

Chapter 2

Logic of Quantified Statements

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- 2.1 Predicates and Quantified Statements
 - 2.2 Multiple Quantified Statements
 - 2.3 Proving Validity of Arguments with Quantified Statements

Introduction

- Logical analysis of compound statements: simple statements joined by the connectives \sim , \wedge , \vee , \rightarrow , \leftrightarrow , cannot be used to determine the validity in the majority.
- To determine the validity for the majority (such as “all” or “some”), it is necessary to separate the statements into parts.
- Predicate calculus: symbolic analysis of predicates and quantified statements.
- Statement calculus / propositional calculus: symbolic analysis of ordinary compound statements.

2.1 Predicates and Quantified Statements

- In logic, a predicate or open statement is a declarative sentence which
 - i. contains one or more variables,
 - ii. is not a statement, but
 - iii. becomes a statement when the variables in it are replaced by certain allowable choices.
- The domain of a predicate variable: set of all values that may be substituted in place of the variables.

- The predicates can be obtained by removing nouns from the statement.
- E.g. Let P be “is a student at TAR UC” and Q be “is a student at”. Then P and Q are predicate symbols.
- The sentences “ x is a student at TAR UC”, “ x is a student at y ” are symbolized as $P(x)$ and as $Q(x, y)$, respectively.
- x and y are predicate variables.

Example 1:

Consider

$p(x)$: The number $(x + 2)$ is an even integer.

$q(x, y)$: The numbers $y + 2$, $x - y$, and $x + 2y$ are even integers.

Domain for x is $\{4, 5\}$ and

domain for y is $\{2, 3\}$

Therefore,

- $p(4)$:
- $p(5)$:

- $q(4, 2)$:
- $q(4, 3)$:
- $q(5, 2)$:
- $q(5, 3)$:

Notes:

- The sets in which predicate variables take their values may be in words or in symbols (set is denoted by upper-case letter and element of set is denoted by lower-case letter).
- E.g. $x \in A$ indicates that “ x is an element of the set A ” or “ x is in A ”,
 $x \notin A$ indicates that “ x is not in A ”.

- E.g. $\{1, 2, 3\}$ refers to the set whose elements are 1, 2, and 3 while $\{1, 2, 3, \dots\}$ indicates the set of all positive integers.
- 2 sets are equal if, and only if, they have exactly same elements.

- Some special symbolic names for certain sets of number:

Symbol	Set
R	set of all real numbers
R^+	set of positive real numbers
Z	set of all integers
Z^{nonneg} / N	set of non-negative integers: 0, 1, 2, ... or natural numbers

Definition:

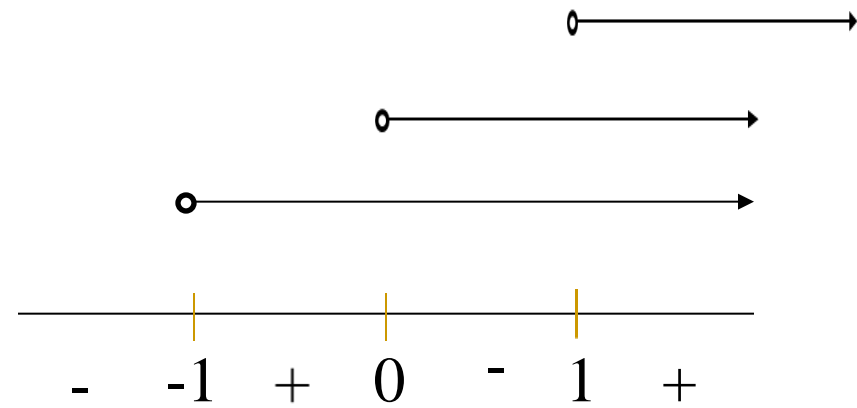
- If $P(x)$ is a predicate and x has domain D , the *truth set* of $P(x)$ is the set of all elements of D that make $P(x)$ true when substituted for x , denoted $\{x \in D \mid P(x)\}$, read as “the set of all x in D such that $P(x)$ ”.
- E.g. Let $P(x)$ be “ x is a factor of 10” and suppose the domain of x is all positive integers. Then the truth set $P(x)$ is $\{1, 2, 5, 10\}$

Example 2:

Find the truth set for the predicate below,

$$x > \frac{1}{x}, \text{ domain: } R$$

$$x - \frac{1}{x} > 0$$



Notation:

- Let $P(x)$ and $Q(x)$ be predicates and suppose the common domain of x is D .
- $P(x) \Rightarrow Q(x)$: every element in the truth set of $P(x)$ is in the truth set of $Q(x)$
- $P(x) \Leftrightarrow Q(x)$: $P(x)$ and $Q(x)$ have identical truth sets.

Example 3:

Let $P(x)$: x is the factor of 8.

$Q(x)$: x is the factor of 4.

$R(x)$: $x < 5$, $x \neq 3$.

Domain for x is \mathbb{Z}^+ .

Example 3:

$P(x)$: x is the factor of 8.

$Q(x)$: x is the factor of 4.

$R(x)$: $x < 5$, $x \neq 3$.

Quantifiers

- Added to predicates to obtain statements.
- Refers to quantities.
- The universal quantifier, \forall , denotes “for all”.
- The existential quantifier, \exists , denotes “there exists”.

Notes:

- $\forall x$ - for all x , for each x , for every x .
- $\exists x$ - for some x , for at least one x , there exists an x such that.

Universal

- Let $Q(x)$ be a predicate and D the domain of x .
- A Universal Statement , “ $\forall x \in D, Q(x)$ ”, is defined to be **true** if, and only if, $Q(x)$ is true for **every** x in D , and **false** if, and only if, $Q(x)$ is false for **at least one** x in D .
- A value for x for which $Q(x)$ is false is called a counterexample to the universal statement.

Existential

- Let $Q(x)$ be a predicate and D the domain of x .
- An Existential Statement, “ $\exists x \in D, Q(x)$ ”, is defined to be **true** if, and only if, $Q(x)$ is true for **at least one** x in D , and **false** if, and only if, $Q(x)$ is false for **every** x in D .

Example 4:

- i) Let $D = \{1, 2, 3, 4, 5\}$, and consider the statement “ $\forall x \in D, x \geq \frac{1}{x}$ ”. Show that this statement is true.

Let $P(x)$ be $x \geq \frac{1}{x}$.

Let $P(x)$ be $x \geq \frac{1}{x}$.

$P(1)$:

$P(2)$:

$P(3)$:

$P(4)$:

$P(5)$:

Since $P(x)$ is true for all $x \in D$, thus
 $\forall x \in D, x \geq \frac{1}{x}$ is a true statement.

ii) Consider the statement “ $\forall x \in R, x > \frac{1}{x}$ ”

Find counterexamples to show that this statement is false.

Let $P(x)$ be $x > \frac{1}{x}$.

Note:

- The technique used in (i) is called method of exhaustion – showing the truth of the predicate separately to each individual element of the domain.
- Only for finite predicate variable.

Example 5:

- i) Consider the statement $\exists x \in \mathbb{Z} \ni x^2 = x$ (\ni means such that). Show that this statement is true.

Let $P(x)$ be $x^2 = x$.

ii) Let $D = \{5, 6, 7, 8, 9, 10\}$ and consider the statement $\exists x \in D \ni x^2 = x$. Show that this statement is false.

Let $P(x)$ be $x^2 = x$.

$P(5)$:

$P(6)$:

$P(7)$:

$P(8)$:

$P(9)$:

$P(10)$:

Note:

- It is important to be able to translate either from formal into informal language.

Example 6:

Rewrite the following formal statements in a variety of equivalent but more informal ways. Do not use the symbol \forall and \exists .

i) $\forall x \in R, x^2 \neq -1$

ii) $\exists x \in Z \ni x^2 = x$

Example 7:

Rewrite each of the following statements formally. Use quantifiers and variables.

- i) Every real number is positive, negative, or zero.
- ii) Some real numbers are rational.

Universal Condition Statements

- $\forall x$, if $P(x)$ then $Q(x)$.

Example 8:

$\forall x \in R$, if $x > 4$, then $x^2 > 16$.

Informal way,

Example 9:

The square of any even integers is even.

Formal way,

Equivalent Forms of Universal Statements

- “ $\forall x \in U$, if $P(x)$ then $Q(x)$ ” can be written in the form “ $\forall x \in D$, $Q(x)$ ” by narrowing U to be the domain D consisting of all values of the variable x that make $P(x)$ true.

Note:

- Conversely, $\forall x \in D$, $Q(x)$ can be rewritten as $\forall x$, if x is in D then $Q(x)$.

Example 10:

“ \forall polygons p , if p is a square, then p is a rectangle.” is equivalent to

Equivalent Forms of Existential Statements

- “ $\exists x \in U$ such that $P(x)$ and $Q(x)$ ” can be written as “ $\exists x \in D$ such that $Q(x)$ ”, provided D is taken to consist of all elements in U that make $P(x)$ true.

Example 11:

“ \exists a number n such that n is prime and n is even.” is equivalent to

Negations of Quantified Statements

- The negation of the statement $\forall x \text{ in } D, Q(x)$ is logically equivalent to $\exists x \text{ in } D \text{ such that } \sim Q(x)$.
- Symbolically,
$$\sim(\forall x \in D, Q(x)) \equiv \exists x \in D \text{ such that } \sim Q(x).$$
- In word, the negation of a universal statement (“all are”) is logically equivalent to an existential statement (“some are not”).

- The negation of the statement $\exists x$ in D such that $Q(x)$ is logically equivalent to $\forall x$ in D , $\sim Q(x)$.
- Symbolically,
$$\sim(\exists x \in D \text{ such that } Q(x)) \equiv \forall x \in D, \sim Q(x).$$
- In word, the negation of an existential statement (“some are”) is logically equivalent to a universal statement (“all are not”).

Example 12:

Write negations for the following statements.

i) \forall irrational numbers x , x is not an integer.

ii) $\exists x \in R$ such that x is rational.

iii) No politicians are honest.

iv) All dinosaurs are extinct.

v) Some exercises have answers.

In summary,

Statement	When is it true?	When is it false?
$\exists x p(x)$	For some (at least one) a in the universe, $p(a)$ is true.	For every a in the universe, $p(a)$ is false.
$\forall x p(x)$	For every replacement a from the universe, $p(a)$ is true.	There is at least one replacement a from the universe for which $p(a)$ is false.

Negations of Universal Conditional Statements

- By the definition of the negation of a “for all” statement,

- $\sim(\forall x, P(x) \rightarrow Q(x))$

$$\equiv \exists x \text{ such that } \sim(P(x) \rightarrow Q(x)) \quad (1)$$

but, the negation of an if-then statement is logically equivalent to an “and” statement,

$$\sim(P(x) \rightarrow Q(x)) \equiv P(x) \wedge \sim Q(x) \quad (2)$$

substitute (2) into (1),

$$\sim(\forall x, P(x) \rightarrow Q(x)) \equiv \exists x \text{ such that } P(x) \wedge \sim Q(x).$$

- Symbolically,
 $\sim(\forall x, \text{ if } P(x) \text{ then } Q(x))$
 $\equiv \exists x \text{ such that } P(x) \text{ and } \sim Q(x).$

Example 13:

Write the negations for the following statements.

i) \forall animals x , if x is a cat then x has whiskers and x has claws.

Example 13:

ii) $\forall x \in R$, if $x > 3$ then $x^2 > 9$.

Example 14:

Negate the following statements and determine their truth values.

i) $r(x): 2x + 1 = 5$ $s(x): x^2 = 9$

$$\exists x [r(x) \rightarrow s(x)].$$

Example 14:

- ii) $p(x)$: x is odd. $q(x)$: $x^2 - 1$ is even.
 $\forall x \in \mathbb{Z}$, if x is odd, then $(x^2 - 1)$ is even.

2.2 Multiple Quantified Statements

- Contain more than 1 quantifier.

Example 15:

Rewrite each of the following without using variables or the symbols \forall or \exists .

- i) \forall colours C , \exists an animal A such that A is coloured C .
- ii) \exists a book b such that \forall people p , p has read b .

Example 16:

Rewrite the following formally using quantifiers and variables.

i) Everybody trusts somebody.

ii) Somebody trusts everybody.

Negations of Multiply Quantified Statements

- The negation of $\forall x, \exists y \ni P(x,y)$ is logically equivalent to $\exists x \ni \forall y, \sim P(x,y)$.
- The negation of $\exists x \ni \forall y, P(x, y)$ is logically equivalent to $\forall x, \exists y \ni \sim P(x,y)$.

Example 17:

Negate the statements below.

i) \exists a book b such that \forall people p , p has read b .

ii) \forall even integers n , \exists an integer k such that $n = 2k$.

iii) \exists a person x such that \forall people y , x loves y .

iv) $\forall x, \exists y [(P(x, y) \wedge Q(x, y)) \rightarrow R(x, y)]$

The Relation Among \forall , \exists , \wedge , and \vee

If $Q(x)$ is a predicate and the domain D of x is the set $\{x_1, x_2, \dots, x_n\}$, then

- “ $\forall x \in D, Q(x)$ ” is logically equivalent to “ $Q(x_1) \wedge Q(x_2) \wedge \dots \wedge Q(x_n)$ ”
- E.g. Let $Q(x)$ be “ $x \cdot x = x$ ” and $D = \{0, 1\}$.

Then $\forall x \in D, Q(x)$ can be rewritten as \forall binary digits $x, x \cdot x = x$.

This is equivalent to $0 \cdot 0 = 0$ and $1 \cdot 1 = 1$, symbolically, $Q(0) \wedge Q(1)$.

- “ $\exists x \in D$ such that $Q(x)$ ” is logically equivalent to “ $Q(x_1) \vee Q(x_2) \vee \dots \vee Q(x_n)$ ”
- E.g. Let $Q(x)$ be “ $x + x = x$ ” and $D = \{0, 1\}$.
Then $\exists x \in D$ such that $Q(x)$ can be rewritten as \exists a binary digit such that $x + x = x$.
This is equivalent to $0 + 0 = 0$ or $1 + 1 = 1$,
symbolically, $Q(0) \vee Q(1)$.

Example 18:

For the universe of natural numbers N , the assertion $\forall x [P(x) \vee Q]$ is equivalent to the infinite conjunction:

which can be rearranged using the distributive laws to form:

which is equivalent to $\forall x P(x) \vee Q$.

Notes:

- The variable x in $P(x)$ for example 18 is bound by quantifiers while the variable in Q is free.
- The universal quantifier, \forall distributes over the logical connective \wedge but not the existential quantifier, \exists .
- The existential quantifier, \exists distributes over the logical connective \vee but the universal quantifier, \forall does not.

Example 19:

For the universe of natural numbers N , the proposition $\forall x [P(x) \wedge Q(x)]$ can be expanded into an infinite conjunction:

which can be rearranged using associative and commutative laws to obtain:

which is equivalent to $\forall x P(x) \wedge \forall x Q(x)$.

Example 20:

Let the universe be the integers,

$P(x)$: x is an even integer.

$Q(x)$: x is an odd integer.

Then $\exists x P(x) \wedge \exists x Q(x)$ is true but

$\exists x [P(x) \wedge Q(x)]$ is false.

Therefore $\exists x P(x) \wedge \exists x Q(x)$ and

$\exists x [P(x) \wedge Q(x)]$ are not equivalent.

However, $\exists x [P(x) \wedge Q(x)]$ implies

$\exists x P(x) \wedge \exists x Q(x)$ is valid.

Variants of Universal and Conditional Statements

- Consider the statement of the form

$\forall x \in D$, if $P(x)$ then $Q(x)$,

1. Its *contrapositive* is the statement

$\forall x \in D$, if $\sim Q(x)$ then $\sim P(x)$.

2. Its *converse* is the statement

$\forall x \in D$, if $Q(x)$ then $P(x)$.

3. Its *inverse* is the statement

$\forall x \in D$, if $\sim P(x)$ then $\sim Q(x)$,

Example 21:

Write the contrapositive, converse and inverse for the following statements.

i) $\forall x \in R, \text{ if } x > 3, \text{ then } x^2 > 9.$

ii) \forall animals A , if A is a cat then A has whiskers and A has claws.

Notes:

- A universal condition statement is logically equivalent to its contrapositive.

$$\begin{aligned}\forall x \in D, \text{ if } P(x) \text{ then } Q(x) \\ \equiv \forall x \in D, \text{ if } \sim Q(x) \text{ then } \sim P(x)\end{aligned}$$

- A universal condition statement is *NOT* logically equivalent to its converse.

$$\begin{aligned}\forall x \in D, \text{ if } P(x) \text{ then } Q(x) \\ \neq \forall x \in D, \text{ if } Q(x) \text{ then } P(x)\end{aligned}$$

- A universal condition statement is *NOT* logically equivalent to its inverse.

$\forall x \in D, \text{ if } P(x) \text{ then } Q(x)$

$\neq \forall x \in D, \text{ if } \sim P(x) \text{ then } \sim Q(x)$

Using Diagrams to Test for Validity

- Helpful and convincing in many situations.
- To test the validity of an argument diagrammatically,
 - i. represent the truth of both premises with diagrams
 - ii. analyze the diagrams to see whether they necessarily represent the truth of the conclusion as well.

Example 22:

Determine the validity of the following argument using diagrams.

All human beings are mortal.

Zeus is not mortal.

\therefore Zeus is not a human being.

Example 22:

Let H : set of human beings

M : set of those who are mortal

Z : Zeus

The two diagrams fit together in only one way, as shown below.

Therefore, Zeus is not a human being
as Z falls outside the disc of H . Hence,
the argument is valid.

Example 23:

Determine the validity of the following argument using diagrams.

All human beings are mortal.

Felix is mortal.

\therefore Felix is a human being.

Let H : set of human beings

M : set of those who are mortal

F : Felix

Major Premise

The disc of H falls entirely
inside the disc of M .

Minor Premise

F falls inside the disc of M .

The possible conclusions are

- (1) Therefore Felix is not a human being as F falls outside the disc of H .
- (2) Therefore Felix is a human being as F falls inside the disc of H .

There is a contradiction between the conclusions, therefore the argument is invalid.

Example 24:

Determine the validity of the following argument using diagrams.

No polynomial functions have horizontal asymptote.

This function has a horizontal asymptote.

\therefore This function is not a polynomial.

Let P : set of polynomial functions

H : set of functions with horizontal asymptote

T : this particular function

Major Premise

The discs of P and H
are separated.

Minor Premise

T falls inside the disc of H .

The possible conclusion is

Therefore this function is not a polynomial function as T falls outside the disc of P .
Hence, the argument is valid.

Example 25:

Determine the validity of the following argument using diagrams.

All discrete mathematics students can tell a valid argument from an invalid one.

All thoughtful people can tell a valid argument from an invalid one.

\therefore All discrete mathematics students are thoughtful.

Let D : set of discrete mathematics students

V : set of students who can tell a valid argument from an invalid one

T : set of thoughtful people

Major Premise

The disc of D falls entirely inside the disc of V .

Minor Premise

The disc of T falls entirely inside the disc of V .

The possible conclusions are

- (1) All thoughtful people are discrete mathematics students as the disc of T falls entirely inside the disc of D .
- (2) All discrete mathematics students are thoughtful as the disc of D falls entirely inside the disc of T .
- (3) Some discrete mathematics students are thoughtful as there is an intersection between the two discs.
- (4) All discrete mathematics students are not thoughtful as both of the discs are separated.

Since the possible conclusions are contradict to each other, thus the argument is invalid.