

14-C3

排序

希尔排序：PS序列

邓俊辉

deng@tsinghua.edu.cn

d-Sorting an $\mathcal{O}(d)$ -Ordered Sequence in $\mathcal{O}(dn)$ Time

❖ If g and h are relatively prime and are both in $\mathcal{O}(d)$

we can d-sort the sequence in $\mathcal{O}(dn)$ time ...

- re-arrange the sequence as a 2D matrix with d columns
- each element is swapped with $\mathcal{O}((g-1) \cdot (h-1)/d) = \mathcal{O}(d)$ elements

inversion free

could be greater than $S[i]$ $[i]$

← $gh - g - h$ →

❖ Since this holds for all elements, $\mathcal{O}(dn)$ steps are enough

PS Sequence

❖ Papernov & Stasevic, 1965 //also called Hibbard's sequence

$$\mathcal{H}_{PS} = \mathcal{H}_{Shell} - 1 = \{ 2^k - 1 \mid k \in \mathcal{N} \} = \{ 1, 3, 7, 15, 31, 63, 127, 255, \dots \}$$

❖ Different items **MAY NOT** be relatively prime, e.g., $h_{2k} = h_k \cdot (h_k + 2)$

But ADJACENT items **MUST** be, since $h_{k+1} - 2 \cdot h_k \equiv 1$

❖ Shell sort with \mathcal{H}_{ps} needs

- $\mathcal{O}(\log n)$ outer iterations and
- $\mathcal{O}(n^{3/2})$ time to sort a sequence of length n //Why ...

$$t < k$$

❖ Let h_t be the h closest to \sqrt{n} and hence $h_t \approx \sqrt{n} = \Theta(n^{1/2})$

1) Consider those iterations for $\{h_k \mid t < k\} = \overline{\{h_{t+1}, h_{t+2}, \dots, h_m\}}$

∴ there would be $\mathcal{O}(n/h_k)$ elements in each of the h_k columns

∴ we can **insertionsort** each column in $\mathcal{O}((n/h_k)^2)$ time

∴ each h_k -sorting costs $\mathcal{O}(n^2/h_k)$ time

∴ all these iterations cost time of
 $\mathcal{O}(2 \times n^2/h_t) = \mathcal{O}(n^{3/2})$

$$k \leq t$$

$$h_k \leq h_t$$

$$k = t$$

$$h_k = h_t$$

$$t < k$$

$$h_t < h_k$$

$$k \leq t$$

2) Consider those iterations for $\{ h_k \mid k \leq t \} = \{ \overleftarrow{h_1, h_2, \dots, h_t} \}$

$\therefore h_{k+1}$ and h_{k+2} are relatively prime and are both in $\mathcal{O}(h_k)$

\therefore each h_k -sorting costs $\mathcal{O}(n \times h_k)$ time

\therefore all these iterations cost $\mathcal{O}(n \times 2 \cdot h_t) = \mathcal{O}(n^{3/2})$ time

❖ This upper bound is TIGHT

❖ What about the average cases?

- $\mathcal{O}(n^{5/4})$ based on simulations
- but not proved yet

$$k \leq t$$
$$h_k \leq h_t$$

$$k = t$$
$$h_k = h_t$$

$$t < k$$
$$h_t < h_k$$