

A Simple Algorithm for Maximal Poisson-Disk Sampling in High Dimensions

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presenter = Scott



Eurographics 2012

UCDAVIS
UNIVERSITY OF CALIFORNIA

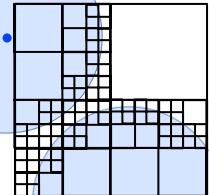
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Overview



- **Classic Dart throwing +**
 - Quadtree
 - Squares track remaining regions
 - Track misses for refinement decisions
 - Avoid refining too deep

[Wei08] Wei L.-Y.: Parallel Poisson disk sampling.
ACM Transactions on Graphics 27, 3 (Aug. 2008), 20:1–20:9.

[BWWM10] Bowers J., Wang R., Wei L.-Y., Maletz D.:
Parallel Poisson disk sampling with spectrum analysis on surfaces.
ACM Transactions on Graphics 29 (Dec. 2010), 166:1– 166:10.

“Make everything as simple as possible, but not simpler.” – A. Einstein

- Flat quadtree – one level of squares active, pool of indices
 - Simpler Datastructure ☺ Less memory ☺
- Globally refine periodically, ignore local misses
 - Simpler Datastructure ☺ More parallel ☺
- Refine to machine precision,
on average it is so rare that memory is not an issue
 - More Maximal ☺

**“This could be the current algorithm of choice for dart throwing.” –
Eurographics reviewer #2**

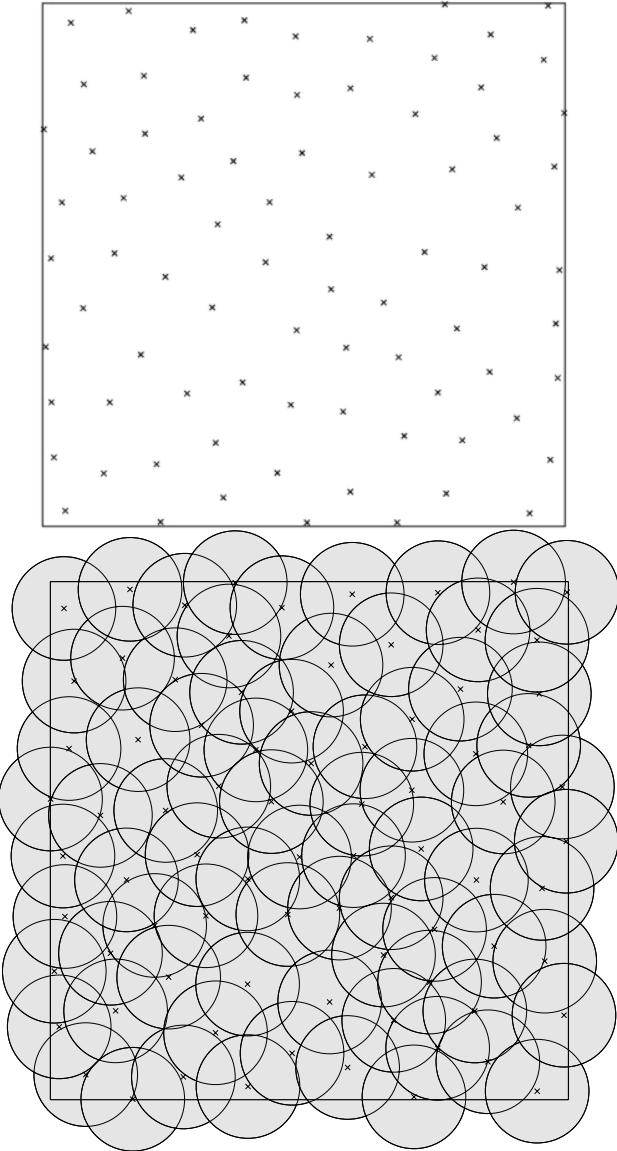
- **Code works great but we can't prove the spatial stats theory.**

Provable:

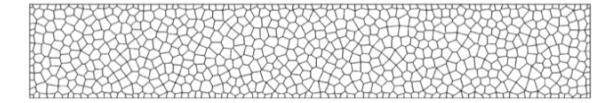
Ebeida M. S., Patney A., Mitchell S. A., Davidson A., Knupp P. M., Owens J. D.:
Efficient maximal Poisson-disk sampling.
ACM Transactions on Graphics 30, 4 (July 2011), 49:1–49:12



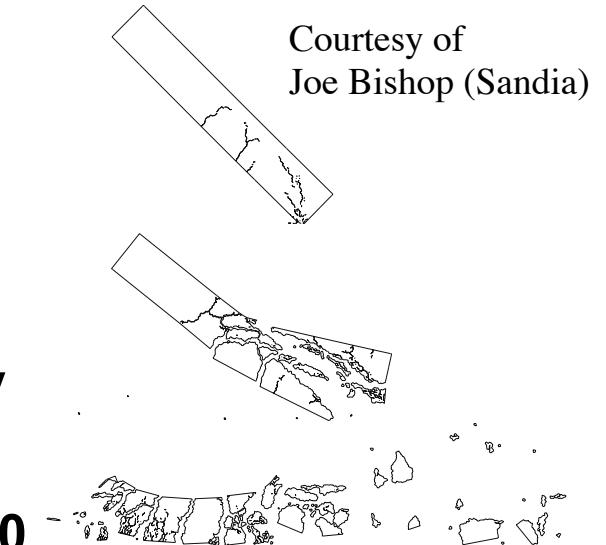
Why MPS? Maximal Poisson-disk Sampling



- **Properties**
 - **Random distribution**
 - Without visible patterns, correlations
 - Blue noise spectrum
 - **Separated-yet-dense**
 - Efficient-yet-quality interpolation
- **Graphics**
 - texture synthesis
- **Mesh generation**
 - Random cracks, quality
- **Design of computer exp.**
 - high dimensions, 10-100



Fracture Simulations





Maximal Poisson-Disk Sampling

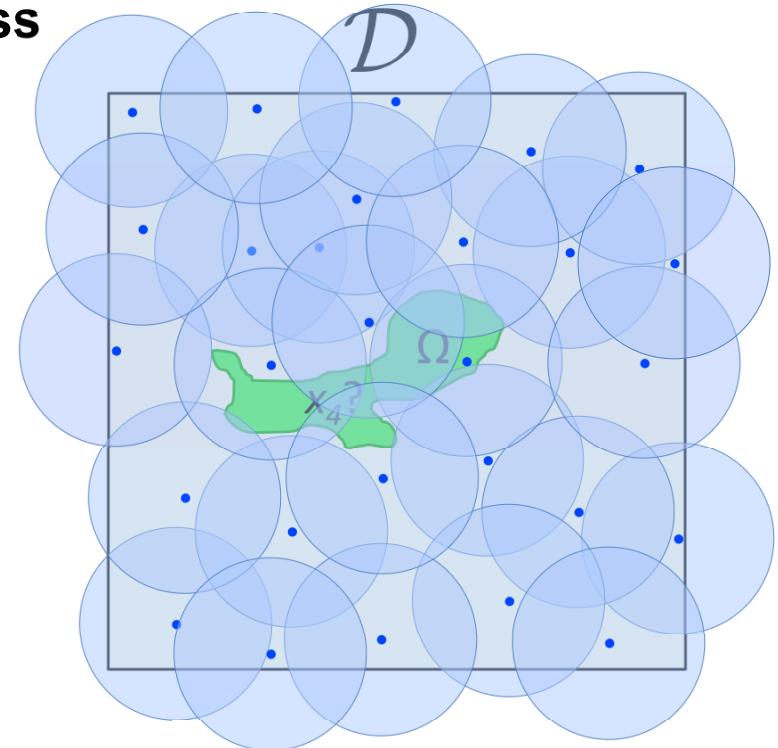
- What is MPS?
 - Dart-throwing
 - Insert random points into a domain, build set X
 - With the “Poisson” process

Empty disk: $\forall x_i, x_j \in X, x_i \neq x_j : \|x_i - x_j\| \geq r$

Bias-free: $\forall x_i \in X, \forall \Omega \subset \mathcal{D}_{i-1} :$

$$P(x_i \in \Omega) = \frac{\text{Area}(\Omega)}{\text{Area}(\mathcal{D}_{i-1})}$$

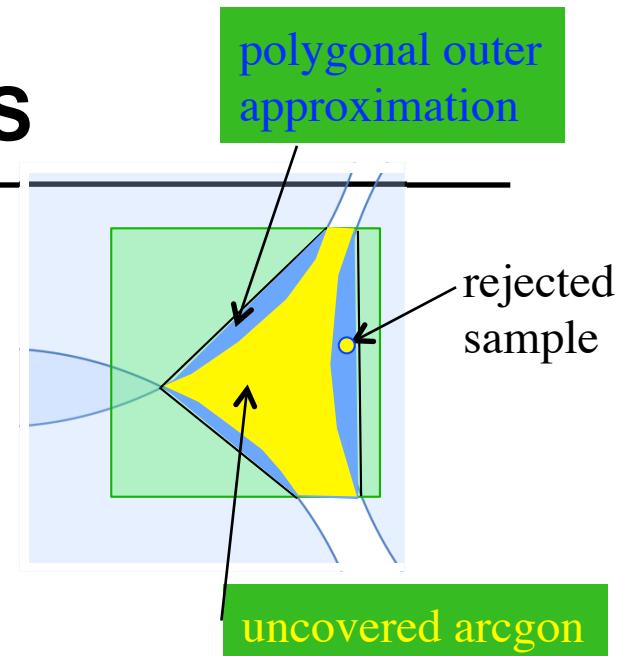
Maximal: $\forall x \in \mathcal{D}, \exists x_i \in X : \|x - x_i\| < r$





Algorithm for MPS

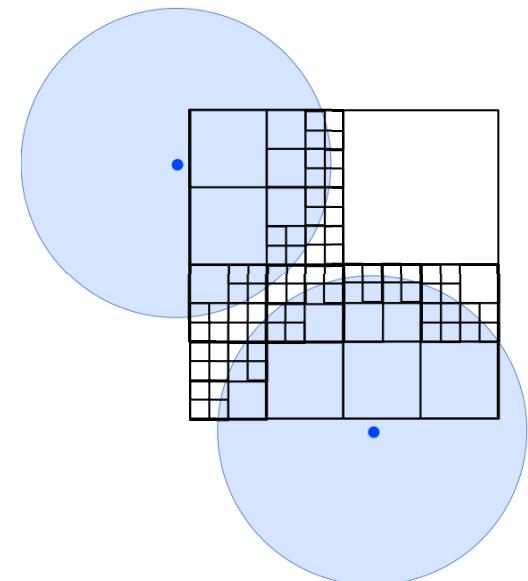
- **Classic algorithm**
 - Throw a point, check if disk overlaps, keep/reject
 - Fast at first, but slows as uncovered area decreases
Can't get maximal.



- **Speedup by targeting just the uncovered area**
 - Polygons Ebeida et al. SIGGRAPH 2011
 - Quadtrees to approximate the uncovered area
 - Discard covered squares
 - Uncovered squares: a sample is always acceptable
 - Partially covered squares: may need to refine

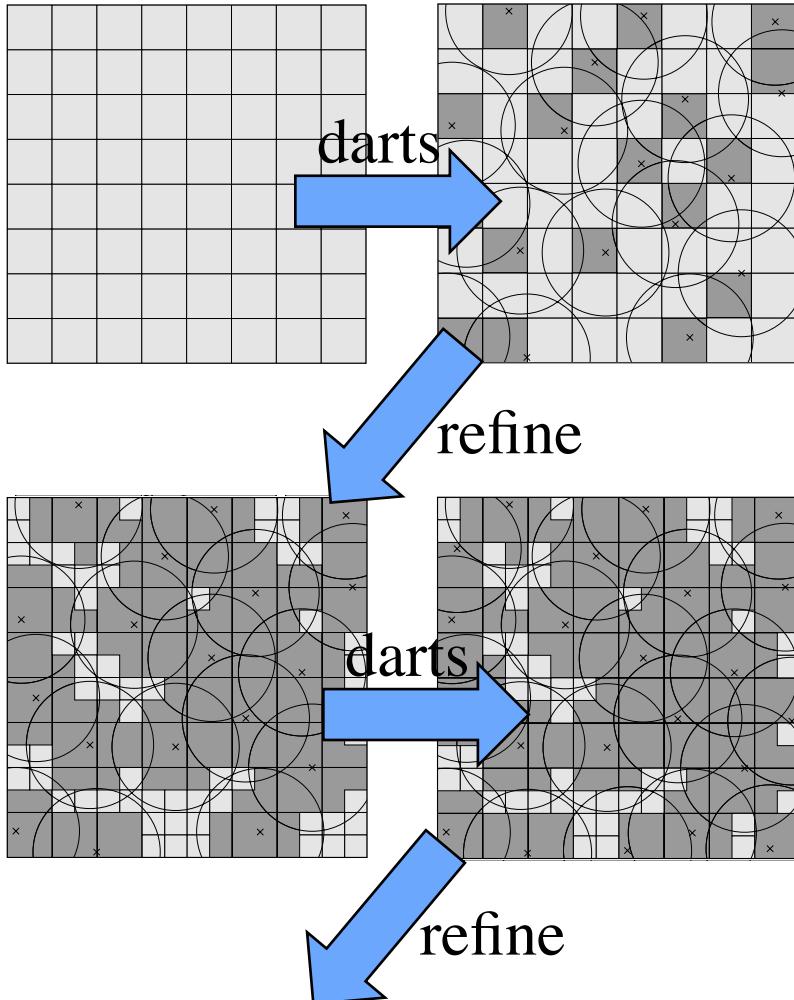
Bias-free: $\forall x_i \in X, \forall \Omega \subset \mathcal{D}_{i-1} :$

$$P(x_i \in \Omega) = \frac{\text{Area}(\Omega)}{\text{Area}(\mathcal{D}_{i-1})}$$





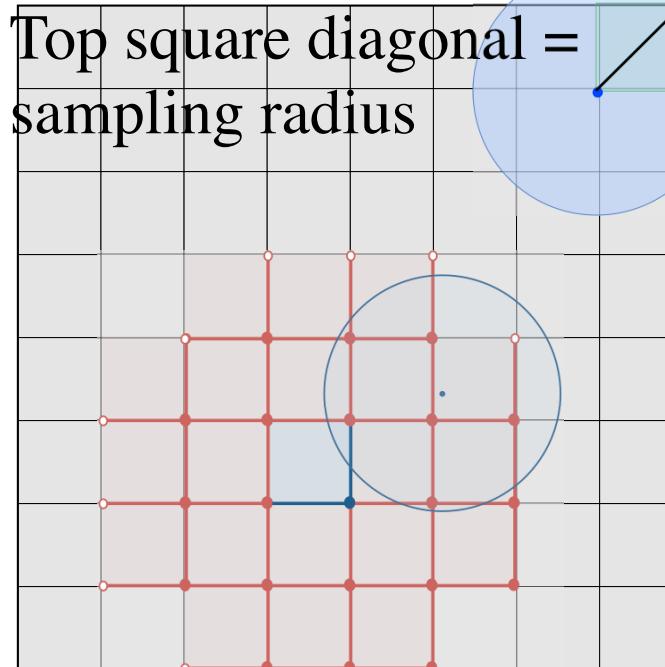
Our Algorithm - Basics



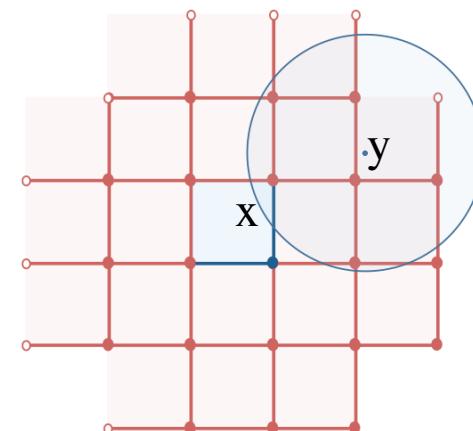
- **Datastructure:**
 - Squares contain uncovered area
- **Throw darts**
 - Pick square, pick point in square
 - If dart is outside nearby circles
 - Accept dart as sample
 - Delete square
- **Refine all squares**
 - Discard subsquares covered by single disks
- **Repeat**



Datastructure: Quadtree Root



- Squares sized so
 - Can fit at most one sample
 - Nearby square template for “Point in disk?” conflict check
 - Pointer from square to its sample

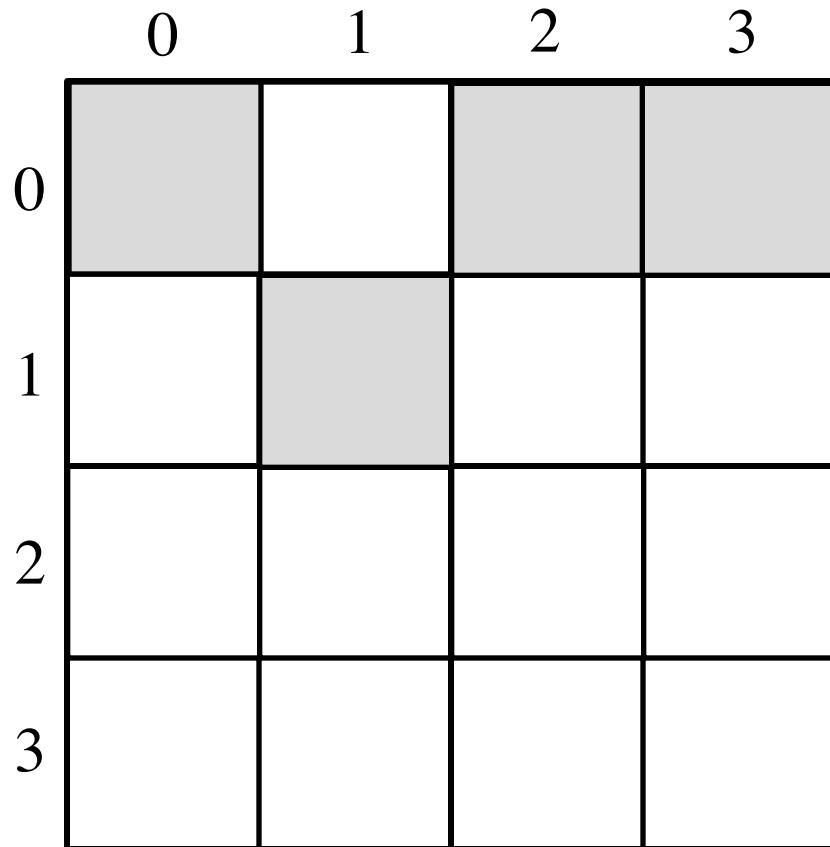


Unpublished extension: use kd-tree for proximity...

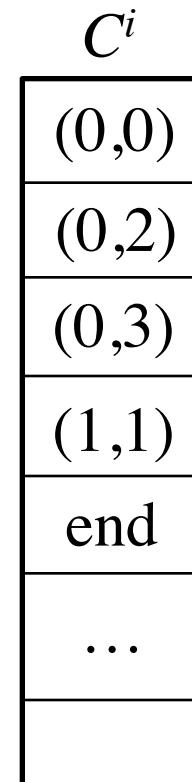


Datastructure: Flat Quadtree Leaves

Flat: Only one level i is used at a time



- Pool of squares
 - Global level i
 - Squares that might accept a sample
 - Array of indices C



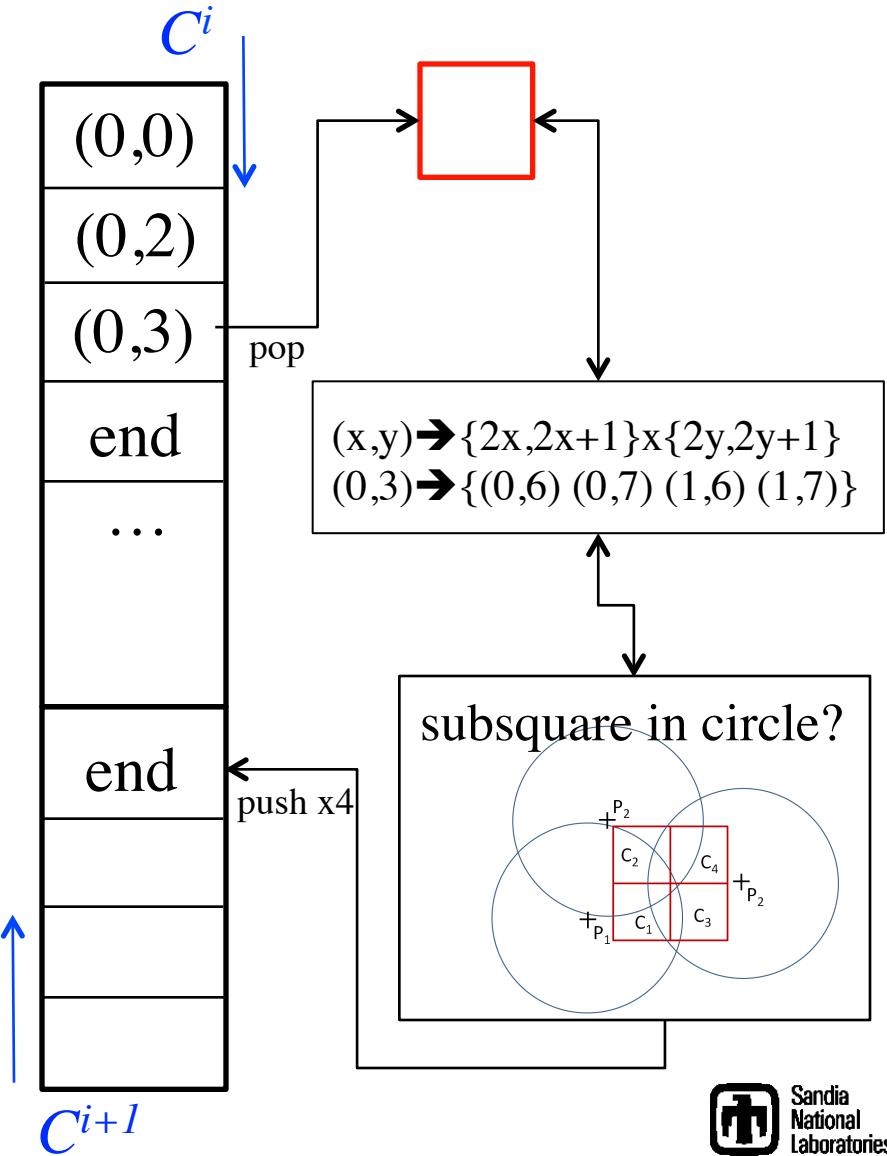
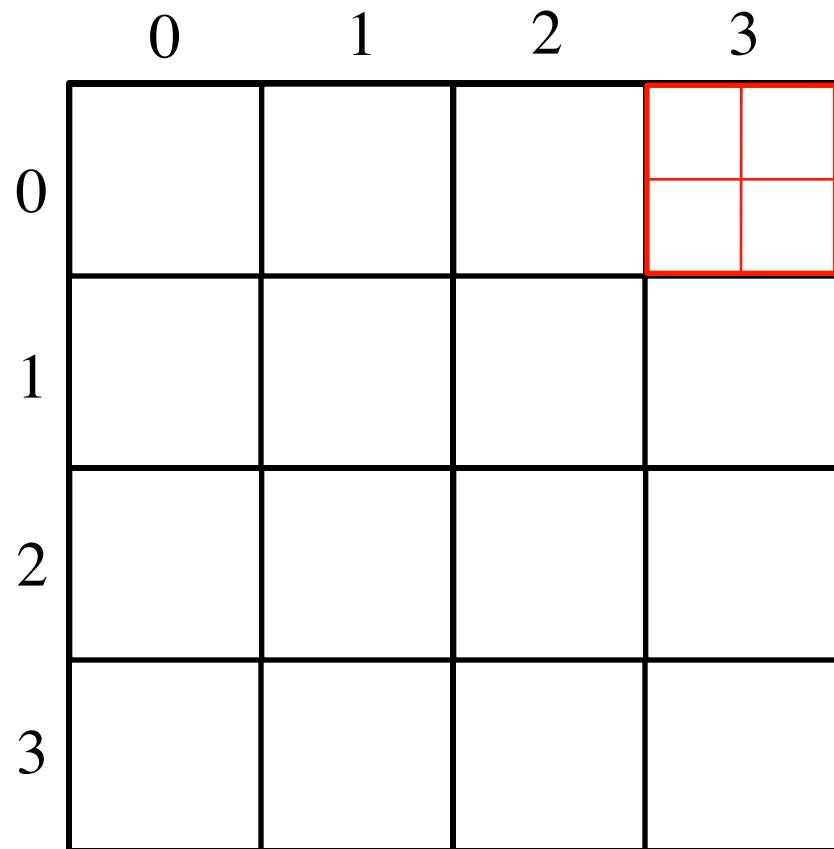
$i=3$
i.e. initial $\times 2^i$
squares per side



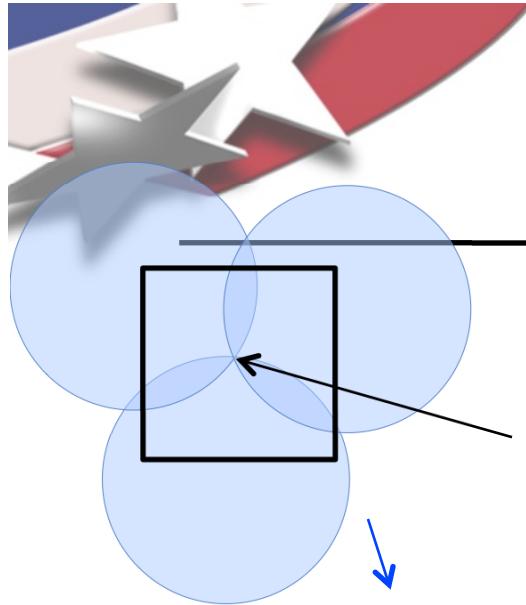
Flat Quadtree Refinement

Update in place.

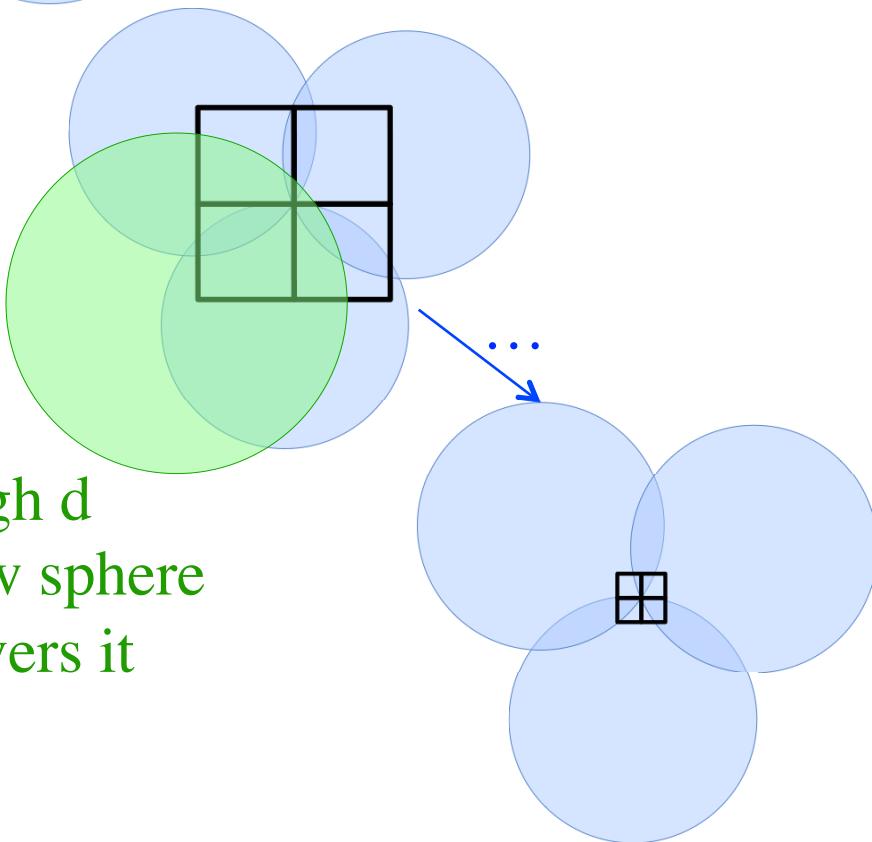
i++



Level Limit?



one uncovered
point
(or is it covered?
let's look closer...)



High d
new sphere
covers it

- **Problem**
 - Test if square in single circle
 - Small voids require infinite refinement
- **Solution: [Wei08], [BWWM10]**
 - Stop early to avoid memory blow-up
- **Solution: Us**
 - Refine to finite-precision
 - **Small voids happen rarely on average so**
 - Memory is fine in practice
 - **Benefit: maximal**



Algorithm – outer loop parameters

Algorithm 1 Simple MPS algorithm, CPU.

```
initialize  $\mathcal{G}^o$ ,  $i = 0$ ,  $\mathcal{C}^i = \mathcal{G}^o$ 
while  $|\mathcal{C}^i| > 0$  do
    {throw darts}
    for all  $A|\mathcal{C}^i|$  (constant) dart throws do
        select an active cell  $\mathcal{C}_c^i$  from  $\mathcal{C}^i$  uniformly at random
        if  $\mathcal{C}_c^i$ 's parent base grid cell  $\mathcal{G}_c^o$  has a sample then
            remove  $\mathcal{C}_c^i$  from  $\mathcal{C}^i$ 
        else
            throw candidate dart  $c$  into  $\mathcal{C}_c^i$ , uniform random
            if  $c$  is disk-free then
                {promote dart to sample}
                add  $c$  to  $\mathcal{G}_c^o$  as an accepted sample  $p$ 
                remove  $\mathcal{C}_c^i$  from  $\mathcal{C}^i$  {additional cells might be
                covered, but these are ignored for now}
            end if
        end if
    end for
    {iterate}
    for all active cells  $\mathcal{C}^i$  do
        if  $i < b$  subdivide  $\mathcal{C}_c^i$  into  $2^d$  subcells
        retain uncovered (sub)cells as  $\mathcal{C}^{i+1}$ 
    end for
    increment  $i$ 
end while
```

Tuning parameter choices: A, B

C^o = number initial cells

C^i = number current squares

How many throws before refining?

$$\text{Throws} = A | C^i |$$

How big does array C need to be
to hold all the refined grid cells?

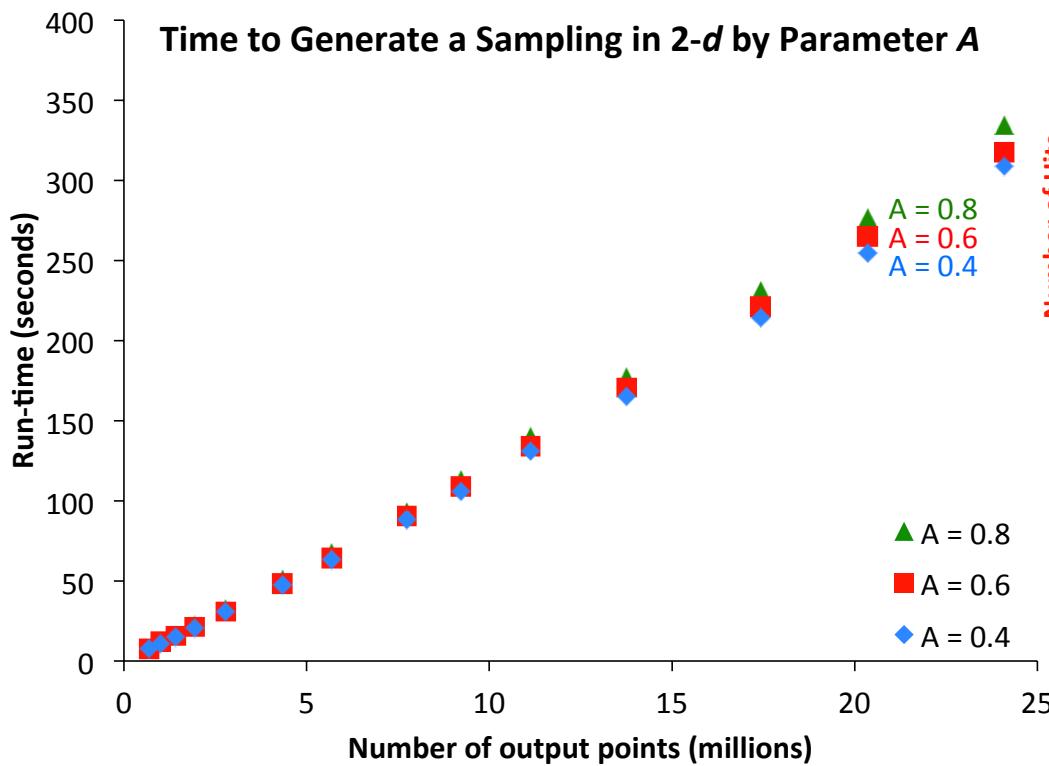
$$C = B | C^o |$$

Big A \leftarrow more time, smaller memory B



A (time) and B (memory) parameters

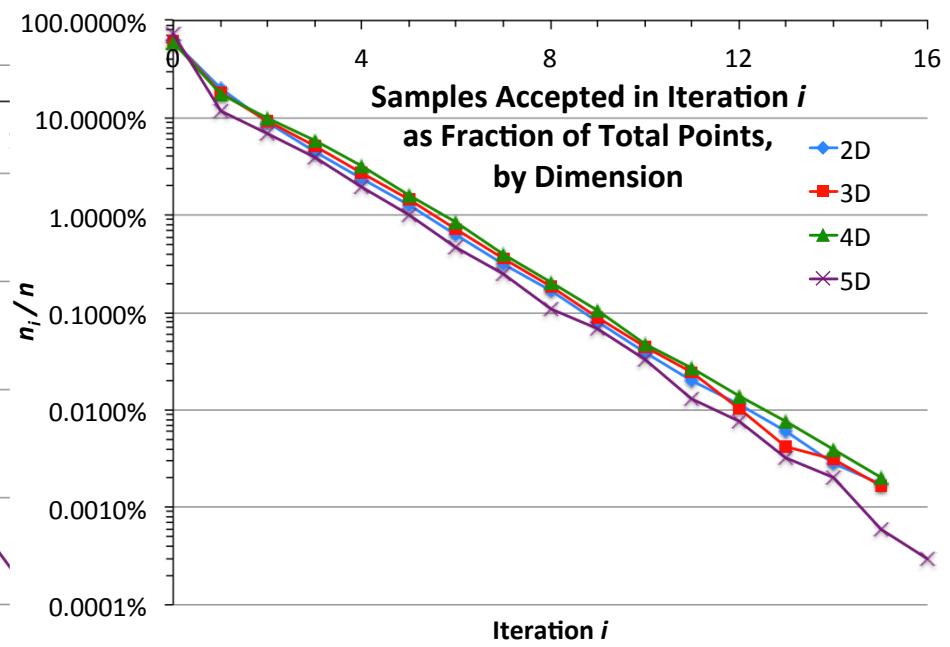
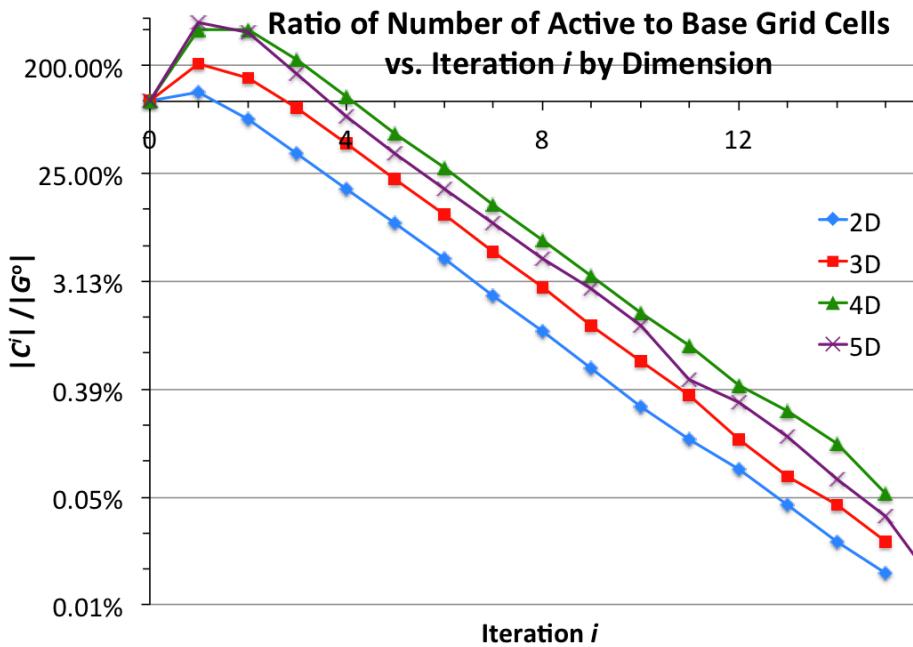
- Big A \leftrightarrow more time, smaller memory B
 - $A \approx 1$, $B \approx \text{dimension}$. (A increases for $d > 4$)
 - Insensitive to value of A above a threshold
 - Intuition: as classical dart throwing,
most hits happen early, no benefit to more throws





Time and Memory Experimental results

- Memory and time peaks in early iterations
 - Exponential convergence thereafter
 - Log y scale



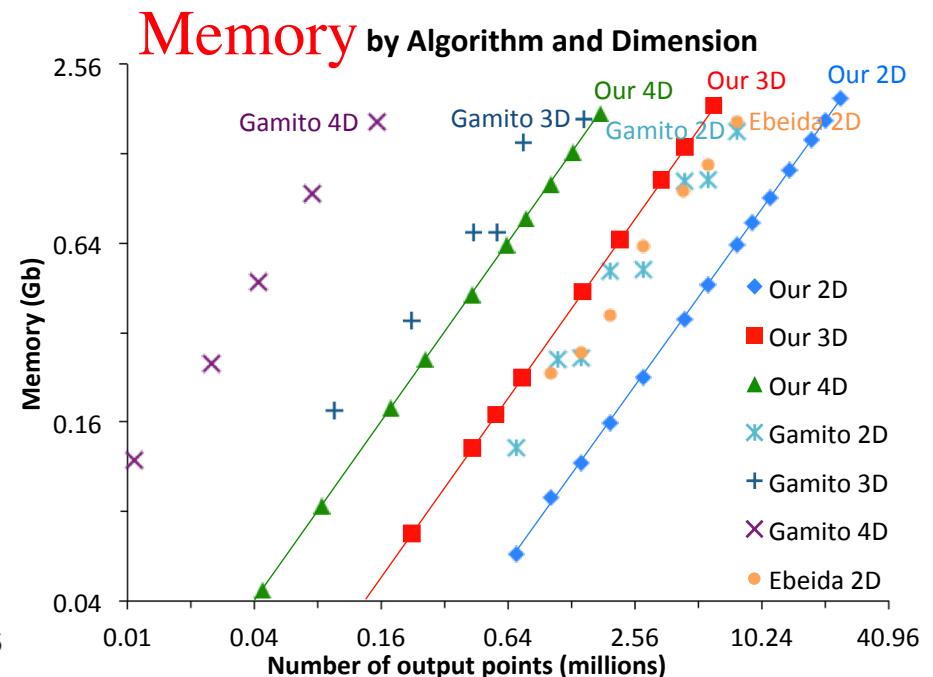
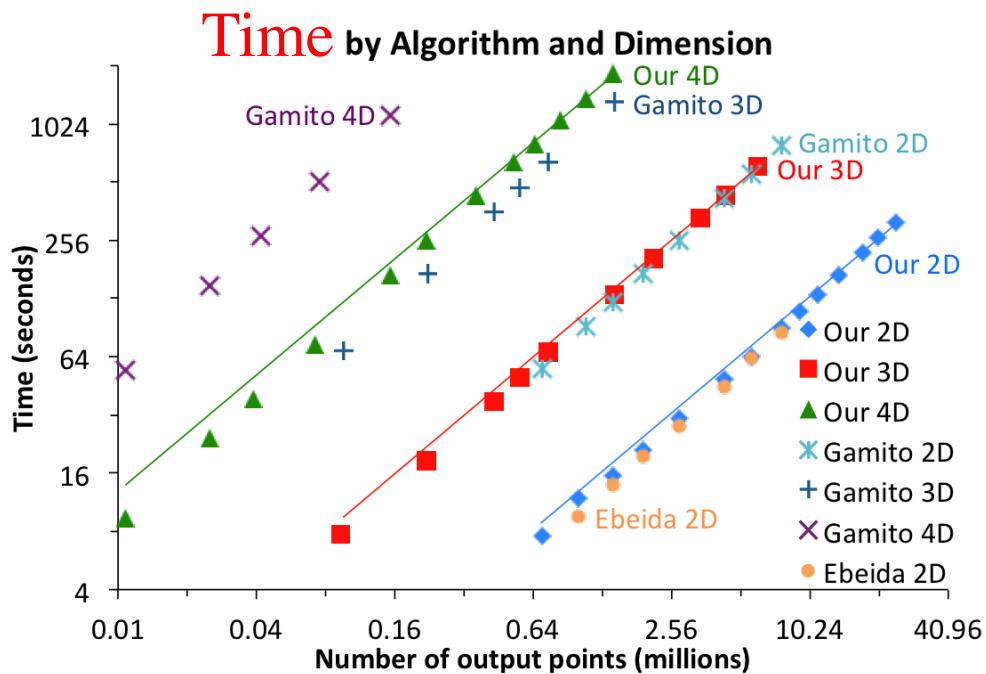
#boxes \approx time, memory,

Time and Memory

vs. true quadtrees (Gamito), polygons (Ebeida 2D)

all linear in both, but constants matter

log-log scales



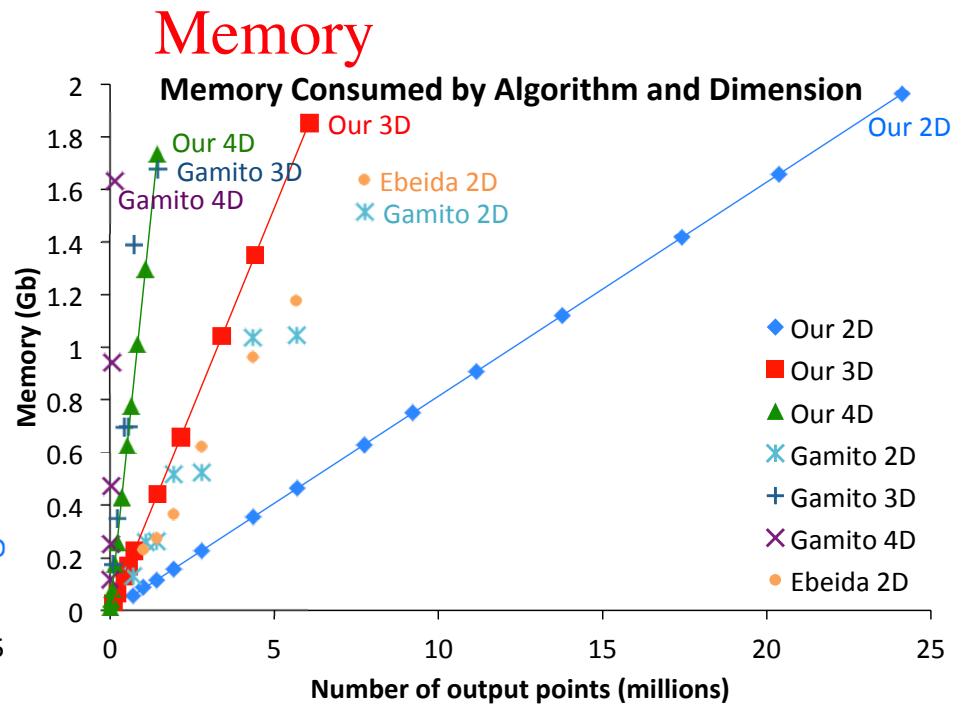
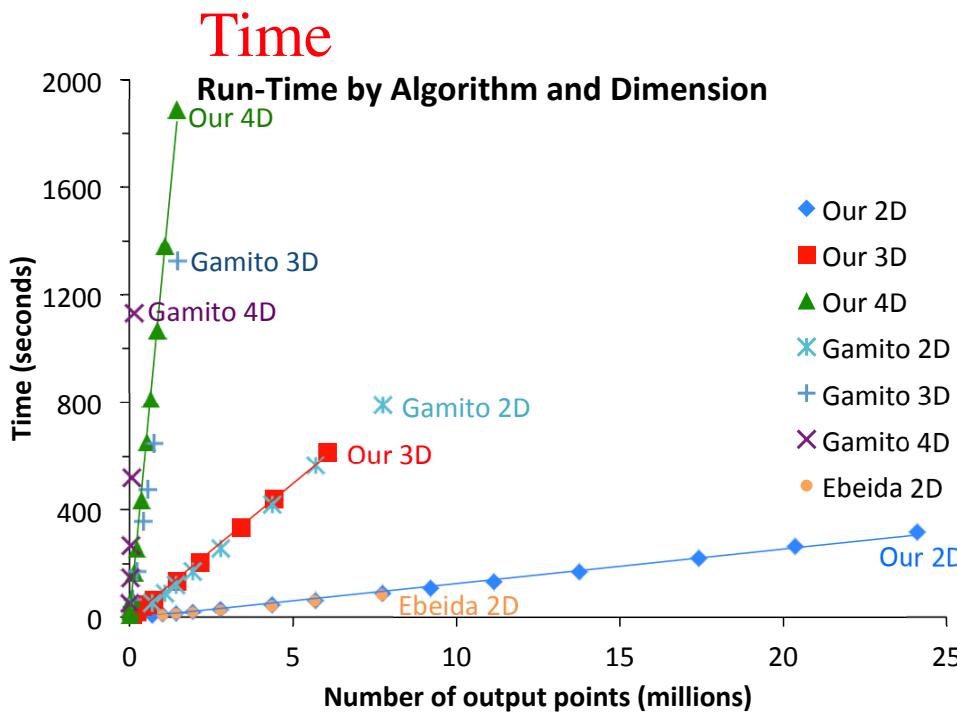
Memory savings from simpler datastructure

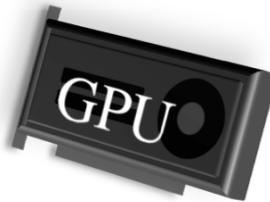
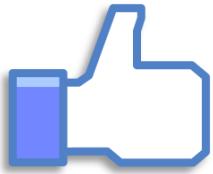
Time savings from that + simpler/fewer checks



Time and Memory Theory

- Run-time
 - Practice: linear in #points, **grows by dimension**
 - Proof: not available
 - Spatial statistics, expected area fraction of cells? And where?
- Memory
 - Linear in #points
 - No dynamic memory allocation



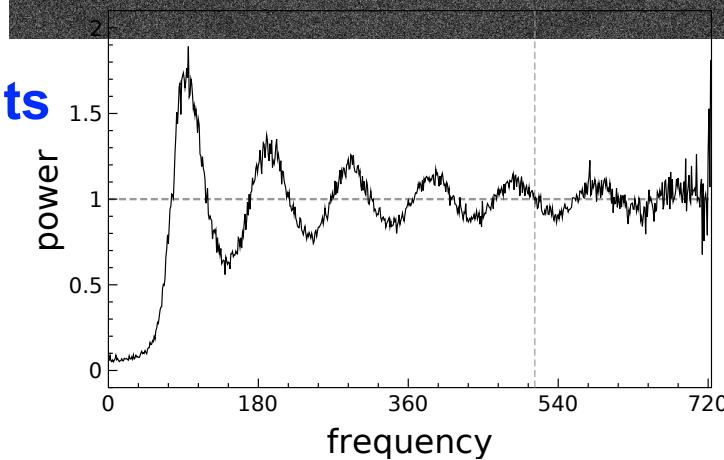
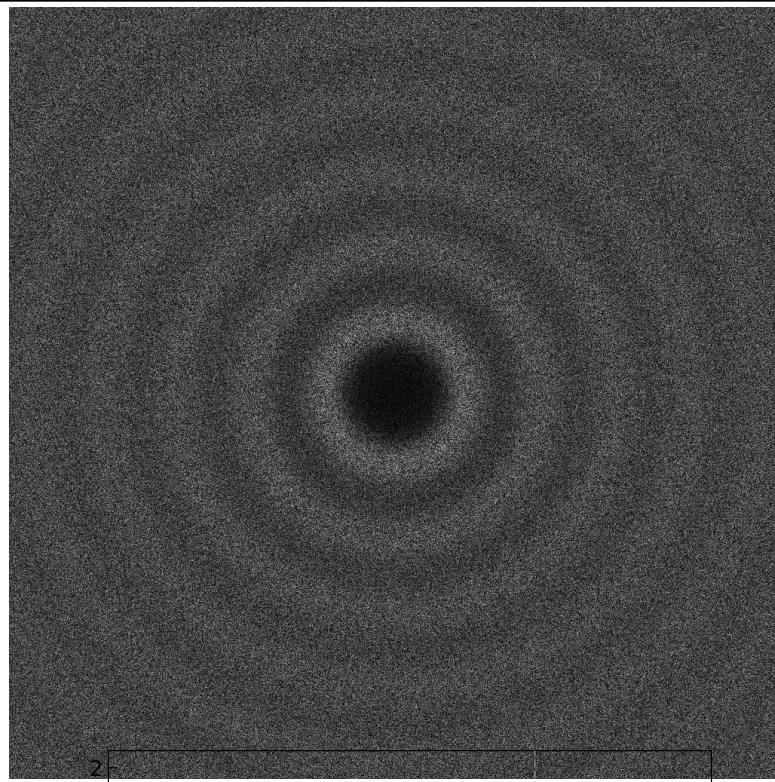
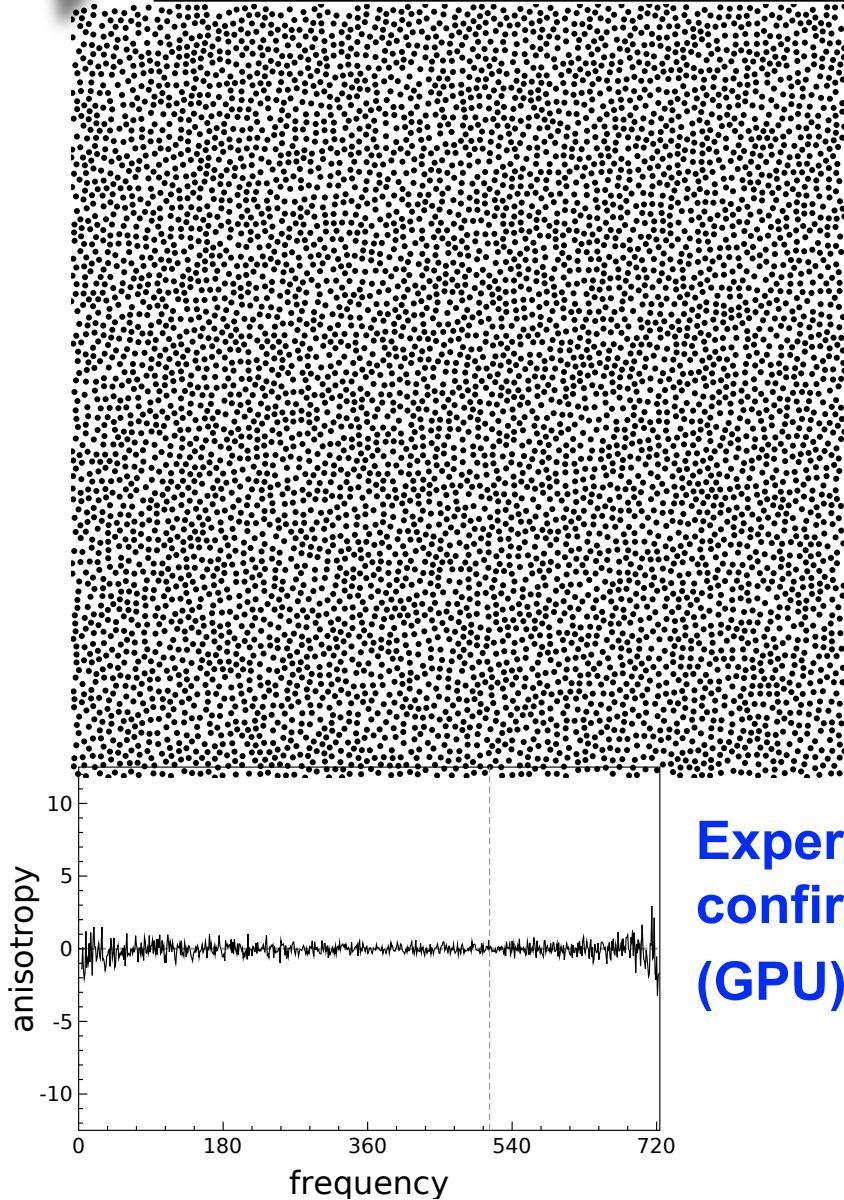


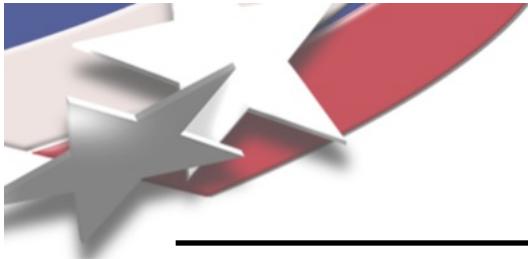
- **Rejection sampling is great on a GPU**
 - Nothing to communicate for a dart miss!
- **10x speedup on NVIDIA GTX 460**
 - Memory-limited to 600k points 2d, 200k in 3d



Point Cloud Quality?

Provably correct bias-free, maximal up to precision



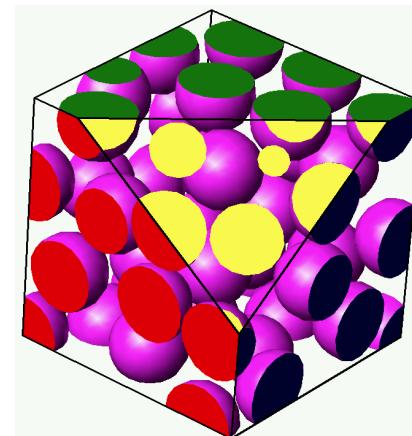
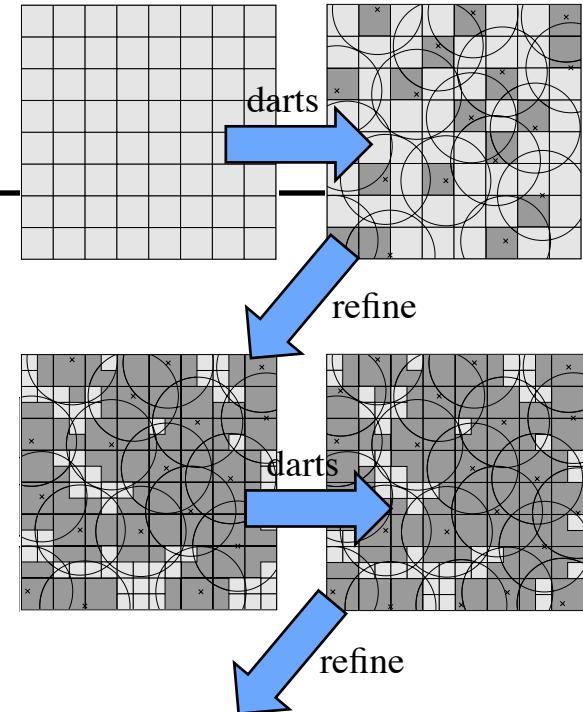


Conclusions

- MPS Maximal Poisson-disk sampling
 - Simpler, faster, less memory
 - Three simple ideas
 - Flat quadtree
 - Constant # throws / ignore misses
 - Global refinement
 - CPU and GPU

Reviewer #0: “The paper is yet another one about faster Poisson-sampling, but I see that it is significantly faster, uses less memory, is just simpler, easier to implement, and works well for higher dimensions.”

- Future, dimensions > 4?
 - Not so great, quadtrees too big
- Two bonus thoughts...





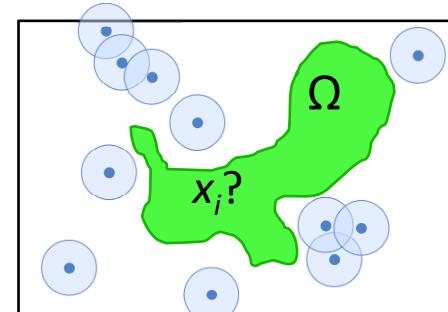
Two bonus thoughts



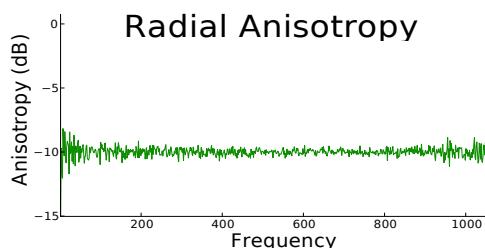
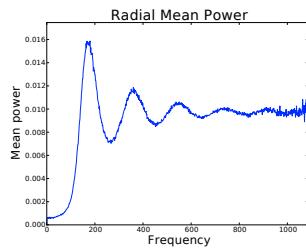
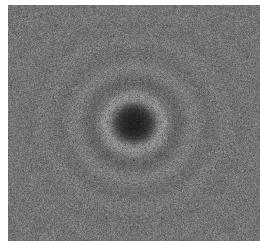
“Unbiased” Opinion

- Unbiased as a description of (serial) process
 - insertion probability independent of location

$$P(x_i \in \Omega) \propto \text{Area}(\Omega)$$



- Unbiased as a description of outcome
 - pairwise distance spectra, blue noise



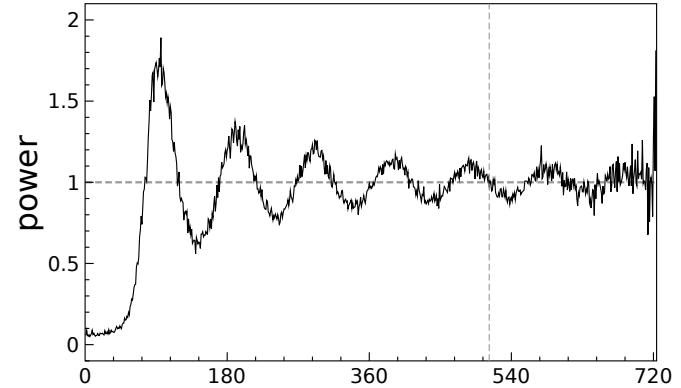
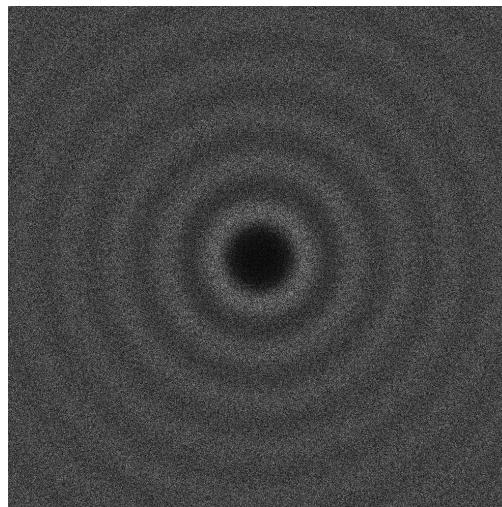
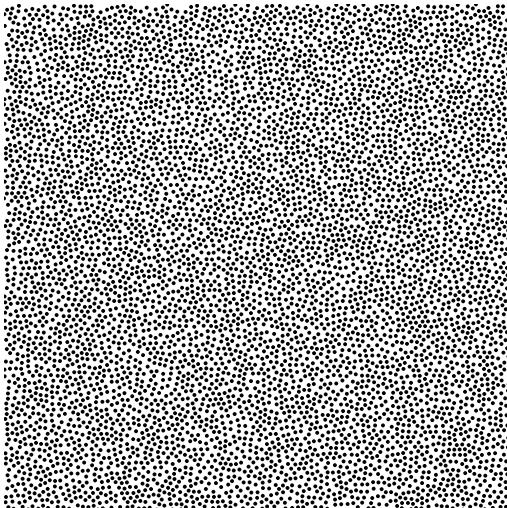
PSA code great
for standard
pictures

- Unbiased process leads to unbiased outcome,
but so might other processes
 - Opinion: need something beyond “viewgraph norm”
 - Need metrics for “how unbiased is it”
 - Define spectrum S that is the limit distribution of unbiased sampling, and standard deviations.
 - Our process generated S' , and $|S-S'| < 0.4$ std dev (S)



What is the real goal?

- Classic MPS – a lot of effort to get maximal



- Two-radii MPS, submitted to CCCG

