

# Matrix exponentiation

# Binary Exponentiation

$$5^n = 5^{n/2} \cdot 5^{n/2}$$

$$\begin{array}{ccccccc} & 1 & 1 & 1 & & & \\ & \swarrow & \swarrow & \swarrow & & & \\ \boxed{5} & | & 5^2 & | & 5^4 & | & 5^8 & | & 5^{16} \\ & & & & & & & & \\ & & & & & & & & = 4 \end{array}$$

# Binary Exponentiation - Recursive

- Int power(int a,int b)
- { if(b==0) return 1;
  - int ans=power(a,b/2);
  - if(b%2) return (ans\*ans\*a);
  - Else return (ans\*ans);
  - }

# Binary Exponentiation - Iterative

- Int power(int a,int b)
  - { int result=1;
    - while(b>0)
    - { if(b%2)           result=result\*a;
      - a\*=a;
      - b/=2;
    - }
  - return result;
  - }

# Fibonacci

- $f(n) = f(n-1) + f(n-2)$
- $f(n) = a*f(n-1)+b*f(n-2)$  where  $a=1,b=1$

- $F(n) = \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} F_{n-1} \\ F_{n-2} \end{bmatrix}$

# Fibonacci – Multiplication property

$$\begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} [1] & \begin{bmatrix} F_{n-2} \\ F_{n-3} \end{bmatrix} \\ [1] & \begin{bmatrix} F_{n-3} \\ F_{n-4} \end{bmatrix} \end{bmatrix}$$

# Fibonacci – Multiplication Property

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^n = \begin{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} F_{n-2} \\ F_{n-3} \end{bmatrix} \\ \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} F_{n-3} \\ F_{n-4} \end{bmatrix} \end{bmatrix} = \begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix}$$

# Fibonacci Formula

$$\begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-1} \begin{bmatrix} F_1 \\ F_0 \end{bmatrix}$$

# Fibonacci code

```
void multiply(int F[2][2], int M[2][2]) {
    int a = F[0][0] * M[0][0] + F[0][1] * M[1][0];
    int b= F[0][0] * M[0][1] + F[0][1] * M[1][1];
    int c = F[1][0] * M[0][0] + F[1][1] * M[1][0];
    int d = F[1][0] * M[0][1] + F[1][1] * M[1][1];
    F[0][0] = a;
    F[0][1] = b;
    F[1][0] = c;
    F[1][1] = d;
}

void power(int F[2][2], int n) {
    if (n == 0 || n == 1)
        return;
    int M[2][2] = {{1,1},{1,0}};
    power(F, n / 2);
    multiply(F, F);
    if (n % 2 != 0)
        multiply(F, M);
}
```

# Bonus

- Binet's formula gives you fibonacci number in  $O(\log n)$  time.
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$$F_n = \left[ \frac{\left( \frac{1+\sqrt{5}}{2} \right)^n}{\sqrt{5}} \right]$$