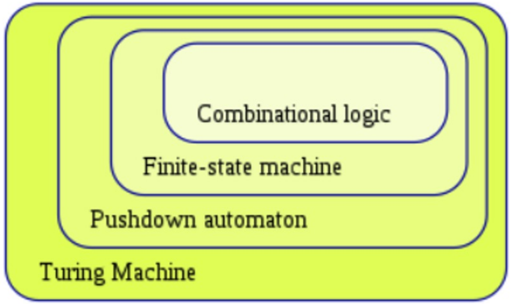
Lecture



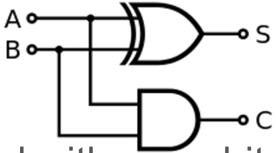
- Last time:
  - Boolean Algebra ⇔ Digital Circuits
    - Point: We can do a lot with just Combinational logic -- all true functions can be evaluated
    - Point: Digital Circuits can be built to evaluate all of these functions.
  - All we need is And (*), Or (+) and Not (')
  - Truth Table → Boolean Algebra
  - Boolean Algebra → Circuits → Boolean Algebra
  - Minimization of Circuits
    - Algebraic Transformations:
    - Karnaugh Maps

- Today: More Combinational Circuits

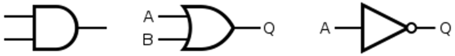
Combinational Logic

- Using just is tedious: AND (*), OR (+), NOT (')
- Solution: Build components and reuse!
    - XOR: A ⊕ B is equivalent to   (A + B) * (A' + B')

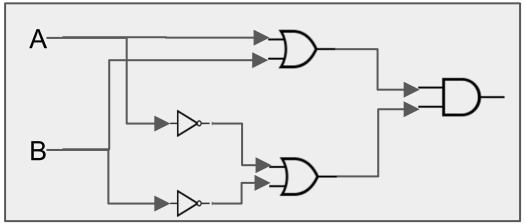
    - Half-Adder: S = A ⊕ B, C = A * B



- Bigger components and with more bits!
    - Binary Addition
    - Binary Subtraction
    - BCD Addition
    - Decoder
    - Multiplexer



XOR:

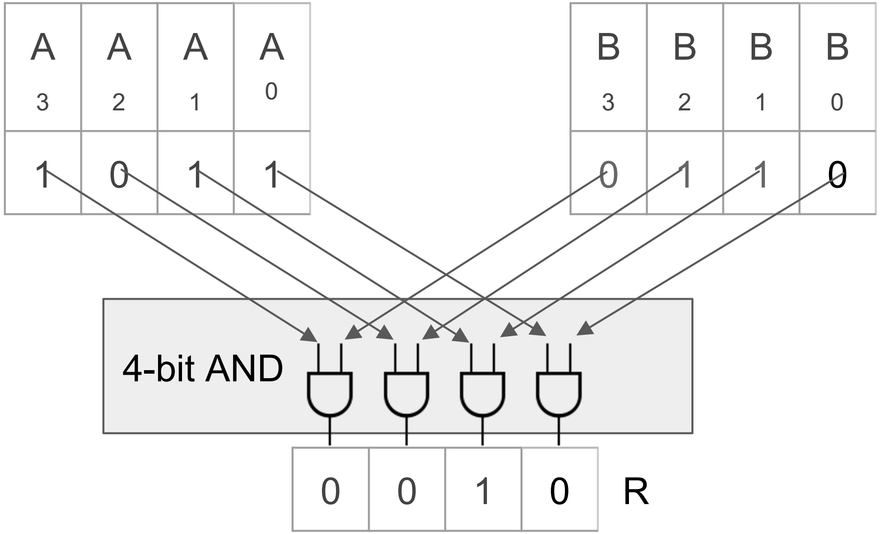


A	B	Carry	Sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

4-bit Bitwise AND

- R = A & B

	1	0	1	1	A
&	0	1	1	0	B
	0	0	1	0	R

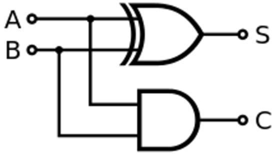


- For a n-bit operation,
  - create n-duplicates of the base circuitry
  - layout duplicates in parallel
  - package it up

1-bit Binary Addition

- Recall:
  - A + B → S, C

				1	A
+				1	B
					S



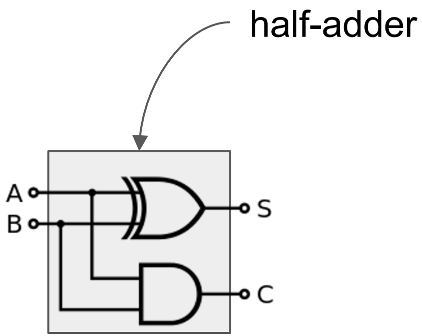
A	B	C_{out}	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

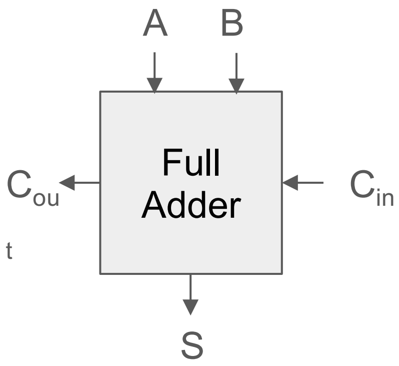
4-bit Binary Addition

- Recall:
  - $C_{in} + A_x + B_x \rightarrow C_{out}, S_x$

	1	0	1	1	A
+	0	1	1	0	B
					S

- Half-Adder is not sufficient!
  We need a Full-Adder

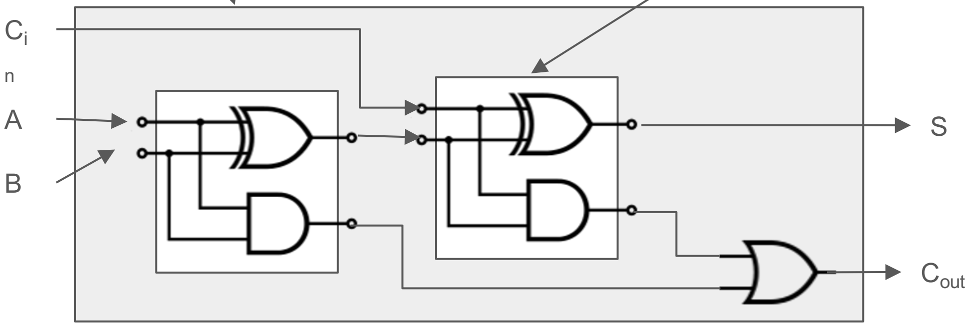




C_{in}	A	B	C_{out}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

Full Adder

half-adder

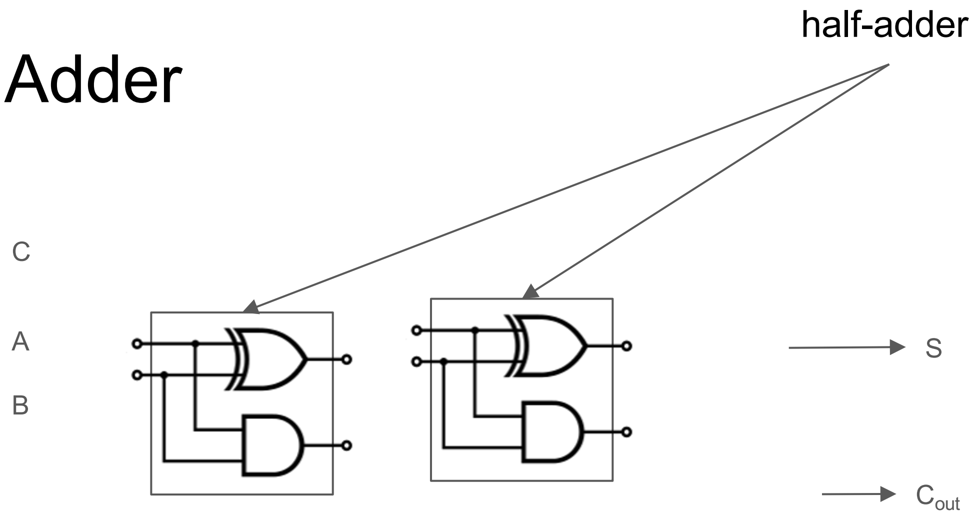


Got It? Any Questions

C_{in}	A	B	C_{out}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

- $C_{out} = AB + C_{in}(A \oplus B)$
- $S = C_{in} \oplus A \oplus B$

Full Adder



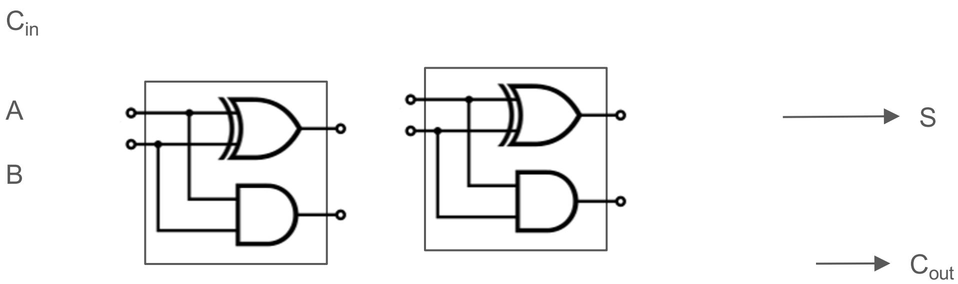
- C_{out} = C'AB +

Note: Renamed C_{in} to be C

C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	
1	1	0	1	0	
1	1	1	1	1	

Use Sum of Products

Full Adder

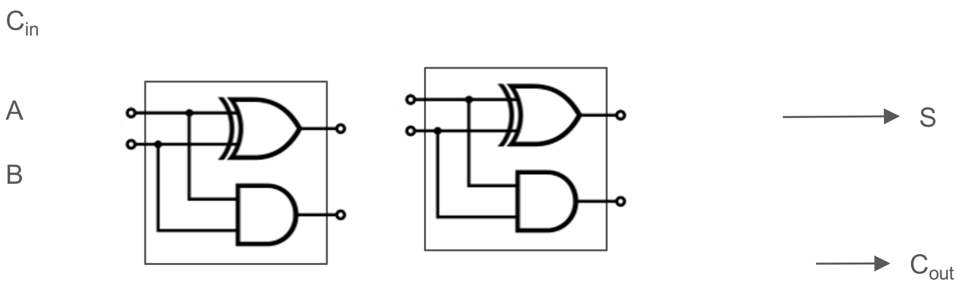


C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	CA'B
1	1	0	1	0	CAB'
1	1	1	1	1	CAB

Use Sum of Products

- C_{out} = C'AB + CA'B + CAB' + CAB

Full Adder

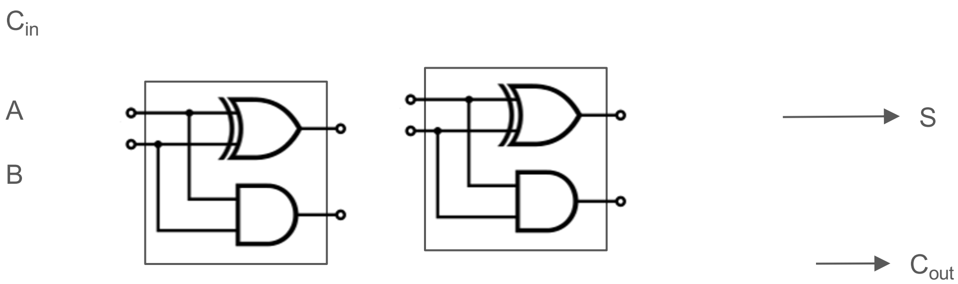


C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	CA'B
1	1	0	1	0	CAB'
1	1	1	1	1	CAB

- $C_{out} = C'AB + \boxed{CA'B + CAB'} + \boxed{CAB}$
   - ➢ $C_{out} = C'AB + CAB + CA'B + CAB'$

Use Commutative Property

Full Adder

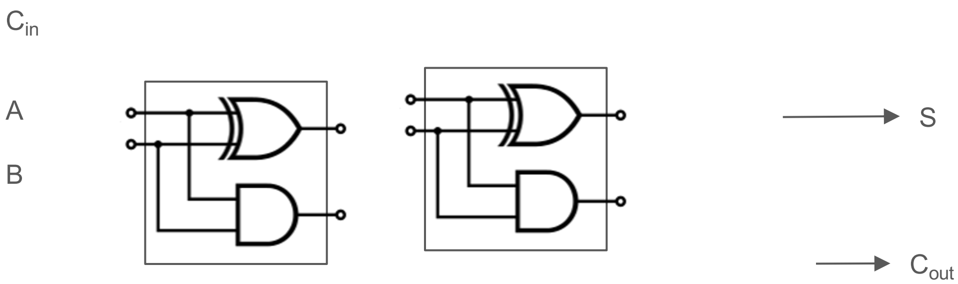


C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	CA'B
1	1	0	1	0	CAB'
1	1	1	1	1	CAB

Use Distributive Property

- $C_{out} = \boxed{C'AB} + \boxed{CAB} + CA'B + CAB'$
  - $C_{out} = (C' + C)AB + CA'B + CAB'$

Full Adder

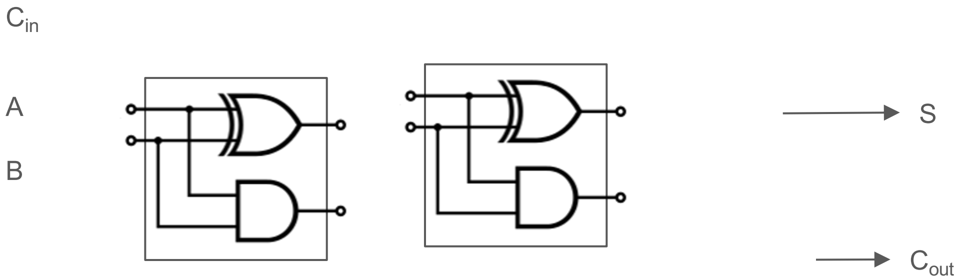


C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	CA'B
1	1	0	1	0	CAB'
1	1	1	1	1	CAB

- $C_{out} = (\boxed{C' + C})AB + CA'B + CAB'$
  - $C_{out} = (\text{true})AB + CA'B + CAB'$

Use Complement Property

Full Adder

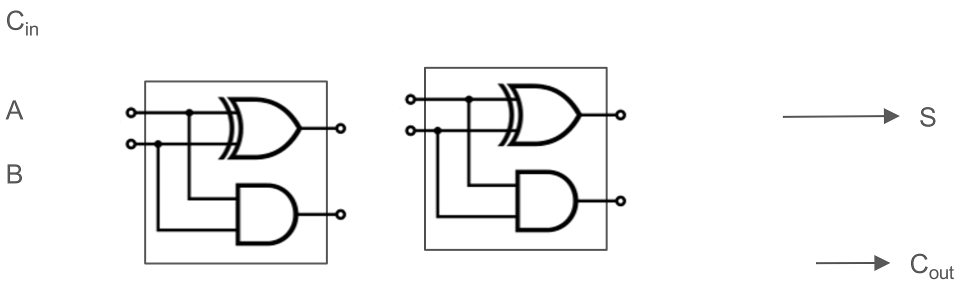


C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	CA'B
1	1	0	1	0	CAB'
1	1	1	1	1	CAB

- $C_{out} = \boxed{(true)AB} + CA'B + CAB'$
  - $C_{out} = AB + CA'B + CAB'$

Use Identity Property

Full Adder

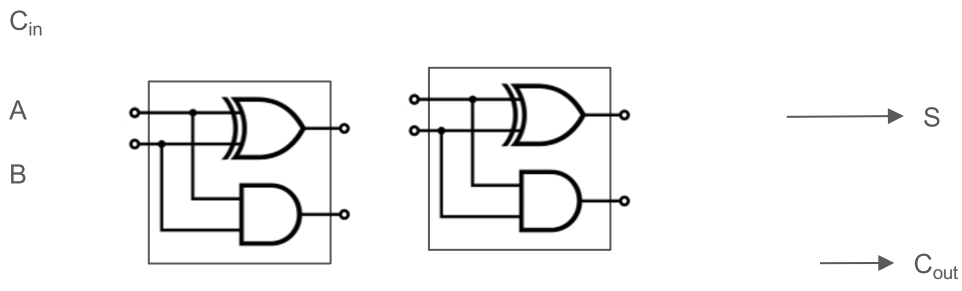


C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	CA'B
1	1	0	1	0	CAB'
1	1	1	1	1	CAB

Use Distributive Property

- $C_{out} = AB + \boxed{CA'B + CAB'}$
  - ➢ $C_{out} = AB + C(A'B + AB')$

Full Adder



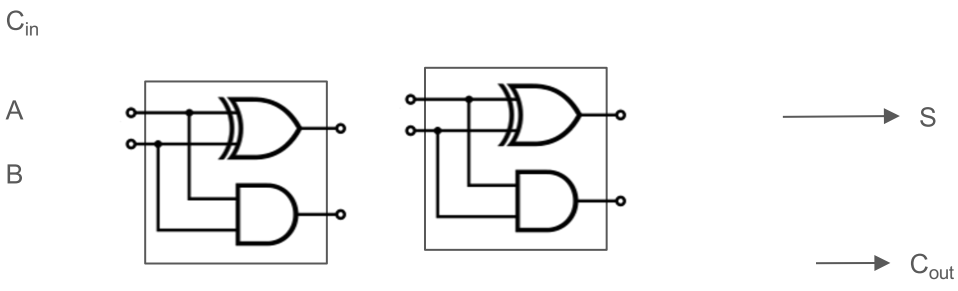
C_{in}	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	
0	1	0	0	1	
0	1	1	1	0	C'AB
1	0	0	0	1	
1	0	1	1	0	CA'B
1	1	0	1	0	CAB'
1	1	1	1	1	CAB

- $C_{out} = AB + \boxed{C(A'B + AB')}$

✓ $C_{out} = AB + C(A \oplus B)$

$A \oplus B \Leftrightarrow A'B + AB'$

Full Adder



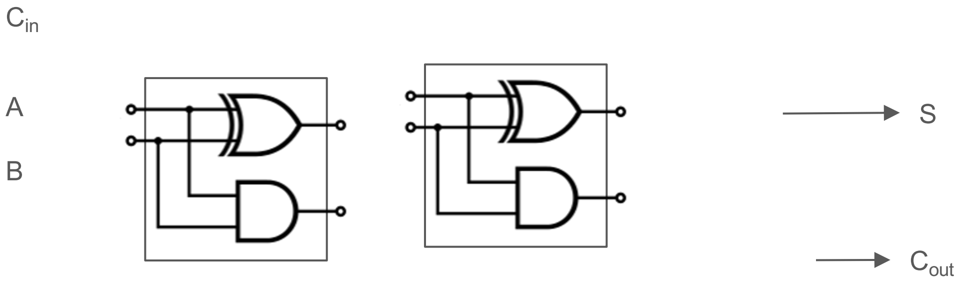
C	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	C'A'B
0	1	0	0	1	C'AB'
0	1	1	1	0	
1	0	0	0	1	CA'B'
1	0	1	1	0	
1	1	0	1	0	
1	1	1	1	1	CAB

Sum of Products:	C'A'B + C'AB' + CA'B' + CAB

- ✓ $C_{out} = AB + C(A \oplus B)$
- $S = C'A'B + C'AB' + CA'B' + CAB$

$A \oplus B \Leftrightarrow A'B + AB'$

Full Adder



C	A	B	C_{out}	S	
0	0	0	0	0	
0	0	1	0	1	C'A'B
0	1	0	0	1	C'AB'
0	1	1	1	0	
1	0	0	0	1	CA'B'
1	0	1	1	0	
1	1	0	1	0	
1	1	1	1	1	CAB

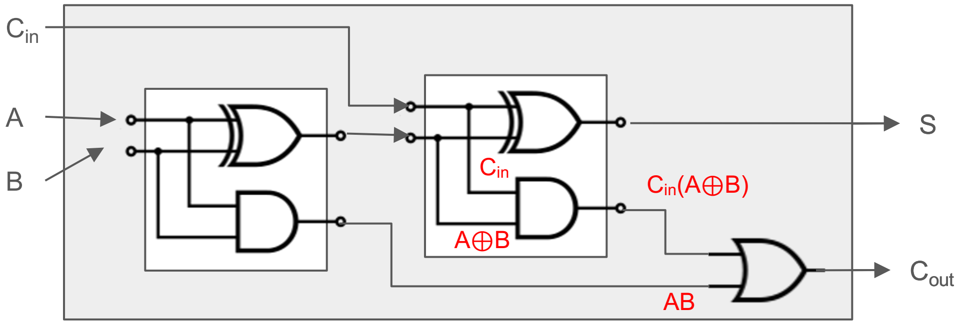
Sum of Products:

$$C'A'B + C'AB' + CA'B' + CAB$$
$$= C'(A'B + AB') + C(A'B' + AB)$$
$$= C'(A\oplus B) + C(A\oplus B)'$$
$$= C\oplus(A\oplus B)$$
$$= C\oplus A\oplus B$$

- ✓ $C_{out} = AB + C(A\oplus B)$
- ● $S = C'A'B + C'AB' + CA'B' + CAB$
- ✓ $S = C\oplus A\oplus B$

$A \oplus B \Leftrightarrow A'B + AB'$

Full Adder



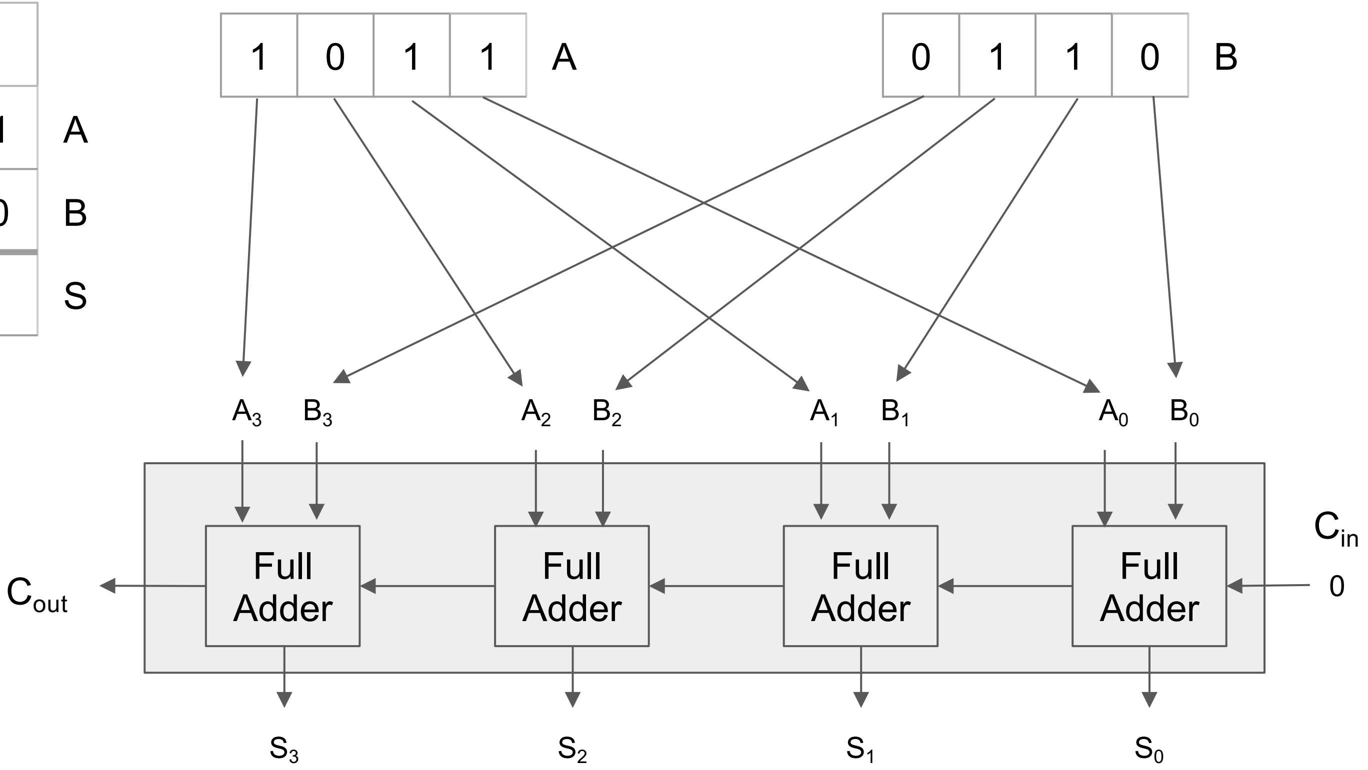
C	A	B	C_{out}	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

Note: Renamed C to be C_{in}

- ✓ $C_{out} = AB + C_{in}(A \oplus B)$
- ✓ $S = C_{in} \oplus A \oplus B$

4-bit Binary Addition (aka: 4-bit Full Adder)

	1	0	1	1	A
+	0	1	1	0	B
					S

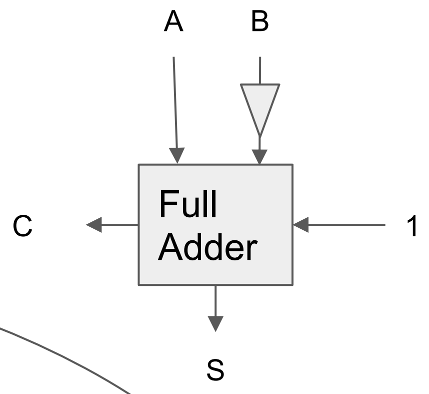


Binary Subtractor

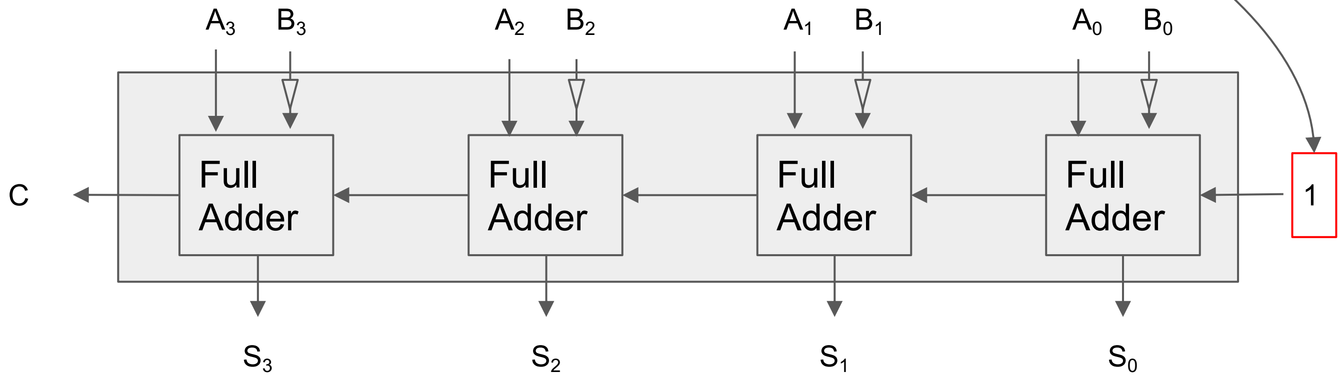
- Recall: $A - B$

  $= A + (\text{2's complement of } B)$

  $= A + (\text{1's complement of } B) + \boxed{1}$



- 4-bit Binary Subtractor



4-bit BCD Addition

- BCD: 35 + 28

$$\begin{array}{r} 35 \\ +\ 28 \end{array}$$

	0	1	1	1	
	0	0	1	1	3
+	0	0	1	0	2

	0	0	0	0	
	0	1	0	1	5
+	1	0	0	0	8

#	Encoding $S_3S_2S_1S_0$		Encoding $S_3S_2S_1S_0$
0	0 0 0 0	8	1 0 0 0
1	0 0 0 1	9	1 0 0 1
2	0 0 1 0	invalid	1 0 1 0
3	0 0 1 1	invalid	1 0 1 1
4	0 1 0 0	invalid	1 1 0 0
5	0 1 0 1	invalid	1 1 0 1
6	0 1 1 0	invalid	1 1 1 0
7	0 1 1 1	invalid	1 1 1 1

- Perform Regular Binary Addition, but account for the invalid patterns
- Add six upon whenever you are in the deadzone or there is overflow
    - Invalid $= S_3 * (S_2 + S_1)$
    - Overflow $= C_{out}$

Overflow = C_{out}

4-bit BCD Addition

	0	1	1	0	6
+	0	1	0	1	5

	0	0	1	0	2
+	0	1	0	1	5