# Lab #3. F - Interpreter

Prof. Jaeseung Choi

Dept. of Computer Science and Engineering

Sogang University



#### **General Information**

- Check the Assignment tab of Cyber Campus
  - Skeleton code (Lab3.tgz) is attached together with this slide
  - Submission will be accepted in the same post, too
- **■** Deadline: 5/10 Friday 23:59
  - Late submission deadline: 5/12 Sunday 23:59 (-20% penalty)
  - Delay penalty is applied uniformly (not problem by problem)
- Please read the instructions in this slide carefully
  - This slide is a step-by-step tutorial for the lab
  - It also contains important submission guidelines
    - If you do not follow the guidelines, you will get penalty

#### **Skeleton Code Structure**

- **■** Copy Lab3.tgz into CSPRO server and decompress it
  - This course will use <a href="mailto:cspro2.sogang.ac.kr">cspro2.sogang.ac.kr</a> (don't miss the 2)
  - Don't decompress-and-copy; copy-and-decompress
- **FMinus:** Only one directory for F- this time
- check.py, config: Script and config file for self-grading (same as before)

```
jschoi@cspro2:~$ tar -xzf Lab3.tgz
jschoi@cspro2:~$ cd Lab3/
jschoi@cspro2:~/Lab3$ ls
FMinus check.py config
```

#### **Directory Structure of FMinus**

- Skeleton code structure under src/ is same as before
  - AST.fs: Syntax definition of the F- language
  - FMinus.fs: You have to implement the semantics here
  - Types.fs: Type definitions needed for semantics
  - Main.fs: Main driver code of the interpreter
  - Lexer.fsl, Parser.fsy: Parser (you don't have to care)
- Do NOT fix any source files other than FMinus.fs

```
jschoi@cspro2:~/Lab3$ cd FMinus/
jschoi@cspro2:~/Lab3/FMinus$ ls
FMinus.fsproj src testcase
jschoi@cspro2:~/Lab3/FMinus$ ls src
AST.fs FMinus.fs Lexer.fsl Main.fs Parser.fsy Types.fs
```

#### F - Language Syntax

- Similar to the extended version of F- in our lecture slide
  - The whole program is an expression

```
E \rightarrow n
    | true | false
   |E+E|E-E
    \mid E < E \mid E > E \mid E = E \mid E <> E
    | if E then E else E
    |  let x = E  in E
    |  let f x = E  in E
    | \mathbf{let} \operatorname{rec} f x = E \mathbf{in} E
    fun x \to E
    \mid E\mid E
```

#### **Expression**

- Relation  $\rho \vdash e \Downarrow v$  defines the evaluation of expression
  - We use the same semantic domain to our lecture slide
  - $lacktriangledown 
    ho \in Env = Var 
    ightarrow Val, v \in Val = Z + B + Func + RecFunc$

$$\overline{\rho \vdash n \Downarrow n} \qquad \overline{\rho \vdash \text{true} \Downarrow true} \qquad \overline{\rho \vdash \text{false} \Downarrow false} \qquad \overline{\rho \vdash x \Downarrow \rho(x)}$$

$$\underline{\rho \vdash e_1 \Downarrow n_1} \qquad \underline{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2} \qquad \underline{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2} \qquad \underline{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2} \qquad \underline{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2} \qquad \underline{\rho \vdash e_1 \vdash e_2 \Downarrow n_1 - n_2}$$

$$\frac{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2}{\rho \vdash e_1 < e_2 \Downarrow true} \quad n_1 < n_2$$

$$\frac{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2}{\rho \vdash e_1 > e_2 \Downarrow true} \quad n_1 > n_2$$

$$\frac{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2}{\rho \vdash e_1 < e_2 \Downarrow false} \quad n_1 \ge n_2$$

$$\frac{\rho \vdash e_1 \Downarrow n_1 \quad \rho \vdash e_2 \Downarrow n_2}{\rho \vdash e_1 > e_2 \Downarrow false} \quad n_1 \le n_2$$

- Relation  $\rho \vdash e \Downarrow v$  defines the evaluation of expression
  - We use the same semantic domain to our lecture slide
  - $lacktriangledown 
    ho \in Env = Var 
    ightarrow Val, v \in Val = Z + B + Func + RecFunc$

$$\frac{\rho \vdash e_1 \Downarrow v_1 \quad \rho \vdash e_2 \Downarrow v_2}{\rho \vdash e_1 = e_2 \Downarrow true} \ (v_1 = v_2 = n) \lor (v_1 = v_2 = b)$$

$$\frac{\rho \vdash e_1 \Downarrow v_1 \quad \rho \vdash e_2 \Downarrow v_2}{\rho \vdash e_1 = e_2 \Downarrow false} (v_1 = n_1 \neq n_2 = v_2) \lor (v_1 = b_1 \neq b_2 = v_2)$$

$$\frac{\rho \vdash e_1 \Downarrow v_1 \quad \rho \vdash e_2 \Downarrow v_2}{\rho \vdash e_1 \iff e_2 \Downarrow false} \ (v_1 = v_2 = n) \lor (v_1 = v_2 = b)$$

$$\frac{\rho \vdash e_1 \Downarrow v_1 \quad \rho \vdash e_2 \Downarrow v_2}{\rho \vdash e_1 <> e_2 \Downarrow true} (v_1 = n_1 \neq n_2 = v_2) \lor (v_1 = b_1 \neq b_2 = v_2)$$

- Relation  $\rho \vdash e \Downarrow v$  defines the evaluation of expression
  - We use the same semantic domain to our lecture slide
  - $lacktriangledown 
    ho \in Env = Var 
    ightarrow Val, v \in Val = Z + B + Func + RecFunc$

$$\frac{\rho \vdash e_1 \Downarrow true \quad \rho \vdash e_2 \Downarrow v}{\rho \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \Downarrow v}$$

$$\frac{\rho \vdash e_1 \Downarrow v_1 \quad \rho[x \mapsto v_1] \vdash e_2 \Downarrow v_2}{\rho \vdash \mathbf{let} \ x = e_1 \ \mathbf{in} \ e_2 \Downarrow v_2}$$

$$\frac{\rho[f \mapsto \langle f, x, e_1, \rho \rangle] \vdash e_2 \Downarrow v}{\rho \vdash \text{let rec } f \ x = e_1 \text{ in } e_2 \Downarrow v}$$

$$\frac{\rho \vdash e_1 \Downarrow false \quad \rho \vdash e_3 \Downarrow v}{\rho \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \Downarrow v}$$

$$\frac{\rho[f \mapsto \langle x, e_1, \rho \rangle] \vdash e_2 \Downarrow v}{\rho \vdash \mathbf{let} f \ x = e_1 \mathbf{in} \ e_2 \Downarrow v}$$

$$\overline{\rho \vdash \text{fun } x \to e \ \Downarrow \langle x, e, \rho \rangle}$$

- Relation  $\rho \vdash e \Downarrow v$  defines the evaluation of expression
  - We use the same semantic domain to our lecture slide
  - $lacktriangledown 
    ho \in Env = Var 
    ightarrow Val, v \in Val = Z + B + Func + RecFunc$

(Application of *non-recursive* function)

$$\frac{\rho \vdash e_1 \Downarrow \langle x, e_b, \rho' \rangle \qquad \rho \vdash e_2 \Downarrow v_{arg} \qquad \rho'[x \mapsto v_{arg}] \vdash e_b \Downarrow v}{\rho \vdash e_1 \quad e_2 \Downarrow v}$$

(Application of *recursive* function)

$$\frac{\rho \vdash e_1 \Downarrow \langle f, x, e_b, \rho' \rangle \quad \rho \vdash e_2 \Downarrow v_{arg} \quad \rho'[x \mapsto v_{arg}][f \mapsto \langle f, x, e_b, \rho' \rangle] \vdash e_b \Downarrow v}{\rho \vdash e_1 \quad e_2 \Downarrow v}$$

## Implementing Semantics

- To complete the interpreter of F-, you must implement the semantics of F- language in FMinus.fs file
  - You have to implement only one function: evalExp()
    - Execution of program = Evaluation of the expression
  - Type definition of Env and Val are provided in Types.fs
  - If the semantics of program is not defined, your interpreter must raise UndefinedSemantics exception defined in Types.fs

```
let rec evalExp (exp: Exp) (env: Env) : Val =
    ...

// This part is given for you.
let run (prog: Program) : Val =
    evalExp prog Map.empty
```

## **Building and Testing**

- In testcase directory, tc-\* and ans-\* files are provided
  - After compiling the interpreter with the dotnet build -o out command, you can run program written in F- language
  - The interpreter will print the evaluation result of the program

## **Tip: Printing AST**

- In F- language, you may feel confused about how the input program is parsed into AST
  - For example, is "f x + 1" parsed into (f x) + 1 or f (x + 1)?
  - You can temporarily add the following printfn() to print out the program AST (don't forget to erase it before the submission)

```
let run (prog: Program) : Val =
    printfn "%A" prog
    evalExp prog Map.empty
```

```
jschoi@cspro2:~/Lab3/FMinus$ cat parsing-test
f x + 1
jschoi@cspro2:~/Lab3/FMinus$ ./out/FMinus parsing-test
Add (App (Var "f", Var "x"), Num 1)
...
```

### **Self-Grading Script**

If you think you have solved all the problems, you can run check.py as a final check

```
'O': Correct, 'X': Incorrect, 'E': Unhandled exception in your code
```

- 'C': Compile error, 'T': Timeout (maybe infinite recursion)
- If you correctly raise UndefinedSemantics exception for an invalid program, it will be graded as '0' (not 'E')
  - If you raise UndefinedSemantics for valid program, it is 'X'

```
jschoi@cspro2:~/Lab3$ ls
FMinus check.py config
jschoi@cspro2:~/Lab3$ $ ./check.py
[*] FMinus : 0000
```

### **Actual Grading**

- I will use different test case set during the real grading
  - So you are encouraged to run you code with your own test cases (try to think of various inputs)
  - Some students ask me to provide more test cases, but it is important to practice this on your own
- You will get the point based on the number of test cases that you pass
  - 100 point in total (but recall that each lab has different weight)

#### **Submission Guideline**

- You should submit only one F# source code file
  - FMinus.fs (from Lab3/FMinus/src/FMinus.fs)
- If the submitted file fails to compile with skeleton code when I type "dotnet build", cannot give you any point
- Submission format
  - Upload this file directly to Cyber Campus (do not zip them)
  - Do not change the file name (e.g., adding any prefix or suffix)
  - If your submission format is wrong, you will get -20% penalty