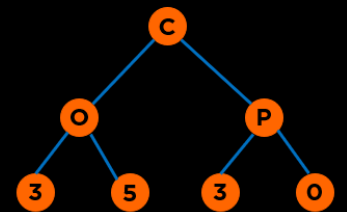


Graphs



Categories of Data Structures

Linear Ordered

Lists

Stacks

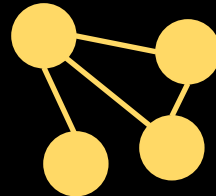
Queues



Non-linear Ordered

Trees

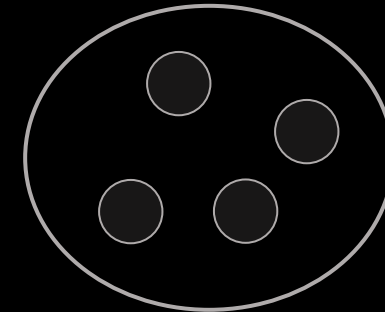
Graphs



Not Ordered

Sets

Tables/Maps



Recap

- **Graphs**

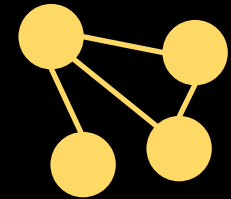
- **Terminology**
- **Types**

- **Graph Implementations**

- **Edge List**
- **Adjacency Matrix**
- **Adjacency List**

Non-linear Ordered

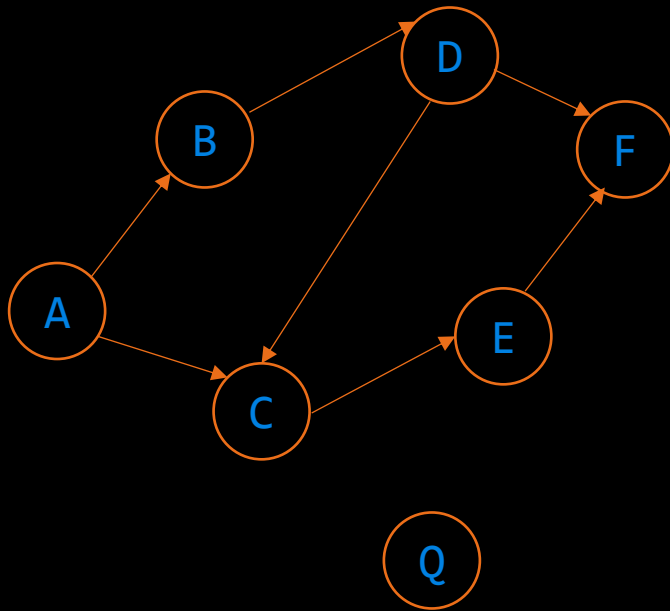
Graphs



Shortest Path

s-t Path

Is there a path between vertices s and t?

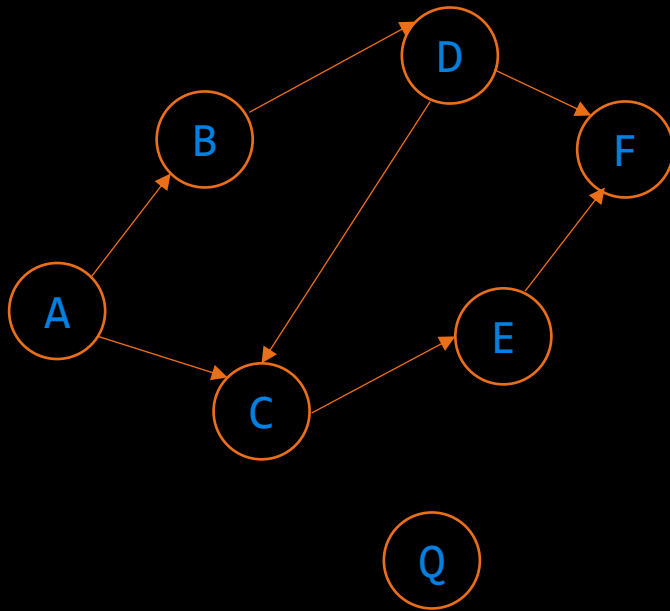


Is there a path between vertices A and C? - Yes

Is there a path between vertices A and Q? - No

s-t Path

Is there a path between vertices s and t?



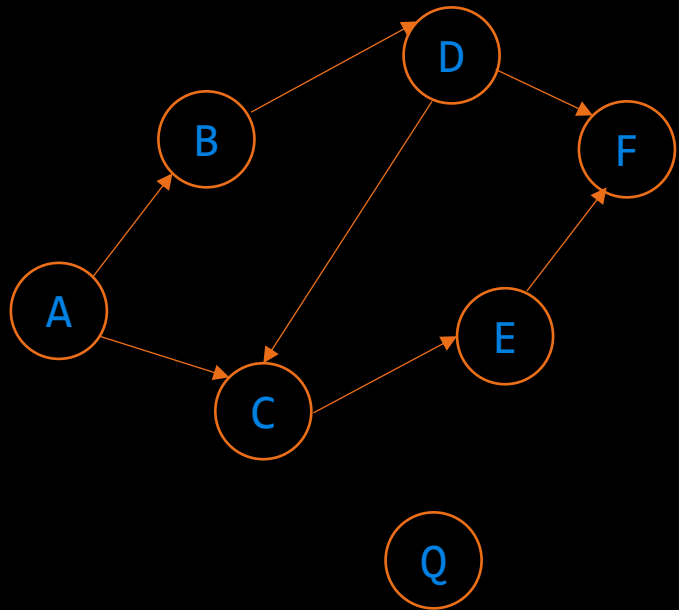
Is there a path between vertices A and C? - Yes

Is there a path between vertices A and Q? - No

Solution

Perform **DFS** or **BFS** with source “s” and if we encounter “t” in the path/traversal, then return True otherwise False

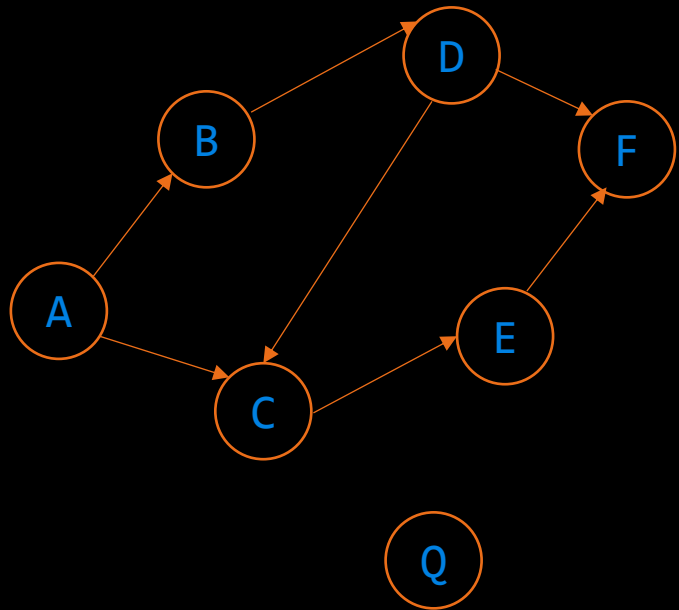
7.2.1 DFS to Find Whether a Given Vertex is Reachable (Iterative)



s-t Path

```
1.  bool dfs(const Graph& graph, int src, int dest)
2.  {
3.      set<int> visited;
4.      stack<int> s;
5.      visited.insert(src);
6.      s.push(src);
7.      while(!s.empty())
8.      {
9.          int u = s.top();
10.         s.pop();
11.         for(auto v: graph.adjList[u])
12.         {
13.             if(v == dest)
14.                 return true;
15.             if ((visited.find(v) == visited.end()))
16.             {
17.                 visited.insert(v);
18.                 s.push(v);
19.             }
20.         }
21.     }
22.     return false;
23. }
```

7.2.1 DFS to Find Whether a Given Vertex is Reachable (Recursive)



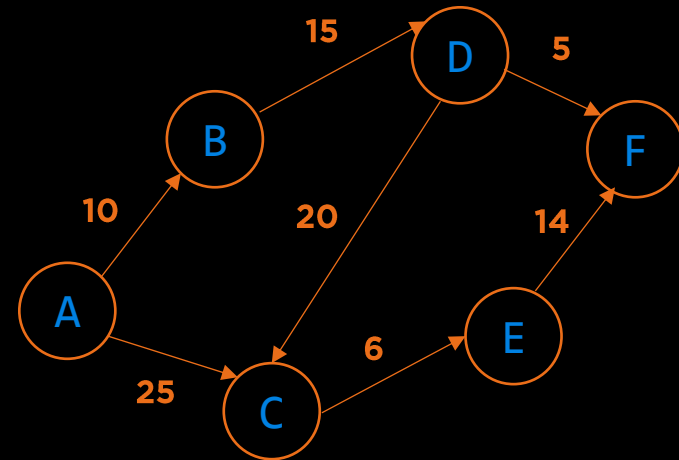
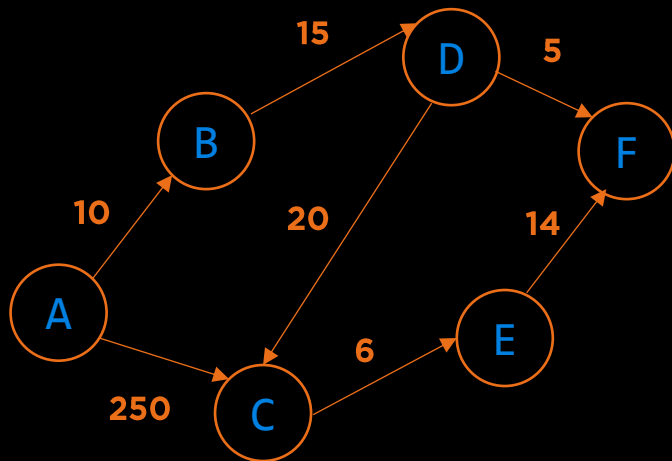
s-t Path: Recursive

```
1. bool dfs_helper(const Graph& graph, int src, int dest, vector<bool>& visited)
2. {
3.     visited[src] = true;
4.
5.     if (src == dest)
6.         return true;
7.
8.     for (int neighbor : graph.adjList[src]) {
9.         if (!visited[neighbor]) {
10.            if (dfs_helper(graph, neighbor, dest, visited))
11.                return true;
12.        }
13.    }
14.    return false;
15. }
16.
17. bool dfs(const Graph& graph, int src, int dest)
18. {
19.     vector<bool> visited(graph.numVertices);
20.     return dfs_helper(graph, src, dest, visited);
21. }
```


Problem with s-t Path

What if the edges are weighted?

The algorithms do not consider the weights.

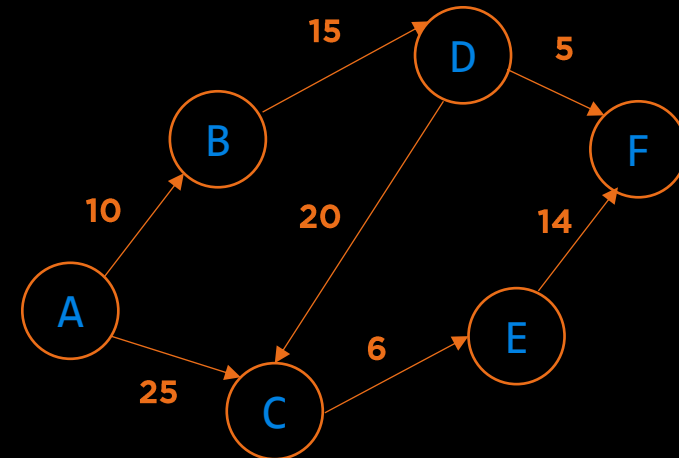
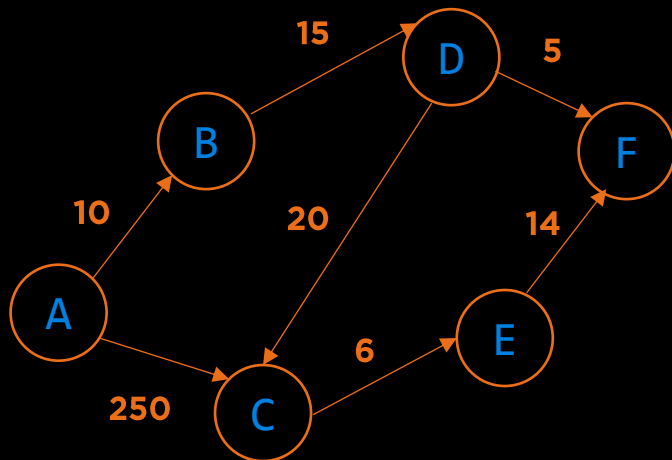


Problem with s-t Path

What if the edges are weighted?

The algorithms do not consider the weights.

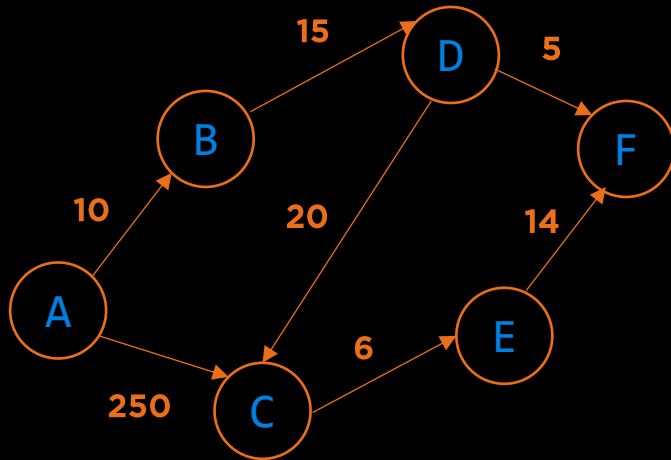
Example 1: Path for A to C will be **A-B-D-C** for a **DFS traversal** which will have a total cost of **45** against **25** for the path directly from **A-C**.



Example 2: Path for A to C will be **A-C** for a **BFS traversal** which might have a total cost of **250** against **45** for the path directly from **A-B-D-C**.

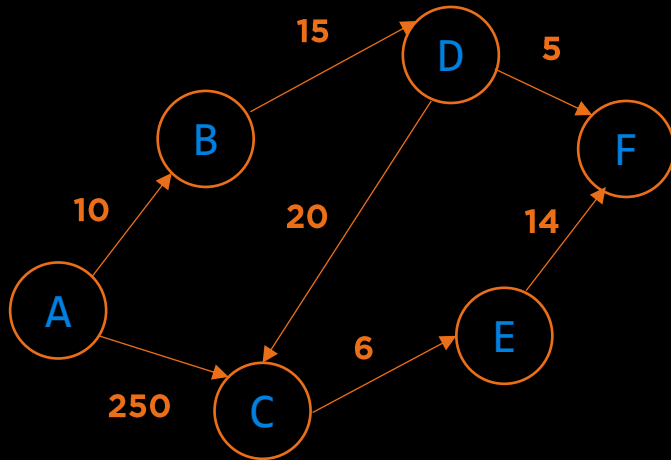
Shortest Weighted s-t Path

What is the shortest weighted path between vertices **s** and **t**?



Shortest Weighted s-t Path

What is the shortest weighted path between vertices **s** and **t**?

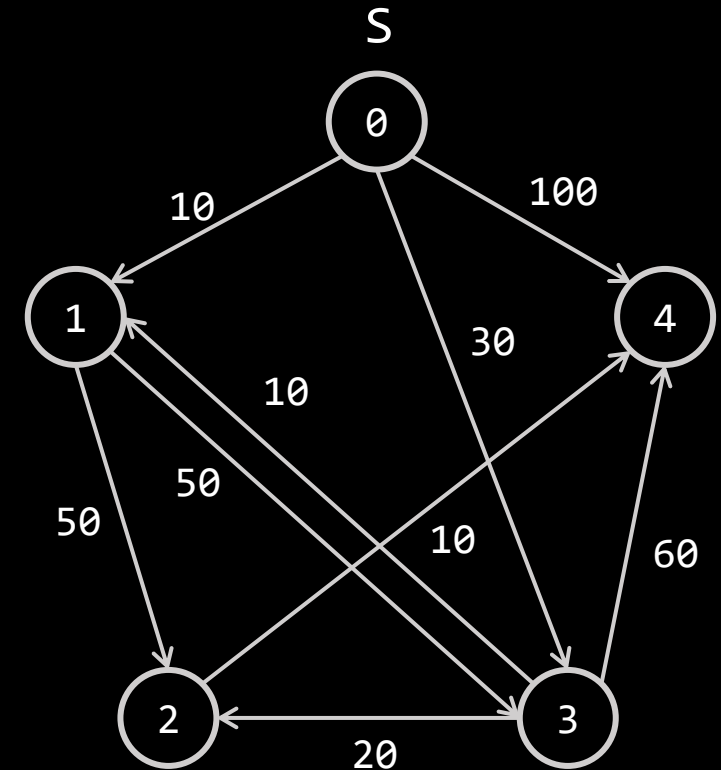


- **Dijkstra's Algorithm**
 - Single Source: Path to all vertices
 - Directed Graphs
 - No negative weights allowed
 - No negative weight cycles allowed
- **Bellman Ford**
 - Single Source: Path to all vertices
 - Negative Weights allowed
 - No negative weight cycles allowed
- **Floyd-Warshall**
 - All pair shortest paths
- **A* Search**

Dijkstra's Shortest Path Algorithm

Example

- Specify a source vertex, S
- Initialize two arrays and two sets
 - Set S will contain the vertices for which we have computed the shortest distance
 - Initially S will be empty
 - Set $V-S$ will contain the vertices we still need to process
 - Initialize $V-S$ by placing all vertices into it
- $d[v]$ will contain shortest distance from s to v
 - Initially all $d[v]$'s will be set to infinity except for source which will be 0
- $p[v]$ will predecessor of v in the path from s to v
 - Initially all $p[v]$'s will be set to -1



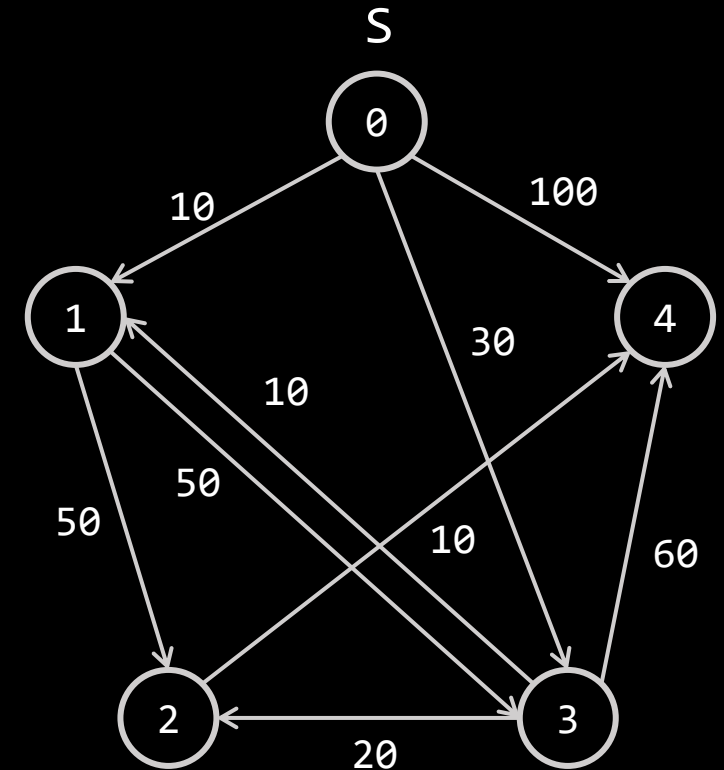
Dijkstra's Shortest Path Algorithm

Example

Computed, $S = \{\}$

Needs processing, $V-S = \{0, 1, 2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	∞	-1
2	∞	-1
3	∞	-1
4	∞	-1



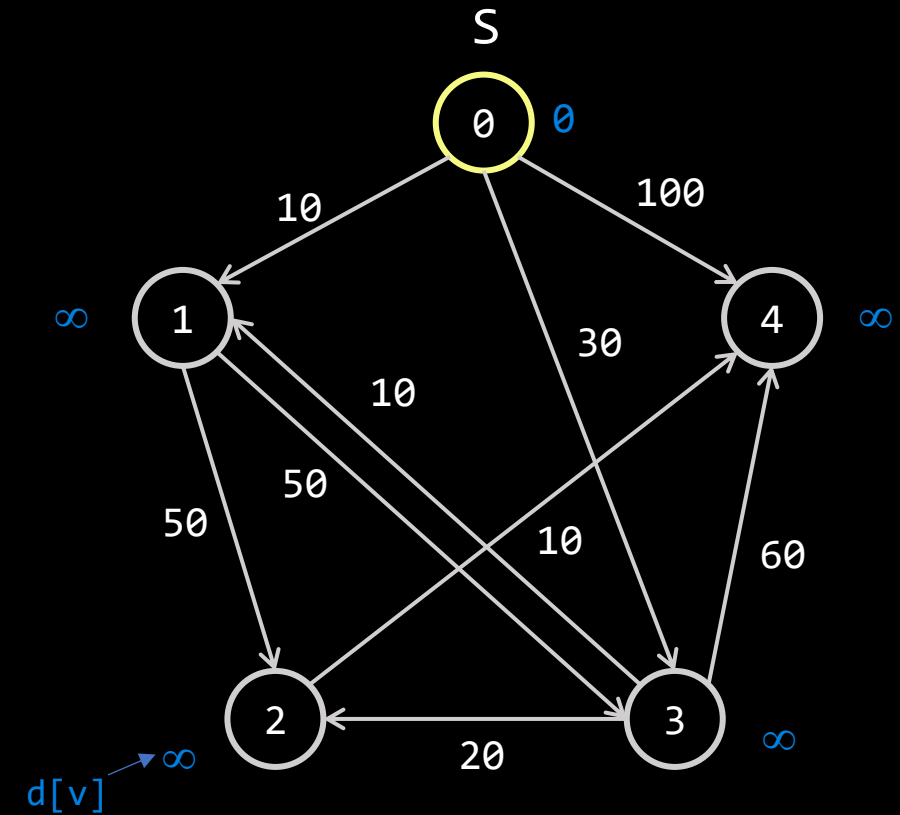
Dijkstra's Shortest Path Algorithm

Example: Start with vertex that has minimum distance in $d[v]$, i.e. 0 and add to Computed

Computed, $S = \{0\}$

Needs processing, $V-S = \{1, 2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	∞	-1
2	∞	-1
3	∞	-1
4	∞	-1



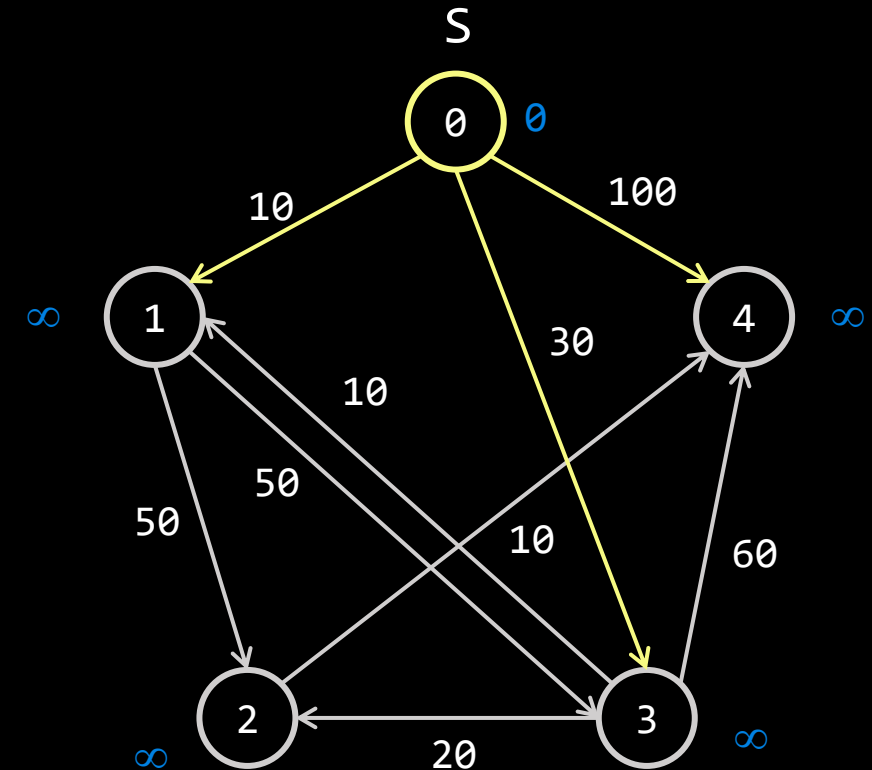
Dijkstra's Shortest Path Algorithm

Example: Process edges adjacent to the vertex 0 and update distances based on relaxation*

Computed, $S = \{0\}$

Needs processing, $V-S = \{1, 2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	∞	-1
2	∞	-1
3	∞	-1
4	∞	-1



* Relaxation

```
if (dist[v] > dist[u] + w)
    dist[v] = dist[u] + w;
```

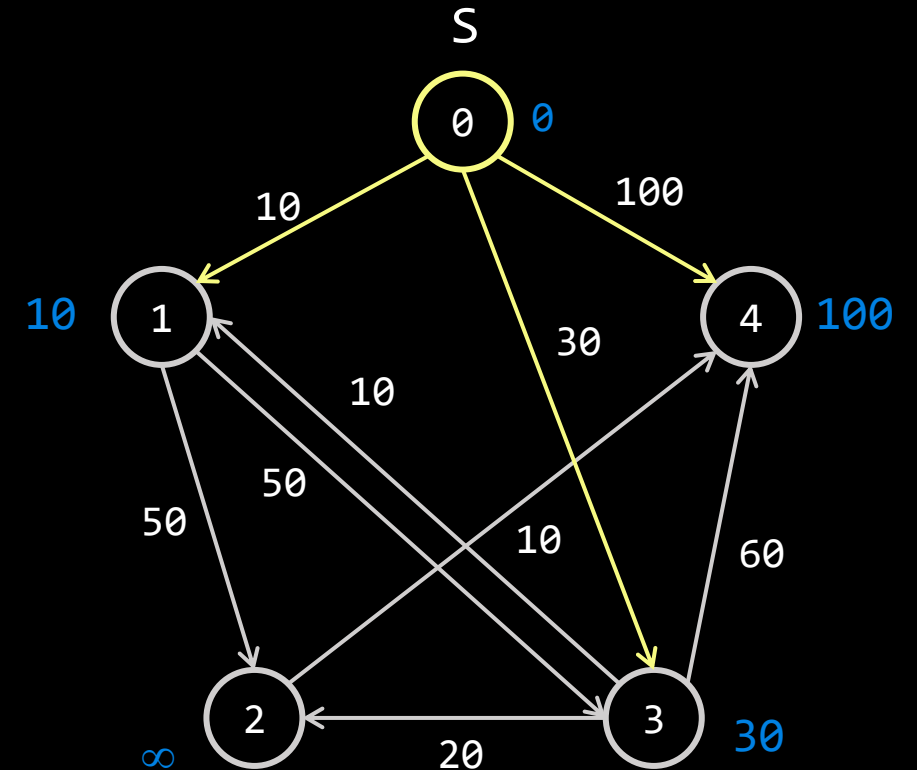

Dijkstra's Shortest Path Algorithm

Example: Process edges adjacent to the vertex 0 and update distances based on relaxation*

Computed, $S = \{0\}$

Needs processing, $V-S = \{1, 2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	∞	-1
3	30	0
4	100	0



* Relaxation

```
if (dist[v] > dist[u] + w)
    dist[v] = dist[u] + w;
```

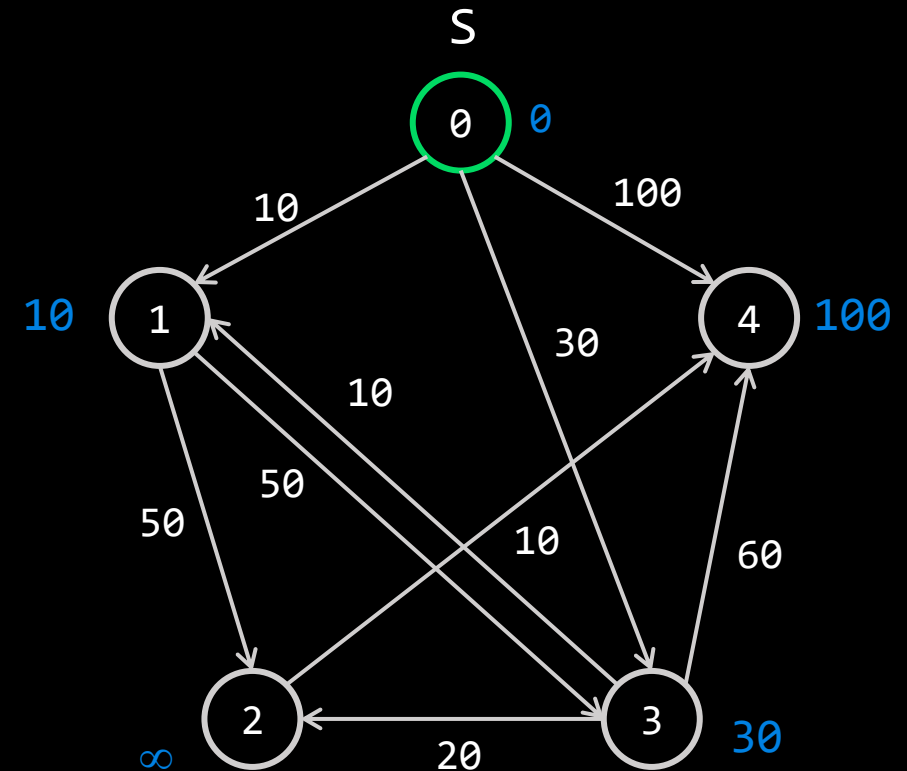
Dijkstra's Shortest Path Algorithm

Example: 0 is now done. Next, repeat the process picking the minimum element in $d[v]$ that has not been computed

Computed, $S = \{0\}$

Needs processing, $V-S = \{1, 2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	∞	-1
3	30	0
4	100	0



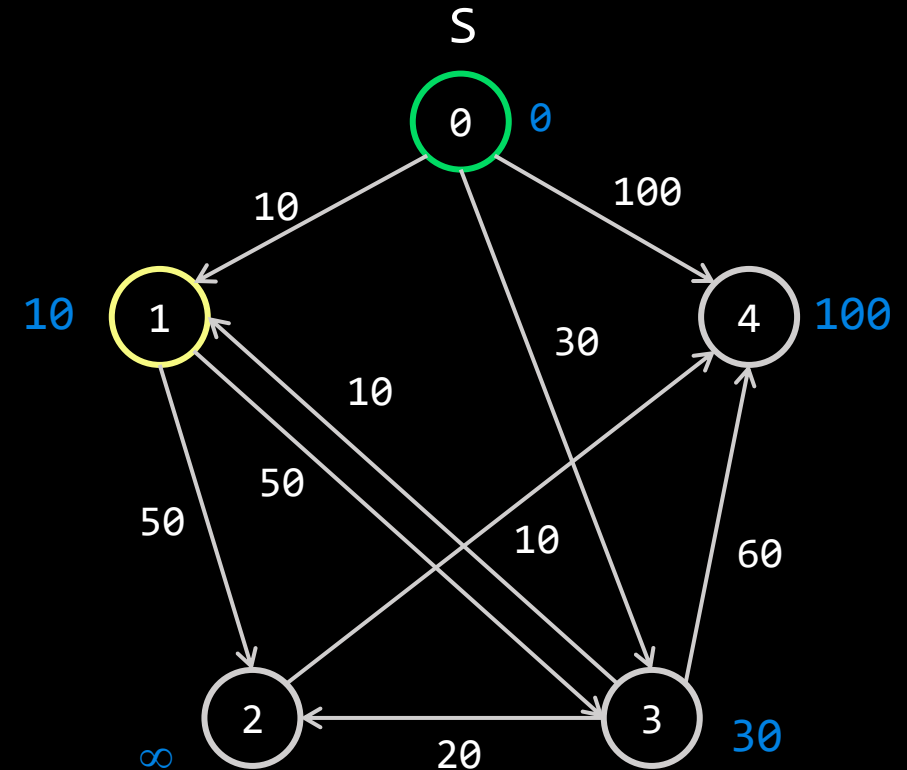
Dijkstra's Shortest Path Algorithm

Example: Pick 1

Computed, $S = \{0\}$

Needs processing, $V-S = \{1, 2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	∞	-1
3	30	0
4	100	0



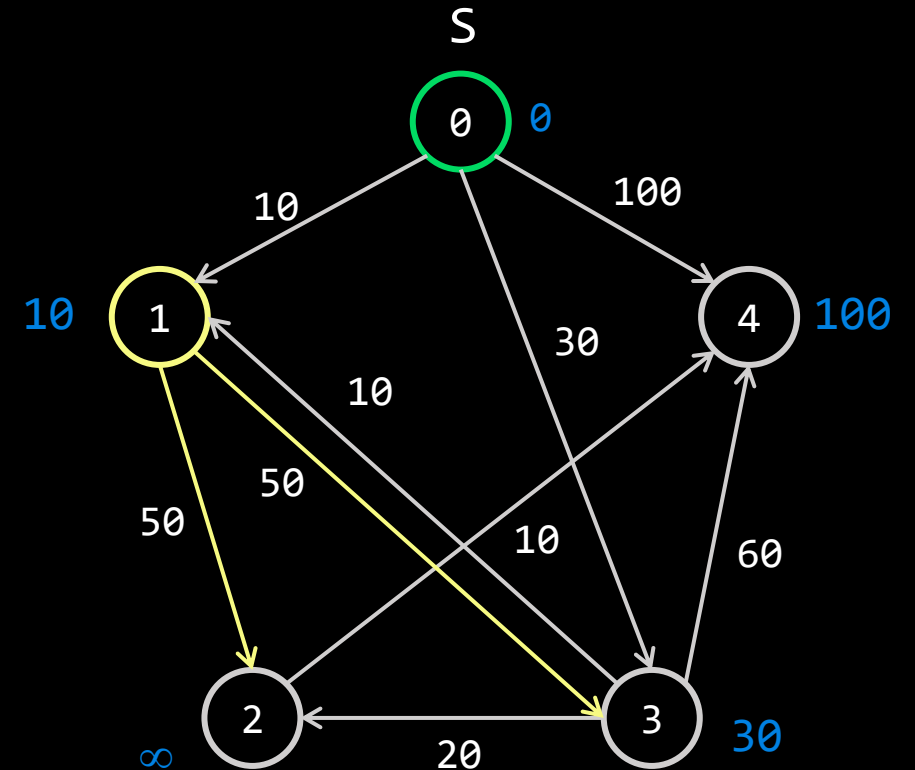
Dijkstra's Shortest Path Algorithm

Example: Process edges adjacent to the vertex 1 and update distances based on relaxation

Computed, $S = \{0\}$

Needs processing, $V-S = \{1, 2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	∞	-1
3	30	0
4	100	0



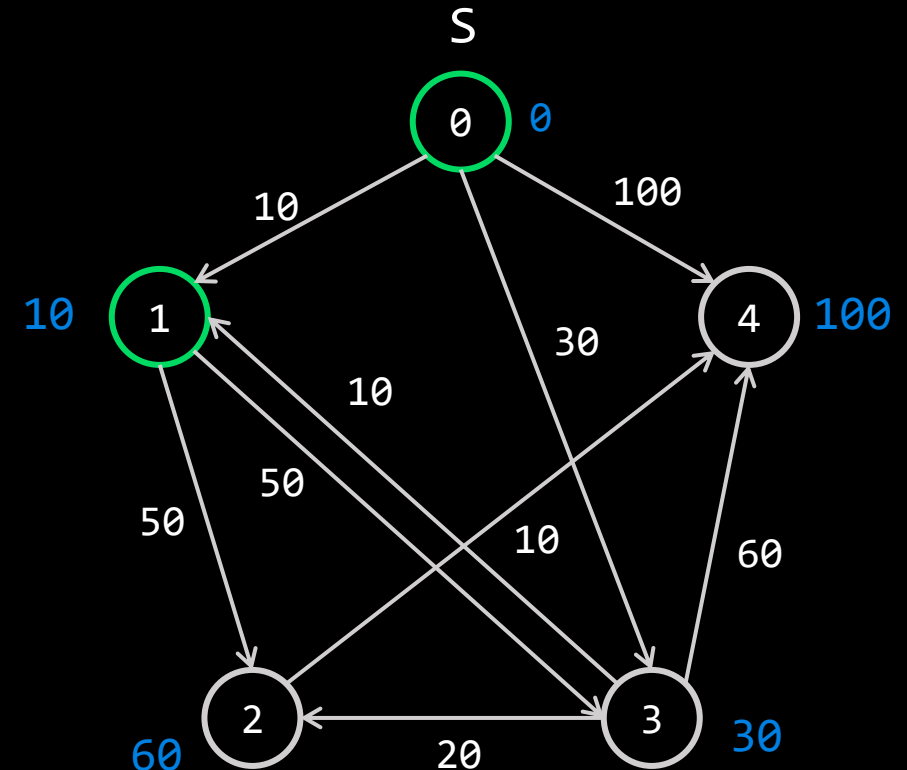
Dijkstra's Shortest Path Algorithm

Example: 1 is now done. Next, repeat the process picking the minimum element in $d[v]$ that has not been computed

Computed, $S = \{0, 1\}$

Needs processing, $V-S = \{2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	60	1
3	30	0
4	100	0



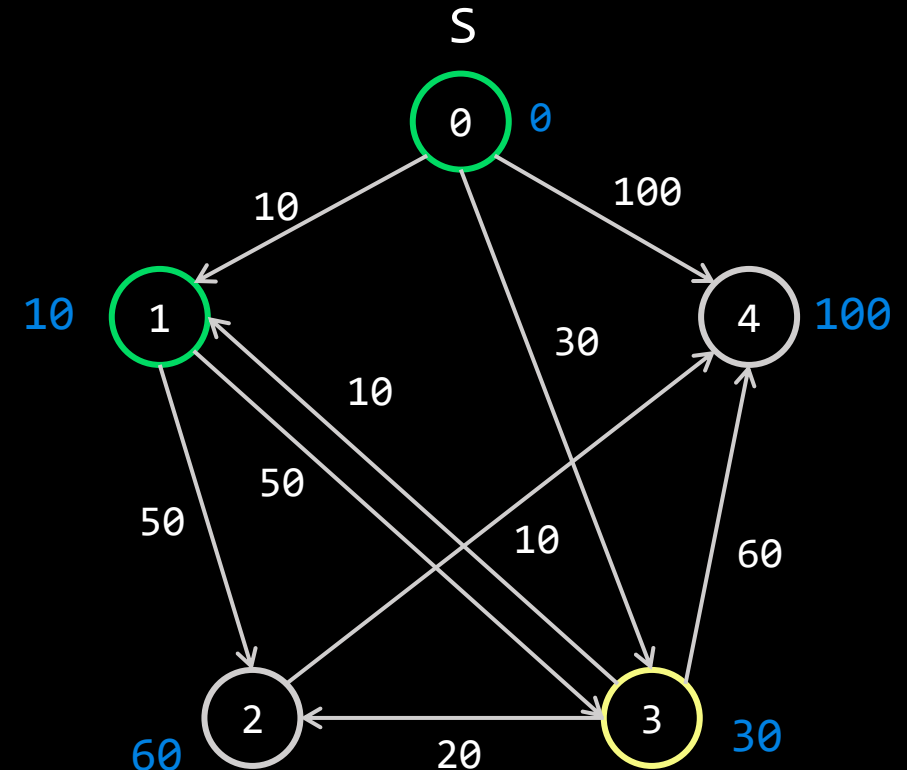
Dijkstra's Shortest Path Algorithm

Example: Pick 3

Computed, $S = \{0, 1\}$

Needs processing, $V-S = \{2, 3, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	60	1
3	30	0
4	100	0



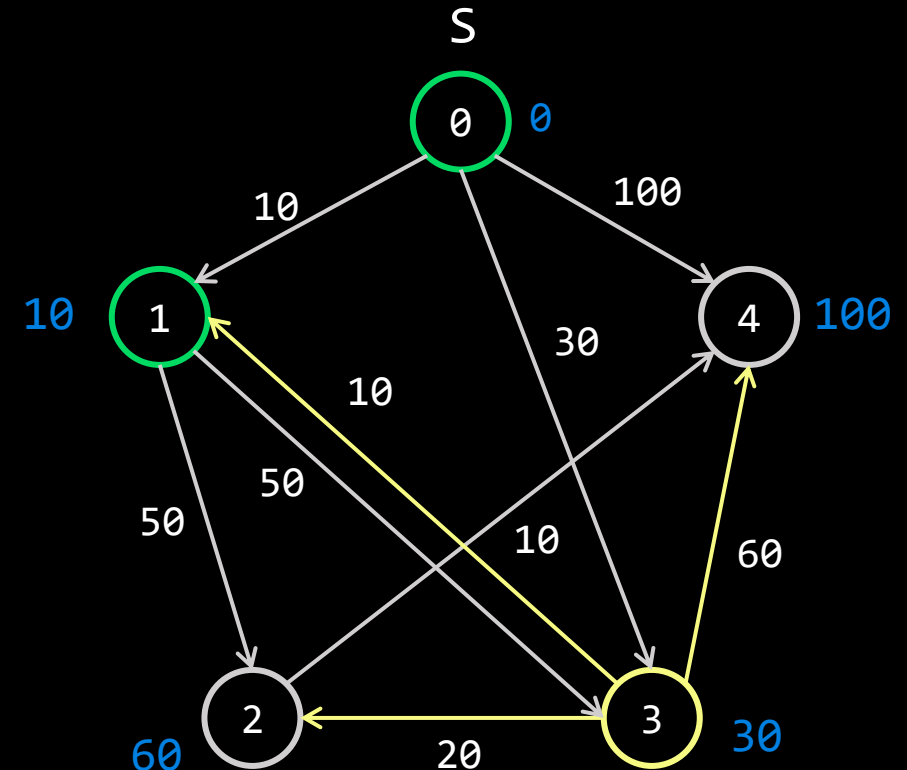
Dijkstra's Shortest Path Algorithm

Example: Process edges adjacent to the vertex 3 and update distances based on relaxation

Computed, $S = \{0, 1, 3\}$

Needs processing, $V-S = \{2, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	60	1
3	30	0
4	100	0



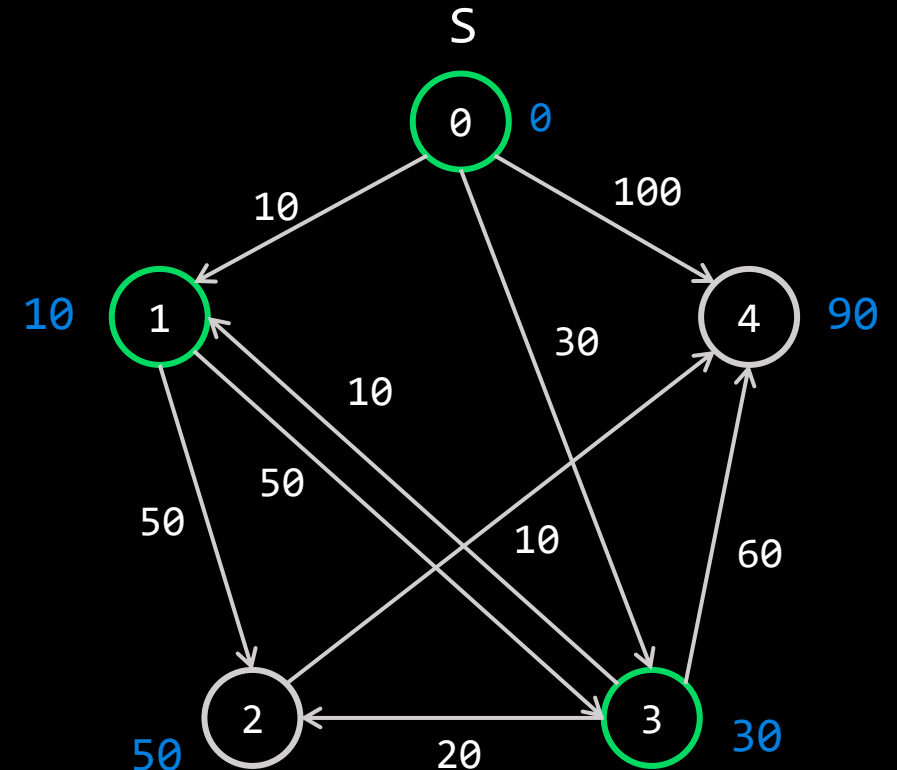
Dijkstra's Shortest Path Algorithm

Example: 3 is now done

Computed, $S = \{0, 1, 3\}$

Needs processing, $V-S = \{2, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	90	3



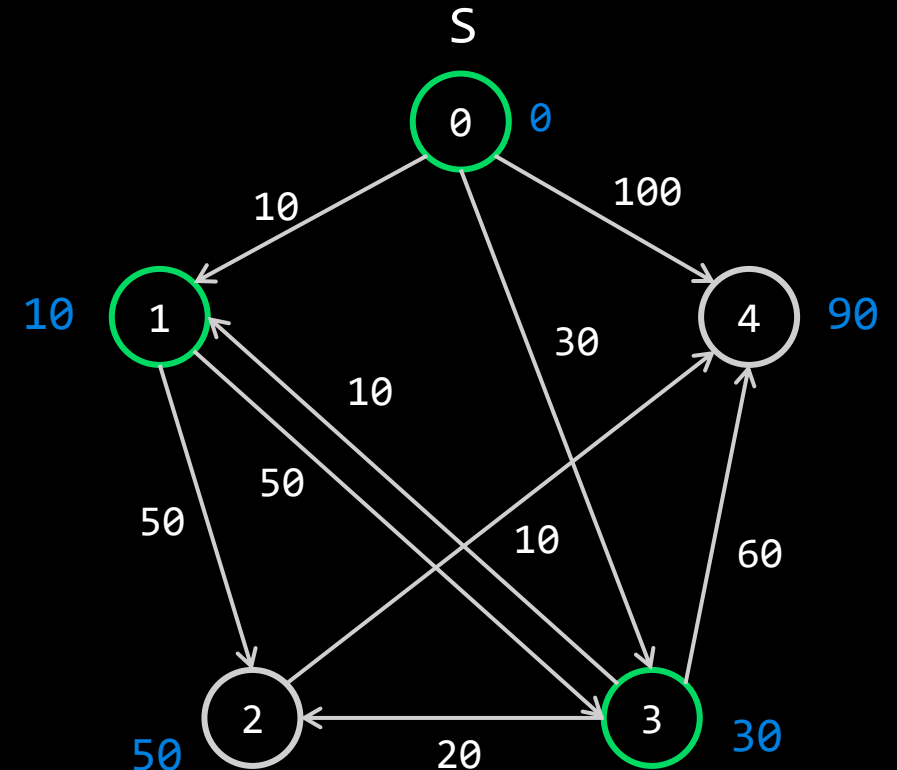
Dijkstra's Shortest Path Algorithm

Example: Next, repeat the process picking the minimum element in $d[v]$ that has not been computed

Computed, $S = \{0, 1, 3\}$

Needs processing, $V-S = \{2, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	90	3



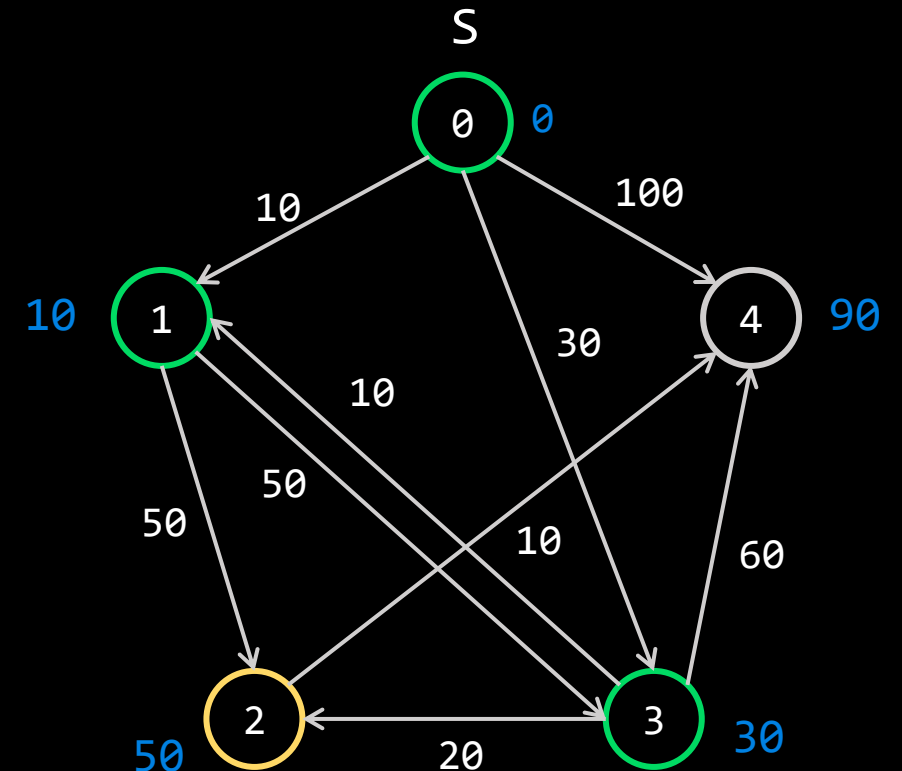
Dijkstra's Shortest Path Algorithm

Example: Pick 2

Computed, $S = \{0, 1, 3\}$

Needs processing, $V-S = \{2, 4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	90	3



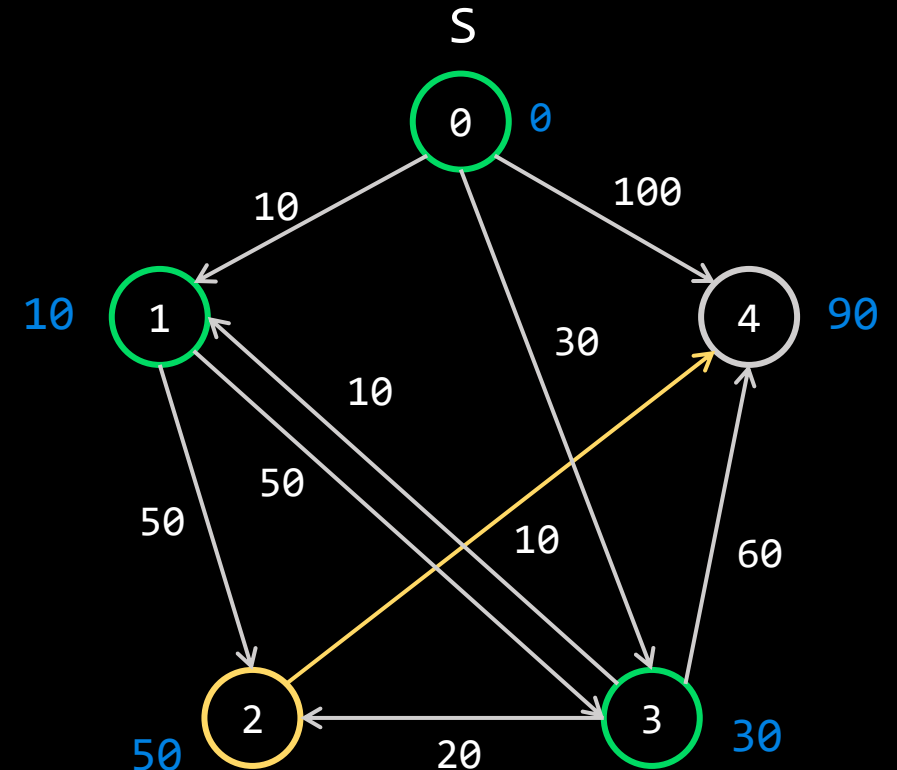
Dijkstra's Shortest Path Algorithm

Example: Process edges adjacent to the vertex 2 and update distances based on relaxation

Computed, $S = \{0, 1, 2, 3\}$

Needs processing, $V-S = \{4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	90	3



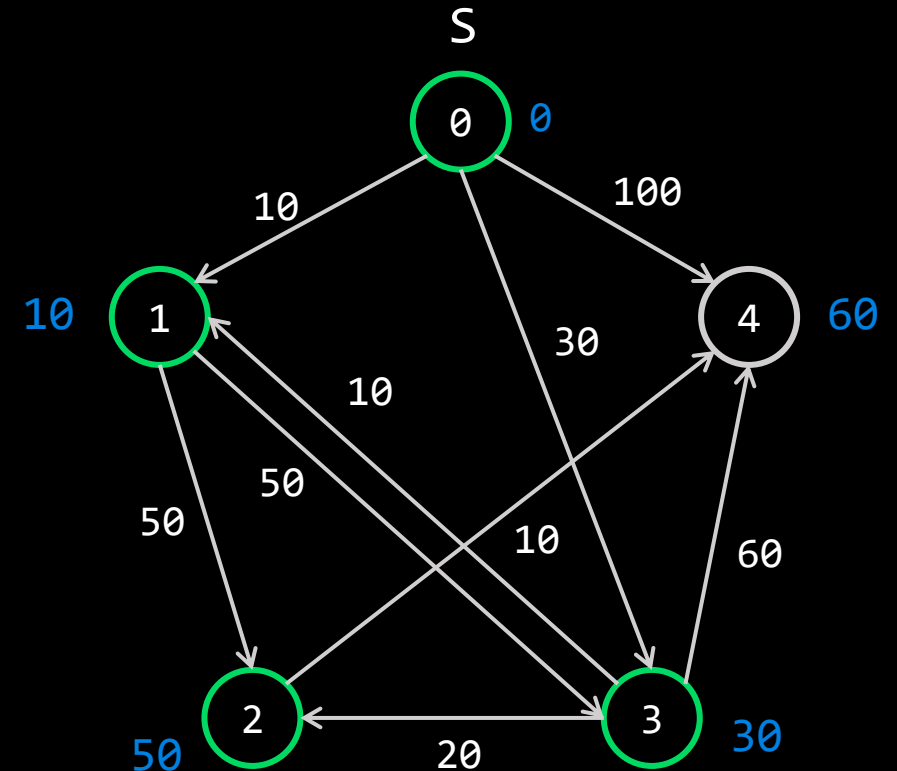
Dijkstra's Shortest Path Algorithm

Example: 2 is now done

Computed, $S = \{0, 1, 2, 3\}$

Needs processing, $V-S = \{4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	60	2



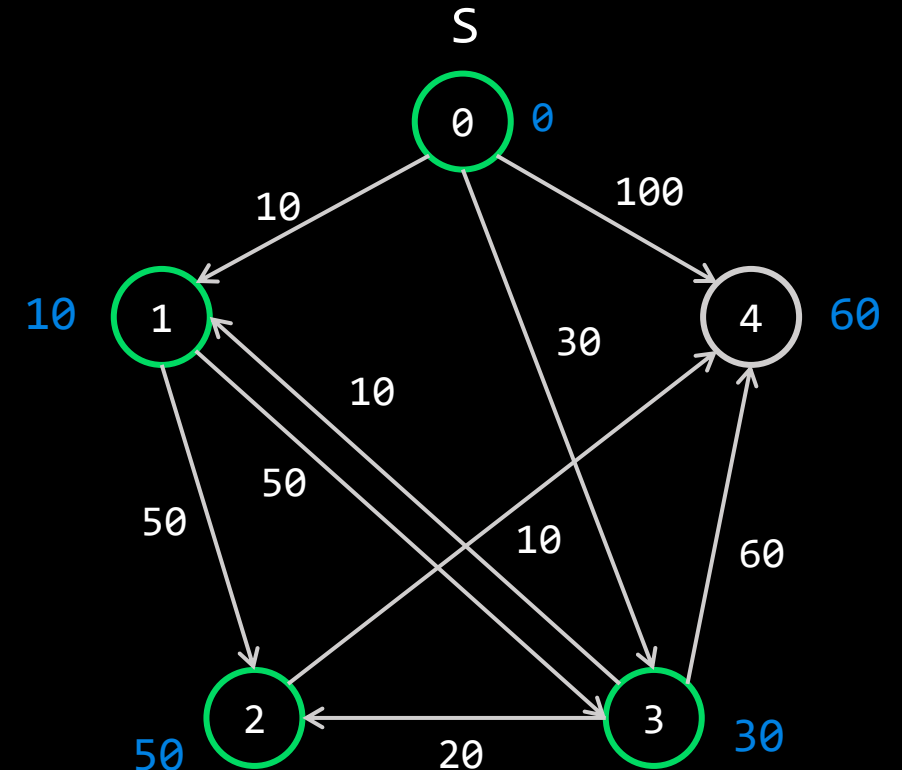
Dijkstra's Shortest Path Algorithm

Example: Next, repeat the process picking the minimum element in $d[v]$ that has not been computed

Computed, $S = \{0, 1, 2, 3\}$

Needs processing, $V-S = \{4\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	60	2



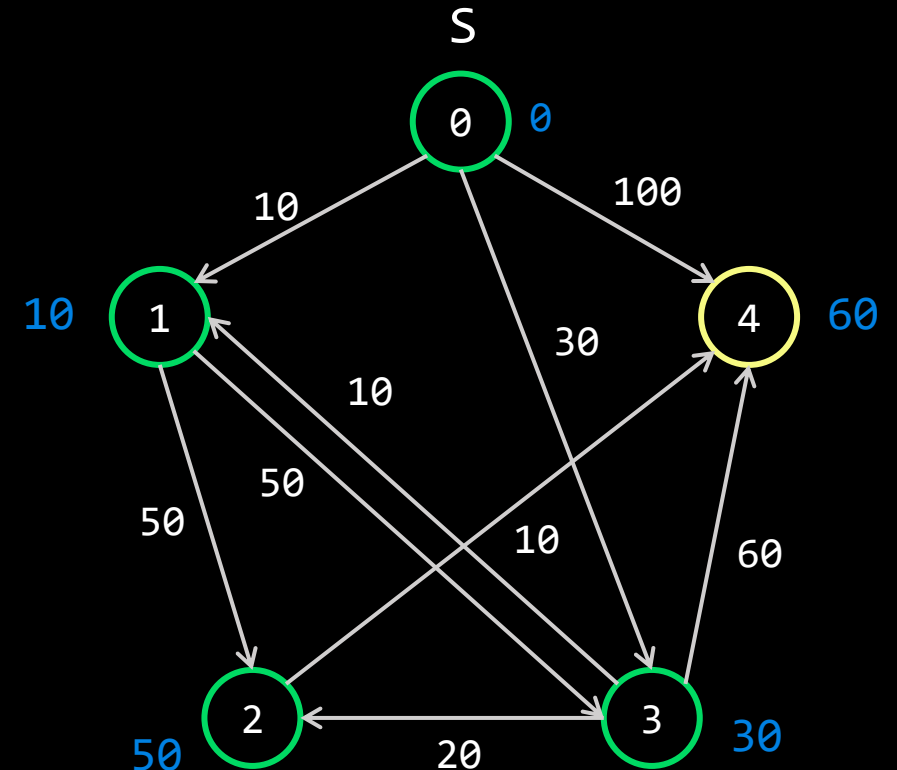
Dijkstra's Shortest Path Algorithm

Example: Pick 4. Process edges adjacent to the vertex 4 and update distances based on relaxation

Computed, $S = \{0, 1, 2, 3, 4\}$

Needs processing, $V-S = \{\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	60	2



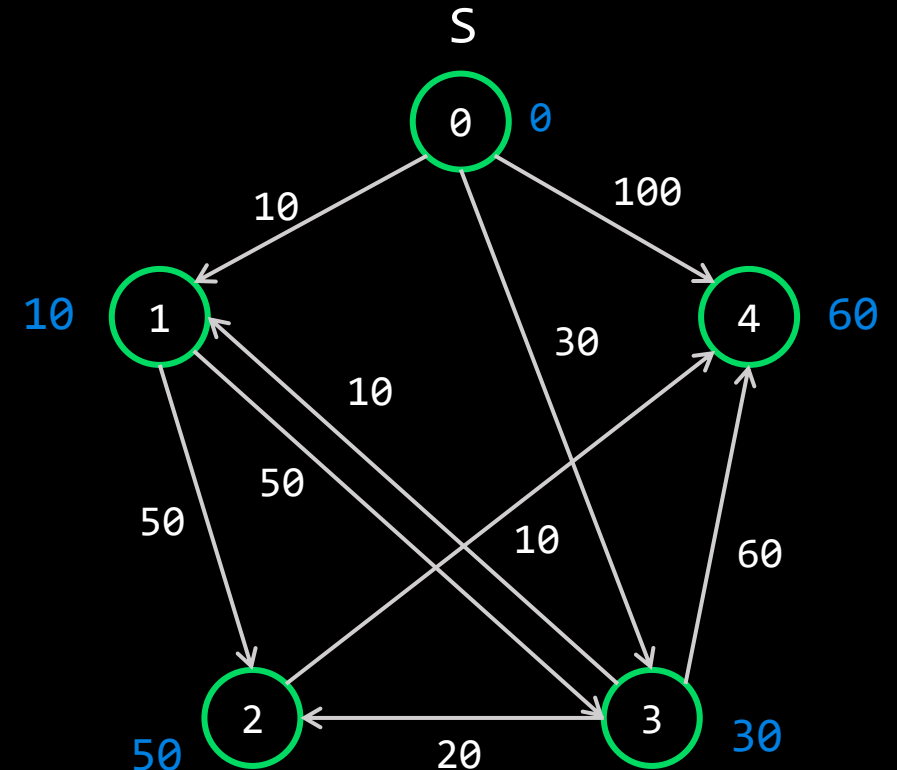
Dijkstra's Shortest Path Algorithm

Example: 4 is now done and V-S is empty. Stop.

Computed, $S = \{0, 1, 2, 3, 4\}$

Needs processing, $V-S = \{\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	60	2



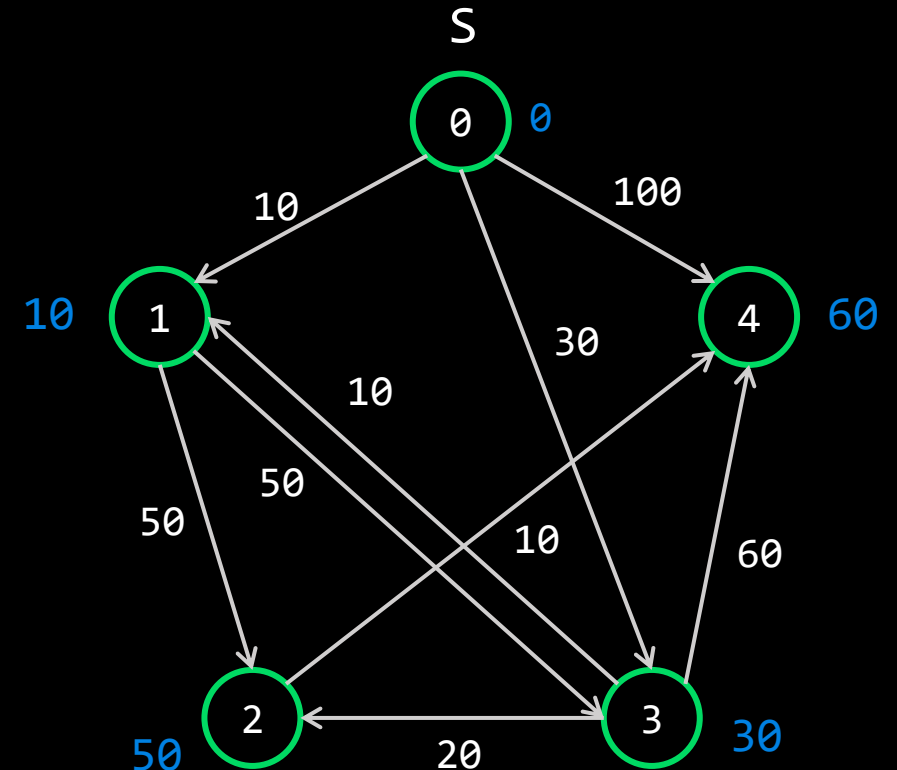
Dijkstra's Shortest Path Algorithm

Example: 4 is now done and V-S is empty. Stop.

Computed, $S = \{0, 1, 2, 3, 4\}$

Needs processing, $V-S = \{\}$

v	$d[v]$	$p[v]$
0	0	-1
1	10	0
2	50	3
3	30	0
4	60	2



Path from 0 to 4: 0 3 2 4 (Cost: 60)

Dijkstra's Shortest Path Algorithm

Dijkstra's Algorithm

1. Initialize S with the start vertex, s , and $V-S$ with the remaining vertices.
2. **for** all v in $V-S$
3. Set $p[v]$ to s .
4. **if** there is an edge (s, v)
5. Set $d[v]$ to $w(s, v)$.
6. **else**
7. Set $d[v]$ to ∞ .
8. **while** $V-S$ is not empty
9. **for** all u in $V-S$, find the smallest $d[u]$.
10. Remove u from $V-S$ and add u to S .
11. **for** all v adjacent to u in $V-S$
12. **if** $d[u] + w(u, v)$ is less than $d[v]$.
13. Set $d[v]$ to $d[u] + w(u, v)$.
14. Set $p[v]$ to u .

Dijkstra's Shortest Path Algorithm

Dijkstra's:

```
PQ.add(source, 0)
For other vertices v, PQ.add(v, infinity)
While PQ is not empty:
    p = PQ.removeSmallest()
    Relax all edges from p
```

Relaxing an edge $u \rightarrow v$ with weight w :

```
If  $d[u] + w < d[v]$ :
     $d[v] = d[u] + w$ 
     $p[v] = u$ 
    PQ.changePriority(v,  $d[v]$ )
```

Dijkstra's Shortest Path Algorithm

Dijkstra's:

```
PQ.add(source, 0)
For other vertices v, PQ.add(v, infinity)
While PQ is not empty:
    p = PQ.removeSmallest()
    Relax all edges from p
```

$O(V \log V)$
 $O(V \log V)$

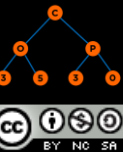
Relaxing an edge $u \rightarrow v$ with weight w :

```
If  $d[u] + w < d[v]$ :
     $d[v] = d[u] + w$ 
     $p[v] = u$ 
    PQ.changePriority(v,  $d[v]$ )
```

$O(E \log V)$

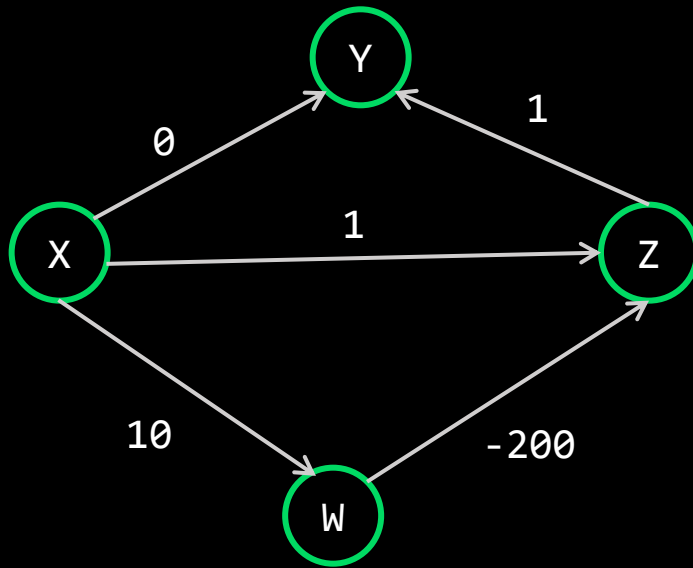
Dijkstra's Properties

- **Greedy Algorithm:** Visit vertices in order of best-known distance from source. On visit, relax every edge from the visited vertex
- Dijkstra's is **guaranteed** to return a **correct result** if all edges are non-negative.
- Dijkstra's is **guaranteed** to be **optimal** so long as there are no negative edges.
- Overall **runtime:** $O(V \cdot \log(V) + V \cdot \log(V) + E \cdot \log V)$.
 - Assuming $E > V$, this is just $O(E \log V)$ for a connected graph.



Dijkstra's Properties

▪ Negative Edges Example:



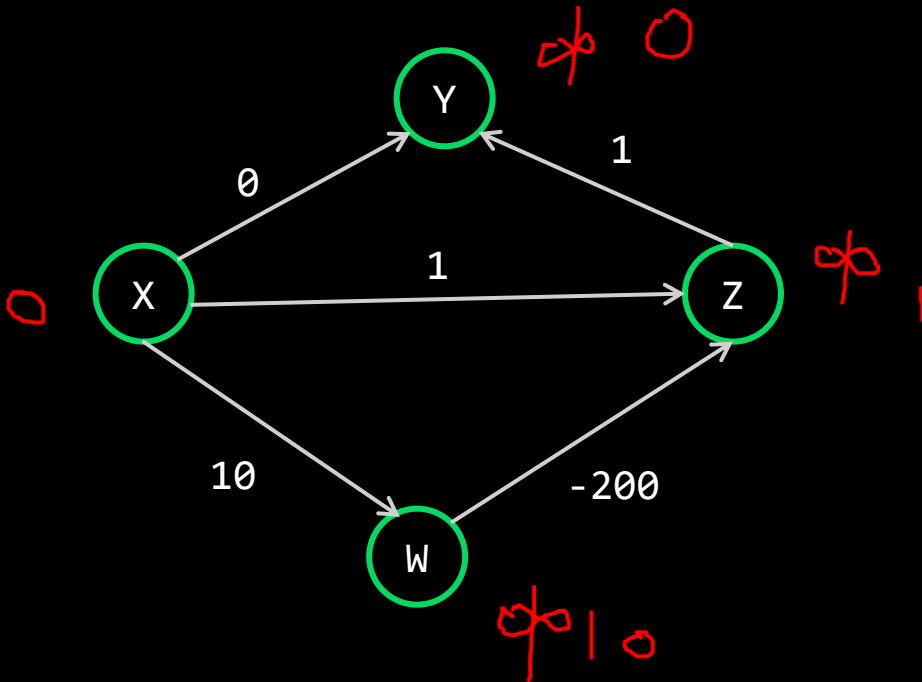
v	$d[v]$	$p[v]$
X		
Y		
Z		
W		

<https://stackoverflow.com/questions/6799172/negative-weights-using-dijkstras-algorithm/6799344#6799344>

<https://stackoverflow.com/questions/6799172/negative-weights-using-dijkstras-algorithm>

Dijkstra's Properties

▪ Negative Edges Example:

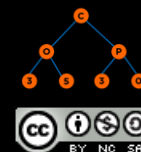


v	d[v]	p[v]
✓ X	0	-1
✓ Y	0	X X
✓ Z	X -190	X X W
✓ W	10	X X

But shortest path between X to Y should be -189;
Once at Z we assume the current distance is the
shortest distance to reach Z at that point

<https://stackoverflow.com/questions/6799172/negative-weights-using-dijkstras-algorithm/6799344#6799344>
<https://stackoverflow.com/questions/6799172/negative-weights-using-dijkstras-algorithm>

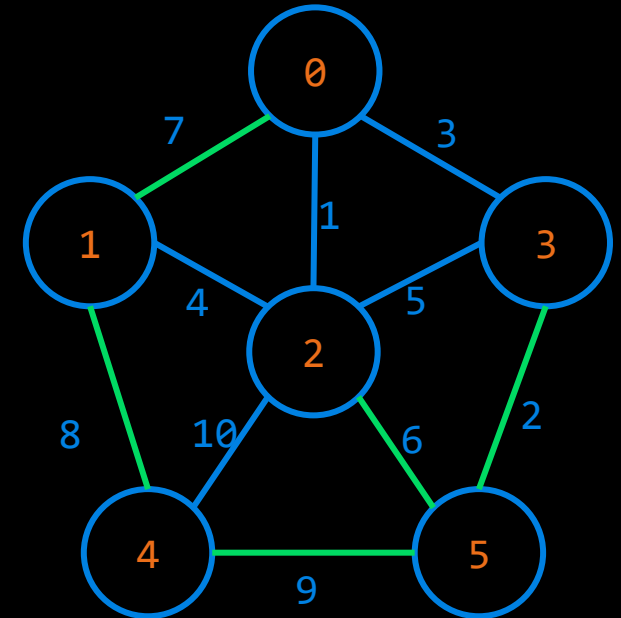
Questions



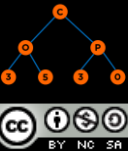
Minimum Spanning Tree

Spanning Tree

- A spanning tree is a subset of the edges of a graph such that there is only one edge between each vertex, and all of the vertices are connected. The tree is connected and acyclic.
- The cost of a spanning tree is the sum of the weights of the edges.
- Minimum spanning tree is the spanning tree with the smallest cost.
- Spanning tree with N vertices will have $N-1$ edges.
- Used in networks, laying wires for electricity/telephones, routing for internet connections, etc.

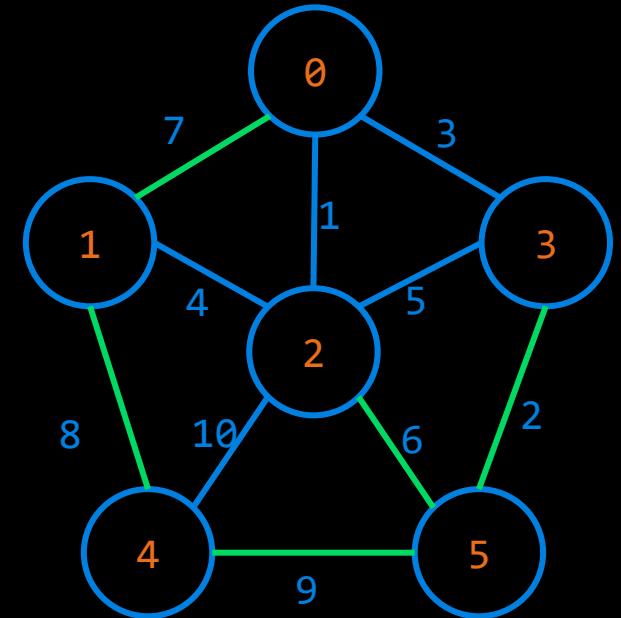


Minimum Spanning Tree – Prim's

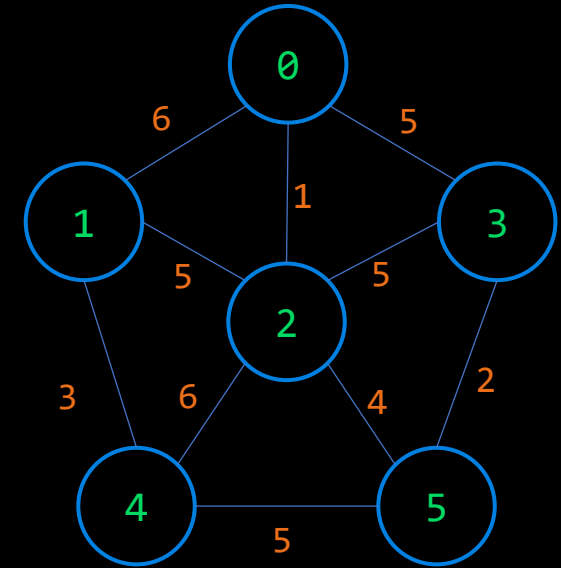
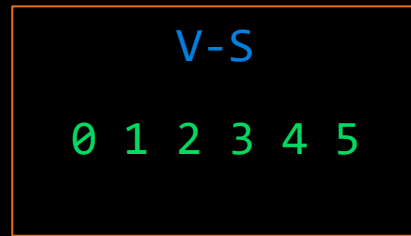
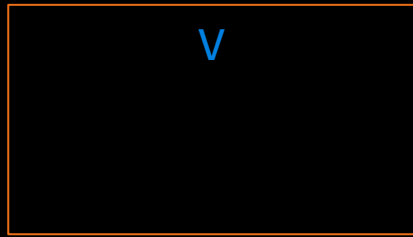


Prim's Algorithm

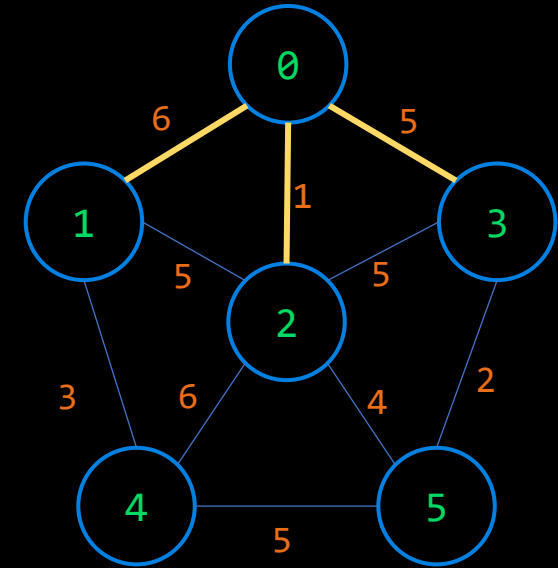
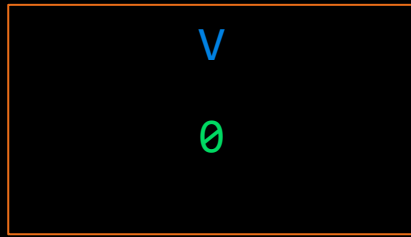
- Prim's algorithm analyzes all the connections between vertices and finds the set with minimum total weight that makes the graph connected.
- The vertices are divided into two sets:
 - S , the set of vertices in the spanning tree
 - And $V-S$, the remaining vertices
- Next, we choose the edge with the smallest weight that connects a vertex in S to a vertex in $V-S$ and add it to the minimum spanning tree.



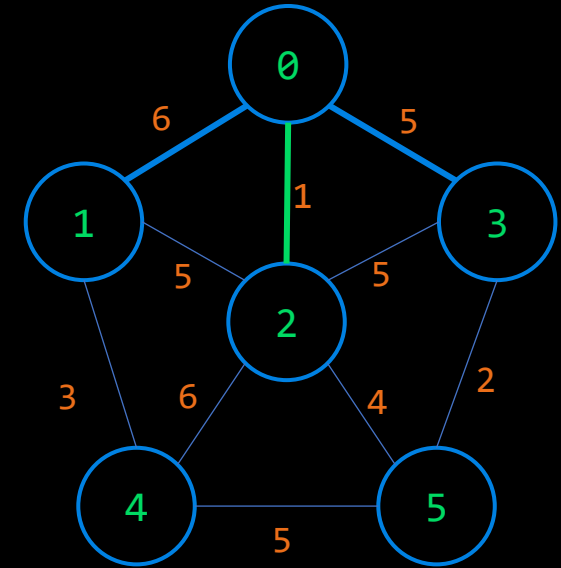
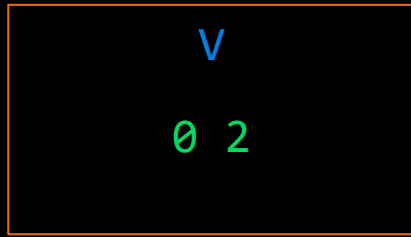
Prim's Algorithm



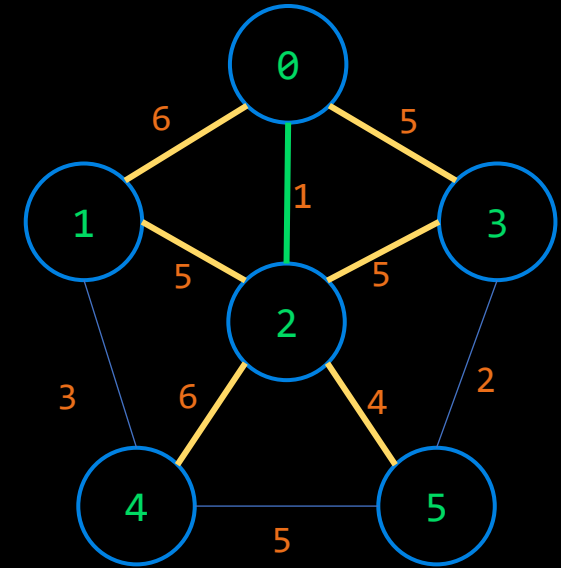
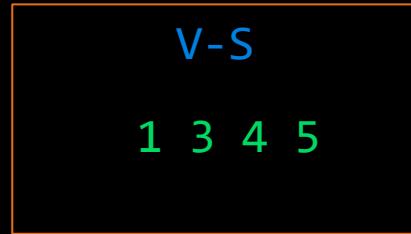
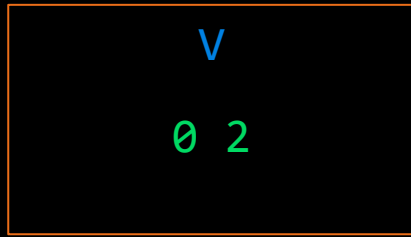
Prim's Algorithm



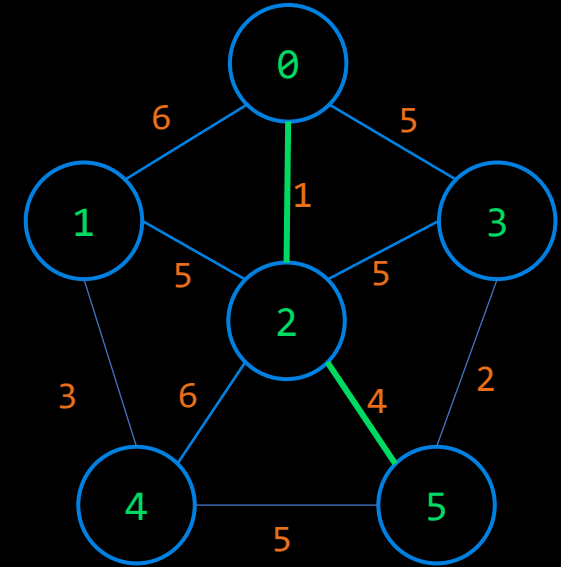
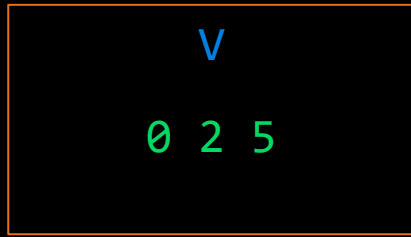
Prim's Algorithm



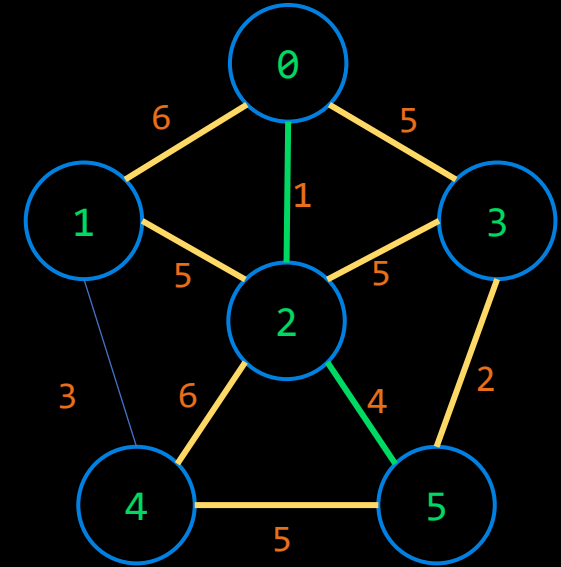
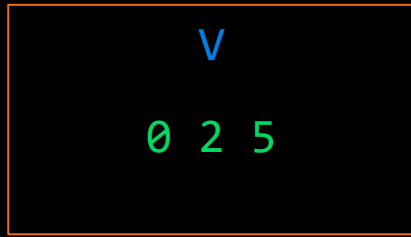
Prim's Algorithm



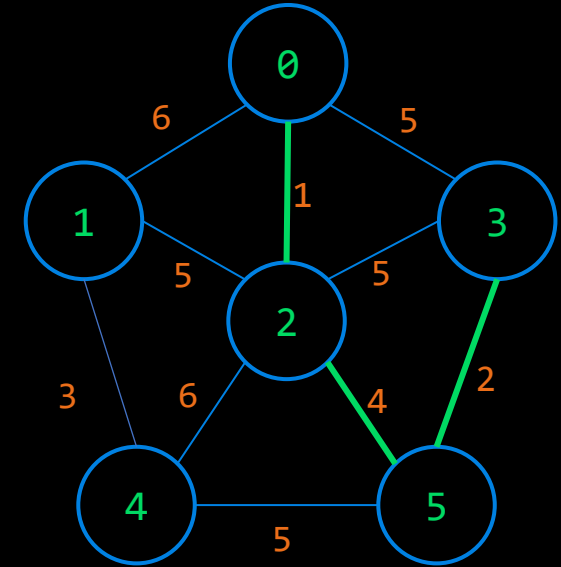
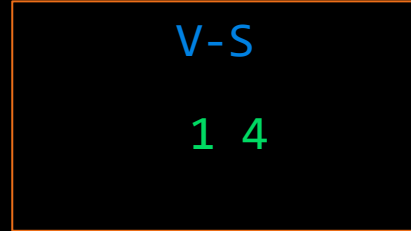
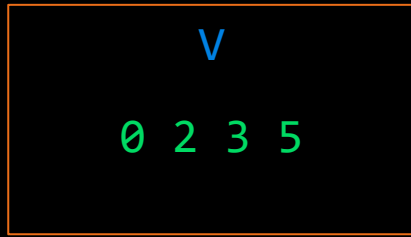
Prim's Algorithm



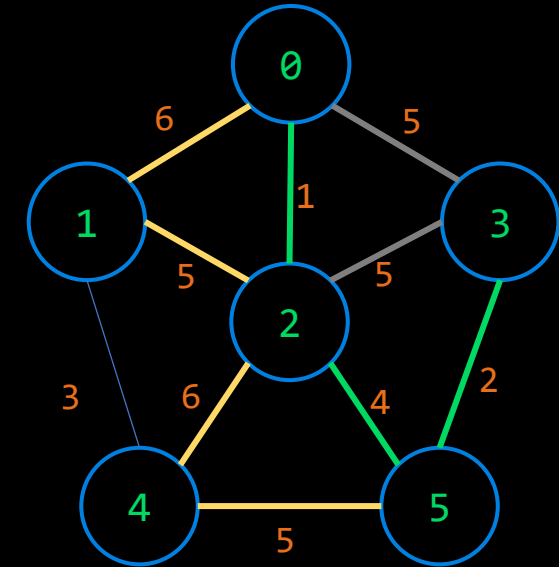
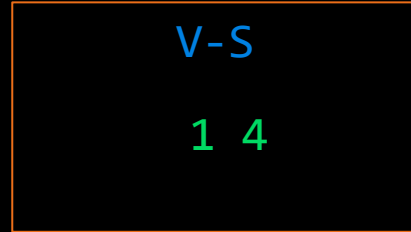
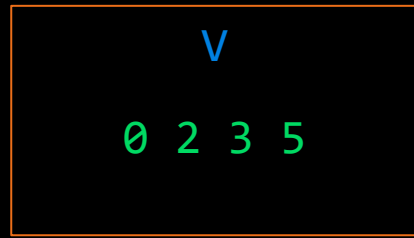
Prim's Algorithm



Prim's Algorithm

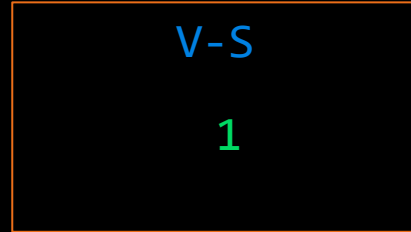
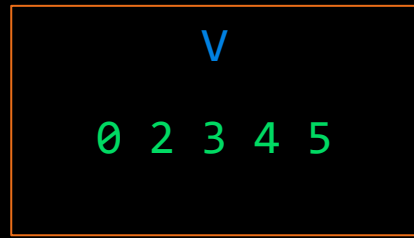


Prim's Algorithm

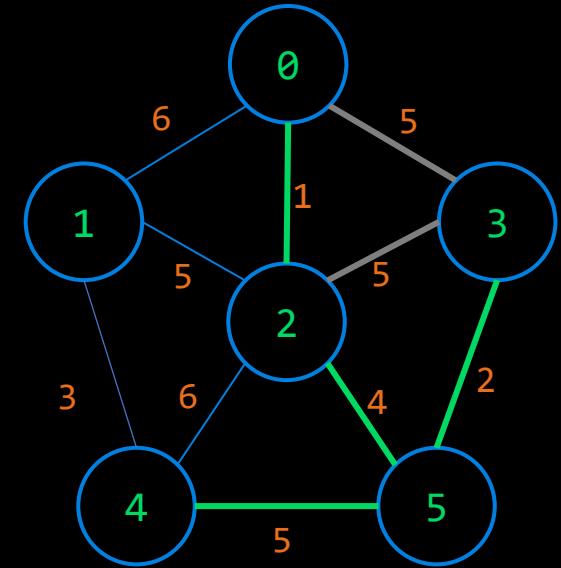


We choose the edge with the smallest weight that connects a vertex in S to a vertex in V-S and add it to the minimum spanning tree. Option to pick 1-2 or 4-5. Pick any.

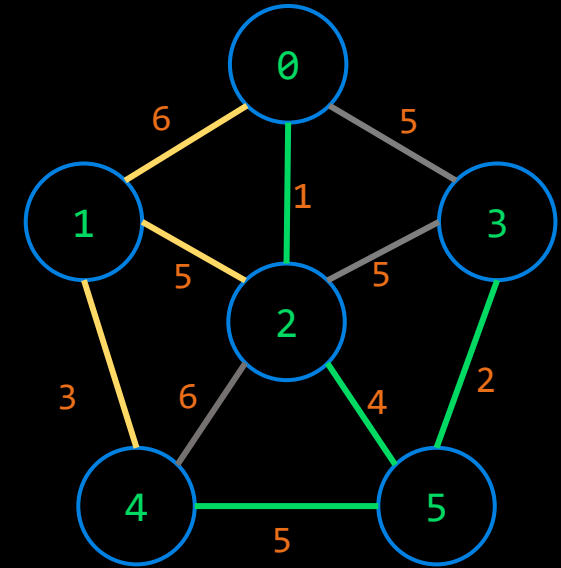
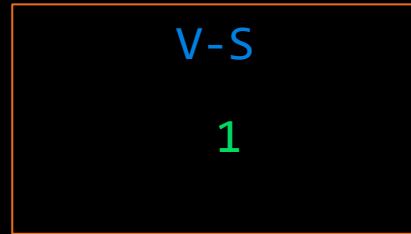
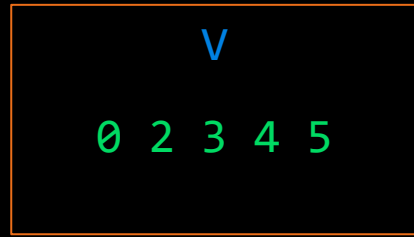
Prim's Algorithm



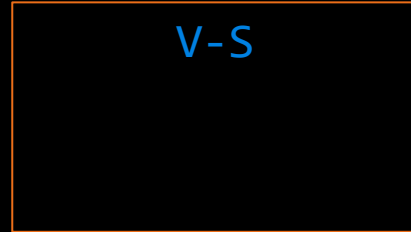
Pick 4-5.



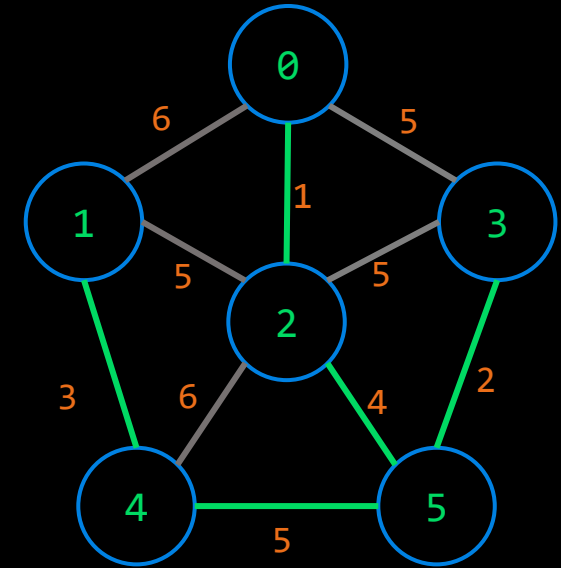
Prim's Algorithm



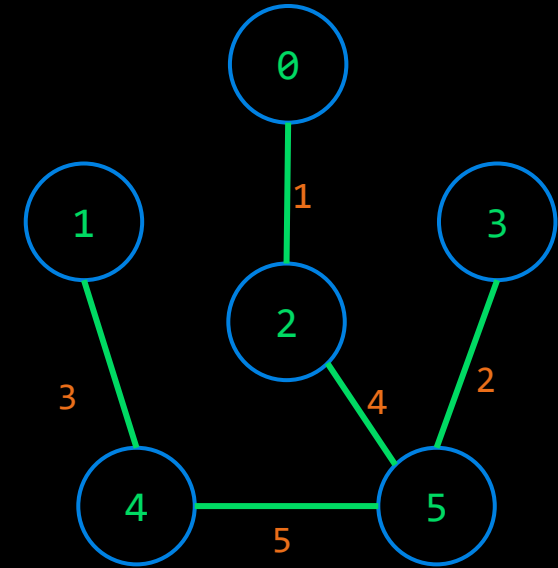
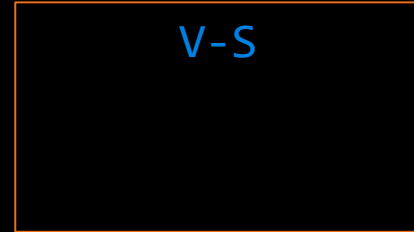
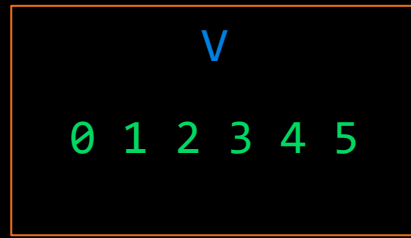
Prim's Algorithm



Pick 1-4.



Prim's Algorithm



Sum of MST = 15.

Prim's Algorithm

Input: An undirected, connected, weighted graph G .

Output: T , a minimum spanning tree for G .

$T := \emptyset$.

Pick any vertex in G and add it to T .

For $j = 1$ to $n-1$

 Let C be the set of edges with one endpoint inside T and one endpoint outside T .

 Let e be a minimum weight edge in C .

 Add e to T .

 Add the endpoint of e not already in T to T .

End-for

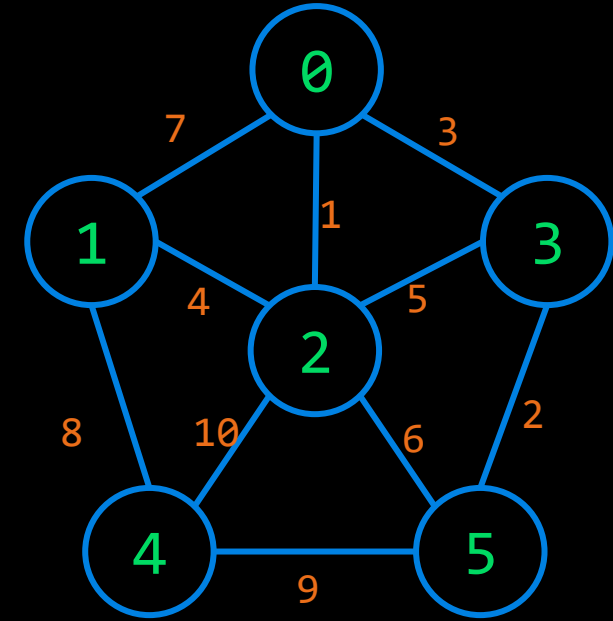
Complexity: $O(EV)$ or $O(E \log V)$ - using priority queues

Minimum Spanning Tree – Kruskal's

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
3-5	2
0-3	3
1-2	4
2-3	5
2-5	6
0-1	7
1-4	8
4-5	9
2-4	10

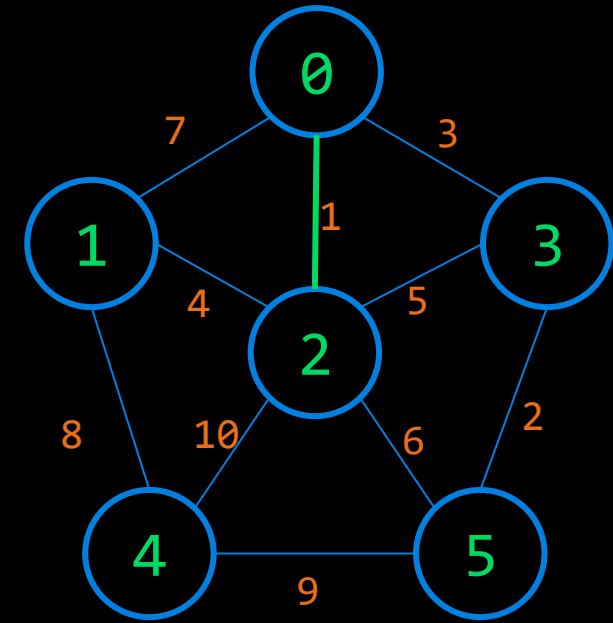


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
3-5	2
0-3	3
1-2	4
2-3	5
2-5	6
0-1	7
1-4	8
4-5	9
2-4	10

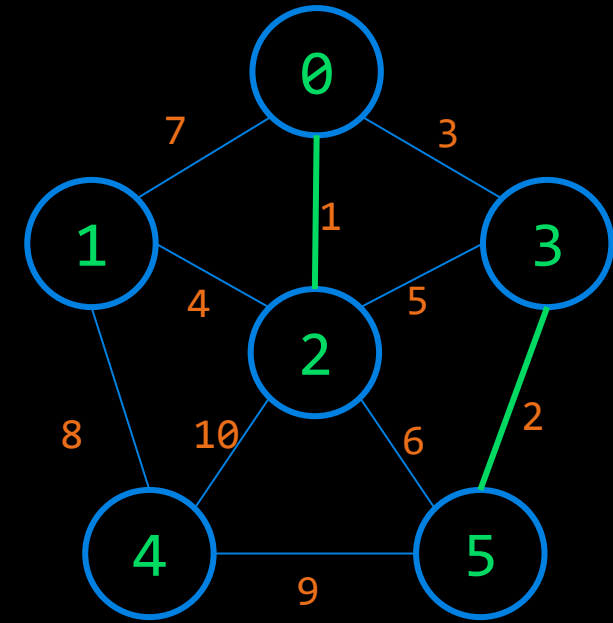


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0-2	1
3-5	2
0-3	3
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1-4	8
4-5	9
2-4	10

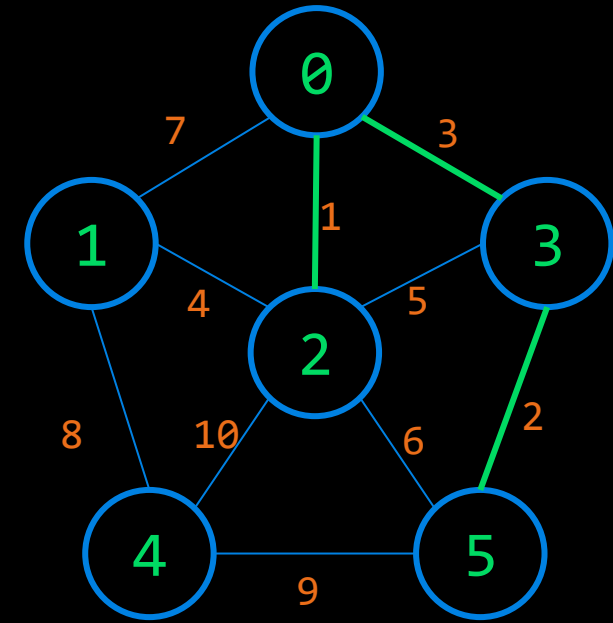


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0-2	1
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1-4	8
4-5	9
2-4	10

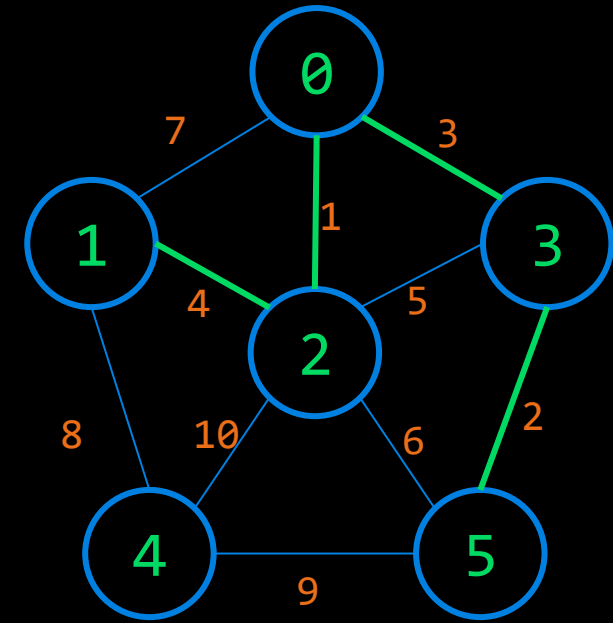


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
3-5	2
0-3	3
1-2	4
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2-5	6
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1-4	8
4-5	9
2-4	10

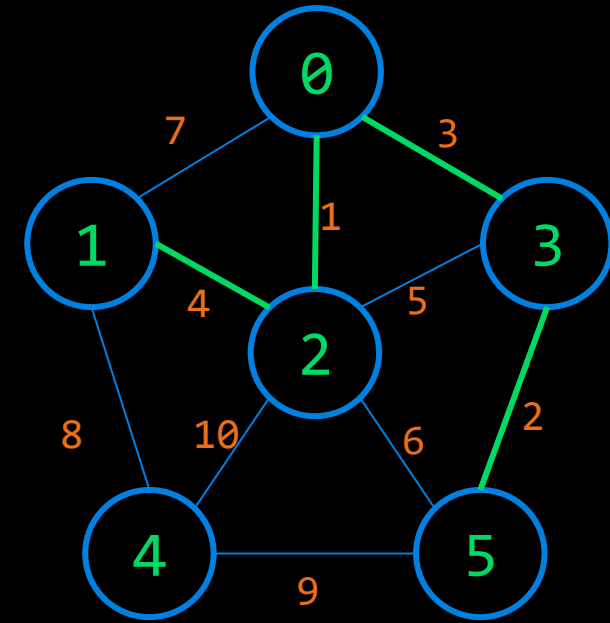


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
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0-3	3
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4-5	9
2-4	10

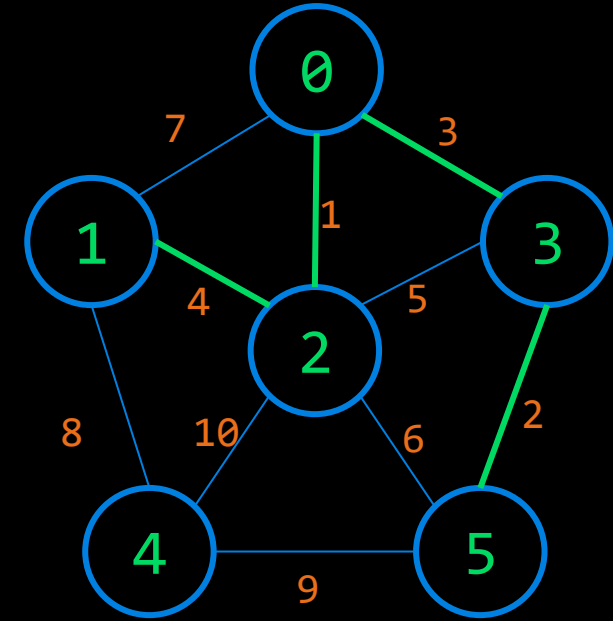


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
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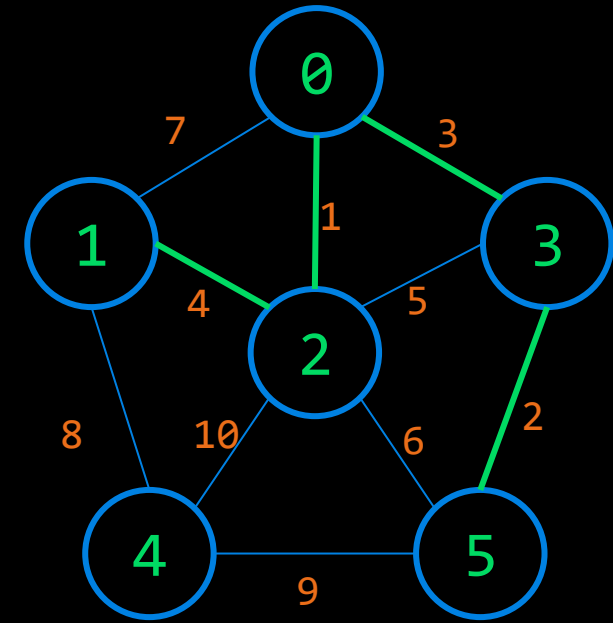


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
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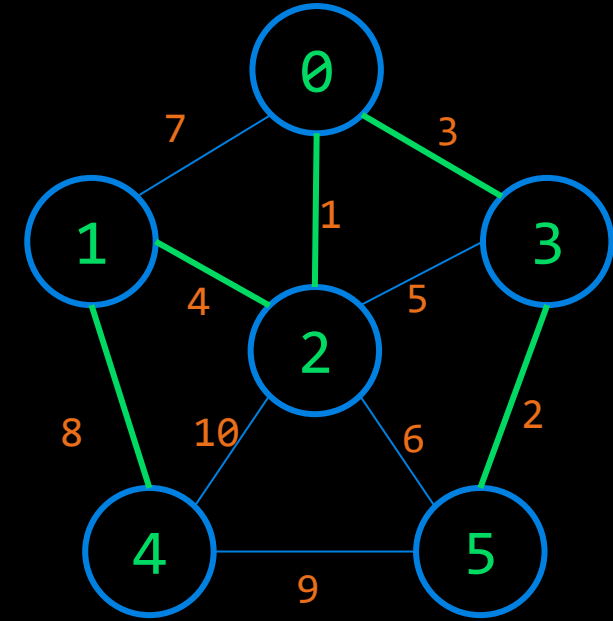


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
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0-3	3
1-2	4
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4-5	9
2-4	10

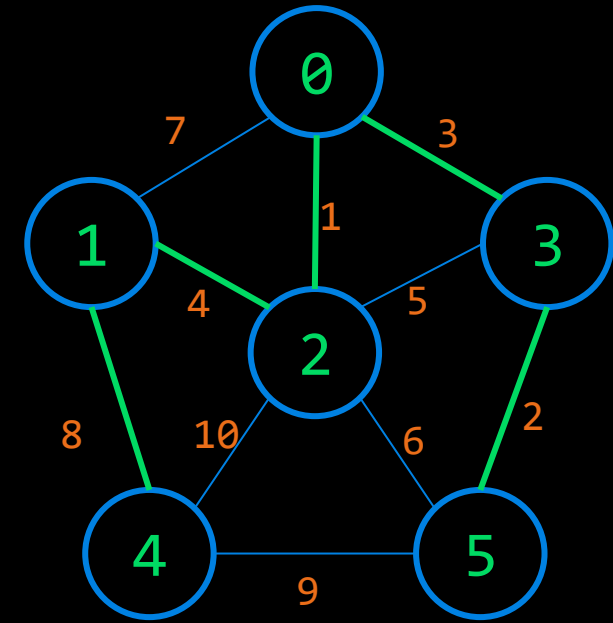


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
3-5	2
0-3	3
1-2	4
2-3	5
2-5	6
0-1	7
1-4	8
4-5	9
2-4	10

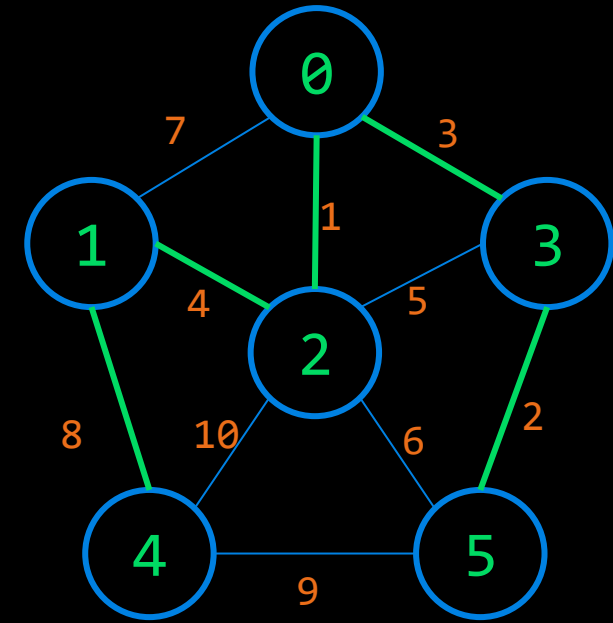


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

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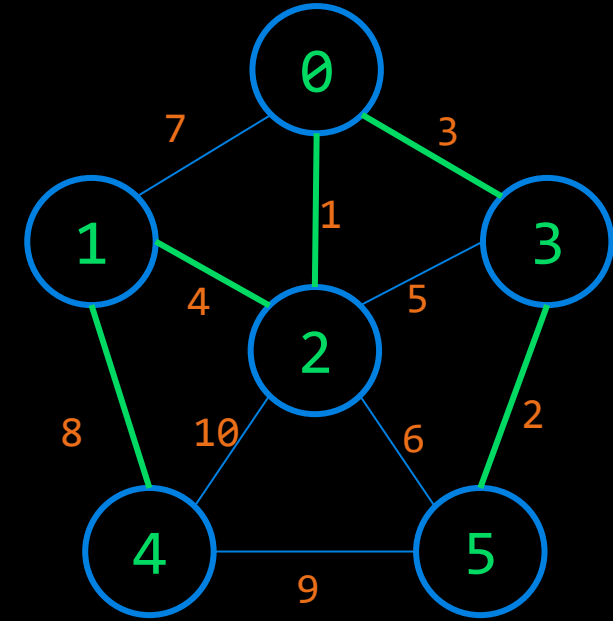
Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

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3-5	2
0-3	3
1-2	4
2-3	5
2-5	6
0-1	7
1-4	8
4-5	9
2-4	10

Minimum Spanning Tree Sum = 18



Questions

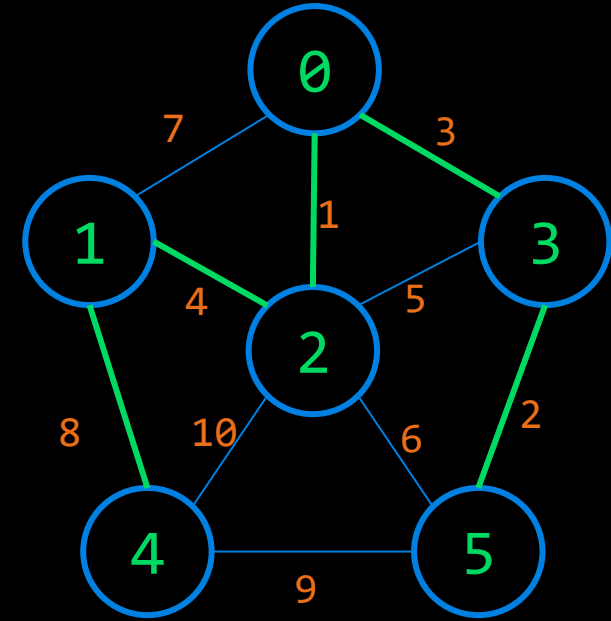
Kruskal's Algorithm

How can we detect a cycle when adding an edge?

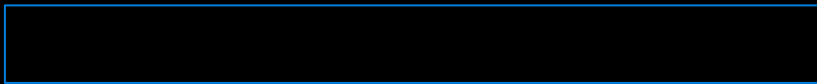
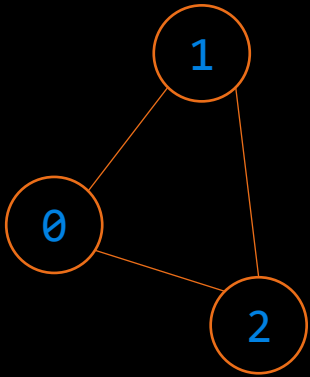
Method 1:

Cycle Detection using DFS. Find back edges.

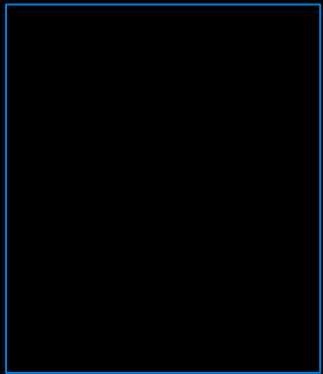
Back Edge: An edge that connects an ancestor during DFS traversal.



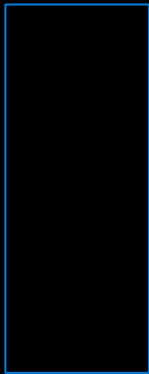
7.3.1 Detect whether there is a Cycle in an Undirected Graph



parent



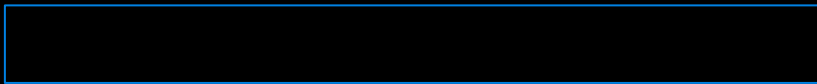
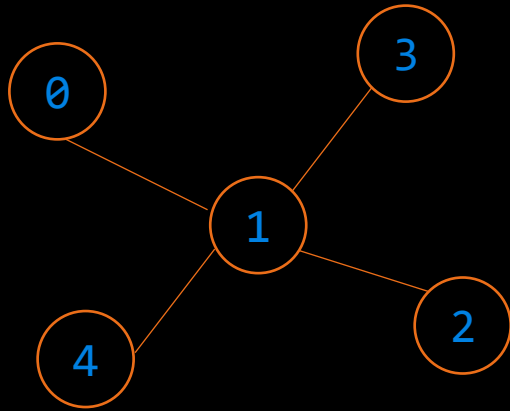
visited



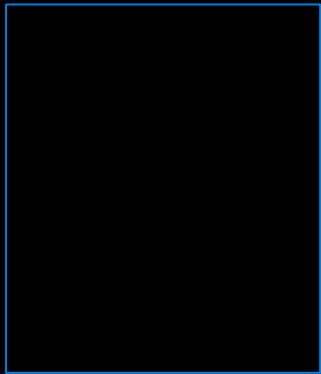
s

```
1.  bool anyCycle(const Graph& graph)
2.  {
3.      set<int> visited;
4.      vector<int> parent(graph.numVertices, -1);
5.      stack<int> s;
6.      visited.insert(0);
7.      s.push(0);
8.      while(!s.empty())
9.      {
10.         int u = s.top();
11.         s.pop();
12.         for(auto v: graph.adjList[u])
13.         {
14.             if ((visited.find(v)==visited.end()))
15.             {
16.                 visited.insert(v);
17.                 s.push(v);
18.                 parent[v] = u;
19.             }
20.             else if (parent[u] != v)
21.                 return true;
22.         }
23.     }
24.     return false;
25. }
```


7.3.1 Detect whether there is a Cycle in an Undirected Graph



parent



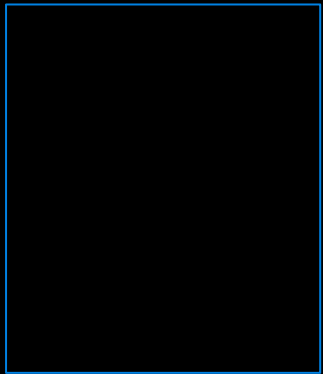
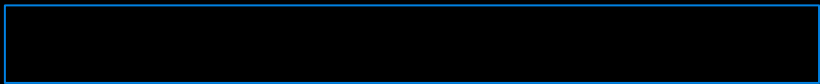
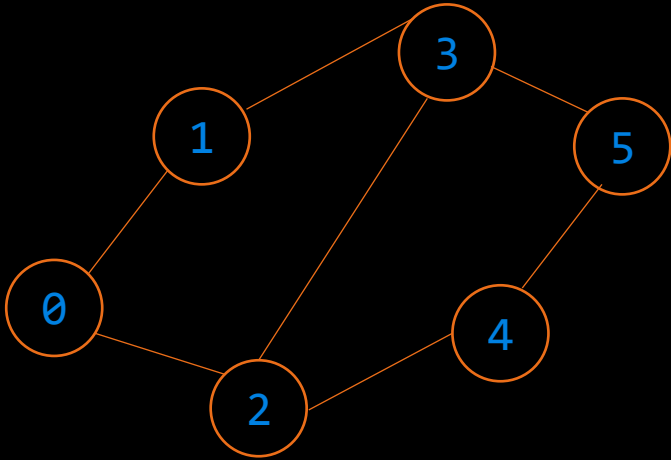
visited



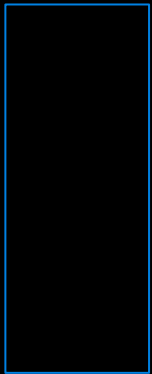
s

```
1.  bool anyCycle(const Graph& graph)
2.  {
3.      set<int> visited;
4.      vector<int> parent(graph.numVertices, -1);
5.      stack<int> s;
6.      visited.insert(0);
7.      s.push(0);
8.      while(!s.empty())
9.      {
10.         int u = s.top();
11.         s.pop();
12.         for(auto v: graph.adjList[u])
13.         {
14.             if ((visited.find(v)==visited.end()))
15.             {
16.                 visited.insert(v);
17.                 s.push(v);
18.                 parent[v] = u;
19.             }
20.             else if (parent[u] != v)
21.                 return true;
22.         }
23.     }
24.     return false;
25. }
```

7.3.1 Detect whether there is a Cycle in an Undirected Graph



visited

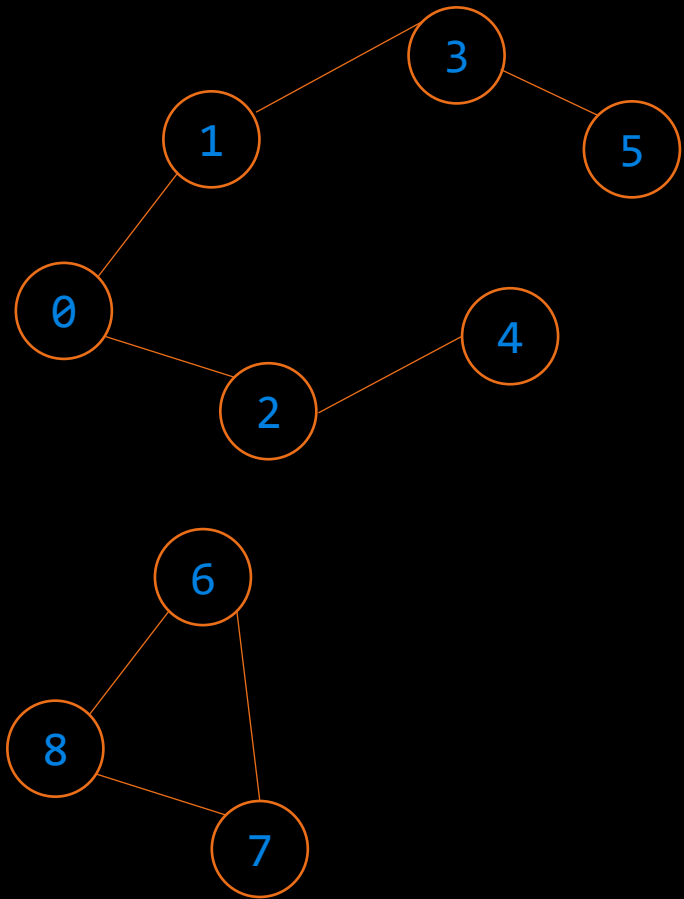


s

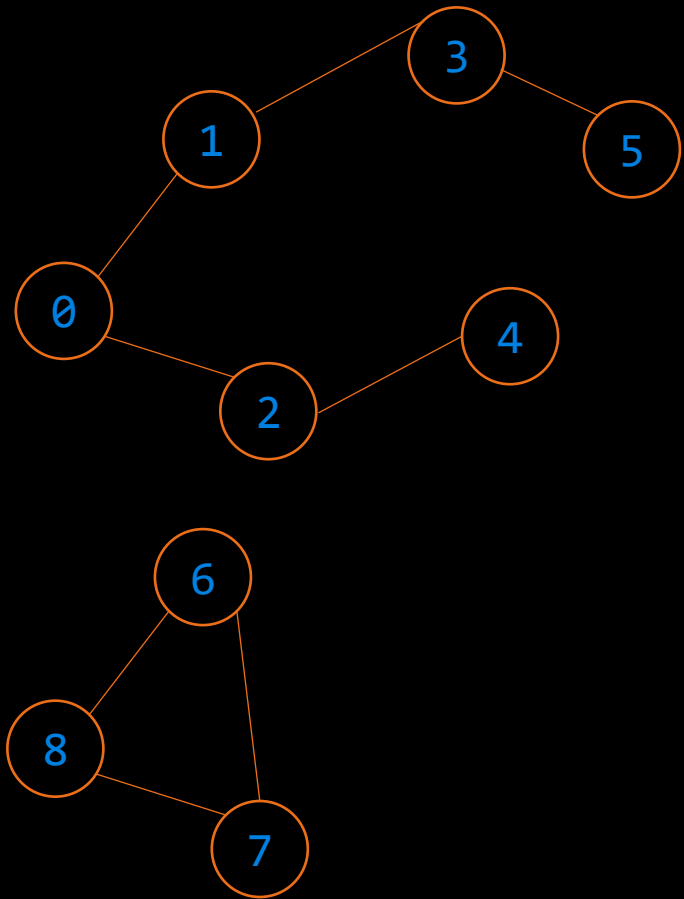
parent

```
1.  bool anyCycle(const Graph& graph)
2.  {
3.      set<int> visited;
4.      vector<int> parent(graph.numVertices, -1);
5.      stack<int> s;
6.      visited.insert(0);
7.      s.push(0);
8.      while(!s.empty())
9.      {
10.         int u = s.top();
11.         s.pop();
12.         for(auto v: graph.adjList[u])
13.         {
14.             if ((visited.find(v)==visited.end()))
15.             {
16.                 visited.insert(v);
17.                 s.push(v);
18.                 parent[v] = u;
19.             }
20.             else if (parent[u] != v)
21.                 return true;
22.         }
23.     }
24.     return false;
25. }
```

7.3.1 Detect whether there is a Cycle in an Undirected Graph



7.3.1 Detect whether there is a Cycle in an Undirected Graph



```
1.  bool anyCycle(const Graph& graph)
2.  {
3.      set<int> visited;
4.      vector<int> parent(graph.numVertices, -1);
5.      stack<int> s;
6.      for(int i=0; i<graph.numVertices; i++)
7.      {
8.          if ((visited.find(i)==visited.end()))
9.          {
10.             visited.insert(i);
11.             s.push(i);
12.             while(!s.empty())
13.             {
14.                 int u = s.top();
15.                 s.pop();
16.                 for(auto v: graph.adjList[u])
17.                 {
18.                     if ((visited.find(v)==visited.end()))
19.                     {
20.                         visited.insert(v);
21.                         s.push(v);
22.                         parent[v] = u;
23.                     }
24.                     else if (parent[u] != v)
25.                         return true;
26.                 }
27.             }
28.         }
29.     }
30.     return false;
31. }
```

Kruskal's Algorithm

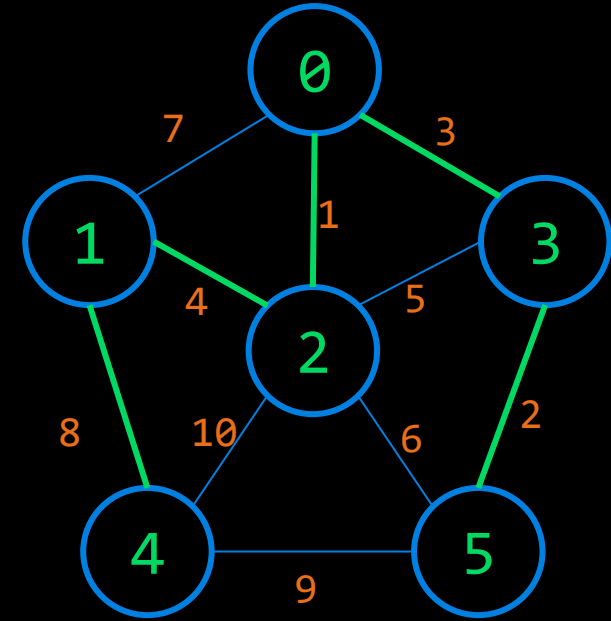
How can we detect a cycle when adding an edge?

Method 1:

Cycle Detection using DFS.

Works correctly but is computationally more expensive.

Complexity: $O(E (V+E))$

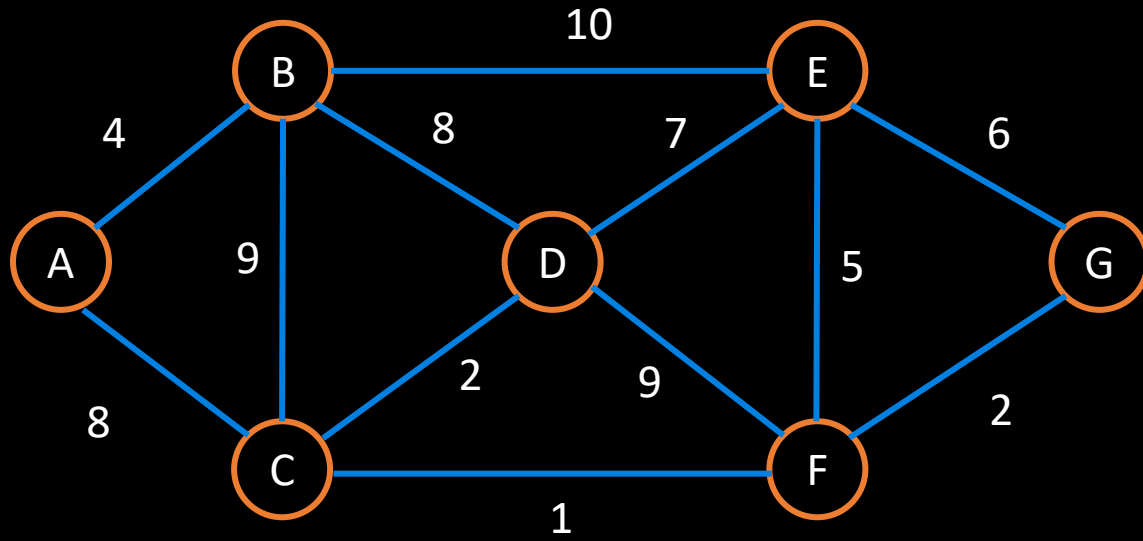


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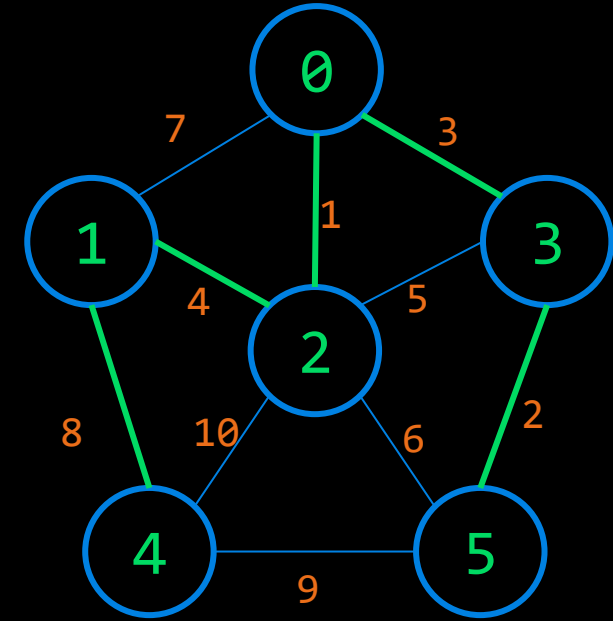


Kruskal's Algorithm

How can we detect a cycle when adding an edge?

Method 2a:

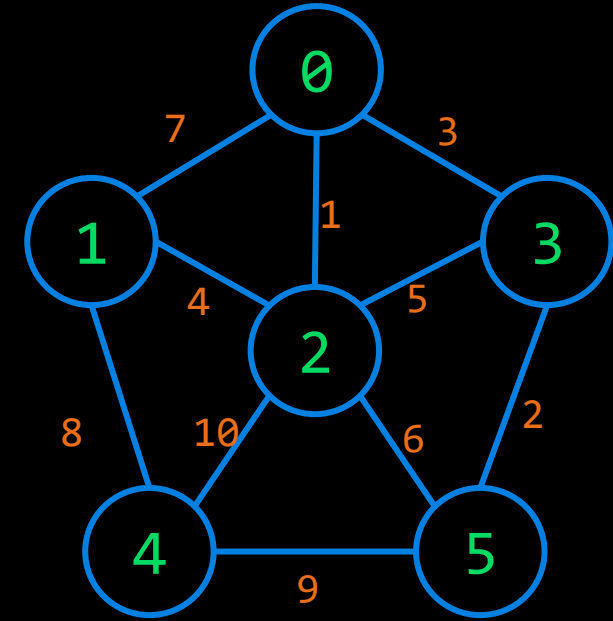
- Create an empty set, S .
- For each edge, E :
 - If either of the vertices connecting E is not a part of the set, add the vertices of E to S
 - If, both the vertices are part of the set S , ignore the edge as it forms a cycle.



Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
3-5	2
0-3	3
1-2	4
2-3	5
2-5	6
0-1	7
1-4	8
4-5	9
2-4	10

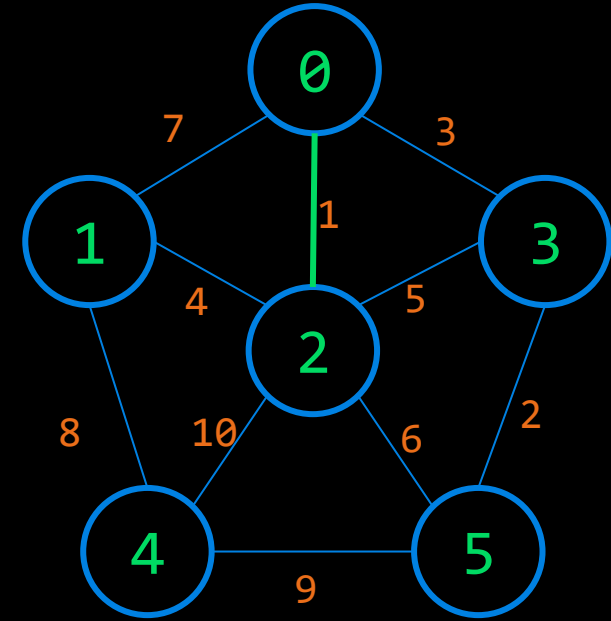


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
3-5	2
0-3	3
1-2	4
2-3	5
2-5	6
0-1	7
1-4	8
4-5	9
2-4	10

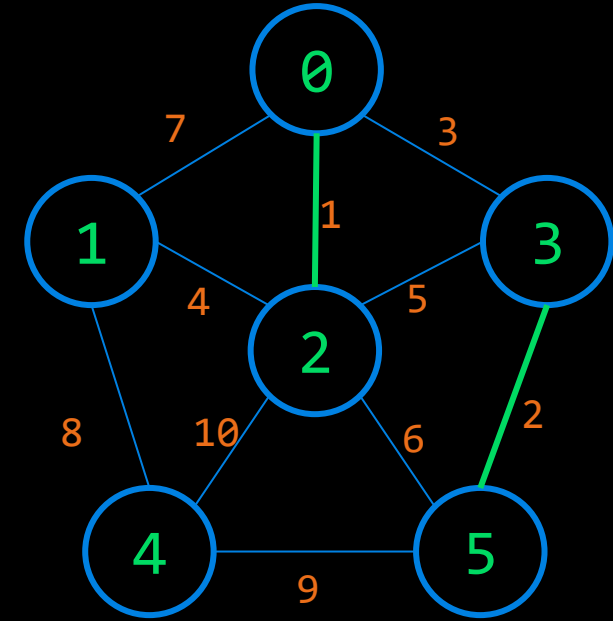
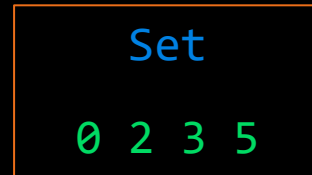


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
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1-4	8
4-5	9
2-4	10

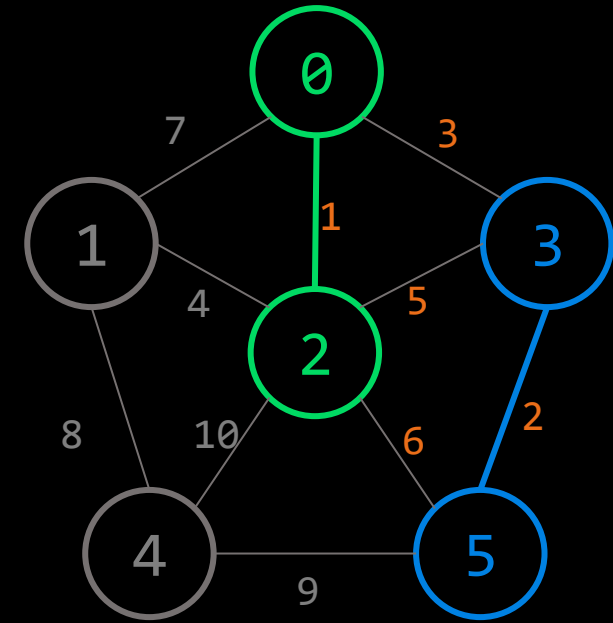
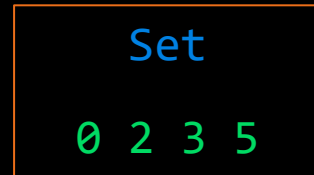


Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

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2-5	6
0-1	7
1-4	8
4-5	9
2-4	10



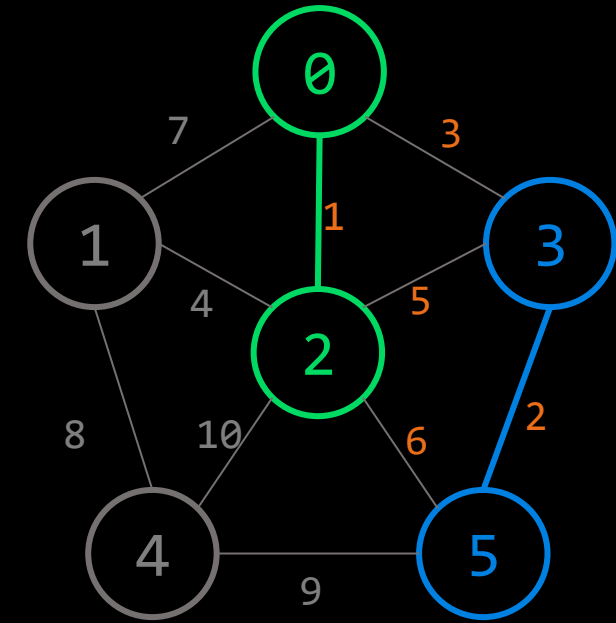
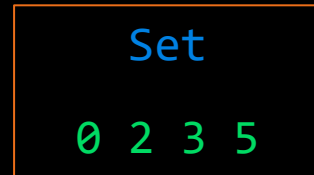
Disconnected Components

Add edges in order as long as they don't create a cycle

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
3-5	2
0-3	3
1-2	4
2-3	5
2-5	6
0-1	7
1-4	8
4-5	9
2-4	10



Disconnected Components
0-3, 2-3, 2-5 will never be
picked in MST

Add edges in order as long as they don't create a cycle

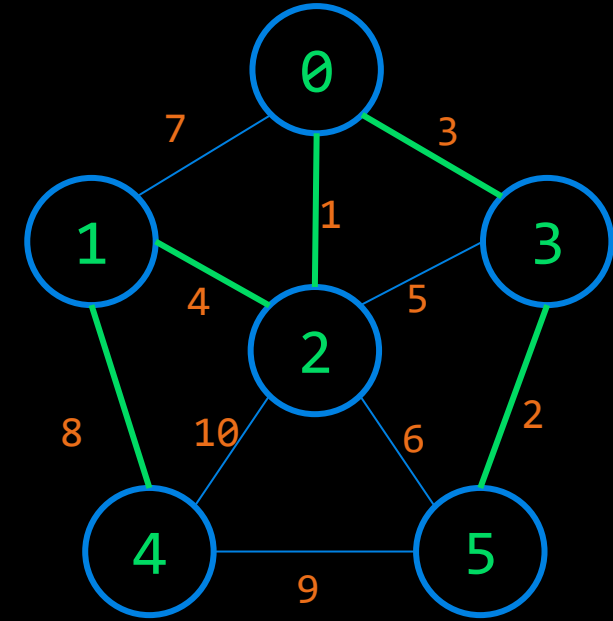
Kruskal's Algorithm

How can we detect a cycle when adding an edge?

Method 2a:

- Create an empty set, S .
- For each edge, E :
 - If either of the vertices connecting E is not a part of the set, add the vertices of E to S
 - If, both the vertices are part of the set S , ignore the edge as it forms a cycle.

This will not work whenever we pick edges in an order such that we have two disconnected components. Adding edges leads to connected components!



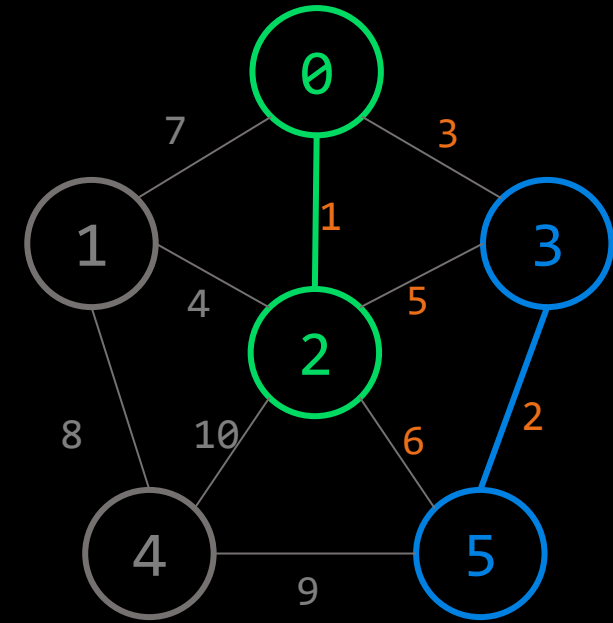
Kruskal's Algorithm

How can we detect a cycle when adding an edge?

Method 2b:

Disjoint Sets - Weighted Union

- A group of sets. There is no item in common in any of the sets.
- Operations:
 - $\text{find}(i)$ identify the set that contains i
 - $\text{union}(i, j)$ merge the set that contains i and the set that contains j
- Disjoint sets represent connected components.
- A cycle is created by adding an edge for which both vertices are in the same connected component.
- Complexity: $O(E \log V)$

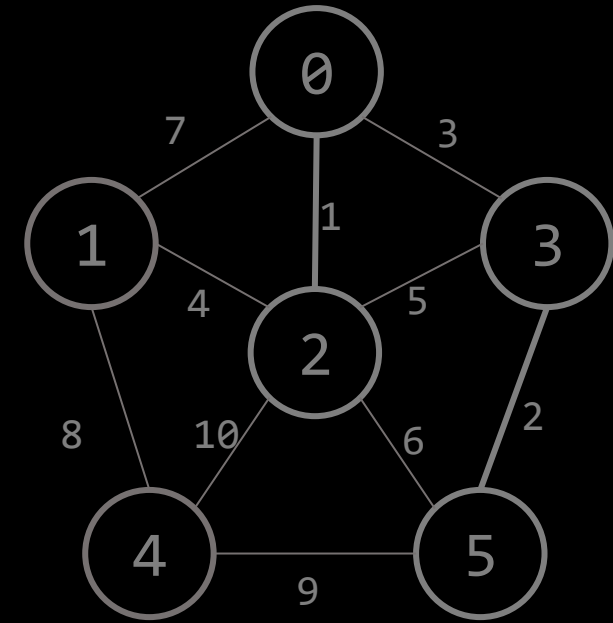


Disconnected Components
Two connected components:
 $\{0, 2\}$ and $\{3, 5\}$

Disjoint Sets

Disjoint Sets - Union/Find

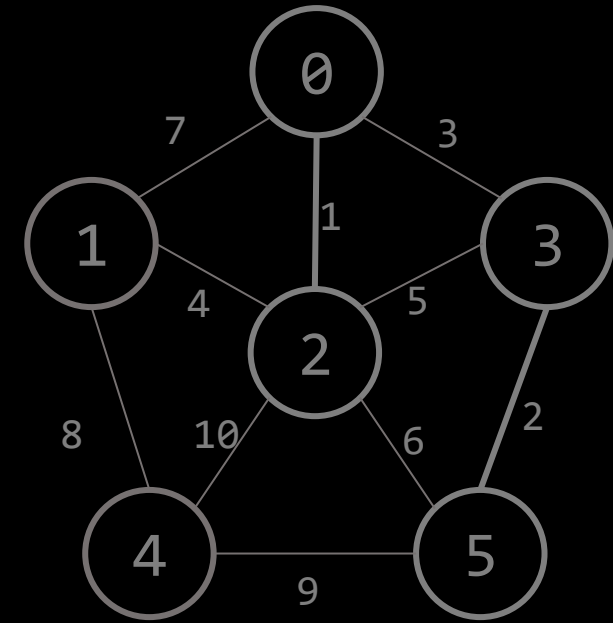
- Optimally represented as an array where each index stores the parent of the “index” vertex. An entire set is represented as a tree.
- Operations:
 - $\text{union}(i, j)$ merge the set that contains i and the set that contains j
 - $\text{find}(i)$ identify the set that contains i



Disjoint Sets

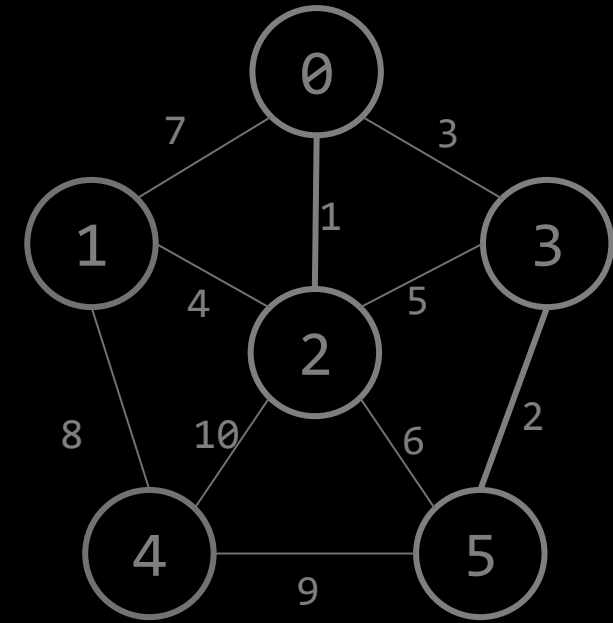
Disjoint Sets - Union/Find

- Optimally represented as an array where each index stores the parent of the “index” vertex. An entire set is represented as a tree.
- Operations:
 - `union(i, j)` merge the set that contains `i` and the set that contains `j`
`pi = find(i)`
`pj = find(j)`
`arr [pi] = pj`
 - `find(i)` identify the set that contains `i`
`if(arr[i]) == -1`
`return i`
`else`
`return find(arr[i])`



Disjoint Sets

- Operations:
 - `union(i, j)` merge the set that contains `i` and the set that contains `j`
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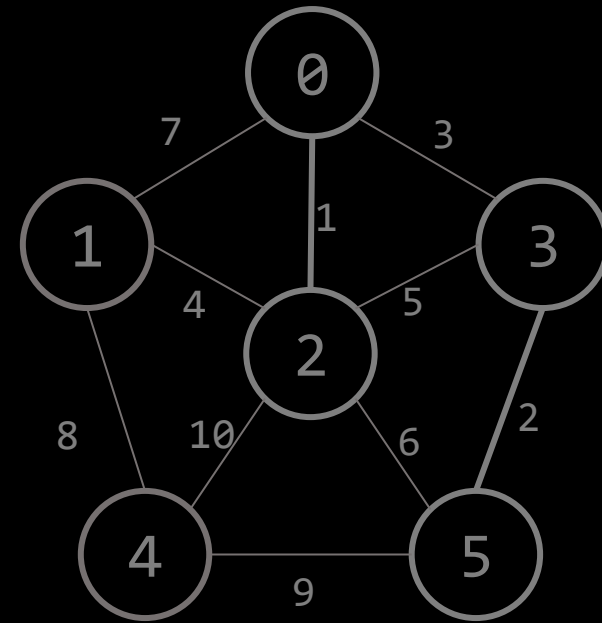


Disjoint Sets

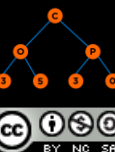
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`if(arr[i] == -1)`
 `return i`
`else`
 `return find(arr[i])`

Cycle detection:

```
choose an edge i-j
if find(i) != find(j) //not in same set
    union(i, j)      //merge sets
else
    "There is a cycle"
```



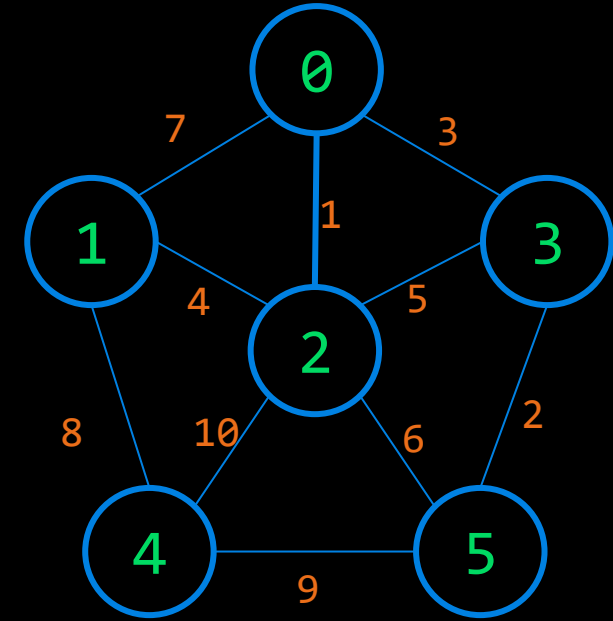
0-2	1
3-5	2
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Kruskal's Algorithm

Arrange edges in ascending order

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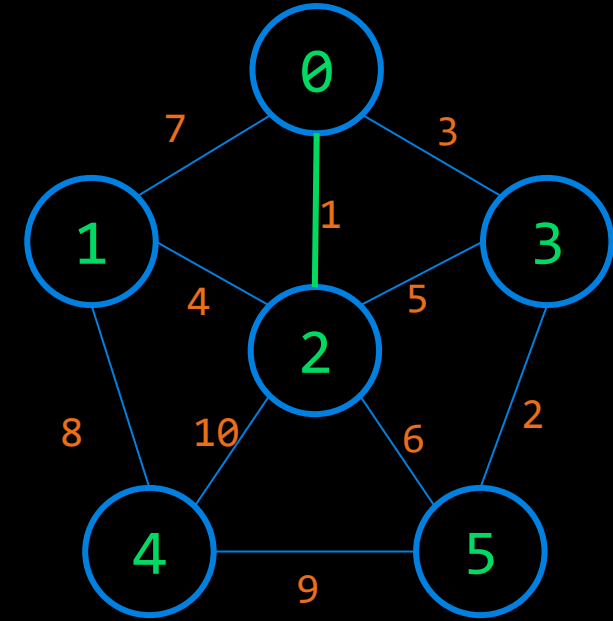
Kruskal's algorithm - choose an edge i - j
if $\text{find}(i) \neq \text{find}(j)$ //not in same disjoint set
union(i, j) //add the edge, which joins the components

Disjoint sets: $\{0\}$, $\{1\}$, $\{2\}$, $\{3\}$, $\{4\}$, $\{5\}$

Kruskal's Algorithm

Arrange edges in ascending order

0-2	1
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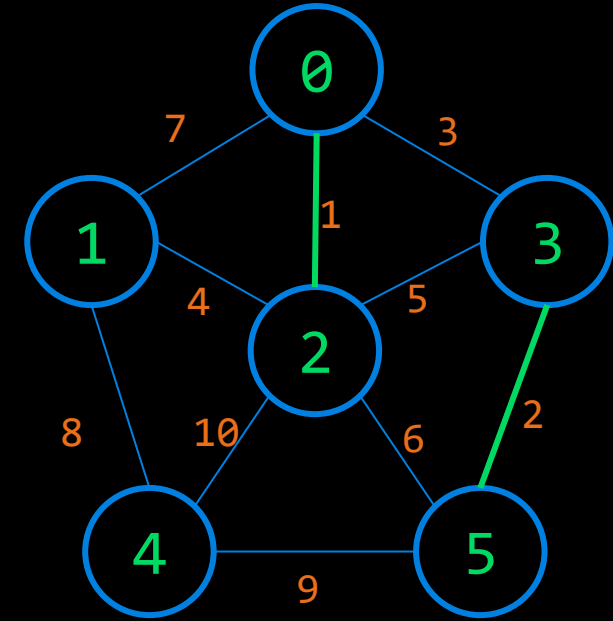
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Disjoint sets: $\{0, 2\}$, $\{1\}$, $\{3\}$, $\{4\}$, $\{5\}$

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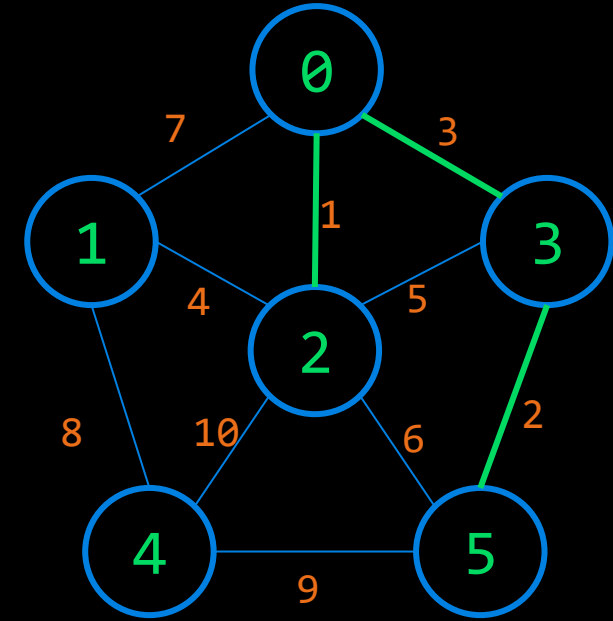
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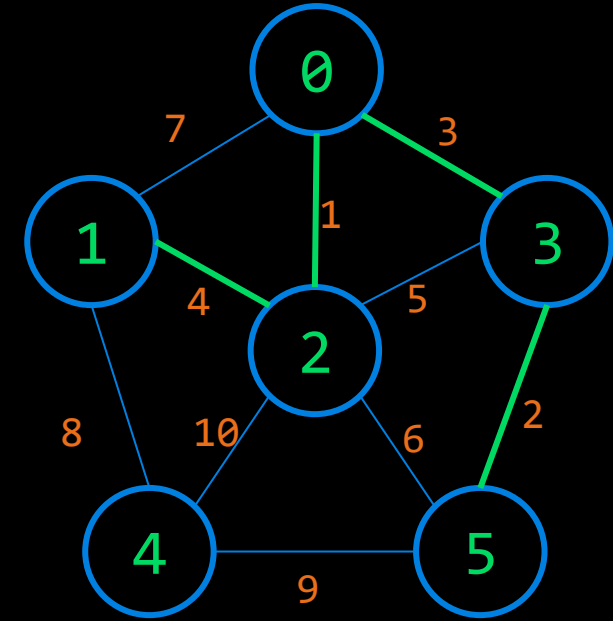
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Disjoint sets: $\{0, 2, 3, 5\}$, $\{1\}$, $\{4\}$

Kruskal's Algorithm

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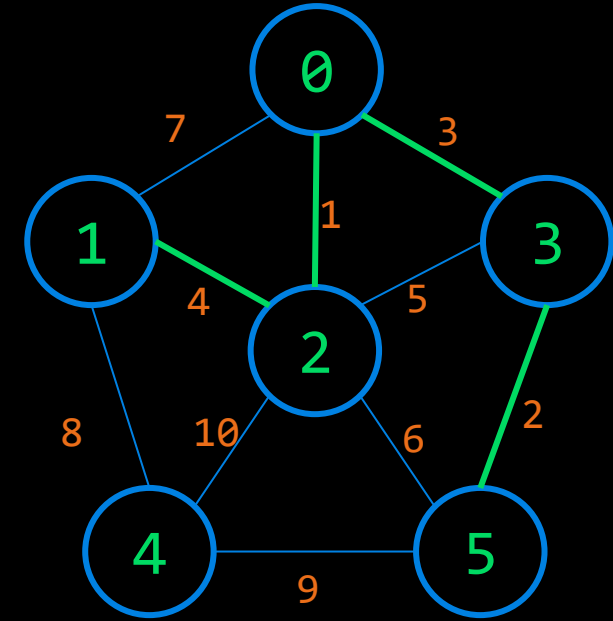
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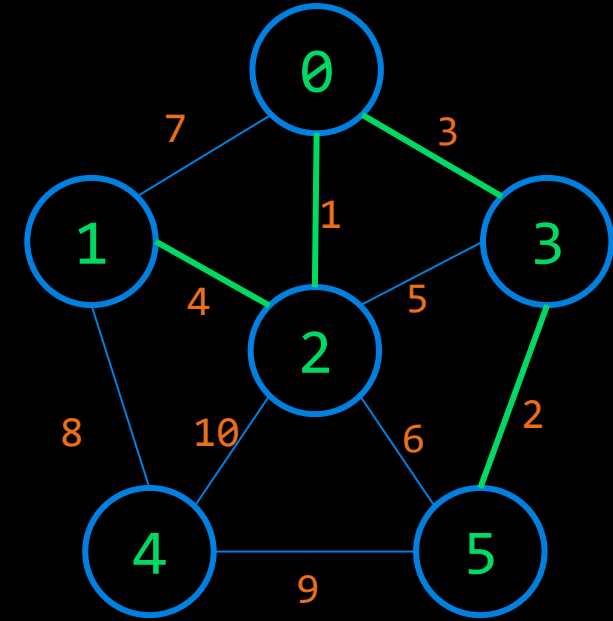
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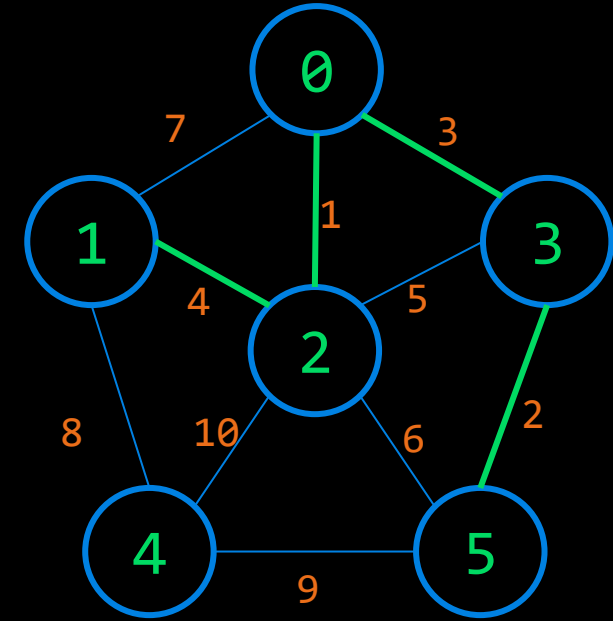
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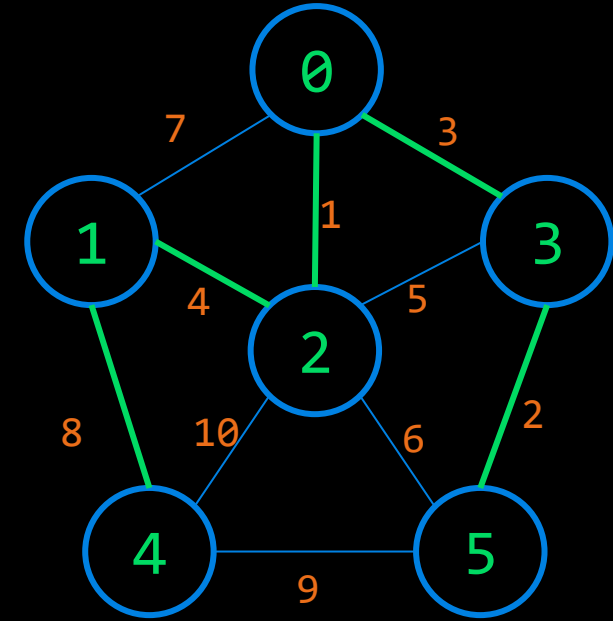
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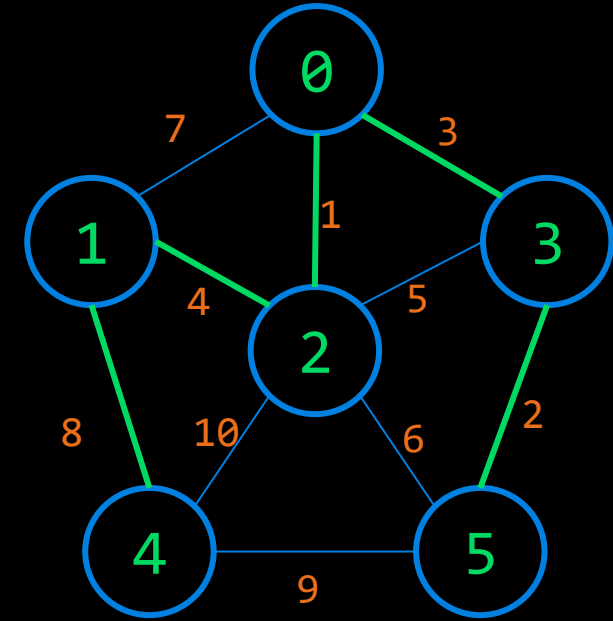
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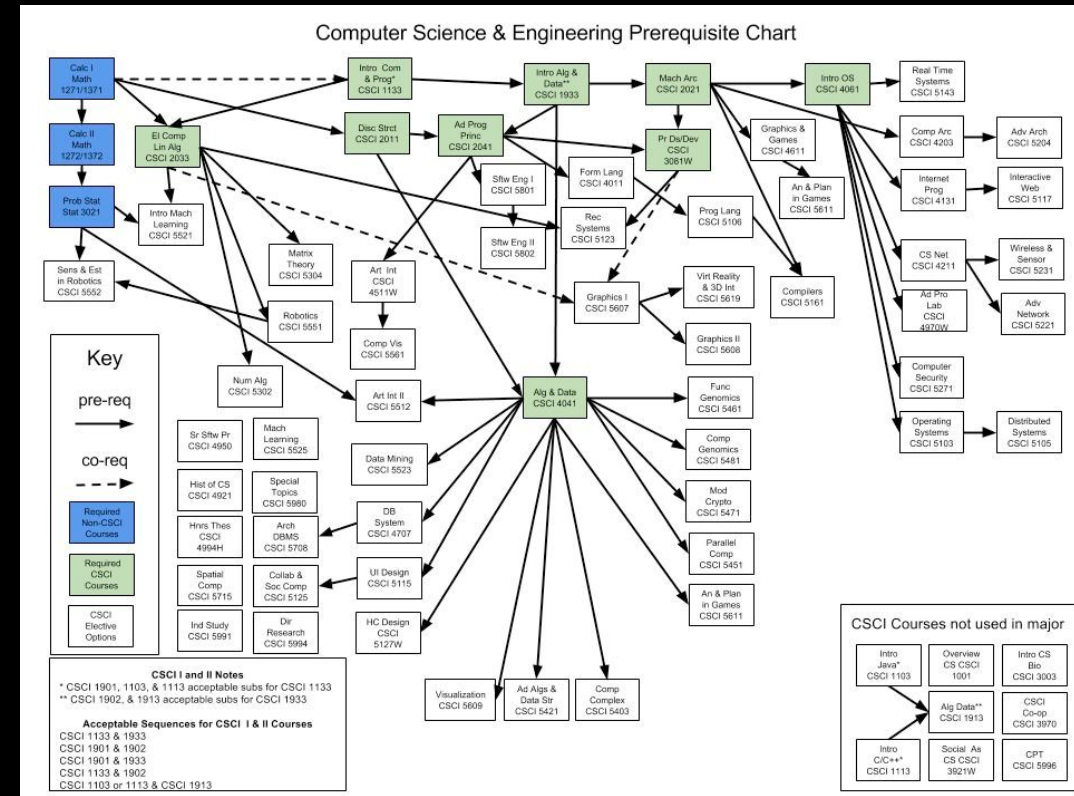
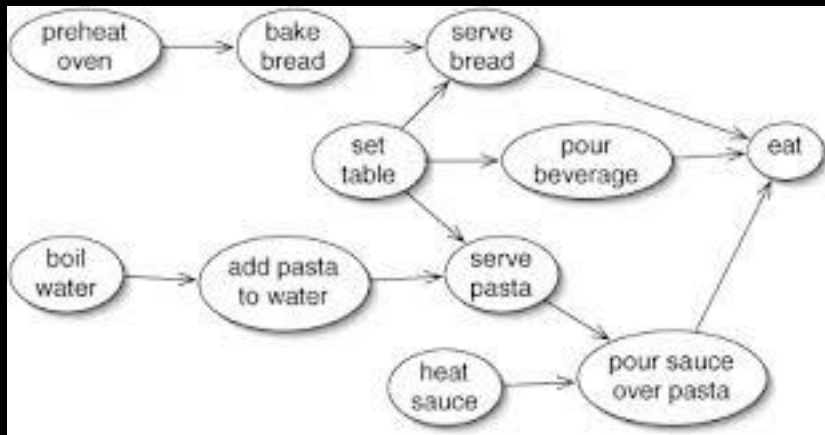
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union(i, j) //add the edge, which joins the components

Disjoint sets: $\{0, 1, 2, 3, 4, 5\}$

Topological Sort

Topological Sort

A topological sort is an ordering of vertices such that if there is an edge from v_i to v_j , then v_j comes after v_i

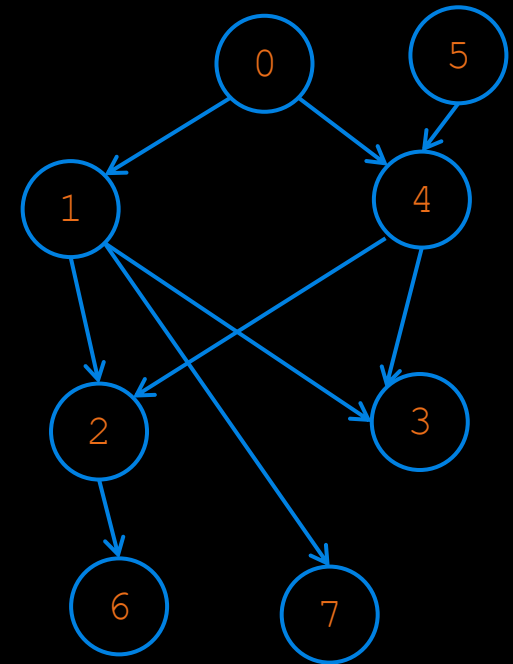


Topological Sort

A simple algorithm to find a topological ordering is first to find any vertex with no incoming edges.

We can then print this vertex, and remove it, along with its edges, from the graph.

Then we apply this same strategy to the rest of the graph.



Topological Sort

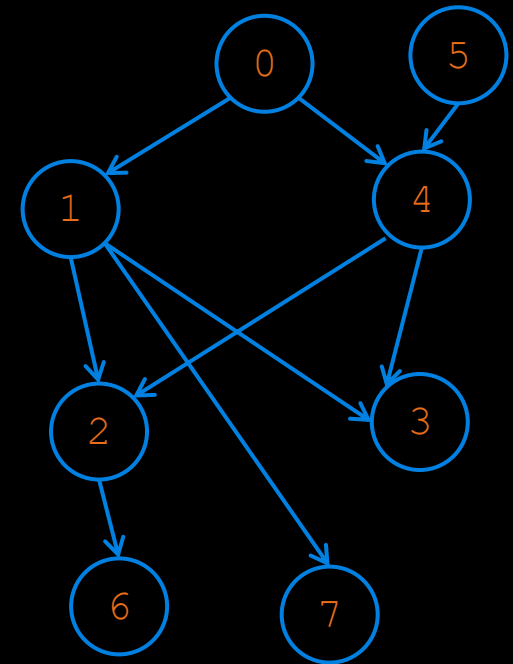
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$V_0 = \{0, 5\}$

Sort Order = $\{\}$



Topological Sort

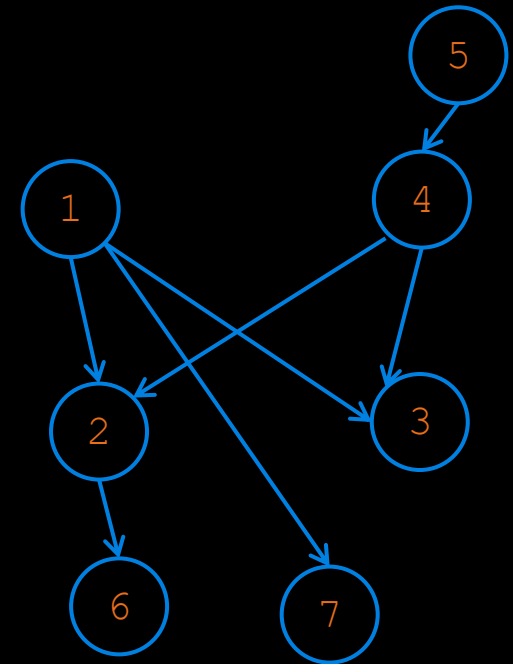
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Then we apply this same strategy to the rest of the graph.

$V_0 = \{5, 1\}$

Sort Order = $\{0\}$



Topological Sort

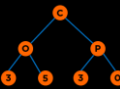
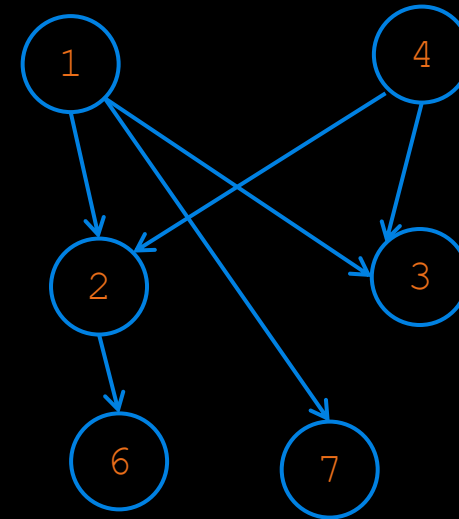
A simple algorithm to find a topological ordering is first to find any vertex with no incoming edges.

We can then print this vertex, and remove it, along with its edges, from the graph.

Then we apply this same strategy to the rest of the graph.

$V_0 = \{1, 4\}$

Sort Order = $\{0, 5\}$



Topological Sort

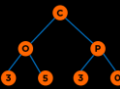
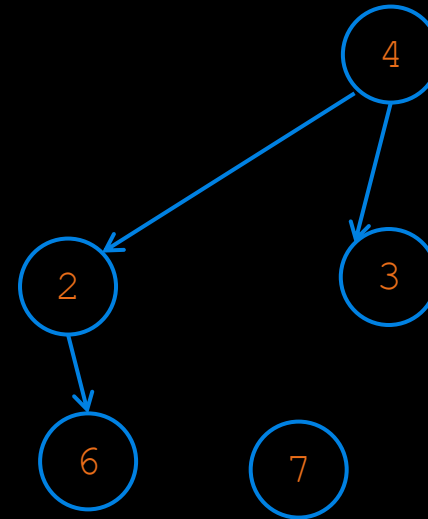
A simple algorithm to find a topological ordering is first to find any vertex with no incoming edges.

We can then print this vertex, and remove it, along with its edges, from the graph.

Then we apply this same strategy to the rest of the graph.

$V_0 = \{4, 7\}$

Sort Order = $\{0, 5, 1\}$



Topological Sort

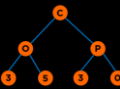
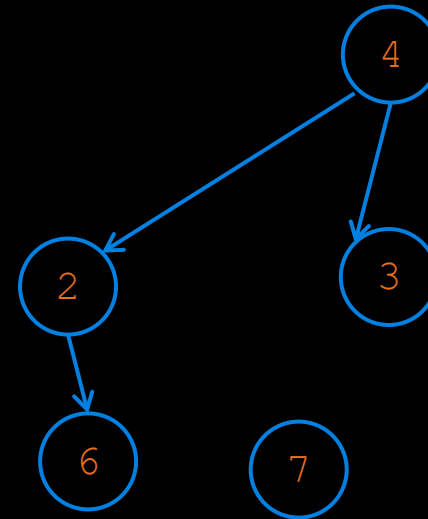
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Then we apply this same strategy to the rest of the graph.

$V_0 = \{4, 7\}$

Sort Order = $\{0, 5, 1\}$



Topological Sort

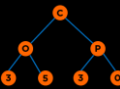
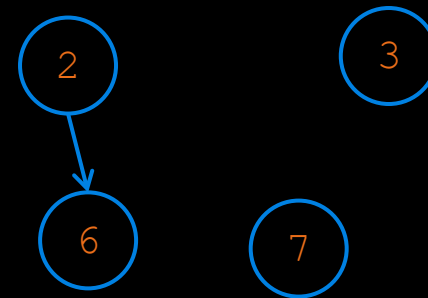
A simple algorithm to find a topological ordering is first to find any vertex with no incoming edges.

We can then print this vertex, and remove it, along with its edges, from the graph.

Then we apply this same strategy to the rest of the graph.

$V_0 = \{7, 2, 3\}$

Sort Order = $\{0, 5, 1, 4\}$



Topological Sort

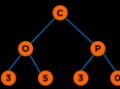
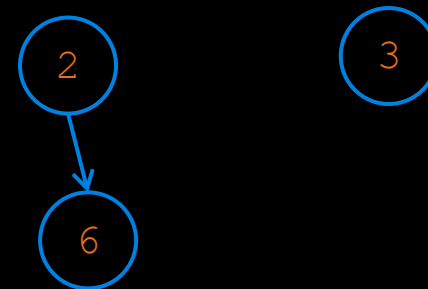
A simple algorithm to find a topological ordering is first to find any vertex with no incoming edges.

We can then print this vertex, and remove it, along with its edges, from the graph.

Then we apply this same strategy to the rest of the graph.

$V_0 = \{2, 3\}$

Sort Order = $\{0, 5, 1, 4, 7\}$



Topological Sort

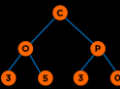
A simple algorithm to find a topological ordering is first to find any vertex with no incoming edges.

We can then print this vertex, and remove it, along with its edges, from the graph.

Then we apply this same strategy to the rest of the graph.

$V_0 = \{3, 6\}$

Sort Order = $\{0, 5, 1, 4, 7, 2\}$



Topological Sort

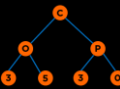
A simple algorithm to find a topological ordering is first to find any vertex with no incoming edges.

We can then print this vertex, and remove it, along with its edges, from the graph.

Then we apply this same strategy to the rest of the graph.

$V_0 = \{6\}$

Sort Order = $\{0, 5, 1, 4, 7, 2, 3\}$



Topological Sort

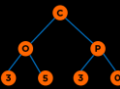
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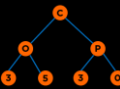
$V_0 = \{\}$

Sort Order = {0, 5, 1, 4, 7, 2, 3, 6}



Topological Sort Pseudocode

```
void Graph::topsort( )
{
    Queue<Vertex> q;
    int counter = 0;
    q.makeEmpty( );
    for each Vertex v
        if( v.indegree == 0 )
            q.enqueue( v );
    while( !q.isEmpty( ) )
    {
        Vertex v = q.dequeue( );
        v.topNum = ++counter; // Assign next number
        for each Vertex w adjacent to v
            if( --w.indegree == 0 )
                q.enqueue( w );
    }
    if( counter != NUM_VERTICES )
        throw CycleFoundException{ };
}
```



Questions

Questions

