# Team notebook

# September 5, 2017

Contents			
1	Data Structure       3        1 BIT       3        2 BigInt       3        3 DisjointSet       3        4 Fraction       9	1	
2	Graph       14         2.1 ArticulationPoint       14         2.2 CentroidDecomposition       14         2.3 EdmondsKarp       15         2.4 EulerianPath       16         2.5 FloyedWarshall       17         2.6 Kosaraju       18         2.7 Kruskal       18	4 4 4 5 6 7 8	
3	Matrix         19           3.1 MatrixExpo         19	_	
4	String       20         4.1 AhoCorasick       26         4.2 KMP       27         4.3 Manacher       25         4.4 SArray(n logn)       25         5.5 SmallestStringRotation       25         4.6 SuffixArray(n logn logn)       26	0 1 2 2	

```
Data Structure
1.1 BIT
using vi = vector < int >;
using vii = vector < vi >;
struct BIT_2D {
   int n;
   vii tree;
   BIT_2D () {}
   BIT_2D ( int _n ) : n( _n ), tree( _n, vi( _n, 0 ) ) {}
   "BIT_2D () {}
   void update_y( int x, int y, int v ) {
      for(; y < n; y + = (y & -y)) {
         tree[x][y] += v;
      }
   }
   void update( int x, int y, int v ) {
      for( ; x<n; x+=(x&-x) ) {</pre>
         update_y( x, y, v );
      }
   int query_y( int x, int y ) {
      int ret = 0;
      for( ; y; y-=(y&-y) ) {
         ret += tree[x][y];
```

}

```
return ret;
   }
   int query( int x, int y ) {
       int ret = 0;
       for(; x; x=(x\&-x)) {
           ret += query_y( x, y );
       return ret;
   }
   int query( int x1, int y1, int x2, int y2 ) {
       return ( query( x2, y2 ) - query( x2, y1-1 ) - query( x1-1, y2 ) +
           query( x1-1, y1-1 ) );
}
struct BIT {
   int n:
   vi tree;
   BIT () {}
   BIT ( int _n ) : n( _n ), tree( _n, 0 ) {}
   "BIT () {}
   void update( int x, int v ) {
       for( ; x<n; y+=(x&-x) ) {</pre>
           tree[x] += v;
       }
   }
   int query( int x ) {
       int ret = 0;
       for( ; x; x-=(x&-x) ) {
           ret += tree[x];
       return ret;
   }
   int query( int x, int y, int x2, int y2 ) {
       return ( query( y ) - query( x-1 ) );
   }
}
```

## 1.2 BigInt

```
using ll = long long;
const int base = 1000000000;
const int base_digits = 9;
struct bigint {
   vector<int> z;
   int sign;
   bigint():
       sign(1) {
   bigint(long long v) {
       *this = v;
   }
   bigint(const string &s) {
       read(s);
   void operator=(const bigint &v) {
       sign = v.sign;
       z = v.z;
   void operator=(long long v) {
       sign = 1;
       if (v < 0)
           sign = -1, v = -v;
       z.clear();
       for (; v > 0; v = v / base)
          z.push_back(v % base);
   }
   bigint operator+(const bigint &v) const {
       if (sign == v.sign) {
          bigint res = v;
          for (int i = 0, carry = 0; i < (int) max(z.size(), v.z.size())</pre>
               || carry; ++i) {
              if (i == (int) res.z.size())
                  res.z.push_back(0);
```

```
res.z[i] += carry + (i < (int) z.size() ? z[i] : 0);
          carry = res.z[i] >= base;
          if (carry)
              res.z[i] -= base;
       }
       return res;
   return *this - (-v);
}
bigint operator-(const bigint &v) const {
   if (sign == v.sign) {
       if (abs() >= v.abs()) {
          bigint res = *this;
          for (int i = 0, carry = 0; i < (int) v.z.size() || carry;</pre>
               ++i) {
              res.z[i] -= carry + (i < (int) v.z.size() ? v.z[i] : 0);
              carry = res.z[i] < 0;
              if (carry)
                  res.z[i] += base;
          }
          res.trim();
          return res;
       }
       return -(v - *this);
   return *this + (-v);
}
void operator*=(int v) {
   if (v < 0)
       sign = -sign, v = -v;
   for (int i = 0, carry = 0; i < (int) z.size() || carry; ++i) {</pre>
       if (i == (int) z.size())
          z.push_back(0);
       long long cur = z[i] * (long long) v + carry;
       carry = (int) (cur / base);
       z[i] = (int) (cur % base);
       //asm("divl %%ecx" : "=a"(carry), "=d"(a[i]) : "A"(cur),
           "c"(base));
   trim();
}
bigint operator*(int v) const {
```

```
bigint res = *this;
   res *= v;
   return res:
friend pair<br/>bigint, bigint> divmod(const bigint &a1, const bigint
    &b1) {
   int norm = base / (b1.z.back() + 1);
   bigint a = a1.abs() * norm;
   bigint b = b1.abs() * norm;
   bigint q, r;
   q.z.resize(a.z.size());
   for (int i = a.z.size() - 1; i >= 0; i--) {
       r *= base:
       r += a.z[i]:
       int s1 = b.z.size() < r.z.size() ? r.z[b.z.size()] : 0;</pre>
       int s2 = b.z.size() - 1 < r.z.size() ? r.z[b.z.size() - 1] : 0;
       int d = ((long long) s1 * base + s2) / b.z.back();
       r -= b * d;
       while (r < 0)
          r += b, --d;
       q.z[i] = d;
   }
   q.sign = a1.sign * b1.sign;
   r.sign = a1.sign;
   q.trim();
   r.trim();
   return make_pair(q, r / norm);
friend bigint sqrt(const bigint &a1) {
   bigint a = a1;
   while (a.z.empty() || a.z.size() % 2 == 1)
       a.z.push_back(0);
   int n = a.z.size();
   int firstDigit = (int) sqrt((double) a.z[n - 1] * base + a.z[n -
   int norm = base / (firstDigit + 1);
   a *= norm;
   a *= norm:
   while (a.z.empty() || a.z.size() % 2 == 1)
```

```
a.z.push_back(0);
   bigint r = (long long) a.z[n - 1] * base + a.z[n - 2];
   firstDigit = (int) sqrt((double) a.z[n - 1] * base + a.z[n - 2]);
   int q = firstDigit;
   bigint res;
   for(int j = n / 2 - 1; j >= 0; j--) {
       for(; ; --q) {
          bigint r1 = (r - (res * 2 * base + q) * q) * base * base +
               (j > 0 ? (long long) a.z[2 * j - 1] * base + a.z[2 * j
               -2]:0);
          if (r1 >= 0) {
              r = r1:
              break;
          }
      }
       res *= base;
       res += q;
       if (j > 0) {
          int d1 = res.z.size() + 2 < r.z.size() ? r.z[res.z.size()</pre>
               + 21 : 0:
          int d2 = res.z.size() + 1 < r.z.size() ? r.z[res.z.size()]
               + 1] : 0:
          int d3 = res.z.size() < r.z.size() ? r.z[res.z.size()] : 0;</pre>
          q = ((long long) d1 * base * base + (long long) d2 * base
               + d3) / (firstDigit * 2);
      }
   }
   res.trim():
   return res / norm;
}
bigint operator/(const bigint &v) const {
   return divmod(*this, v).first;
}
bigint operator%(const bigint &v) const {
   return divmod(*this, v).second;
}
void operator/=(int v) {
   if (v < 0)
```

```
sign = -sign, v = -v;
   for (int i = (int) z.size() - 1, rem = 0; i \ge 0; --i) {
       long long cur = z[i] + rem * (long long) base;
       z[i] = (int) (cur / v);
       rem = (int) (cur % v);
   }
   trim();
}
bigint operator/(int v) const {
   bigint res = *this;
   res /= v;
   return res;
}
int operator%(int v) const {
   if (v < 0)
       v = -v:
   int m = 0:
   for (int i = z.size() - 1; i >= 0; --i)
       m = (z[i] + m * (long long) base) % v;
   return m * sign;
}
void operator+=(const bigint &v) {
   *this = *this + v;
void operator-=(const bigint &v) {
   *this = *this - v;
void operator*=(const bigint &v) {
   *this = *this * v;
void operator/=(const bigint &v) {
   *this = *this / v;
bool operator<(const bigint &v) const {</pre>
   if (sign != v.sign)
       return sign < v.sign;</pre>
   if (z.size() != v.z.size())
       return z.size() * sign < v.z.size() * v.sign;</pre>
   for (int i = z.size() - 1; i >= 0; i--)
       if (z[i] != v.z[i])
           return z[i] * sign < v.z[i] * sign;</pre>
```

```
return false;
}
bool operator>(const bigint &v) const {
   return v < *this;</pre>
}
bool operator<=(const bigint &v) const {</pre>
   return !(v < *this);</pre>
}
bool operator>=(const bigint &v) const {
   return !(*this < v);</pre>
}
bool operator==(const bigint &v) const {
   return !(*this < v) && !(v < *this);</pre>
}
bool operator!=(const bigint &v) const {
   return *this < v || v < *this;
}
void trim() {
   while (!z.empty() && z.back() == 0)
       z.pop_back();
   if (z.empty())
       sign = 1;
}
bool isZero() const {
   return z.empty() || (z.size() == 1 && !z[0]);
}
bigint operator-() const {
   bigint res = *this;
   res.sign = -sign;
   return res;
}
bigint abs() const {
   bigint res = *this;
   res.sign *= res.sign;
   return res;
}
long longValue() const {
   long long res = 0;
   for (int i = z.size() - 1; i >= 0; i--)
```

```
res = res * base + z[i];
   return res * sign;
}
friend bigint gcd(const bigint &a, const bigint &b) {
   return b.isZero() ? a : gcd(b, a % b);
friend bigint lcm(const bigint &a, const bigint &b) {
   return a / gcd(a, b) * b;
void read(const string &s) {
   sign = 1;
   z.clear();
   int pos = 0;
   while (pos < (int) s.size() && (s[pos] == '-' || s[pos] == '+')) {</pre>
       if (s[pos] == '-')
           sign = -sign;
       ++pos;
   for (int i = s.size() - 1; i >= pos; i -= base_digits) {
       int x = 0:
       for (int j = max(pos, i - base_digits + 1); j <= i; j++)</pre>
           x = x * 10 + s[j] - '0';
       z.push_back(x);
   }
   trim();
friend istream& operator>>(istream &stream, bigint &v) {
   string s;
   stream >> s:
   v.read(s);
   return stream;
}
friend ostream& operator<<(ostream &stream, const bigint &v) {</pre>
   if (v.sign == -1)
       stream << '-';
   stream << (v.z.empty() ? 0 : v.z.back());
   for (int i = (int) v.z.size() - 2; i >= 0; --i)
       stream << setw(base_digits) << setfill('0') << v.z[i];</pre>
   return stream;
```

```
static vector<int> convert_base(const vector<int> &a, int old_digits,
    int new_digits) {
   vector<long long> p(max(old_digits, new_digits) + 1);
   p[0] = 1;
   for (int i = 1; i < (int) p.size(); i++)</pre>
       p[i] = p[i - 1] * 10;
   vector<int> res;
   long long cur = 0;
   int cur_digits = 0;
   for (int i = 0; i < (int) a.size(); i++) {</pre>
       cur += a[i] * p[cur_digits];
       cur_digits += old_digits;
       while (cur_digits >= new_digits) {
           res.push_back(int(cur % p[new_digits]));
           cur /= p[new_digits];
           cur_digits -= new_digits;
       }
   }
   res.push_back((int) cur);
   while (!res.empty() && res.back() == 0)
       res.pop_back();
   return res;
}
typedef vector<long long> vll;
static vll karatsubaMultiply(const vll &a, const vll &b) {
   int n = a.size();
   vll res(n + n);
   if (n <= 32) {
       for (int i = 0; i < n; i++)</pre>
           for (int j = 0; j < n; j++)
              res[i + j] += a[i] * b[j];
       return res:
   }
   int k = n >> 1;
   vll a1(a.begin(), a.begin() + k);
   vll a2(a.begin() + k, a.end());
   vll b1(b.begin(), b.begin() + k);
   vll b2(b.begin() + k, b.end());
   vll a1b1 = karatsubaMultiply(a1, b1);
   vll a2b2 = karatsubaMultiply(a2, b2);
```

```
for (int i = 0; i < k; i++)</pre>
       a2[i] += a1[i];
   for (int i = 0; i < k; i++)</pre>
       b2[i] += b1[i];
   vll r = karatsubaMultiply(a2, b2);
   for (int i = 0; i < (int) a1b1.size(); i++)</pre>
       r[i] -= a1b1[i]:
   for (int i = 0; i < (int) a2b2.size(); i++)</pre>
       r[i] -= a2b2[i];
   for (int i = 0; i < (int) r.size(); i++)</pre>
       res[i + k] += r[i];
   for (int i = 0; i < (int) a1b1.size(); i++)</pre>
       res[i] += a1b1[i];
   for (int i = 0; i < (int) a2b2.size(); i++)</pre>
       res[i + n] += a2b2[i];
   return res;
}
bigint operator*(const bigint &v) const {
   vector<int> a6 = convert_base(this->z, base_digits, 6);
   vector<int> b6 = convert_base(v.z, base_digits, 6);
   vll a(a6.begin(), a6.end());
   vll b(b6.begin(), b6.end());
   while (a.size() < b.size())</pre>
       a.push_back(0);
   while (b.size() < a.size())</pre>
       b.push_back(0);
   while (a.size() & (a.size() - 1))
       a.push_back(0), b.push_back(0);
   vll c = karatsubaMultiply(a, b);
   bigint res;
   res.sign = sign * v.sign;
   for (int i = 0, carry = 0; i < (int) c.size(); i++) {</pre>
       long long cur = c[i] + carry;
       res.z.push_back((int) (cur % 1000000));
       carry = (int) (cur / 1000000);
   res.z = convert_base(res.z, 6, base_digits);
   res.trim():
   return res;
long long sod() {
   long long res = 0;
```

```
long long ret = 0;
       for (int i = z.size() - 1; i >= 0; i--) {
           res = res * base + z[i];
           while( res ) {
              ret += res % 10;
              res /= 10;
           }
       }
       return ret;
   }
};
bigint random_bigint(int n) {
   string s;
   for (int i = 0; i < n; i++) {</pre>
       s += rand() \% 10 + 0;
   }
   return bigint(s);
}
```

### 1.3 DisjointSet

```
/**
   Implementation of Disjoint-Set Union Data Structure
   Running time:
       O(nlog(n))
   Usage:
       - call make_set() to reset the set
       - call find_rep() to get the set of the vertex
       - call union_() to merge to sets
   Input:
       - n, number of sets
   Tested Problems:
     TJVA:
       10608 - Friends
       11503 - Virtual Friends
       10583 - Ubiquitous Religions
**/
struct Disjoint_Set {
   int n;
   vector < int > par, cnt, rnk;
```

```
Disjoint_Set( int n ) : n(n), rnk(n), par(n), cnt(n) {}
   void make_set() {
       for(int i=0; i<n; i++) {</pre>
           par[i] = i;
           cnt[i] = 1;
           rnk[i] = 0;
       }
   }
   int find_rep( int x ) {
       if(x != par[ x ]) {
          par[x] = find_rep(par[x]);
       }
       return par[ x ];
   int union_( int u, int v ) {
       if( ( u = find_rep( u ) ) != ( v = find_rep( v ) ) ) {
           if( rnk[ u ] < rnk[ v ] ) {</pre>
              cnt[ v ] += cnt[ u ];
              par[ u ] = par[ v ];
              return cnt[v];
          } else {
              rnk[ u ] = max( rnk[ u ], rnk[ v ] + 1 );
              cnt[ u ] += cnt[ v ];
              par[ v ] = par[ u ];
          }
       }
       return cnt[u];
} DS( sz );
```

#### 1.4 Fraction

```
struct fraction {
    11 up, down;
    11 gcd, lcm;
    void thik_kor() {
        gcd = __gcd( up, down );
        up /= gcd;
        down /= gcd;
}
```

```
fraction operator += ( const fraction &rhs ) {
    gcd = __gcd( down, rhs.down );
    lcm = ( down / gcd ) * rhs.down;
    up = ( ( lcm / down ) * up ) + ( ( lcm / rhs.down ) * rhs.up );
    down = lcm;
    thik_kor();
    return *this;
}
ans, inp;
```

#### 1.5 LCA

```
/**
   Implementation of LCA ( Lowest Common Ancestor ) with sparse table
   Running time:
       O(n * log(n))
   Usage:
       - call dfs(0,0,0)
       - call init()
       - call query() to get output
   Input:
       - Graph
       - n, nodes
   Output:
       - Lowest Common Ancestor
   Tested Problems:
     SPOJ:
       LCA - Lowest Common Ancestor
       QTREE2 - Query on a tree II
**/
#include <bits/stdc++.h>
using namespace std;
const int sz = 1005:
int lg[sz];
int lvl[sz];
int table[sz][12];
int par[sz];
vector < int > G[sz];
void dfs( int fr, int u, int dep ) {
   lvl[u] = dep;
```

```
par[u] = fr;
   for( int v: G[u] ) {
       if( fr == v ) continue;
       dfs( u, v, dep+1 );
   }
}
int init( int n ) {
    memset( table, -1, sizeof table );
    for( int i=0; i<n; i++ ) {</pre>
       table[i][0] = par[i];
    for( int j=1; ( 1 << j ) < n; j++ ) {</pre>
       for( int i=0; i<n; i++ ) {</pre>
           if( table[i][j-1] == -1 ) continue;
           table[i][j] = table[ table[i][j-1] ][j-1];
       }
    for( int i=0; i<10; i++ ) lg[1 << i] = i;</pre>
    for( int i=1; i<sz; i++ ) if( !lg[i] ) lg[i] = lg[i-1];</pre>
}
int query( int n, int p, int q ) {
    int log;
    if( lvl[p] < lvl[q] ) swap( p, q );</pre>
   log = lg[ lvl[ p ] ];
   for( int i=log; i>=0; i-- ) {
       if( lvl[p] - ( 1 << i ) >= lvl[q] ) {
           p = table[ p ][ i ];
       }
    if( p == q ) return p;
    for( int i=log; i>=0; i-- ) {
       if( table[ p ][ i ] != -1 && table[ p ][ i ] != table[ q ][ i ] ) {
           p = table[ p ][ i ];
           q = table[ q ][ i ];
       }
    return par[p];
int dist( int n, int p, int q ) {
    int lca = query( n, p, q );
   return lvl[p] + lvl[q] - 2 * lvl[lca];
```

## 1.6 Mo'sAlgo

```
/**
   Implementation of Mo's Algo with SQRT-Decomposition Data Structure
   Running time:
       O( ( n + q ) * sqrt( n ) * f() )
   Mo's Algo is a algorithm to process queries offline
   For it to work, this condition must be satisified:
       1) There can be no updates in the array
       2) All queries must be known beforehand
   Tested Problems:
     CF:
       220B - Little Elephant and Array
**/
#include <bits/stdc++.h>
using namespace std;
using piii = pair < pair < int, int >, int >;
const int mx = 1e5 + 1;
int BLOCK_SIZE;
int n, m;
int calc;
int ar[mx];
int ans[mx];
unordered_map < int, int > cnt;
piii query[mx];
struct {
   bool operator()( const piii &a, const piii &b ) {
       int block_a = a.first.first / BLOCK_SIZE;
       int block_b = b.first.first / BLOCK_SIZE;
       if( block_a != block_b ) {
           return block_a < block_b;</pre>
       return a.first.second < b.first.second;</pre>
   }
} cmp;
void add( int x ) {
   calc -= (cnt[x] == x ? 1 : 0);
   cnt[x]++:
   calc += (cnt[x] == x ? 1 : 0);
}
void remove( int x ) {
```

```
calc -= (cnt[x] == x ? 1 : 0);
    cnt[x]--;
    calc += (cnt[x] == x ? 1 : 0);
int main() {
    #ifdef LU_SERIOUS
       freopen( "in.txt", "r", stdin );
//
         freopen( "out.txt", "w+", stdout );
    #endif // LU_SERIOUS
    while( ~scanf( "%d %d", &n, &m ) ) {
       BLOCK_SIZE = sqrt( n );
       cnt.clear();
       calc = 0;
       for( int i=0; i<n; i++ ) scanf( "%d", ar+i );</pre>
       for( int i=0; i<m; i++ ) {</pre>
           scanf( "%d %d", &query[i].first.first, &query[i].first.second
               ):
           query[i].second = i;
       }
       sort( query, query+m, cmp );
       int mo 1 = 0, mo r = -1:
       for( int i=0; i<m; i++ ) {</pre>
           int left = query[i].first.first - 1;
           int right = query[i].first.second - 1;
           while( mo_r < right ) {</pre>
               mo_r++;
               add( ar[mo_r] );
           }
           while( mo_r > right ) {
               remove( ar[mo_r] );
               mo_r--;
           while( mo_l < left ) {</pre>
               remove( ar[mo_1] );
               mo_1++;
```

## 1.7 PersistantSegmentTree

```
const int sz = 1e5 + 10;
int ar[sz];
struct node {
   node *left:
   node *right;
   int val;
   node( int val = 0, node *left = nullptr, node *right = nullptr ) {
       this->val = val;
       this->left = left;
       this->right = right;
   }
   void build( int 1, int r ) {
       if( 1 == r ) {
          this->val = ar[1];
          return;
       }
       int mid = (1 + r) >> 1;
       this->left = new node();
       this->right = new node();
       this->left->build( 1, mid );
       this->right->build( mid + 1, r );
```

```
this->val = this->left->val + this->right->val;
   node *update( int 1, int r, int idx, int x ) {
       if( r < idx || idx < 1 ) {</pre>
           return this:
       }
       if( 1 == r ) {
           node *ret = new node( this->val, this->left, this->right );
          ret->val += x;
          return ret:
       int mid = (1 + r) >> 1;
       node *ret = new node( this->val );
       ret->left = this->left->update( 1, mid, idx, x );
       ret->right = this->right->update( mid + 1, r, idx, x );
       ret->val = ret->left->val + ret->right->val;
       return ret;
   }
   int query( int 1, int r, int i, int j ) {
       if(r < i || 1 > j ) {
          return 0;
       if( i <= 1 && r <= j ) {</pre>
          return this->val;
       int mid = (1 + r) >> 1;
       return this->left->query( l, mid, i, j ) + this->right->query( mid
           + 1, r, i, j);
} *root[sz]:
int main() {
   ar[] = \{ 1, 2, 3 \};
   root[0] = new node();
   root[0]->build(0, 2);
   root[1] = root[0]->update(0, 2, 1, 1);
   root[1] = root[1]->update(0, 2, 1, 1);
   printf( "%d\n", root[0].query( 0, 2, 0, 2 ) );
   printf( "%d\n", root[1].query( 0, 2, 0, 2 ) );
```

#### 1.8 SegmentTree

```
struct info {
   int prop, sum;
} tree[ mx * 3 ];
void update( int node, int b, int e, int i, int j, int x ) {
// cerr << b << " " << e << " " << i << " " << j << " " << x << "\n";
   if( i > e || j < b ) {</pre>
       return:
   }
   if( b >= i && e <= j ) {</pre>
       tree[node].sum = (e - b + 1) * x;
       tree[node].prop = x;
       return;
   }
   int left = node << 1;</pre>
   int right = left | 1;
   int mid = (b + e) >> 1;
   if( tree[node].prop != -1 ) {
       tree[left].sum = ( mid - b + 1 ) * tree[node].prop;
       tree[right].sum = ( e - mid ) * tree[node].prop;
       tree[node].sum = tree[left].sum + tree[right].sum;
       tree[left].prop = tree[node].prop;
       tree[right].prop = tree[node].prop;
       tree[node].prop = -1;
   }
   update(left, b, mid, i, j, x);
   update(right, mid + 1, e, i, j, x);
   tree[node].sum = tree[left].sum + tree[right].sum;
}
int query( int node, int b, int e, int i, int j ) {
   if( i > e || j < b ) {
       return 0;
   }
   if(b \ge i and e \le j) {
       return tree[node].sum;
   }
   int left = node << 1;</pre>
```

```
int right = left | 1;
int mid = (b + e) >> 1;

if( tree[node].prop != -1 ) {
    tree[left].sum = ( mid - b + 1 ) * tree[node].prop;
    tree[right].sum = ( e - mid ) * tree[node].prop;
    tree[node].sum = tree[left].sum + tree[right].sum;
    tree[left].prop = tree[node].prop;
    tree[right].prop = tree[node].prop;
    tree[node].prop = -1;
}

int p1 = query( left, b, mid, i, j );
int p2 = query( right, mid + 1, e, i, j );
return p1 + p2;
}
```

### 1.9 SqrtDecomposition

```
/**
   Implementation of SQRT-Decomposition Data Structure
   Running time:
       O((n+q) * sqrt(n) * f())
   Usage:
       - call int() to initialize the array
       - call update() to update the element in a position
       - call query() to get ans from segment [L...R]
   Input:
       - n. number of elements
       - n elements
       - q queries
   Tested Problems:
     lightOJ:
       1082 - Array Queries
#include <bits/stdc++.h>
using namespace std;
const int mx = 1e5 + 1;
const int sz = 1e3 + 1;
const int inf = 1e9;
```

```
int BLOCK_SIZE;
int n, q, t, cs, x, y;
int BLOCKS[sz];
int ar[mx];
int getID( int idx ) {
    return idx / BLOCK_SIZE;
}
void init() {
    for( int i=0; i<sz; i++ ) BLOCKS[i] = inf;</pre>
}
void update( int idx, int val ) {
    int id = getID( idx );
    BLOCKS[id] = min( val, BLOCKS[id] );
}
int query( int 1, int r ) {
    int le = getID( 1 );
    int ri = getID( r );
    int ret = inf;
    if( le == ri ) {
       for( int i=1; i<=r; i++ ) {</pre>
           ret = min( ret, ar[i] );
       }
       return ret;
    }
    for( int i=1; i<(le+1)*BLOCK_SIZE; i++ ) ret = min( ret, ar[i] );</pre>
    for( int i=le+1; i<ri; i++ ) ret = min( ret, BLOCKS[i] );</pre>
    for( int i=ri*BLOCK_SIZE; i<=r; i++ ) ret = min( ret, ar[i] );</pre>
    return ret;
}
int main() {
    #ifdef LU SERIOUS
       freopen( "in.txt", "r", stdin );
//
         freopen( "out.txt", "w+", stdout );
    #endif // LU_SERIOUS
    scanf( "%d", &t );
    for( cs=1; cs<=t; cs++ ) {</pre>
       scanf( "%d %d", &n, &q );
```

```
BLOCK_SIZE = sqrt( n );
init();
for( int i=0; i<n; i++ ) {
      scanf( "%d", &ar[i] );
      update( i, ar[i] );
}
printf( "Case %d:\n", cs );
for( int i=0; i<q; i++ ) {
      scanf( "%d %d", &x, &y );
      printf( "%d\n", query( x-1, y-1 ) );
}
return 0;
}</pre>
```

# 2 Graph

#### 2.1 ArticulationPoint

```
/**
An O(V+E) approach:
- perform a DFS on the graph.
- compute d(i) and low(i) for each vertex 1 ... i
       d(i): dfs number of i, represents the discovery time.
       low(i): the least dfn reachable from i through a path consisting
           of zero or
              more edges follwoing by zero or one back edges.
- vertex u is an AP if and only if:
       - u is the root of the dfs tree and has at least two children.
       - u is not the root and has a child v for which low(v) >= d(u).
**/
#include <bits/stdc++.h>
using namespace std;
const int mx = 1e4 + 10;
vector < int > G[mx];
int tim, root, n, m, a, b;
int ap[mx], vis[mx], low[mx], d[mx], par[mx];
void ap_dfs(int u) {
   tim++;
```

```
int cnt = 0;
   low[u] = tim;
   d[u] = tim;
   vis[u] = 1;
   int v;
   for(int i=0; i<G[u].size(); i++) {</pre>
       v = G[u][i];
       if( v == par[u] ) continue;
       if( !vis[v] ) {
           par[u] = v;
           ap_dfs( v );
           low[u] = min( low[u], low[v] );
           /// d[u] < low[v] if bridge is needed
           if( d[u] <= low[v] && u != root ) {</pre>
               ap[u] = 1;
           }
           cnt++;
       } else {
           low[u] = min(low[u], d[v]);
       if( u == root && cnt > 1 ) ap[u] = 1;
   }
}
int main() {
       return 0:
```

# 2.2 CentroidDecomposition

```
321C/322E - Ciel the Commander
**/
const int sz = 1e5 + 10;
vector < int > G[sz];
char ans[sz];
int tr[sz], fl[sz];
void dfs( int u, int p ) {
   tr[u] = 1;
   for( int v: G[u] ) {
       if( v != p && !fl[v] ) {
          dfs( v, u );
          tr[u] += tr[v];
       }
}
int centroid( int u ) {
   dfs(u, -1);
   int ret = u;
   int found = 0, par = -1;
   while( 1 ) {
       found = 0;
       for( int v: G[ret] ) {
           if( !fl[v] && v != par && tr[v] >= ( tr[u] + 1 ) / 2 ) {
              found = 1;
              par = ret;
              ret = v;
              break;
          }
       if( !found ) break;
   return ret;
}
void rec( int u, char a ) {
   u = centroid( u );
   fl[u] = 1;
   ans[u] = a;
   for( int v: G[u] ) {
       if( !fl[v] ) rec( v, a + 1 );
   return;
```

## 2.3 EdmondsKarp

```
/**
   Implementation of Edmonds-Karp max flow algorithm
   Running time:
       0(|V|*|E|^2)
   Usage:
       - add edges by add_edge()
       - call max_flow() to get maximum flow in the graph
   Input:
       - n, number of nodes
       - directed, true if the graph is directed
       - graph, constructed using add_edge()
       - source, sink
   Output:
       - Maximum flow
   Tested Problems:
     CF:
       653D - Delivery Bears
     UVA:
       820 - Internet Bandwidth
       10330 - Power Transmission
**/
#include <bits/stdc++.h>
using namespace std;
const int INF = 1e9;
struct edmonds_karp {
   int n;
   vector < int > par;
   vector < bool > vis;
   vector < vector < int > > graph;
   edmonds_karp () {}
   edmonds_karp( int _n ) : n( _n ), par( _n ), vis( _n ), graph( _n,
       vector< int > ( _n, 0 ) ) {}
   ~edmonds_karp() {}
   void add_edge( int from, int to, int cap, bool directed ) {
       this->graph[ from ][ to ] += cap;
       this->graph[ to ][ from ] = directed ? graph[ to ][ from ] + cap :
           graph[ to ][ from ];
   }
```

```
bool bfs( int src, int sink ) {
   int u:
   fill( vis.begin(), vis.end(), false );
   fill( par.begin(), par.end(), -1 );
   vis[ src ] = true;
   queue < int > q;
   q.push( src );
   while( !q.empty() ) {
       u = q.front();
       q.pop();
       if( u == sink ) return true;
       for(int i=0; i<n; i++) {</pre>
           if( graph[u][i] > 0 and not vis[i] ) {
              q.push( i );
              vis[ i ] = true;
              par[ i ] = u;
          }
       }
   }
   return par[ sink ] != -1;
int min val( int i ) {
   int ret = INF;
   for( ; par[ i ] != -1; i = par[ i ] ) {
       ret = min( ret, graph[ par[i] ][ i ] );
   }
   return ret;
void augment_path( int val, int i ) {
   for( ; par[ i ] != -1; i = par[ i ] ) {
       graph[ par[i] ][ i ] -= val;
       graph[ i ][ par[i] ] += val;
   }
}
int max_flow( int src, int sink ) {
   int min_cap, ret = 0;
   while( bfs( src, sink ) ) {
       augment_path( min_cap = min_val( sink ), sink );
       ret += min_cap;
   return ret;
```

```
};
```

#### 2.4 EulerianPath

```
Implementation of Hierholzer's algorithm for finding Euler Path /
       Circuit
   Running time:
       O( | E | )
   Input:
       - adj, graph
   Tested Problems:
       CodeChef:
           TOURISTS - Tourists in Mancunia
**/
struct Edge;
typedef list< Edge >::iterator iter;
struct Edge {
   int next_vertex;
   iter reverse_edge;
   Edge( int next_vertex )
       :next_vertex(next_vertex)
       { }
};
const int sz = 1e5 + 10;
int num_vertices;
list < Edge > adj[ max_vertices ];
vector< int > path;
void find_path( int v ) {
   while( adj[v].size() > 0 ) {
       int vn = adj[v].front().next_vertex;
       adj[vn].erase( adj[v].front().reverse_edge );
       adj[v].pop_front();
       find_path( vn );
   path.push_back( v );
}
```

```
void add_edge( int a, int b ) {
   adj[ a ].push_front( Edge( b ) );
   iter ita = adj[ a ].begin();
   adj[ b ].push_front( Edge( a ) );
   iter itb = adj[ b ].begin();
   ita->reverse_edge = itb;
   itb->reverse_edge = ita;
}
```

## 2.5 FloyedWarshall

```
/**
   Implementation of Floyd Warshall Alogrithm
   Running time:
       O( |v| ^ 3 )
   Input:
       - n, number vertex
       - graph, inputed as an adjacency matrix
   Tested Problems:
     UVA:
       544 - Heavy Cargo - MaxiMin path
       567 - Risk - APSP
**/
using vi = vector < int >;
using vvi = vector < vi >;
/// mat[i][i] = 0, mat[i][j] = distance from i to j, path[i][j] = i
void APSP( vvi &mat, vvi &path ) {
   int V = mat.size();
   for( int via=0; via; via<V; via++ ) {</pre>
       for( int from=0; from<V; from++ ) {</pre>
           for( int to=0; to<V; to++ ) {</pre>
               if( mat[ from ][ via ] + mat[ via ][ to ] < mat[ from ][</pre>
                  mat[ from ][ to ] = mat[ from ][ via ] + mat[ via ][ to
                       1:
                  path[ from ][ to ] = path[ via ][ to ];
```

```
}
           }
       }
   }
}
/// prints the path from i to j
void print( int i, int j ) {
    if( i != j ) {
       print( i, path[i][j] );
   }
    cout << j << "\n";
}
/// check if negative cycle exists
bool negative_cycle( vvi &mat ) {
    APSP( mat );
    return mat[0][0] < 0;</pre>
}
void transtitive_closure( vvi &mat ) {
    int V = mat.size();
    for( int via=0; via; via<V; via++ ) {</pre>
       for( int from=0; from<V; from++ ) {</pre>
           for( int to=0; to<V; to++ ) {</pre>
               mat[ from ][ to ] |= ( mat[ from ][ via ] & mat[ via ][ to
                   ]);
           }
       }
   }
}
/// finding a path between two nodes that maximizes the minimum cost
void mini_max( vvi &mat ) {
    int V = mat.size();
    for( int via=0; via; via<V; via++ ) {</pre>
       for( int from=0; from<V; from++ ) {</pre>
           for( int to=0; to<V; to++ ) {</pre>
```

```
mat[ from ][ to ] = min( mat[ from ][ to ], max( mat[ from
                   ][ via ], mat[ via ][ to ] ) );
           }
       }
/// finding a path between two nodes that minimizes the maximum cost
/// eg: max load a truck can carry from one node to another node where
/// the paths have weight limit
void maxi_min( vvi &mat ) {
   int V = mat.size();
   for( int via=0; via; via<V; via++ ) {</pre>
       for( int from=0; from<V; from++ ) {</pre>
           for( int to=0: to<V: to++ ) {</pre>
              mat[ from ][ to ] = max( mat[ from ][ to ], min( mat[ from
                   [ via ], mat[ via ][ to ] );
          }
       }
   }
```

# 2.6 Kosaraju

```
#include <bits/stdc++.h>
using namespace std;
int p, t;
bool vis[1001];
vector<int> G[1001], gT[1001];
map<string,int> mp;
stack < int > top_sorted;

void dfs_top_sort(int u) {
    vis[u] = true;
    for(int v: G[u]) {
        if(!vis[v]) {
            dfs_top_sort( v );
        }
}
```

```
top_sorted.push( u );
}
void top_sort() {
       for(int i=1; i<=p; i++) {</pre>
              if(!vis[i]) {
                      dfs_top_sort(i);
       }
}
void dfs_kosaraju(int u) {
       vis[u] = true;
       for(int v: gT[u]) {
              if(!vis[v]) {
                      dfs_kosaraju( v );
              }
       }
}
int kosaraju() {
       memset( vis, false, sizeof(vis) );
       top_sort();
       int u, ret = 0;
       memset( vis, false, sizeof(vis) );
       while(!top_sorted.empty()) {
              u = top_sorted.top();
              top_sorted.pop();
              if(!vis[u])
                      dfs_kosaraju( u ), ret++;
       }
       return ret;
}
```

#### 2.7 Kruskal

```
/**
   Implementation of Kruskal's minimum spanning tree algorithm
   Running time:
        O(|E|log|V|)
   Usage:
        - initialize by calling init()
```

```
- add edges by add_edge()
       - call kruskal() to generate minimum spanning tree
   Input:
       - n, number of nodes, provided when init() is called
       - graph, constructed using add_edge()
   Output:
       - weight of minimum spanning tree
       - prints the mst
   Tested Problems:
       UVA:
           1208 - Oreon
#include <bits/stdc++.h>
using namespace std;
struct edge {
   int u, v, cost;
   bool operator < (const edge& other) const{</pre>
       if( other.cost == this->cost ) {
       if( other.u == this->u ) {
           return other.v > this->v;
       } else {
           return other.u > this->u;
       } else {
           return other.cost > this->cost;
};
vector< edge > edges;
vector < int > par, cnt, rank;
int N;
void init( int n ) {
   N = n;
   par.resize( n );
   cnt.resize( n );
   rank.resize( n );
}
void add_edge( int u, int v, int c ) {
   edges.push_back( { u, v, c } );
```

```
void make_set() {
   for(int i=0; i<N; i++) {</pre>
       par[i] = i;
       cnt[i] = 1;
       rank[i] = 0;
   }
}
int find_rep( int x ) {
   if(x != par[ x ]) {
       par[ x ] = find_rep( par[ x ] );
   }
   return par[ x ];
}
int kruskal() {
   int ret = 0;
   make_set();
   sort( edges.begin(), edges.end() );
   cout << "Case " << ++cs << ":\n";
   for( edge e : edges ) {
       int u = e.u;
       int v = e.v:
       if( ( u = find_rep( u ) ) != ( v = find_rep( v ) ) ) {
           if( rank[ u ] < rank[ v ] ) {</pre>
               cnt[v] += cnt[u]:
              par[ u ] = par[ v ];
           } else {
              rank[ u ] = max( rank[ u ], rank[ v ] + 1 );
               cnt[ u ] += cnt[ v ];
              par[ v ] = par[ u ];
           cout << city[ e.u ] << "-" << city[ e.v ] << " " << e.cost <<</pre>
               "\n";
           ret += e.cost;
       }
   return ret;
}
```

### 3 Matrix

### 3.1 MatrixExpo

```
Implementation of Matrix Exponentiation
   Running time:
       O( log( n ) )
   Input:
       - n, exponent
       - recurrence matrix, power of which to be determined
   Tested Problems:
     UVA:
       10229 - Modular Fibonacci
       10518 - How Many Calls?
       12470 - Tribonacci
**/
const int mat_sz = 2;
struct Matrix {
   int a[mat_sz] [mat_sz];
   void clear() {
       memset(a, 0, sizeof(a));
   void one() {
       for( int i=0; i<mat_sz; i++ ) {</pre>
           for( int j=0; j<mat_sz; j++ ) {</pre>
              a[i][j] = i == j;
          }
       }
   }
   Matrix operator + (const Matrix &b) const {
       Matrix tmp;
       tmp.clear();
       for (int i = 0; i < mat_sz; i++) {</pre>
           for (int j = 0; j < mat_sz; j++) {</pre>
              tmp.a[i][j] = a[i][j] + b.a[i][j];
              if (tmp.a[i][j] >= mod) {
                  tmp.a[i][j] -= mod;
           }
       }
       return tmp;
   Matrix operator * (const Matrix &b) const {
```

```
Matrix tmp;
       tmp.clear();
       for (int i = 0; i < mat_sz; i++) {</pre>
           for (int j = 0; j < mat_sz; j++) {</pre>
               for (int k = 0; k < mat_sz; k++) {</pre>
                   tmp.a[i][k] += (long long)a[i][j] * b.a[j][k] % mod;
                  if (tmp.a[i][k] >= mod) {
                       tmp.a[i][k] -= mod;
                  }
               }
           }
       return tmp;
    Matrix pw(int x) {
       Matrix ans, num = *this;
       ans.one();
       while (x > 0) {
           if (x & 1) {
               ans = ans * num;
           }
           num = num * num;
           x >>= 1;
       }
       return ans;
   }
};
```

# 4 String

## 4.1 AhoCorasick

```
const int sz = 1e6 + 10;
const int MAX = 150 * 70 + 100;
char inp[sz], s[155][75];
int cnt[155];

struct AhoCorasick {
   vector < int > mark[MAX + 7];
   int state, failure[MAX + 7];
   int trie[MAX + 7][ 26 ];
```

```
AhoCorasick() {
   init();
}
void init() {
   mark[0].clear();
   fill( trie[0], trie[0] + 26, -1 );
   state = 0;
}
int value( char c ) {
   return c - 'a';
}
void add( char *s, int t ) {
   int root = 0, id;
   for( int i=0; s[i]; i++ ) {
       id = value( s[i] );
       if( trie[ root ][ id ] == -1 ) {
          trie[ root ][ id ] = ++state;
           mark[state].clear();
           fill( trie[state], trie[state + 1] + 26, - 1 );
       root = trie[ root ][ id ];
   mark[ root ].push_back( t );
}
void computeFailure() {
   queue < int > Q;
   failure[0] = 0;
   for( int i=0; i<26; i++ ) {</pre>
       if( trie[ 0 ][ i ] != -1 ) {
          failure[ trie[ 0 ][ i ] ] = 0;
           Q.push( trie[ 0 ][ i ]);
       else trie[ 0 ][ i ] = 0;
   while( !Q.empty() ) {
       int u = Q.front();
       Q.pop();
       for( int v: mark[ failure[ u ] ] ) mark[ u ].push_back( v );
       for( int i=0; i<26; i++ ) {</pre>
           if( trie[ u ][ i ] != -1 ) {
              failure[ trie[ u ][ i ] ] = trie[ failure[ u ] ][ i ];
```

```
Q.push( trie[ u ][ i ] );
}
else trie[ u ][ i ] = trie[ failure[ u ] ][ i ];
}
}
automata;

void countFreq() {
  for( int i=0,root=0,id; inp[i]; i++ ) {
    id = automata.value( inp[i] );
    root = automata.trie[ root ][ id ];
    if( root == 0 ) continue;
    for( int v: automata.mark[ root ] ) cnt[v]++;
}
}
```

#### 4.2 KMP

```
/// complexity : o( n + m )
///solution reference loj 1255 Substring Frequency
#include <bits/stdc++.h>
using namespace std;
int t;
const int mx = 1e6 + 10;
char a[mx], b[mx];
int table[mx], lenA, lenB;
void hash_table( char *s ) {
       table[ 0 ] = 0;
       int i = 1, j = 0;
       while( i < lenB ) {</pre>
              if( s[i] == s[j] ) {
                      j++;
                      table[i] = j;
                      i++;
              } else {
                      if( j ) {
                             j = table[j - 1];
                      } else {
                             table[i] = 0;
                             i++;
```

```
}
              }
       }
}
int kmp( char *s, char *m ) {
       hash_table( m );
       int i = 0, j = 0;
       int ans = 0;
       while( i < lenA ) {</pre>
               while( i < lenA && j < lenB && s[i] == m[j] ) {
                      i++;
                      j++;
              if( j == lenB ) {
                      j = table[j - 1];
                      ans++;
              } else if( i < lenA && s[i] != m[j] ) {</pre>
                      if( j ) {
                              j = table[j - 1];
                      } else {
                             i++;
                      }
              }
       return ans;
}
int main() {
#ifdef LU_SERIOUS
       freopen("in.txt", "r", stdin);
#endif // LU_SERIOUS
       scanf( "%d", &t );
       for(int cs=1; cs<=t; cs++) {</pre>
               lenA = 0; lenB = 0;
               scanf("%s", &a);
               scanf("%s", &b);
               lenA = strlen( a );
              lenB = strlen( b );
               printf( "Case %d: %d\n", cs, kmp( a, b ) );
       }
       return 0;
}
```

#### 4.3 Manacher

```
const int sz = 2e5 + 10:
char inp[sz], str[sz];
int LPS[sz];
int call(){
   int len = 0;
   str[ len++ ]='*';
   for( int i=0; inp[i]; i++ ) {
       str[ len++ ] = inp[i];
       str[ len++ ] = '*';
   }
   str[len] = '\0';
   int c = 0, r = 0, ans = 0;
   for( int i=1; i<len-1; i++ ) {</pre>
       int _i = c - ( i - c );
       if( r > i ) LPS[i] = min( LPS[_i] , r - i );
       else LPS[i] = 0;
       while( i - 1 - LPS[i] >= 0 && str[ i - 1 - LPS[i] ] == str[ i + 1
           + LPS[i] ] ) {
          LPS[i]++;
       }
       if( i + LPS[i] > r ) {
          r = i + LPS[i];
          c = i:
       ans = max( ans, LPS[i] );
   }
   return ans;
```

# 4.4 SArray( n logn )

```
int n;
                                          // the length of input string
int RA[MAX_N], tempRA[MAX_N];  // rank array and temporary rank array
int SA[MAX_N], tempSA[MAX_N]; // suffix array and temporary suffix array
int c[MAX_N];
                                             // for counting/radix sort
char P[MAX_N];
                             // the pattern string (for string matching)
                                        // the length of pattern string
int m;
int Phi[MAX_N];
                                  // for computing longest common prefix
int PLCP[MAX_N];
int LCP[MAX_N]; // LCP[i] stores the LCP between previous suffix T+SA[i-1]
// and current suffix T+SA[i]
bool cmp(int a, int b)
   return strcmp(T + a, T + b) < 0; // compare</pre>
void constructSA slow()
                                    // cannot go beyond 1000 characters
   for (int i = 0; i < n; i++) SA[i] = i; // initial SA: {0, 1, 2, ...,
   sort(SA, SA + n, cmp); // sort: O(n log n) * compare: O(n) = O(n^2)
                                                              // O(n)
void countingSort(int k)
   int i, sum, maxi = max(300, n); // up to 255 ASCII chars or length of
   memset(c, 0, sizeof c);
                                                // clear frequency table
   for (i = 0; i < n; i++) // count the frequency of each integer rank</pre>
       c[i + k < n ? RA[i + k] : 0]++;
   for (i = sum = 0; i < maxi; i++)</pre>
       int t = c[i];
       c[i] = sum;
       sum += t;
   for (i = 0; i < n; i++)
                                 // shuffle the suffix array if necessary
       tempSA[c[SA[i]+k < n ? RA[SA[i]+k] : 0]++] = SA[i];
   for (i = 0; i < n; i++)</pre>
                                           // update the suffix array SA
       SA[i] = tempSA[i];
}
```

```
void constructSA()
                         // this version can go up to 100000 characters
{
   int i, k, r;
   for (i = 0; i < n; i++) RA[i] = T[i];</pre>
                                                  // initial rankings
   for (i = 0; i < n; i++) SA[i] = i; // initial SA: {0, 1, 2, ..., n-1}
   for (k = 1; k < n; k <<= 1) // repeat sorting process log n times
       countingSort(k); // actually radix sort: sort based on the second
           item
                              // then (stable) sort based on the first
       countingSort(0);
       tempRA[SA[O]] = r = 0;
                                      // re-ranking; start from rank r =
           0
       for (i = 1; i < n; i++)</pre>
                                             // compare adjacent suffixes
           tempRA[SA[i]] = // if same pair => same rank r; otherwise,
               increase r
              (RA[SA[i]] == RA[SA[i-1]] &\& RA[SA[i]+k] == RA[SA[i-1]+k])
                  ? r : ++r:
       for (i = 0; i < n; i++)</pre>
                                              // update the rank array RA
          RA[i] = tempRA[i];
       if (RA[SA[n-1]] == n-1) break;
                                              // nice optimization trick
   }
}
void computeLCP_slow()
   LCP[0] = 0:
                                                      // default value
   for (int i = 1; i < n; i++)
                                           // compute LCP by definition
       int L = 0;
                                                   // always reset L to 0
       while (T[SA[i] + L] == T[SA[i-1] + L]) L++; // same L-th char, L++
       LCP[i] = L;
   }
}
void computeLCP()
{
   int i, L;
   Phi[SA[0]] = -1;
                                                      // default value
   for (i = 1; i < n; i++)
                                                 // compute Phi in O(n)
       Phi[SA[i]] = SA[i-1]; // remember which suffix is behind this
           suffix
   for (i = L = 0; i < n; i++)
                                       // compute Permuted LCP in O(n)
       if (Phi[i] == -1)
```

```
{
          PLCP[i] = 0; // special case
           continue;
       while (T[i + L] == T[Phi[i] + L]) L++; // L increased max n times
       PLCP[i] = L;
       L = \max(L-1, 0);
                                               // L decreased max n times
   for (i = 0; i < n; i++)</pre>
                                                 // compute LCP in O(n)
       LCP[i] = PLCP[SA[i]]; // put the permuted LCP to the correct
}
ii stringMatching()
                                      // string matching in O(m log n)
   int lo = 0, hi = n-1, mid = lo;
                                            // valid matching = [0..n-1]
   while (lo < hi)</pre>
                                                    // find lower bound
                                                    // this is round down
       mid = (lo + hi) / 2:
       int res = strncmp(T + SA[mid], P, m); // try to find P in suffix
           'mid'
       if (res >= 0) hi = mid;
                                   // prune upper half (notice the >=
           sign)
                    lo = mid + 1:
                                         // prune lower half including
       else
           mid
                                       // observe '=' in "res >= 0" above
   if (strncmp(T + SA[lo], P, m) != 0) return ii(-1, -1); // if not found
   ii ans;
   ans.first = lo;
   10 = 0:
   hi = n - 1;
   mid = lo:
   while (lo < hi)</pre>
                             // if lower bound is found, find upper bound
       mid = (lo + hi) / 2;
       int res = strncmp(T + SA[mid], P, m);
       if (res > 0) hi = mid;
                                                      // prune upper half
                                         // prune lower half including
       else
                   lo = mid + 1;
           mid
                             // (notice the selected branch when res == 0)
   if (strncmp(T + SA[hi], P, m) != 0) hi--;
                                                        // special case
   ans.second = hi:
   return ans;
} // return lower/upperbound as first/second item of the pair,
    respectively
```

```
ii LRS()
                         // returns a pair (the LRS length and its index)
   int i, idx = 0, maxLCP = -1;
   for (i = 1; i < n; i++)</pre>
                                                // O(n), start from i = 1
       if (LCP[i] > maxLCP)
           maxLCP = LCP[i], idx = i;
   return ii(maxLCP, idx);
}
int owner(int idx)
   return (idx < n-m-1) ? 1 : 2;</pre>
ii LCS()
                         // returns a pair (the LCS length and its index)
   int i, idx = 0, maxLCP = -1;
   for (i = 1; i < n; i++)</pre>
                                                // O(n), start from i = 1
       if (owner(SA[i]) != owner(SA[i-1]) && LCP[i] > maxLCP)
           maxLCP = LCP[i], idx = i;
   return ii(maxLCP, idx);
}
int main()
   //printf("Enter a string T below, we will compute its Suffix
        Array:\n");
   strcpy(T, "GATAGACA");
   //T = "ABCDE"
   n = (int)strlen(T);
   T[n++] = '$':
   // if '\n' is read, uncomment the next line
   //T[n-1] = '$'; T[n] = 0;
   constructSA_slow();
                                                         // O(n^2 \log n)
   printf("The Suffix Array of string T = \%s' is shown below (O(n^2 \log T))
        n) version):\n", T);
   printf("i\tSA[i]\tSuffix\n");
   for (int i = 0; i < n; i++) printf("%2d\t%2d\t%s\n", i, SA[i], T +
        SA[i]);
   constructSA();
                                                           // O(n log n)
   printf("\nThe Suffix Array of string T = '%s' is shown below (O(n log
        n) version):\n", T);
```

```
printf("i\tSA[i]\tSuffix\n");
for (int i = 0; i < n; i++) printf("%2d\t%2d\t%s\n", i, SA[i], T +
    SA[i]);
computeLCP();
                                                             // O(n)
// LRS demo
ii ans = LRS():
                             // find the LRS of the first input string
char lrsans[MAX_N];
strncpy(lrsans, T + SA[ans.second], ans.first);
printf("\nThe LRS is '\%s' with length = \%d\n\n", lrsans, ans.first);
// stringMatching demo
//printf("\nNow, enter a string P below, we will try to find P in
    T: n");
strcpy(P, "A");
m = (int)strlen(P);
// if '\n' is read, uncomment the next line
//P[m-1] = 0; m--;
ii pos = stringMatching();
if (pos.first != -1 && pos.second != -1)
   printf("%s is found SA[%d..%d] of %s\n", P, pos.first, pos.second,
        T):
   printf("They are:\n");
   for (int i = pos.first; i <= pos.second; i++)</pre>
       printf(" %s\n", T + SA[i]);
else printf("%s is not found in %s\n", P, T);
// LCS demo
//printf("\nRemember, T = \frac{\%}{n} \nNow, enter another string P: n, T);
// T already has '$' at the back
strcpy(P, "CATA");
m = (int)strlen(P);
// if '\n' is read, uncomment the next line
//P[m-1] = 0; m--;
strcat(T, P);
                                                         // append P
strcat(T, "#");
                                              // add '$' at the back
n = (int)strlen(T);
                                                         // update n
// reconstruct SA of the combined strings
constructSA();
                                                       // O(n log n)
computeLCP();
                                                             // O(n)
printf("\nThe LCP information of 'T+P' = '%s':\n", T);
```

### 4.5 SmallestStringRotation

```
/**
   Implementation of Lexicographically smallest string rotation
   Running time:
       0(2 * s.size())
   Input:
       - s, string
   Tested Problems:
       UVA:
          719 - Glass Beads
       DevSkill:
          DCP-207: Mina and Raju Part 2
**/
const int sz = 1e5 + 10;
int f[sz];
int calc( const string& s ) {
   int n = s.size();
   string t = s + s;
   memset( f, -1, sizeof f );
   int k = 0:
   for( int j = 1; j < 2 * n; ++j ) {
       int i = f[j - k - 1];
       while ( i != -1 \&\& t[j] != t[k + i + 1] ) {
          if(t[i] < t[k+i+1]) {
              k = j - i - 1;
          }
          i = f[i];
```

```
}
if( i == -1 && t[j] != t[k + i + 1] ) {
    if( t[j] < t[k + i + 1] ) {
        k = j;
    }
    f[j - k] = -1;
} else {
    f[j - k] = i + 1;
}
return k;
}</pre>
```

### 4.6 SuffixArray( n logn logn )

```
/**
   Implementation of Suffix Array
   Running time:
       O(n \log(n) \log(n))
       - s, string for that suffix array to be completed
   Output:
       - Suffix Array
   Tested Problems:
     SPOJ:
       SARRAY - Suffix Array
**/
#include <bits/stdc++.h>
using namespace std;
struct suffix {
   int index;
   int rank[2];
   bool operator < ( const suffix &other ) const {</pre>
       if( this->rank[0] == other.rank[0] ) {
           return this->rank[1] < other.rank[1];</pre>
       return this->rank[0] < other.rank[0];</pre>
};
vector < int > buildSuffixArray( const string &s ) {
   int n = int( s.size() );
```

```
vector < int > sufArray;
   vector < suffix > suffixes( n );
   for( int i=0; i<n; i++ ) {</pre>
       suffixes[i].index = i;
       suffixes[i].rank[0] = s[i];
       suffixes[i].rank[1] = i + 1 < n ? s[i+1] : -1;
   }
   vector < int > ind( n );
   int nextIndex, rank, prev_rank, n_2 = n << 1;</pre>
   sort( suffixes.begin(), suffixes.end() );
   for( int k=4; k<n_2; k<<=1 ) {</pre>
       rank = 0;
       prev_rank = suffixes[0].rank[0];
       suffixes[0].rank[0] = rank;
       ind[ suffixes[0].index ] = 0;
       for( int i=1: i<n: i++ ) {</pre>
           if( suffixes[i].rank[0] == prev_rank && suffixes[i].rank[1] ==
               suffixes[i-1].rank[1] ) {
              prev_rank = suffixes[i].rank[0];
              suffixes[i].rank[0] = rank;
          } else {
              prev_rank = suffixes[i].rank[0];
              suffixes[i].rank[0] = ++rank;
           }
           ind[ suffixes[i].index ] = i;
       for( int i=0: i<n: i++ ) {</pre>
           nextIndex = suffixes[i].index + k / 2;
           suffixes[i].rank[1] = nextIndex < n ? suffixes[ ind[ nextIndex</pre>
               ] ].rank[0] : -1;
       sort( suffixes.begin(), suffixes.end() );
   }
   for( const suffix suf: suffixes ) {
       sufArray.push_back( suf.index );
   }
   return sufArray;
int main() {
   #ifdef CLown1331
       freopen( "in.txt", "r", stdin );
   #endif /// CLown1331
   string s;
   while( cin >> s ) {
```

}

```
vector < int > sufArray = buildSuffixArray( s );
       for( const int ind: sufArray ) {
           cout << ind << "\n";</pre>
       cerr << "---\n";
    return 0;
}
```

### 4.7 Zalgo

```
int L = 0, R = 0:
for( int i = 1; i < n; i++ ) {</pre>
       if ( i > R ) {
              L = R = i;
               while ( R < n \&\& s[R-L] == s[R] ) R++;
              z[i] = R-L; R--;
       } else {
       int k = i-L;
              if (z[k] < R-i+1) z[i] = z[k];
               else {
                      L = i;
                      while ( R < n \&\& s[R-L] == s[R] ) R++;
                      z[i] = R-L; R--;
              }
       }
int maxz = 0. res = 0:
for ( int i = 1; i < n; i++ ) {</pre>
       if ( z[i] == n-i && maxz >= n-i ) { res = n-i; break; }
       maxz = max(maxz, z[i]);
}
```