM testverbatim
```

```
1 PROGRAM testverbatim
2 
3 IF (.FALSE.) THEN
4 PRINT*,'These lines'
```

```
PRINT*, 'are not printed'
PRINT*,'if the the CLAW compiler has processed'
PRINT*,'the file.'
END IF

END PROGRAM testverbatim
```

A Code examples

A.1 kcache 2

Assignment is created during the transformation.

Original code

```
DO j3 = ki3sc+1, ki3ec
!$claw kcache data(array2)
var1 = x * y - array2(j1, j3)
var2 = z * array2(j1, j3)
END DO
```

Transformed code

```
DO j3 = ki3sc+1, ki3ec
    array2_k = array2(j1, j3)
    var1 = x * y - array2_k
    var2 = z * array2_k
    END DO
```

A.2 kcache 3

Using init and private clause

```
9 END DO 10 !$acc end parallel
```

```
!$acc parallel private(array1_k_m1)
   DO j3 = ki3sc+1, ki3ec
     IF (j3 == ki3sc+1) THEN
3
4
       array1_k_m1 = array1(j1, j3-1)
5
     END IF
6
     var3 = array1_k_m1
7
     array1_k_m1 = x * y * z
8
     array1(j1,ki3sc) = array1_k_m1
9
     var1 = x * y - array1_k_m1
10
     var2 = z * array1_k_m1
11
   END DO
   !$acc end parallel
```

A.3 loop-interchange 2

Example with 3 levels of loop.

Original code

```
!$claw loop-interchange (k,i,j)
DO i=1, iend ! loop at depth 0
DO j=1, jend ! loop at depth 1
DO k=1, kend ! loop at depth 2
! loop body here
END DO
END DO
END DO
END DO
```

```
1
  ! CLAW transformation (loop-interchange (k,i,j))
2
  DO k=1, kend! loop at depth 2
3
    DO i=1, iend
                   ! loop at depth 0
4
      DO j=1, jend ! loop at depth 1
5
       ! loop body here
6
      END DO
7
    END DO
  END DO
```

A.4 loop-interchange 3

Example with OpenACC directives.

Original code

```
!$acc parallel
1
  !$acc loop gang
  !$claw loop-interchange
  DO i=1, iend
    !$acc loop vector
5
6
    DO k=1, kend
7
      ! loop body here
8
    END DO
9
  END DO
  !$acc end parallel
```

Transformed code

```
1
  ! CLAW transformation (loop-interchange i < -- > k)
  !$acc parallel
  !$acc loop gang
  DO k=1, kend
    !$acc loop vector
5
6
    DO i=1, iend
7
      ! loop body here
8
    END DO
9
  END DO
  !$acc end parallel
```

A.5 loop-fusion 2

Example using the *group* clause.

```
1
   DO k=1, iend
     !$claw loop-fusion group(g1)
3
     DO i=1, iend
4
       ! loop #1 body here
5
     END DO
6
7
     !$claw loop-fusion group(g1)
8
     DO i=1, iend
9
       ! loop #2 body here
10
     END DO
11
```

```
12
     !$claw loop-fusion group(g2)
13
     DO i=1, jend
14
       ! loop #3 body here
15
     END DO
16
17
     !$claw loop-fusion group(g2)
18
     DO i=1, jend
19
       ! loop #4 body here
20
     END DO
   END DO
```

```
DO k=1, iend
1
2
     ! CLAW transformation (loop-fusion group g1)
3
     DO i=1, iend
4
       ! loop #1 body here
5
       ! loop #2 body here
6
     END DO
7
8
     ! CLAW tranformation (loop-fusion group g2)
9
     DO i=1, jend
10
       ! loop #3 body here
       ! loop #4 body here
11
12
     END DO
   END DO
13
```

A.6 loop-fusion 3

Example showing behavior with other directives (OpenACC).

```
!$acc parallel
   !$acc loop gang
   DO k=1, iend
4
     !$acc loop seq
5
     !$claw loop-fusion
6
     DO i=1, iend
7
       ! loop #1 body here
8
     END DO
9
10
     !$acc loop vector
     !$claw loop-fusion
11
     DO i=1, iend
12
13
       ! loop #2 body here
     END DO
15 END DO
```

```
16 | !$acc end parallel
```

```
!$acc parallel
   !$acc loop gang
   DO k=1, iend
     ! CLAW transformation (loop-fusion same block group)
5
     !$acc loop seq
6
     DO i=1, iend
7
       ! loop #1 body here
8
       ! loop #2 body here
9
     END DO
10
   END DO
   !$acc end parallel
```

A.7 loop-fusion 4

Example using the *collapse* clause.

Original code

```
1
   DO k=1, iend
2
     !$claw loop-fusion collapse(2)
3
     DO i=0, iend
       DO j=0, jend
 4
5
         ! nested loop #1 body here
6
       END FO
7
     END DO
8
9
      !$claw loop-fusion collapse(2)
10
     DO i=0, iend
       DO j=0, jend
11
12
         ! loop #2 body here
13
       END DO
     END DO
14
   END DO
15
```

```
DO k=1, iend
! CLAW transformation (loop-fusion collapse(2))
DO i=0, iend
DO j=0, jend
! nested loop #1 body here
! nested loop #2 body here
END DO
```

```
8 END DO END DO
```

A.8 loop-extract 2

Example with fusion clause.

Original code

```
PROGRAM main
     !$claw loop-extract(i=istart,iend) map(value1,value2:i) fusion group(g1)
3
     CALL xyz (value1, value2)
4
5
     !$claw loop-fusion group(g1)
6
     DO i = istart, iend
 7
       ! some computation here
8
       print*,'Inside loop', i
9
     END DO
10
   END PROGRAM main
11
   SUBROUTINE xyz (value1, value2)
12
13
     REAL, INTENT (IN) :: value2(x:y), value2(x:y)
14
15
     DO i = istart, iend
       ! some computation with value1(i) and value2(i) here
16
17
     END DO
   END SUBROUTINE xyz
```

```
1
   PROGRAM main
2
     !CLAW extracted loop
     DO i = istart, iend
3
 4
       CALL xyz_claw(value1(i), value2(i))
 5
       ! some computation here
       print*,'Inside loop', i
6
     END DO
7
8
   END PROGRAM main
9
10
   SUBROUTINE xyz(value1, value2)
11
     REAL, INTENT (IN) :: value2(x:y), value2(x:y)
12
13
     DO i = istart, iend
14
       ! some computation with value1(i) and value2(i) here
15
     END DO
16
   END SUBROUTINE xyz
17
18 !CLAW extracted loop new subroutine
```

```
19 SUBROUTINE xyz_claw(value1, value2)
20 REAL, INTENT (IN) :: value1, value2
21 ! some computation with value1 and value2 here
22 END SUBROUTINE xyz_claw
```

A.9 loop-extract 3

Example with more complicated mapping.

Original code

```
1
   PROGRAM main
2
     !$claw loop-extract(i=istart,iend) map(value1,value2:i/j)
3
     CALL xyz (value1, value2)
   END PROGRAM main
4
5
   SUBROUTINE xyz (value1, value2, j)
7
     INTGER, INTENT(IN) :: j
8
     REAL , INTENT(IN) :: value2(x:y), value2(x:y)
9
10
     DO i = istart, iend
       ! some computation with value1(j) and value2(j) here
11
12
     END DO
   END SUBROUTINE xyz
```

```
1
   PROGRAM main
2
     !CLAW extracted loop
3
     DO i = istart, iend
4
       CALL xyz_claw(value1(i), value2(i))
5
     END DO
6
   END PROGRAM main
7
8
   SUBROUTINE xyz (value1, value2, j)
9
     INTGER, INTENT(IN) :: j
10
     REAL , INTENT(IN) :: value2(x:y), value2(x:y)
11
     DO i = istart, iend
12
13
       ! some computation with value1(j) and value2(j) here
14
     END DO
15
   END SUBROUTINE xyz
16
17
   !CLAW extracted loop new subroutine
18
   SUBROUTINE xyz_claw(value1, value2, j)
19
     INTGER, INTENT(IN) :: j
20
     REAL, INTENT (IN) :: value1, value2
    ! some computation with value1 and value2 here
```

A.10 array-transform 2

Example with induction and acc clauses

Original code

```
SUBROUTINE vector_add
1
2
     INTEGER :: i = 10
3
     INTEGER, DIMENSION(0:9) :: vec1
4
5
     !$acc parallel
6
     !$claw array-transform induction(myinduc) acc(loop)
7
     vec1(0:i) = vec1(0:i) + 10;
8
9
     !$claw array-transform acc(loop)
10
     vec1(0:i) = vec1(0:i) + 1;
11
     !$acc end parallel
   END SUBROUTINE vector_add
```

```
1
   SUBROUTINE vector_add
2
     INTEGER :: claw_i
3
     INTEGER :: myinduc
     INTEGER :: i = 10
 4
 5
     INTEGER, DIMENSION(0:9) :: vec1
 6
 7
      !$acc parallel
 8
 9
      ! CLAW transformation array notation vec1(0:i) to do loop
10
      !$acc loop
     DO myinduc = 0, i
11
12
       vec1(myinduc) = vec1(myinduc) + 10;
13
14
15
      ! CLAW transformation array notation vec1(0:i) to do loop
      !$acc loop
16
     DO claw_i = 0, i
17
18
       vec1(claw_i) = vec1(claw_i) + 1;
19
     END DO
20
21
      !$acc end parallel
   END SUBROUTINE vector_add
```

A.11 array-transform 3

Example with fusion clause

Original code

```
1
   SUBROUTINE vector_add
2
     INTEGER :: i = 10
     INTEGER, DIMENSION(0:9) :: vec1
3
     INTEGER, DIMENSION(0:9) :: vec2
4
5
6
     !$claw array-transform induction(claw_i) fusion
7
     vec1(0:i) = vec1(0:i) + 10;
8
9
     !$claw array-transform induction(claw_i) fusion
10
     vec2(0:i) = vec2(0:i) + 1;
   END SUBROUTINE vector_add
```

Transformed code

```
1
   SUBROUTINE vector_add
     INTEGER :: claw_i
3
     INTEGER :: i = 10
4
     INTEGER, DIMENSION(0:9) :: vec1
     INTEGER, DIMENSION(0:9) :: vec2
5
6
7
     ! CLAW transformation array notation vec1(0:i) to do loop
8
     ! CLAW transformation array notation vec2(0:i) to do loop
9
     ! CLAW transformation fusion
10
     DO claw_i=0, i
11
       vec1(claw_i) = vec1(claw_i) + 10;
12
       vec2(claw_i) = vec2(claw_i) + 1;
13
     END DO
14
   END SUBROUTINE vector_add
```

A.12 array-transform 4

Example with 2-dimensional arrays

```
SUBROUTINE vector_add
INTEGER :: i = 10
INTEGER, DIMENSION(0:10,0:10) :: vec1
INTEGER, DIMENSION(0:10,0:10) :: vec2

vec1(0:i,0:i) = 0;
```

```
7  vec2(0:i,0:i) = 100;
8 
9  !$claw array-transform
10  vec1(0:i,0:i) = vec2(0:i,0:i) + 10
11  END SUBROUTINE vector_add
```

```
SUBROUTINE vector_add
   INTEGER :: i = 10
3
   INTEGER, DIMENSION(0:10,0:10) :: vec1
4
   INTEGER, DIMENSION(0:10,0:10) :: vec2
5
6
   vec1(0:i,0:i) = 0;
7
   vec2(0:i,0:i) = 100;
8
9
   DO claw_i = 0, i, 1
10
    DO claw_j = 0, i, 1
       vec1(claw_i,claw_j) = vec2(claw_i, claw_j) + 10
11
12
    END DO
13
   END DO
   END SUBROUTINE vector_add
```

A.13 remove 2

Original code

```
1
   DO k=1, kend
2
     DO i=1, iend
3
      ! loop #1 body here
 4
     END DO
5
6
     !$claw remove
7
     PRINT*, k
8
     PRINT*, k+1
9
     !$claw end remove
10
11
     DO i=1, iend
12
       ! loop #2 body here
     END DO
13
14
   END DO
```

```
1 DO k=1, kend
2 DO i=1, iend
3 ! loop #1 body here
```

```
4 END DO
5
6 DO i=1, iend
7 ! loop #2 body here
8 END DO
9 END DO
```