Q1 Answer is the same in all of a, b, c, d. Refer to notes tecture 14, page 8. Q2 (a) Gain in path shown. (L-)R) (H) (H) (WN) (-1) (WN2) (1) but $W_N = 1$ $\left(e^{-j02T}\right)$ => gain - WN overall on this path. (6) from x[7] to X[2], there is only a sigle path. In general, looking at the outputs X(k), it may seem at first there are two paths arriving at any sample, but in fact if you look at the inputs, there is only one path between each input and ench output.

Tutorial 6

Note this Q is Problem 6 in Chapter 9 of the O&S text

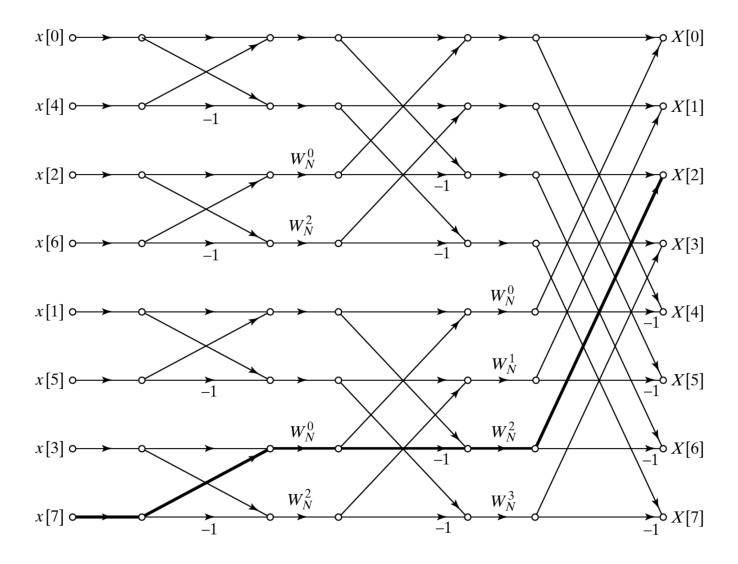


Figure P6 shows the graph representation of a decimation-in-time FFT algorithm for N=8. The heavy line shows a path from sample x[7] to DFT sample X[2].

- (a) What is the "gain" along the path that is emphasized in Figure P6?
- **(b)** How many other paths in the flow graph begin at x[7] and end at X[2]? Is this true in general? That is, how many paths are there between each input sample and each output sample?
- (c) Now consider the DFT sample X[2]. By tracing paths in the flow graph of Figure P6, show that each input sample contributes the proper amount to the output DFT sample; i.e., verify that

$$X[2] = \sum_{n=0}^{N-1} x[n]e^{-j(2\pi/N)2n}.$$

TutoRid6.

QQDFT sample X[2] Here need to trace path for all riputs x[0] --- x[7], look at sains along each path, then sum. (1) *(1) *(1) +(1) +(1) +(1) contributes Input $\times(0)$ (1)+(1)+(1)+(1)+(1)+WN --1 W8 X CI] x CJ (1)(1)(1)(1)(1)(1)(1)(1)x [2] $-1 \times [2]$ -W8 X[3] 人[3] $(1)(1)(W_{N})(-1)W_{N}^{2}(1)$ x [4] 1 X(4) (1)(1)(1)(1)(1)(1)W8 X[5] (1)(1)(1)(1) WN 1 x [5] x [6] $(1)(1) \sim_{N}^{0} (-1)(1)(1)$ -1 X[6] - W8 X[7] X [7] (1)(1) WN. -1 WN 1

From equation for DFT (as sivening)

X[2] = \(\times \times \left(2\mathbb{T}/\times) 2n \)

X[2] = \(\times \times \left(n) \)

\[= 0 \ \times \times \times \]

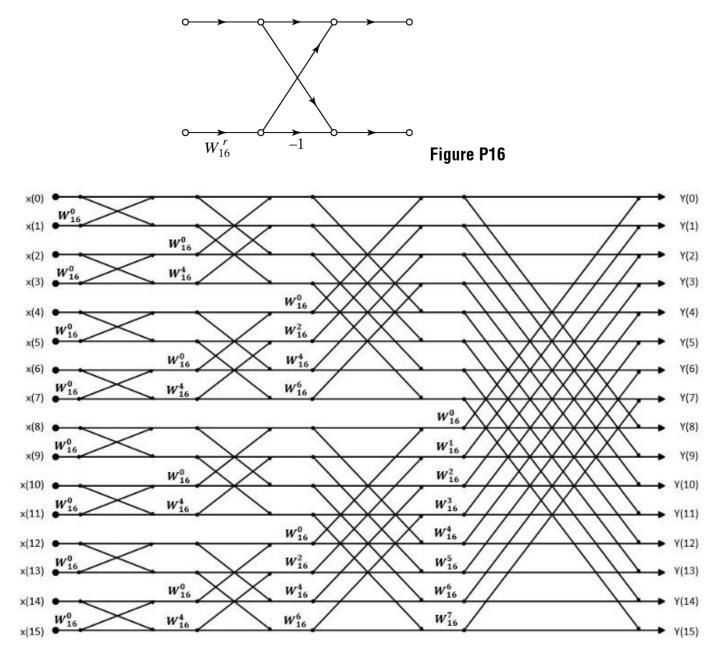
\[= \frac{1}{2\mathbb{T}^2} \]

 $X(2) = X[0]e^{-j}\frac{2F^{2}(0)}{8} + X[j]e^{-j}\frac{2F^{2}}{8} + X[2]e^{-j}\frac{2F^{2}(0)}{8} + X[3]e^{-j}\frac{2F^{2}(0)}{8} + X[4]e^{-j}\frac{2F^{2}(0)}{8} + X[5]e^{-j}\frac{2F^{2}(0)}{8} + X[5]e^{-j}\frac{2F^{2}(0)}{8$

Tutorial6. (c) (c+d) term is X (3) j 27.8 e e jet 2.5 X(e]

Note this Q is Problem 16 from Chapter 9 of the O&S text

The butterfly in Figure P16 was taken from a decimation-in-time FFT with N = 16. Assume that the four stages of the signal flow graph are indexed by m = 1, ..., 4. What are the possible values of r for each of the four stages?

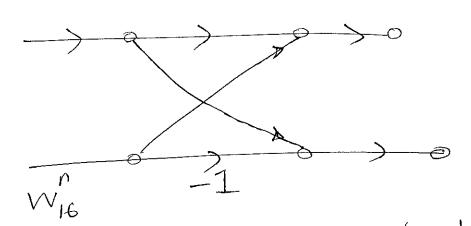


Source:

 $https://www.researchgate.net/publication/304297645_FPGA_design_and_implementation_of_radix-2_Fast_Fourier_Transform_algorithm_with_16_and_32_points/figures?lo=1$

The diagram above is useful for reference. With extra samples, each stage will have more samples to process. Note changes in twiddle factors.

Q3.



Decimation in time with N=16

Get 4 stages in our signal from

graph (necall 3 for N=8=2, hence

now 4 stages, always V stages with 2)

Stane M	values of r for W16
Stage M	
2	0,4 iè W16 W16
3	02,4,6
4	0,1,2,3,4,5,6,7

Q4.

In lectures, we showed that for a sequence x[n] which is zero for n < 0 and for n > N-1 and with $N = 2^v$, where v is a positive integer that if we et g[n] = x[2n] and h[n] = x[2n+1] then the N-point discrete Fourier transform (DFT) of the sequence x[n] can be obtained by appropriately combining the N/2-point DFTs of the sequences g[n] and h[n].

Extend this, using a similar method, to show that the sequences g[n] and h[n] can be again decimated.

Using your findings, draw the flow graph for the case N=8, showing how the 4-point (N/2) DFTs can be performed using instead the 2-point (N/4) point DFTs.

Tutorial 6. 94. In the Rectures we showed that we can write N-point DFT of x(n) in texms of N/2 point DFIs of g[n] and h[n] where they are the even I odd samples in X(n) as definied in the question. Proof was done in lectures and we Decimale g(n) and h(n) further=) divide by factor of 2 again, Lets look at gCnJ first. G[R] = N/2-1 G[R] = S[n] W_{N/2} Divide into even and odd like before nk

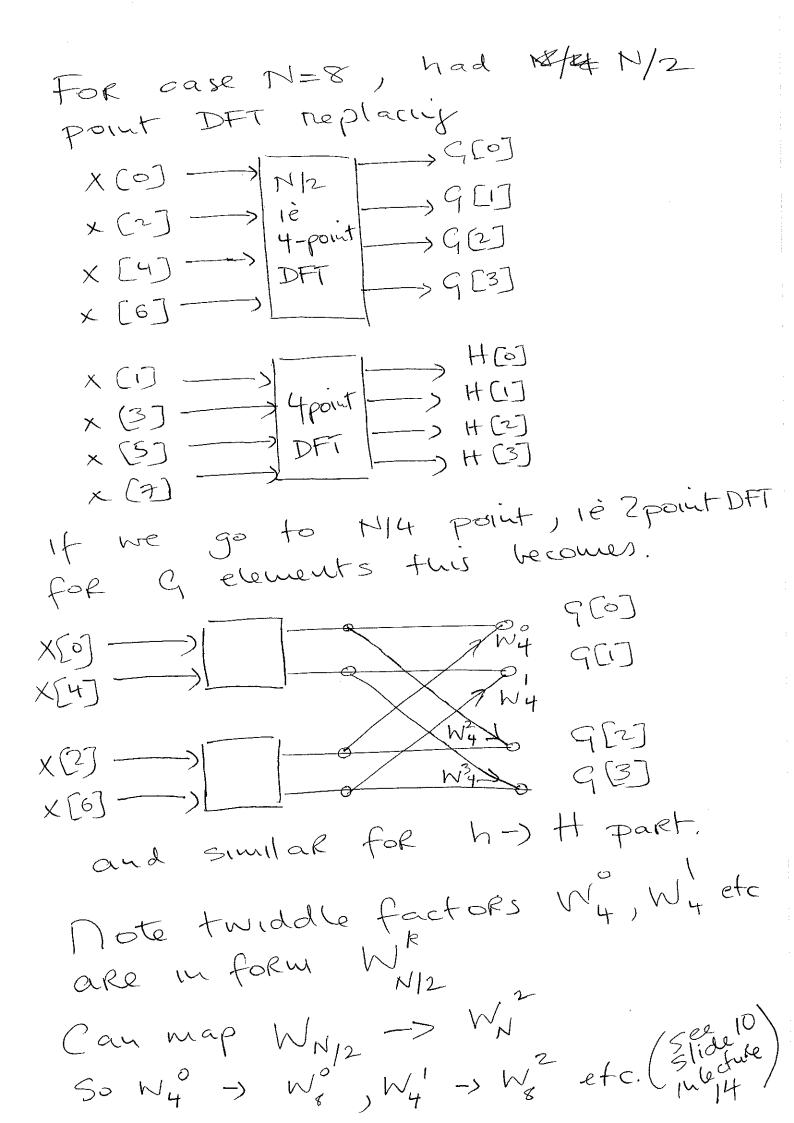
N/2-1

N/2-1

Seven and odd like before

N/2-1

N/2-Which just like in original proof give's $G(k) = \frac{N/4-1}{5} g[2n]W_{N/4}^{nk} + W_{N/2} \frac{5}{n=0} g[2n+1]W_{N/4}^{nk}$



You need to analyse the frequency content of a signal with 1024 samples. How many complex multiplications are required for a straightforward DFT calculation? Compare this with the number of complex multiplications needed for a 1024-point FFT.

Q5. 1024 samples.

Direct DFT => N2 complex mults.

FFT O(N) complex mults.

-> 1,048,576 Vs 1024