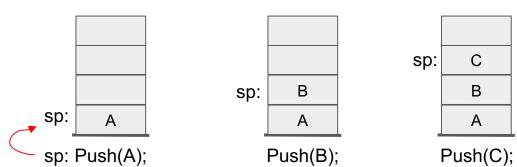
#### Macros

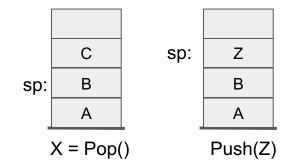
- Uses:
  - to improve readability of your program
  - to eliminate repetitive code construction
  - o to reduce overhead associated with subroutines
- Register usage must be carefully considered
- Register Approaches:
  - New Pseudo Instruction: use the \$at register -- but don't use any pseudo instructions
  - Marshalling: use only the registers provided as arguments
  - Inlined- subroutine: Utilize: \$a0, \$a1, \$a2, \$a3, \$v0, and \$v1
- Syntax:

```
.macro <name>(%arg1 .. %argn)
  list of native instructions>
.end_macro
```

### **Stack Operations**

- Push(a)  $\Leftrightarrow$  sp = sp + 1 sp[0] = a  $x = Pop() \Leftrightarrow$
- Stack is an abstract data structure
- The stack is an array of words
- Operations:
  - Push: Push(A), Push(B), Push(C)
  - $\circ$  Pop: X = Pop();
  - Push: Push(Z);





x = sp[0]

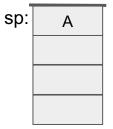
sp = sp - 1

### But the MIPS Way

 $Push(a) \Leftrightarrow sp = sp - 1$ sp[0] = a  $x = Pop() \Leftrightarrow$  x = sp[0] sp = sp + 1

- Stack is an abstract data structure
- Operations:
  - Push: Push(A), Push(B), Push(C)
  - $\circ$  Pop: X = Pop();
- sp: points to the current top of stack

Push(a) ⇔
subi \$sp, \$sp, 4
sw \$a0, 0(\$sp)



sp: B

A B C C

sp: A B C

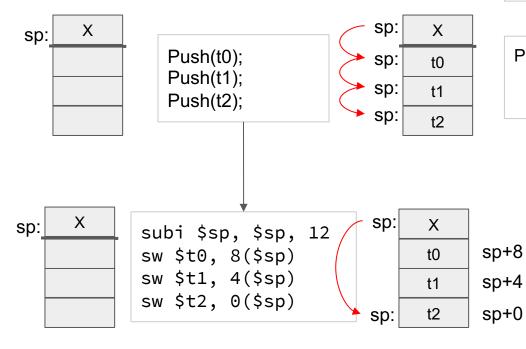
Push(A);

Push(B);

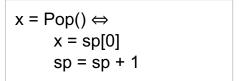
Push(C);

X = Pop()

### Multiple Pushes / Pops



```
Push(a) ⇔
sp = sp - 1
sp[0] = a
```



Push(a) ⇔
subi \$sp, \$sp, 4
sw \$a0, 0(\$sp)

x = Pop() ⇔ lw \$v0, 0(\$sp) addi \$sp, \$sp, 4

```
t0 = Pop();
t1 = Pop();
t2 = Pop();
lw $t0, 8($sp)
lw $t1, 4($sp)
lw $t2, 0($sp)
addi $sp, $sp, 12
```

## Printf: print(variable); println("string");

- Java Prototype

   public PrintStream printf(String format, Object... args)

   Java Example

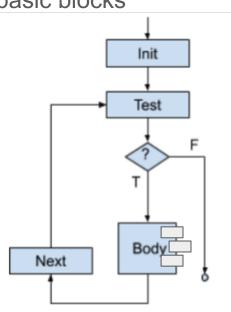
   printf("the value of x is %d", x);
- Format Specifier,
  - %conversion
- Format Conversions:
  - o d: decimal, u: unsigned decimal
  - o o: octal
  - x: hexadecimal
  - o c: character
  - o s: string
  - o f: floating point
  - ← t: binary (bi"t")

- MIPS Macros:
  - print\_d, print\_u
  - o print\_o
  - o print\_x
  - o print\_c
  - o print\_s
  - <del>○ print\_f</del>
  - o print\_t

#### **Control Flow Graph**

- A graphic representation of the representation between basic blocks
- A basic block:
  - a list of instructions with
  - a single entry point (starting point)
  - a single exit point (last instruction)
- Such representations model the behavior of our code
- Recall the while loop, and other control structures
- What about subroutines calls

(subroutine: general term for ... methods, functions, procedures, etc.)



While Loop

### Three Address Code (TAC)

- A generic assembly language in which all instructions have at most three addresses
- An address references either
  - a register location
  - a memory location
  - o an immediate value, e.g., "2"

#### Examples:

1. 
$$a = y + x$$

2. 
$$a = y$$

3. 
$$a = x + 2$$

4. 
$$b = d * 2 + y$$

$$\circ$$
 t0 = d \* 2

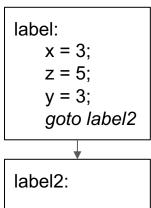
$$\circ$$
 t1 = t0 + y

$$\circ$$
 b = t1

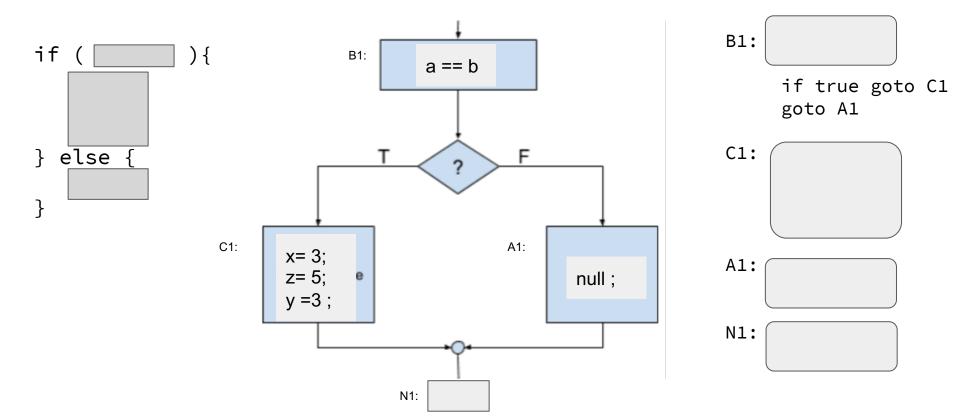
Immediate values are stored as part of the instruction in a location

#### **Basic Blocks**

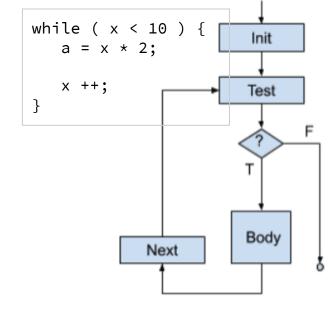
- A number of instructions in which there is
  - o a single entry point (via a label), and
  - a single exit point (via a goto)
- All programs can be broken down into a set of basic blocks
- A control flow graph determines which a basic block is executed.
- Standard control flow graphs
  - o if-then-else and all other variants (e.g., switch)
  - o while, do-while and all other variants
  - o for loop and all other variants
  - call-return

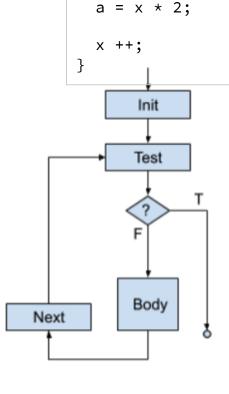


#### Code Flow: If-then-else

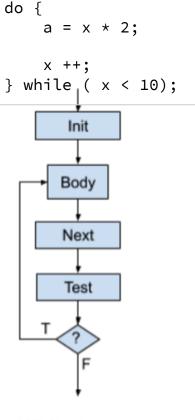


# Control Flow: Loops





until ( x>= 10 {



While Loop

Do While Loop

for ( i = 0; i < 10; i++) { a = x \* 2

Until Loop