

Architecturally-defined Attestation, Attested TLS and Open Challenges

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Based on joint works with Arto Niemi, Hannes Tschofenig, Thomas Fossati, Simon Frost, Ned Smith, Mariam Moustafa, Tuomas Aura, Yaron Sheffer, Ionut Mihalcea, Jean-Marie Jacquet and Henk Birkholz

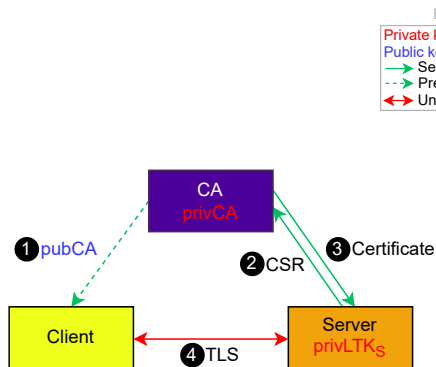
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Outline

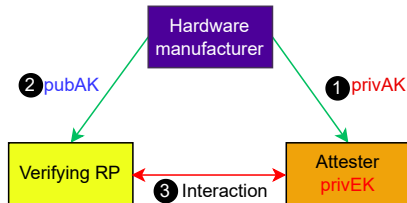
- 1 Attested TLS
- 2 Open Problems

TLS vs. Architecturally-defined Attestation



- CA as Trust Anchor

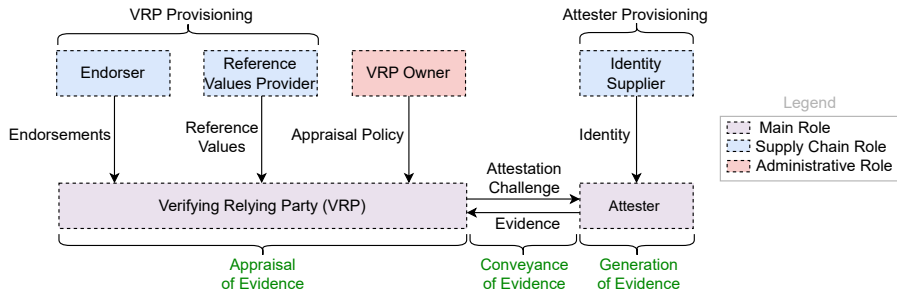
$$Cert = \text{sign}(\text{privCA}, ID \parallel \text{pubLTK})$$



- HW manufacturer as Trust Anchor

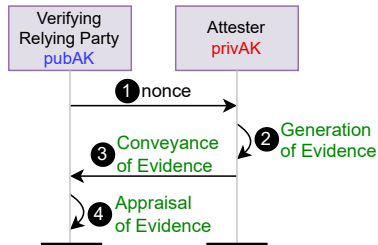
$$Evidence = \text{sign}(\text{privAK}, m \parallel \text{pubEK})$$

Architecturally-defined Attestation: Architecture Perspective

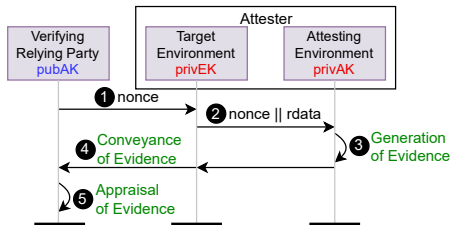


Architecturally-defined Attestation: Protocol Perspective

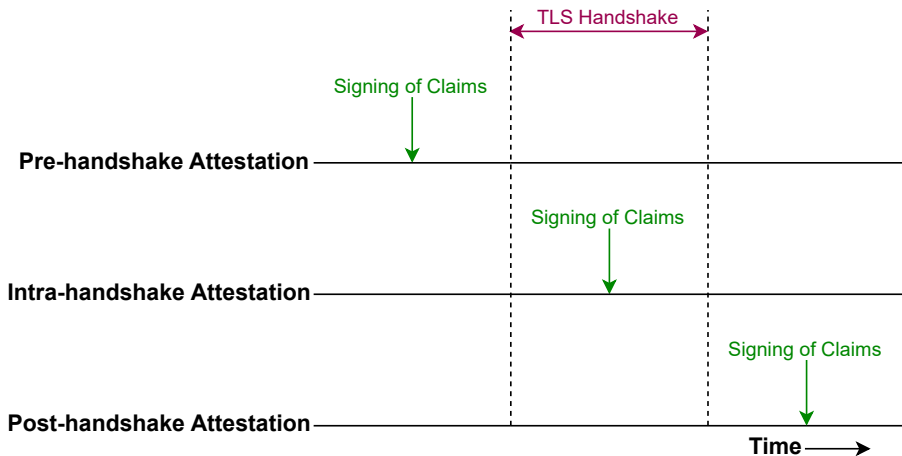
- Integrated Attester



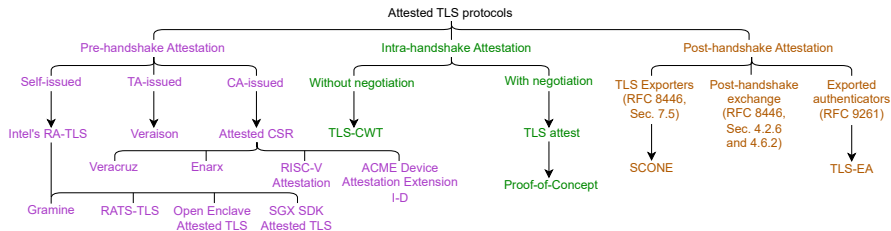
- Attester with separate env.



Attested TLS

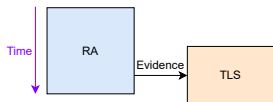


Main Design Options

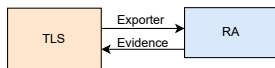


Channel Bindings

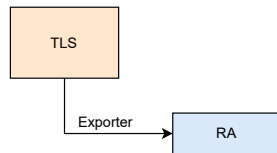
- Pre-HS



- Intra-HS



- Post-HS

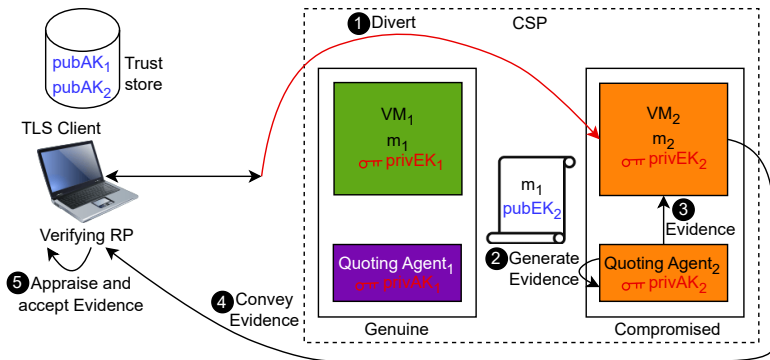


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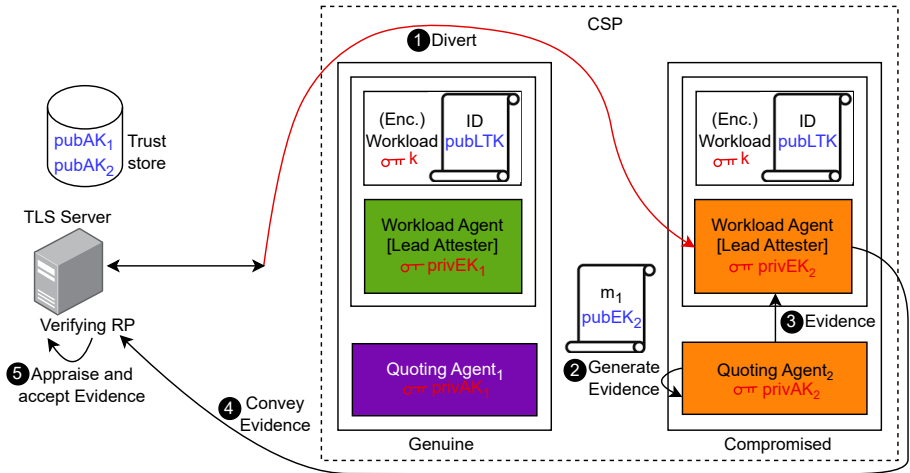
A Diversion Attack Within CSP

- AK of a specific machine may be compromised. (e.g., privAK_2)
 - Transient execution attacks, as demonstrated by Foreshadow¹
- VM_2 impersonates VM_1



¹Van Bulck, Minkin, Weisse, Genkin, Kasikci, Piessens, Silberstein, Wenisch, Yarom, and Strackx, "Foreshadow: Extracting the Keys to the Intel SGX Kingdom with Transient Out-of-Order Execution", 2018.

B How to Provision ID and LTK?



- CSP gets the key k

Questions for Discussion

- Threat model: CSP as fully untrusted vs. “Honest but curious”? [Fritz]

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- Equivalence class? [Keith]
- Temporal: How has pubAK become untrusted? [Jim]