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DOMAIN WINTER WINNING CAMP ASSIGNMENT

Student Name: Kriti Saxena

UID: 22BCS14603

Branch: BE-CSE

Section/Group: 22BCS_FL_IOT-603/B

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STACK AND QUEUE STANDARD QUESTION

VERY EASY:

Q. Design a stack that supports push, pop, top, and retrieving the minimum element in constant time.

Implement the MinStack class:

- MinStack() initializes the stack object.
- void push(int val) pushes the element val onto the stack.
- void pop() removes the element on the top of the stack.
- int top() gets the top element of the stack.
- int getMin() retrieves the minimum element in the stack.

You must implement a solution with O(1) time complexity for each function.

Example 1:

Input

["MinStack","push","push","push","getMin","pop","top","getMin"]

[[],[-2],[0],[-3],[,],[,],[,]]

Output

[null,null,null,null,-3,null,0,-2]

Explanation

MinStack minStack = new MinStack();

minStack.push(-2);

minStack.push(0);



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```
minStack.push(-3);
```

```
minStack.getMin(); // return -3
```

```
minStack.pop();
```

```
minStack.top(); // return 0
```

```
minStack.getMin(); // return -2
```

CODE:

```
#include <iostream>
```

```
#include <stack>
```

```
#include <limits.h>
```

```
class MinStack {
```

```
private:
```

```
    std::stack<int> stack;
```

```
    std::stack<int> minStack;
```

```
public:
```

```
    MinStack() {
```

```
        // Constructor initializes an empty stack
```

```
    }
```

```
    void push(int val) {
```

```
        stack.push(val);
```

```
        // Push onto minStack if it's empty or val is less than or equal to the current minimum
```

```
        if (minStack.empty() || val <= minStack.top()) {
```

```
            minStack.push(val);
```

```
        }
```

```
    }
```

```
    void pop() {
```

```
        if (stack.top() == minStack.top()) {
```

```
            minStack.pop();
```

```
        }
```

```
        stack.pop();
```

```
    }
```

```
    int top() {
```

```
        return stack.top();
```

```
    }
```

```
    int getMin() {
```

```
        return minStack.top();
```

```
    }
```

```
};
```

```
int main() {
```

```
    MinStack minStack;
```

```
    minStack.push(-2);
```

```
minStack.push(0);
minStack.push(-3);
std::cout << minStack.getMin() << std::endl; // Output: -3
minStack.pop();
std::cout << minStack.top() << std::endl;    // Output: 0
std::cout << minStack.getMin() << std::endl; // Output: -2
return 0;
}
```

OUTPUT:

```
-3
0
-2
```

Example 2:

Input:

["MinStack", "push", "push", "push", "push", "getMin", "pop", "getMin", "top", "getMin"]

[[], [5], [3], [7], [3], [], [], [], [], []]

Output

[null, null, null, null, null, 3, null, 3, 7, 3]

Explanation:

MinStack minStack = new MinStack();

minStack.push(5); # Stack: [5], MinStack: [5]

minStack.push(3); # Stack: [5, 3], MinStack: [5, 3]

minStack.push(7); # Stack: [5, 3, 7], MinStack: [5, 3]

minStack.push(3); # Stack: [5, 3, 7, 3], MinStack: [5, 3, 3]

minStack.getMin(); # Returns 3

minStack.pop(); # Removes 3; Stack: [5, 3, 7], MinStack: [5, 3]

minStack.getMin(); # Returns 3

minStack.top(); # Returns 7

minStack.getMin(); # Returns 3

- Minimum values are maintained as: [5] → [5, 3] → [5, 3] → [5, 3]
- After pops, the minimum values update accordingly.

Example 3:



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Input:

["MinStack", "push", "push", "push", "getMin", "pop", "getMin", "pop", "getMin"]

[[], [2], [1], [4], [], [], [], [], []]

Output:

[null, null, null, null, 1, null, 1, null, 2]

Explanation:

```
minStack = MinStack()
minStack.push(2)
minStack.push(1)
minStack.push(4)
minStack.push(1)
print(minStack.getMin()) # Output: 1
minStack.pop()
print(minStack.getMin()) # Output: 1
minStack.pop()
print(minStack.getMin()) # Output: 1
minStack.pop()
print(minStack.getMin()) # Output: 2
```

- Minimum values are maintained as: $[2] \rightarrow [2, 1] \rightarrow [2, 1] \rightarrow [2, 1]$
- After pops, the minimum values update accordingly.

Constraints:

- $-2^{31} \leq \text{val} \leq 2^{31} - 1$
- Methods pop, top and getMin operations will always be called on non-empty stacks.
- At most $3 * 10^4$ calls will be made to push, pop, top, and getMin.

Sources: <https://leetcode.com/problems/min-stack/>

CODE:

```
#include <iostream>
#include <stack>
#include <limits.h>
```

```
class MinStack {
private:
    std::stack<int> stack;
    std::stack<int> minStack;
```



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```
public:

    MinStack() {
        // Constructor initializes an empty stack
    }

    void push(int val) {
        stack.push(val);
        // Push onto minStack if it's empty or val is less than or equal to the current minimum
        if (minStack.empty() || val <= minStack.top()) {
            minStack.push(val);
        }
    }

    void pop() {
        if (stack.top() == minStack.top()) {
            minStack.pop();
        }
        stack.pop();
    }

    int top() {
        return stack.top();
    }

    int getMin() {
        return minStack.top();
    }
};

int main() {
    MinStack minStack;
    minStack.push(5); // Stack: [5], MinStack: [5]
```



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```
minStack.push(3);    // Stack: [5, 3], MinStack: [5, 3]
minStack.push(7);    // Stack: [5, 3, 7], MinStack: [5, 3]
minStack.push(3);    // Stack: [5, 3, 7, 3], MinStack: [5, 3, 3]

std::cout << minStack.getMin() << std::endl; // Output: 3
minStack.pop();      // Removes 3; Stack: [5, 3, 7], MinStack: [5, 3]
std::cout << minStack.getMin() << std::endl; // Output: 3
std::cout << minStack.top() << std::endl;   // Output: 7
std::cout << minStack.getMin() << std::endl; // Output: 3

return 0;
}
```

OUTPUT:

```
3
3
7
3
```

EASY LEVEL:-

Q. The school cafeteria offers circular and square sandwiches at lunch break, referred to by numbers 0 and 1 respectively. All students stand in a queue. Each student either prefers square or circular sandwiches.

The number of sandwiches in the cafeteria is equal to the number of students. The sandwiches are placed in a stack. At each step:

If the student at the front of the queue prefers the sandwich on the top of the stack, they will take it and leave the queue.

Otherwise, they will leave it and go to the queue's end.

This continues until none of the queue students want to take the top sandwich and are thus unable to eat.

You are given two integer arrays `students` and `sandwiches` where `sandwiches[i]` is the type of the `i`th sandwich in the stack (`i = 0` is the top of the stack) and `students[j]` is the preference of the `j`th student in the initial queue (`j = 0` is the front of the queue). Return the



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number of students that are unable to eat.

Example 1:

Input: students = [1,1,0,0], sandwiches = [0,1,0,1]

Output: 0

Explanation:

- Front student leaves the top sandwich and returns to the end of the line making students = [1,0,0,1].
 - Front student leaves the top sandwich and returns to the end of the line making students = [0,0,1,1].
 - Front student takes the top sandwich and leaves the line making students = [0,1,1] and sandwiches = [1,0,1].
 - Front student leaves the top sandwich and returns to the end of the line making students = [1,1,0].
 - Front student takes the top sandwich and leaves the line making students = [1,0] and sandwiches = [0,1].
 - Front student leaves the top sandwich and returns to the end of the line making students = [0,1].
 - Front student takes the top sandwich and leaves the line making students = [1] and sandwiches = [1].
 - Front student takes the top sandwich and leaves the line making students = [] and sandwiches = [].
- Hence all students are able to eat.

Example 2:

Input: students = [1,1,1,0,0,1], sandwiches = [1,0,0,0,1,1]

Output: 3

Constraints:

- $1 \leq \text{students.length}, \text{sandwiches.length} \leq 100$
- $\text{students.length} == \text{sandwiches.length}$
- $\text{sandwiches}[i]$ is 0 or 1.
- $\text{students}[i]$ is 0 or 1.

Approach

- Create two queues of students and sandwiches
- And a count variable to check if is loop in left student
- If students in the queue cannot have their ordered sandwiches, it makes a loop. If it is a loop, just break and return the result
- Then implement the program like the given rules.

Time complexity: $O(n)$

Space complexity: $O(n)$

Reference : <https://leetcode.com/problems/number-of-students-unable-to-eat-lunch/description/>

CODE:

```
#include <iostream>
#include <queue>
#include <vector>
```

```
int countStudentsUnableToEat(std::vector<int>& students, std::vector<int>& sandwiches) {
    std::queue<int> studentQueue;
```



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```
std::queue<int> sandwichStack;

// Populate the queues
for (int student : students) {
    studentQueue.push(student);
}
for (int sandwich : sandwiches) {
    sandwichStack.push(sandwich);
}

int count = 0; // Counter for unsuccessful rotations

while (!studentQueue.empty() && count < studentQueue.size()) {
    if (studentQueue.front() == sandwichStack.front()) {
        // Student eats the sandwich
        studentQueue.pop();
        sandwichStack.pop();
        count = 0; // Reset counter
    } else {
        // Student moves to the back of the queue
        studentQueue.push(studentQueue.front());
        studentQueue.pop();
        count++;
    }
}

// Remaining students are those unable to eat
return studentQueue.size();
}

int main() {
    std::vector<int> students1 = {1, 1, 0, 0};
    std::vector<int> sandwiches1 = {0, 1, 0, 1};
    std::cout << countStudentsUnableToEat(students1, sandwiches1) << std::endl; // Output: 0

    std::vector<int> students2 = {1, 1, 1, 0, 0, 1};
    std::vector<int> sandwiches2 = {1, 0, 0, 0, 1, 1};
    std::cout << countStudentsUnableToEat(students2, sandwiches2) << std::endl; // Output: 3

    return 0;
}
```

OUTPUT:

```
0
3
```




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MEDIUM:-

Q. Given a circular integer array `nums` (i.e., the next element of `nums[nums.length - 1]` is `nums[0]`), return the next greater number for every element in `nums`.

The next greater number of a number `x` is the first greater number to its traversing-order next in the array, which means you could search circularly to find its next greater number. If it doesn't exist, return `-1` for this number.

Example 1:

Input: `nums = [1,2,1]`

Output: `[2,-1,2]`

Explanation:

- The first 1's next greater number is 2;
- The number 2 can't find next greater number.
- The second 1's next greater number needs to search circularly, which is also 2.

Example 2:

Input: `nums = [1,2,3,4,3]`

Output: `[2,3,4,-1,4]`

Constraints:

- $1 \leq \text{nums.length} \leq 10^4$
- $-10^9 \leq \text{nums}[i] \leq 10^9$

Reference : <https://leetcode.com/problems/next-greater-element-ii/description/>

CODE:

```
#include <iostream>
#include <vector>
#include <stack>
```

```
std::vector<int> nextGreaterElements(const std::vector<int>& nums) {
    int n = nums.size();
    std::vector<int> result(n, -1); // Initialize result array with -1
    std::stack<int> stk; // Stack to store indices
```

```
    // Iterate through the array twice (to simulate circular behavior)
    for (int i = 0; i < 2 * n; ++i) {
        int num = nums[i % n];
```

```
        // Check for next greater element
        while (!stk.empty() && nums[stk.top()] < num) {
            result[stk.top()] = num;
            stk.pop();
        }
```

```
}

// Push index onto stack for first pass only
if (i < n) {
    stk.push(i);
}
}

return result;
}

int main() {
    std::vector<int> nums1 = {1, 2, 1};
    std::vector<int> result1 = nextGreaterElements(nums1);
    for (int num : result1) {
        std::cout << num << " ";
    }
    std::cout << std::endl; // Output: 2 -1 2

    std::vector<int> nums2 = {1, 2, 3, 4, 3};
    std::vector<int> result2 = nextGreaterElements(nums2);
    for (int num : result2) {
        std::cout << num << " ";
    }
    std::cout << std::endl; // Output: 2 3 4 -1 4

    return 0;
}
```

OUTPUT:

```
2 -1 2
2 3 4 -1 4
```

HARD

Q. You are given an array of integers nums, there is a sliding window of size k which is moving from the very left of the array to the very right. You can only see the k numbers in the window. Each time the sliding window moves right by one position. Return the max sliding window.

Example 1:

Input: nums = [1,3,-1,-3,5,3,6,7], k = 3

Output: [3,3,5,5,6,7]

Explanation:

Window position	Max
[1 3 -1] -3 5 3 6 7	3
1 [3 -1 -3] 5 3 6 7	3



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```
1 3 [-1 -3 5] 3 6 7 5
1 3 -1 [-3 5 3] 6 7 5
1 3 -1 -3 [5 3 6] 7 6
1 3 -1 -3 5 [3 6 7] 7
```

Example 2:

Input: nums = [1], k = 1

Output: [1]

Constraints:

- $1 \leq \text{nums.length} \leq 105$
- $-104 \leq \text{nums}[i] \leq 104$
- $1 \leq k \leq \text{nums.length}$

Reference : <https://leetcode.com/problems/sliding-window-maximum/description/>

CODE:

```
#include <iostream>
#include <vector>
#include <deque>
```

```
std::vector<int> maxSlidingWindow(const std::vector<int>& nums, int k) {
    std::vector<int> result;
    std::deque<int> dq; // Will store indices of nums

    for (int i = 0; i < nums.size(); ++i) {
        // Remove elements out of the current window
        if (!dq.empty() && dq.front() == i - k) {
            dq.pop_front();
        }

        // Remove elements smaller than the current element from the back
        while (!dq.empty() && nums[dq.back()] < nums[i]) {
            dq.pop_back();
        }

        // Add the current element's index to the deque
        dq.push_back(i);

        // Add the maximum element of the current window to the result
        if (i >= k - 1) {
            result.push_back(nums[dq.front()]);
        }
    }

    return result;
}
```

```
int main() {
    std::vector<int> nums1 = {1, 3, -1, -3, 5, 3, 6, 7};
    int k1 = 3;
    std::vector<int> result1 = maxSlidingWindow(nums1, k1);
}
```



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```
for (int num : result1) {
    std::cout << num << " ";
}
std::cout << std::endl; // Output: 3 3 5 5 6 7

std::vector<int> nums2 = {1};
int k2 = 1;
std::vector<int> result2 = maxSlidingWindow(nums2, k2);
for (int num : result2) {
    std::cout << num << " ";
}
std::cout << std::endl; // Output: 1

return 0;
}
```

OUTPUT:

```
3 3 5 5 6 7
1
```

VERY HARD

21. There are a number of plants in a garden. Each of the plants has been treated with some amount of pesticide. After each day, if any plant has more pesticide than the plant on its left, being weaker than the left one, it dies.

You are given the initial values of the pesticide in each of the plants. Determine the number of days after which no plant dies, i.e. the time after which there is no plant with more pesticide content than the plant to its left.

Example 1

$p = [3, 6, 2, 7, 5]$

// pesticide levels

Use a 1-indexed array. On day 1, plants 2 and 4 die leaving $p' = [3, 2, 5]$. On day 2, plant 3 in p' dies leaving $p'' = [3, 2]$. There is no plant with a higher concentration of pesticide than the one to its left, so plants stop dying after day 2.

Function Description

Complete the function `poisonousPlants` in the editor below.

`poisonousPlants` has the following parameter(s):

`int p[n]`: the pesticide levels in each plant

Returns

- `int`: the number of days until plants no longer die from pesticide

Input Format



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- The first line contains an integer n , the size of the array p .
- The next line contains n space-separated integers $p[i]$.

Constraints

- $1 \leq n \leq 10^5$
- $1 \leq p[i] \leq 10^9$

Example 2:

Sample Input

```
7
6 5 8 4 7 10 9
```

Sample Output

```
2
```

Explanation

Initially all plants are alive.

Plants = {(6,1), (5,2), (8,3), (4,4), (7,5), (10,6), (9,7)}

Plants[k] = (i,j) \Rightarrow j^{th} plant has pesticide amount = i.

After the 1st day, 4 plants remain as plants 3, 5, and 6 die.

Plants = {(6,1), (5,2), (4,4), (9,7)}

After the 2nd day, 3 plants survive as plant 7 dies.

Plants = {(6,1), (5,2), (4,4)}

Plants stop dying after the 2nd day.

References: <https://www.hackerrank.com/challenges/poisonous-plants/problem?isFullScreen=true>

CODE:

```
#include <iostream>
#include <vector>
#include <stack>
#include <algorithm>
```

```
int poisonousPlants(const std::vector<int>& p) {
    int n = p.size();
    std::vector<int> days(n, 0); // Days each plant takes to die
    std::stack<int> s; // Stack to keep track of plant indices

    int maxDays = 0; // Maximum days required for plants to stop dying
```



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```
for (int i = 0; i < n; ++i) {
    // Check if current plant will die
    int day = 0;
    while (!s.empty() && p[s.top()] >= p[i]) {
        day = std::max(day, days[s.top()]);
        s.pop();
    }

    // If stack is not empty, current plant will die
    if (!s.empty()) {
        days[i] = day + 1;
    }

    maxDays = std::max(maxDays, days[i]);
    s.push(i);
}

return maxDays;
}

int main() {
    std::vector<int> p1 = {3, 6, 2, 7, 5};
    std::cout << poisonousPlants(p1) << std::endl; // Output: 2

    std::vector<int> p2 = {6, 5, 8, 4, 7, 10, 9};
    std::cout << poisonousPlants(p2) << std::endl; // Output: 2

    return 0;
}
```

OUTPUT:

```
2
2
```