RTEMS Thread Queue Simulation

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1 Introduction

This document provides simulations of the various thread queue designs used in the RTEMS operating system. These range from simple first-in first-out (FIFO) designs, through priority-based approaches, to algorithms for SMP that support schedulability analysis.

We focus on thread queues for RTEMS Tasks waiting to gain access to a shared resource, via a mutex lock.

The basic building blocks are a simple FIFO queue, and a priority queue that assumes all threads have a unique priority.

The most complex form is a round-robin/FIFO queue whose contents are priority queues, one for each scheduling cluster. This is used to implement the Multiprocessor Resource Sharing Protocol(MrsP) thread queue algorithm.

Key papers are: [BW13], [Bra13], [CBHM15], [GZBW17], [ZGBW17], [Gom19], [ZGW $^+$ 20], [ZCW $^+$ 21]. Most relevent RTEMS documents: [RTEa], [RTEb], and [RTEc].

2 RTEMS Thread Queues

```
module Queues
   (FIFOQ
   , isEmptyFIFOQ, viewFIFOQ, enqueueFIFO, dequeueFIFO
   , Priority, PRIOQ
   , isEmptyPRIOQ, viewPRIOQ, enqueuePRIO, dequeuePRIO
   , Cluster, CLSTRQ
   , isEmptyCLSTRQ, viewCLSTRQ, enqueueCLSTR, dequeueCLSTR
) where
```

Here we define the three types of queues used for thread synchronisation and scheduling.

2.1 FIFO Queues

See [RTEa, §3.5].

We model a FIFO queue as a Haskell list, parameterised by content object type, with enqueue and dequeue operations, and an emptiness check.

```
type FIFOQ obj = [obj]
isEmptyFIFOQ :: FIFOQ obj -> Bool
isEmptyFIFOQ = null

viewFIFOQ :: Show obj => FIFOQ obj -> String
viewFIFOQ = show

enqueueFIFO :: obj -> FIFOQ obj -> FIFOQ obj
enqueueFIFO thing fifoq = fifoq ++ [thing]

dequeueFIFO :: MonadFail m => FIFOQ obj -> m (obj,FIFOQ obj)
dequeueFIFO [] = fail "empty FIFO queue"
dequeueFIFO (thing:restq) = return (thing,restq)
```

We have a variant of a FIFO queue called round-robin (RR).

In this, the dequeue operation immediately performs an enqueue operation with the item just dequeued. Typically the queue is initially setup by enqueuing all desired items, and subsequent operations consists solely of dequeing.

```
dequeueRR :: MonadFail m => FIFOQ obj -> m (obj,FIFOQ obj)
dequeueRR [] = fail "empty RR queue"
dequeueRR (thing:restq) = return (thing,restq++[thing])
```

2.2 Priority Queues

See [RTEa, §3.5].

We model a priority queue as a Haskell list, parameterised by content object type, with enqueue and dequeue operations, and an emptiness check.

```
type Priority = Int
type PRIOQ obj = [(Priority,obj)]
isEmptyPRIOQ :: PRIOQ obj -> Bool
isEmptyPRIOQ = null
viewPRIOQ :: Show obj => PRIOQ obj -> String
viewPRIOQ = show
enqueuePRIO :: obj -> Priority -> PRIOQ obj -> PRIOQ obj
enqueuePRIO thing p [] = [(p,thing)]
enqueuePRIO thing p prioq@(first@(q,_):restq)
  | p < q
               = (p,thing) : prioq
  -- p == q, insert as per FIFO, after those of same priority (c-user 3.5)
  | otherwise = first
                           : enqueuePRIO thing p restq
dequeuePRIO :: MonadFail m => PRIOQ obj -> m (obj,Priority,PRIOQ obj)
dequeuePRIO [] = fail "empty PRIO queue"
dequeuePRIO ((p,thing):restq) = return (thing,p,restq)
```

2.3 Clustered Scheduling Queues (SMP)

See [RTEa, §3.5,§5.4].

For cluster scheduling, each scheduler has its own priority queue, and these queues are themselves placed in a global round-robin queue.

```
type Cluster = Int
type CLSTRQ obj = FIFOQ (Cluster, PRIOQ obj)
isEmptyCLSTRQ :: CLSTRQ obj -> Bool
isEmptyCLSTRQ = all isEmptyPRIOQ . map snd
viewCLSTRQ :: Show obj => CLSTRQ obj -> String
viewCLSTRQ = show
enqueueCLSTR :: obj -> Priority -> Cluster -> CLSTRQ obj -> CLSTRQ obj
enqueueCLSTR thing p c [] = [(c,[(p,thing)])]
enqueueCLSTR thing p c (first@(c',prioq):rest)
  c == c'
              = (c',enqueuePRIO thing p prioq):rest
  | otherwise = first : enqueueCLSTR thing p c rest
dequeueCLSTR :: MonadFail m => CLSTRQ obj -> m (obj,Priority,Cluster,CLSTRQ obj)
dequeueCLSTR [] = fail "empty CLSTR queue"
dequeueCLSTR ((c,prioq):restq)
  = do (thing,p,prioq') <- dequeuePRIO prioq</pre>
       if isEmptyPRIOQ prioq' -- delete empty queues (???)
         then return (thing,p,c,restq)
         else return (thing,p,c,restq ++ [(c,prioq')])
```

3 RTEMS Simulation Runner

```
module Runner
( run
) where
import Queues
```

Here we provide a simple mechanism/language to run simulations

The basic idea is to declare some objects, and then invoke actions upon them.

The language is line based, each line starting with a keyword.

We segregate objects by their type.

```
data Object
```

Uninterpreted arbitrary test objects

```
= Unint
```

FIFO FIFO queue objects

```
| FIFO Object
```

PRIO Priority queue objects

```
| PRIO Object
```

CLSTR Cluster queue objects

```
| ClusterQ Object
```

The queue objects are themselves parameterised with a content object.

We define the state to be a collection of named objects:

4 Program Mainline

5 Test Program

```
main :: IO ()
main = putStrLn "Thread Q Sim Test suite not yet implemented"
```

References

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