DOMAIN WINTER WINNING CAMP 2024

Student Name: Ayush Raj UID :22BCS10240

Branch: CSE Section/Group: 22BCS_FL_IOT-603/A

Semester: 6th

VERY EASY

Flood Fill - link

You are given an image represented by an m x n grid of integers image, where image[i][j] represents the pixel value of the image. You are also given three integers sr, sc, and color. Your task is to perform a flood fill on the image starting from the pixel image[sr][sc].

To perform a flood fill:

Begin with the starting pixel and change its color to color.

Perform the same process for each pixel that is directly adjacent (pixels that share a side with the original pixel, either horizontally or vertically) and shares the same color as the starting pixel.

Keep repeating this process by checking neighboring pixels of the updated pixels and modifying their color if it matches the original color of the starting pixel.

The process stops when there are no more adjacent pixels of the original color to update.

Return the modified image after performing the flood fill.

Example 1:

Input: image = [[1,1,1],[1,1,0],[1,0,1]], sr = 1, sc = 1, color = 2

Output: [[2,2,2],[2,2,0],[2,0,1]]

Explanation:

From the center of the image with position (sr, sc) = (1, 1) (i.e., the red pixel), all pixels connected by a path of the same color as the starting pixel (i.e., the blue pixels) are colored with the new color.

Note the bottom corner is not colored 2, because it is not horizontally or vertically connected to the starting pixel.

Example 2:

Input: image = [[0,0,0],[0,0,0]], sr = 0, sc = 0, color = 0

Output: [[0,0,0],[0,0,0]]

Explanation:

The starting pixel is already colored with 0, which is the same as the target color. Therefore, no changes are made to the image.

Constraints:

```
• m == image.length
```

- n == image[i].length
- $1 \le m, n \le 50$
- $0 \le \text{image}[i][j], \text{color} \le 2^16$
- $0 \le sr \le m$
- $0 \le sc \le n$

CODE:

```
def floodFill(image, sr, sc, color):
  rows, cols = len(image), len(image[0])
  original color = image[sr][sc]
  if original color == color:
     return image
  def dfs(r, c):
     if r < 0 or r >= rows or c < 0 or c >= cols or image[r][c]!= original color:
       return
     image[r][c] = color
     dfs(r+1, c)
     dfs(r - 1, c)
     dfs(r, c + 1)
     dfs(r, c - 1)
  dfs(sr, sc)
  return image
image1 = [[1, 1, 1], [1, 1, 0], [1, 0, 1]]
sr1, sc1, color1 = 1, 1, 2
print(floodFill(image1, sr1, sc1, color1))
```

Discover. Learn. Empower.

Output

[[2, 2, 2], [2, 2, 0], [2, 0, 1]]

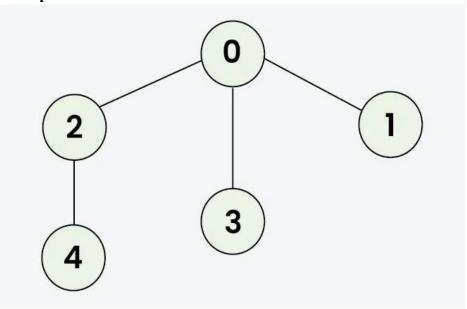
=== Code Execution Successful ===

EASY DFS of Graph

Given a connected undirected graph represented by an adjacency list adj, which is a vector of vectors where each adj[i] represents the list of vertices connected to vertex i. Perform a Depth First Traversal (DFS) starting from vertex 0, visiting vertices from left to right as per the adjacency list, and return a list containing the DFS traversal of the graph.

Note: Do traverse in the same order as they are in the adjacency list.

Example 1:



Input: adj = [[2,3,1], [0], [0,4], [0], [2]]

Output: [0, 2, 4, 3, 1]

Explanation: Starting from 0, the DFS traversal proceeds as follows:

Visit $0 \rightarrow \text{Output: } 0$

Visit 2 (the first neighbor of 0) \rightarrow Output: 0, 2

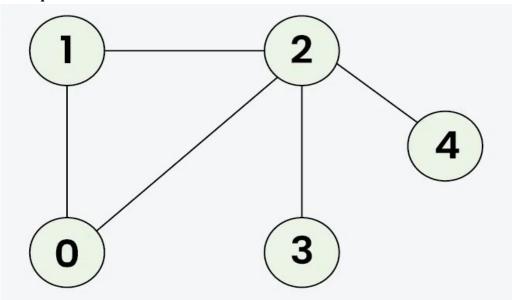
Visit 4 (the first neighbor of 2) \rightarrow Output: 0, 2, 4

Backtrack to 2, then backtrack to 0, and visit $3 \rightarrow \text{Output: } 0, 2, 4, 3$

Discover. Learn. Empower.

Finally, backtrack to 0 and visit $1 \rightarrow$ Final Output: 0, 2, 4, 3, 1

Example 2:



Input: adj = [[1, 2], [0, 2], [0, 1, 3, 4], [2], [2]]

Output: [0, 1, 2, 3, 4]

Explanation: Starting from 0, the DFS traversal proceeds as follows:

Visit $0 \rightarrow \text{Output: } 0$

Visit 1 (the first neighbor of 0) \rightarrow Output: 0, 1

Visit 2 (the first neighbor of 1) \rightarrow Output: 0, 1, 2

Visit 3 (the first neighbor of 2) \rightarrow Output: 0, 1, 2, 3

Backtrack to 2 and visit $4 \rightarrow$ Final Output: 0, 1, 2, 3, 4

Constraints:

- $1 \le adj.size() \le 1e4$
- $1 \le adj[i][j] \le 1e4$

CODE:

```
def dfs_traversal(adj):
    def dfs(node, visited, result):

    visited[node] = True
    result.append(node)

    for neighbor in adj[node]:
        if not visited[neighbor]:
            dfs(neighbor, visited, result)

    visited = [False] * len(adj)
    result = []
```

```
dfs(0, visited, result)
```

```
return result

adj = [[2, 3, 1], [0], [0, 4], [0], [2]]

output = dfs_traversal(adj)

print(output)
```

```
Output

[0, 2, 4, 3, 1]

=== Code Execution Successful ===
```

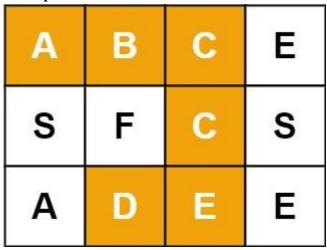
Medium

Word Search

Given an m x n grid of characters board and a string word, return true if word exists in the grid.

The word can be constructed from letters of sequentially adjacent cells, where adjacent cells are horizontally or vertically neighboring. The same letter cell may not be used more than once.

Example 1:



Input: board = [["A","B","C","E"],["S","F","C","S"],["A","D","E","E"]], word = "ABCCED"

Output: true

Example 2:

Α	В	С	Е
S	F	С	S
Α	D	Е	Е

Input: board = [["A","B","C","E"],["S","F","C","S"],["A","D","E","E"]], word = "SEE"

Output: true

Example 3:

Α	В	С	Е
S	F	С	S
Α	D	Е	Е

Input: board = [["A","B","C","E"],["S","F","C","S"],["A","D","E","E"]], word = "ABCB"

Output: false

Constraints:

- m == board.length
- n = board[i].length
- $1 \le m, n \le 6$
- 1 <= word.length <= 15
- board and word consists of only lowercase and uppercase English letters.

Follow up: Could you use search pruning to make your solution faster with a larger board?

CODE:

```
def exist(board, word):
  rows, cols = len(board), len(board[0])
  def dfs(r, c, index):
     if index == len(word):
       return True
     if r < 0 or r \ge rows or c < 0 or c \ge rows or board[r][c] != word[index]:
       return False
     temp, board[r][c] = board[r][c], '#'
     found = (dfs(r + 1, c, index + 1)) or
          dfs(r - 1, c, index + 1) or
          dfs(r, c + 1, index + 1) or
          dfs(r, c - 1, index + 1))
     board[r][c] = temp
     return found
  for row in range(rows):
     for col in range(cols):
       if board[row][col] == word[0] and dfs(row, col, 0):
          return True
  return False
board1 = [["A", "B", "C", "E"],
      ["S", "F", "C", "S"],
      ["A", "D", "E", "E"]]
word1 = "ABCCED"
print(exist(board1, word1)) # Output: True
   Output
 True
 === Code Execution Successful ===
```

Hard

Rotting Oranges

You are given an m x n grid where each cell can have one of three values:

0 representing an empty cell,

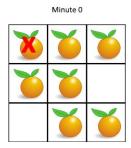
1 representing a fresh orange, or

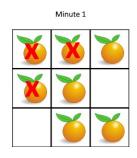
2 representing a rotten orange.

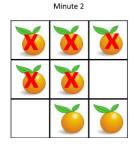
Every minute, any fresh orange that is 4-directionally adjacent to a rotten orange becomes rotten.

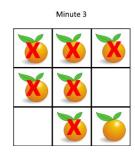
Return the minimum number of minutes that must elapse until no cell has a fresh orange. If this is impossible, return -1.

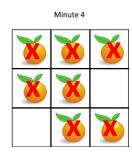
Example 1:











Input: grid = [[2,1,1],[1,1,0],[0,1,1]]

Output: 4

Example 2:

Input: grid = [[2,1,1],[0,1,1],[1,0,1]]

Output: -1

Explanation: The orange in the bottom left corner (row 2, column 0) is never rotten, because rotting only happens 4-directionally.

Example 3:

Input: grid = [[0,2]]

 $\textbf{Output:}\ 0$

Explanation: Since there are already no fresh oranges at minute 0, the answer is just 0.

Constraints:

- m == grid.length
- n == grid[i].length
- $1 \le m, n \le 10$
- grid[i][j] is 0, 1, or 2.

CODE:

from collections import deque

```
def orangesRotting(grid):
  rows, cols = len(grid), len(grid[0])
  queue = deque()
  fresh count = 0
  for r in range(rows):
    for c in range(cols):
       if grid[r][c] == 2:
         queue.append((r, c))
       elif grid[r][c] == 1:
         fresh count += 1
  if fresh count == 0:
    return 0
  minutes passed = 0
  directions = [(1, 0), (-1, 0), (0, 1), (0, -1)]
  while queue:
    for in range(len(queue)):
      x, y = queue.popleft()
       for dx, dy in directions:
         nx, ny = x + dx, y + dy
         grid[nx][ny] = 2
           queue.append((nx, ny))
           fresh count -= 1
    minutes passed += 1
  return minutes passed - 1 if fresh count == 0 else -1
grid1 = [[2, 1, 1], [1, 1, 0], [0, 1, 1]]
print(orangesRotting(grid1)) # Output: 4
```

Output

4

=== Code Execution Successful ===

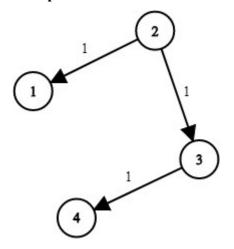
Very Hard

Network Delay Time

You are given a network of n nodes, labeled from 1 to n. You are also given times, a list of travel times as directed edges times[i] = (ui, vi, wi), where ui is the source node, vi is the target node, and wi is the time it takes for a signal to travel from source to target.

We will send a signal from a given node k. Return the minimum time it takes for all the n nodes to receive the signal. If it is impossible for all the n nodes to receive the signal, return -1.

Example 1:



Input: times = [[2,1,1],[2,3,1],[3,4,1]], n = 4, k = 2

Output: 2

Example 2:

Input: times = [[1,2,1]], n = 2, k = 1

Output: 1

CHANDIGARH UNIVERSITY Discover. Learn. Empower.

Example 3:

Input: times = [[1,2,1]], n = 2, k = 2 **Output**: -1

Constraints:

- $1 \le k \le n \le 100$
- 1 <= times.length <= 6000
- times[i].length == 3
- $1 \le ui, vi \le n$
- ui != vi
- $0 \le wi \le 100$
- All the pairs (ui, vi) are unique. (i.e., no multiple edges.)

CODE:

```
import heapq
from collections import defaultdict
def networkDelayTime(times, n, k):
  graph = defaultdict(list)
  for u, v, w in times:
     graph[u].append((v, w))
  min heap = [(0, k)]
  shortest times = \{\}
  while min heap:
     current time, node = heapq.heappop(min heap)
     if node in shortest times:
       continue
     shortest times[node] = current time
     for neighbor, weight in graph[node]:
        if neighbor not in shortest times:
          heapq.heappush(min_heap, (current_time + weight, neighbor))
  return max(shortest times.values()) if len(shortest times) == n else -1
times1 = [[2, 1, 1], [2, 3, 1], [3, 4, 1]]
n1, k1 = 4, 2
print(networkDelayTime(times1, n1, k1)) # Output: 2
```

Output

2

=== Code Execution Successful ===