C Language Concepts

CS 231
System Programming Topics

Process Control

- Each process has a unique process ID
- getpid() returns ID of process
- getppid() returns ID of parent process
- petuid() returns real user ID of process (getgid for group)
- geteuid() returns effective process ID (getegid for group)

Fork Function

- Create a process by calling fork()
 - fork is called once, returns twice
 - once in the calling process
 - once in the created process

Parent Process

- A process calls fork() to create a child process
- ➡ If parent process dies before child, child is orphan and is adopted by init (process id 1)
- fork() returns value of o to child, child PID to parent
- Look at code forkPID.c

Child and Parent are Nearly Identical

- Child process inherits
 - real & effective user & group IDs
 - current working directory
 - root directory
 - open files

Child Parent Differences

- return value from fork()
- process ID
- parent process ID
- pending alarms cleared for child
- child pending signals is empty set

Uses for fork()

- ◆ Common for server program to wait for request, then fork() child to serve request, parent continues waiting for requests
- Process wants to execute different program: child calls exec to replace its code with new program.

wait and waitpid functions

- A parent waits on its child to get exit status
- If child dies and parent does not wait, child becomes a zombie process
- wait() blocks parent until a child dies
- waitpid() can wait for a specific child, or not block parent at all

Wait Functions API

```
#include <sys/types.h>
    #include <sys/wait.h>

pid_t wait(int *wstatus);

pid_t waitpid(pid_t pid, int *wstatus, int options);

int waitid(idtype_t idtype, id_t id, siginfo_t *infop, int options);

example code in parentWait.c
```

exec system calls

- int execl(const char * path, const char * argo,...,const char* argn, char * /*NULL*/)
- int execv(const char * path, const char * argv[])
- int execle(const char * path, const char * argo,...,const char * argn, char * /*NULL*/, char * const envp[])

exec calls continued

- int execve(const char * path, const char * argv[], char *
 const envp[])
- int execlp(const char * file, const char * argo,...,const char * argn, char * /*NULL*/)
- int execup(const char * file, const char * argv[])

exec details

- names are execl, execv, execle, execve, execlp, execvp
- l stands for list arguments are character strings
- v stands for vector, array of pointers to arguments
- e stands for environment, array of pointers to environment strings
- ▶ p first argument is a file name, if '/' in name, it is a path, otherwise,
 PATH environment variable lists paths to search for file name
- see forkexecv.c and env.c

IPC - Pipe

- * kernel data structure for interprocess communication
- information flows in one direction only
- processes must share a common ancestor
- there is a file descriptor for each pipe end

pipe function

- int pipe(int fileDes[2])
- fd[o] open for read
- fd[1] open for write
- call pipe in a process, then fork one or more children to establish communication
- close read end in writing process, and write end in reading process

file descriptor

- The operating system keeps track of all open files
- Everything the operating system can affect (in UNIX and variants) is a file
- Essentially, there is an array of open file information
- file descriptor is an index into array
- file descriptors may be duplicated with functions dup and dup2, see example code: pipe.c

Signals

- A signal is a message sent by OS to process
- The default response to most signals is for process to die
- The process can ignore (most) signals, or handle them with code
- See the code in signal1.c to signal4.c, illustrates use of signal function
- ◆ A more powerful function (gives programmer more control) is sigaction it is also much more complicated