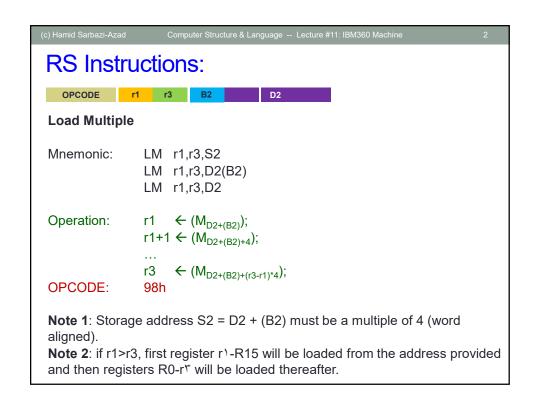
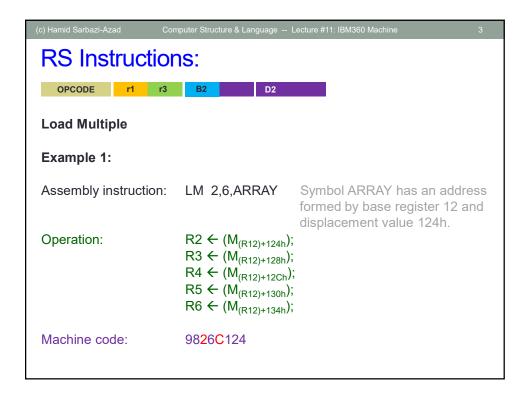
## Computer Structure and Language

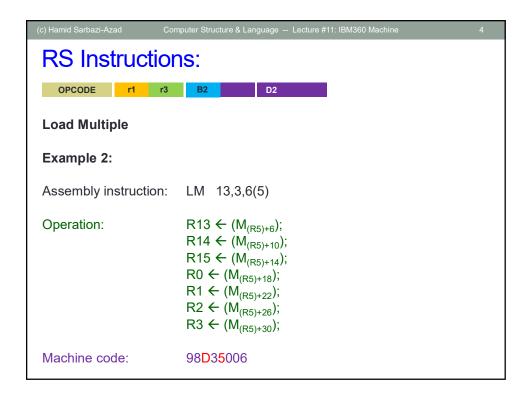
## Hamid Sarbazi-Azad

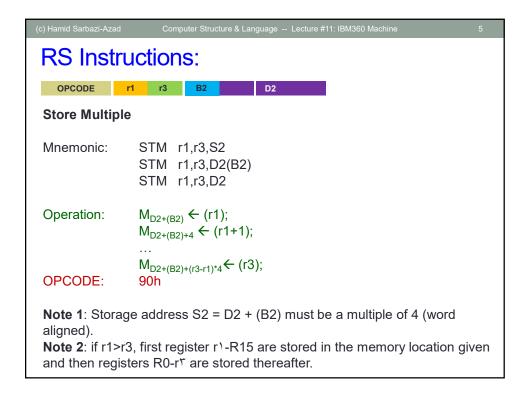
Department of Computer Engineering Sharif University of Technology (SUT) Tehran, Iran

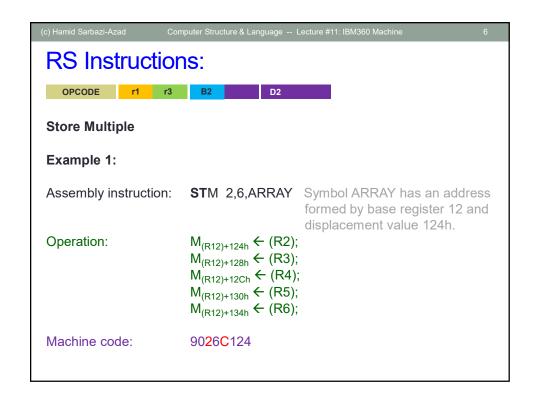


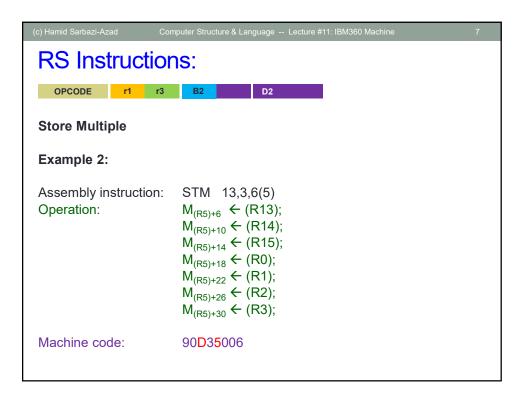


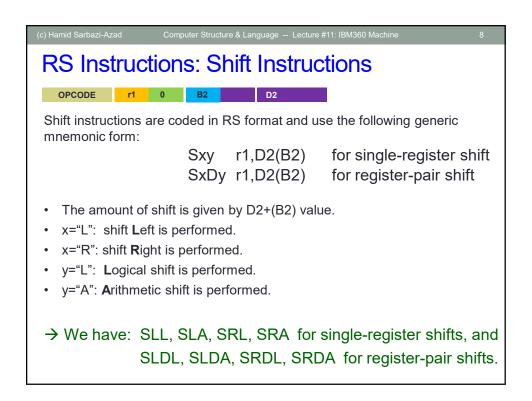


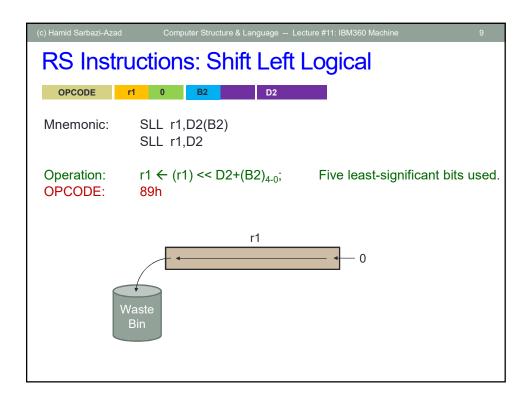


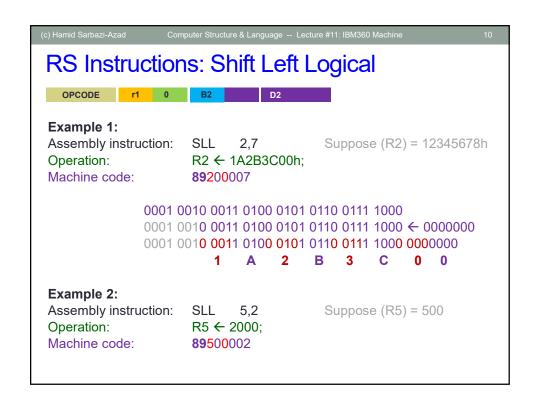


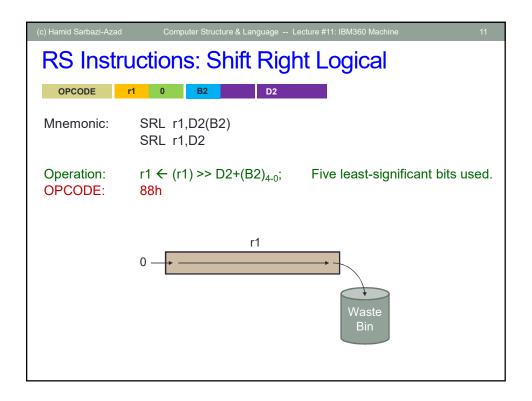


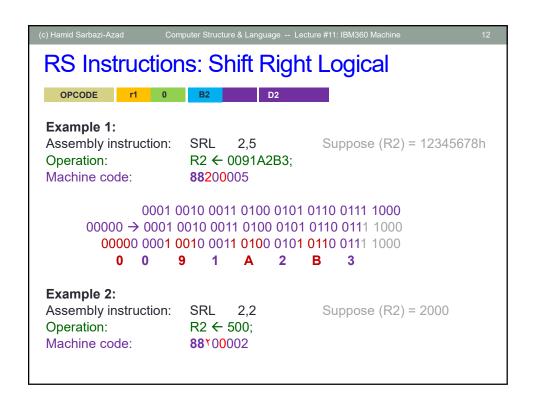


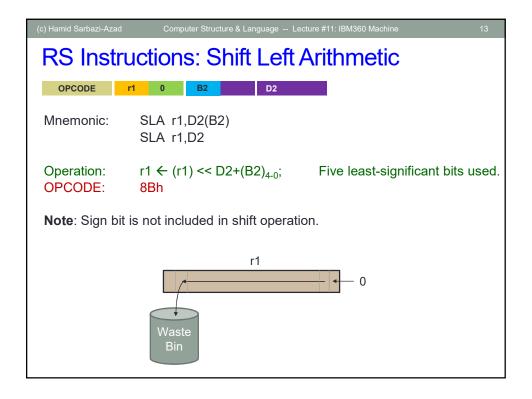


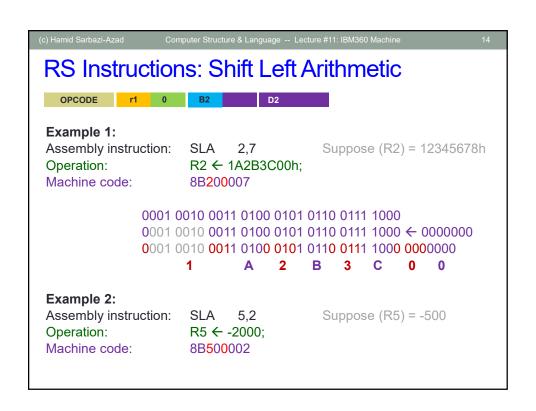


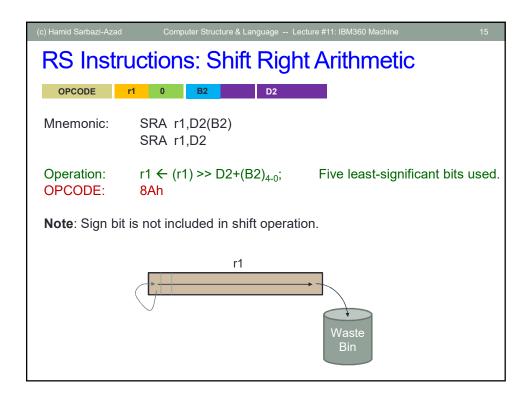


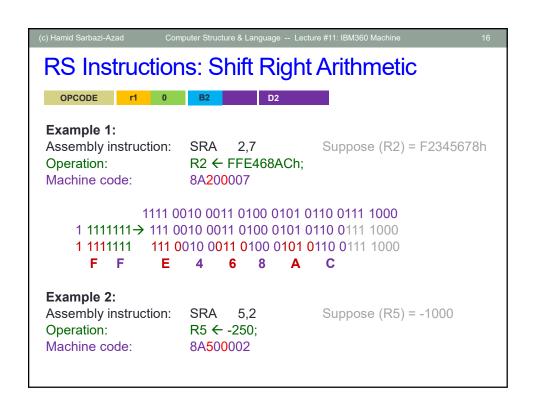


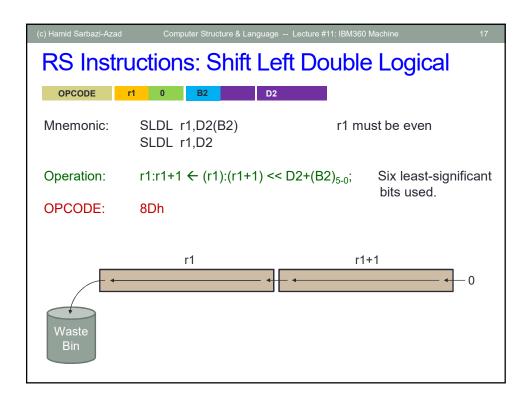


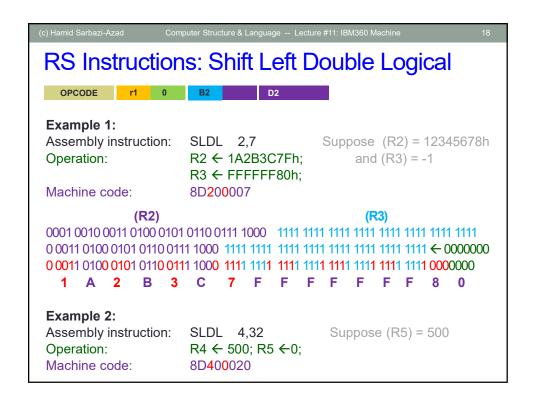


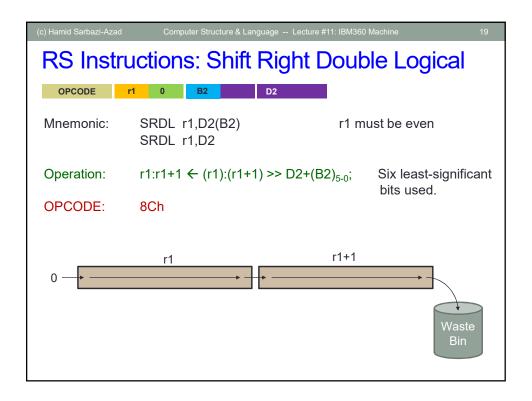


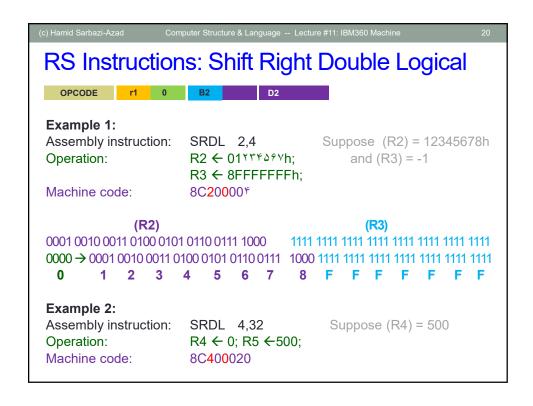


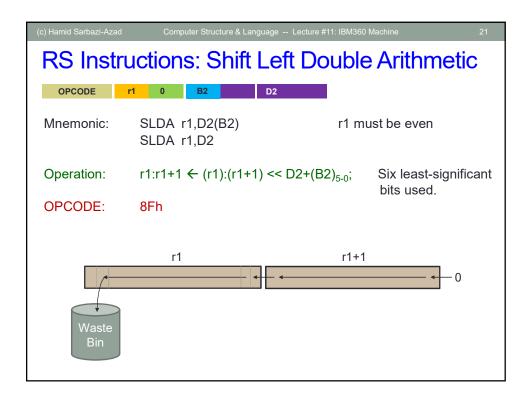


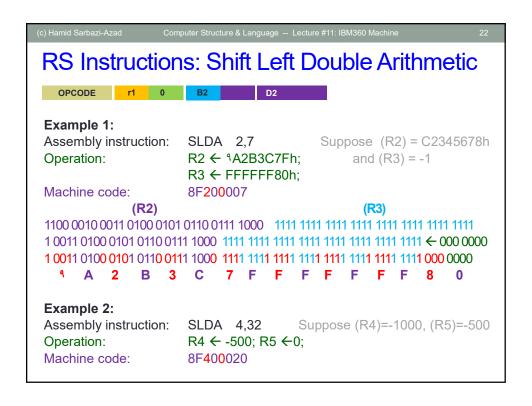


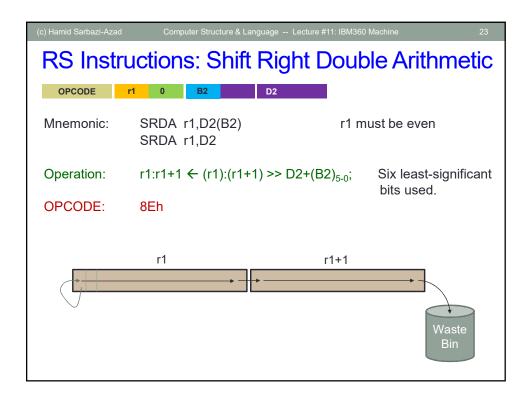


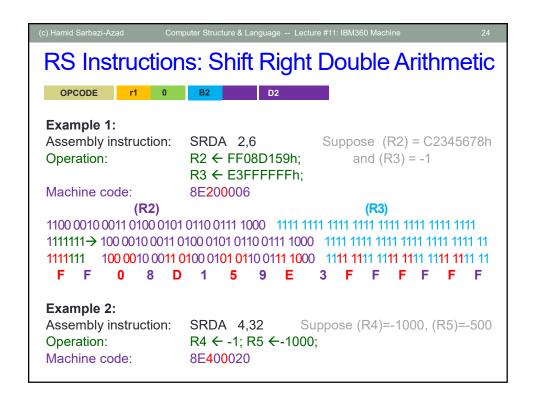


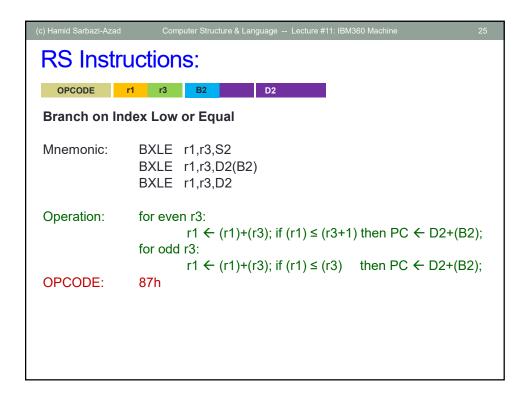


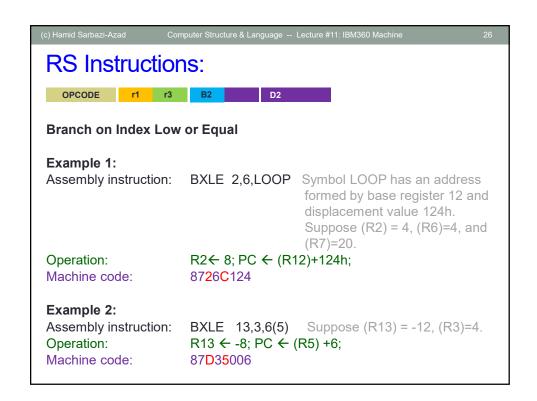


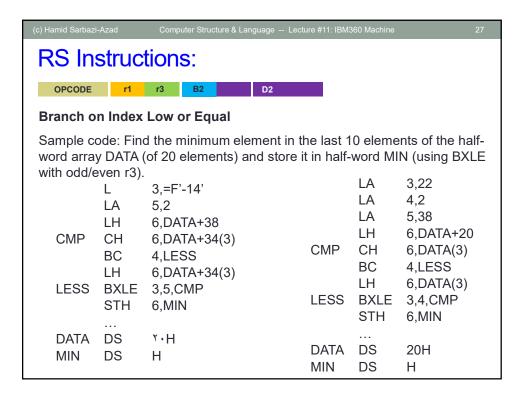


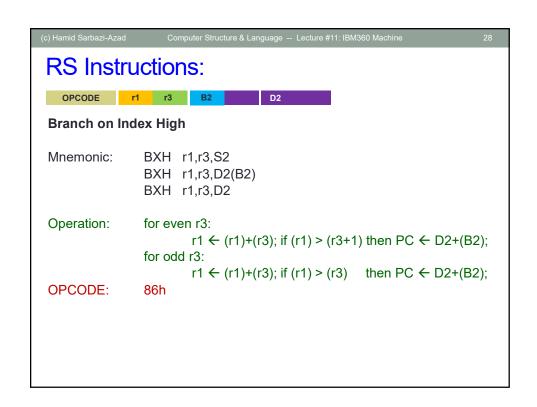


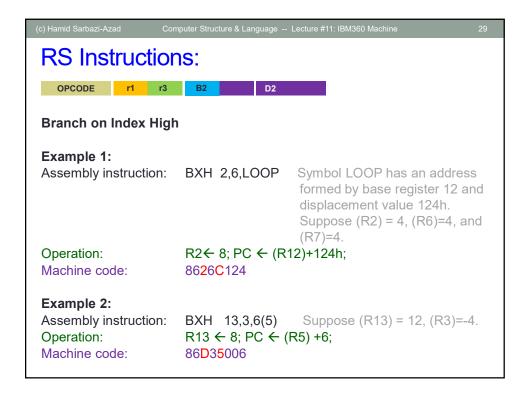


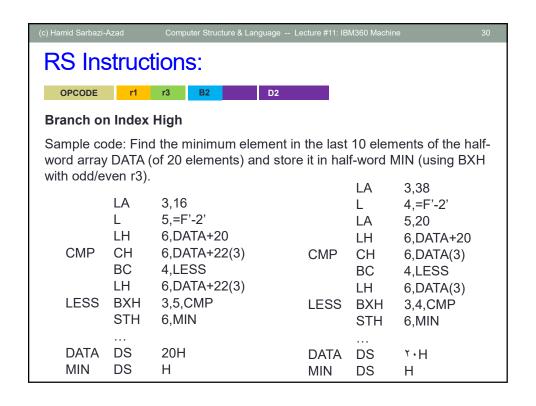












(c) Hamid Sarbazi-Azad

Computer Structure & Language -- Lecture #11: IBM360 Machine

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## **Binary Multiplication**

Recall the pen and paper technique we use to multiply n-bit number A and m-bit number B to generate (m+n)-bit number  $P = A \times B$ .

$$A = a_{n-1} \cdots a_1 a_0$$

$$B = b_{m-1} \cdots b_1 b_0$$

$$M_0 = a_{n-1} b_0 \cdots a_1 b_0 a_0 b_0$$

$$M_1 = a_{n-1} b_1 \cdots a_1 b_1 a_0 b_1$$

$$A = a_{n-1} b_1 \cdots a_1 b_1 a_0 b_1$$

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Computer Structure & Language -- Lecture #11: IBM360 Machin

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## **Binary Multiplication**

Multiplication can be done in a serial manner.

Consider n-bit numbers  $A(a_{n-1}a_{n-2}...a_0)$  and  $B(b_{n-1}b_{n-2}...b_0)$ .

We can write:

$$A \times B = (A \times 2^{n-1} \times b_{n-1}) + (A \times 2^{n-2} \times b_{n-2}) + \dots + (A \times 2^{1} \times b_{1}) + (A \times b_{0})$$

$$= [A \ll (n-1)] \times b_{n-1} + [A \ll (n-2)] \times b_{n-2} + \dots + [A \ll 1] \times b_{1} + [A \ll 0] \times b_{0}$$

$$= \sum_{i=0..n-1} [A \ll i] \times b_{i}$$

Shifting A for n-1 bits to the left we need a 2n-bit register to keep the partial products.

Alternatively, we can shift the accumulator, that accumulates the partial products, to the right.  $\rightarrow$  a tricky circuit to save hardware  $\odot$ 

