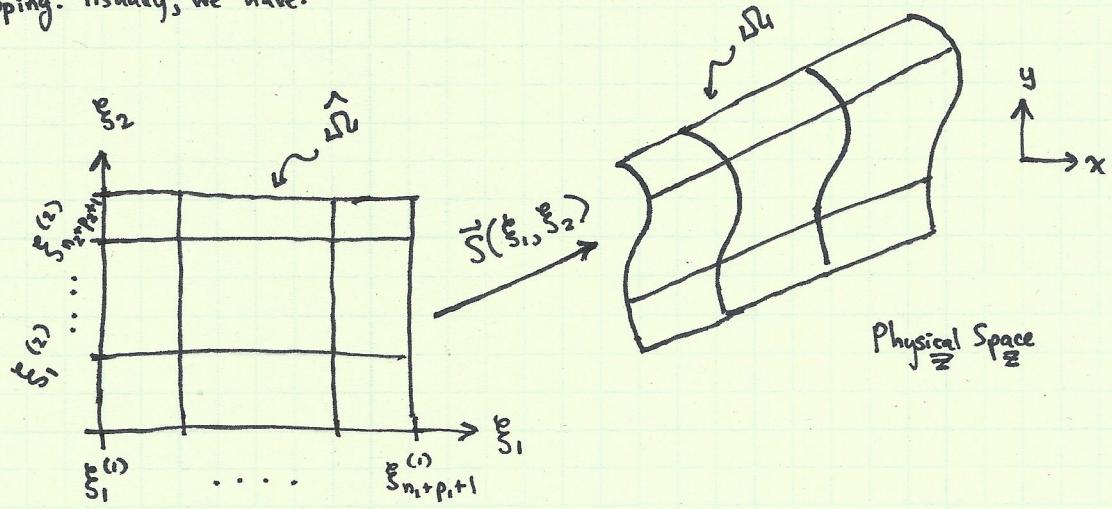


Multi-Patch Geometries:

In almost all practical circumstances, it is practically impossible to represent the geometry of interest with a single NURBS mapping. Thus, it becomes necessary to describe the geometry using multiple NURBS mappings. We refer to such geometries as multi-patch geometries.

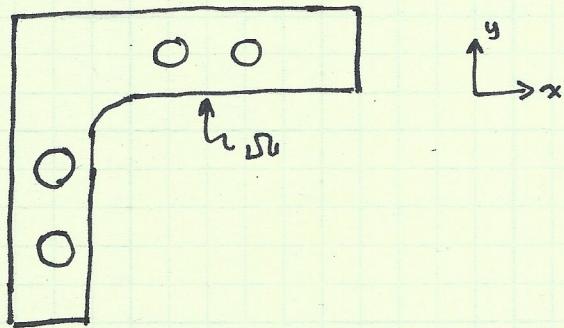
To describe what we mean by "patch", we return to the setting of a single NURBS mapping. Visually, we have:



Parametric Space

The single NURBS mapping $\vec{S}(\xi_1, \xi_2)$ takes the patch from parametric space to physical space. We interchangeably refer to Δu and its image $\hat{\Delta}u$ as the NURBS patch. Note that for a single-patch geometry, the patch Δu is the domain.

Now consider the following domain:

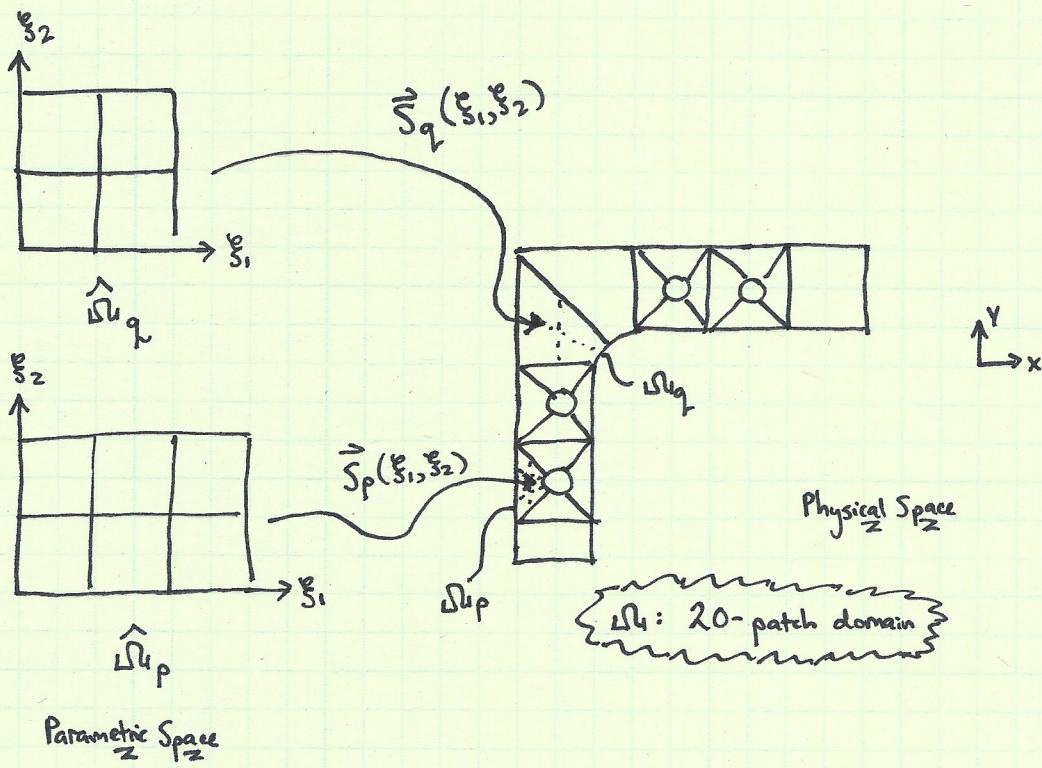


For this configuration, it is impossible to use a single NURBS map to parameterize the geometry. Instead, we must use multiple NURBS patches and mappings:

$$\vec{S}_p(\xi_1, \xi_2) : \hat{\Delta}u_p \rightarrow \Delta u_p \quad p=1, \dots, n \text{ patch}$$

Patch p in Patch p in
parametric space physical space

Visually:



To ensure a watertight geometry, we must ensure the patches $\{S_i\}_{i=1}^{n_{\text{patch}}}$ meet parametrically and geometrically on the internal faces where they meet.

To be more precise, let S_i and S_j be two disjoint patches such that $\partial S_i \cap \partial S_j \neq \emptyset$. Then, we require that S_i and S_j parametrize the shared face $\partial S_i \cap \partial S_j$ exactly up to changes in orientation. This requires that the control nets associated with ∂S_i and ∂S_j are coincident.

identically!

All of the aforementioned Bézier extraction machinery may be easily extended to the multi-patch setting. All we need is a mechanism to associate patch element numbers and basis function numbers to global element numbers and basis function numbers. We use the following two arrays in this regard:

$$\text{PTG element } (p, e_p) = e$$

Patch Number Local Element Number Global Element Number

$$\text{PTG basis } (p, i_p) = i$$

Patch Number Local Basis Function Number Global Basis Function Number

To construct these arrays, we order elements and basis functions patch-by-patch. That is, the elements associated with patch 1 have global element numbers $1, \dots, n_{\text{el}}$, where

n_{elp} is the number of elements associated with patch 1, the elements associated with patch 2 have global element numbers $n_{el_1}+1, \dots, n_{el_1}+n_{el_2}$, and so forth. As control nets associated with shared faces must be coincident, we also mate basis functions at shared interfaces. This requires a master-slave relationship between such basis functions. If \mathcal{W}_{li} and \mathcal{W}_{lj} are two disjoint patches with $\partial\mathcal{W}_i \cap \partial\mathcal{W}_j \neq \emptyset$, we enforce that \mathcal{W}_{li} is slave to \mathcal{W}_{lj} if $i > j$ and vice versa.

With the above conventions defined, we are ready to present a function which, given the number of elements n_{elp} and number of basis functions n_p associated with each patch, returns the PTG element and PTG basis arrays:

```

function Patch Connectivity ( { $n_{elp}$ }^{n_{patch}}, { $n_p$ }^{n_{patch}} )
begin
    Set itr = 0

    for p = 1, ..., n_patch
        for e_p = 1, ..., n_{elp}
            set PTG element (p, e_p) = e_p + itr ↳ Construct PTG element
        endloop
        Set itr = itr + n_{elp}
    endloop

    Set itr = 0

    for p = 1, ..., n_patch
        for i_p = 1, ..., n_p
            Set flag = 0
            Set q = 1
            while (q < p)
                for i_q = 1, ..., n_q
                    if  $\vec{P}_{i_p, p} = \vec{P}_{i_q, q}$  ↳ Control Point i_p of Patch p ↳ Control Point i_q of Patch q
                        Set PTG basis(p, i_p) = PTG basis(q, i_q)
                        Set flag = 1
                        break out of while loop
                endif
            endloop
            if flag = 0
                Set PTG basis(p, i_p) = itr + 1
                Set itr = itr + 1
            endif
        endloop
    endloop

    return PTG element and PTG basis
end

```

Using the PTG element and PTG basis arrays, we can now easily compute the element extraction operators and IEN array for the global domain from the patch-wise element extraction operators and IEN arrays. We restrict our attention to the 2D setting.

```
function Extract And Localize ( {n1,p}p=1npatch, {n2,p}p=1npatch, {P1,p}p=1npatch, {P2,p}p=1npatch,
```

begin

for p = 1, ..., n_{patch}

Call Extraction 2D (n_{1,p}, n_{2,p}, P_{1,p}, P_{2,p}, E_{1,p}, E_{2,p}) to obtain:

{n_{elp}} : The number of elements of patch p

{C_{e,p}}_{e=1}^{n_{elp}} : The element extraction operators for patch p

Call IEN 2D (n_{1,p}, n_{2,p}, P_{1,p}, P_{2,p}, E_{1,p}, E_{2,p}) to obtain the IEN array for patch p, which is denoted here as IEN_p

Set n_p as the maximum entry of IEN_p

endloop

Call Patch Connectivity ({n_{elp}}_{p=1}^{n_{patch}}, {n_p}_{p=1}^{n_{patch}}) to obtain

{PTG element}
{PTG basis}

for p = 1, ..., n_{patch}

for e_p = 1, ..., n_{elp}

Set e = PTG element (p, e_p)

Set C_e = C_{e,p}

Extraction Operators

for a = 1, ..., n_{loc,p} = (P_{1,p}+1)(P_{2,p}+1)

Set i_p = IEN_p(a, e_p)

Set i = PTG basis (p, i_p)

Set IEN(a, e) = i

IEN Array

endloop

endloop

endloop

return n_{el} = $\sum_{p=1}^{n_{patch}} n_{elp}$, {C_e}_{e=1}^{n_{el}} and IEN

end

- It should be noted that even in cases where a single-patch geometry may be employed, doing so might introduce extreme mesh distortion which may adversely affect finite element analysis. The picture below demonstrates such a situation. Alternatively, a multi-patch geometry may introduce far less distortion, leading to improved finite element results.

