





## Lenguajes y Autómatas II Árbol de Derivación

**REPASO** 

Unidad V. Análisis Sintáctico





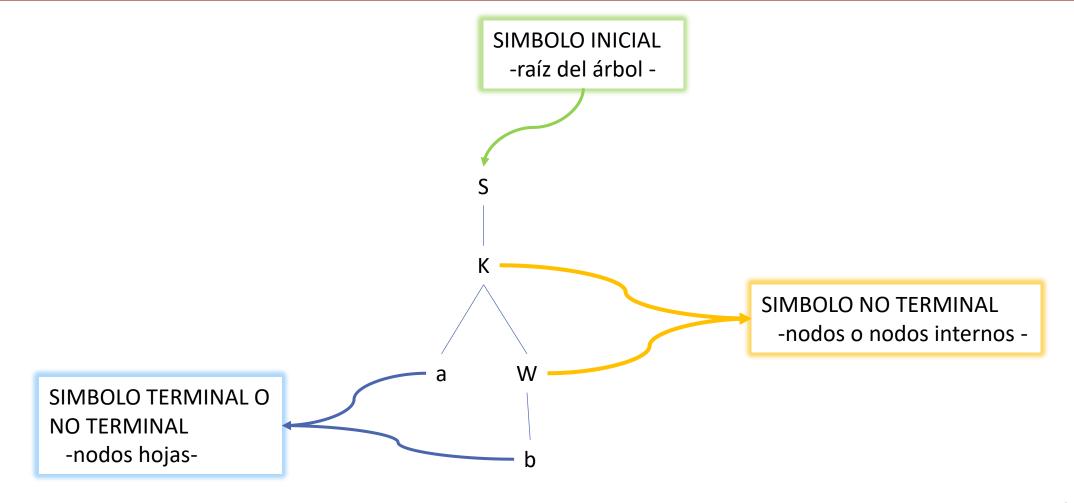


Es una técnica que se utiliza para revisar que una palabra cumpla con la expresión regular que representa la Gramática Libre del Contexto.















## ¿CÓMO SE QUE LA PALABRA DE ENTRADA ES CORRECTA O PERTENECE A LA EXPRESION REGULAR?

- En las hojas del árbol de derivación, de izquierda a derecha, al concatenar los símbolos, coincida con la palabra de entrada, son iguales
- En las hojas del árbol de derivación, no haya ningún símbolo NO TERMINAL y hayamos terminado de recorrer toda la palabra de entrada.







## ¿CÓMO SE QUE LA PALABRA DE ENTRADA ES INCORRECTA O NO PERTENECE A LA EXPRESION REGULAR?

- En las hojas del árbol de derivación, de izquierda a derecha, al concatenar los símbolos, NO coincide con la palabra de entrada, son diferentes
- En las hojas del árbol de derivación, hay un símbolo NO TERMINAL
- No terminamos de recorrer toda la palabra de entrada.







# Expresión regular a







Sea  $\Sigma$ = {a} y G =(T,N,I,P) una gramática libre del contexto, que representa el lenguaje a  $_{W=a}$ 

$$T = \{a\}$$

$$N=\{S\}$$

$$I = S$$

**P**:

S | | | |

**Correcta** 

$$W = a$$







E.R. = 
$$a = \{a\}$$

$$T = \{a\}$$

$$N=\{S\}$$

$$I = S$$

**P**:

W = b

S

#### Incorrecta.

- 1. Las palabras son diferentes
- 2. La hoja del árbol es la raíz
- La hoja del árbol es un símbolo NO TERMINAL

GLC

W = S







E.R. = 
$$a = \{a\}$$

$$T = \{a\}$$

$$N=\{S,A\}$$

Р:





Correcta







E.R. = 
$$a = \{a\}$$

**P**:

A -> a

$$W = \frac{a}{b}$$

A |

GLC W = a

#### Incorrecta

- 1. Son diferentes las palabras
- 2. No terminamos de recorrer la palabra de entrada







# Expresión regular ab



w = ab





#### Árbol de derivación

Sea  $\Sigma$ = {a,b} y G =(T,N,I,P) una gramática libre del contexto, que representa el lenguaje ab

```
E.R. = ab = {ab}
```

```
T = {a,b}
N={S, K, W}
I = S
P:
S -> K
K -> aW
W → b
```







E.R. = 
$$ab = {ab}$$

$$S \rightarrow K$$
  
 $K \rightarrow aW$   
 $W \rightarrow b$ 

$$w = \frac{a}{b}$$

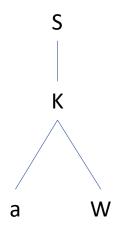






$$E.R. = ab = {ab}$$

$$w = \frac{a}{b}$$



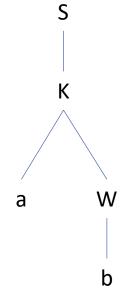






$$E.R. = ab = {ab}$$

$$w = ab$$



#### **Correcta**







E.R. = 
$$ab = {ab}$$

 $W \rightarrow b$ 

$$w = a$$

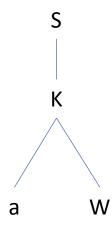






E.R. = 
$$ab = {ab}$$

$$w = a$$





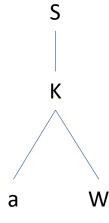




$$E.R. = ab = {ab}$$

 $W \rightarrow b$ 





#### Incorrecta

- 1. Las palabras son diferentes
- 2. Hay una hoja del árbol que tiene un símbolo NO terminal (W)



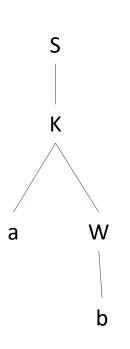




$$E.R. = ab = {ab}$$

K -> aW

 $W \rightarrow b$ 



w = a

Otra forma de generar el árbol.

Técnica no es recomendable para un compilador

#### Incorrecta

1. Las palabras son diferentes

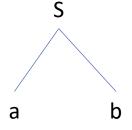






$$E.R. = ab = {ab}$$

w=ab



**Correcta** 







$$E.R. = ab = {ab}$$

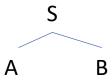
$$T = \{a,b\}$$
  
 $N=\{S,A,B\}$ 

$$I = S$$

P:

$$B \rightarrow b$$

$$W = b$$







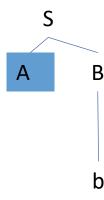


$$E.R. = ab = {ab}$$

P:

$$B \rightarrow b$$

$$W = b$$



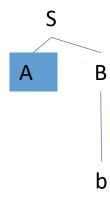






$$E.R. = ab = {ab}$$

$$W = b$$



INCORRECTA.

- 1. Una hoja del árbol es NO TERMINAL
- 2. Las palabras son incorrectas

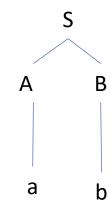
$$GLC$$
  $W = Ab$ 







$$W = b$$



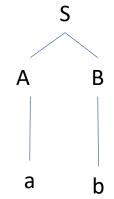
Otra forma de hacerlo







$$W = b$$



INCORRECTA.

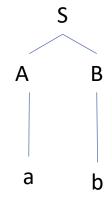
Las palabras son diferentes







$$W = ba$$



INCORRECTA.

Las palabras son diferentes

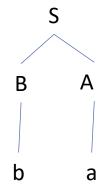
$$W = ab$$







$$W = ba$$



#### MAL HECHO EL ARBOL.

No es correcto cambiar el orden de los símbolos que está en la producción S para forzar el generar la palabra de entrada

$$GLC = ba$$







# Expresión regular a U b







E.R. = $\mathbf{a} \cup \mathbf{b}$	= {a, b}
$T = \{a,b\}$	
$N=\{S,A,B\}$	
I = S	
P:	
S -> A   B	
A -> a	
B -> b	

W = a

S







```
E.R. = a \cup b = \{a, b\}
T = \{a,b\}
N=\{S,A,B\}
I = S
   S -> A | B
   A \rightarrow a
   B \rightarrow b
```

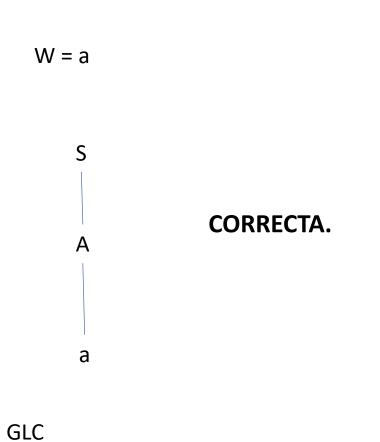


W = a





```
E.R. = a \cup b = \{a, b\}
T = \{a,b\}
N=\{S,A,B\}
I = S
   S -> A | B
   A \rightarrow a
   B \rightarrow b
```









$$W = b$$

S







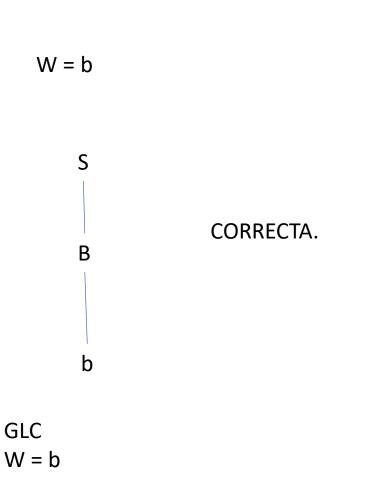
```
E.R. = a \cup b = \{a, b\}
T = \{a,b\}
N=\{S,A,B\}
I = S
   S -> A | B
   A \rightarrow a
   B \rightarrow b
```







```
E.R. = a \cup b = \{a, b\}
T = \{a,b\}
N=\{S,A,B\}
I = S
   S -> A | B
   A \rightarrow a
   B \rightarrow b
```













#### INCORRECTA.

- 1. Las palabras son diferentes
- No terminamos de revisar la palabra





















w = c

S

#### **INCORRECTA**

La raíz del árbol es a la Vez nodo hoja y es un NO TERMINAL. La palabra que genera es Diferente a la entrada.

GLC W =S







$$W = \frac{a}{b}$$

S







$$W = \frac{a}{a}b$$

S |







$$W = ab$$

S | | |

#### **INCORRECTA**

No son iguales las palabras - Faltó de generar b-

$$W = a$$







```
e.r. = a U b = \{a, b\}
T = \{a,b\}
N=\{S,A,B\}
I = S
   S \rightarrow A \mid B
   A -> a
   B \rightarrow b
```

W = b

S







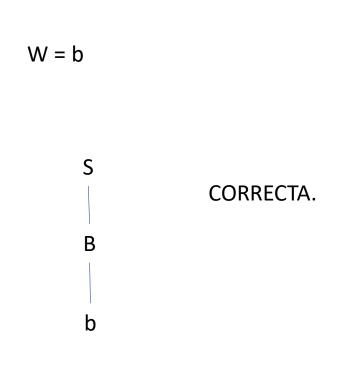
```
e.r. = a U b = \{a, b\}
T = \{a,b\}
N=\{S,A,B\}
I = S
   S \rightarrow A \mid B
   A -> a
   B \rightarrow b
```







```
e.r. = a U b = \{a, b\}
T = \{a,b\}
N=\{S,A,B\}
I = S
   S \rightarrow A \mid B
   A -> a
   B \rightarrow b
```









# Expresión regular (ab) U c







```
E.R. = (ab) \cup c = {ab, c}

T={a,b,c}

N={S, A, B,C}

I = S

P:

S \rightarrow A | C

A \rightarrow aB

B \rightarrow b

C \rightarrow c
```

W = ab

S







```
E.R. = (ab) \cup c = {ab, c}
T={a,b,c}
                                                                     W = ab
N={S, A, B,C}
I = S
P:
 S \rightarrow A \mid C
  A \rightarrow aB
  B \rightarrow b
  C \rightarrow c
```







```
E.R. = (ab) \cup c = {ab, c}

T={a,b,c}

N={S, A, B,C}

I = S

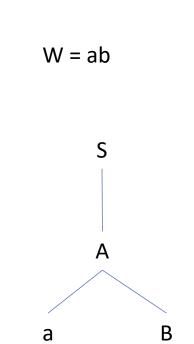
P:

S \rightarrow A | C

A \rightarrow aB

B \rightarrow b

C \rightarrow c
```

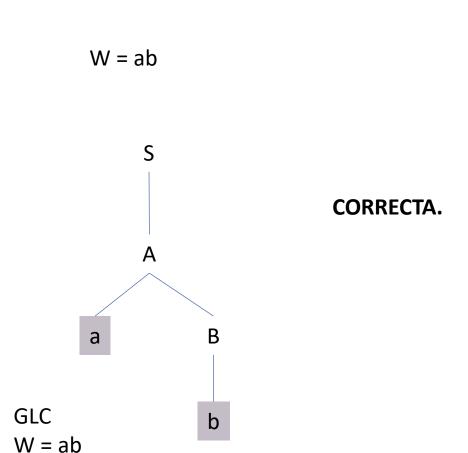








```
E.R. = (ab) \cup c = {ab, c}
T={a,b,c}
N={S, A, B,C}
I = S
P:
 S \rightarrow A \mid C
  A \rightarrow aB
  B \rightarrow b
  C \rightarrow c
```









```
E.R. = (ab) \cup c = {ab, c}

T={a,b,c}

N={S, A, B,C}

I = S

P:

S \rightarrow A | C

A \rightarrow aB

B \rightarrow b

C \rightarrow c
```

W = c

S







```
E.R. = (ab) \cup c = {ab, c}
T={a,b,c}
                                                                      W = c
N={S, A, B,C}
I = S
P:
 S \rightarrow A \mid C
  A \rightarrow aB
  B \rightarrow b
  C \rightarrow c
```







```
E.R. = (ab) \cup c = {ab, c}

T={a,b,c}

N={S, A, B,C}

I = S

P:

S \rightarrow A | C

A \rightarrow aB

B \rightarrow b

C \rightarrow c
```

W = c

CORRECTA.

GLCW = c

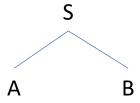






```
E.R. = (ab) \cup c = {ab, c}
T={a,b,c}
N={S, A, B,C}
I = S
 S \rightarrow AB \mid C
 A \rightarrow a
 B-> b
C -> c
```

$$W = ab$$



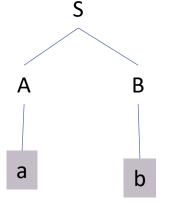






```
E.R. = (ab) \cup c = {ab, c}
T={a,b,c}
N={S, A, B,C}
I = S
 S \rightarrow AB \mid C
 A \rightarrow a
 B-> b
 C -> c
```

$$W = ab$$



#### CORRECTA.

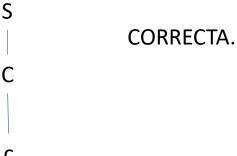






```
E.R. = (ab) \cup c = {ab, c}
T={a,b,c}
N={S, A, B,C}
I = S
 S \rightarrow AB \mid C
 A \rightarrow a
 B-> b
C -> c
```

$$W = c$$

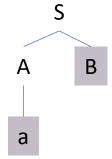












INCORRECTA.

Porque tiene un símbolo

No terminal B.

Son diferentes palabras

$$W = aB$$

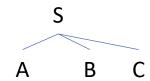






```
E.R. = (ab) \cup c = {ab, c}
T={a,b,c}
N={S, A, B,C}
I = S
 S \rightarrow AB | C
 A \rightarrow a
 B-> b
 C \rightarrow c
```

$$W = ab$$



#### **INCORRECTA**

Representación del árbol

En S hay 2 términos AB y C NO se deben colocar todos. Hay que seleccioanr sólo Uno.

$$W =$$







# Expresión regular a\*







```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
        S \rightarrow A
       A \rightarrow aA \mid \epsilon
```

W=aaa







```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
        S \rightarrow A
       A \rightarrow aA \mid \epsilon
```

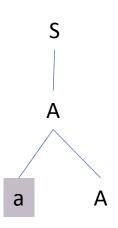
W=aaa







```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
        S \rightarrow A
       A \rightarrow aA \mid \epsilon
```



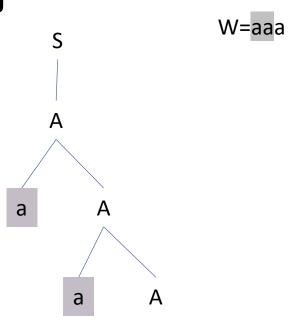
W=aaa







```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
         S \rightarrow A
        A \rightarrow aA \mid \epsilon
```

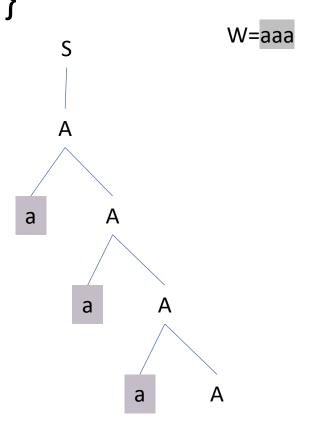








```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
        S \rightarrow A
       A \rightarrow aA \mid \epsilon
```

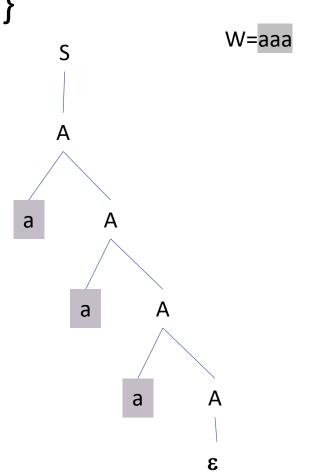








```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
        S \rightarrow A
       A \rightarrow aA \mid \epsilon
```



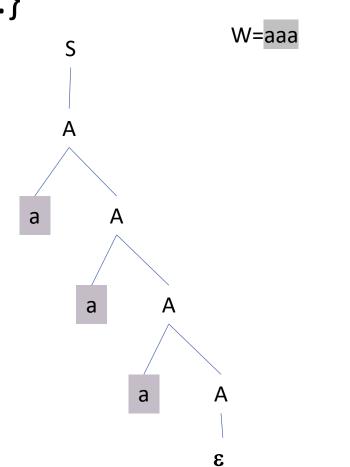






```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
        S \rightarrow A
       A \rightarrow aA \mid \epsilon
                                 GLC
```

 $W = aaa\varepsilon = aaa$ 



Correcta





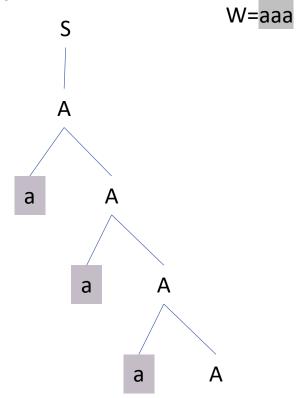


```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
```

**P**:

$$S \rightarrow A$$
  
  $A \rightarrow aA \mid \epsilon$ 

GLC W = aaaA



#### **Incorrecta**

Las palabras no coinciden Mala derivación

Faltó generar un nodo hijo a A con la cadena vacía

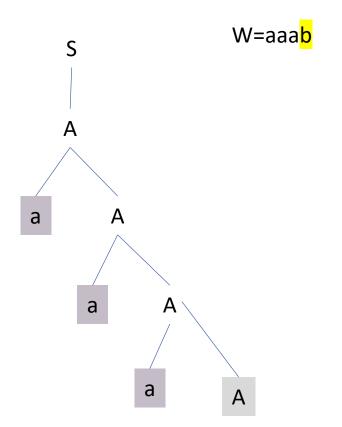






```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
```

```
T={a}
N={S,A}
I=S
P:
S \rightarrow A
A \rightarrow aA \mid \epsilon
```



Incorrecta

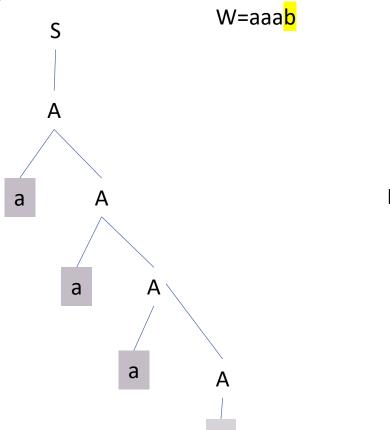






```
E.R. = a^* = \{\varepsilon, a, aa, aaa, aaaa,....\}
T={a}
N=\{S,A\}
I=S
        S \rightarrow A
       A \rightarrow aA \mid \epsilon
                                 GLC
```

W = aaa



Incorrecta







```
E.R. =a*= {\epsilon, a, aa, aaa, aaaa,...}

T={a}

N={S}

I=S

P:

S \rightarrow aS | \epsilon
```





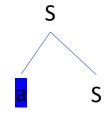


```
E.R. =a^*= {\epsilon, a, aa, aaa, aaaa,...}
T={a}
N={S}
```

**P**:

I=S

$$S \rightarrow aS \mid \epsilon$$



$$W =$$



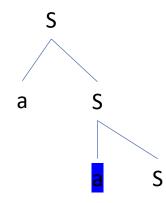




```
E.R. =a*= {ε, a, aaa, aaaa, ....}
T={a}
N={S}
I=S
```

 $S \rightarrow aS \mid \epsilon$ 

$$W = aaa$$



$$W =$$

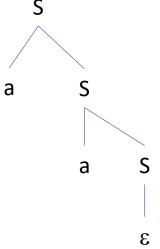






```
E.R. =a^* = \{\varepsilon, a, aa, aaa, aaaa,...\}
T={a}
N=\{S\}
I=S
        S \rightarrow aS \mid \epsilon
```

$$W = aa$$



Correcta



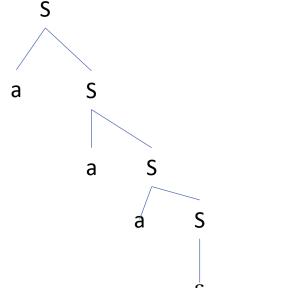




```
E.R. =a^* = \{\varepsilon, a, aa, aaa, aaaa,...\}
T={a}
N=\{S\}
I=S
        S \rightarrow aS \mid \epsilon
```

GLC

 $W = aaa\epsilon = a$ 



**Correcta** 







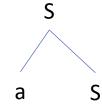
E.R. 
$$=a^* = \{ \varepsilon, a, aa, aaa, aaaa, ... \}$$

$$N=\{S\}$$

**P**:

$$S \rightarrow aS \mid \epsilon$$





#### **INCORRECTA**

Diferentes palabras
Hay nodos hojas con
SIMBOLOS NO TERMINALES







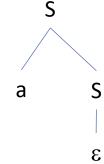
E.R. 
$$=a^* = \{ \varepsilon, a, aa, aaa, aaaa, ... \}$$

$$N=\{S\}$$

**P**:

$$S \rightarrow aS \mid \epsilon$$

$$W = ab$$



**INCORRECTA** 

Diferentes palabras







# Expresión regular a<sup>+</sup>







E.R. =  $a^+$  ={ a, aa, aaa, aaaa,...}

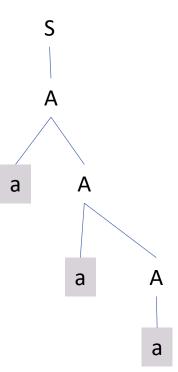
$$N=\{S,A\}$$

Р:

$$S \rightarrow A$$

$$A \rightarrow aA \mid a$$





**Correcta** 



a





### Árbol de derivación

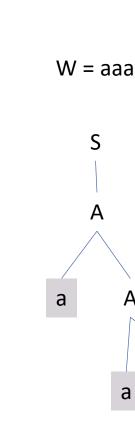
E.R.  $=a^{+}= \{ a, aa, aaa, aaaa,... \}$ 

 $N=\{S,A\}$ 

I=S

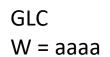
 $S \rightarrow A$ 

 $A \rightarrow aA \mid a$ 



#### **Incorrecto**

Las palabras son diferentes Árbol de derivación incorrecta









```
E.R. =a^+=\{a, aa, aaa, aaaa, ...\}
T=\{a\}
N=\{S\}
I=S
P:
S \rightarrow aS \mid a
V=a
S \rightarrow aS \mid a
```







 $S \rightarrow aS \mid a$ 

$$W = a$$



#### **INCORRECTA**

EL ARBOL DE DERIVACION No corresponde a la palabra de entrada







## Expresión regular (0-3)\*







```
333, 1230, 3322113,...}
T=\{0,1,2,3\}
N=\{S,A\}
I=S
P:
     S \rightarrow \varepsilon \mid 0S \mid 1S \mid 2S \mid 3S
```

W =21

S

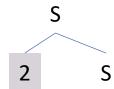






```
333, 1230, 3322113,...}
T=\{0,1,2,3\}
N=\{S,A\}
I=S
P:
     S \rightarrow \varepsilon \mid 0S \mid 1S \mid 2S \mid 3S
```

W = 21



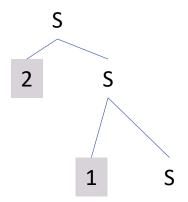






```
333, 1230, 3322113,...}
T=\{0,1,2,3\}
N=\{S,A\}
I=S
P:
    S \rightarrow \epsilon \mid 0S \mid 1S \mid 2S \mid 3S
```

W = 21



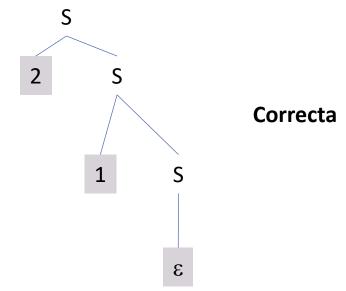






```
333, 1230, 3322113,...}
T=\{0,1,2,3\}
N=\{S,A\}
I=S
P:
    S \rightarrow \epsilon | 0S | 1S | 2S | 3S
```

W = 21



GLC  $W = 21\varepsilon = 21$ 

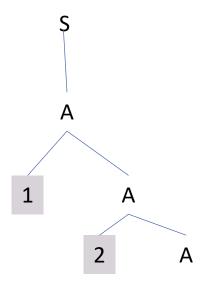






```
E.R. = (0-3)*
T=\{0,1,2,3\}
N={ S,A }
I=S
P:
        S \rightarrow A
        A \rightarrow \epsilon \mid 0A \mid 1A \mid 2A \mid 3A
```

$$W = 123$$



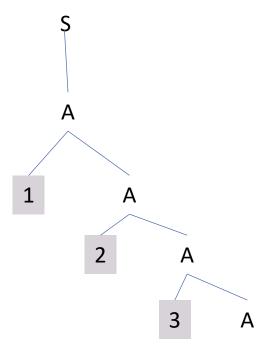






```
E.R. = (0-3)*
T=\{0,1,2,3\}
N={ S,A }
I=S
P:
        S \rightarrow A
        A \rightarrow \epsilon \mid 0A \mid 1A \mid 2A \mid 3A
```



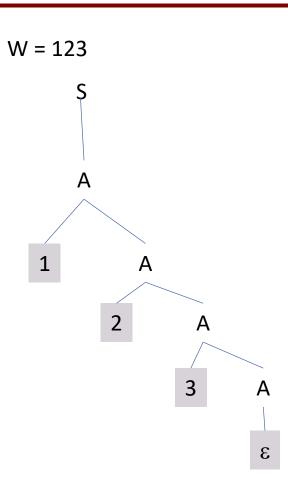








```
E.R. = (0-3)*
T=\{0,1,2,3\}
N={ S,A }
I=S
P:
        S \rightarrow A
        A \rightarrow \epsilon \mid 0A \mid 1A \mid 2A \mid 3A
```

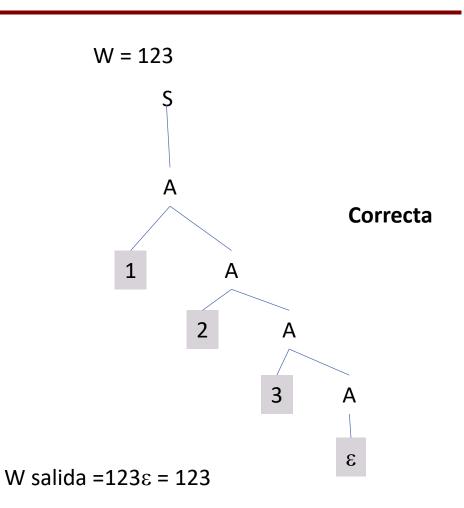








```
E.R. = (0-3)*
T=\{0,1,2,3\}
N=\{S,A\}
I=S
P:
        S \rightarrow A
        A \rightarrow \epsilon \mid 0A \mid 1A \mid 2A \mid 3A
```

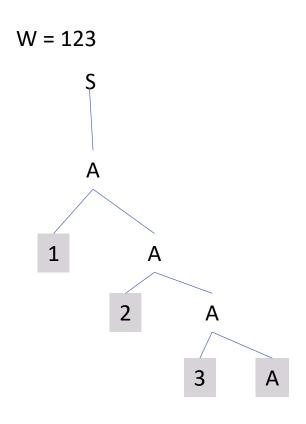








```
E.R. = (0-3)*
T=\{0,1,2,3\}
N=\{S,A\}
                                           Incorrecto
                                           Árbol de derivación
                                           incompleto
I=S
P:
       S \rightarrow A
       A \rightarrow \epsilon | 0A | 1A | 2A | 3A
```



W salida =123A







# Expresión regular a+b







```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```

w=aaab

S







```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

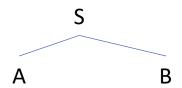
P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```

w=aaab





w=aaab





```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

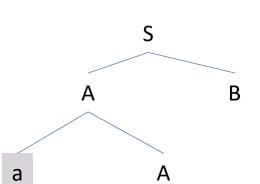
I=S

P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```









```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

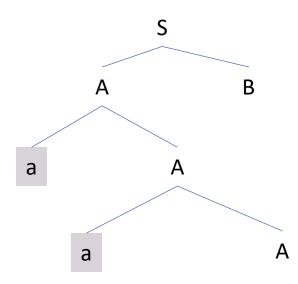
P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```











```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

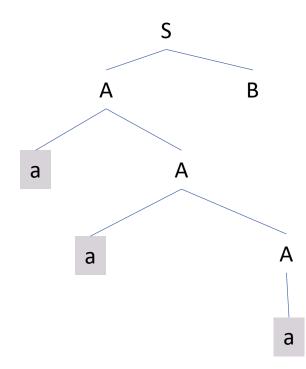
P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```











```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

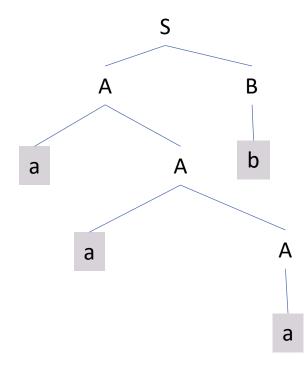
P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```











```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

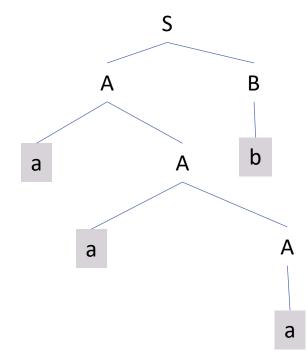
P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```

w=aaab



Correcta







```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

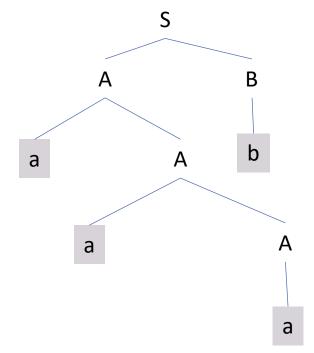
P:

S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```

w=aaabaa



#### Incorrecta

- Las palabras son diferentes
- Faltaron
   símbolos en la
   palabra de
   entrada por
   revisar

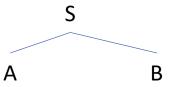






E.R. =  $a^+b$  ={ab, aab, aaab, a...ab,....} T={a,b} N={S,A,B} I=S P: S  $\rightarrow$  AB A  $\rightarrow$  aA | a B  $\rightarrow$  b w=ba

Opción 1



Incorrecta

- Las palabras son diferentes
- Las hojas del árbol son símbolos no terminales







```
E.R. = a^+b ={ab, aab, aaab, a...ab,....}

T={a,b}

N={S,A,B}

I=S

P:

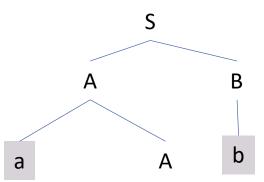
S \rightarrow AB

A \rightarrow aA | a

B \rightarrow b
```



w=ba



#### Opción 2

Incorrecta

- Las palabras son diferentes
- Las hojas del árbol tiene símbolos no terminales







E.R. =  $a^+b$  ={ab, aab, aaab, a...ab,....} T={a,b} N={S,A,B} I=S P: S  $\rightarrow$  AB A  $\rightarrow$  aA | a B  $\rightarrow$  b w=ba

S

Α

a

Opción 3

Incorrecta

 Las palabras son diferentes







## Bibliografía

• Teoría de Autómatas. Un enfoque práctico

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