

Deadlock

(And how you will likely encounter it)

Deadlock

- Classic Deadlock Definition
 - 2 threads- each have a resource that the other wants

```
1  mutex m1;
2  mutex m2;
3  void f1(int i){
4      while(i>0){
5          lock_guard<mutex> lg1(m1);
6          lock_guard<mutex> lg2(m2);
7      }
8  }
9
10 void f2(int i){
11     while(i>0){
12         lock_guard<mutex> lg1(m2); //mutexes aquired out of order
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14     }
15 }
16
17 int main() {
18     //the threaded way with 2 mutexes
19     thread t1(f1, NUMB_TIMES);
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21     t1.join();
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← If t1 is interrupted here

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If t1 is interrupted here

Then t2 runs and is interrupted here

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If t1 is interrupted here

Then t2 runs and is interrupted here

Then the program will stop making forward progress

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Demo: [simple_deadlock project](#)

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Demo: simple_deadlock project

To fix: always acquire lock objects in same order

Deadlock- Possible Solution

```
25 //the prevention lock is used to protect
26 //critical sections where locks are being aquired
27 prevention.lock();
28
29 //then aquire all necessary locks
30 //knowing that no other thread can be
31 //in a critical section protected by
32 //the prevention lock
33 L1.lock();
34 L2.lock();
35
36 prevention.unlock()
37
38
39 //grab this lock first
40 prevention.lock();
41
42 L2.lock(); //ordering does not matter now
43 L1.lock(); //given the prevention lock
44
45 prevention.unlock()
```

- Put all lock acquisition in critical sections
- Good
 - Cannot be interrupted while acquiring locks
- Bad
 - Additional lock (prevention) to manage
 - Have to know ahead of time what locks we need
 - Decreases concurrency as we are likely acquiring locks early. Critical sections are larger than needed

Deadlock- Possible Solution 2

```
mutex m1;
mutex m2;
void f1(int i){
    while(i>0){
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        lock_guard<mutex> lg2(m2);
    }
}

void f2(int i){
    while(i>0){
        lock_guard<mutex> lg2(m1);
        lock_guard<mutex> lg1(m2);
    }
}

int main() {
    //the threaded way with 2 mutexes
    thread t1(f1, NUMB_TIMES);
    thread t2(f2, NUMB_TIMES);
    t1.join();
    t2.join();
}
```

- Always acquire locks in same order
- Good
 - Cannot deadlock
- Not as good
 - Locking is not so simple. Locks are spread out among classes, functions and libraries with many conditional statements.
 - Its often hard to predict what the ordering will be.
- But probably about as good as it will get

Things that act like deadlock

- Locking twice on the same thread without an intervening unlock

```
48 mutex m;
49 void fun2(){
50     m.lock();
51 }
52 void fun1(){
53     m.lock();
54     fun2();
55 }
56 int main() {
57     // superficial blocking? Its actually undefined
58     //from the C++ 11 standard
59     //30.4.1.2.1/4 [Note: A program may deadlock if the
60     //thread that owns a mutex object calls lock() on that object.]
61     //it MAY, or it may not block. It may work on 1 compiler and
62     //not another
63     m.lock();
64     m.lock(); ← This causes thread
65                 to block and wait
66                 to acquire a lock
67                 that it already
68                 owns
69     //the way locking twice without
70     //really happens
71     fun1();
72 }
```

Things that act like deadlock

```
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49 void fun2(){
50 |     m.lock();
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52 void fun1(){
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66 |     //the way locking twice without
67 |     //really happens
68 |     fun1(); ←
69 }
```

This is how it
happens in real
code

- Locking twice on the same thread without an intervening unlock

Demo: simple_deadlock project

Deadlock – the real world

- Locking is not so simple, especially if you have more than 1 lock.
- Locks are spread out among classes, functions and libraries with many conditional statements
- Its often hard to see what the lock ordering will be
- Once a program deadlocks it stops, it does not consume processor cycles, or any more memory than it had when the deadlock occurred. But it will never exit. It must be killed and restarted.