

### Quiz4 (ADA-2023)

#### Answer keys

Q1. Given a graph  $G = (V, E)$  with positive edge weights, the Bellman-Ford algorithm and Dijkstra's algorithm can produce different shortest-path trees despite always producing the same shortest-path weights. True or False ?

Solution: **True**.

Q2. Dijkstra's Algorithm has lesser time complexity compared to Bellman Ford's algorithm, not considering the space complexity.

a) **True**

b) False

Q3. Given  $G = (N, M)$  where  $N$  is the number of vertices and  $M$  is the number of edges, it is known that  $G$  is a complete undirected graph. What is the total number of **spanning trees** for  $G$  ?

a.  $N * 2$

b.  $N^2$

**c.  $N^{(N-2)}$**

d.  $N^{(M-1)}$

Q4. Read the following statements -

I) Let  $G = (V, E)$  be a flow network, with a source  $s$ , a sink  $t$ , and a positive integer capacity  $c$  on every edge  $e$ . If  $f$  is a maximum  $(s, t)$ -flow in  $G$ , then for all  $e \in E$  that goes out of  $s$ ,  $f(e) = c$ .

II) Let  $G = (V, E)$  be a flow network with source  $s$ , sink  $t$  and all the edges of  $G$  have odd capacities. Then for any maximum  $(s, t)$ -flow  $f$  and for any edge  $e \in E$ ,  $f(e)$  is either zero or odd.

Which of the following statements is/are **incorrect**?

A. I only

B. II only

**C. Both I and II**

D. Neither I or II

Q5. Read the following statements -

I) Let  $G$  be a directed graph with positive edge weights  $w : E(G) \rightarrow \mathbb{R}$ .

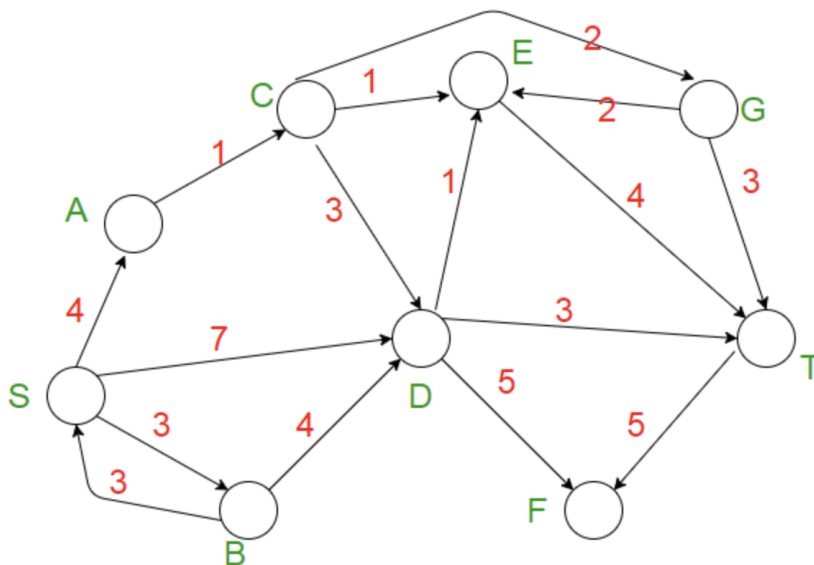
Suppose that we modify the graph  $G$  into  $G'$  as follows. For every edge  $e \in E(G)$ , we set  $w'(e) = w(e)/2$  to be the modified weights in  $G'$ . Then, every shortest path from  $s$  to  $t$  in  $G$  is a shortest path from  $s$  to  $t$  in  $G'$ .

II) Let  $G = (V, E)$  be a flow network, with a source  $s$ , a sink  $t$ , and a positive capacity  $c$  on every edge  $e$ . Suppose  $(A, B)$  is the unique  $s$ - $t$  minimum cut in  $G$  w.r.t these capacities. Consider a modified graph  $G'$  where the capacity of every edge is increased by 1. Then  $(A, B)$  is still a minimum  $(s, t)$ -cut for  $G'$ .

Which of the following statements is/are correct?

- A. I only
- B. II only
- C. Both I and II
- D. Neither I or II

Q6. Consider the directed graph shown in the figure below. There are multiple shortest paths between vertices  $S$  and  $T$ . Which one will be reported by Dijkstra's shortest path algorithm? Assume that, in any iteration, the shortest path to a vertex  $v$  is updated only when a strictly shorter path to  $v$  is discovered.



- A. SDT
- B. SBDT

C. SACDT

D. SACET

Q7. Consider the following statements given below:

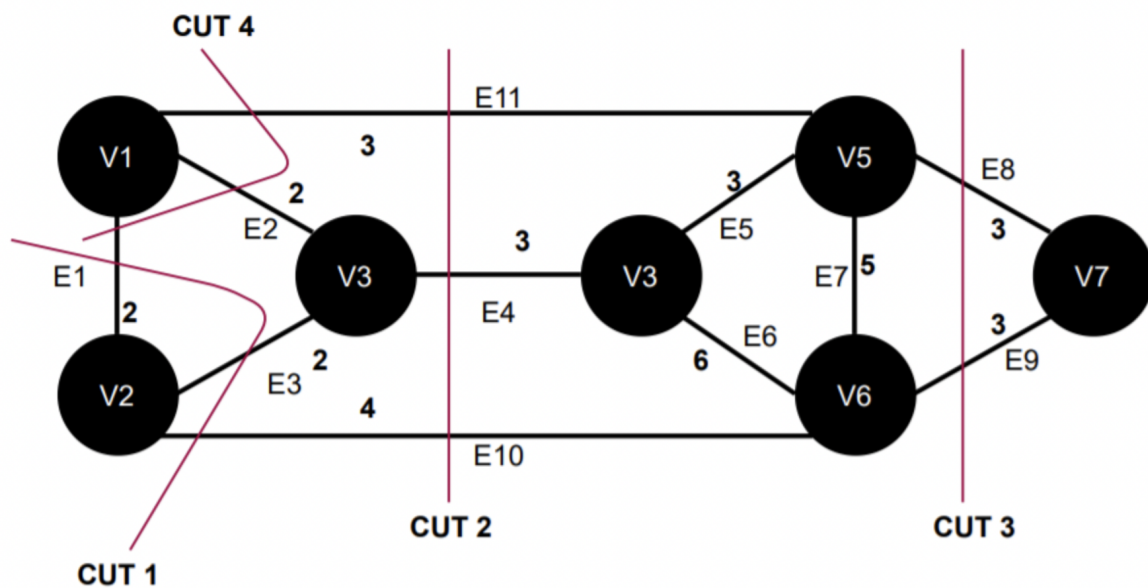
S1 : If a graph contains a negative weight cycle then Dijkstra's algorithm may or may not terminate.

S2 : The Bellman-Ford algorithm guarantees that it will always produce a shortest path between two given vertices  $u$  and  $v$  in any weighted graph.

Which of the above statements are **incorrect**?

1. Only S1
2. Only S2
3. Both S1 and S2
4. None of these

Q8. Which cut is the minimum cut of the following network?



A. CUT 3

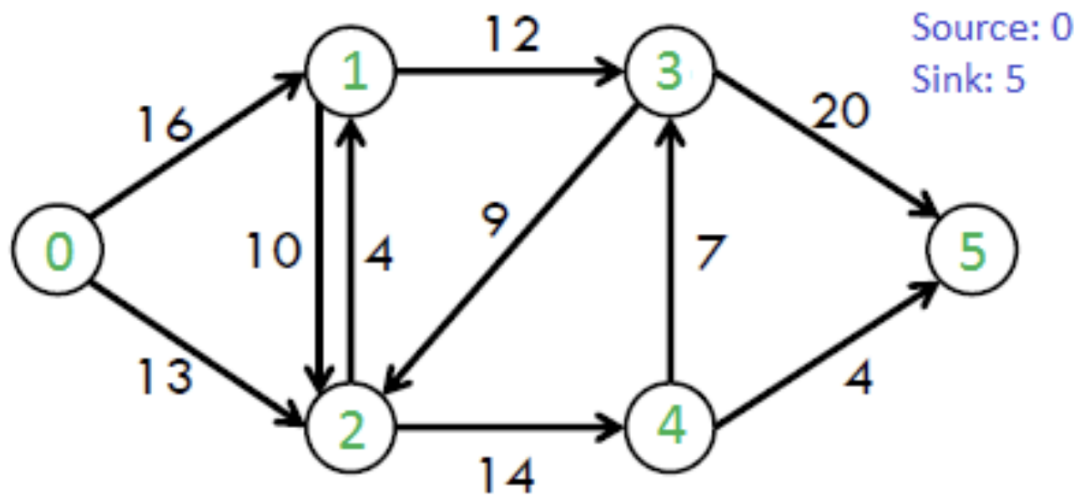
B. CUT 1

C. CUT 2

D. CUT 4

E. None of the above

Q9. The value of max-flow in the following graph is:



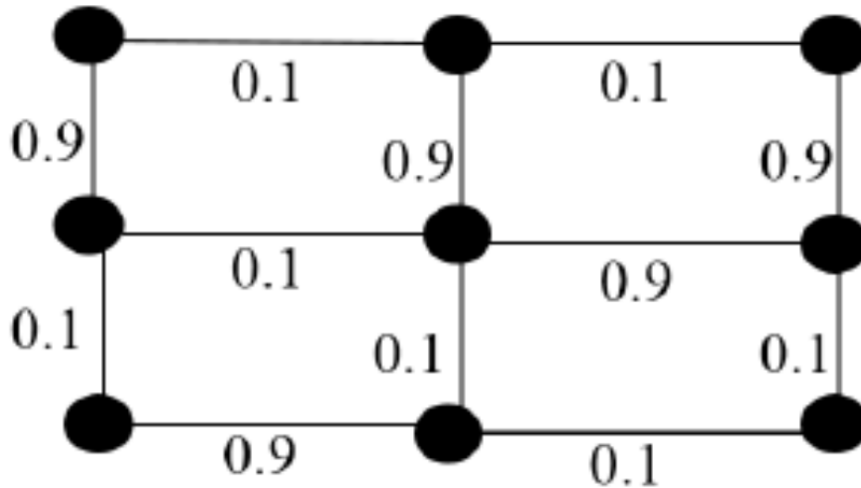
- A. 22
- B. 23**
- C. 24
- D. 25

Q10. Suppose that  $f$  is a feasible but not maximum  $s, t$ -flow in a network with nonnegative capacity. Read the statements below and choose the correct statement.

- S1. Residual graph  $G_f$  will never have a path from  $s$  to  $t$  that was not a path from  $s$  to  $t$  in the original graph  $G$ .
- S2. The residual graph  $G_f$  can have a path  $P^*$  from  $s$  to  $t$  with  $\min\{\text{capacity}(e) : e \in P^*\} > 0$  that is not there at all in the original graph.
- S3. The residual graph  $G_f$  will always have a path  $P^*$  from  $s$  to  $t$  with  $\min\{\text{capacity}(e) : e \in P^*\}$  that is not there at all in the original graph.
- S4. None of the above

- A. S1
- B. S2**
- C. S3
- D. S4

Q11. Consider the following undirected graph with edge weights as shown:



The total number of possible minimum spanning trees are-

- A. 2
- B. 3**
- C. 4
- D. 6

Q12. Read the following statements:

- I) In a flow network, the maximum amount of flow passing from the source to the sink is equal to the total weight of the edges in a minimum cut.
- II) Let  $f$  be a flow on a network  $(G, c)$  with net flow  $v$  and let  $C$  be a vertex cut  $(S, T)$  with capacity  $k$ . Then  $v > k$ .

Which of the above statements is/are correct:

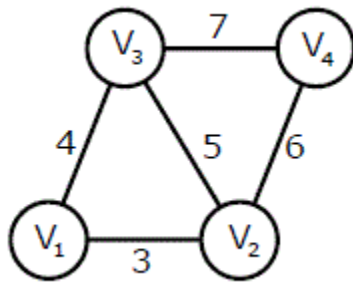
- A. Only I**
- B. Only II
- C. Both I and II
- D. Neither I nor II

Q13. Each maximum flow defines a minimum capacity cut, using the method discussed in class. Claim: This minimum capacity cut is unique, i.e., for a given maximum flow, there could not be more than one minimum capacity cut for this flow.

- A. True

B. False

Q14. An undirected graph  $G(V, E)$  contains  $n$  ( $n > 2$ ) nodes named  $v_1, v_2, \dots, v_n$ . Two nodes  $v_i, v_j$  are connected if and only if  $0 < |i - j| \leq 2$ . Each edge  $(v_i, v_j)$  is assigned a weight  $i + j$ . A sample graph with  $n = 4$  is shown below. What will be the cost of the minimum spanning tree (MST) of such a graph with  $n$  nodes?



a.  $(11n^2 - 5n) / 12$

b.  $n^2 - n + 1$

c.  $6n - 11$

d.  $2n + 1$

Q15. Let  $G$  be an undirected connected graph with distinct edge weight. Let  $e_{\max}$  be the edge with maximum weight and  $e_{\min}$  the edge with minimum weight. Which of the following statements is false?

A. Every minimum spanning tree of  $G$  must contain  $e_{\min}$

B. If  $e_{\max}$  is in a minimum spanning tree, then its removal must disconnect  $G$

C. No minimum spanning tree contains  $e_{\max}$

D.  $G$  has a unique minimum spanning tree