

GoSeed: Generating an Optimal Seeding Plan for Deduplicated Storage

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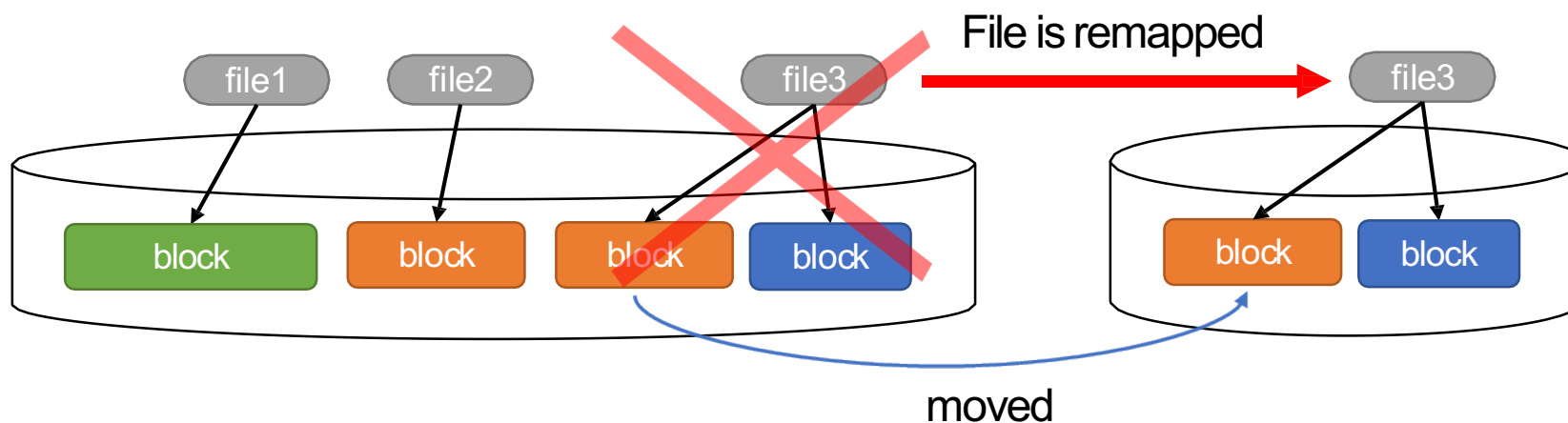
Speaker: Liang, Jiacheng

Background

➤ Data Migration:

- Moving data between volume
- Typically performed for load balancing and resizing

➤ Seeding : Data migration with empty destination



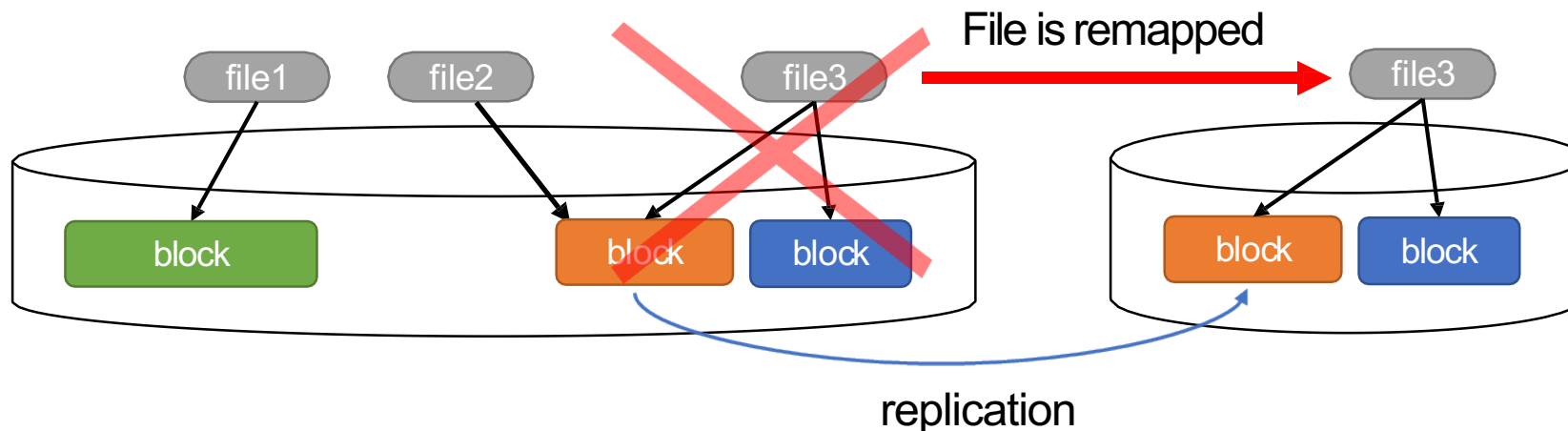
Motivation

➤ Deduplication complicates the migration:

- File remapped will cause replication.

➤ Current status:

- Greedy algorithms for determining the set of migrated files, but the efficiency has never been systematically compared.



Background :What is ILP

➤Satisfies

$a_0x_0 + a_1x_1 + a_2x_3 + \cdots + a_{n-1}x_{n-1} \leq c$, where $a_0, \dots, a_{n-1}, c \in \mathbb{Z}$

➤Maximizes

$Tx = a_0x_0 + a_1x_1 + a_2x_3 + a_0x_0 + \cdots + a_{n-1}x_{n-1}$

➤How to combine data migration

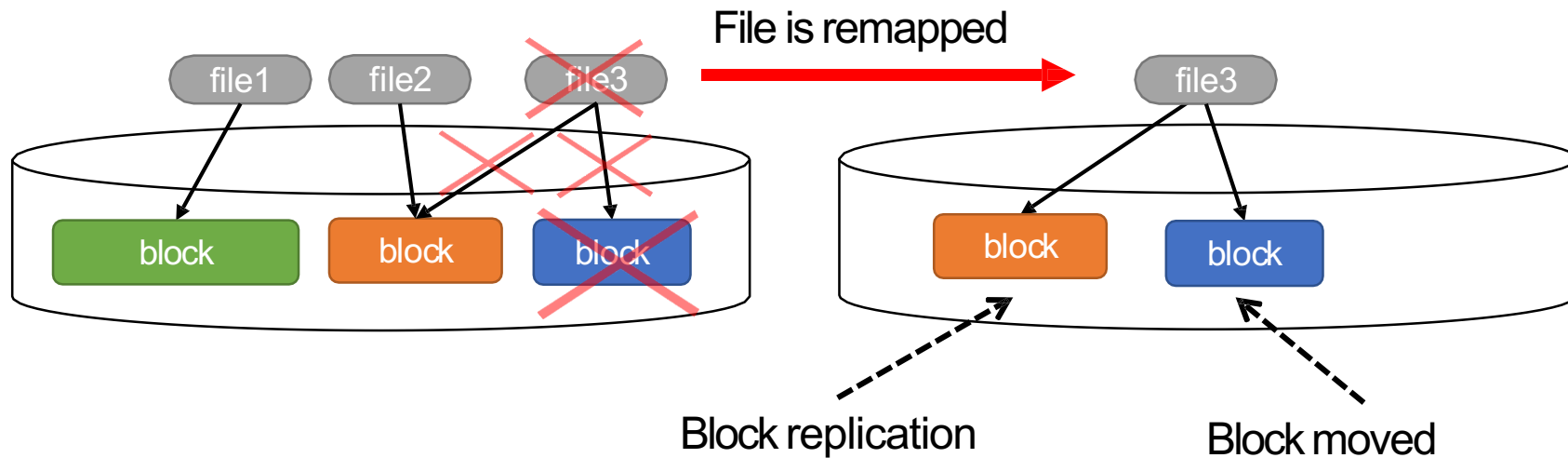
c as the the migration amount, make $a_0x_0 + a_1x_1 + a_2x_3 + a_0x_0 + \cdots + a_{n-1}x_{n-1}$ closest to c .

Design: ILP combine data migration

- Migrate $M\%$ of physical occupancy
 - Actually $M \pm \epsilon$
 - R as the cost of the migration
 - m_i : block i migrated (0/1)
 - r_i : block i replicated (0/1)
 - x_i : file i remmapped (0/1)
 - s : the size of the block
 - M : $\sum b_i \in B_{V_1} s(b_i) \cdot m_i = M$, The total size of migrated blocks
 - R : $\sum b_i \in B_{V_1} s(b_i) \cdot r_i$
- target : minimize R ,closest M

Constraint : blocks & file are always bound

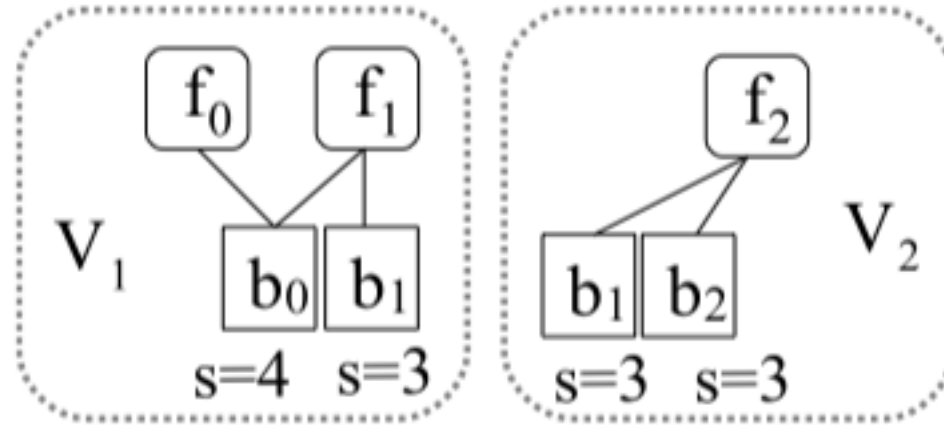
- Blocks are copied or moved with their files
- Blocks cannot move without their files



Design: ILP formulation

- m_i : block i migrated (0/1)
- r_i : block i replicated (0/1)
- x_i : file i remapped (0/1)
- s : the size of the block
- M : $\sum b_i \in B_{V_1} s(b_i) \cdot m_i = M$,
The total size of migrated blocks
- R : $\sum b_i \in B_{V_1} s(b_i) \cdot r_i$

target : minimize R ,closest M



- Block b_2 is migrated: $m_2 = 1$
- File f_2 is remapped: $x_2 = 1$
- Block b_1 is replicated: $r_1 = 1$
- The remaining files and blocks are untouched:
 $x_0 = x_1 = m_0 = m_1 = r_0 = r_2 = 0$
- The total cost is $R = 3 \cdot r_1 = 3$

Theory vs. Practice

- Dedup: ~10 logical TB \Rightarrow ~billion variables and constraints
- ILP Solvers: Solvers are efficient with several 100K

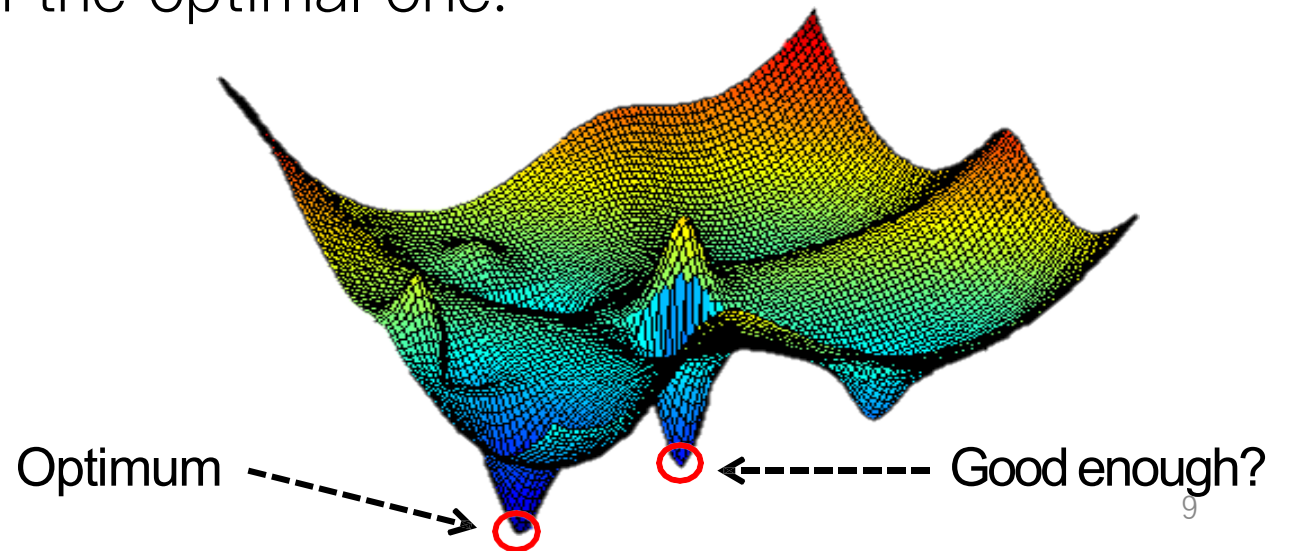
Method I: Solver timeout

Limited the runtime of ILP solver

- Specifying a timeout value
- Return the best feasible solution found thus far

Problem:

- When the solver is timed out, we cannot necessarily tell how far the suboptimal solution is from the optimal one.



Method II: fingerprint sampling

- Sample a subset of the fingerprints [Harnik et al. FAST19]
- Sample degree k : each sample $\frac{1}{.2^k}$ of the blocks

Problem:

- Smaller sample
 - ⇒ a larger ILP instance
 - ⇒ more accurately represents the sampled system
 - ⇒ more likely to time out

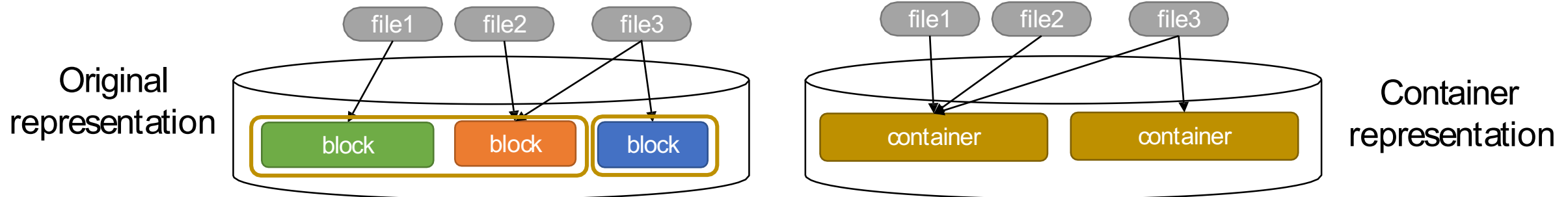
Method III: container-based aggregation

- Deduplication system stores blocks in containers
- We treat each container as a block
- use only the container IDs for generating the variables and constraints.

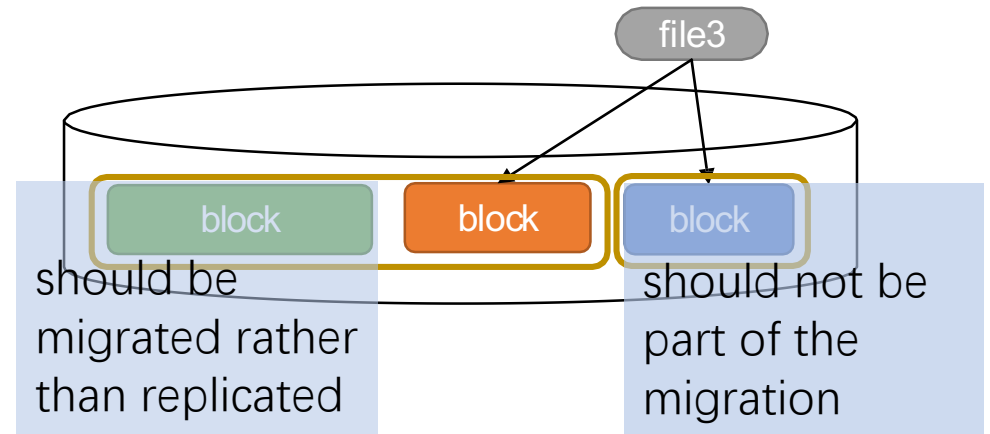
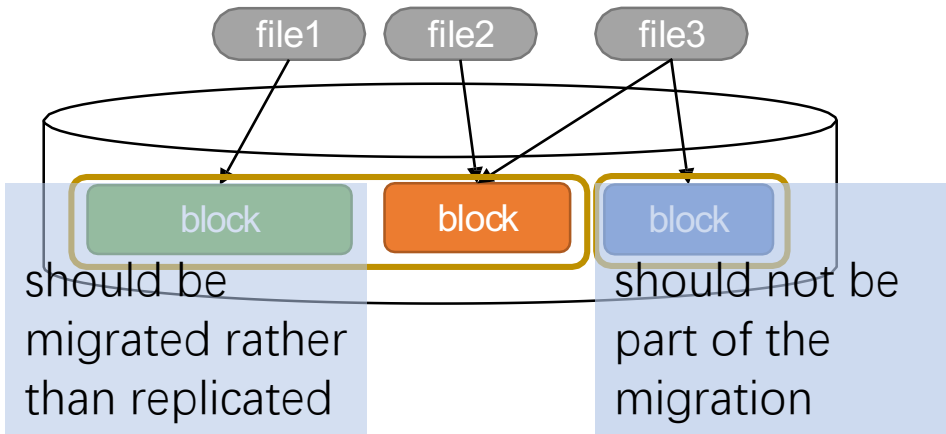
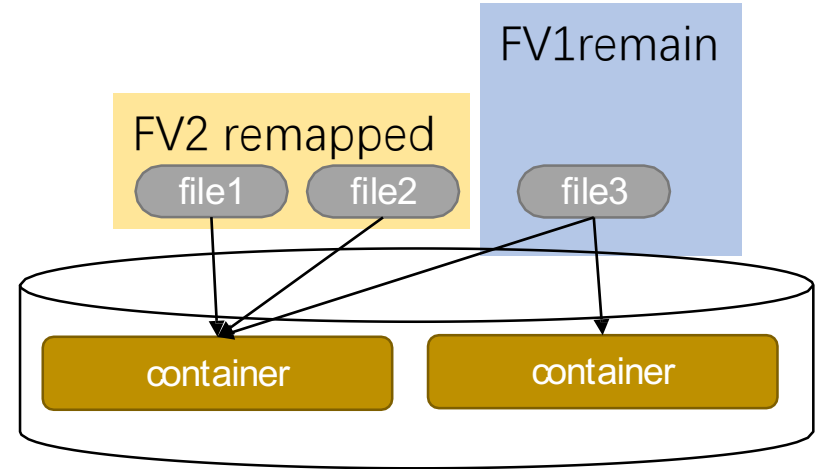
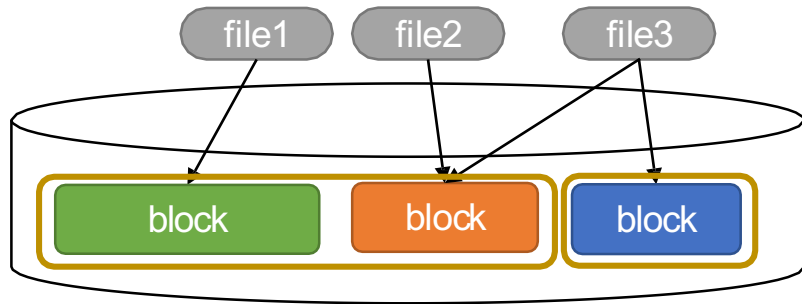
+smaller ILP problem

+no need to decompress/unpack containers

–false positives



-false positives



Experimental setup

- We compare:
- GoSeed: our ILP-based approach
- Rangoli [Nagesh & Kathpal Systor'13]:
 - Sort & divide the blocks into bins
 - Migrate the “best” bin
- SGreedy [Harnik et al. FAST19]:
 - Iterative method
 - Each iteration remap the “best” file

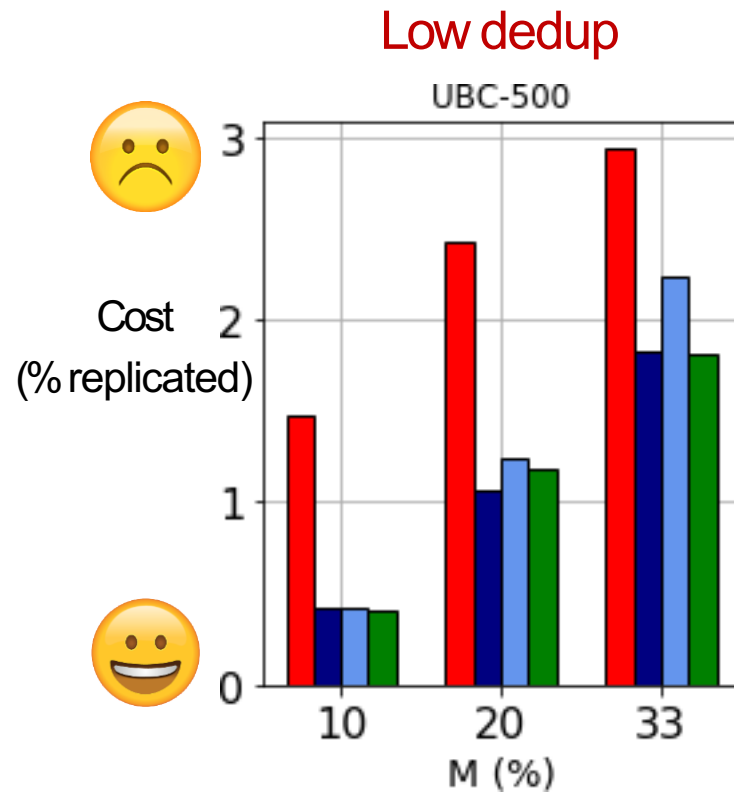


	UBC 500	Homes	MacOS -day
Logical size	19.5TB	8.9TB	43TB
Dedup ratio	0.31	0.13	0.01

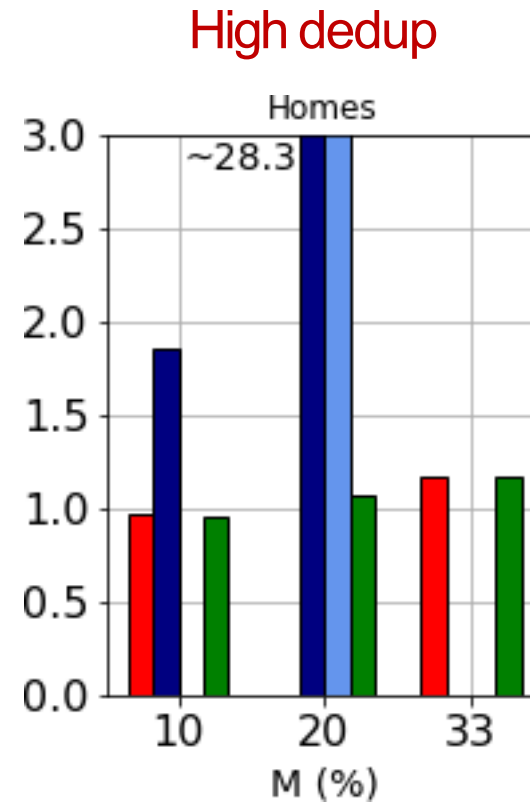
Cost



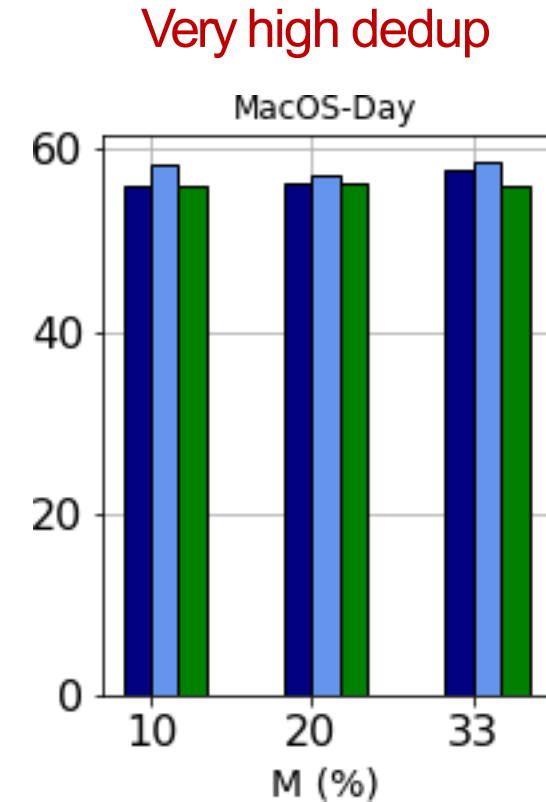
- UBC-500 is considered “easy”



- Homes is harder

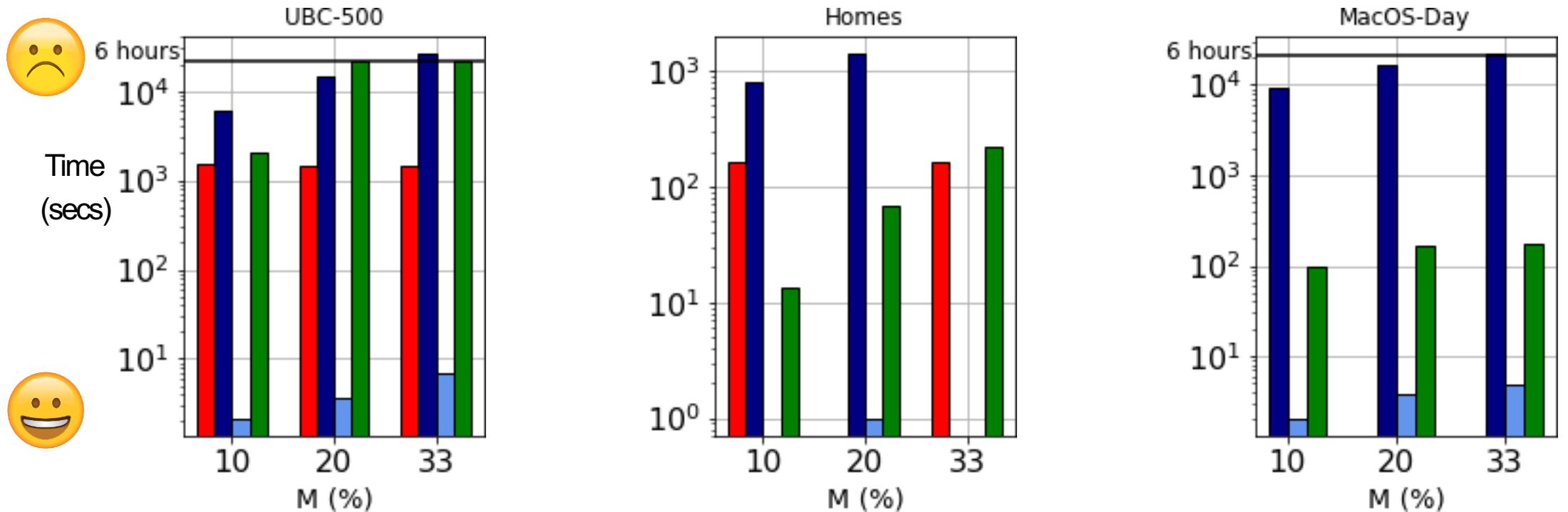


- MacOS-Day consists daily backups



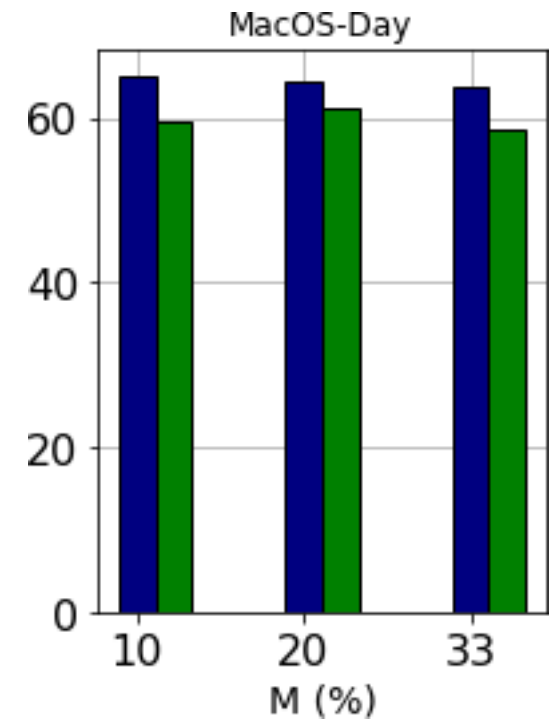
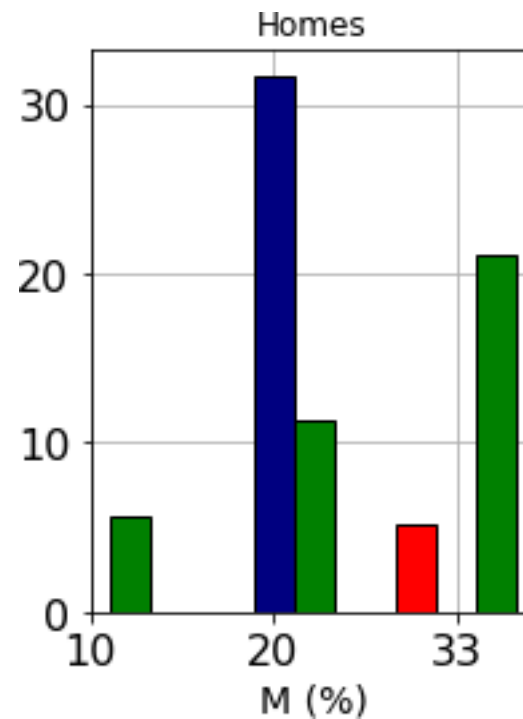
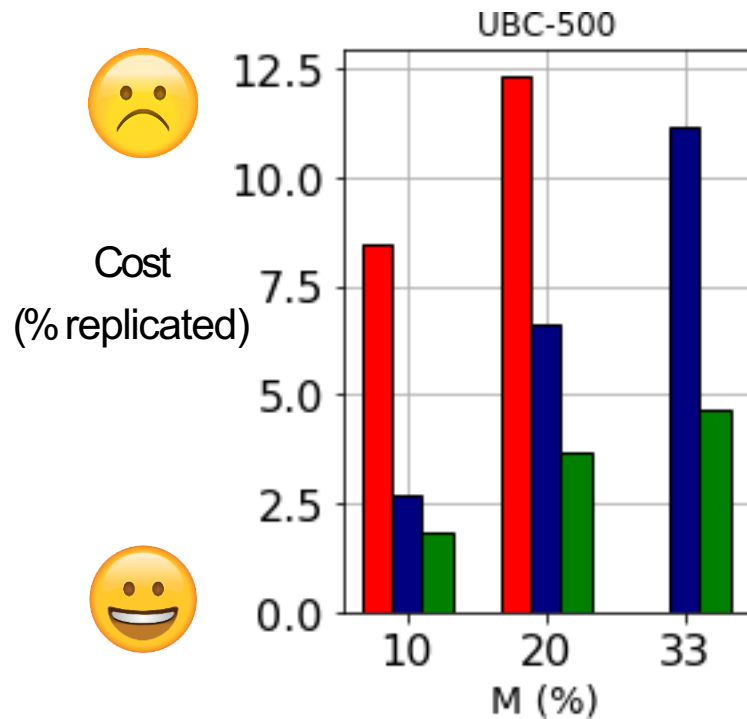
- GoSeed always finds a solution, it's often the best
- Greedy algorithms sometimes fail to find a solution

Time



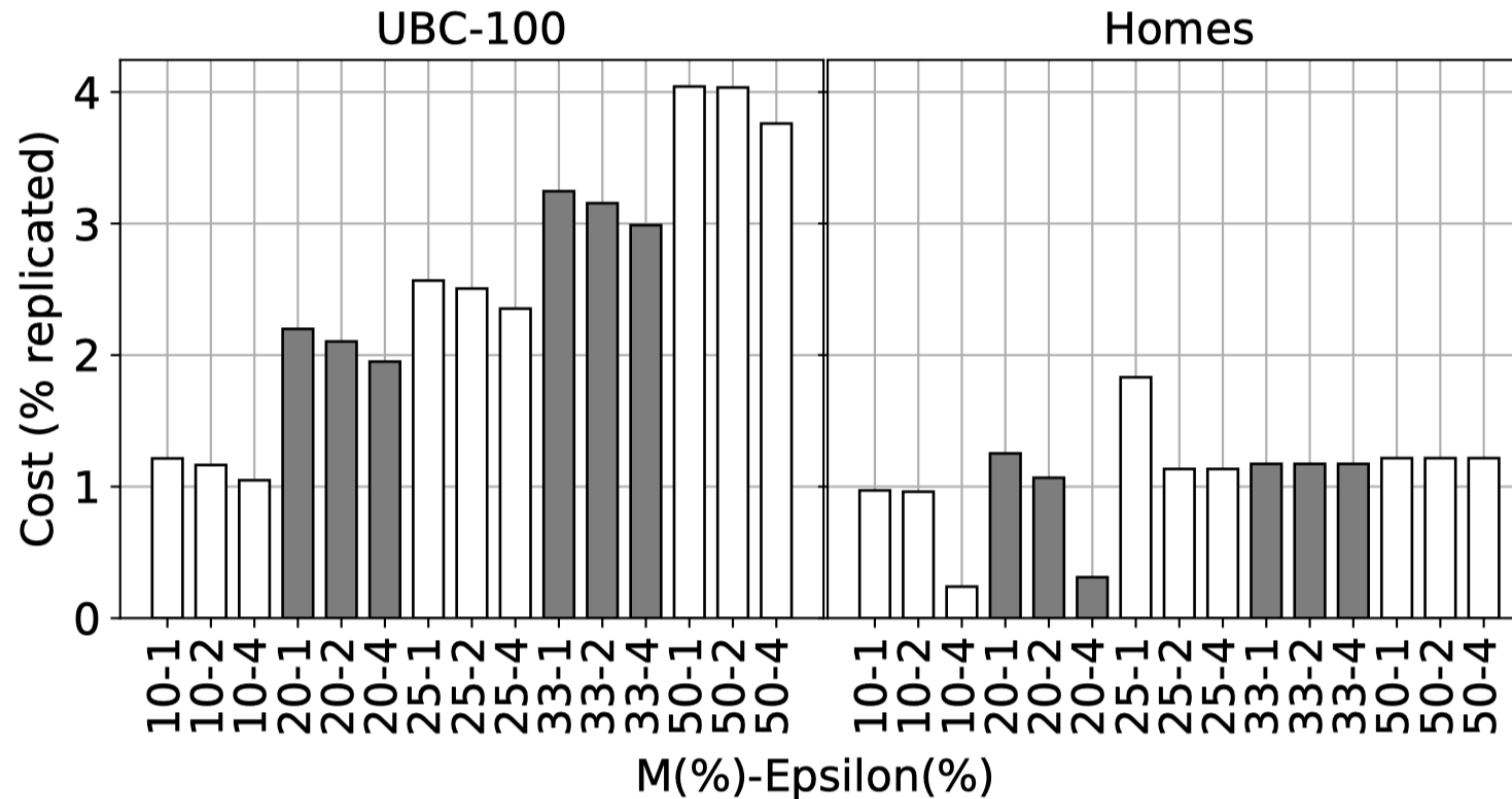
➤ Runtimes and resources are acceptable.

Container cost



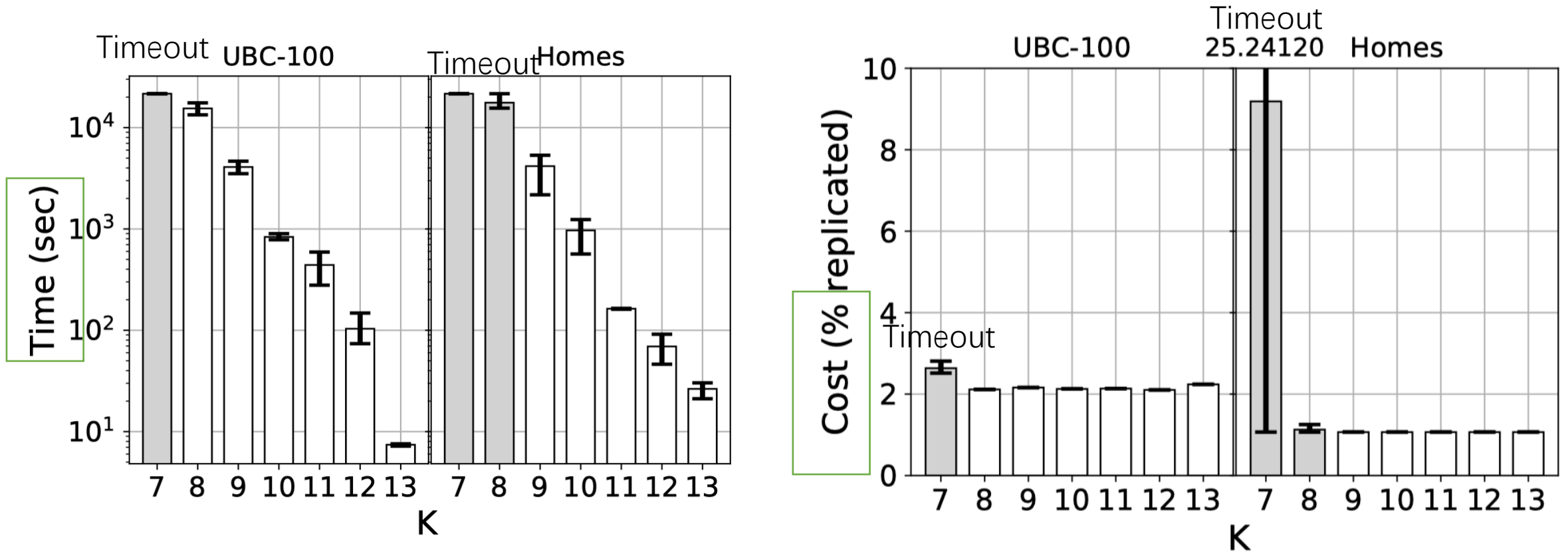
➤ GoSeed always gives a solution, almost always better

ILP parameter M-Epsilon Evaluation



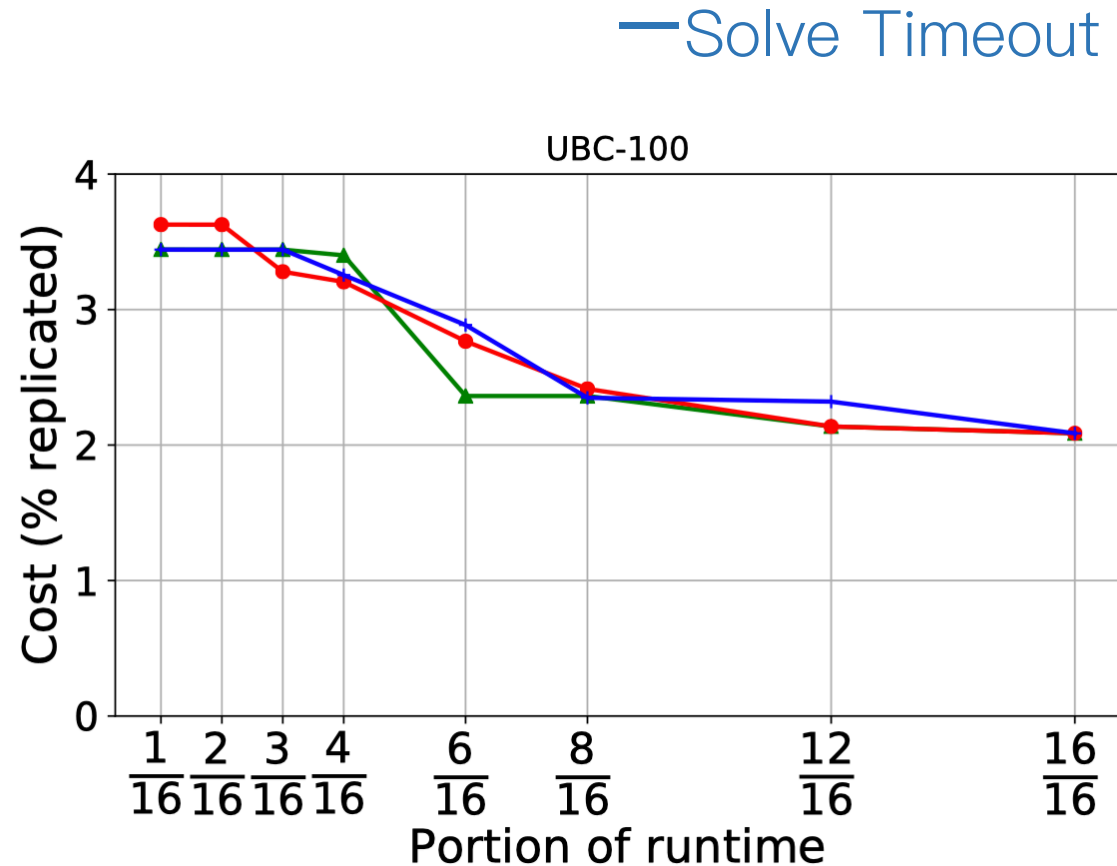
- The increase of M leads to higher cost, and the increase of ϵ leads to the decrease of cost

Effect of fingerprint sampling



- Solving time increases exponentially with instance size
- Cost is hardly affected by the sampling degree, unless the instance becomes too large.

Suboptimal solution



- Migration cost decreases when timeout increases

Conclusions

- Formulate of data migration as an ILP problem and solve it with ILP solvers
- Use three acceleration methods and they are essential and effective
- GoSeed often outperforms greedy approaches, and always find a feasible solution to the problem.
- The container-based approach is separate from the sample-based approach, applicable to different deduplication systems