

Programming GPUs for non-graphics workloads – *from General Purpose GPU (GPGPU) to GPU compute*



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Disclaimer

The author's views expressed in this presentation do not necessarily reflect the views of IBM.

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I would also like to thank Patrick Cozzi for inviting me to teach in his classes multiple years in a row and for letting me re-use his introductory material.

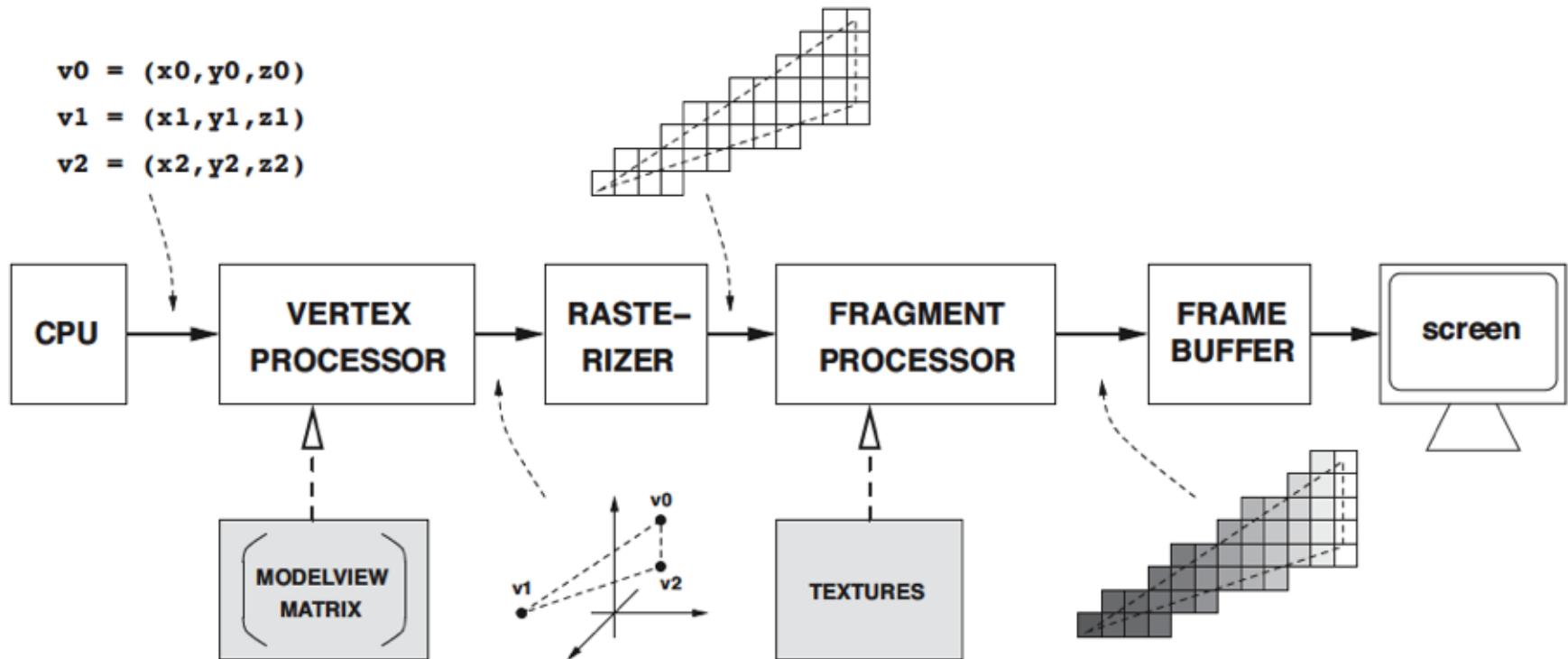
Agenda

Programming GPUs for non-graphics workloads

- GPGPU
 - A brief introduction
 - Index Search implemented using openGL and Cg
- Modern GPU computing with CUDA
 - A very brief intro to CUDA
 - Index Search in CUDA
 - Performance optimizations
 - A new GPU-optimal (index) search algorithm

GPU Programming pre-CUDA

■ The graphics pipeline



- Vertex Processor - geometric transformations of vertices in 3D space
- Rasterizer - transforms geometric primitives (triangles) into pixels
- Fragment Processor – colors the pixels
- Programmable were only vertex and fragment processors

GPGPU Programming

- GPGPU(.org) started in 2002 by Mark Harris
- Using Graphics APIs to solve non-graphics tasks
 - E.g. OpenGL & Cg
- Required use of graphics APIs
 - OpenGL for data transfers
 - Cg to “program”
 - Operations:
 - geometric transformations using the vertex processor (scatter)
 - coloring using the fragment processor (gather)
 - Vertices are stored as float4 (x,y,z,w)
 - Textures = 2D arrays of float4 vectors (r,g,b,a)
 - Compute = drawing

GPGPU Programming

Steps for GPGU compute:

1. Organize data in a screen size array
2. Set up a viewport with 1:1 pixel:texel ratio
3. Create and bind texture of the same size
4. Download input data into texture
5. Bind (load) fragment program (computational kernel)
6. Render a screen size quad to perform computation, i.e. run fragment program on each Pixel
7. Read back results

Let's Pick a Simple, but Omnipresent Task ... Search

- Why Search?
- Honestly, how many times a day do you visit:

  ?

Let's Pick a Simple, but Omnipresent Task ... Search

- Why Search?
- Honestly, how many times a day do you visit:



- How do you search (millions of) documents efficiently?
- Use an inverted index

	Keyword	DocID
Sorted ↓	Adam	1, 2, 3
	Bethlehem	4, 5
	Character	1, 2, 3, 301, 5790
	Drachenflieger	301, 317, 5790
	Eva	1, 2
	Flughafenbahnhof	5790
	Grabdenkmal	2, 5790
	Haubentaucher	300, 5790

Searching an Index

- The task: search an inverted (document) index

sorted ↓

Keyword	DocID
Adam	1, 2, 3
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Character	1, 2, 3, 301, 5790
Drachenflieger	301, 317, 5790
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Haubentaucher	300, 5790

{ 16 characters max.

Can be stored separately. Lookup by position.

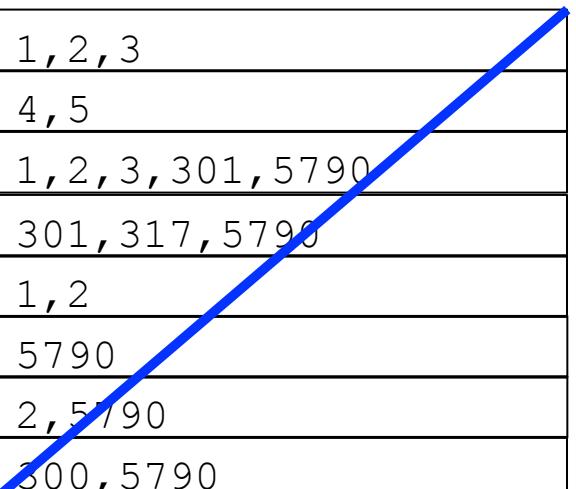
Searching an Index

- The task: search an inverted (document) index

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Keyword	DocID
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Grabdenkmal	2, 5790
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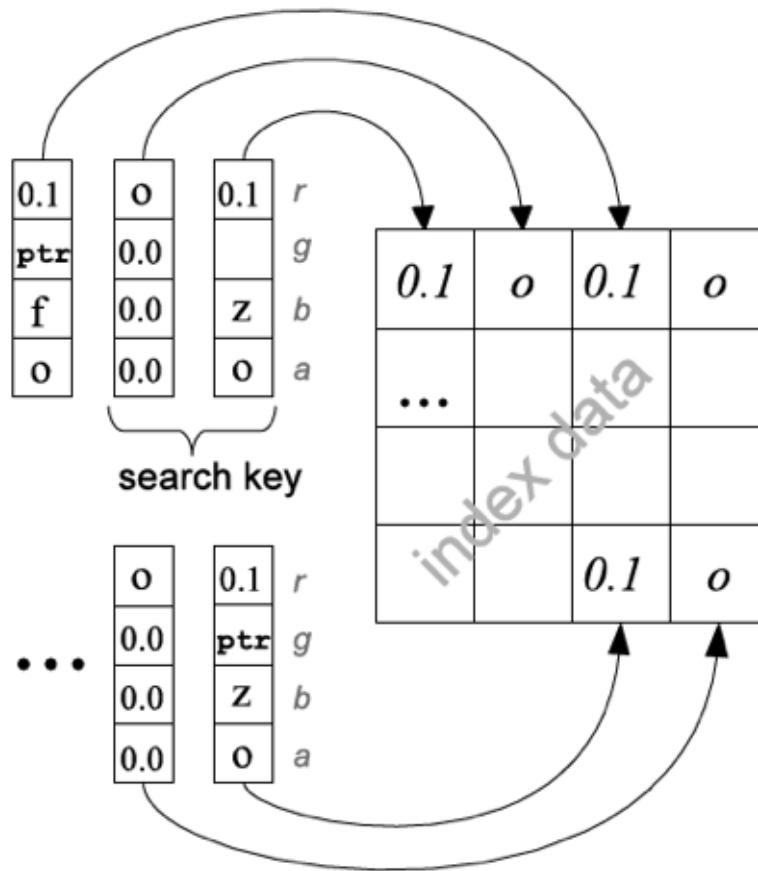
- On the CPU we use a few library calls and we are done

```
char searchkey[16] = "Flughafenbahnhof";
result = bsearch( (void*) &searchkey, index,
                  numentries, sizeof(char)*16,
                  (int(*)(const void*, const void*)) strcmp);
```

GPGPU Search – Data Format

- Storing data

- Obviously you want 1:1 pixel-to-texture element (texel) ratio
unless you would like to play Scrabble ;-)



- Ascii mapped to 0.0 to 255.0
- 1 pixel stores 4 chars (better?)
- Mark beginning of words with 0.1
- Need to store pointer/position in document index ptr
- Align word boundaries with pixel boundaries r
- Null-terminated strings 0.0

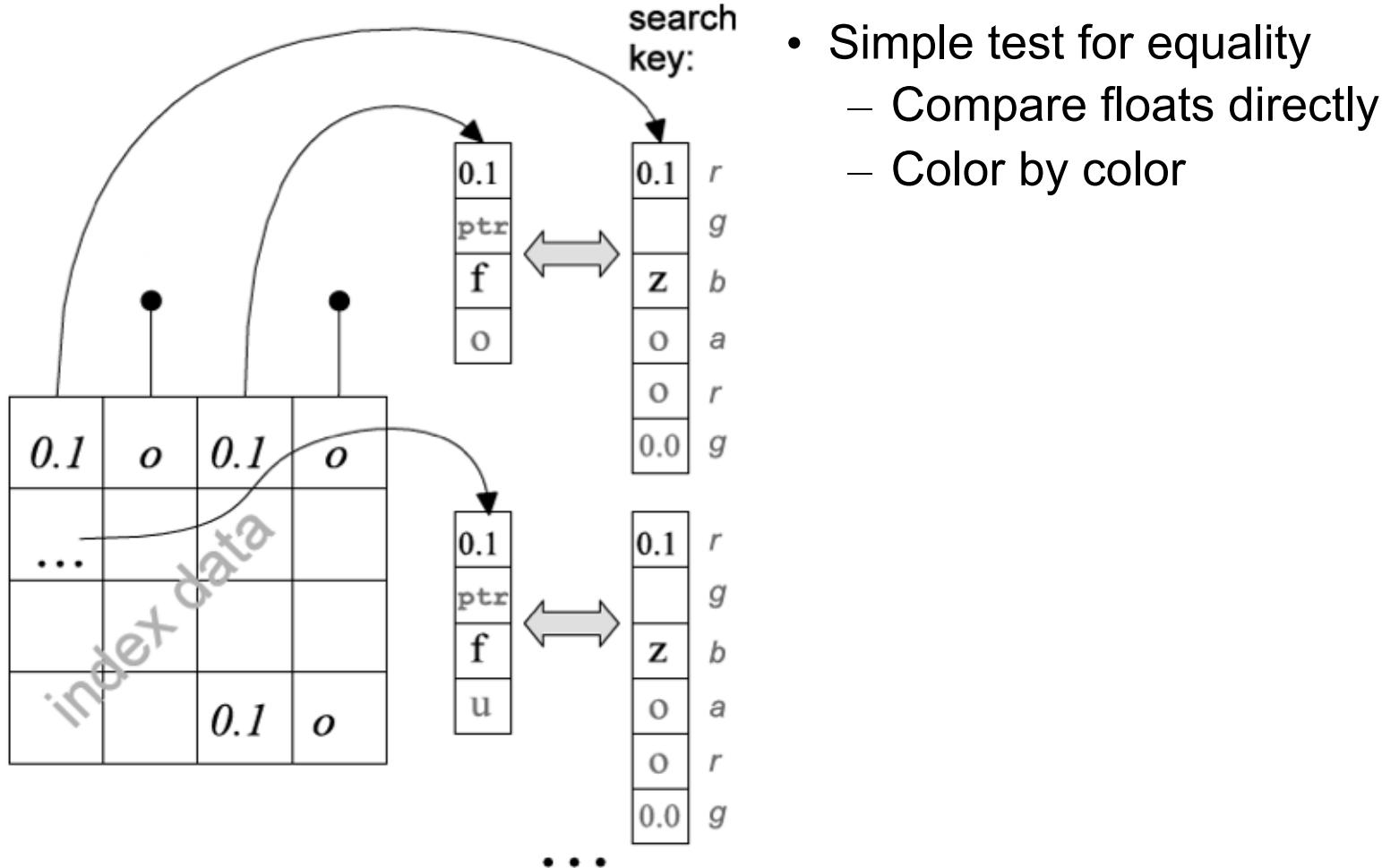
GPGPU Search – Data Storage

- Store data in texture

```
float* data = malloc(sizeof(float)*1200*1200*4);  
...  
data[pos++] = 0.1;  
data[pos++] = *(float*)&docindex;  
for (i=0;i<=strlen(currentString);i++) {  
    data[pos++] = (float)currentString[i];  
}  
...  
glTexSubImage2D(GL_TEXTURE_RECTANGLE_ARB,  
                 0,0,0, // detail level, x-, y- offset  
                 1200, 1200, // size  
                 GL_RGBA, // texture format  
                 GL_FLOAT, // data format  
                 data); // data pointer
```

GPGPU Search in Action

- Comparing search key with stored strings



GPGPU Search Code

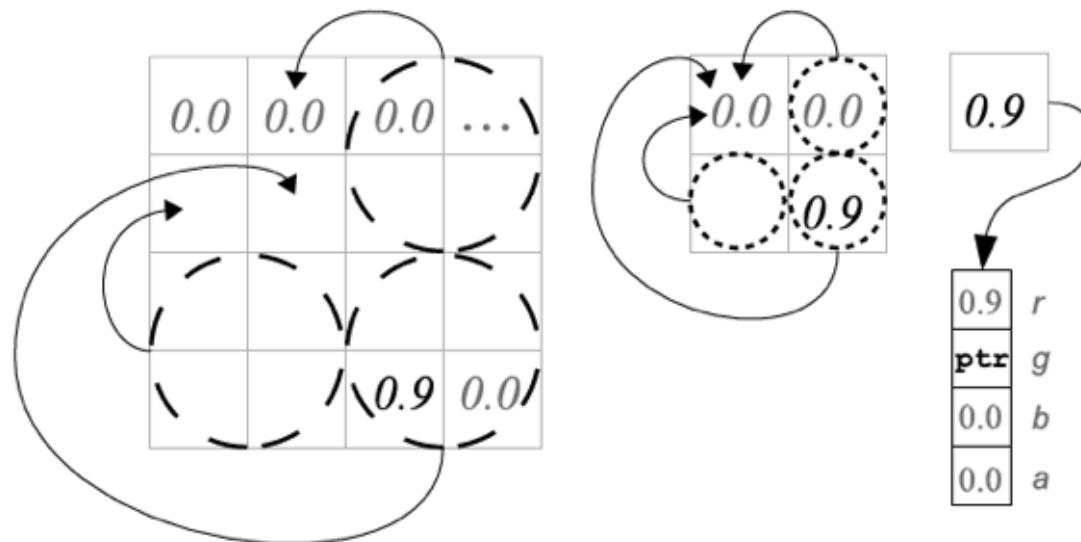
```
float4 search(float2 coords: WPOS,
              uniform samplerRECT texCgFrag) : COLOR {
    float2 data_coords = coords;
    float2 searchkey_coords = float2(0.5,0.5);
    float4 data = texRECT(texCgFrag, data_coords );
    float4 searchkey = texRECT(texCgFrag, searchkey_coords);
    float done =0.0;
    if (data.r == 0.1) {
        if (done == 0.0) {
            if (data.b != searchkey.b) done = -1.0;
            if (data.b == searchkey.b)
                if (data.b== 0.0) done = 1.0;
        }
        if (done == 0.0) {
            if (data.a != searchkey.a) done = -1.0;
            ...
        }
    }
}
```

GPGPU Search – Code Execution

- To execute the code:
`drawQuad(1200,1200);`
- Result uses a magic number (not used for ASCII mapping) **0.9**
- After completion Result is anywhere in the texture
- Copying whole texture back to main memory inefficient
- Reduction

GPGPU Search – Reduction

- To execute the code:
`drawQuad(1200,1200);`
- Result uses a magic number (not used for ASCII mapping) **0.9**
- After completion Result is anywhere in the texture
- Copying whole texture back to main memory inefficient
- Reduction:



GPGPU Search – Reduction

- Reduction means gathering “neighborhood” data

```
float4 reduce (float2 coords: WPOS,
               uniform samplerRECT texCgFrag2) : COLOR {
    float2 topleft = ((coords-0.5)*2.0)+0.5;
    float4 val1 = texRECT(texCgFrag2, topleft);
    float4 val2 = texRECT(texCgFrag2, topleft+float2(1,0));
    float4 val3 = texRECT(texCgFrag2, topleft+float2(1,1));
    float4 val4 = texRECT(texCgFrag2, topleft+float2(0,1));
    float4 result = (0.0,0.0,0.0,0.0);
    if (val4.r == 0.9) result = val4;
    if (val3.r == 0.9) result = val3;
    if (val2.r == 0.9) result = val2;
    if (val1.r == 0.9) result = val1;
    return result;
}
```

GPGPU Search – Reduction

- Repeat until we end up with a single pixel
- Search result will be in top left pixel

```
numPasses = (int)(log((double)width)/log(2.0));
```

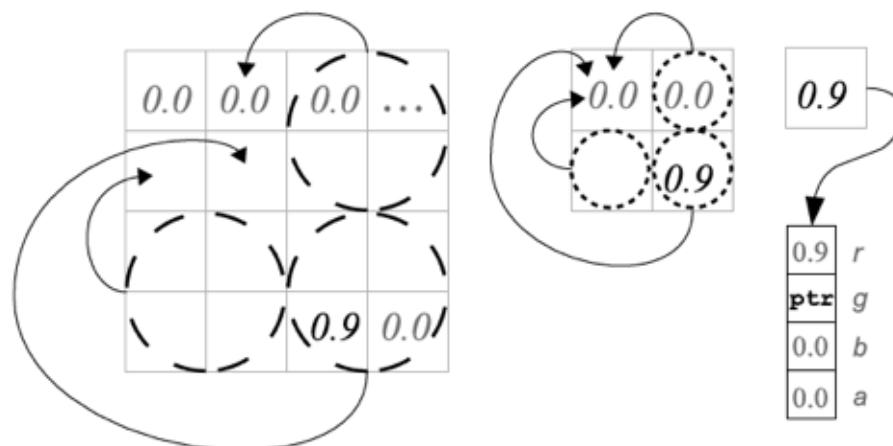
```
for (i=0; i<numPasses; i++) {
```

```
...
```

```
    outputWidth = outputWidth / 2;
```

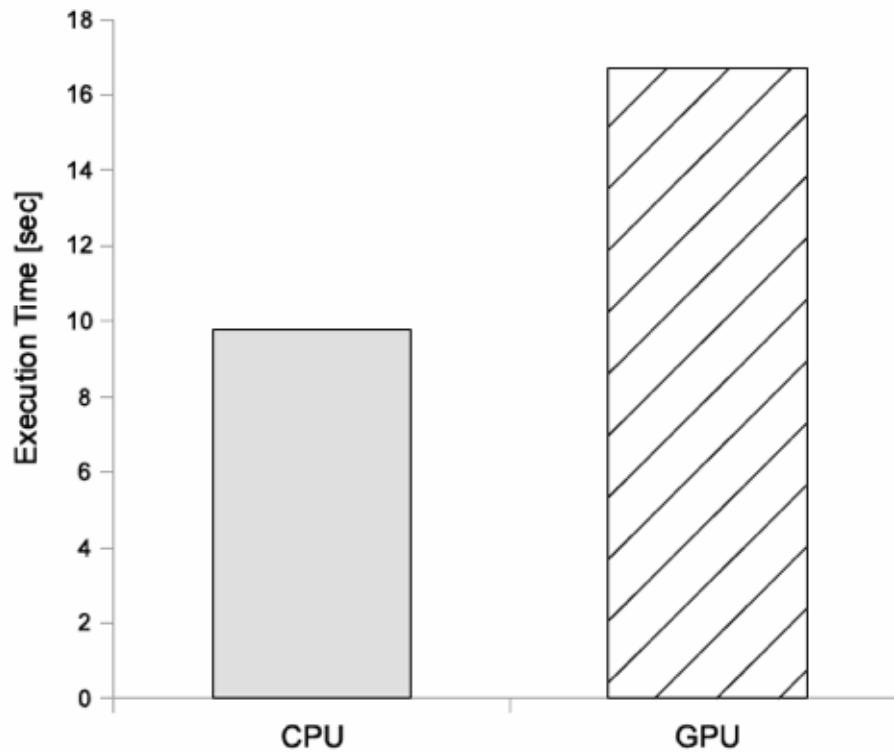
```
    drawQuad(outputWidth,outputWidth);
```

```
...
```



GPGPU Search - Performance

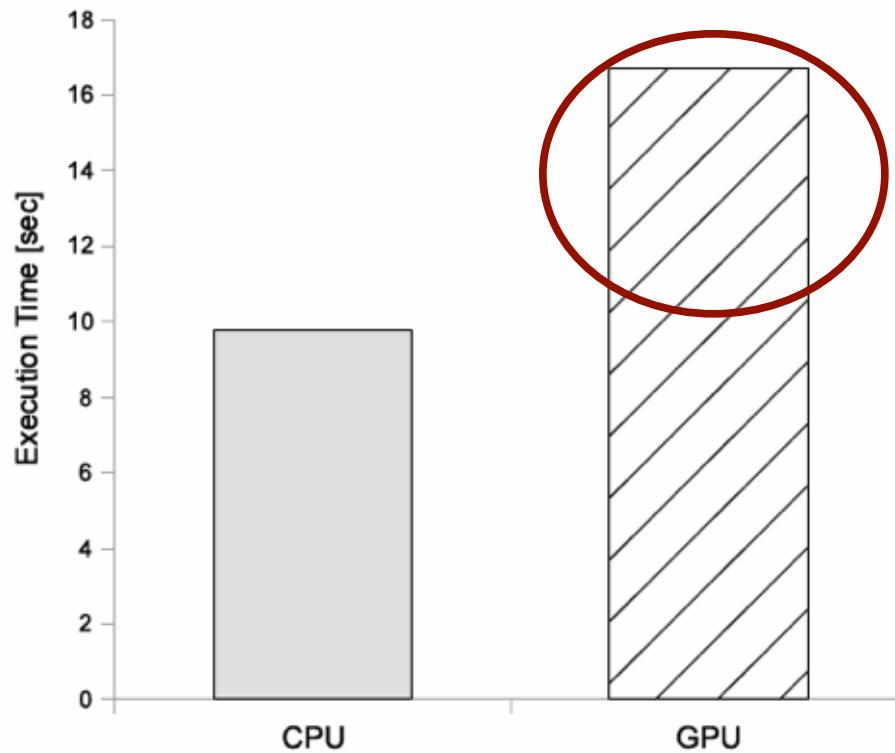
- 10k Berkeley DB index operations (insert delete), all require searching the index first, Test001.tcl
- Berkley DB uses B-trees, which needed to be flattened for the GPU



Time required for 10k insert/delete operations using a dual-core 2.2ghz AMD Opteron vs. an nVidia 7900GS with 7 vertex and 20 fragment processors.

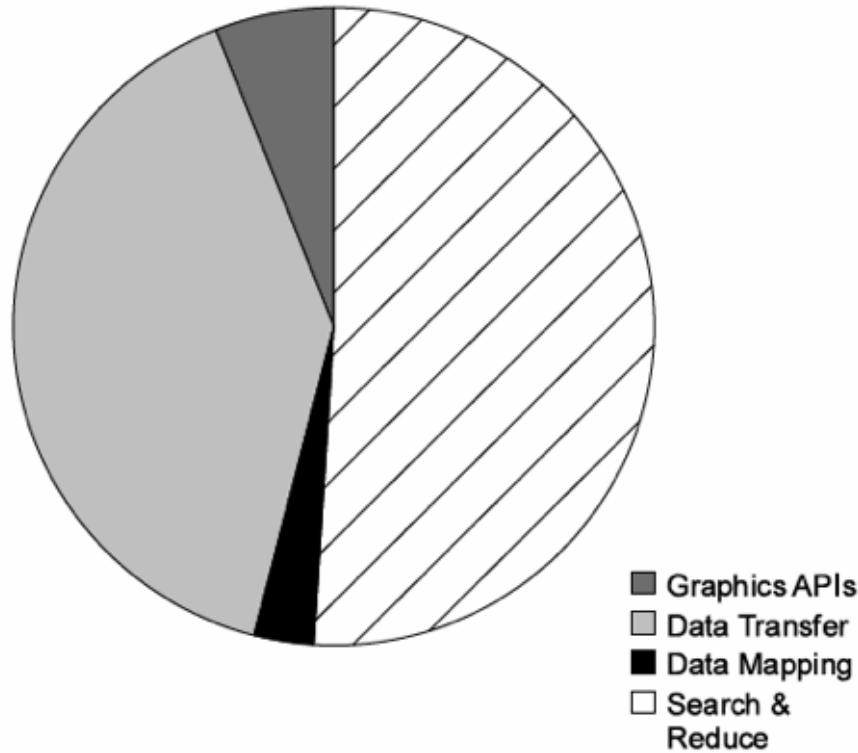
GPGPU Search - Performance

- 10k Berkeley DB index operations (insert delete), all require searching the index first, Test001.tcl
- Berkley DB uses B-trees, which needed to be flattened for the GPU



Time required for 10k insert/delete operations using a dual-core 2.2ghz AMD Opteron vs. an nVidia 7900GS with 7 vertex and 20 fragment processors.

GPGPU Search – Where does time go ?



Time required for 10k insert/delete operations using a dual-core 2.2ghz AMD Opteron vs. an nVidia 7900GS with 7 vertex and 20 fragment processors.

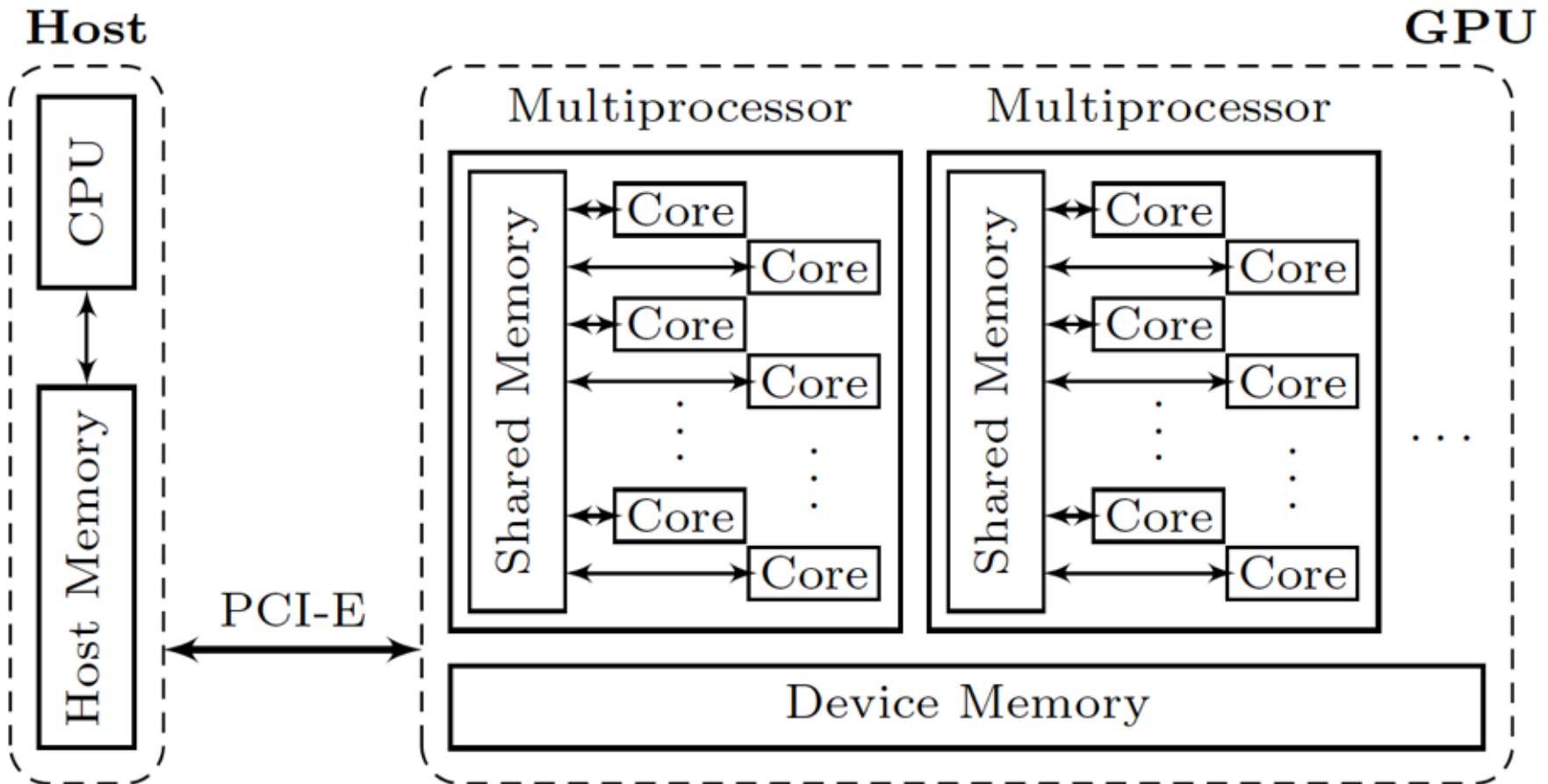
- Data Transfer ~40%
 - More efficient data mapping, e.g. 4 char = 1 float problematic?
- CUDA made GPGPU obsolete ...

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CUDA Key Concepts – Architecture



CUDA Key Concepts – Function Classifiers

- `__global__`
 - callable from host
 - must return void
- `__device__`
 - callable only from device
 - function inlined by default (newer CUDA versions)
- Global and device functions
 - No recursion (except Fermi)
 - No static variables
 - No malloc()
 - Careful with function calls through pointers (Fermi)
 - Cannot access host memory “directly”

CUDA Key Concepts – Memory address spaces

- Host (CPU) and Device (GPU) have separate (memory) address spaces
 - Data needs to be “transferred” to/from the GPU
 - Simplest way is to explicitly copy data to/from device memory
 - Data copy always initiated by host

```
cudaMemcpy(void* dst,  
          const void* src,  
          size_t count,  
          cudaMemcpyHostToDevice | cudaMemcpyDeviceToHost  
)
```

- Specify **direction** of data copy
- **ToDevice** for input data
- **ToHost** for results
- When calling **__global__** function pass **dst** pointer

CUDA Key concepts – Vector types

- `char[1-4], uchar[1-4], short[1-4], ushort[1-4], int[1-4], uint[1-4], long[1-4], ulong[1-4], longlong[1-2], ulonglong[1-2]`
- `float[1-4], double[1-2]`
- `dim3`
- **Available in host and device code**
- **Construct with `make_<type name>`**

```
int2 i2 = make_int2(1, 2);
float4 f4 = make_float4(
    1.0f, 2.0f, 3.0f, 4.0f);
```
- **Access with `.x, .y, .z, and .w`**

```
int2 i2 = make_int2(1, 2);
int x = i2.x;
int y = i2.y
```
- **No `.r, .g, .b, .a, etc.` like OpenGL, Cg**

CUDA Key Concepts – Invoking GPU Functions (Kernels)

```
__global__ void gpu_Kernel(int a, ...) {  
    ...  
}  
...  
  
dim3 grid(14,0,0);  
Dim3 block(192,0,0);  
gpu_Kernel<<<grid,block>>>(42,...);
```

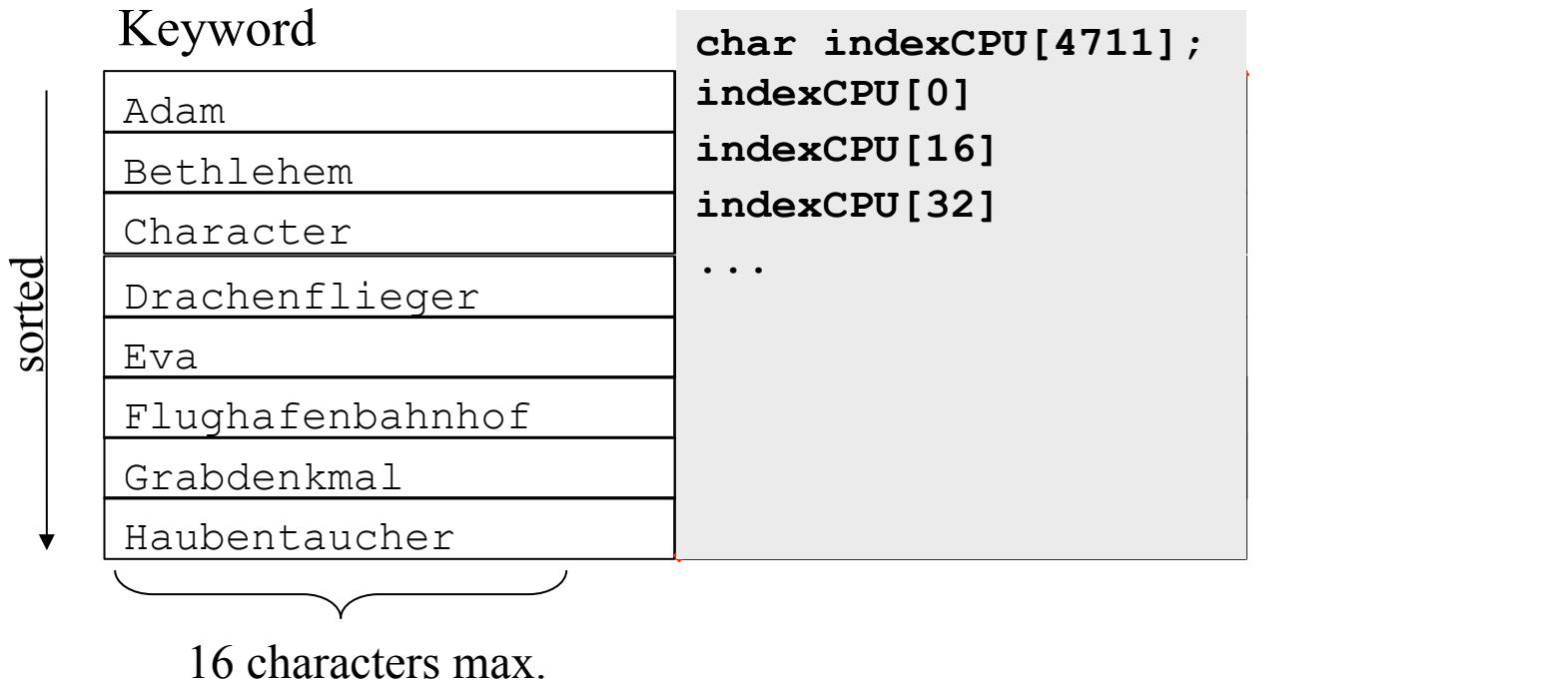
- Calling GPU (`__global__`) function requires to specify
 - grid dimensions – How many blocks of threads to launch
 - 1 block executes on 1 streaming multiprocessor to completion
 - block dimensions – How many threads are in a block
 - threads execute in groups of 32 (warps) in SIM[T/D] fashion
 - #threads > warp can be synchronized with `__syncthreads()`

CUDA Key Concepts – “Global” Variables

```
__device__ int a_dev;  
...  
__shared__ int a_smem;
```

- **__device__ variables**
 - stored in device memory
 - accessible from all blocks
- **__shared__ variables**
 - stored in shared on-chip memory (space constraints?)
 - accessible only within a block

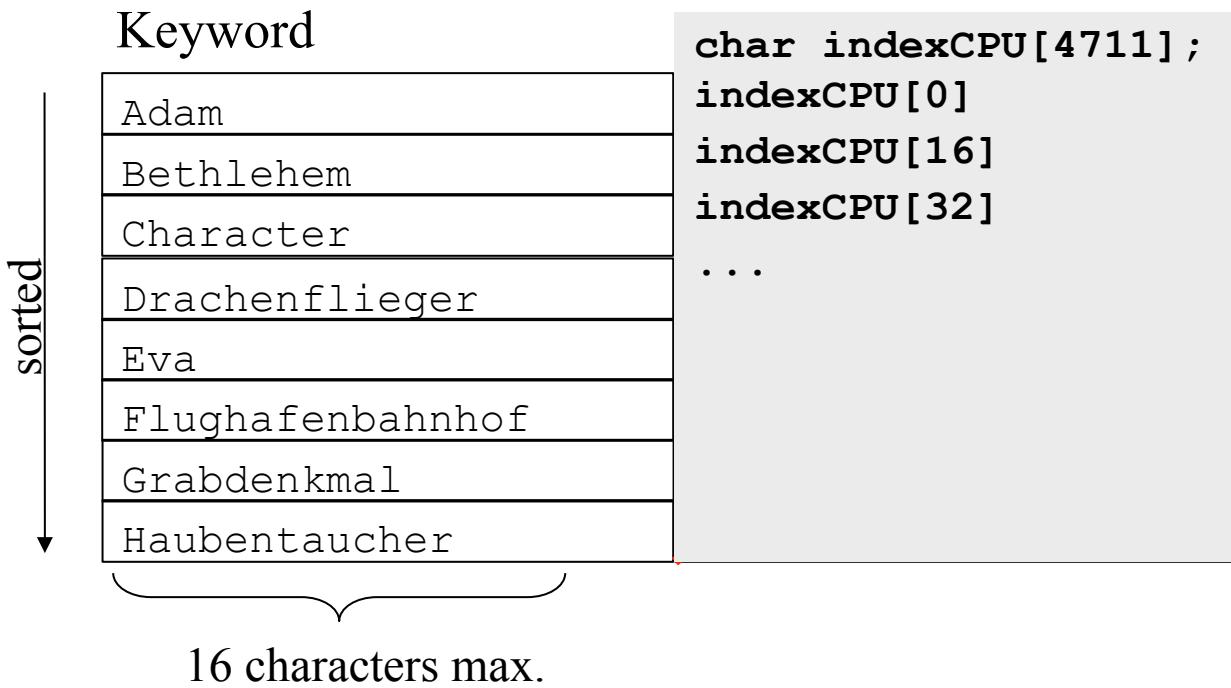
Index search on the CPU



- On the CPU we use a few library calls and we are done

```
char searchkey[16] = "Flughafenbahnhof";
result = bsearch((void*) searchkey, indexCPU,
                 numentries, sizeof(char)*16,
                 (int(*)(const void*, const void*)) strcmp);
```

A Simple implementation of (index) search



- On the CPU we use a few library calls and we are done

```
char searchkey[16] = "Flughafenbahnhof";
result = bsearch((void*) searchkey, indexCPU,
                 numentries, sizeof(char)*16,
                 (int(*)(const void*, const void*)) strcmp);
```

- Can we just port a CPU implementation?

Index search on the CPU

- Get the data to the GPU

```
char* indexGPU;
char* searchkeysGPU;
char* resultsGPU;
// copy the data
cudaMalloc((void**)&indexGPU, sizeof(char)*wordlength*entries);
cudaMemcpy(indexGPU, indexCPU, sizeof(char)*wordlength*entries,
           CudaMemcpyHostToDevice);
// copy the searchkey(s)
cudaMalloc((void**)&searchkeysGPU, ...
cudaMemcpy(searchkeysGPU, searchkeysCPU,
           sizeof(char)*wordlength*numsearches,
           CudaMemcpyHostToDevice);
// make room for the results
cudaMalloc((void**)&resultsGPU, ...
```

A Simple GPU implementation

- Get the data to the GPU

```
char* indexGPU;
char* searchkeysGPU;
char* resultsGPU;
// copy the data
cudaMalloc((void**)&indexGPU, sizeof(char)*wordlength*entries);
cudaMemcpy(indexGPU, indexCPU, sizeof(char)*wordlength*entries,
           CudaMemcpyHostToDevice);
// copy the searchkey(s)
cudaMalloc((void**)&searchkeysGPU, ...
cudaMemcpy(searchkeysGPU, searchkeysCPU,
           sizeof(char)*wordlength*numsearches,
           CudaMemcpyHostToDevice);
// make room for the results
cudaMalloc((void**)&resultsGPU, ...
```

- Know your hardware (GTX 285, 30 SMs, 8 cores each, 240 cores)
 - Set up an execution configuration & call global function

```
dim3 Dg = dim3(30,0,0);
dim3 Db = dim3(8,0,0);
searchGPU<<<Dg,Db>>>(indexGPU, entries...;
```

A Simple GPU implementation

- The GPU kernel

A Simple GPU implementation

- The GPU kernel

```
__global__ void searchGPU(char* index, int entries, int wordlength,
                           char* search_keys, int* results) {
    char* res;
    // use block and thread numbers for indexing
    res = bsearch(&search_keys[((blockIdx.x*BLOCK_SIZE)+threadIdx.x)*wordlength],
                  index,
                  entries,
                  wordlength);
    // use block and thread numbers for indexing
    results[(blockIdx.x*BLOCK_SIZE)+threadIdx.x] = (res-data)/
                                                MAX_WORD_LENGTH;
}
```

- There is no libc on the GPU =(
- Just stick __device__ in front of the libc code?
- “bsearch” is recursive, but there is no recursion on the GPU
- ➔ Write a iterative one ...

A Simple GPU binary search

```
__device__ char* bsearchGPU(char *key, char *base, int n, int size) {
    char *mid_point;
    int cmp;

    while (n > 0) {
        mid_point = (char *)base + size * (n >> 1);
        if ((cmp = strcasecmpGPU(key, mid_point)) == 0)
            return (char *)mid_point;
        if (cmp > 0) {
            base = (char *)mid_point + size;
            n = (n - 1) >> 1;
        } // cmp < 0
        else n >>= 1;
    }
    return (char *)NULL;
}
```

- Still need strcmp

A Simple GPU binary search

```
__device__ char* bsearchGPU(char *key, char *base, int n, int size) {
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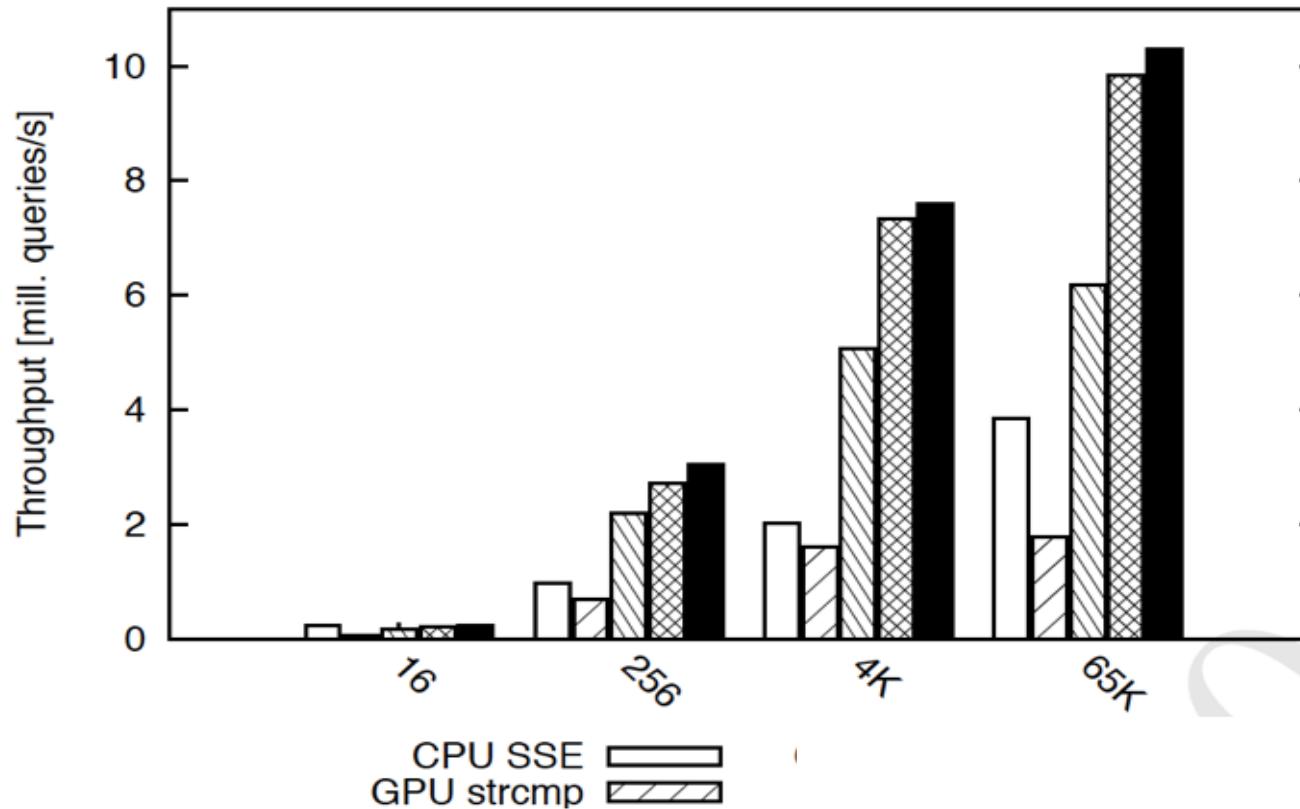
    while (n > 0) {
        mid_point = (char *)base + size * (n >> 1);
        if ((cmp = strcmpGPU(key, mid_point)) == 0)
            return (char *)mid_point;
        if (cmp > 0) {
            base = (char *)mid_point + size;
            n = (n - 1) >> 1;
        } // cmp < 0
        else n >>= 1;
    }
    return (char *)NULL;
}
```

- Still need strcmp
- Again, stick `__device__` in front of the libc code

```
__device__ int strcmpGPU(char* s1, char* s2) {
    while (*s1 == *s2++)
        if (*s1++ == 0) return 0;
    return (*s1 - *(s2 - 1));
}
```

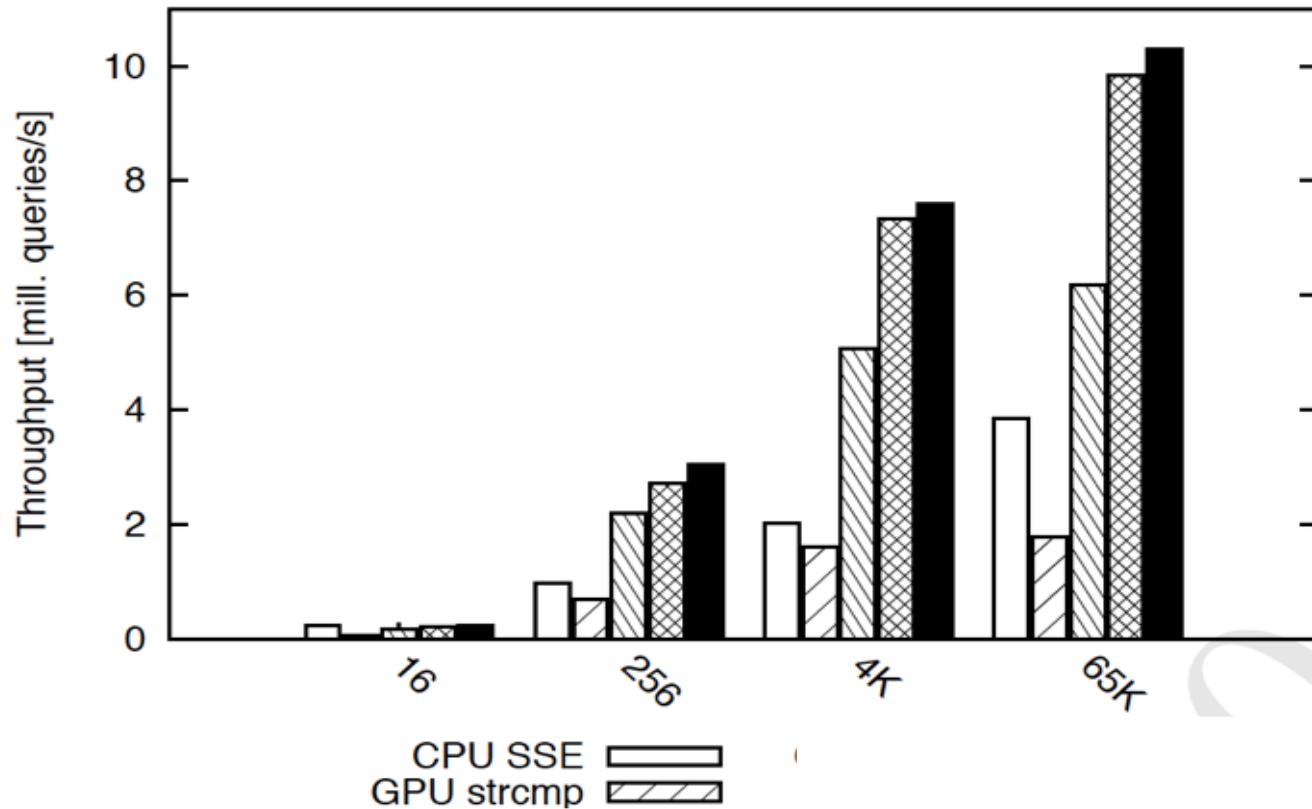
Binary Search on the GPU

- Searching a large data set (512MB) with 33 million (225) 16-character strings



Binary Search on the GPU – Why is it slow?

- Searching a large data set (512MB) with 33 million (225) 16-character strings

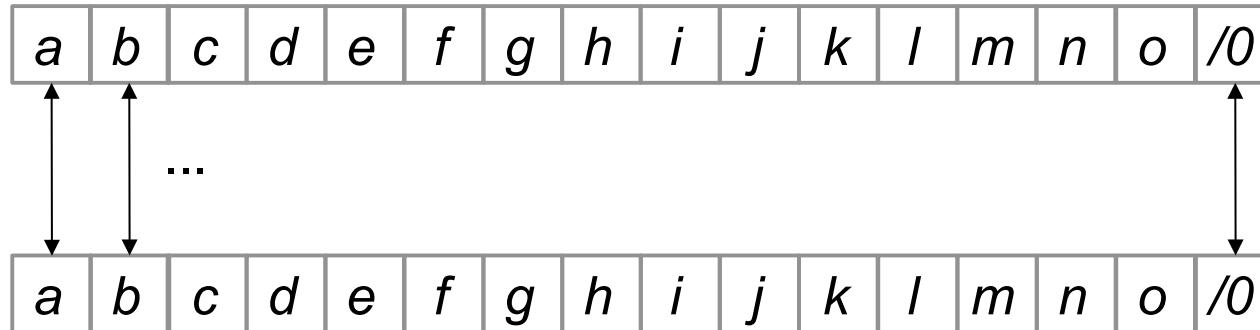


- It's slower than a CPU implementation for all data set sizes!
 - Let's try some optimizations ...

Search requires to compare

- Search naturally requires MANY comparisons
- The strcmp() library function:

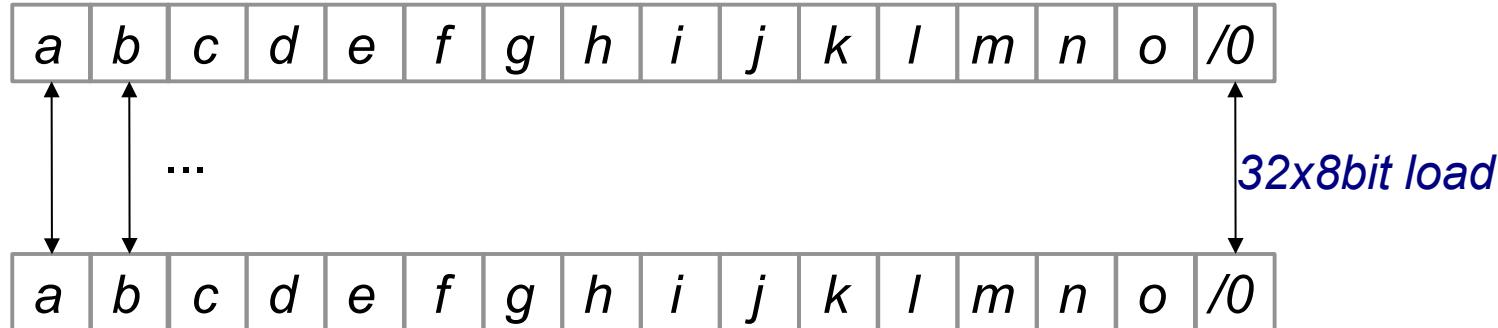
```
int strcmp(const char* s1, const char* s2) {  
    while (*s1 == *s2++)  
        if (*s1++ == 0) return 0;  
    return (*s1 - *(s2 - 1));  
}
```



Search requires to compare

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- The strcmp() library function:

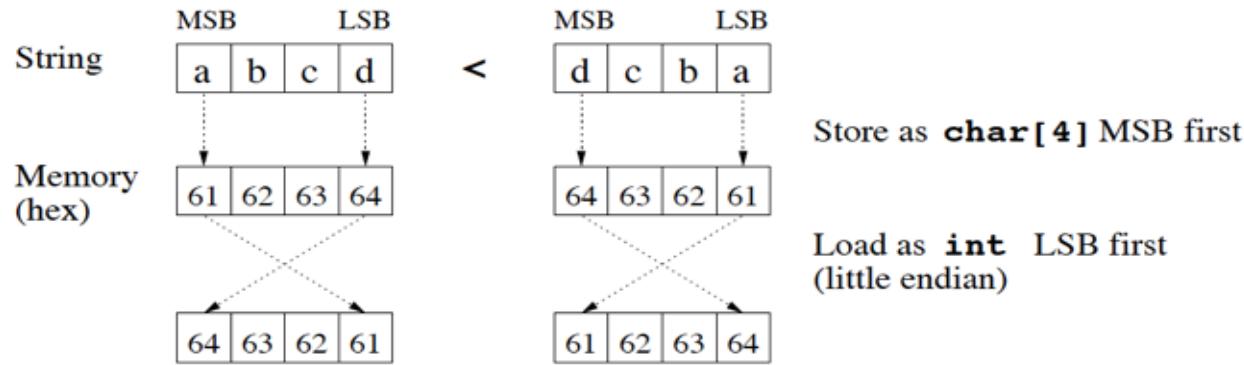
```
int strcmp(const char* s1, const char* s2) {  
    while (*s1 == *s2++)  
        if (*s1++ == 0) return 0;  
    return (*s1 - *(s2 - 1));  
}
```



- Byte-wise memory access is known to be slow

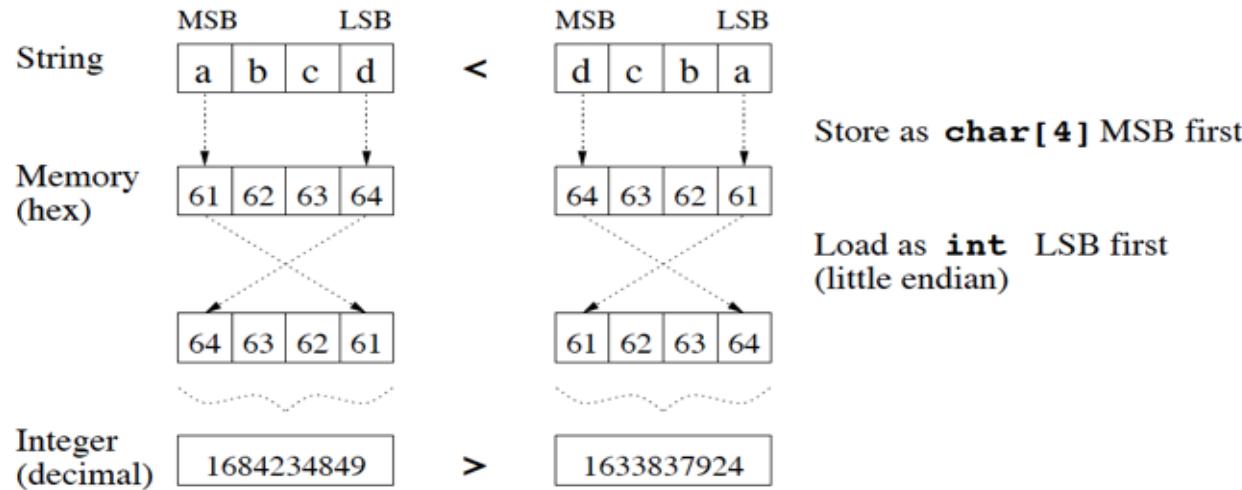
Optimizing compare operations

- How about vector string comparison, a la SSE?
- No Byte vectors on the GPU ... but Integer vectors



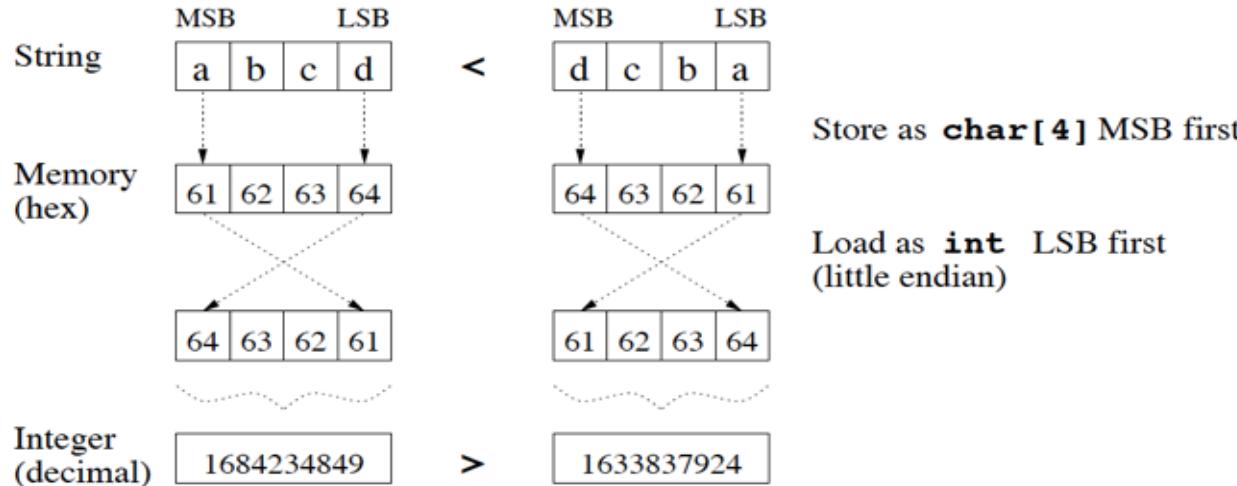
Optimizing compare operations

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Optimizing compare operations

- How about vector string comparison, a la SSE?
- No Byte vectors on the GPU ... but Integer vectors



- Loading character strings as int changes endianness
- CPU has bswap, on the GPU we have to write it:

```
#define BSWP( x ) ; \
temp = ( x ) << 24 ; \
temp = temp | ( ( ( x ) << 8) & 0x00FF0000 ) ; \
temp = temp | ( ( ( unsigned ) ( x ) >> 8) & 0x0000FF00 ) ; \
x = temp | ( ( unsigned ) ( x ) >> 24 ) ;
```

Optimizing compare operations

- Comparing integer vectors (bswap for <> skipped for clarity)

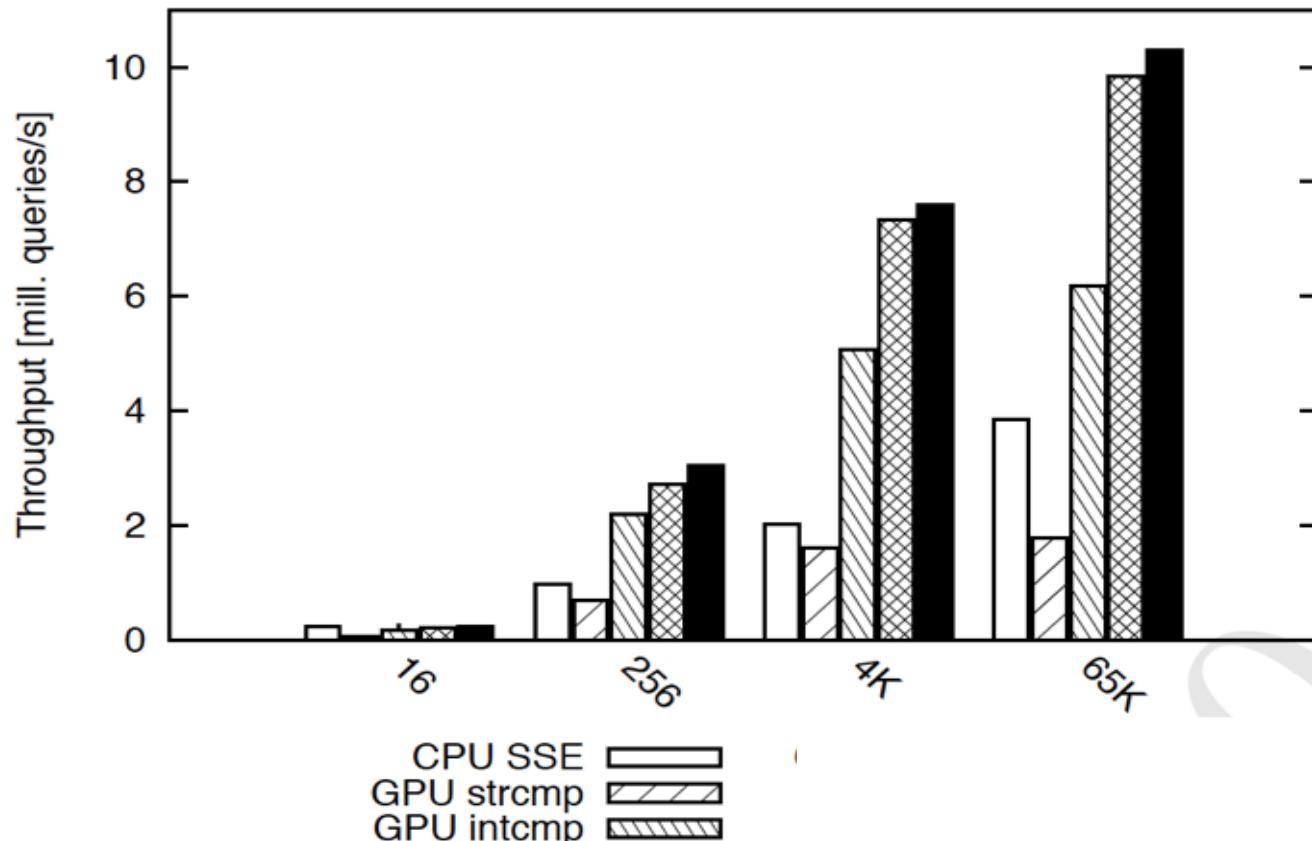
```
__device__ int intcmp(uint4* a, uint4* b) {

    int r =1;
    if ((*a).x < (*b).x)
        r=-1;
    else if ((*a).x == (*b).x) {
        if ((*a).y < (*b).y)
            r=-1;
        else if ((*a).y == (*b).y) {
            if ((*a).z < (*b).z)
                r=-1;
            else if ((*a).z == (*b).z) {
                if ((*a).w < (*b).w)
                    r=-1;
                else if ((*a).w == (*b).w)
                    r=0;
            }
        }
    }
    return r;
}
```

- Still dereferencing 16 memory pointers ...

Binary Search on the GPU – Why is it slow?

- Searching a large data set (512MB) with 33 million (225) 16-character strings



- With intcmp it's only marginally faster than a CPU implementation
- We still do pointer chasing, i.e. roundtrips to memory ...

Reducing global memory access

- Intcmp is memory latency sensitive

Processor	L1 [cyc]	L2 [cyc]	L3 [cyc]	mem [cyc]
Intel Core i7 2.6GHz	4	10	40	350
nVidia GT200b 1.5 GHz	4	n/a	n/a	500

- We can use shared memory like L1

x 16 for each comparison !!!

Reducing global memory access

- Intcmp is memory latency sensitive

Processor	L1 [cyc]	L2 [cyc]	L3 [cyc]	mem [cyc]
Intel Core i7 2.6GHz	4	10	40	350
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x 16 for each comparison !!!

- We can use shared memory like L1

```

__shared__ uint4 cache[NUM_THREADS*2];

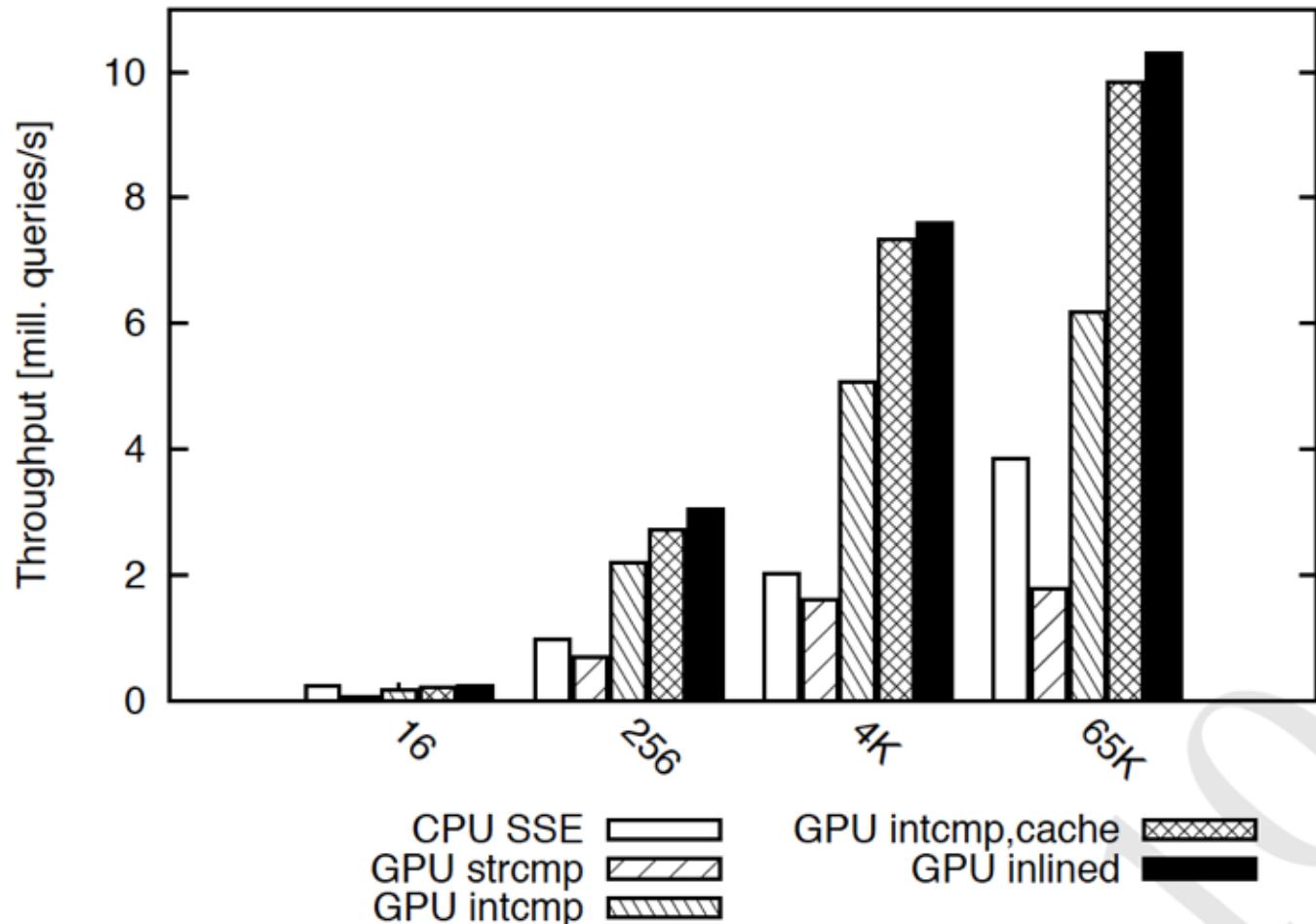
__device__ uint4* bsearchGPU( uint4 *key,  uint4 *base,
                             size_t nmemb,   size_t size)
{
    uint4 *mid_point;
    int cmp;
    cache[threadIdx.x*2] = *key;

    while (nmemb > 0) {
        mid_point = (uint4 *)base + size * (nmemb >> 1);
        cache[threadIdx.x*2+1] = *mid_point;
        if ((cmp = intcmp(&cache[threadIdx.x*2],
                          &cache[threadIdx.x*2+1])) == 0)
            return (uint4 *)mid_point;
    }
}

```

Binary Search on the GPU – optimized

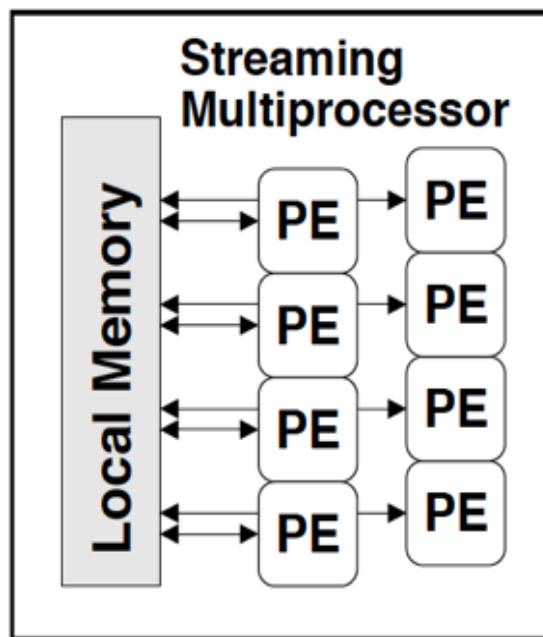
- Searching a large data set (512MB) with 33 million (225) 16-character strings



Is binary search optimal for a SIM[D/T] architecture ?

GPU architecture reminder – SIMD/SIMT

- Inside Streaming Multiprocessor
 - Single Instruction Multiple Threads/Data (SIMT/SIMD)
 - All PEs in 1SM execute same instruction or no-op (SIMD threads)
 - Warps of 32 threads (or more to hide memory latency)



What happens during Multi-threaded Binary Search ?

- Index: a sorted char array 32 entries
- 4 queries: **t , 8 , f , r**
- 4 processor cores: P1–P4
- 1 processor core – 1 search: P0:**t** , P1:**8** , P2:**f** , P3:**r**
- Theoretical worst-case execution time: $\log_2(32)=5$

4	5	6	7	8	9	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
---	---	---	---	----------	---	---	---	---	---	---	----------	---	---	---	---	---	---	---	---	---	---	---	----------	---	----------	---	---	---	---	---	---

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Iter. 1)
P0:t, P1:8, P2:f, P3:r

Iter. 2)

4	5	6	7	8	9	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

P1:8, P2:f P0:t, P3:r

What happens during Multi-threaded Binary Search ?

Iter. 2) 

The diagram illustrates string slicing operations on a character array. On the left, a horizontal line represents memory with boxes labeled from 4 to j. Red arrows point to index 8 ('8') and index f ('f'). Brackets below the line indicate ranges: P1:8 covers indices 4 to 8, and P2:f covers indices 4 to f. On the right, another horizontal line represents memory with boxes labeled from r to z. A red arrow points to index t ('t'), and a bracket below it indicates the range P0:t, covering indices r to t.

Iter. 4)  

Iter. 5) 
P1:8

Multi-threaded Binary Search - Analysis

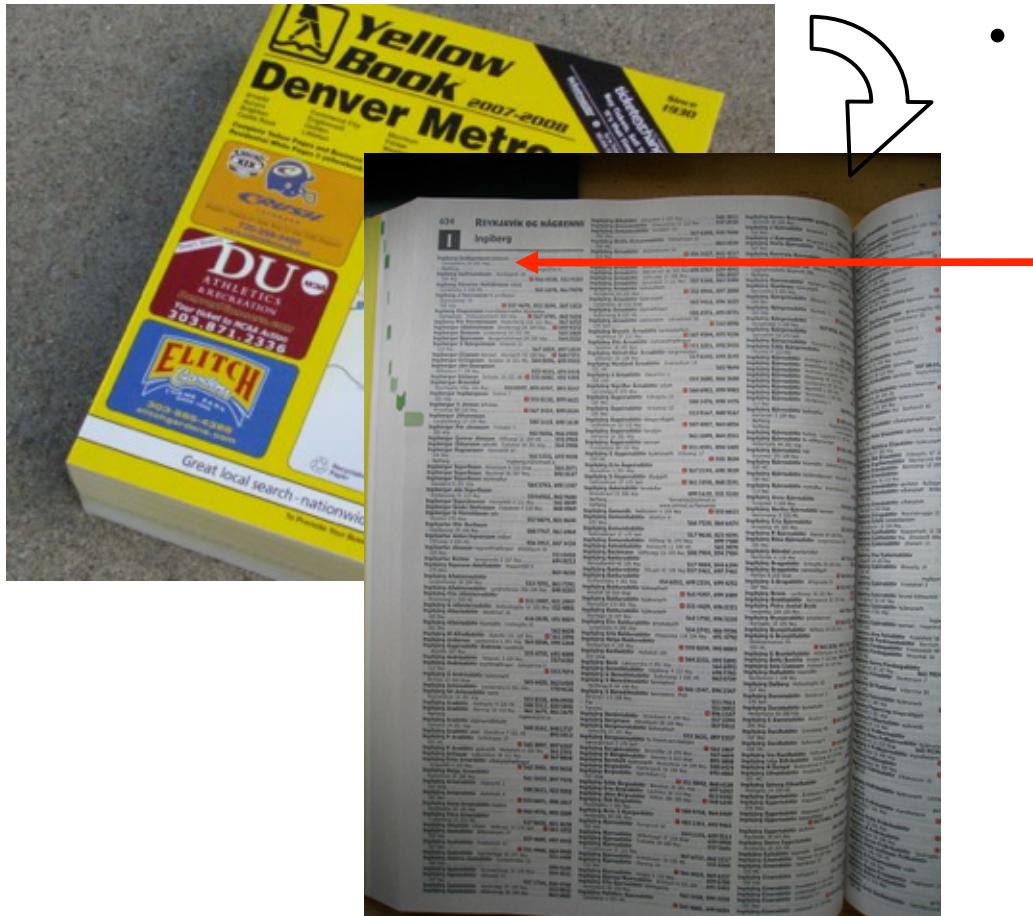
- 100% utilization requires $\# \text{cores concurrent queries}$
- Queries finishing early
→ utilization < 100%
- Memory access collisions
→ serialized memory access
- $\# \text{memory accesses } \log_2(n)$
- More threads
→ more results
→ response time likely to be worst case: $\log_2(n)$



Can we improve the worst case?

Binary Search

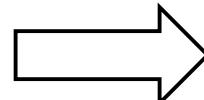
- How Do you (efficiently) search an index?



- Open phone book ~middle
- 1st name = whom you are looking for?
- < , > ?
- Iterate
 - Each iteration: $\# \text{entries}/2$ ($n/2$)
 - Total time:
→ $\log_2(n)$

Parallel (Binary) Search

- What if you have some friends (3) to help you ?

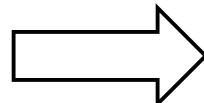


- Divide et impera !
 - Each is using binary search takes $\log_2(n/4)$
 - All can work in parallel → faster: $\log_2(n/4) < \log_2(n)$
- Give each of them $\frac{1}{4}$ *

* You probably want to tear it a little more intelligent than that, e.g. at the binding ;-)

Parallel (Binary) Search

- What if you have some friends (3) to help you ?

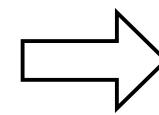
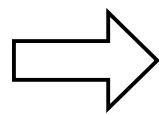


- Divide et impera !
 - Each is using binary search takes $\log_2(n/4)$
 - All can work in parallel → faster: $\log_2(n/4) < \log_2(n)$
 - 3 of you are **wasting time** !
- Give each of them $\frac{1}{4}$ *

* You probably want to tear it a little more intelligent than that, e.g. at the binding ;-)

P-ary Search

- Divide et impera !!



...

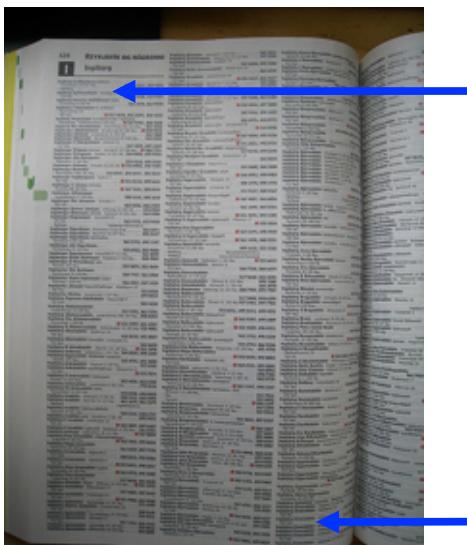
- How do we know who has the right piece ?

P-ary Search

- Divide et impera !!



- How do we know who has the right piece ?



- It's a sorted list:
 - Look at first and last entry of a subset
 - If **first entry** < searched name < **last entry**
 - Redistribute
 - Otherwise ... throw it away
 - Iterate

P-ary Search

- What do we get?



+

- Each iteration: $n/4$
→ $\log_4(n)$
- Assuming redistribution time is negligible:
 $\log_4(n) < \log_2(n/4) < \log_2(n)$
- But each does 2 lookups !
- How time consuming are lookup and redistribution ?

P-ary Search

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||

||

memory synchronization
access

P-ary Search

- What do we get?



+

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II

memory access

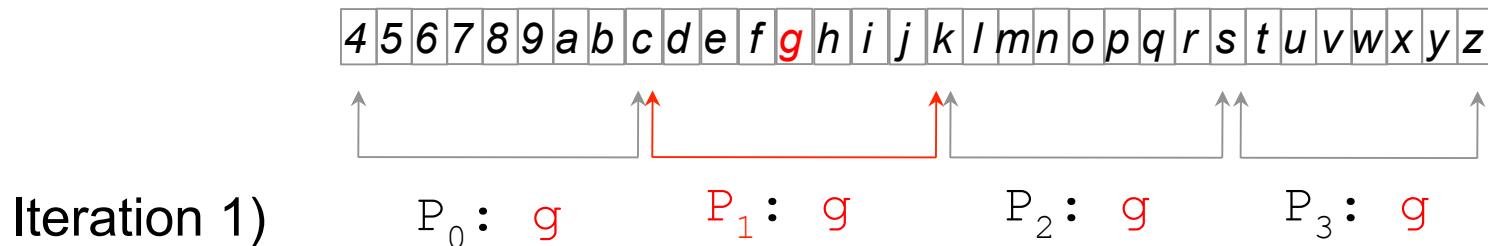
II

synchronization

- Searching a database index can be implemented the same way
 - Friends = Processor cores (threads)
 - Without destroying anything ;-)

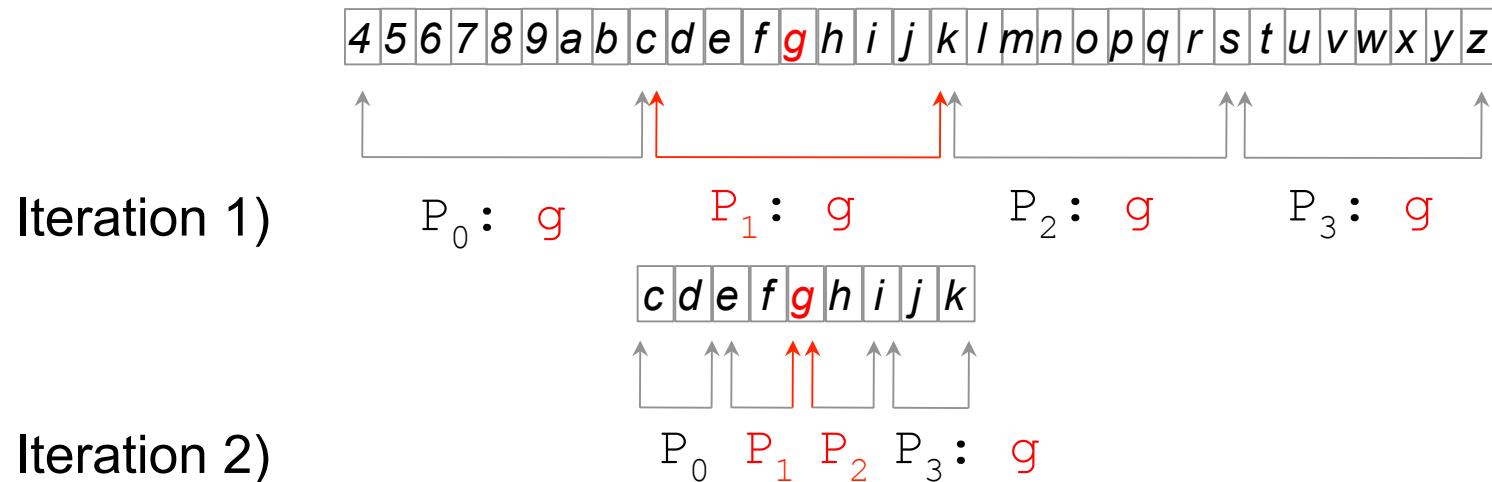
P-ary Search - Implementation

- Strongly relies on fast synchronization
 - friends = threads / vector elements



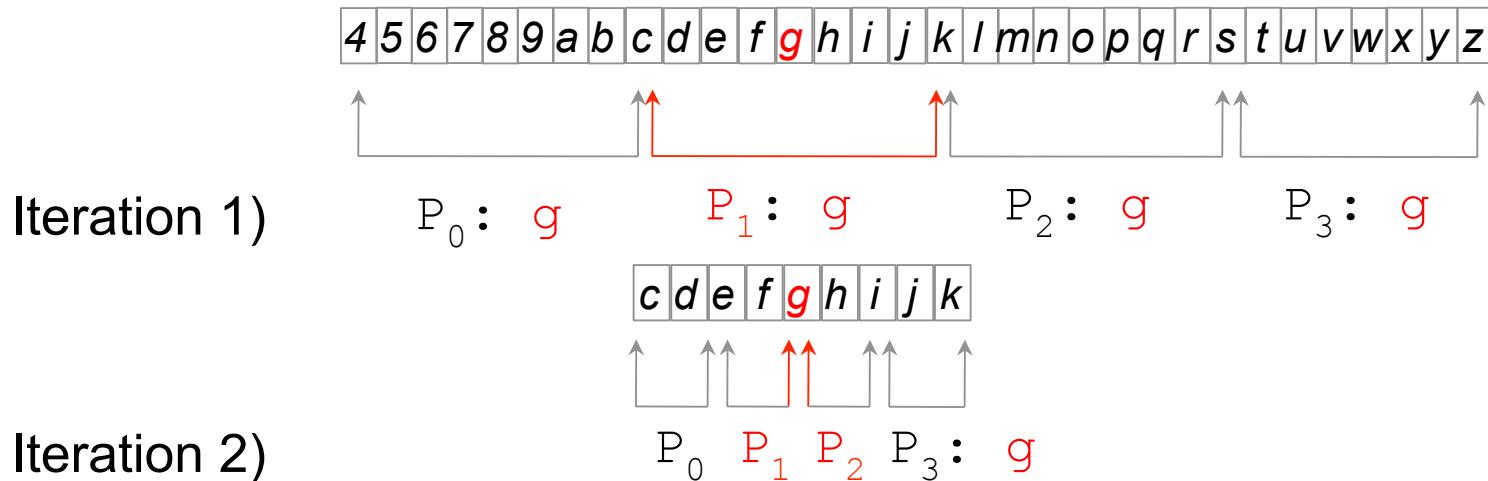
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P-ary Search - Implementation

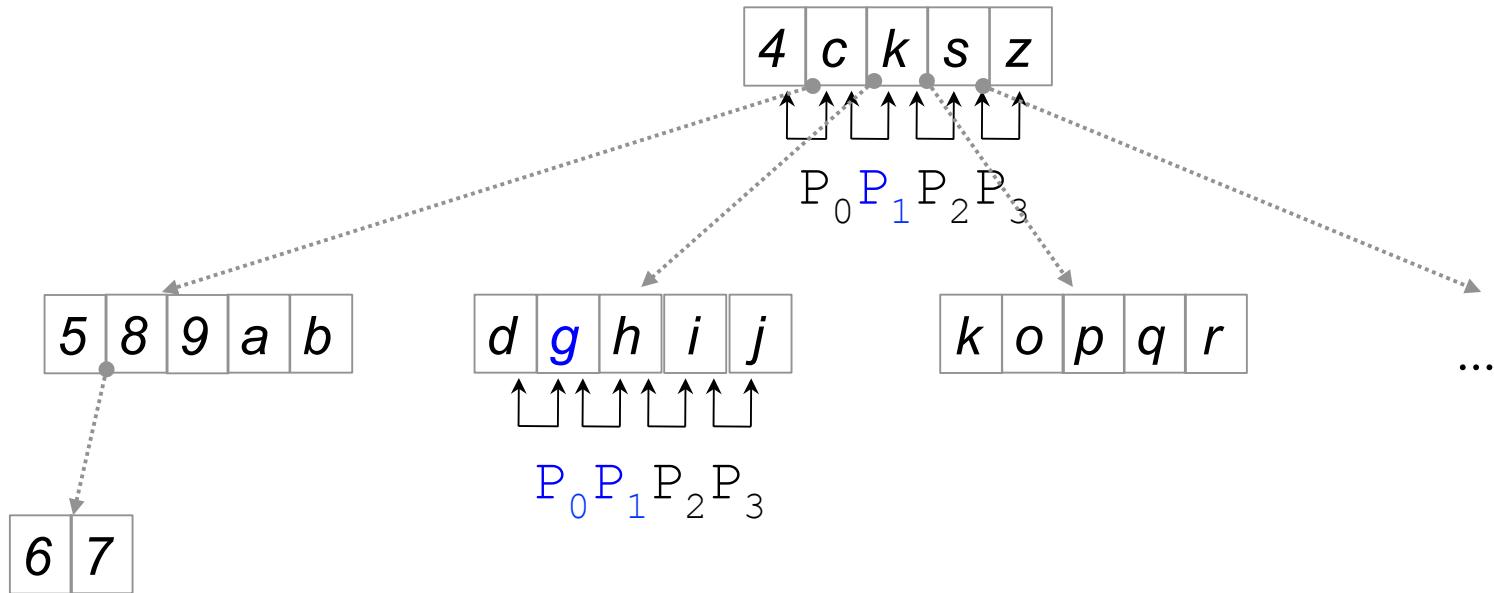
- Strongly relies on fast synchronization
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- Synchronization ~ repartition cost
- pthreads (\$\$), **cmpxchng(\$)**
- SIMD SSE-vector, GPU threads via shared memory (~0)
- Implementation using a B-tree is similar and (obviously) faster

P-ary Search - Implementation

- B-trees group pivot elements into nodes



- Access to pivot elements is coalesced instead of a gather
- Nodes can also be mapped to
 - Cache Lines (CSB+ trees)
 - Vectors (SSE)
 - #Threads per block

P-ary Search on a sorted integer list – Implementation (1)

```
__shared__ int offset;
__shared__ int cache[BLOCKSIZE+2]

__global__ void parySearchGPU(int* data, int length,
                           int* list_of_search_keys, int* results)

int start, sk;
int old_length = length;
// initialize search range starting with the whole data set
if (threadIdx.x == 0) {
    offset = 0;
    // cache search key and upper bound in shared memory
    cache[BLOCKSIZE] = 0x7FFFFFFF;
    cache[BLOCKSIZE+1] = list_of_search_keys[blockIdx.x];
    results[blockIdx.x] = -1;
}
__syncthreads();
//
sk = cache[BLOCKSIZE+1];
```

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}
syncthreads();
//
sk = cache[BLOCKSIZE+1];
```

Why?

P-ary Search on a sorted list – Implementation (2)

```
// repeat until the #keys in the search range < #threads
while (length > BLOCKSIZE) {
    // calculate search range for this thread
    length = length/BLOCKSIZE;
    if (length * BLOCKSIZE < old_length) length += 1;
    old_length = length;
    // why don't we just use floating point?
    start = offset + threadIdx.x * length;
    // cache the boundary keys
    cache[threadIdx.x] = data[start];
    __syncthreads();
    // if the searched key is within this thread's subset,
    // make it the one for the next iteration
    if (sk >= cache[threadIdx.x] && sk < cache[threadIdx.x+1]) {
        offset = start;
    }
    __syncthreads();
    // all threads start next iteration with the new subset
}
```

P-ary Search on a sorted list – Implementation (2)

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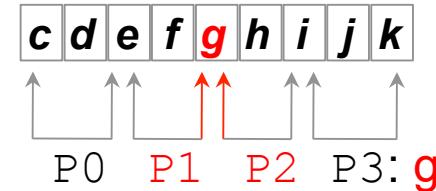
Why?

P-ary Search on a sorted list – Implementation (3)

```
// last iteration
start = offset + threadIdx.x;
if (sk == data[start])
    results[blockIdx.x] = start;
}
```

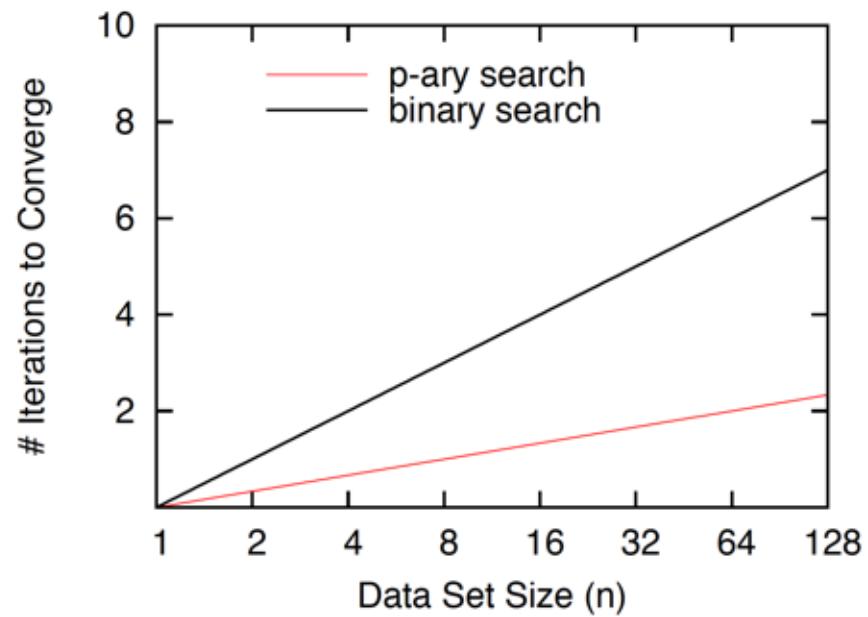
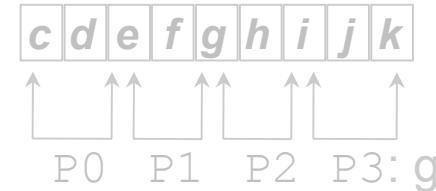
P-ary Search – Analysis

- 100% processor utilization for each query
- Multiple threads can find a result
 - How does this impact correctness?



P-ary Search – Analysis

- 100% processor utilization for each query
- Multiple threads can find a result
 - How does this impact correctness?
- Convergence depends on #threads
- GTX285: 1 SM, 8 cores(threads) → $p=8$
- Better Response time
 - $\log_p(n)$ vs $\log_2(n)$

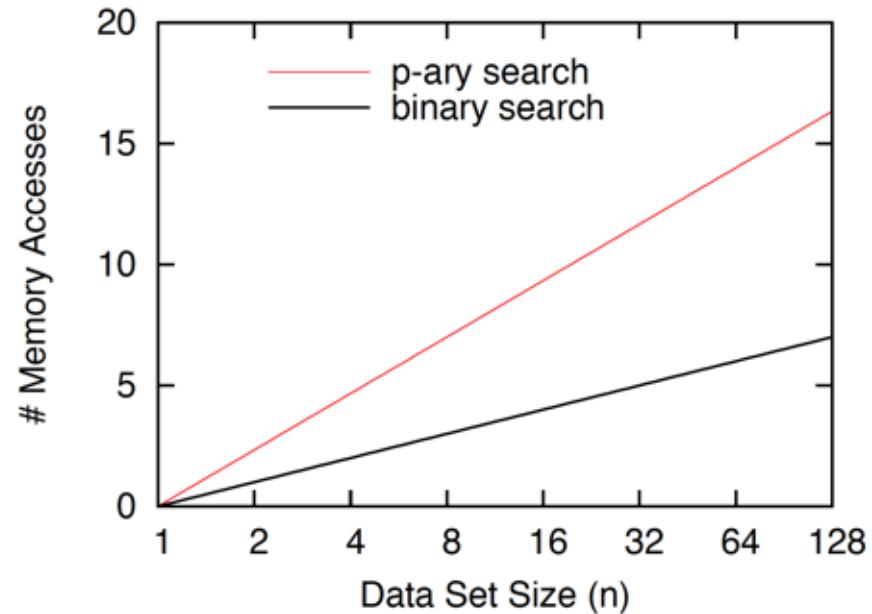
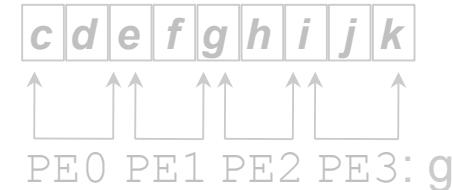


P-ary Search – Analysis

- 100% processor utilization for each query
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 - Does not change correctness
- Convergence depends on #threads

GTX285: 1 SM, 8 cores(threads) → p=8

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 - $\log_p(n)$ vs $\log_2(n)$
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 - $(p^2 \text{ per iteration}) * \log_p(n)$
 - Caching
 - $(p-1) * \log_p(n)$ vs. $\log_2(n)$

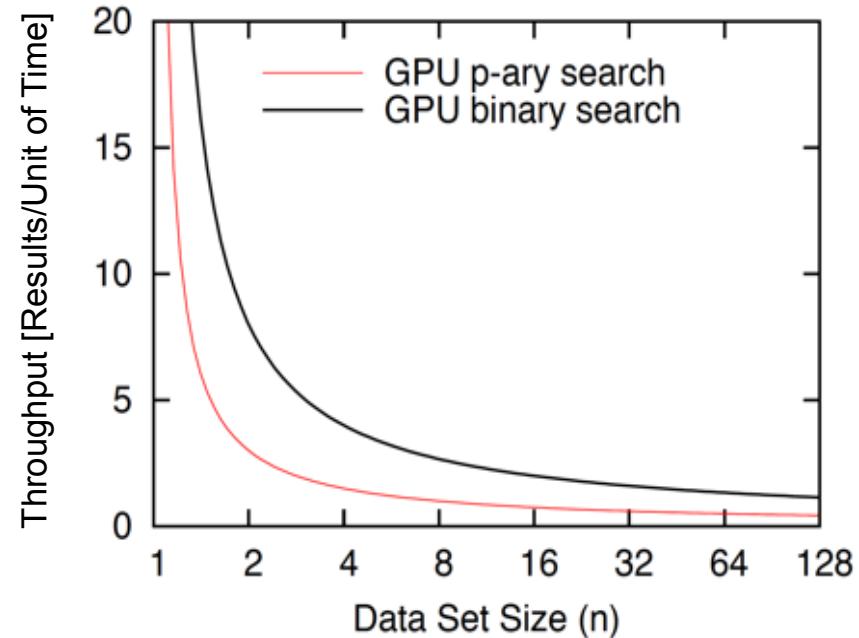
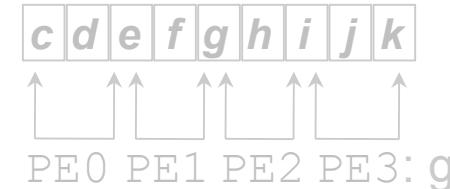


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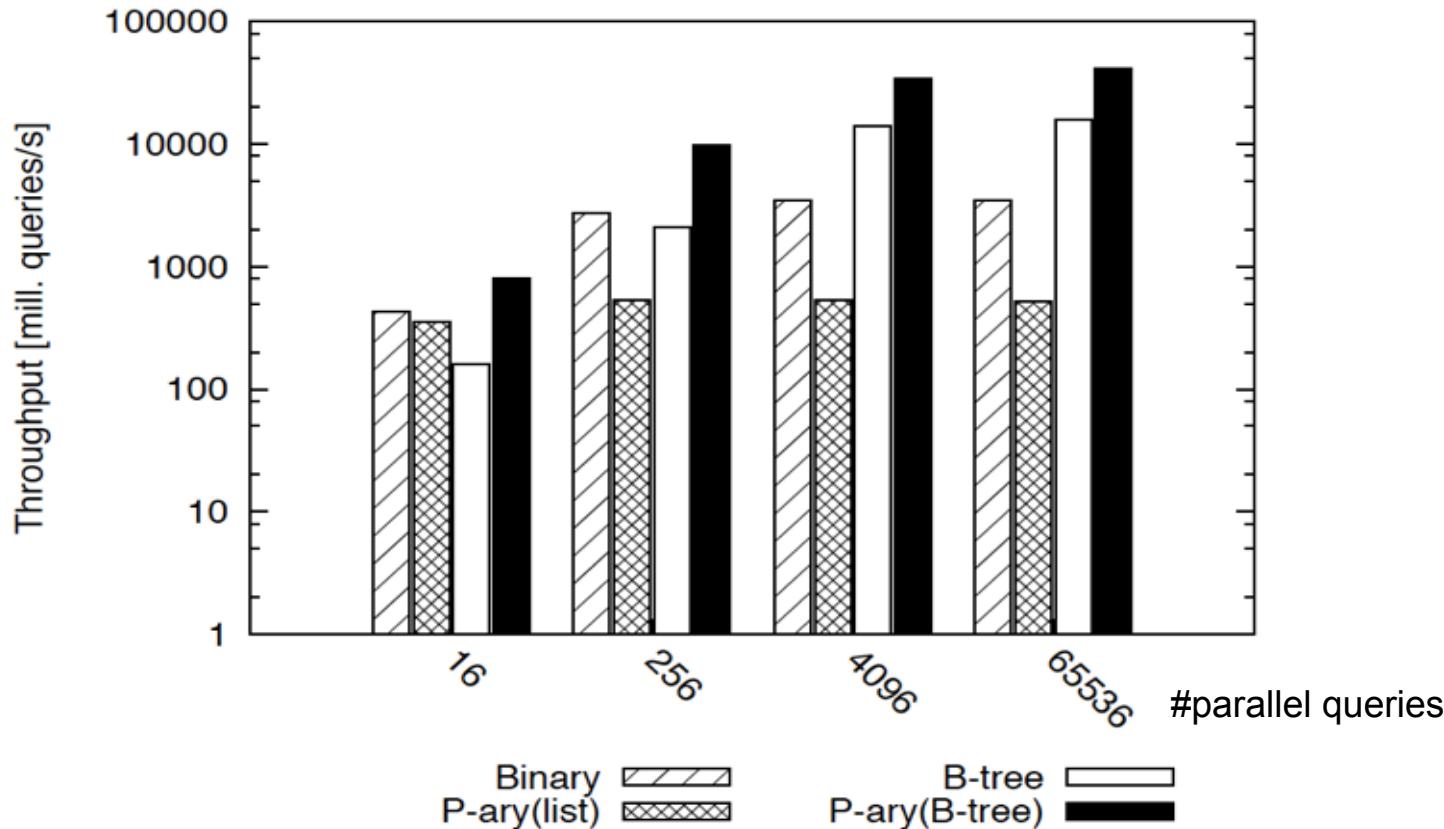
GTX285: 1 SM, 8 cores(threads) $\rightarrow p=8$

- Better Response time
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 - p^2 per iteration * $\log_p(n)$
 - Caching
 - $(p-1) \cdot \log_p(n)$ vs. $\log_2(n)$
- Lower Throughput
 - $1/\log_p(n)$ vs $p/\log_2(n)$



P-ary Search (GPU) – Throughput

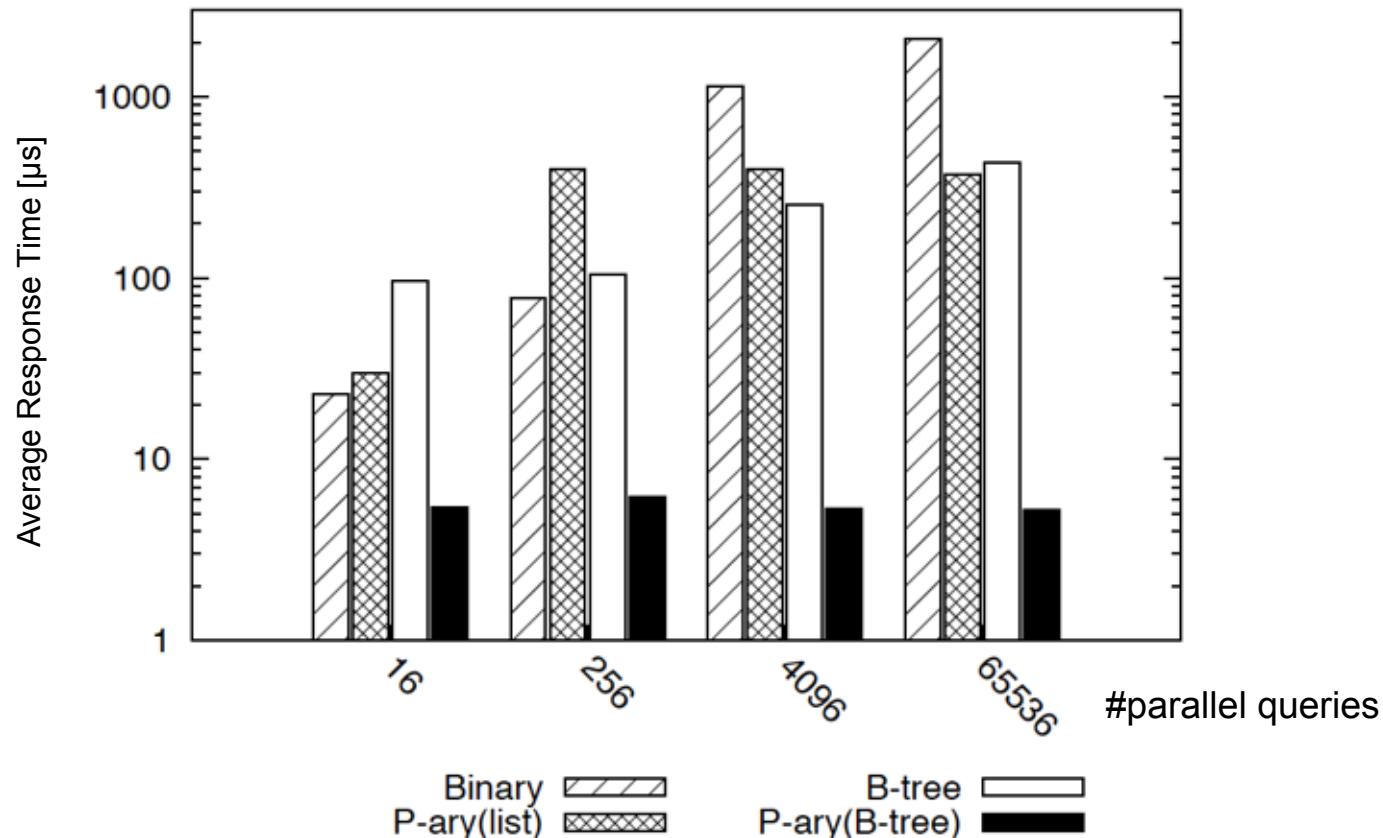
- Superior throughput compared to conventional algorithms



Searching a 512MB data set with 134mill. 4-byte integer entries,
Results for a nVidia GT200b, 1.5GHz, GDDR3 1.2GHz.

P-ary Search (GPU) – Response Time

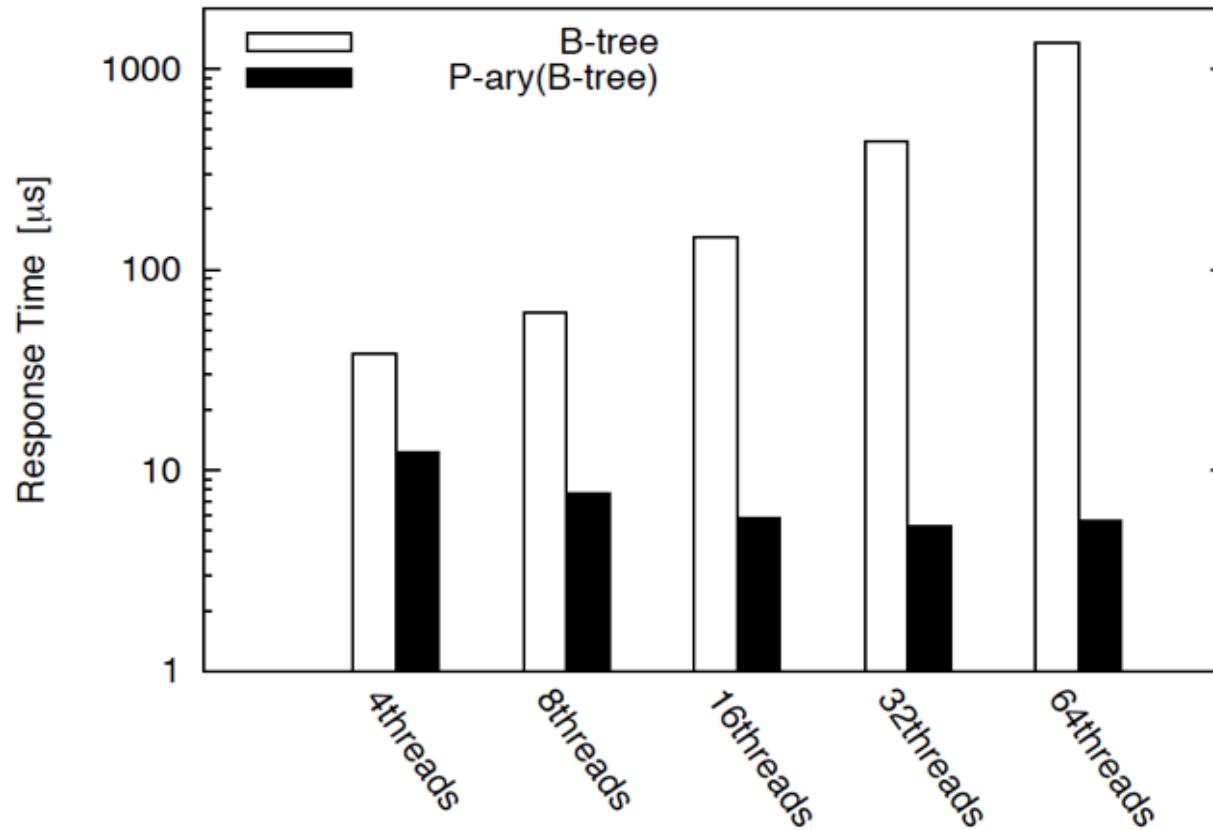
- Response time is workload independent for B-tree implementation



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P-ary Search (GPU) – Scalability

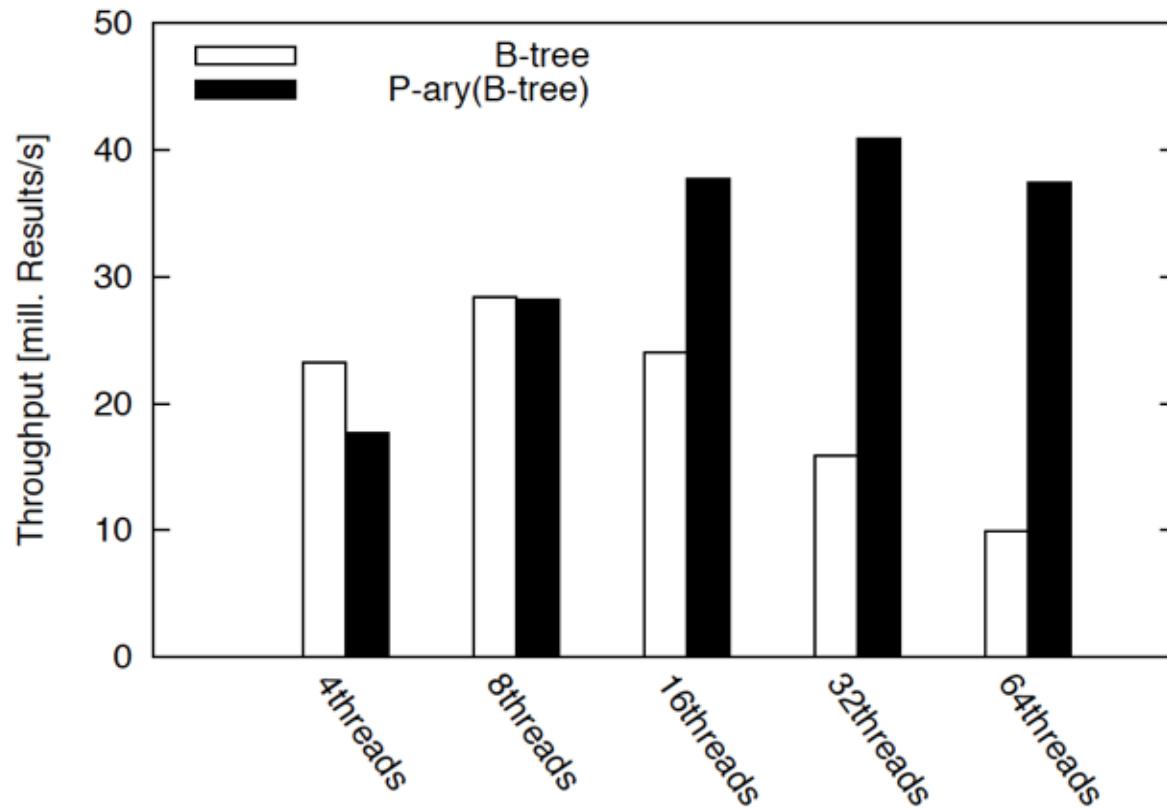
- GPU Implementation using SIMT (SIMD threads)
- Scalability with increasing #threads (P)



64K search queries against a 512MB data set with 134mill. 4-byte integer entries,
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Questions?