

P2P

Peer 2 Peer

No always a server, all end-systems want to each other
they request and serve each other

Self-Scalability - New peers bring in new capability, but also new demands
Complex Management (changing IPs, churn)

BitTorrent - File sharing using 256 bit chunks

Server with a tracker, registers all peers

Request
Chunks

Alice may ask for a list of chunks peers have

Downloading Rarest First

Alice sends chunks to 4 peers at the highest rate, chokes others
updates friends every 30s, reveals every 10s

Higher upload rate, find better partners, get file faster

Distribution
Time

Distribute F-sized file to N clients

u_s - server upload

u_i - peer upload

d_i - peer download

For Server

Time to upload copies

NF/u_s

$$D_{cs} \geq \max \left\{ \frac{NF}{u_s}, \frac{F}{d_{min}} \right\}$$

/ increases linear

Lower Bound

$\frac{u_s}{N} \leq d_{min}$ - Server sends file in parallel

$\frac{u_s}{N} > d_{min}$ - Server sends file in parallel at rate d_{min}

Too long if N is big

For Client

Each Client must download a copy

d_{min} - min download rate
 F/d_{min}

$$D_{cs} = \frac{NF}{u_s}$$

$$D_{cs} = \frac{F}{d_{min}}$$

P2P
(lower)
Bounded

Server need to upload a copy
 F/u_s

Client downloads a copy
 F/d_{min}

Aggregate copies - NF bits
 $u_s + \sum_i u_i$

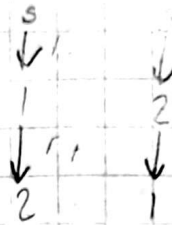
$$D_{P2P} \geq \max \left\{ \frac{F}{u_s}, \frac{F}{d_{min}}, \frac{NF}{u_s + \sum_i u_i} \right\}$$

Only for N

P2P lower bound



Find $\max \sum r_i$



Trees



$\sum r_i \leq d_i \leftarrow$ transmission rate is dependent on d_i of nodes

$$\sum_{j \in \text{children}(i)} r_j \leq u_i$$

Children Nodes

$$u_s \leq \min \left(d_{\min}, \frac{u_s + \sum_i u_i}{N} \right) \rightarrow r_i = \frac{u_i}{\sum_i u_i} u_s$$

$$\text{Download rate } r_i + \sum r_j = \sum_i r_i = u_s < d_{\min}$$

$$\text{Upload rate } \sum_i r_i = u_s$$

$$\text{Total Upload } (N-1)r_i = \frac{u_i(N-1)}{N} u_s < u_i \quad \left. \begin{array}{l} \text{Download rate} \\ \text{Upload rate} \\ \text{Total Upload} \end{array} \right\} \frac{F}{u_s}$$

$$\frac{u_s + \sum_i u_i}{N} \leq \min(d_{\min}, u_i)$$

$$r_i = \frac{u_i}{N+1}$$

$$d_{\min} \leq \min \left(u_s, \frac{u_s + \sum_i u_i}{N} \right) \quad \frac{F}{d_{\min}}$$