

# A Metric on the Set of Elliptic Curves over $\mathbf{F}_p$

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## Abstract

Elliptic Curves over finite field have found application in many areas including cryptography. In the current article we define a metric on the set of elliptic curves defined over a prime field  $\mathbf{F}_p, p > 3$ .

**Keywords:** Elliptic Curves, Elliptic Curve Cryptosystems, Metric, Isomorphism Classes of Elliptic Curves.

## 1 Introduction

Elliptic curves are beautiful geometric entities which have fascinated mathematicians for more than a century. The curves have been studied at length and many of their interesting properties have been unearthed. In last two decades, the study of the curves received a new impetus when many of their applications were discovered. Particularly elliptic curve cryptosystems (ECC) (proposed jointly by Koblitz [2] and Miller [3] in 1985) built on the strength of elliptic curve discrete logarithm problem (ECDLP) integrated the study of the curves to the mainstream of cryptographic research. In the current article we propose a simple metric on the set of elliptic curves over a prime field  $\mathbf{F}_p, p > 3$ . For details about elliptic curve or elliptic curve cryptography the readers can refer to [1].

## 2 The Metric

The metric we propose is based on the concept of isomorphic classes of elliptic curves. Two curves on the same isomorphic class will have a finite distance between them. The distance of a curve from all the curves in an isomorphism class different than its own will be defined to be infinity.

**Elliptic Curves Over Prime Fields  $\mathbf{F}_p, p > 3$ :** an elliptic curve is represented by an equation of the form

$$C : y^2 = x^3 + ax + b$$

where  $a, b \in \mathbf{F}_p$  and  $4a^3 + 27b^2 \neq 0$ . The set of rational points over  $\mathbf{F}_p$  are the set of all points over  $\mathbf{F}_p \times \mathbf{F}_p$  which satisfy this equation together with a special point, called the point at infinity.

Isomorphism on the set of elliptic curves over  $\mathbf{F}_p$  is an equivalence relation defined as follows.

**Isomorphic Curves:** Let

$$C_i : y^2 = x^3 + a_i x + b_i, i = 1, 2$$

be two curves over  $\mathbf{F}_p, p > 3$ .  $C_1$  is said to be isomorphic to  $C_2$  if there exists a  $t \in \mathbf{F}_p$  such that  $a_2 = t^4 a_1$  and  $b_2 = t^6 b_1$ .

Let  $g$  be a generator of the field  $\mathbf{F}_p$ . Then, given any non-zero element  $z \in \mathbf{F}_p$  there exists an integer  $k \in \{0, 1, \dots, p-2\}$  such that  $z = g^k$ . We will refer to the set  $\{0, 1, \dots, p-2\}$  as an index set of  $g$ . Note that the index set of  $g$  is not unique. Any residue class of  $p-1$  can act as an index set. For defining the metric we will always use the index set  $\{-\frac{p-1}{2}+1, -\frac{p-1}{2}+2, \dots, -1, 0, 1, \dots, \frac{p-1}{2}\}$ . We will refer to this index set of a generator  $g$  as the *standard* index set of  $g$ .

Let  $C_1$  and  $C_2$  be any two curves over the  $\mathbf{F}_p$ . If  $C_1$  and  $C_2$  are not isomorphic we define the distance between them to be infinite. Otherwise let  $t \in \mathbf{F}_p$  be the field element which transforms the parameter of  $C_1$  to those of  $C_2$  (or parameters of  $C_2$  to those of  $C_1$ ) (see the definition of isomorphic curve). Let  $t = g^r$ , where  $r$  is in the standard index set of  $g$ . Then we define the distance between  $C_1$  and  $C_2$  to be  $|r|^1$ . That is

$$\begin{aligned} d_g(C_1, C_2) &= |r| \text{ if } C_1 \text{ and } C_2 \text{ are isomorphic and } t = g^r, \\ d_g(C_1, C_2) &= \infty \text{ otherwise.} \end{aligned}$$

Now we claim that  $d_g$  as defined above is a metric.

Clearly,  $d_g \geq 0$ . Also, if  $C_1$  and  $C_2$  are the same curve, then they are isomorphic and for them  $t = 1$  and  $r = 0$ . Hence it follows that  $d_g(C_1, C_2) = 0$  if  $C_1 = C_2$ . To prove the converse is equally simple.

Next we will show that  $d_g(C_1, C_2) = d_g(C_2, C_1)$ . If these curves are not isomorphic then there is nothing to prove as both of these distances are  $\infty$ . So let us assume that they are isomorphic. Let  $t = g^r$  be the element in  $\mathbf{F}_p$  which transforms parameters of  $C_1$  to those of  $C_2$  (i.e.  $a_2 = t^4 a_1, b_2 = t^6 b_1$ ). Then  $t^{-1} = g^{-r}$  transforms parameters of  $C_2$  to those of  $C_1$  (i.e.  $a_1 = t^{-1^4} a_2, b_1 = t^{-1^6} b_2$ ). Hence  $d_g(C_1, C_2) = |r|$  and  $d_g(C_2, C_1) = |-r|$ , which are the same.

Finally, we have to prove the *triangle inequality*, i.e. we have to show that for any three curves  $C_i, i = 1, 2, 3$ ,

$$d_g(C_1, C_2) + d_g(C_2, C_3) \geq d_g(C_1, C_3).$$

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<sup>1</sup>there may be several  $t$ 's which define the same isomorphism. Let  $t_1, \dots, t_l$  'define' the same isomorphism. Write  $t_i = g^{\alpha_i}; 1 \leq i \leq l$ . Choose that  $i$  for which  $\alpha_i$  is minimum.

Clearly, this is obvious if  $C_1$  is not isomorphic to  $C_2$  or  $C_2$  is not isomorphic to  $C_3$ . In that case both sides of the inequality are  $\infty$ .

So let us assume that  $C_1$  is isomorphic to  $C_2$  and  $C_2$  is isomorphic to  $C_3$ . As isomorphism is an equivalence relation  $C_1$  is also isomorphic to  $C_3$ . Let

$$C_i : y^2 = x^3 + a_i x + b_i$$

$i = 1, 2, 3$ . Then there exist  $t_1, t_2 \in F_p$  and indices  $r_1, r_2$  in the standard index set of  $g$  such that

$$a_2 = t_1^4 a_1, b_2 = t_1^6 b_1, t_1 = g^{r_1}$$

and

$$a_3 = t_2^4 a_2, b_3 = t_2^6 b_2, t_2 = g^{r_2}$$

Now

$$a_3 = (t_1 t_2)^4 a_1, b_3 = (t_1 t_2)^6 b_1$$

Let  $t_1 t_2 = t_3 = g^{r_3}$ . Then  $r_3 = r_1 + r_2 \pmod{(p-1)}$ . Hence  $r_3 \leq r_1 + r_2$ . We have now,

$$d_g(C_1, C_2) = r_1,$$

$$d_g(C_2, C_3) = r_2,$$

$$d_g(C_1, C_3) = r_3,$$

Hence

$$d_g(C_1, C_2) + d_g(C_2, C_3) \geq d_g(C_1, C_3).$$

This establishes the triangle inequality.

### 3 Conclusion

In this article, we have defined a metric on the set of elliptic curves over  $\mathbf{F}_p$ . The metric is dependent on the choice of the generator of the underlying field. A better metric will be the one which is independent over all generators. A candidate for such a metric can be  $d(C_1, C_2) = \Sigma_g d_g(C_1, C_2)$  or we can take the average over all the generator dependent distances. One interesting open question is: does there exist one generator whose metric agrees with the average metric? Or is there a special class of fields for which there exist a generator whose corresponding metric agrees with the average metric?

### References

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