

Information-theoretic inequalities.

1. Lemma [Log sum inequalities].

Let P, Q be two discrete distributions. Then

$$D_{KL}(P, Q) \geq p \cdot \log \frac{p}{q}, \quad (1)$$

where $p = \|P\|_1$, $q = \|Q\|_1$.

Proof.

To prove Eq (1), we only need to show that

$$\sum_{i=1}^n p_i \cdot \log \frac{p_i}{q_i} \geq p \cdot \frac{\log p}{\log q},$$

which is equivalent to prove that

$$\sum_{i=1}^n p_i \cdot \log \frac{q_i \cdot p_i}{p \cdot q_i} \geq 0.$$

To see this, observe that $\log \frac{1}{x} \geq 1 - x$ for any $x > 0$, and let $\lambda = \frac{q}{p}$. Then

$$\begin{aligned} \sum_{i=1}^n p_i \cdot \log \frac{q_i \cdot p_i}{p \cdot q_i} &= \sum_{i=1}^n p_i \cdot \log \lambda \cdot \frac{p_i}{q_i} \geq \sum_{i=1}^n p_i \cdot \left(1 - \frac{1}{\lambda} \cdot \frac{q_i}{p_i}\right) = \sum_{i=1}^n p_i \cdot \left(1 - \frac{p}{q} \cdot \frac{q_i}{p_i}\right) \\ &= \sum_{i=1}^n p_i - \frac{p}{q} \cdot q_i \\ &= 0. \end{aligned}$$

□

2. Information Processing Inequality.

Let X, Y be two random variables, which take values on Σ . Let $f: \Sigma \rightarrow f(\Sigma)$. Then

$$D_{KL}(P_{f(X)}, P_{f(Y)}) \leq D_{KL}(X, Y).$$

$$\text{Proof. } D_{KL}(x, y) = \sum_{w \in \Sigma} p_x(w) \log \frac{p_x(w)}{p_y(w)}$$

$$= \sum_{i \in f(x)} \sum_{w \in f^{-1}(i)} p_x(w) \log \frac{p_x(w)}{p_y(w)}$$

$$\geq \sum_{i \in f(x)} p_{f(x)}(i) \log \frac{p_{f(x)}(i)}{p_{f(y)}(i)}$$

$$= D_{KL}(f(x), f(y)).$$

□