معماری کامپیوتر

فصل شش

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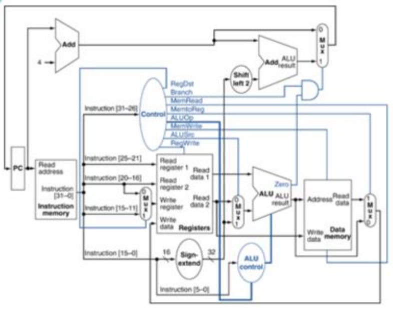
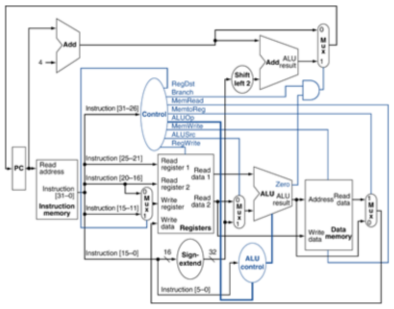


Computer Architecture

Chapter Six

Datapath & Control



Computer Architecture

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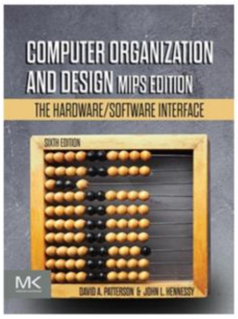
- D. Patterson & J. Hennessey, “Computer Organization & Design, The Hardware/Software Interface”, 6th Ed., MK publishing, 2021


COMPUTER ORGANIZATION AND DESIGN MIPS EDITION
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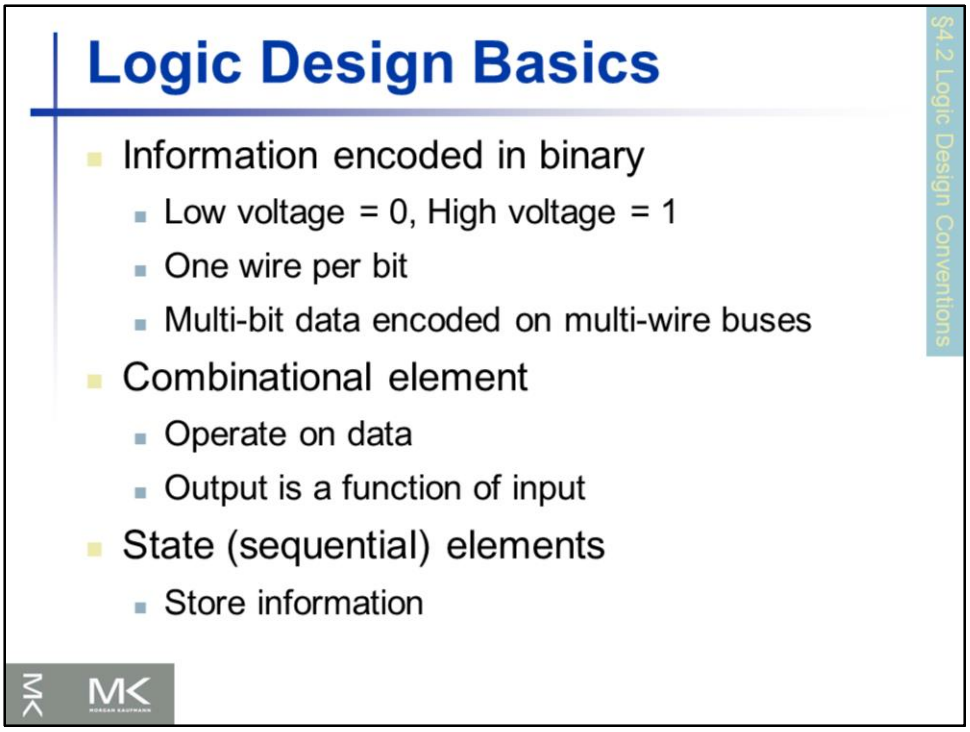
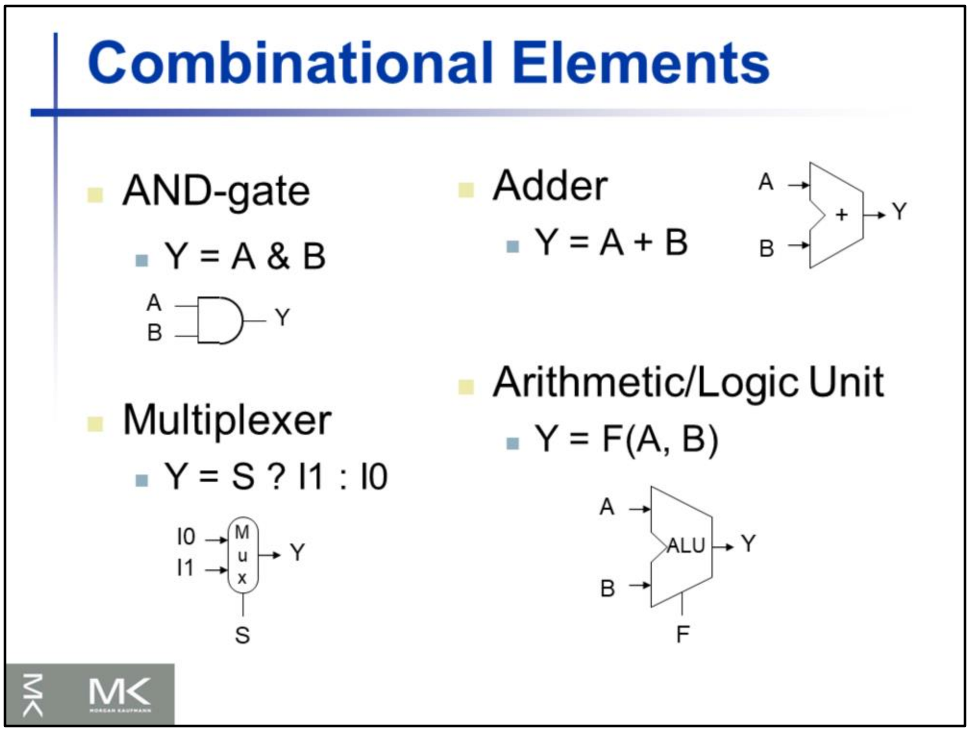


Logic Design Basics

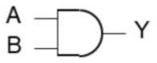
§4.2 Logic Design Conventions

- Information encoded in binary
  - Low voltage = 0, High voltage = 1
  - One wire per bit
  - Multi-bit data encoded on multi-wire buses
- Combinational element
  - Operate on data
  - Output is a function of input
- State (sequential) elements
  - Store information

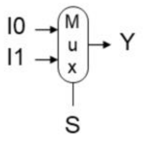


Combinational Elements

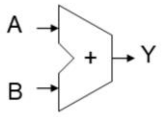
- AND-gate
  - Y = A & B



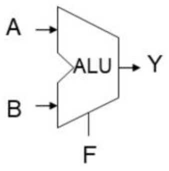
- Multiplexer
  - Y = S ? I1 : I0



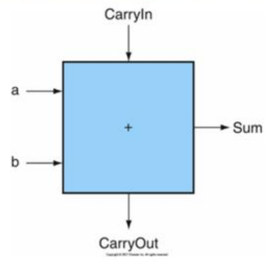
- Adder
  - Y = A + B



- Arithmetic/Logic Unit
  - Y = F(A, B)



1-bit Adder

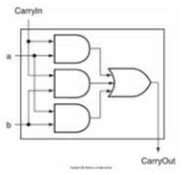


Inputs			Outputs		Comments
a	b	CarryIn	CarryOut	Sum	
0	0	0	0	0	$0+0+0=00_{two}$
0	0	1	0	1	$0+0+1=01_{two}$
0	1	0	0	1	$0+1+0=01_{two}$
0	1	1	1	0	$0+1+1=10_{two}$
1	0	0	0	1	$1+0+0=01_{two}$
1	0	1	1	0	$1+0+1=10_{two}$
1	1	0	1	0	$1+1+0=10_{two}$
1	1	1	1	1	$1+1+1=11_{two}$

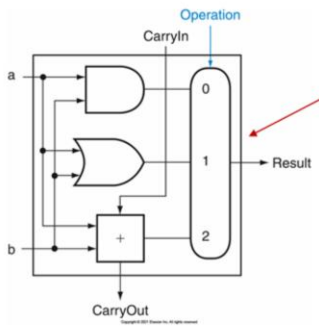
$$Sum = \left(a.\bar{b}.\overline{CarryIn}\right) + \left(\bar{a}.b.\overline{CarryIn}\right) + \left(\bar{a}.\bar{b}.CarryIn\right) + \left(\bar{a}.\bar{b}.\overline{CarryIn}\right)$$

Inputs		
a	b	CarryIn
0	1	1
1	0	1
1	1	0
1	1	1

Values of the inputs when CarryOut is a 1.



Simple ALU


Operation
CarryIn
a
b
0
1
2
+
Result
CarryOut
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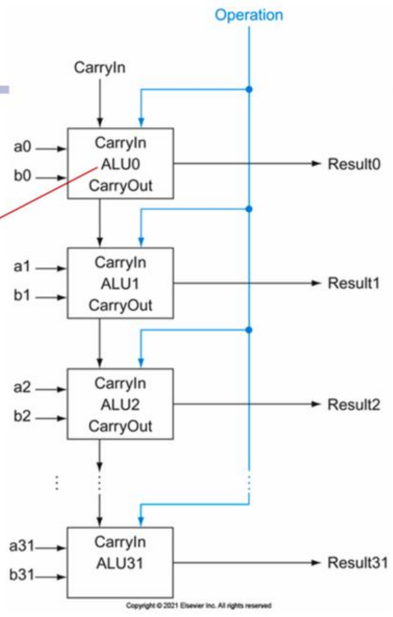
Operation
CarryIn
a0
b0
CarryIn
ALU0
CarryOut
Result0
a1
b1
CarryIn
ALU1
CarryOut
Result1
a2
b2
CarryIn
ALU2
CarryOut
Result2
a31
b31
CarryIn
ALU31
Result31
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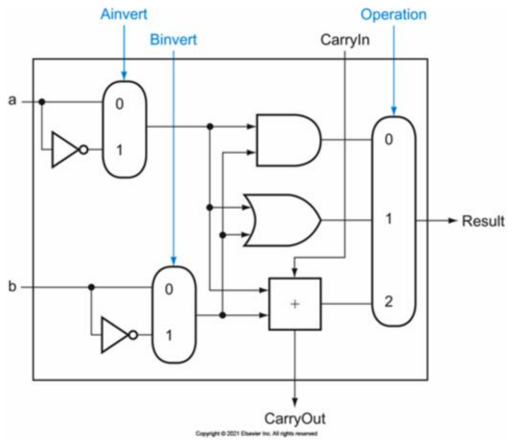
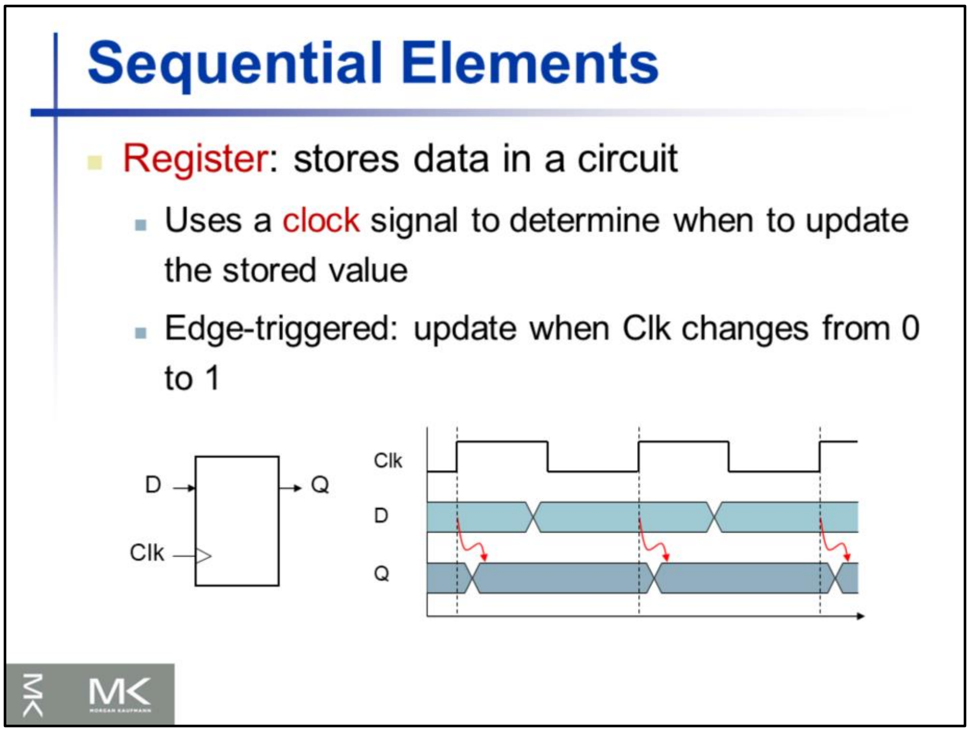
More Complex ALU

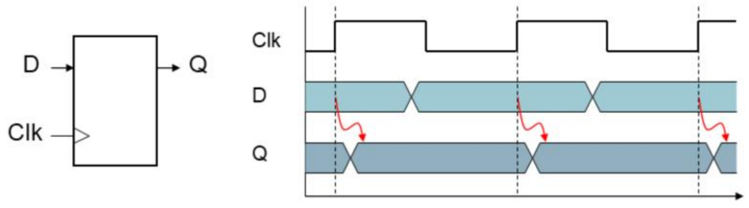


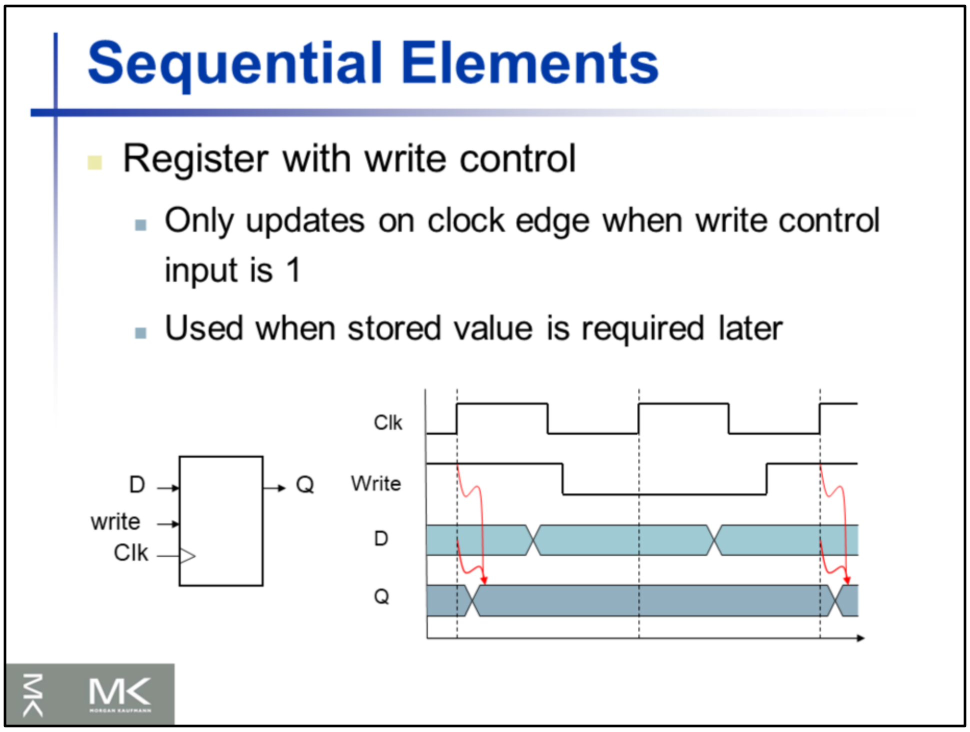
Find more in Appendix B



Sequential Elements

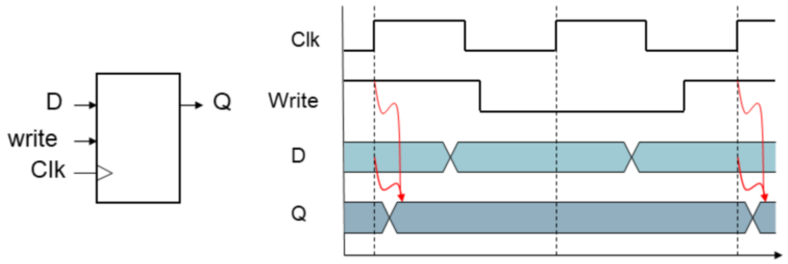
- Register: stores data in a circuit
  - Uses a clock signal to determine when to update the stored value
  - Edge-triggered: update when Clk changes from 0 to 1

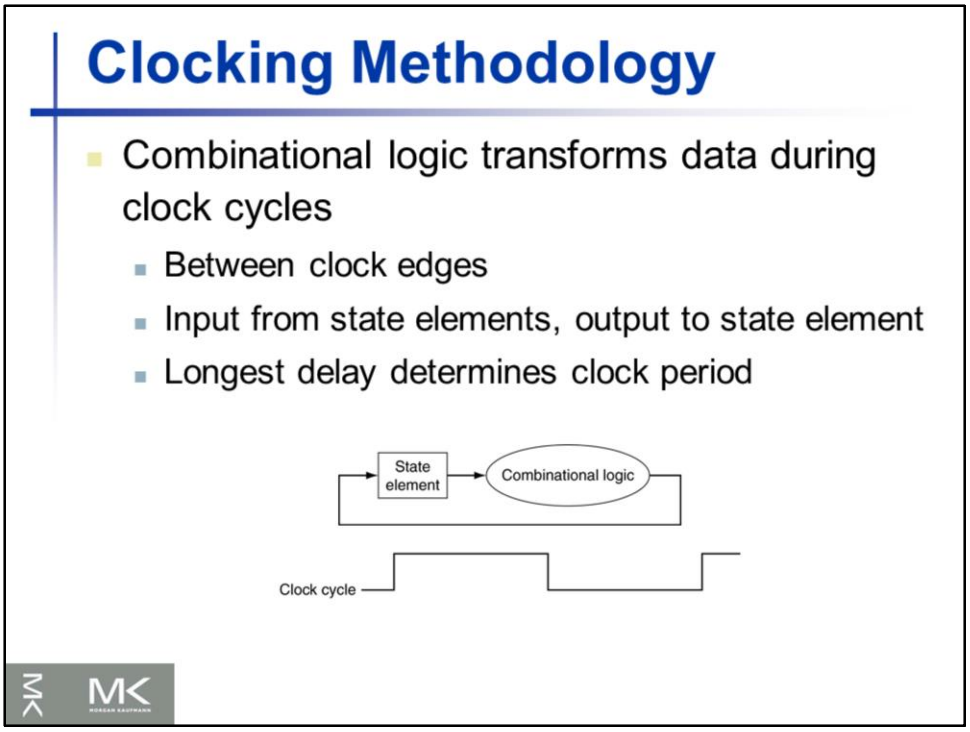




Sequential Elements

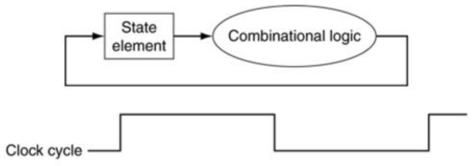
- Register with write control
  - Only updates on clock edge when write control input is 1
  - Used when stored value is required later

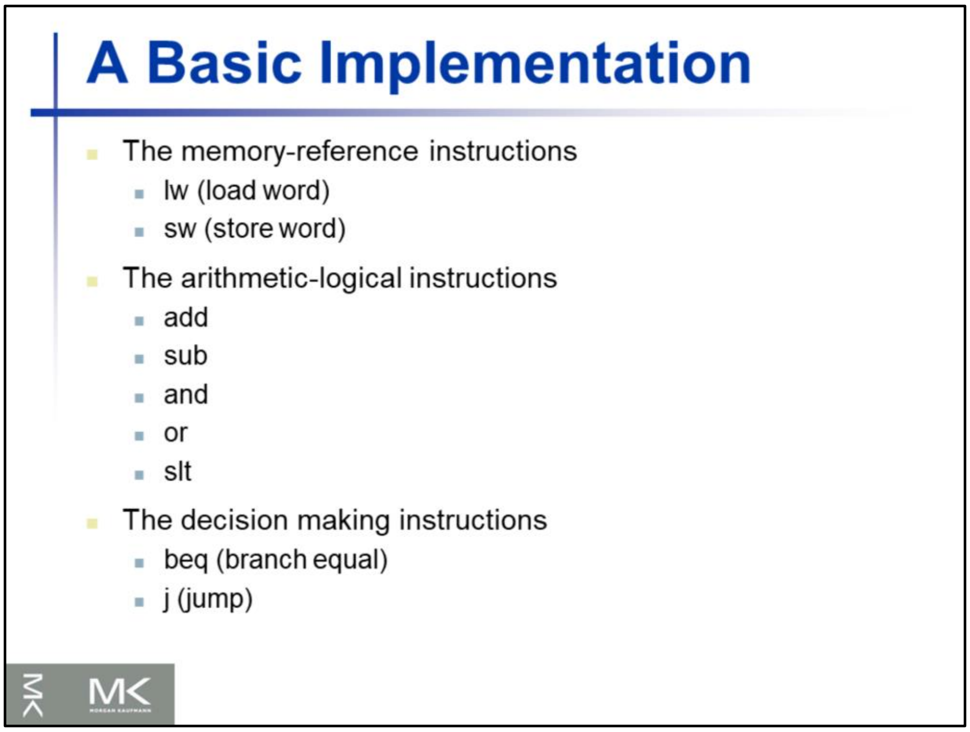




Clocking Methodology

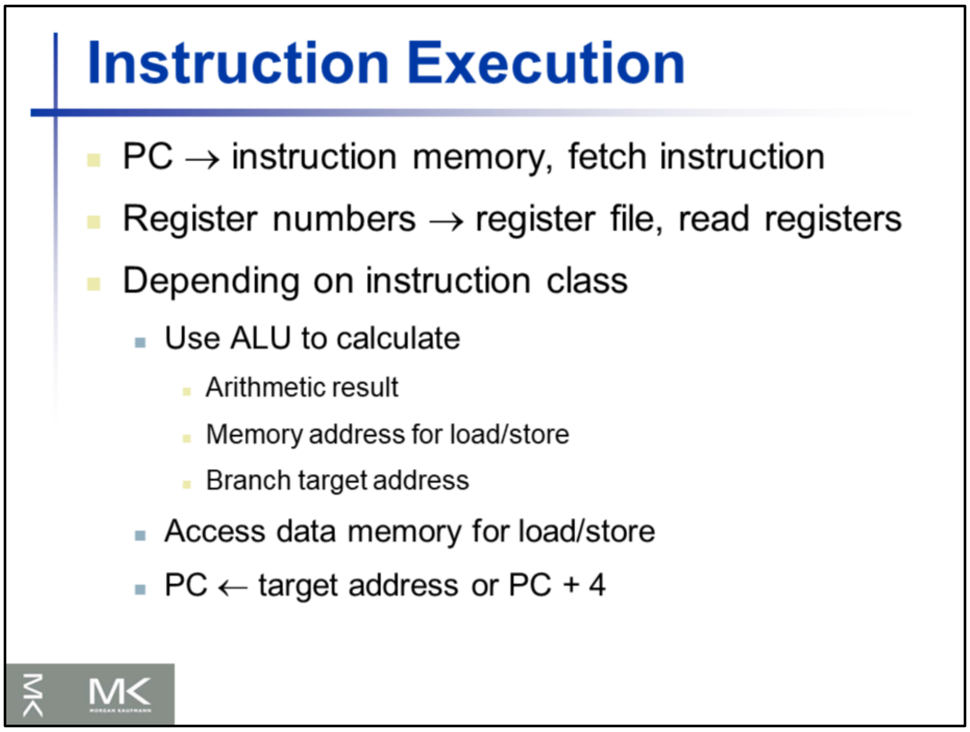
- Combinational logic transforms data during clock cycles
  - Between clock edges
  - Input from state elements, output to state element
  - Longest delay determines clock period





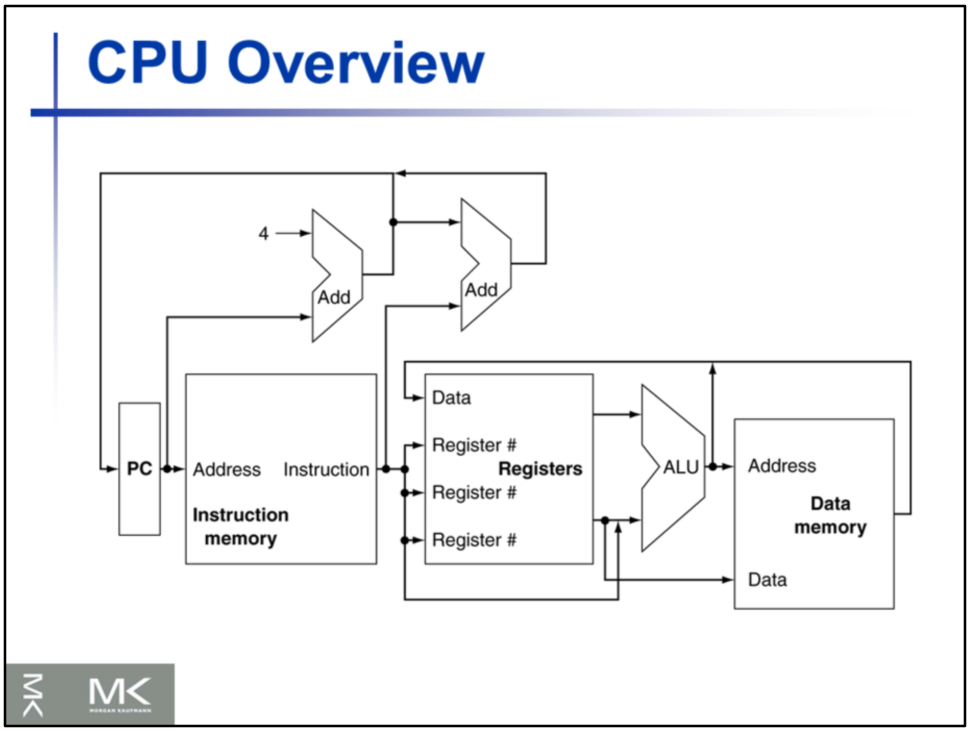
A Basic Implementation

- The memory-reference instructions
  - lw (load word)
  - sw (store word)
- The arithmetic-logical instructions
  - add
  - sub
  - and
  - or
  - slt
- The decision making instructions
  - beq (branch equal)
  - j (jump)



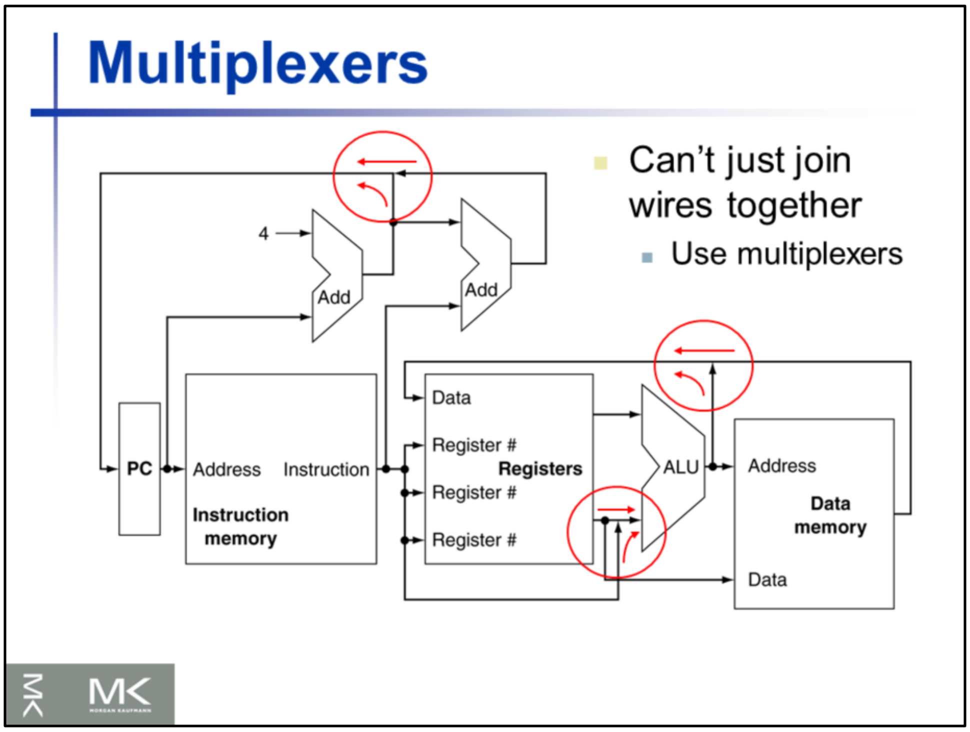
Instruction Execution

- PC → instruction memory, fetch instruction
- Register numbers → register file, read registers
- Depending on instruction class
  - Use ALU to calculate
    - Arithmetic result
    - Memory address for load/store
    - Branch target address
  - Access data memory for load/store
  - PC ← target address or PC + 4



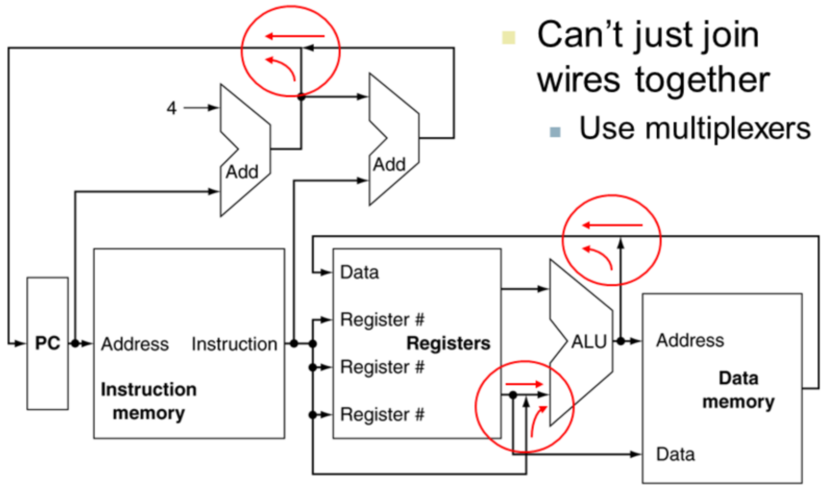
CPU Overview

FIGURE 4.1 An abstract view of the implementation of the MIPS subset showing the major functional units and the major connections between them.
All instructions start by using the program counter to supply the instruction address to the instruction memory. After the instruction is fetched, the register operands used by an instruction are specified by fields of that instruction. Once the register operands have been fetched, they can be operated on to compute a memory address (for a load or store), to compute an arithmetic result (for an integer arithmetic-logical instruction), or a compare (for a branch). If the instruction is an arithmetic-logical instruction, the result from the ALU must be written to a register. If the operation is a load or store, the ALU result is used as an address to either load a value from memory into the registers or store a value from the registers. The result from the ALU or memory is written back into the register file. Branches require the use of the ALU output to determine the next instruction address, which comes either from the ALU (where the PC and branch offset are summed) or from an adder that increments the current PC by 4.



Multiplexers

- Can't just join wires together
  - Use multiplexers



Control

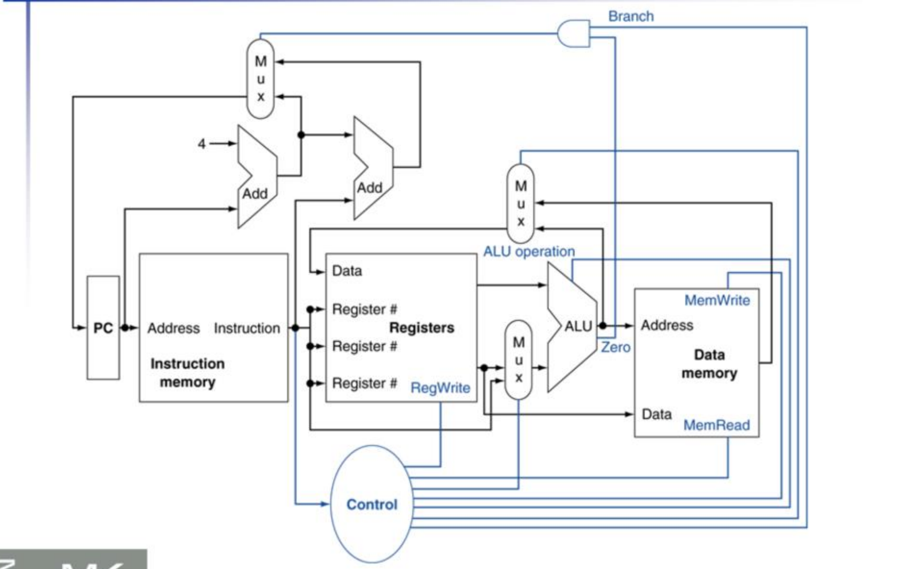
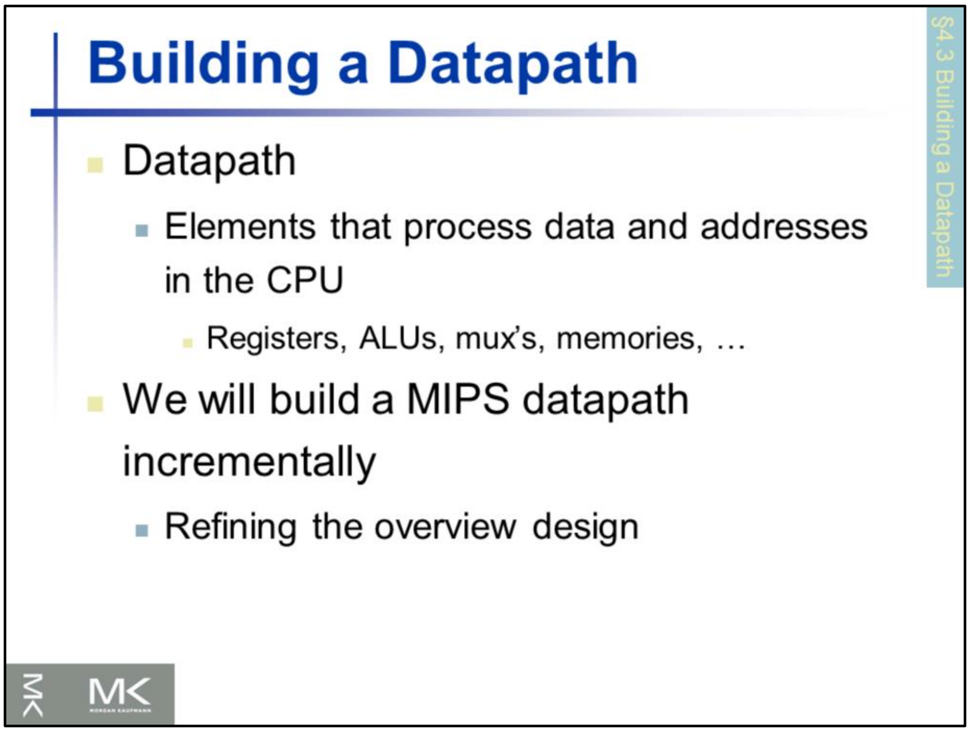


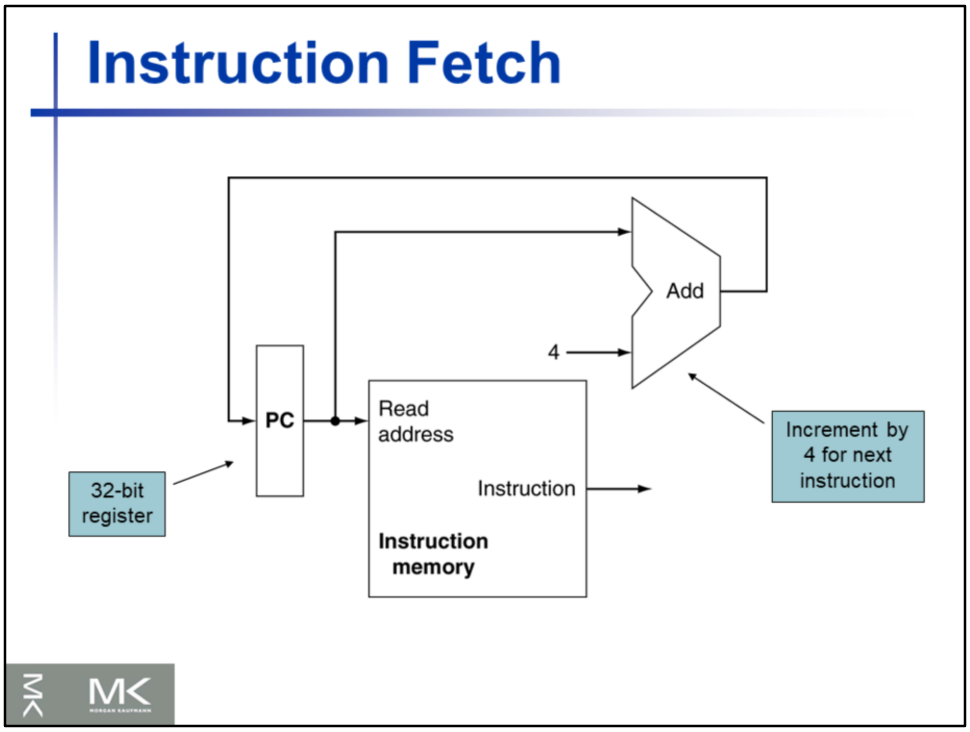
FIGURE 4.2 The basic implementation of the MIPS subset, including the necessary multiplexors and control lines.
The top multiplexor ("Mux") controls what value replaces the PC (PC + 4 or the branch destination address); the
multiplexor is controlled by the gate that "ANDs" together the Zero output of the ALU and a control signal that indicates that
the instruction is a branch. The middle multiplexor, whose output returns to the register file, is used to steer the output
of the ALU (in the case of an arithmetic-logical instruction) or the output of the data memory (in the case of a load) for
writing into the register file. Finally, the bottommost multiplexor is used to determine whether the second ALU input is from the registers (for an arithmetic-logical instruction or a branch) or from the offset field of the instruction (for a load or store).
The added control lines are straightforward and determine the operation performed at the ALU, whether the data memory should read or write, and whether the registers should perform a write operation. The control lines are shown in color to make them easier to see.



Building a Datapath

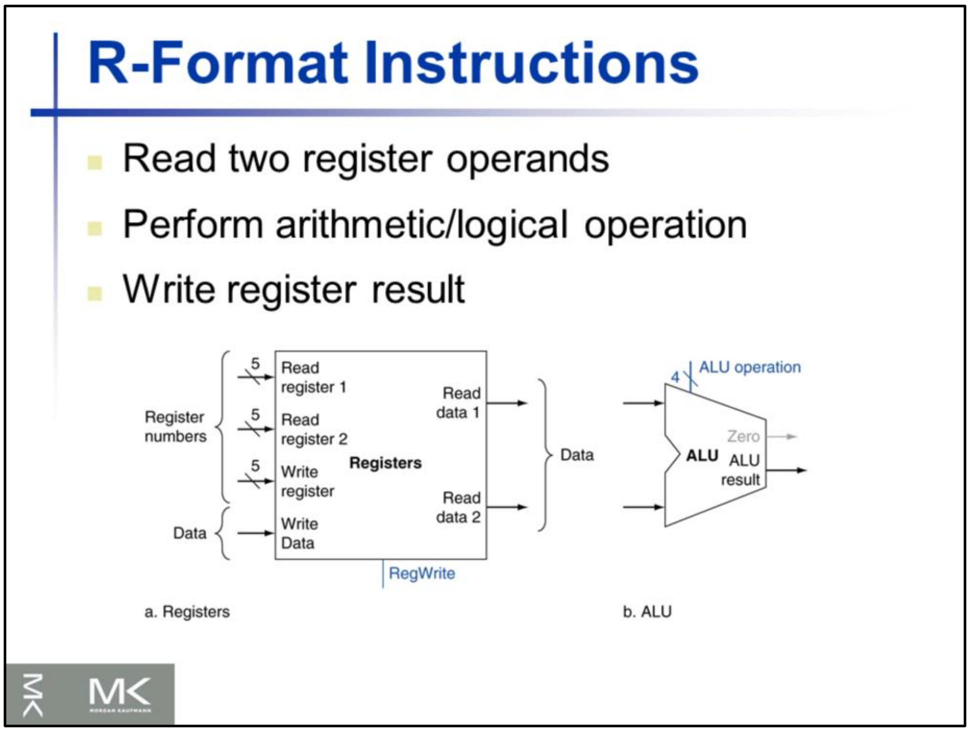
§4.3 Building a Datapath

- Datapath
  - Elements that process data and addresses in the CPU
    - Registers, ALUs, mux's, memories, …
- We will build a MIPS datapath incrementally
  - Refining the overview design



Instruction Fetch

FIGURE 4.6 A portion of the datapath used for fetching instructions and incrementing the program counter.
The fetched instruction is used by other parts of the datapath. The state elements are the instruction memory and the program counter.
The instruction memory need only provide read access because the datapath does not write instructions. Since the instruction memory only reads, we treat it as combinational logic: the output at any time reflects the contents of the location specified by the address input, and no read control signal is needed. (We will need to write the instruction memory when we load the program; this is not hard to add, and we ignore it for simplicity.)
The program counter is a 32-bit register that is written at the end of every clock cycle and thus does not need a write control signal. The adder is an ALU wired to always add its two 32-bit inputs and place the sum on its output.



R-Format Instructions

- Read two register operands
- Perform arithmetic/logical operation
- Write register result

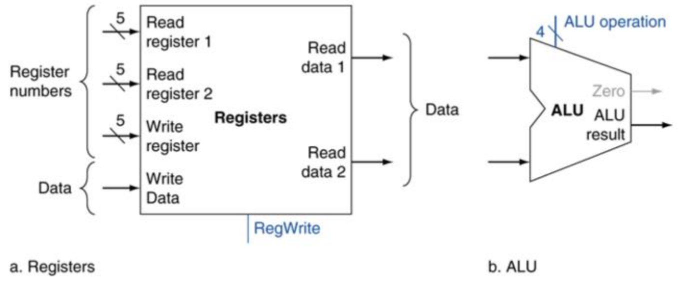
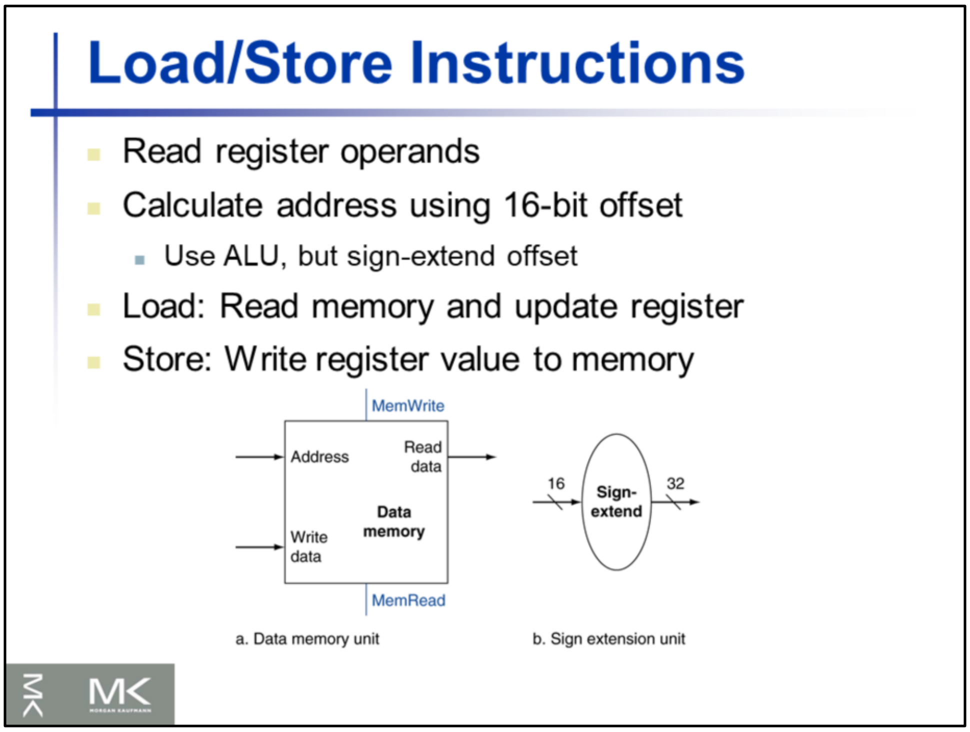


FIGURE 4.7 The two elements needed to implement R-format ALU operations are the register file and the ALU.
The register file contains all the registers and has two read ports and one write port. The design of multiported register files is discussed in Section B.8 of **Appendix B**. The register file always outputs the contents of the registers corresponding to the Read register inputs on the outputs; no other control inputs are needed. In contrast, a register write must be explicitly indicated by asserting the write control signal. Remember that writes are edge-triggered, so that all the write inputs (i.e., the value to be written, the register number, and the write control signal) must be valid at the clock edge. Since writes to the register file are edge-triggered, our design can legally read and write the same register within a clock cycle: the read will get the value written in an earlier clock cycle, while the value written will be
available to a read in a subsequent clock cycle. The inputs carrying the register number to the register file are all 5 bits wide, whereas the lines carrying data values are 32 bits wide.
The operation to be performed by the ALU is controlled with the ALU operation signal, which will be 4 bits wide, using the ALU designed in **Appendix B**. We will use the Zero detection output of the ALU shortly to implement branches. The overflow output will not be needed until Section 4.10, when we discuss exceptions; we omit it until then.

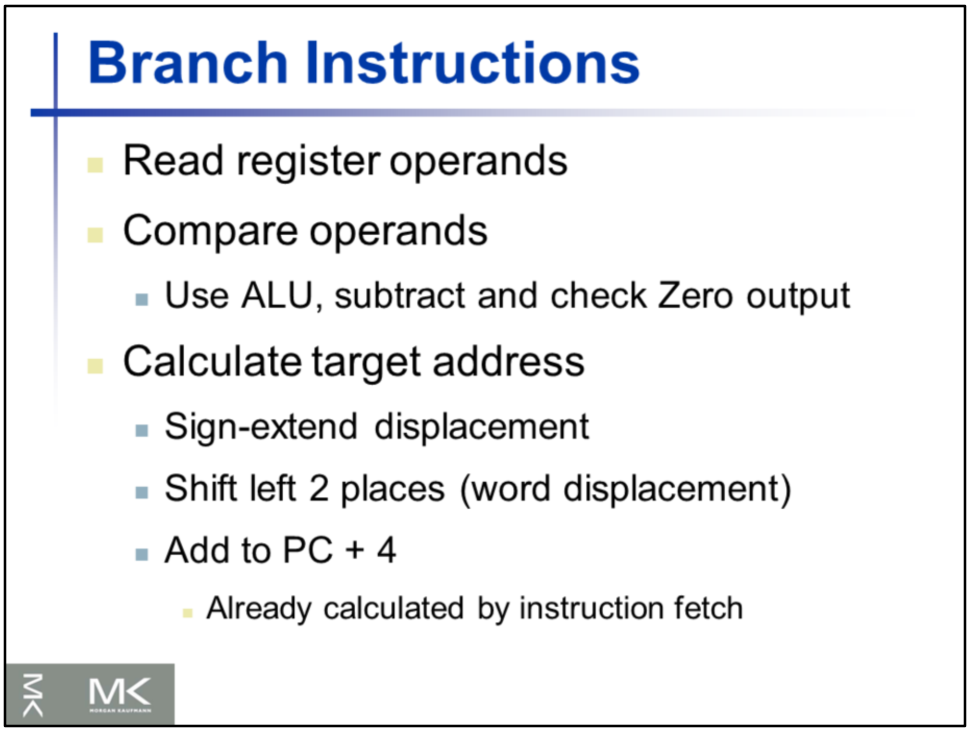


Load/Store Instructions

- Read register operands
- Calculate address using 16-bit offset
  - Use ALU, but sign-extend offset
- Load: Read memory and update register
- Store: Write register value to memory

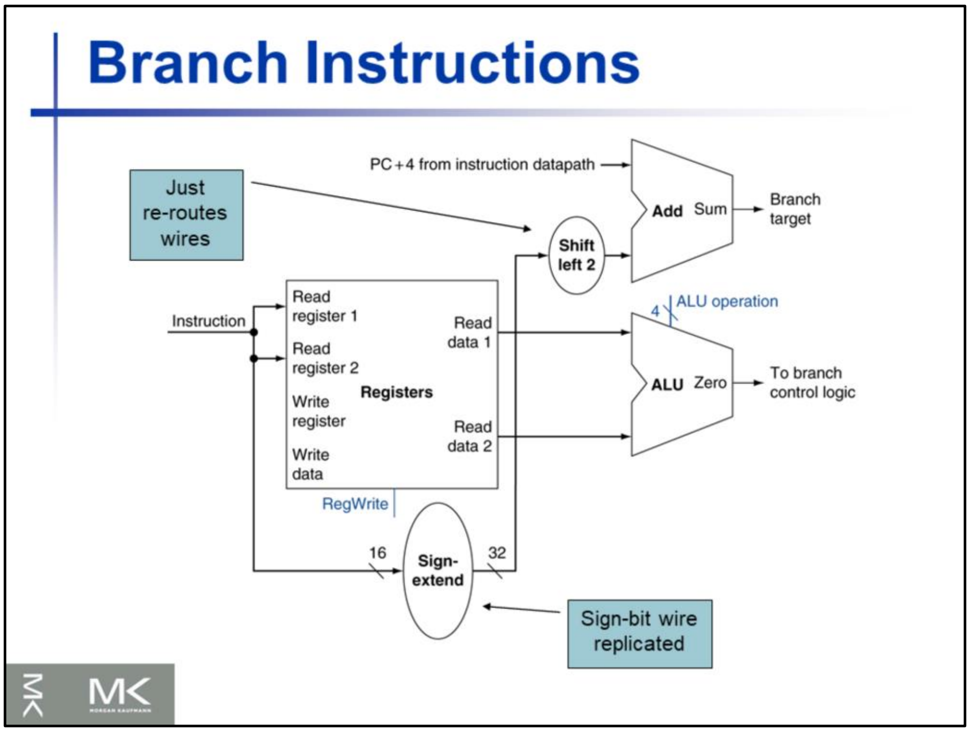
MemWrite
Address
Read data
Data memory
Write data
MemRead
a. Data memory unit
16
Sign-extend
32
b. Sign extension unit

FIGURE 4.8 The two units needed to implement loads and stores, in addition to the register file and ALU of Figure 4.7, are the data memory unit and the sign extension unit.
The memory unit is a state element with inputs for the address and the write data, and a single output for the read result. There are separate read and write controls, although only one of these may be asserted on any given clock. The memory unit needs a read signal, since, unlike the register file, reading the value of an invalid address can cause problems, as we will see in Chapter 5. The sign extension unit has a 16-bit input that is sign-extended into a 32-bit result appearing on the output (see Chapter 2). We assume the data memory is edge-triggered for writes. Standard memory chips actually have a write enable signal that is used for writes. Although the write enable is not edge-triggered, our edge-triggered design could easily be adapted to work with real memory chips. See Section B.8 of **Appendix B** for further discussion of how real memory chips work.



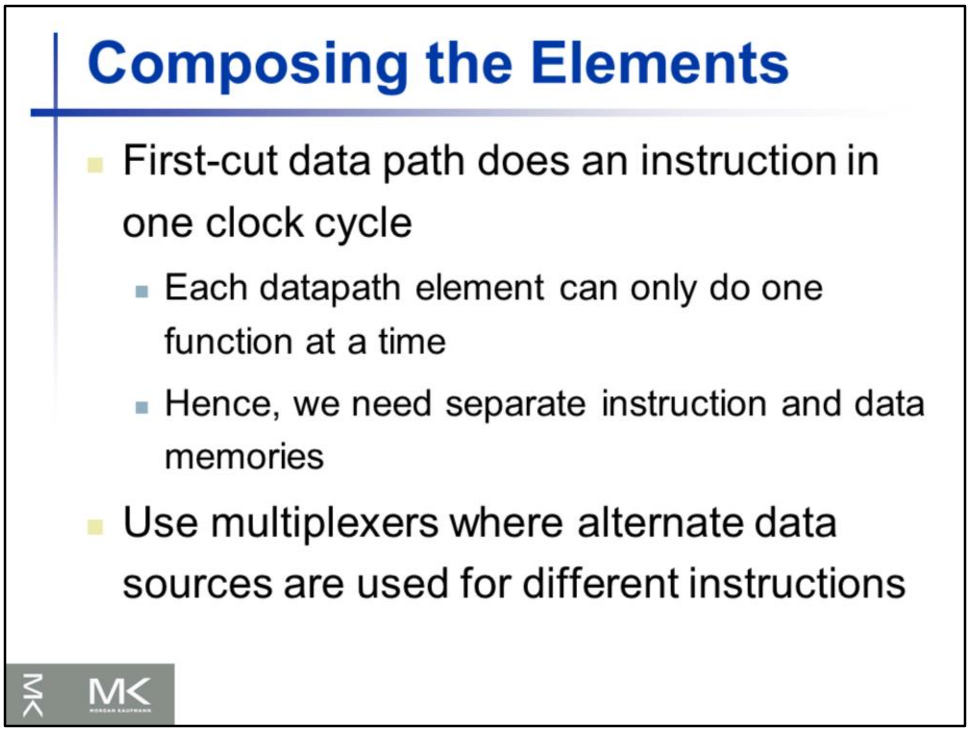
Branch Instructions

- Read register operands
- Compare operands
  - Use ALU, subtract and check Zero output
- Calculate target address
  - Sign-extend displacement
  - Shift left 2 places (word displacement)
  - Add to PC + 4
    - Already calculated by instruction fetch



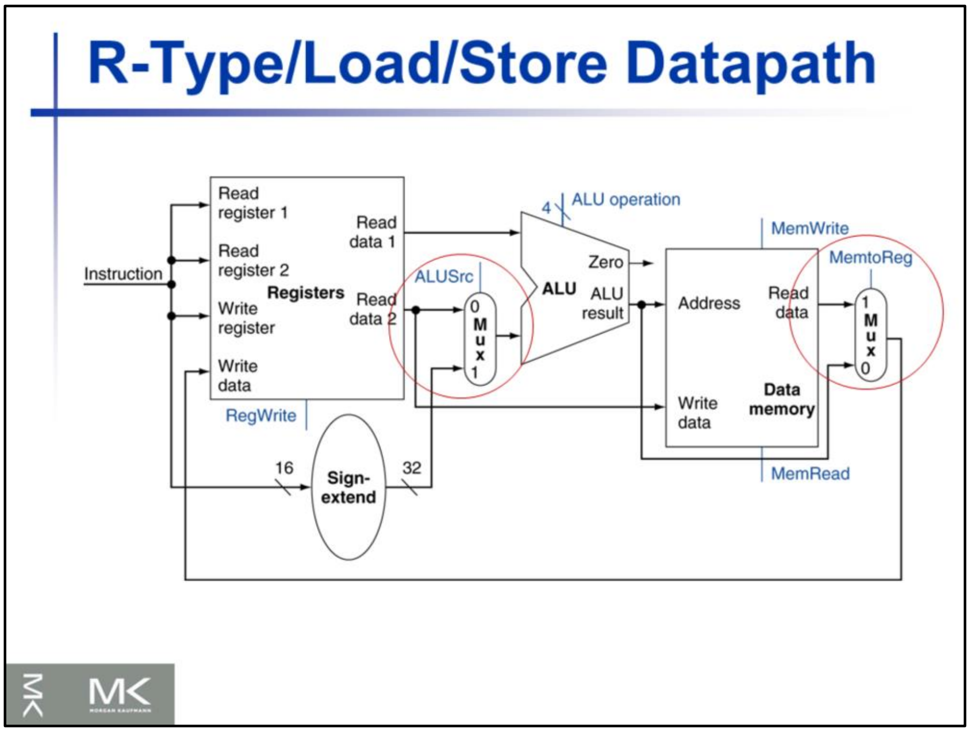
Branch Instructions

FIGURE 4.9 The portion of a datapath for a branch uses the ALU to evaluate the branch condition and a separate adder to compute the branch target as the sum of the incremented PC and the sign-extended, lower 16 bits of the instruction (the branch displacement), shifted left 2 bits.
The unit labeled *Shift left 2* is simply a routing of the signals between input and output that adds 00_{two} to the low-order end of the sign-extended offset field; no actual shift hardware is needed, since the amount of the "shift" is constant. Since we know that the offset was sign-extended from 16 bits, the shift will throw away only "sign bits." Control logic is used to decide whether the incremented PC or branch target should replace the PC, based on the Zero output of the ALU.



Composing the Elements

- First-cut data path does an instruction in one clock cycle
  - Each datapath element can only do one function at a time
  - Hence, we need separate instruction and data memories
- Use multiplexers where alternate data sources are used for different instructions



R-Type/Load/Store Datapath

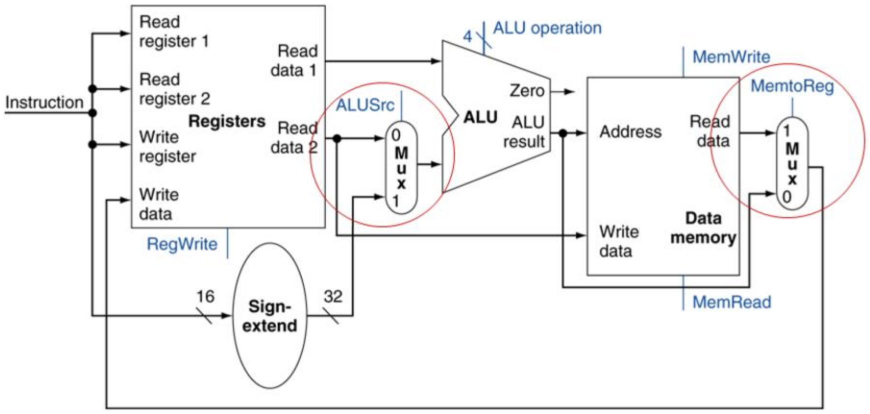


FIGURE 4.10 The datapath for the memory instructions and the R-type instructions.
This example shows how a single datapath can be assembled from the pieces in Figures 4.7 and 4.8 by adding multiplexors.

Full Datapath

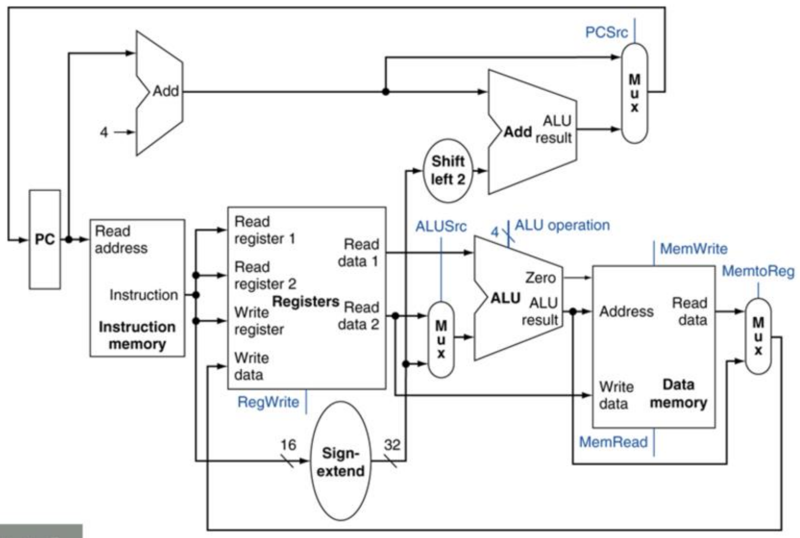
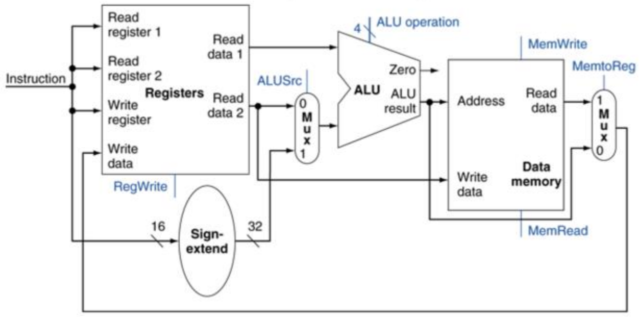
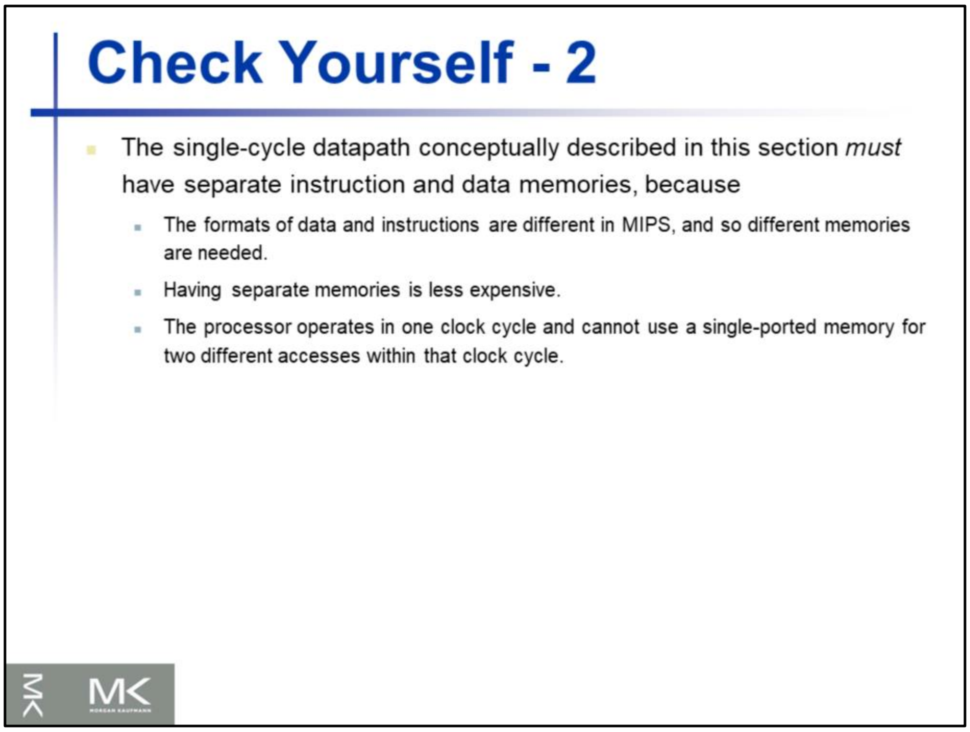


FIGURE 4.11 The simple datapath for the core MIPS architecture combines the elements required by different instruction classes.
The components come from Figures 4.6, 4.9, and 4.10. This datapath can execute the basic instructions (load-store word, ALU operations, and branches) in a single clock cycle. Just one additional multiplexor is needed to integrate branches.
The support for jumps will be added later.

Check Yourself - 1

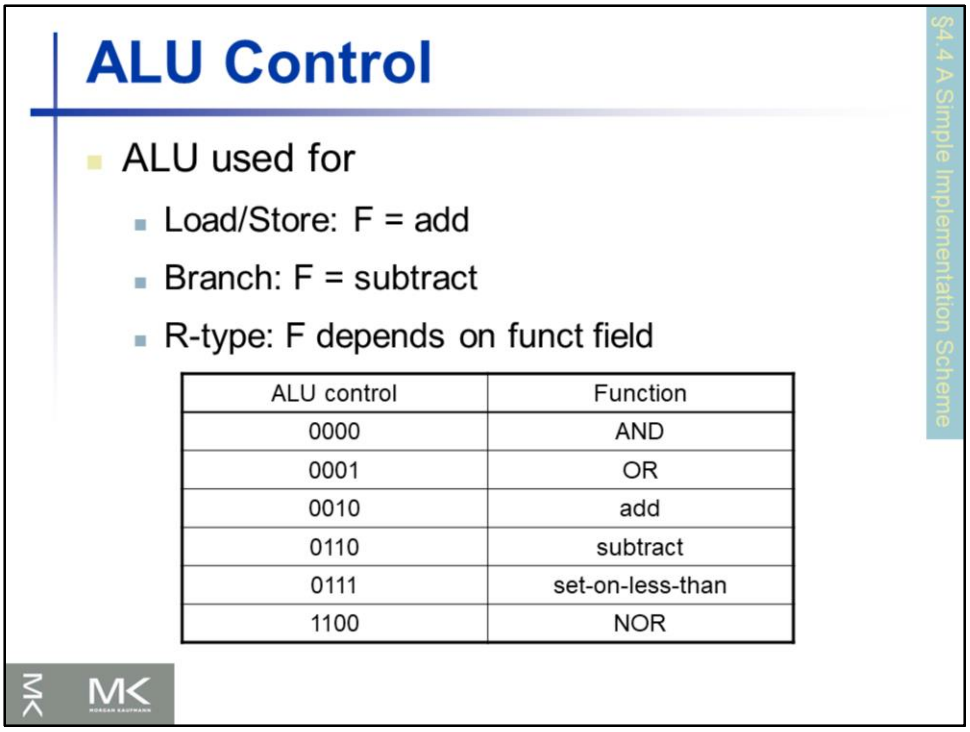
- Which of the following is correct for a load instruction?
  - MemtoReg should be set to cause the data from memory to be sent to the register file.
  - MemtoReg should be reset to cause the correct register destination to be sent to the register file.
  - We do not care about the setting of MemtoReg for loads.





Check Yourself - 2

- The single-cycle datapath conceptually described in this section *must* have separate instruction and data memories, because
  - The formats of data and instructions are different in MIPS, and so different memories are needed.
  - Having separate memories is less expensive.
  - The processor operates in one clock cycle and cannot use a single-ported memory for two different accesses within that clock cycle.

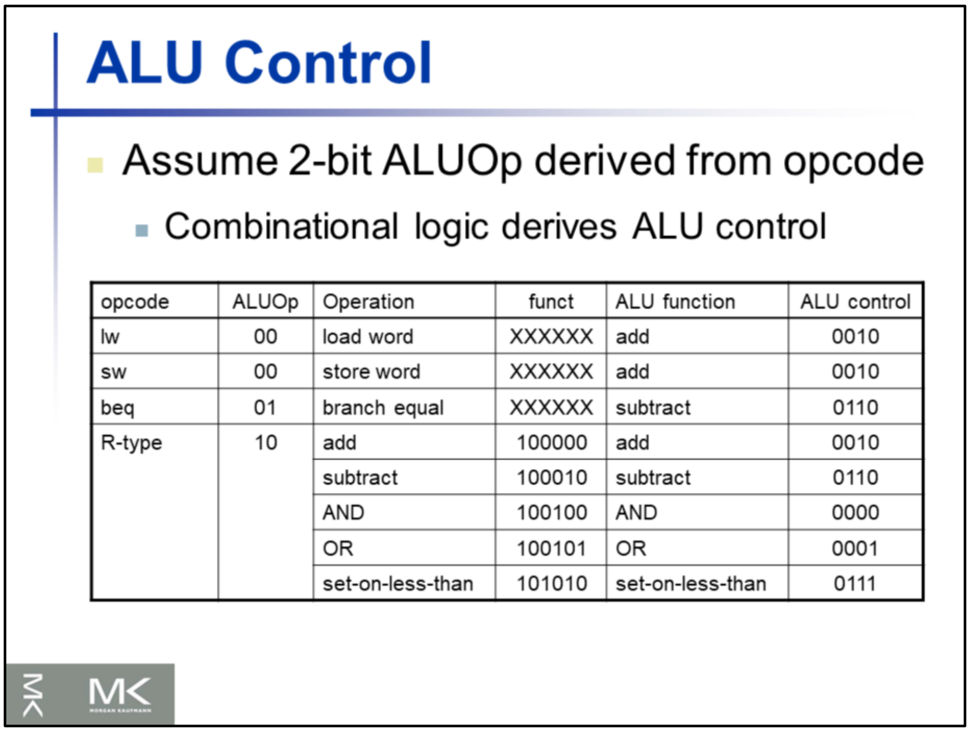


ALU Control

§4.4 A Simple Implementation Scheme

- ALU used for
  - Load/Store: F = add
  - Branch: F = subtract
  - R-type: F depends on funct field

ALU control	Function
0000	AND
0001	OR
0010	add
0110	subtract
0111	set-on-less-than
1100	NOR

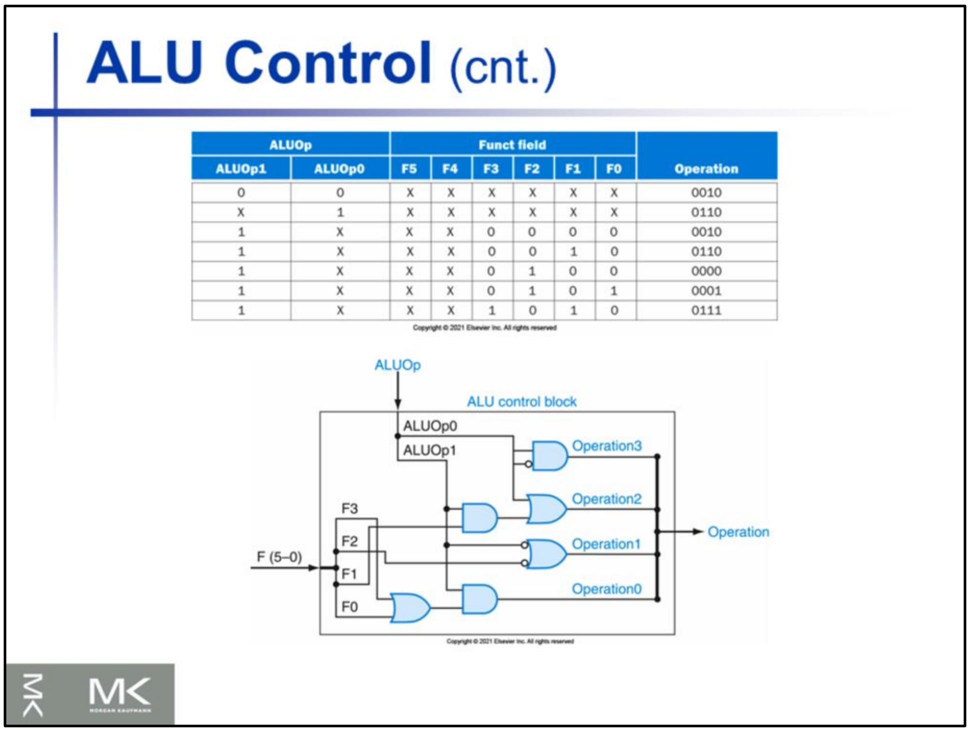


ALU Control

- Assume 2-bit ALUOp derived from opcode
  - Combinational logic derives ALU control

opcode	ALUOp	Operation	funct	ALU function	ALU control
lw	00	load word	XXXXXX	add	0010
sw	00	store word	XXXXXX	add	0010
beq	01	branch equal	XXXXXX	subtract	0110
R-type	10	add	100000	add	0010
		subtract	100010	subtract	0110
		AND	100100	AND	0000
		OR	100101	OR	0001
		set-on-less-than	101010	set-on-less-than	0111

FIGURE 4.12 How the ALU control bits are set depends on the ALUOp control bits and the different function codes for the R-type instruction. The opcode, listed in the first column, determines the setting of the ALUOp bits. All the encodings are shown in binary.
Notice that when the ALUOp code is 00 or 01, the desired ALU action does not depend on the function code field; in this case, we say that we “don’t care” about the value of the function code, and the funct field is shown as XXXXXX. When the ALUOp value is 10, then the function code is used ALUOp indicates whether the operation to be performed should be add (00) for loads and stores, subtract (01) for beq, or determined by the operation encoded in the funct field (10).



ALU Control (cnt.)

ALUOp		Funct field						
ALUOp1	ALUOp0	F5	F4	F3	F2	F1	F0	Operation
0	0	X	X	X	X	X	X	0010
X	1	X	X	X	X	X	X	0110
1	X	X	X	0	0	0	0	0010
1	X	X	X	0	0	1	0	0110
1	X	X	X	0	1	0	0	0000
1	X	X	X	0	1	0	1	0001
1	X	X	X	1	0	1	0	0111

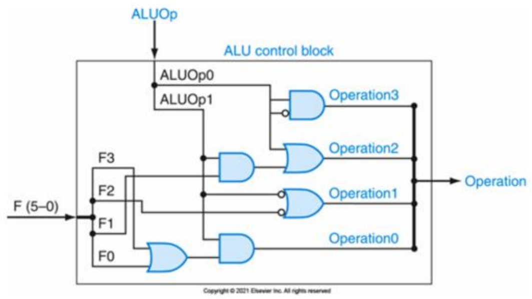
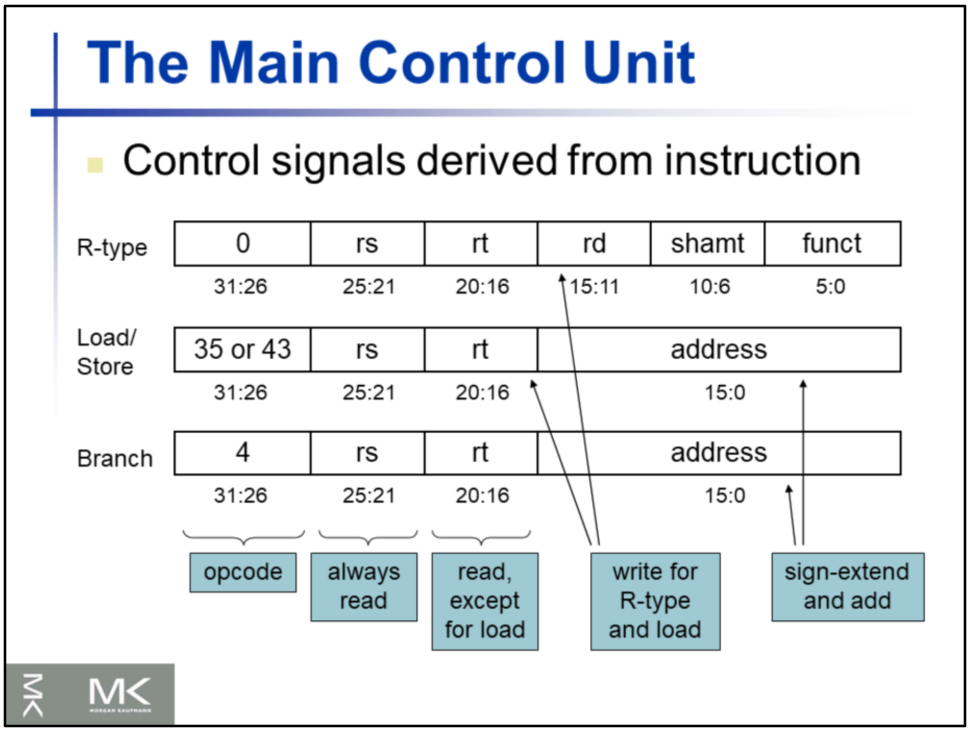


FIGURE 4.13 The truth table for the 4 ALU control bits (called Operation). The inputs are the ALUOp and function code field. Only the entries for which the ALU control is asserted are shown. Some don't-care entries have been added. For example, the ALUOp does not use the encoding 11, so the truth table can contain entries 1X and X1, rather than 10 and 01. Note that when the function field is used, the first 2 bits (F5 and F4) of these instructions are always 10, so they are don't-care terms and are replaced with XX in the truth table.



The Main Control Unit

- Control signals derived from instruction

FIGURE 4.14 The three instruction classes (R-type, load and store, and branch) use two different instruction formats.

(a) Instruction format for R-format instructions, which all have an opcode of 0. These instructions have three register operands: rs, rt, and rd. Fields rs and rt are sources, and rd is the destination. The ALU function is in the funct field and is decoded by the ALU control design in the previous section. The R-type instructions that we implement are add, sub, AND, OR, and slt. The shamt field is used only for shifts; we will ignore it in this chapter.

(b) Instruction format for load (opcode = 35_{ten}) and store (opcode = 43_{ten}) instructions. The register rs is the base register that is added to the 16-bit address field to form the memory address. For loads, rt is the destination register for the loaded value. For stores, rt is the source register whose value should be stored into memory.

(c) Instruction format for branch equal (opcode = 4). The registers rs and rt are the source registers that are compared for equality. The 16-bit address field is sign-extended, shifted, and added to the PC + 4 to compute the branch target address.

Datapath With Control

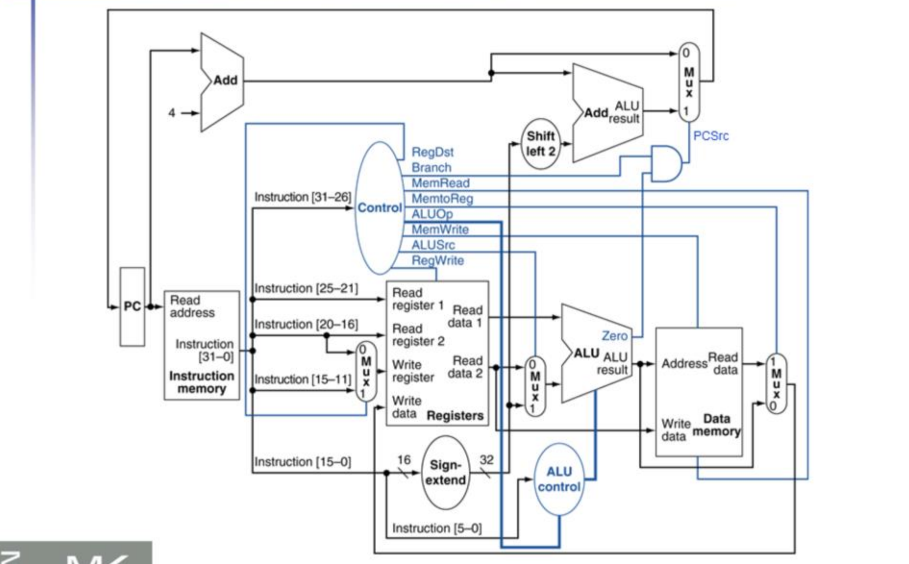
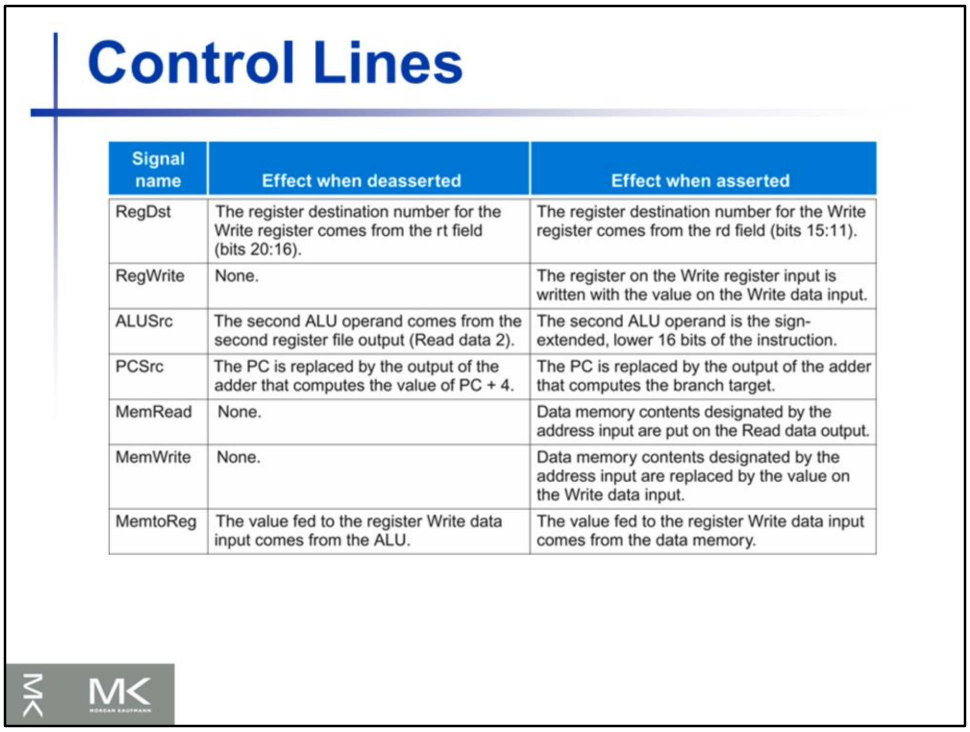


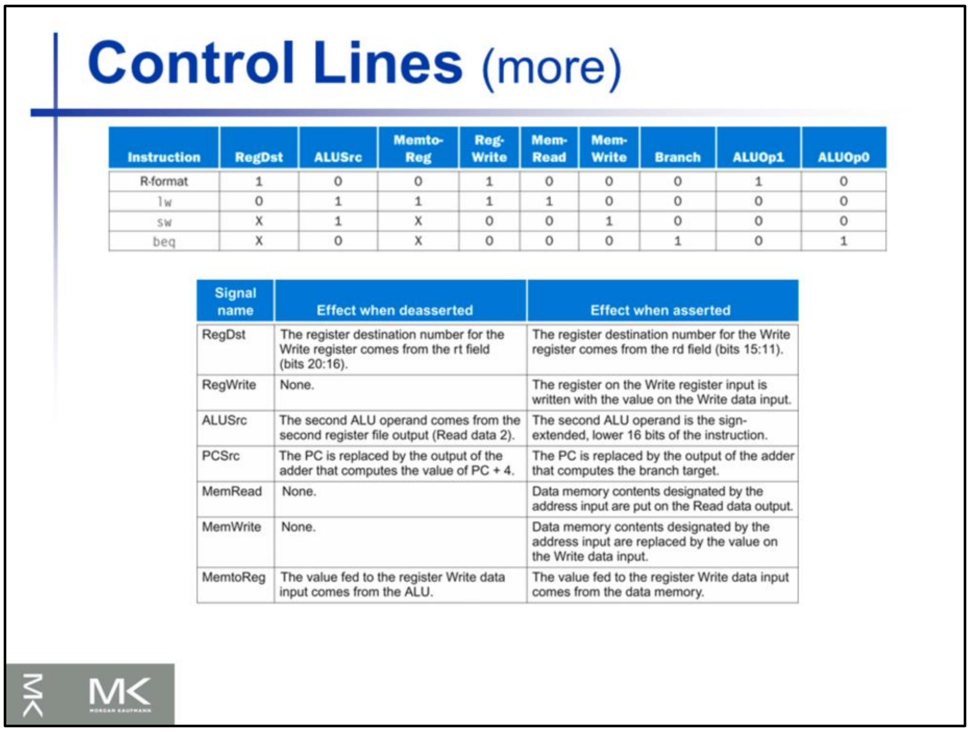
FIGURE 4.17 The simple datapath with the control unit.
The input to the control unit is the 6-bit opcode field from the instruction. The outputs of the control unit consist of three 1-bit signals that are used to control multiplexors (RegDst, ALUSrc, and MemtoReg), three signals for controlling reads and writes in the register file and data memory (RegWrite, MemRead, and MemWrite), a 1-bit signal used in determining whether to possibly branch (Branch), and a 2-bit control signal for the ALU (ALUOp).
An AND gate is used to combine the branch control signal and the Zero output from the ALU; the AND gate output controls the selection of the next PC. Notice that PCSrc is now a derived signal, rather than one coming directly from the control unit. Thus, we drop the signal name in subsequent figures.



Control Lines

Signal name	Effect when deasserted	Effect when asserted
RegDst	The register destination number for the Write register comes from the rt field (bits 20:16).	The register destination number for the Write register comes from the rd field (bits 15:11).
RegWrite	None.	The register on the Write register input is written with the value on the Write data input.
ALUSrc	The second ALU operand comes from the second register file output (Read data 2).	The second ALU operand is the sign-extended, lower 16 bits of the instruction.
PCSrc	The PC is replaced by the output of the adder that computes the value of PC + 4.	The PC is replaced by the output of the adder that computes the branch target.
MemRead	None.	Data memory contents designated by the address input are put on the Read data output.
MemWrite	None.	Data memory contents designated by the address input are replaced by the value on the Write data input.
MemtoReg	The value fed to the register Write data input comes from the ALU.	The value fed to the register Write data input comes from the data memory.

FIGURE 4.16 The effect of each of the seven control signals.
When the 1-bit control to a two-way multiplexor is asserted, the multiplexor selects the input corresponding to 1.
Otherwise, if the control is deasserted, the multiplexor selects the 0 input.
Remember that the state elements all have the clock as an implicit input and that the clock is used in controlling writes. Gating the clock externally to a state element can create timing problems.

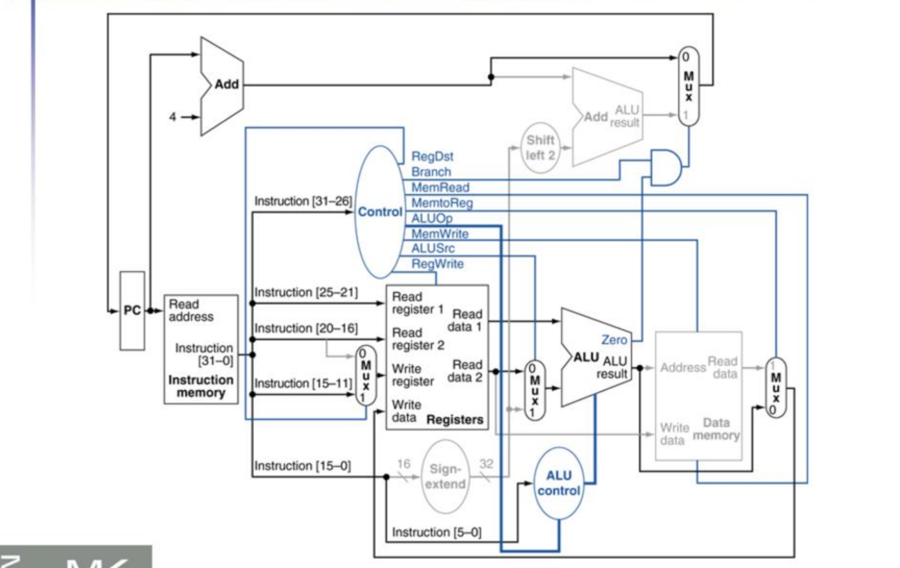


Control Lines (more)

Instruction	RegDst	ALUSrc	Memto-Reg	Reg-Write	Mem-Read	Mem-Write	Branch	ALUOp1	ALUOp0
R-format	1	0	0	1	0	0	0	1	0
lw	0	1	1	1	1	0	0	0	0
sw	X	1	X	0	0	1	0	0	0
beq	X	0	X	0	0	0	1	0	1

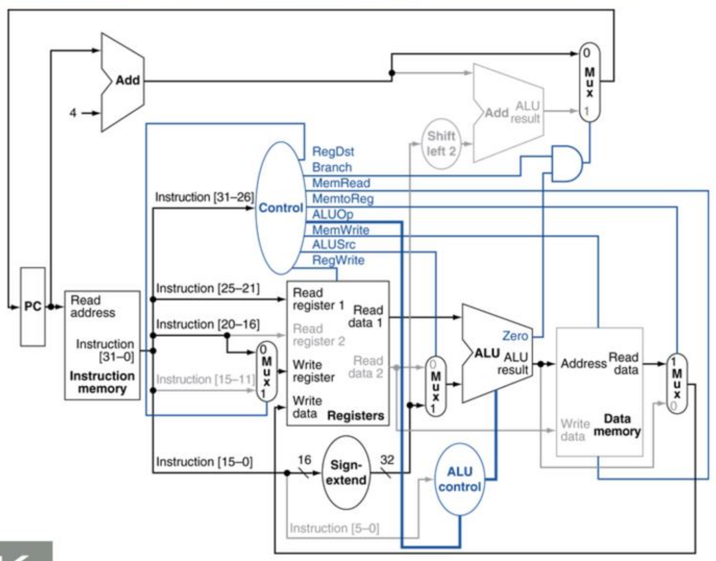
Signal name	Effect when deasserted	Effect when asserted
RegDst	The register destination number for the Write register comes from the rt field (bits 20:16).	The register destination number for the Write register comes from the rd field (bits 15:11).
RegWrite	None.	The register on the Write register input is written with the value on the Write data input.
ALUSrc	The second ALU operand comes from the second register file output (Read data 2).	The second ALU operand is the sign-extended, lower 16 bits of the instruction.
PCSrc	The PC is replaced by the output of the adder that computes the value of PC + 4.	The PC is replaced by the output of the adder that computes the branch target.
MemRead	None.	Data memory contents designated by the address input are put on the Read data output.
MemWrite	None.	Data memory contents designated by the address input are replaced by the value on the Write data input.
MemtoReg	The value fed to the register Write data input comes from the ALU.	The value fed to the register Write data input comes from the data memory.

R-Type Instruction



R-type instruction, such as add **$t1,$t2,$t3**
Although everything occurs in one clock cycle, we can think of four steps to execute the instruction:
1. The instruction is fetched, and the PC is incremented.
2. Two registers, $t2 and $t3, are read from the register file; also, the main control unit computes the setting of the control lines during this step.
3. The ALU operates on the data read from the register file, using the function code (bits 5:0, which is the funct field, of the instruction) to generate the ALU function.
4. The result from the ALU is written into the register file using bits 15:11 of the instruction to select the destination register ($t1).

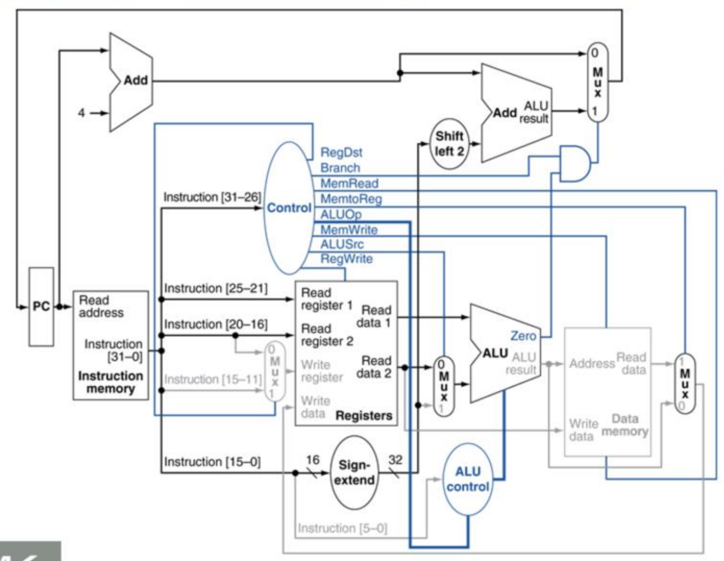
Load Instruction



We can think of a load instruction such as **lw $t1, offset($t2)** operating in five steps:

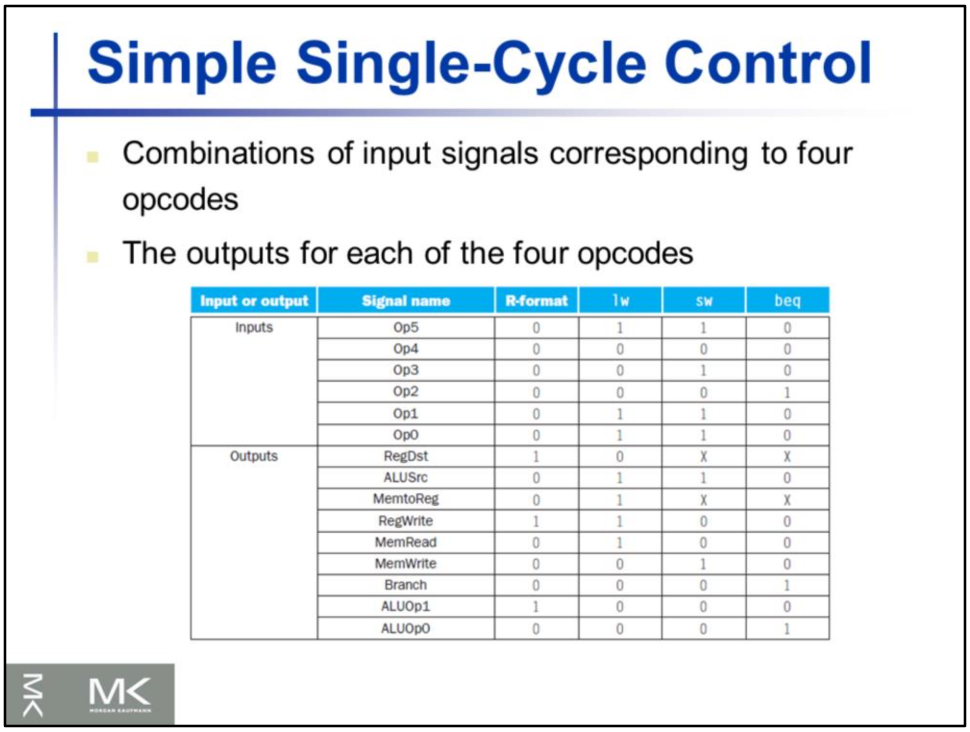
1. An instruction is fetched from the instruction memory, and the PC is incremented.
2. A register ($t2) value is read from the register file.
3. The ALU computes the sum of the value read from the register file and the sign-extended, lower 16 bits of the instruction (offset).
4. The sum from the ALU is used as the address for the data memory.
5. The data from the memory unit is written into the register file; the register destination is given by bits 20:16 of the instruction ($t1).

Branch-on-Equal Instruction



The operation of the branch-on-equal instruction, such as beq $t1, $t2, offset occurs in four steps:

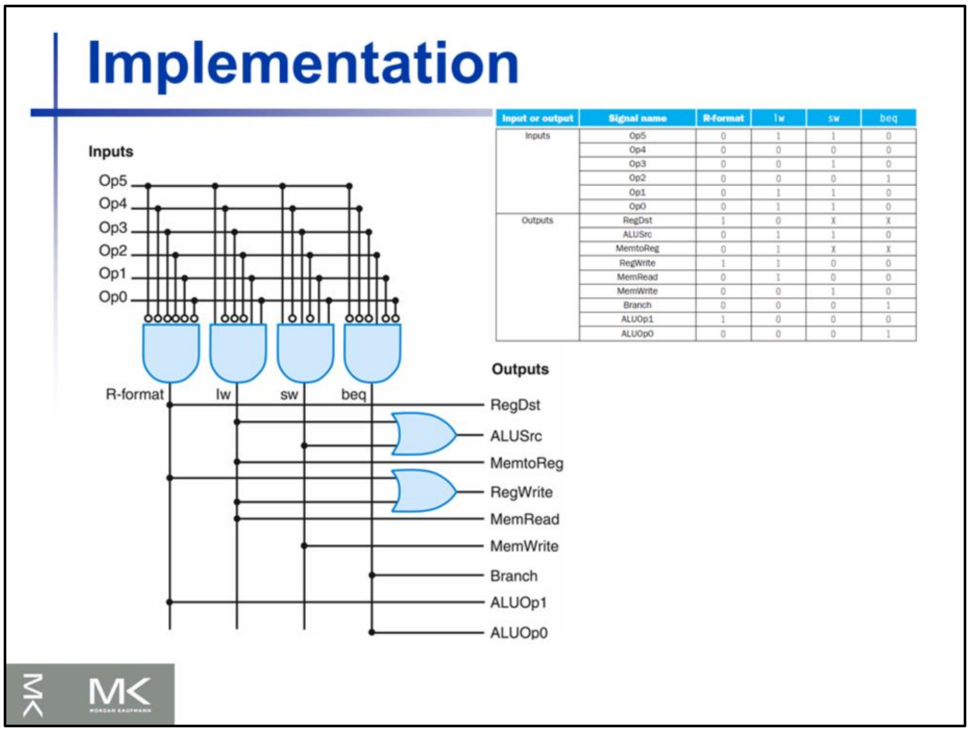
1. An instruction is fetched from the instruction memory, and the PC is incremented.
2. Two registers, $t1 and $t2, are read from the register file.
3. The ALU performs a subtract on the data values read from the register file. The value of PC + 4 is added to the sign extended, lower 16 bits of the instruction (offset) shifted left by two; the result is the branch target address.
4. The Zero result from the ALU is used to decide which adder result to store into the PC.



Simple Single-Cycle Control

- Combinations of input signals corresponding to four opcodes
- The outputs for each of the four opcodes

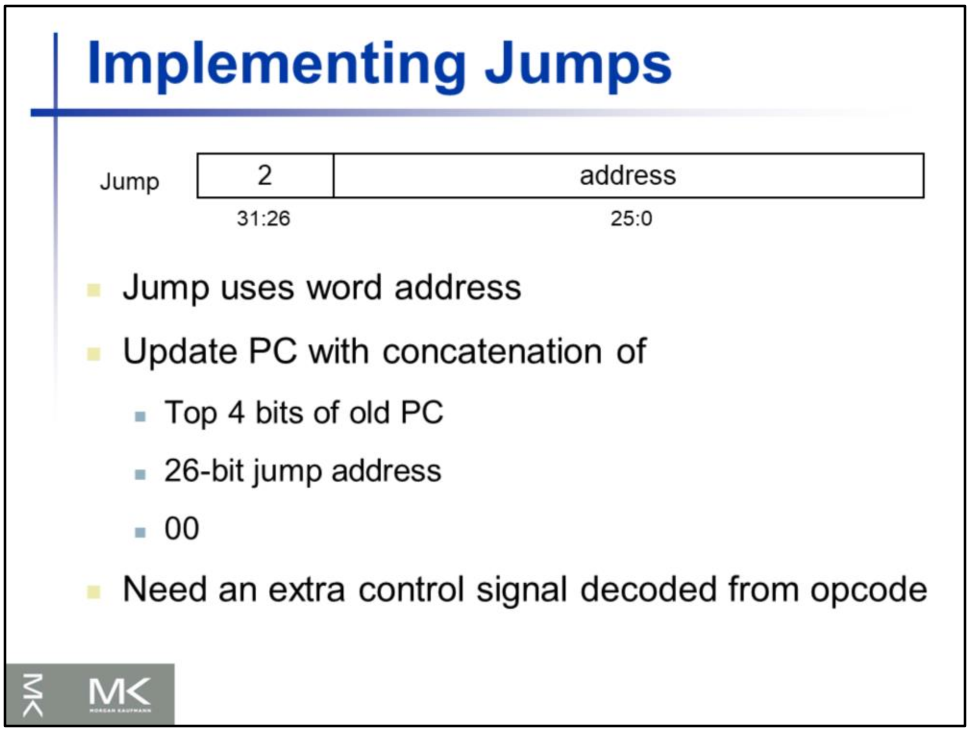
Input or output	Signal name	R-format	lw	sw	beq
Inputs	Op5	0	1	1	0
	Op4	0	0	0	0
	Op3	0	0	1	0
	Op2	0	0	0	1
	Op1	0	1	1	0
	Op0	0	1	1	0
Outputs	RegDst	1	0	X	X
	ALUSrc	0	1	1	0
	MemtoReg	0	1	X	X
	RegWrite	1	1	0	0
	MemRead	0	1	0	0
	MemWrite	0	0	1	0
	Branch	0	0	0	1
	ALUOp1	1	0	0	0
	ALUOp0	0	0	0	1



Implementation

Input or output	Signal name	R-format	lw	sw	beq
Inputs	Op5	0	1	1	0
	Op4	0	0	0	0
	Op3	0	0	1	0
	Op2	0	0	0	1
	Op1	0	1	1	0
	Op0	0	1	1	0
Outputs	RegDst	1	0	X	X
	ALUSrc	0	1	1	0
	MemtoReg	0	1	X	X
	RegWrite	1	1	0	0
	MemRead	0	1	0	0
	MemWrite	0	0	1	0
	Branch	0	0	0	1
	ALUOp1	1	0	0	0
	ALUOp0	0	0	0	1

The structure, called a programmable logic array (PLA), uses an array of AND gates followed by an array of OR gates. The inputs to the AND gates are the function inputs and their inverses (bubbles indicate inversion of a signal). The inputs to the OR gates are the outputs of the AND gates (or, as a degenerate case, the function inputs and inverses). The output of the OR gates is the function outputs.



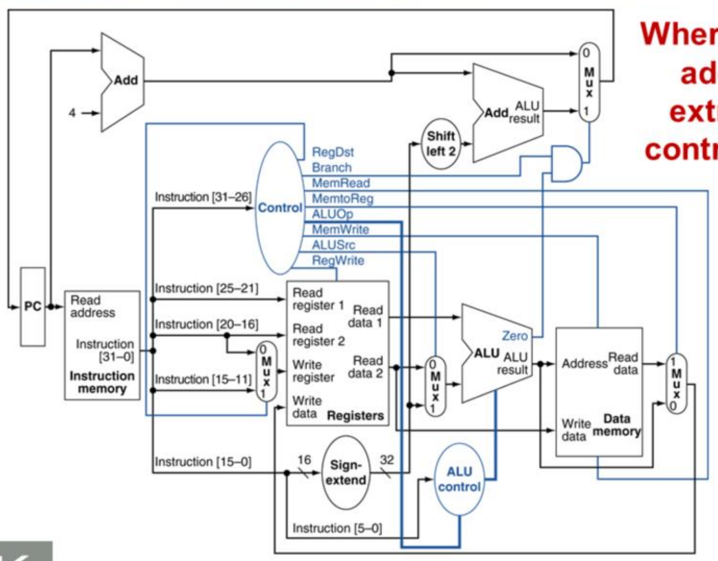
Implementing Jumps

Jump	2	address
	31:26	25:0

- Jump uses word address
- Update PC with concatenation of
  - Top 4 bits of old PC
  - 26-bit jump address
  - 00
- Need an extra control signal decoded from opcode

FIGURE 4.23 Instruction format for the jump instruction (opcode = 2). The destination address for a jump instruction is formed by concatenating the upper 4 bits of the current PC + 4 to the 26-bit address field in the jump instruction and adding 00 as the 2 low-order bits.

Datapath With Control



Where to add extra control?

Datapath With Jumps Added

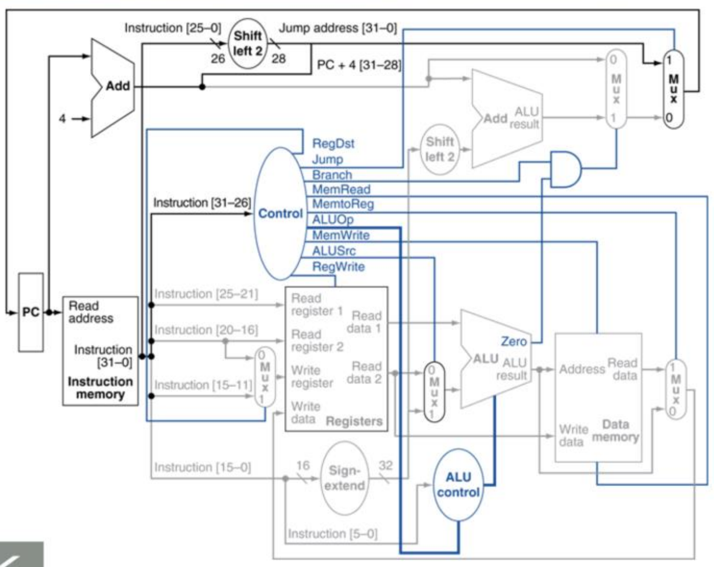
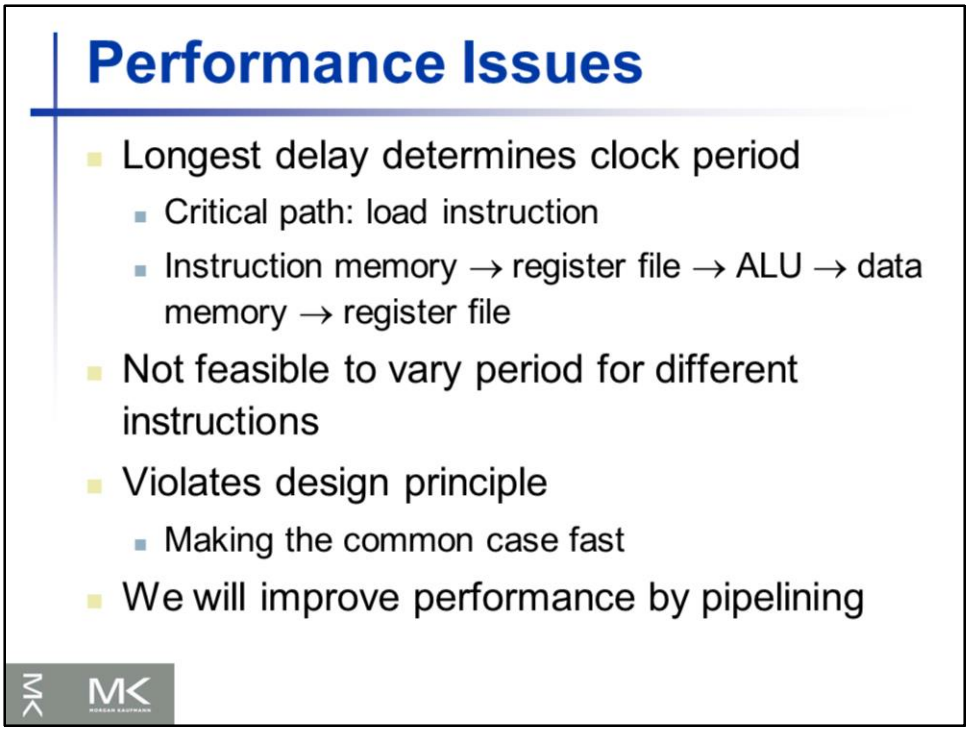


FIGURE 4.24 The simple control and datapath are extended to handle the jump instruction.
An additional multiplexor (at the upper right) is used to choose between the jump target and either the branch target or the sequential instruction following this one. This multiplexor is controlled by the jump control signal. The jump target address is obtained by shifting the lower 26 bits of the jump instruction left 2 bits, effectively adding 00 as the low order bits, and then concatenating the upper 4 bits of PC + 4 as the high-order bits, thus yielding a 32-bit address.



Performance Issues

- Longest delay determines clock period
  - Critical path: load instruction
  - Instruction memory → register file → ALU → data memory → register file
- Not feasible to vary period for different instructions
- Violates design principle
  - Making the common case fast
- We will improve performance by pipelining

Because we must assume that the clock cycle is equal to the worst case delay for all instructions, it's useless to try implementation techniques that reduce the delay of the common case but do not improve the worst-case cycle time. A single-cycle implementation thus violates the great idea of making the **common case fast**.