



Authenticated Encryption

CBC paddings attacks

Recap

Authenticated encryption: CPA security + ciphertext integrity

- Confidentiality in presence of **active** adversary
- Prevents chosen-ciphertext attacks

Limitation: cannot help bad implementations ... (this segment)

Authenticated encryption modes:

- Standards: GCM, CCM, EAX
- General construction: encrypt-then-MAC

The TLS record protocol (CBC encryption)

Decryption: $\text{dec}(k_{b \rightarrow s}, \text{record}, \text{ctr}_{b \rightarrow s})$:

step 1: CBC decrypt record using k_{enc}

step 2: check pad format: abort if invalid

step 3: check tag on $[++\text{ctr}_{b \rightarrow s} \parallel \text{header} \parallel \text{data}]$
abort if invalid

Two types of error:

- padding error
- MAC error



Padding oracle

Suppose attacker can differentiate the two errors
(pad error, MAC error):

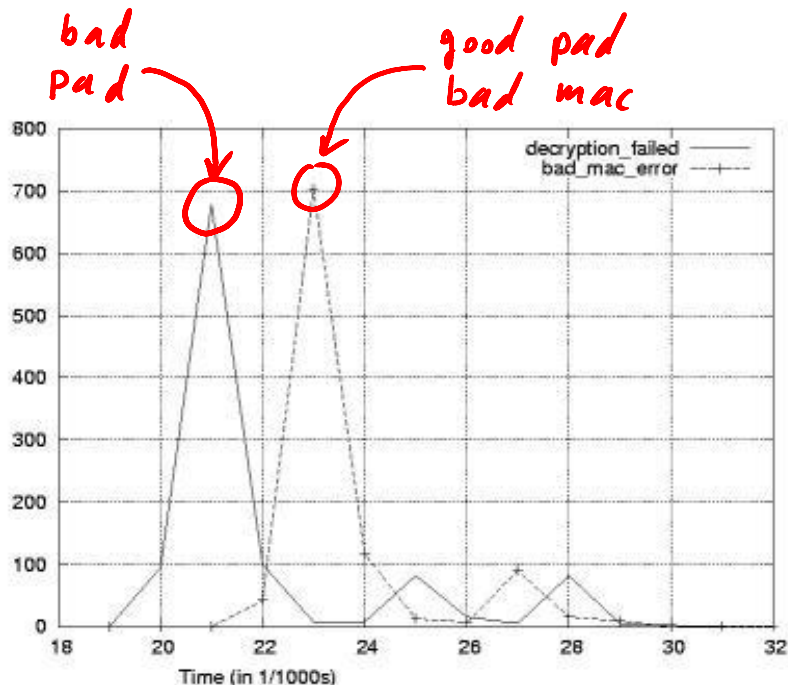
⇒ **Padding oracle:**

attacker submits ciphertext and learns if
last bytes of plaintext are a valid pad

Nice example of a
chosen ciphertext attack



Padding oracle via timing OpenSSL



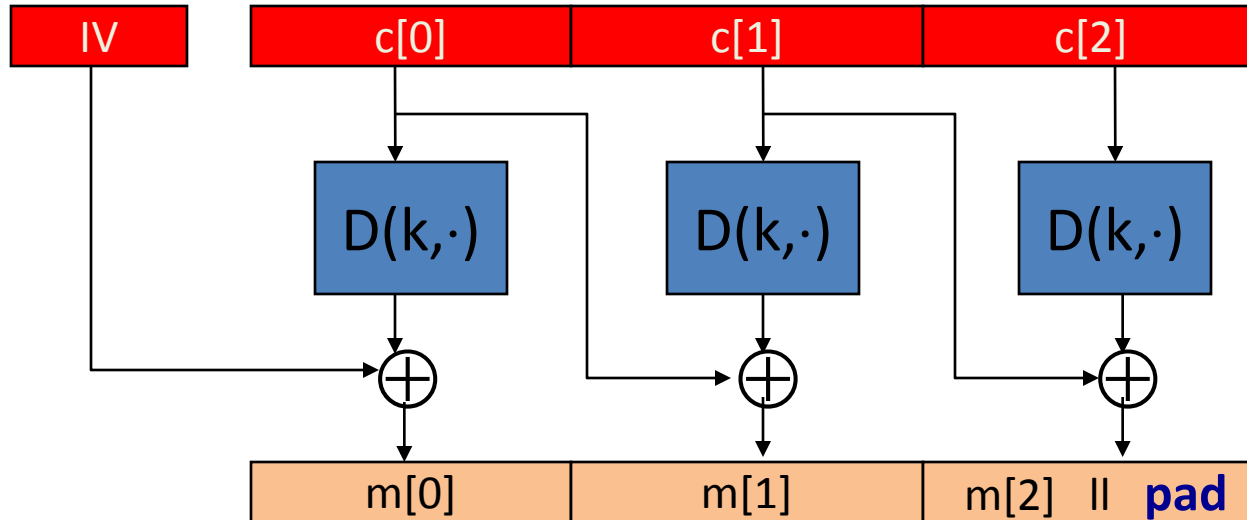
Credit: Brice Canvel

(fixed in OpenSSL 0.9.7a)

In older TLS 1.0: padding oracle due to different alert messages.

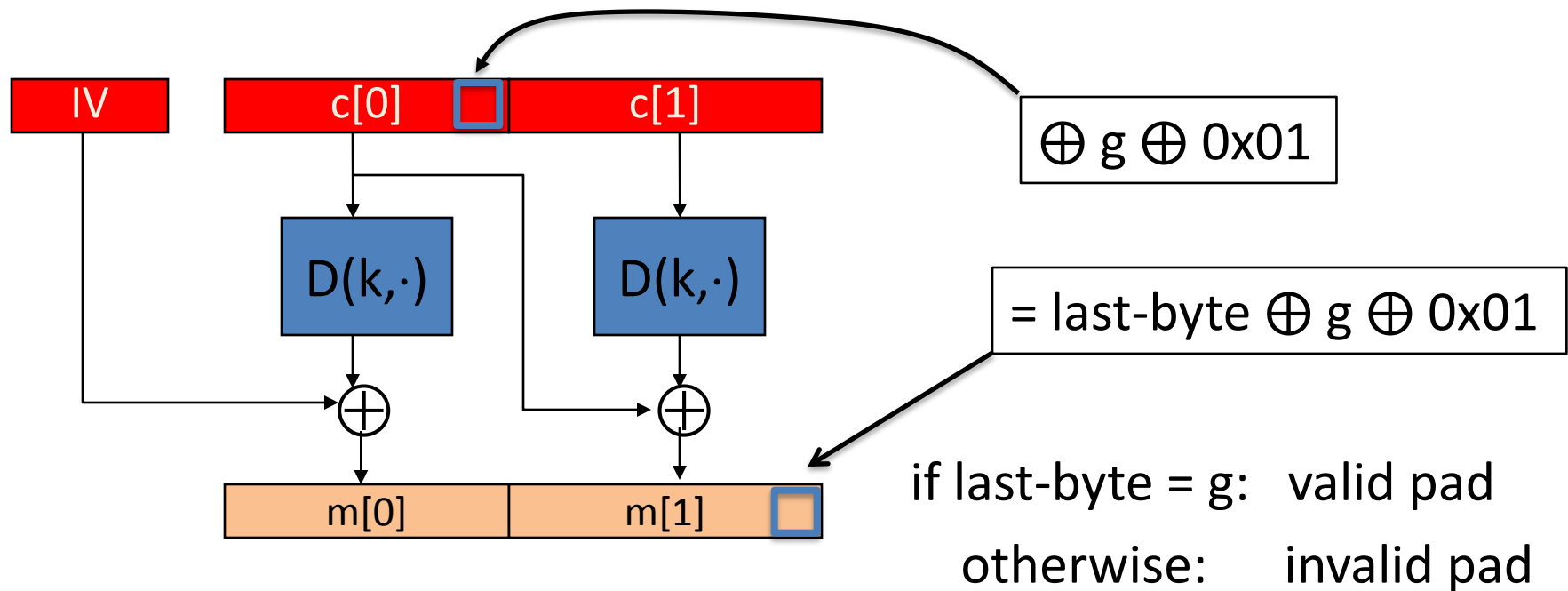
Using a padding oracle (CBC encryption)

Attacker has ciphertext $\mathbf{c} = (c[0], c[1], c[2])$ and it wants $\mathbf{m}[1]$



Using a padding oracle (CBC encryption)

step 1: let **g** be a guess for the last byte of $m[1]$



Using a padding oracle (CBC encryption)

Attack: submit $(IV, c'[0], c[1])$ to padding oracle

\Rightarrow attacker learns if last-byte = g

Repeat with $g = 0, 1, \dots, 255$ to learn last byte of $m[1]$

Then use a $(02, 02)$ pad to learn the next byte and so on ...

IMAP over TLS

Problem: TLS renegotiates key when an invalid record is received

Enter IMAP over TLS: (protocol for reading email)

- Every five minutes client sends login message to server:
LOGIN "username" "password"
- Exact same attack works, despite new keys
⇒ recovers password in a few hours.


Lesson

1. Encrypt-then-MAC would completely avoid this problem:

MAC is checked first and ciphertext discarded if invalid

2. MAC-then-CBC provides A.E., but padding oracle destroys it

Will this attack work if TLS used counter mode instead of CBC?
(i.e. use MAC-then-CTR)

- ☐ Yes, padding oracles affect all encryption schemes
- ☐ It depends on what block cipher is used
- ☐ No, counter mode need not use padding 
- ☐

End of Segment