## 3.1.6 Conditional expectations

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In the discrete case, the conditional probability mass function of  $\xi$  for given  $\eta$  is

$$p_{\xi|\eta}(x_i|y_j) = \frac{P(\xi = x_i, \eta = y_j)}{P(\eta = y_j)} = \frac{p_{ij}}{p_{ij}}.$$

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In the continuous case, the conditional probability density function of  $\xi$  for given  $\eta$  is

$$p_{\xi|\eta}(x|y) = \frac{p(x,y)}{p_{\eta}(y)}.$$

In general, if the limit

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$$F_{\xi|\eta}(x|y) = \lim_{\epsilon \to 0} \frac{P(\xi \le x, -\epsilon + y < \eta < \epsilon + y)}{P(-\epsilon + y < \eta < \epsilon + y)}$$

exists for all x, we call it the conditional distribution of  $\xi$  for given  $\eta=y.$ 

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The mathematical expectation of a conditional distribution is called the conditional mathematical expectation:

$$E[\eta|\xi=x] = \int_{-\infty}^{\infty} y dF_{\eta|\xi}(y|x).$$

为了强调y是积分变量,上述积分也常写为 $\int_{-\infty}^{\infty} y F_{\eta|\xi}(dy|x)$ .

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If  $\eta$  has conditional density  $p_{\eta|\xi}(y|x)$  given  $\xi = x$ , then

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## Example

$$(\xi,\eta) \sim N(a,b,\sigma_1^2,\sigma_2^2,r)$$
, then  $\eta|_{\xi=x} \sim$ 

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, then  $\eta|_{\xi=x} \sim N(b + r\sigma_2(x - a)/\sigma_1, \sigma_2^2(1 - r^2))$ .

In turn,

$$E(\eta|\xi=x) = b + r\frac{\sigma_2}{\sigma_1}(x-a).$$

Denote by  $E(\eta|\xi)$ : when  $\xi=x$  the function takes value  $E(\eta|\xi=x).$   $E(\eta|\xi)$  is a r.v. and a function of  $\xi$ .

$$E(E(\eta|\xi)) = E\eta.$$

Proof.

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**Proof.** We give the proof only for continuous random variables below. Suppose that  $(\xi, \eta)$  has the pdf p(x, y). In this case,

$$p_{\xi}(x) = \int_{-\infty}^{\infty} p(x, y) dy,$$

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### and then

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$$= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} y p(x,y) dy dx = E\eta.$$

When  $\xi$  is a discrete random variable, letting  $p_i = P(\xi = x_i)$ , then

$$E\eta = \sum_{i} p_i E(\eta | \xi = x_i).$$

This is similar to the total probability formula, called the total expectation formula.

A miner is trapped in a mine containing 3 doors. The first door leads to a tunnel that will take him to safety after 3 hours of travel. The second door leads to a tunnel that will return him to the mine after 5 hours of travel. The third door leads to a tunnel that will return him to the mine after 7 hours. If we assume that the miner is at all times equally likely to choose any one of the doors, what is the expected length of time until he reaches safety?

$$E[\xi|\eta=1]=3,$$

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#### Now

$$E\xi = E[\xi|\eta = 1]P(\eta = 1)$$
  
  $+E[\xi|\eta = 2]P(\eta = 2)$   
  $+E[\xi|\eta = 3]P(\eta = 3)$ 

Now

$$E\xi = E[\xi|\eta = 1]P(\eta = 1) + E[\xi|\eta = 2]P(\eta = 2) + E[\xi|\eta = 3]P(\eta = 3) = \frac{1}{3}(3 + 5 + E\xi + 7 + E\xi) = 5 + \frac{2}{3}E\xi.$$

$$E\xi = 15.$$

### Example

$$\{\xi_i, i \geq 1\}$$
 i.i.d  $\sim B(n, p), \nu \sim P(\lambda)$ .  $\nu$  is independent of  $\{\xi_i, i \geq 1\}$ . Find  $E(\sum_{i=1}^{\nu} \xi_i)$ .

### Solution.

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Solution. Let  $\eta = \sum_{i=1}^{\nu} \xi_i$ , then small

$$E(\eta|\nu=r) = E(\sum_{i=1}^{\nu} \xi_i|\nu=r)$$

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$$= E(\sum_{i=1}^{r} \xi_i) = \sum_{i=1}^{r} E\xi_i = rnp.$$

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$$E\eta = \sum_{r=0}^{\infty} E(\eta|\nu=r)P(\nu=r)$$

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$$= np\sum_{r=1}^{\infty} rP(\nu=r) = npE\nu = np\lambda.$$

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$$E[g(\eta)|\xi = x] = \sum_{j} g(y_j) p_{\eta|\xi}(y_j|x)$$

in the discrete case,

$$E[g(\eta)|\xi=x] = \int_{-\infty}^{\infty} g(y)p_{\eta|\xi}(y|x)dy$$

in the continuous case,

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Example

The quick-sort algorithm(快速排序法)

The quick-sort algorithm(快速排序法)

设有n个不同的数 $x_1, x_2, \ldots, x_n$ .我们要将它们按从小到大的次序排列起来 $x_{(1)} < x_{(2)} < \ldots < x_{(n)}$ . 进行这样的排列需要对这n个数两两进行比较, 如果全部进行, 共需要比较 $\frac{n(n-1)}{2}$ 次, 这就是用计算机排列这些数所需要的计算量. 是否有一种算法可以减少比较次数呢?

一种称为快速排序算法(quick-sort algorithm)是这样进行的: 从集合 $\{x_1, x_2, \ldots, x_n\}$ 中随机地取一个数 $x_J$ ,将其它数与 $x_J$ 进行比较,把小于 $x_J$ 的数放在其左边,这样的数构成集合L,把大于 $x_J$ 的数放在其右边,这样的数构成集合R. 然后,对L和R进行同样处理,依此类推,直到最后每个集合中只有一个数为止. 我们用 $\xi$ 表示进行比较的总次数,求 $q_n = E\xi$ .

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**解:** 设用快速排序法排 $L(x_J$ 左边的数)所需要比较进行次数为 $\xi_L$ , 排 $R(x_J$ 右边的数)所需要进行比较次数为 $\xi_R$ . 则

$$\xi = \xi_L + \xi_R + (n-1).$$

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$$\xi = \xi_L + \xi_R + (n - 1).$$

在 $x_J = x_{(i)}$ 的条件下, L和R中分别有i - 1, n - i个元素. 所以

$$E[\xi_L|x_J = x_{(i)}] = q_{i-1}, \ E[\xi_R|x_J = x_{(i)}] = q_{n-i}.$$

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而

$$P(x_J = x_{(i)}) = \frac{1}{n}$$
.

所以

$$q_n = E\xi = \sum_{i=1}^n E[\xi | x_J = x_{(i)}] P(x_J = x_{(i)})$$
$$= n - 1 + \frac{1}{n} \sum_{i=1}^n (q_{i-1} + q_{n-i}) = n - 1 + \frac{2}{n} \sum_{i=1}^n q_{i-1}.$$

### 所以

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$$= n - 1 + \frac{1}{n} \sum_{i=1}^n (q_{i-1} + q_{n-i}) = n - 1 + \frac{2}{n} \sum_{i=1}^n q_{i-1}.$$

$$nq_n = n(n-1) + 2 \sum_{i=1}^n q_{i-1}.$$

$$nq_n - (n-1)q_{n-1} = n(n-1) - (n-1)(n-2) + 2q_{n-1}.$$

 $na_n = 2(n-1) + (n+1)a_{n-1}$ 

$$\frac{q_n}{n+1} = \frac{q_{n-1}}{n} + \frac{2(n-1)}{n(n+1)}$$

$$= \frac{q_{n-1}}{n} + \frac{2}{n} + 4\left(\frac{1}{n+1} - \frac{1}{n}\right)$$

$$= \dots = 2\sum_{k=1}^{n} \frac{1}{k} + \frac{4}{n+1} - 4$$

$$= 2\left(\log n + \gamma + \frac{1}{2n} + O(n^{-2})\right) - \frac{4n}{n+1}.$$

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因此

$$q_n = 2(n+1)\log n + n(2\gamma - 4) + 2\gamma + 1 + O(n^{-1}).$$