

Q1. $180.110.10.50 \rightarrow 200.18.30.22$

a. Destination address in case of indirect routing will be the mobile node's permanent address. It will be $180.110.10.50$.

Source address will be mobile node's permanent address : $180.110.10.50$

b. In case of direct routing, dest address of a data packet will be mobile node's COA as the correspondent will contact with home agent and gets in COA. — $200.18.30.22$.

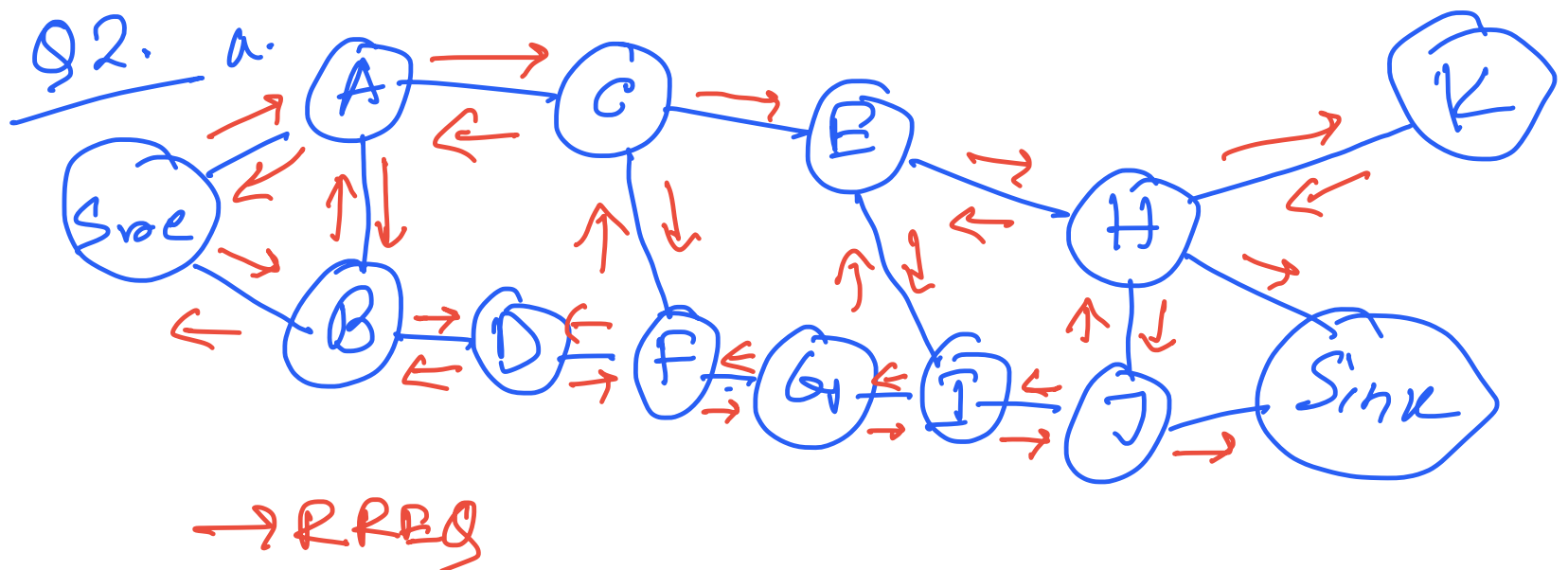
The Source address will also be $200.18.30.22$.

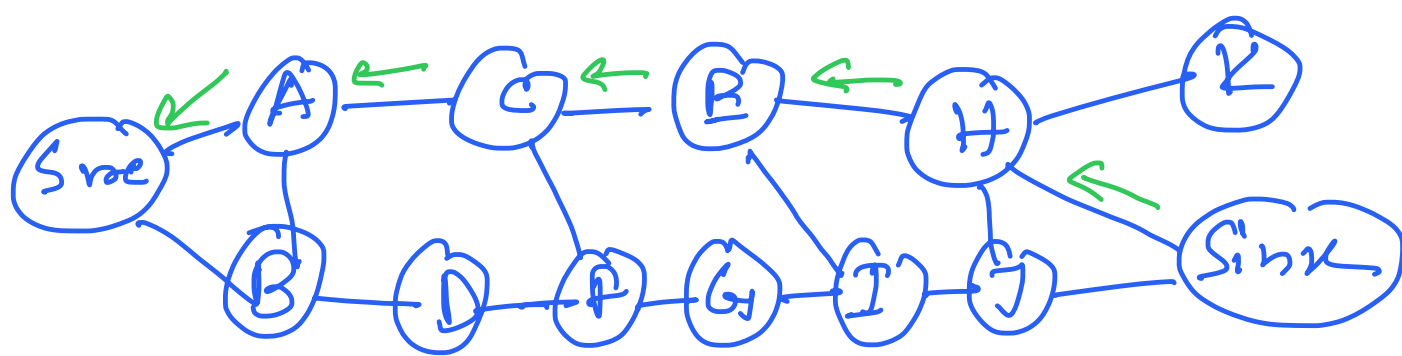
c. In this case as the mobile node is in correspondent node's network, direct routing is more efficient. In case of indirect, the packets will go to home agent then again to correspondent's network, will have to travel extra hops that is why direct is more efficient.

d. Firewall stops any packets whose source / dest address does not belong to correspondent's network.

Neither direct routing / Indirect routing will be impacted by this. As, for indirect, the packets will go to home network from that home agent sends to foreign agent with COA address only. Then mobile node sends directly to correspondent with PA but since this is intended not incoming / outgoing firewall won't cause a problem.

For direct routing, the packets from directly between conn. & mobile in the same network so no problem is caused.





← RREP

b. Routing table entry B, C, D, G, J, E
 Every node that receives RREP from Src.
 Sets up reverse path routing table entry
 pointing to Src.

Similarly, every node that receives RREP sets
 up routing table entry pointing to dest
 on the forward path

B

Dest	Next hop	Seq. No.
Src	Src	S ₁

C

Dest	Next hop	Seq No.
Src	A	S ₁
Sink	E	S ₂

D

Dest	Next hop	Seq No.
Src	B	S ₁

G

Dest	Next hop	Seq. No.
Src	F	S ₁

J

Dest	Next hop	Seq No.
Src	H	S ₁

E

Dest	Next hop	Seq. No.
Src	C	S ₁
Sink	H	S ₂

Assuming S_1 Seqno. generated by Src
 S_2 " " " Sink.

Note that a node who was not in RREP path, will not have entry for Sink.
 Similarly a node " receives RREQ from which over path earliest will put Routing table entry accordingly.

d. Best route

Src - A - C - E - H - Sink

Q3.

Src \rightarrow 2

A \rightarrow 2

C \rightarrow 3

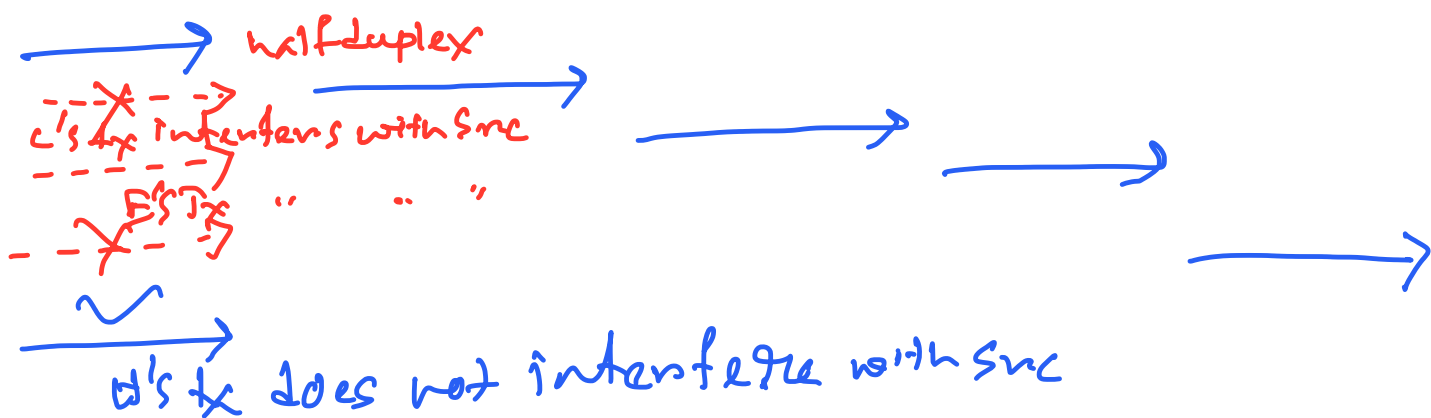
E \rightarrow 4

H \rightarrow 2

Sink \rightarrow 1

a.

Src A C E H Sink



J₁
 J₂
 J₃
 J₄
 J₅

Thus, effective data rate = $\frac{R}{4}$

b.

Time Slot₁

Src \rightarrow A (Ch₁)

C \rightarrow E (Ch₂)

H \rightarrow Sink (Ch₁)

Time Slot₂

A \rightarrow C (Ch₁)

E \rightarrow H (Ch₂)

Thus, effective data rate = $\frac{R}{2}$

4.

In the PIP protocol, the Tear down message is sent by the source only not by the sink.

When the Src determines that all packets have been acknowledged, sends a Tear down message. After forwarding Tear down nodes switch from SRC to vlc mode. If Src.

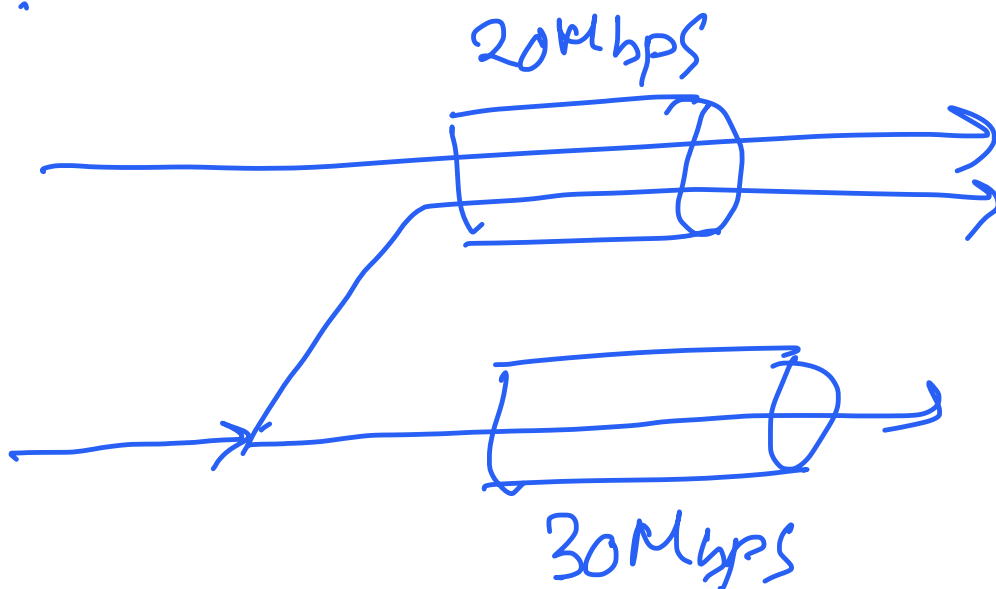
sends the Tear down message and even if the link between two nodes is 100% lossy even then there is no inconsistency between

Src & Sink, as Src sends Tear down only when it knows all packets have been received successfully.

5. We can't use data sequence numbers as it is because two issues

- ① Since the latency characteristics on two links would be different, there will be reordering of packets
- ② Each path will contain holes in the sequence numbers, this will upset some middle boxes & they won't allow such packets to flow.

6.



a. MPTCP + ECN TCP congestion control.
with ECN TCP each link will be shared evenly between subflows that use it

Thus, on the 20Mbps path
Regular TCP would get - 10Mbps

one subflow gets $\sim 10\text{Mbps}$

The other subflow is not shared with anyone, hence gets $\sim 30\text{Mbps}$

Thus, regular TCP gets $\sim 10\text{Mbps}$

MPTCP gets $\sim 10 + 30$
 $\sim 40\text{Mbps}$

b. For Coupled congestion control,
It is fair to regular TCP & utilizes paths efficiently to put more traffic on less congested paths.

Thus, it will utilize 30Mbps link fully & push more traffic there & very less traffic on the 20Mbps link.

Thus, regular TCP gets $\sim 20\text{Mbps}$
MPTCP gets $\sim 30\text{Mbps}$

Q7. b & c.

Q8.

a. SampleRate

Starts with 54 →
Samples 10th put out 48 as lossless Tx
Current avg

moves to 48 as avg with $48 < 54$

at 48 Samples at 36 as 36 lossless Tx
Current avg.

moves to 36 as avg with $36 < 48$

10th packet it Samples at 24, 48, 54

24 as lossless Tx
it does not move to 24
Current avg.
as its avg > 36 's avg.

So, SampleRate stays at 36 & keeps sampling
at 24, 48 & 54.

SampleRate →

54 →
Samples 10th put at 48 & lossless Tx
Current avg

moves to 48 as avg at 48 & diff-time $<$ threshold
at 48 → does not sample at 36 as diff $>$ threshold

So, SampleRate will stay at 48 & will keep
Sampling at 54

b. SampleRate

54 \rightarrow 48 lossless < avg.

moves to 48 as avg < 54's avg

at 48 samples at 54, 36

moves to 36 as avg < 48's avg

Does not sample at 24 as lossless > ^{crossed} avg

So, SampleRate will stay at 36 &

keeps sampling at 48, 54.

SampleRate

54

↓

48

↓

36

Stays at 36 & keeps sampling at 48, 54