

Clock Synchronization for Interactive Music Systems

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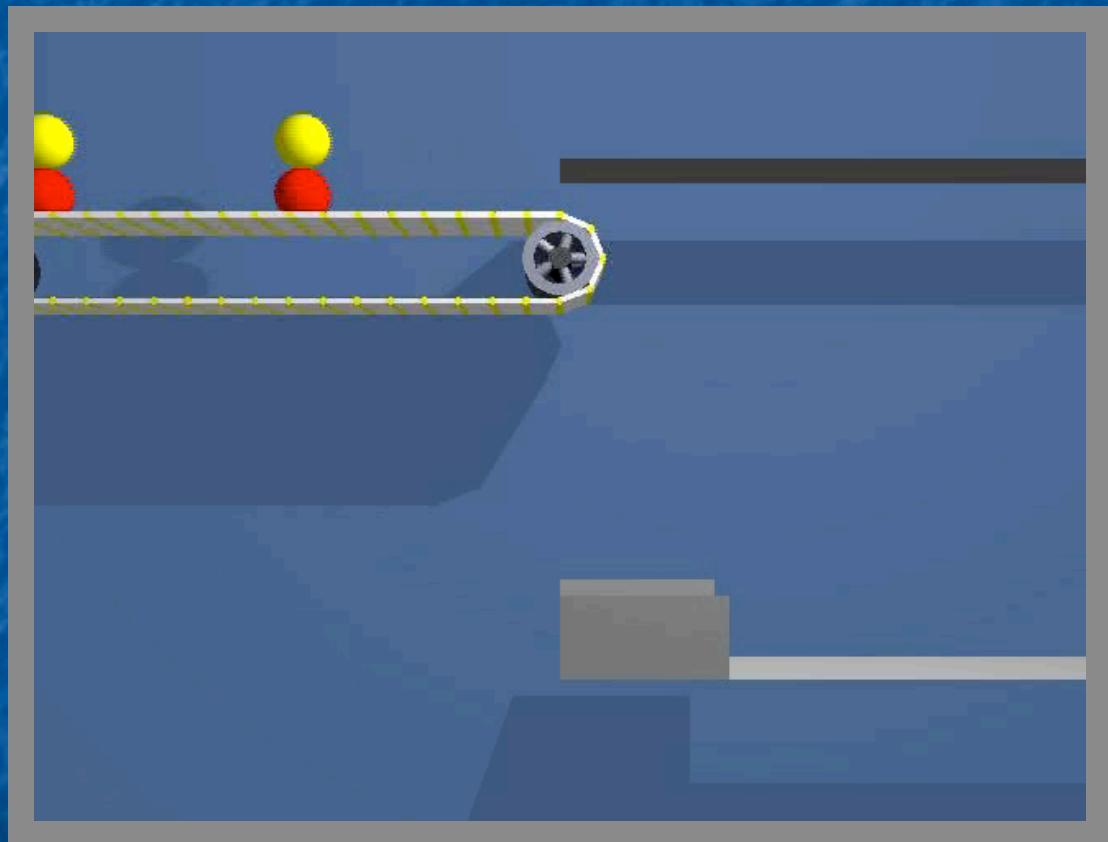
Overview

- Why clock synchronization?
- Characterize the problem
- Simple solution
- Some more elaborate approaches
- What next?

Why Clock Synchronization?

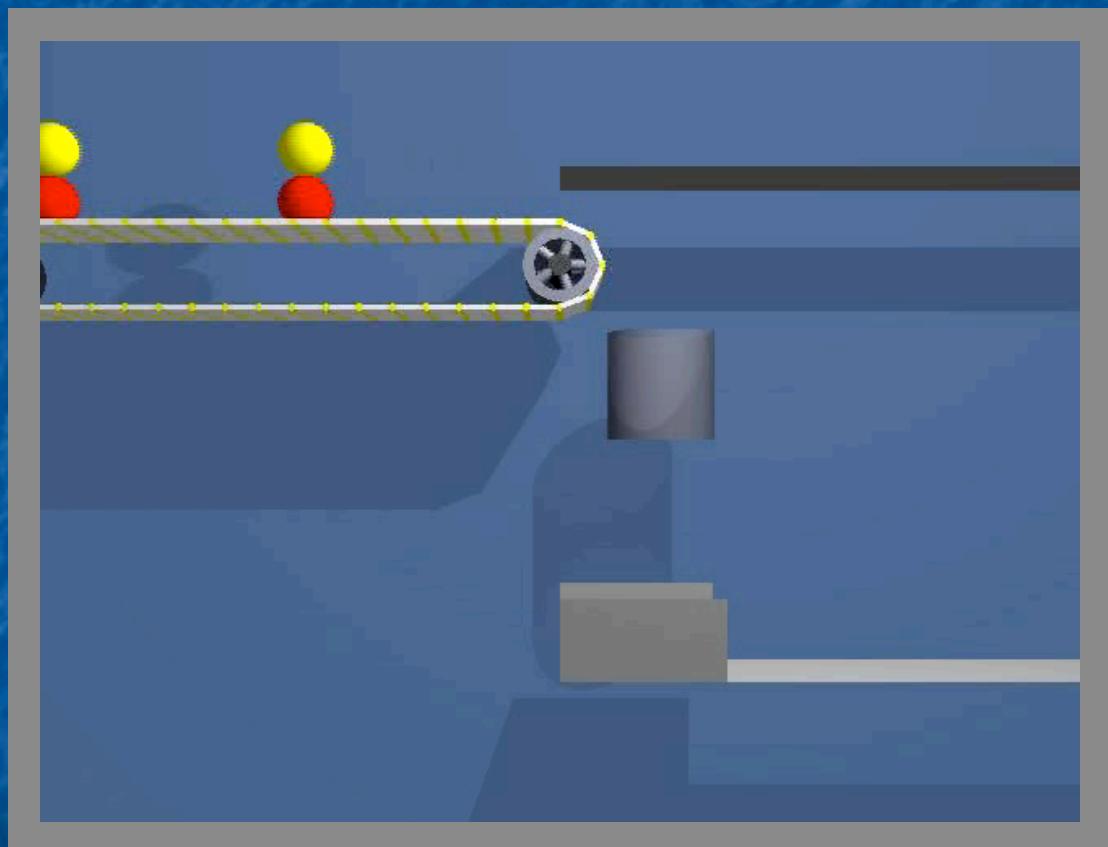
If you have
low-latency
communication,
you do *not*
need clock
synchronization

...



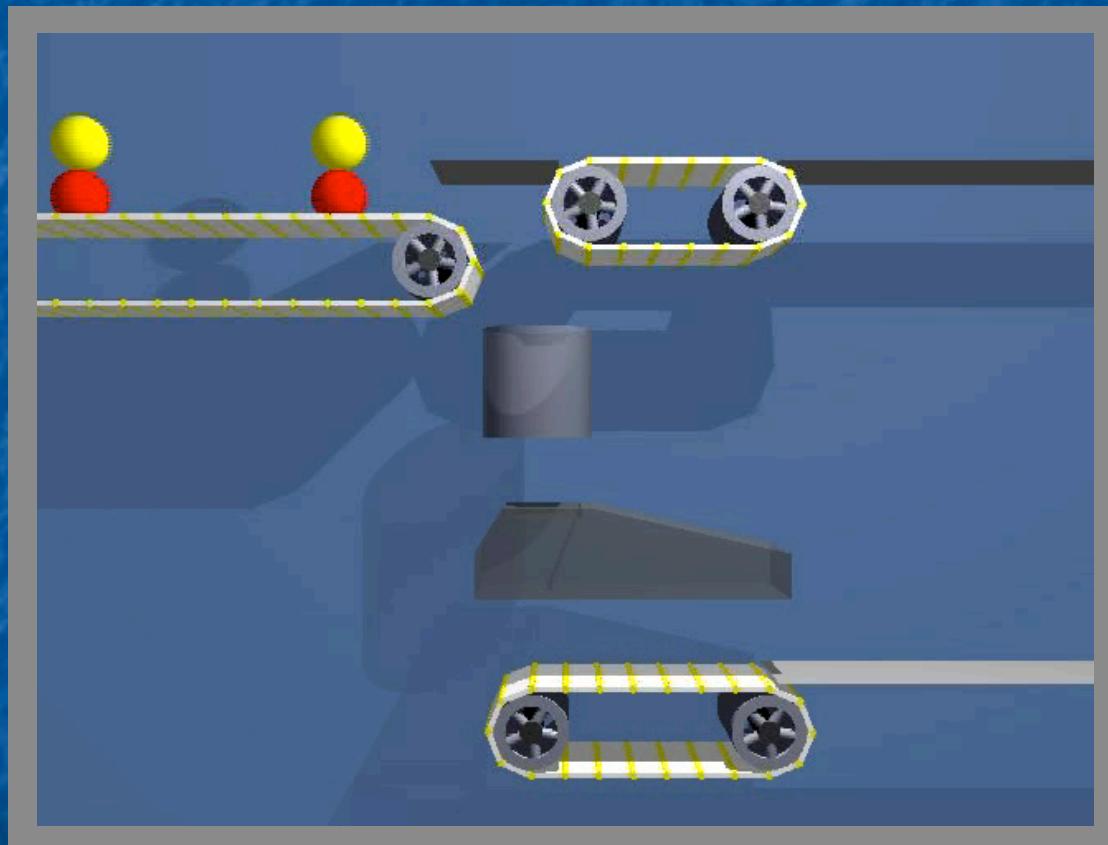
Why Clock Synchronization? (2)

If network communication *sometimes* has high delays (latency), then event synchronization is difficult...



Why Clock Synchronization? (3)

Scheduling according to timestamps can overcome some synchronization problems (but not latency problems)...



Why Clock Synchronization? (4)

- Timestamps are only as good as the local clock...
- ...therefore the goal is:
Synchronize clocks to a precision that is much better than network latency and jitter.

The Design Space

- What do we synchronize to?
 - Global consensus (internal synchronization)
 - Master reference clock (external synch.)
- Who's in charge?
 - No one (symmetric)
 - Master (asymmetric, master-controlled)
 - Slave (asymmetric, slave-controlled)
- Special synchronization hardware?
 - Yes: hardware synchronization
 - No: software synchronization

Clock and Network Characteristics

- Crystal clock accuracy: +/-0.02%
- Frequency drift: low
- Network latency: <1ms
- Network jitter: long tail (0.5s)
- Jitter reading clock or frame #: <1ms
- This should be easy...

Network Latency and Jitter

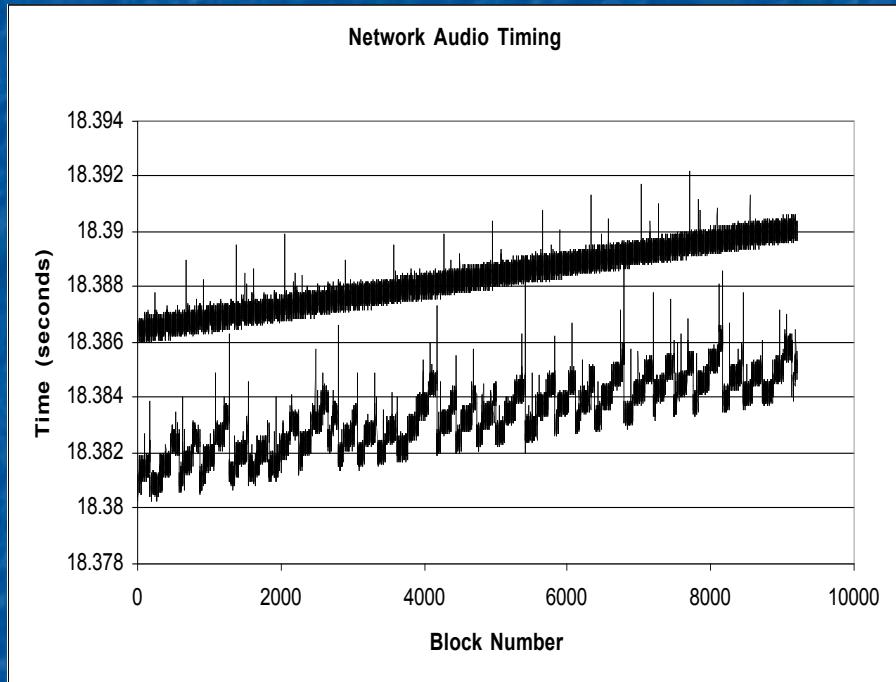
- Interactive music systems
 - not compute bound
 - short or empty network and task queues
 - Messages *usually* get through quickly
- To read remote system time:
 - send message; wait for reply
 - quick reply => low latency and jitter
 - add half of transit time to compensate for latency
 - result should be well below 1ms error

Logical Clock Model

- Assume that time is a linear function of the local clock or sample count:
$$\text{LogicalTime} = \text{offset} + \text{rate} * \text{LocalTime}$$
- Clock synchronization amounts to updating *offset* and *rate*.

Simple Solution

- Periodically read remote “master” clock
- If reply returns quickly, update local time
- Otherwise, continue with previous model until next period.



Audio block arrival time (lower), and block write time (upper).

More Elaborate Approaches

- Dominique Fober:
 - Use window of recent timestamp messages
 - Reject outliers, estimate offset and rate
 - Use exponential smoothing
- Brandt and Dannenberg:
 - Treat logical clock as feedback control system
 - In simulation, achieved 1.1ms clock error with 5ms error reading sample clock.

What Next?

- How do you handle dropped frames?
 - If time is measured in frames, time can jump.
 - You could inform the slaves when time jumps.
 - Or slaves could try to guess when time jumps.
 - In general, fast recovery is in conflict with stability and low error.
- How do you deal with unmatched sample rates?
 - Resample?
 - Ignore it and work at control level?

Conclusions

- Clock synchronization is critical for networked interactive systems
 - *Assuming that network latency is significant!*
- Clocks and networks have almost ideal properties.
- Simple approaches work well to $\sim 1\text{ms}$.
- Advanced techniques can achieve near-frame accuracy over ordinary networks.