排序

希尔排序: Pratt序列

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Pratt's Sequence, 1971

$$\mathcal{H}_{pratt} = \{ 2^p \cdot 3^q \mid p, q \in \mathcal{N} \}$$

$$= \{ 1, 2, 3, 4, 6, 8, 9, 12, 16, 18, 24, 27, 32, 36, \dots \}$$

❖ Note that

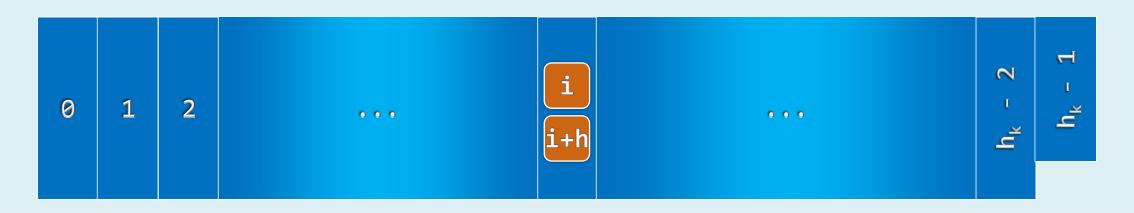
- adjacent items are NOT always relatively prime and
- there are $\mathcal{O}(\log^2 n)$ items no greater than n
- ❖ With Apratt,

Shellsort sorts a sequence of length n in $\mathcal{O}(n \cdot \log^2 n)$ time ...

From (2,3)-ordered to 1-ordered

$$x(2,3) = 1$$

- ∴ To the right of each element in a (2,3)-ordered sequence, only the NEXT element can be smaller
- \therefore It costs $\mathcal{O}(n)$ time to sort such a sequence



From $(2*h_k, 3*h_k)$ -ordered to h_k -ordered

- \clubsuit Divide S into h_k subsequences, each of which is (2,3)-ordered
 - \therefore it costs altogether $\mathcal{O}(n)$ time to sort them resp.
- * : there are altogether $\mathcal{O}(\log^2 n)$ iterations
 - ightharpoonup we need $\mathcal{O}(n \cdot \log^2 n)$ time

