

Lecture 2: End of Syntactic Analysis, start of Program Semantics

17-355/17-665/17-819: Program Analysis Claire Le Goues and Fraser Brown Jan 19, 2023

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Administrivia

- HW1 out; due next week.
 - Lots of references online, many linked in assignment.
 - Recitation will have some practice problems
 - Submit via Canvas/Gradescope. Copy/paste your queries + screenshots of analysis results.
- Office hours are up on website.
- Make sure you're on Piazza.



Learning objectives

- Describe the function of an AST and outline the principles behind AST walkers and declarative languages for simple bug-finding analyses.
- Recognize the basic WHILE demonstration language and translate between WHILE and While3Addr.
- Define the meaning of programs using operational semantics
- Read and write inference rules and derivation trees
- Use big- and small-step semantics to show how WHILE programs evaluate
- Use structural induction to prove things about program semantics



Remember this code?

```
void copy_bytes(char dest[], char source[], int n) {
   for (int i = 0; i < n; ++i)
      dest[i] = source[i];
}</pre>
```



Here's the AST Clang makes:

```
-FunctionDecl 0x13900a3d8 <foo.c:1:1, line:4:1> line:1:6 copy_bytes 'void (char *, char *, int)'
  -ParmVarDecl 0x13900a1b0 <col:17, col:27> col:22 used dest 'char *':'char *'
  -ParmVarDecl 0x13900a238 <col:30, col:42> col:35 used source 'char *':'char *'
  -ParmVarDecl 0x13900a2b8 <col:45, col:49 col:49 used n 'int'
  -CompoundStmt 0x13900a7e8 <col:52, line:4:1>
   `-ForStmt 0x13900a7b0 <line:2:3, line:3:23>
     -DeclStmt 0x13900a578 e:2:8, col:17>
       `-VarDecl 0x13900a4f0 <col:8, col:16> col:12 used i 'int' cinit
         `-IntegerLiteral 0x13900a558 <col:16> 'int' 0
     1-<<<NULL>>>
     |-BinaryOperator 0x13900a600 <col:19, col:23> 'int' '<'
      |-ImplicitCastExpr 0x13900a5d0 <col:19> 'int' <LValueToRValue>
        `-DeclRefExpr 0x13900a590 <col:19> 'int' lvalue Var 0x13900a4f0 'i' 'int'
       `-ImplicitCastExpr 0x13900a5e8 <col:23> 'int' <LValueToRValue>
         `-DeclRefExpr 0x13900a5b0 <col:23> 'int' lvalue ParmVar 0x13900a2b8 'n' 'int'
     |-UnaryOperator 0x13900a640 <col:26, col:28> 'int' prefix '++'
       `-DeclRefExpr 0x13900a620 <col:28> 'int' lvalue Var 0x13900a4f0 'i' 'int'
      -BinaryOperator 0x13900a790 <line:3:5, col:23> 'char' '='
       |-ArraySubscriptExpr 0x13900a6c8 <col:5, col:11> 'char' lvalue
        |-ImplicitCastExpr 0x13900a698 <col:5> 'char *':'char *' <LValueToRValue>
           `-DeclRefExpr 0x13900a658 <col:5> 'char *':'char *' lvalue ParmVar 0x13900a1b0 'dest' 'char *':'char *'
         `-ImplicitCastExpr 0x13900a6b0 <col:10> 'int' <LValueToRValue>
           `-DeclRefExpr 0x13900a678 <col:10> 'int' lvalue Var 0x13900a4f0 'i' 'int'
        -ImplicitCastExpr 0x13900a778 <col:15, col:23> 'char' <LValueToRValue>
         `-ArraySubscriptExpr 0x13900a758 <col:15, col:23> 'char' lvalue
           |-ImplicitCastExpr 0x13900a728 <col:15> 'char *':'char *' <LValueToRValue>
            `-DeclRefExpr 0x13900a6e8 <col:15> 'char *':'char *' lvalue ParmVar 0x13900a238 'source' 'char *':'char *'
           `-ImplicitCastExpr 0x13900a740 <col:22> 'int' <LValueToRValue>
             `-DeclRefExpr 0x13900a708 <col:22> 'int' lvalue Var 0x13900a4f0 'i' 'int'
```

Our first static analysis: AST walking

- One way to find "bugs" is to walk the AST, looking for particular patterns.
 - o Traverse the AST, look for nodes of a particular type
 - Check the neighborhood of the node for the pattern in question.
- Various frameworks, some more language-specific than others.
 - o Tradeoffs between language agnosticism and semantic information available.
 - Consider "grep": very language agnostic, not very smart.
 - Python's "astor" package designed for Python ASTs. Clean API; highly specific.
- Classic architecture based on Visitor pattern:
 - o class Visitor has a visitX method for each type of AST node X
 - Default Visitor code just descends the AST, visiting each node
 - To do something interesting for AST element of type X, override visitX
- More recent approaches based on semantic search, declarative logic programming, or query languages.



CodeQL

- A language for querying code. Developed by GitHub.
- Supports many common languages.
- Library of common programming patterns and optimizations.

Inefficient empty string test

```
ID: java/inefficient-empty-string-test
Kind: problem
Severity: recommendation
Precision: high
Tags:
- efficiency
- maintainability
Query suites:
- java-security-and-quality.qls
```

Click to see the query in the CodeQL repository

When checking whether a string s is empty, perhaps the most obvious solution is to write something like s.equals("") (or "".equals(s)). However, this actually carries a fairly significant overhead, because String.equals performs a number of type tests and conversions before starting to compare the content of the strings.

Recommendation

The preferred way of checking whether a string s is empty is to check if its length is equal to zero. Thus, the condition is s.length() == 0. The length method is implemented as a simple field access, and so should be noticeably faster than calling equals.

Note that in Java 6 and later, the String class has an isEmpty method that checks whether a string is empty. If the codebase does not need to support Java 5, it may be better to use that method instead.

Example

In the following example, class InefficientDBClient uses equals to test whether the strings user and pw are empty. Note that the test "".equals(pw) quards against NullPointerException, but the test user.equals("") throws a NullPointerException if user is null.

In contrast, the class EfficientDBClient uses length instead of equals. The class preserves the behavior of InefficientDBClient by guarding pw.length() == 0 with an explicit test for null. Whether or not this guard is desirable depends on the intended behavior of the program.



Example: Java string compare with ""

```
// Inefficient version
     class InefficientDBClient {
         public void connect(String user, String pw) {
             if (user.equals("") || "".equals(pw))
                 throw new RuntimeException();
 8
     // More efficient version
     class EfficientDBClient {
11
                                                                           Hint: doub
         public void connect(String user, String pw) {
12
             if (user.length() == 0 || (pw != null && pw.length() == 0))
13
                 throw new RuntimeException();
14
15
             . . .
16
17
```



CodeQL query for empty string comparison

```
Query: InefficientEmptyStringTest.ql

    Collapse source

/**
 * @name Inefficient empty string test
 * @description Checking a string for equality with an empty string is inefficient.
 * @kind problem
 * @problem.severity recommendation
 * @precision high
 * @id java/inefficient-empty-string-test
 * @tags efficiency
         maintainability
 */
import java
from MethodAccess mc
where
  mc.getQualifier().getType() instanceof TypeString and
  mc.getMethod().hasName("equals") and
    mc.getArgument(0).(StringLiteral).getRepresentedString() = "" or
    mc.getQualifier().(StringLiteral).getRepresentedString() = ""
select mc, "Inefficient comparison to empty string, check for zero length instead."
```



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Together: String concatenation in a loop

- Pseudocode for a simple syntactic analysis that warns when string concatenation occurs in a loop
 - Why? In Java and .NET it may be more efficient to use a StringBuffer

```
class StringConcatLoopAnalysis extends Visitor {
  private int loopLevel = 0;

  void visitStringConcat(StringConcat e) {
    if (loopLevel > 0)
       warn("Use StringBuffer instead)")
    super.visitStringConcat(e); // visits children
  }

  void visitWhile(While e) {
    loopLevel++;
    super.visitWhile(e); // visits children
    loopLevel--;
  }

// similar for other looping constructs
```

