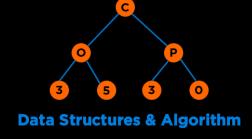
Graphs



Categories of Data Structures

Linear Ordered

Non-linear Ordered

Not Ordered

Lists

Trees

Sets

Stacks

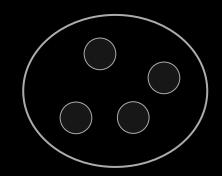
Graphs

Tables/Maps

Queues







Categories of Data Structures

Linear Ordered

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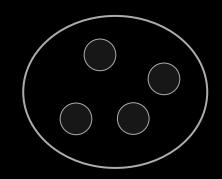
Graphs

Tables/Maps

Queues







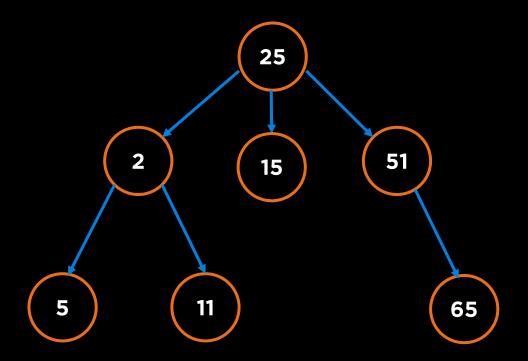
Agenda

- Graphs
 - Terminology
 - Types
 - Use cases
- Graph Implementations
 - Edge List
 - Adjacency Matrix
 - Adjacency List



Trees

Hierarchical, Acyclic, and Exactly one path between two nodes

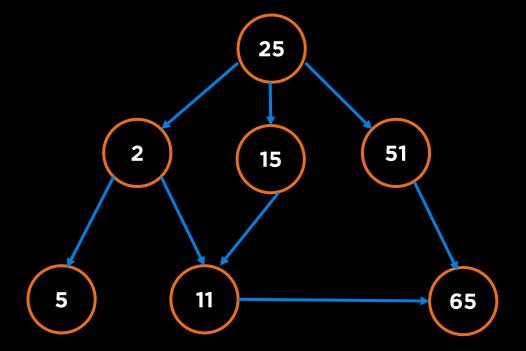




Graphs

An ordered pair of a set of nodes and a set of edges.

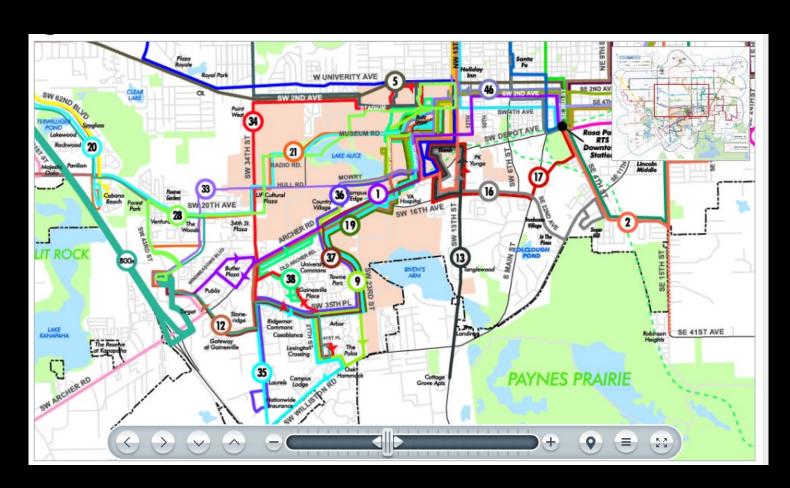
$$G = (V, E)$$





Graphs

Example



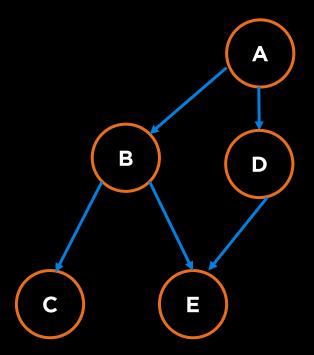


Vertex

Each node in a Graph is called a Vertex

```
V = {A, B, C, D, E}

|V| is the number of vertices in the graph
|V| = 5
```

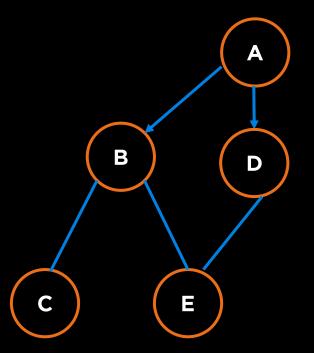




Edge

The connections between two nodes is called an edge.

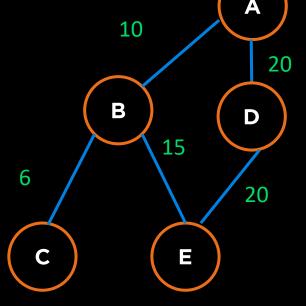
```
E = {(A,B), (A,D), {B,C}, {B,E}, {D,E}}
|E| is the number of edges in the graph
|E| = 5
```





Weight

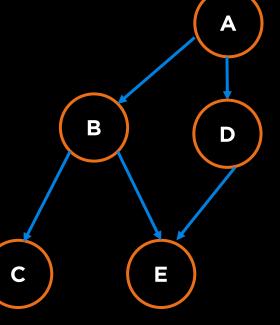
The edges in a graph may have associated values known as their weights. A weight is like a cost to travel from one vertex to the other over the edge.



Adjacent Vertices

A vertex is adjacent to another vertex if there is an edge to it from that other vertex.

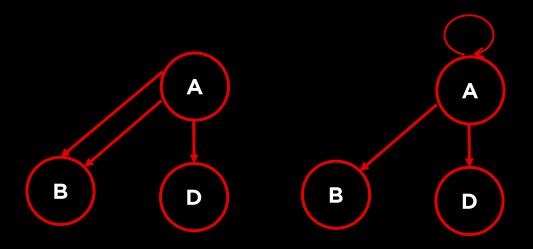
B is adjacent to A but A is not adjacent to B

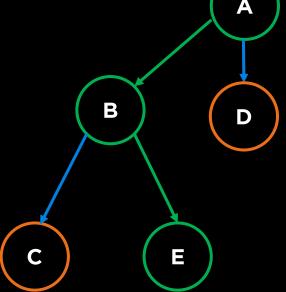




Simple Graph

A simple graph is a graph with no edges that connect a vertex to itself, i.e. no "loops" and no two edges that connect the same vertices, i.e. no "parallel edges".



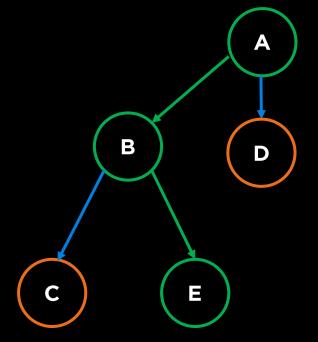




Path

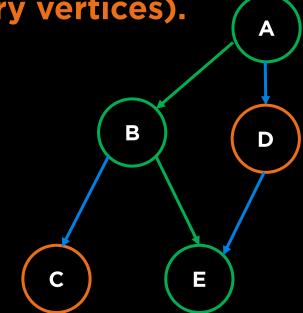
A path is a sequence of vertices in which each successive vertex is adjacent to its predecessor.

Path from A to E: A, B, E



Simple Path

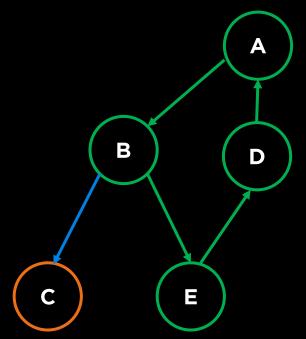
In a simple path, the vertices and edges are distinct except that the first and last vertex may be the same (no repeated intermediatory vertices).



Cycle

A cycle is a simple path in which only the first and final vertices are the same.

A - B - E - D - A is a cycle.

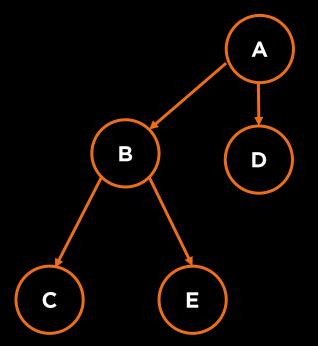




Connected Vertex

Two vertices are connected if there is a path between them.

A and C are connected
D and C are not connected

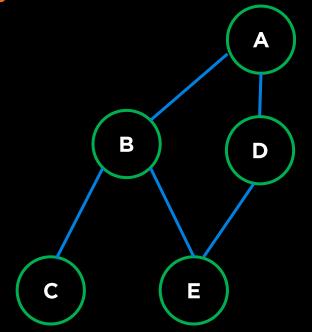




Connected Graph

An undirected graph is called a connected graph if there is a path from every vertex to every other vertex.

This is a connected graph



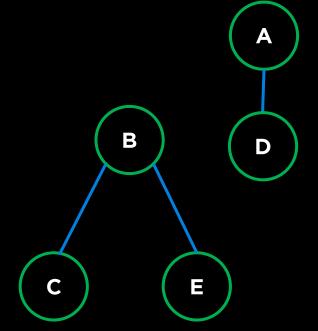


Connected Graph

An undirected graph is called a connected graph if there is a path from every vertex to every other vertex.

This is not a connected graph.

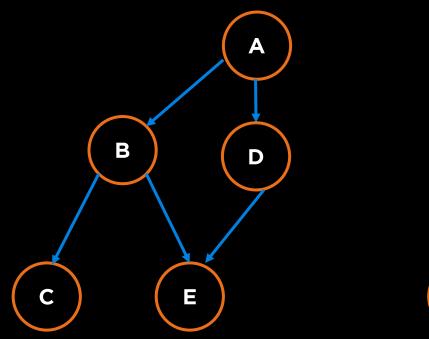
Connected components:
{A,D} and {B,C,E}

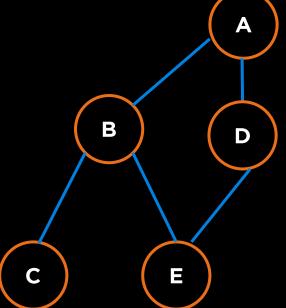




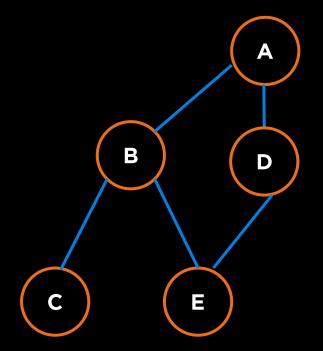


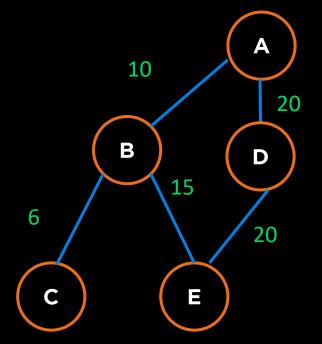
Directed (Digraph) vs Undirected





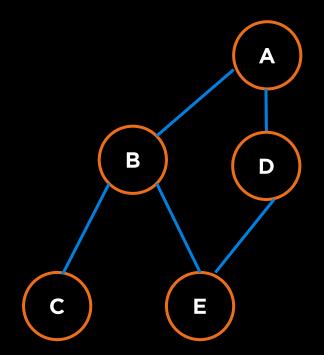
Weighted vs Unweighted

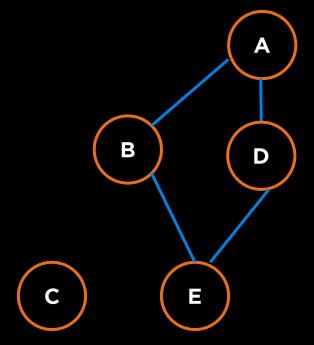




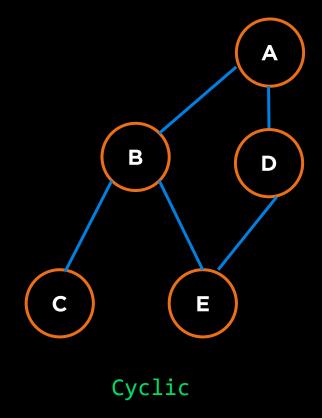


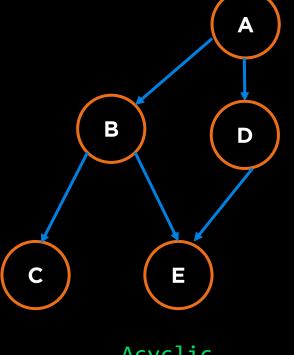
Connected vs Unconnected





Cyclic vs Acyclic

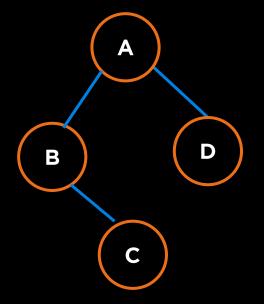






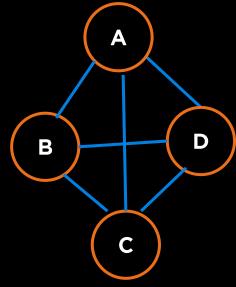
Dense vs Sparse

- The density of a graph is the ratio of |E| to |V|²
- We can assume that |E| is
 - ~ |V|² for a dense graph [Density ~ 1]
 - ~ |V| for a sparse graph [Density ~ 0]



```
Directed Graphs:
    0 <= |E| <= |V|(|V|-1)

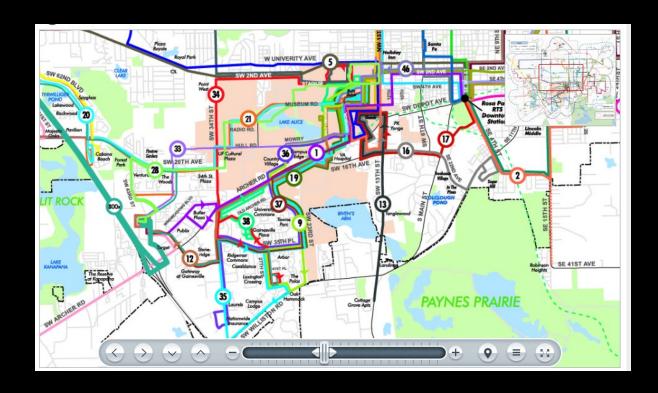
Undirected Graphs:
    0 <= |E| <= |V|(|V|-1)/2</pre>
```





Graphs

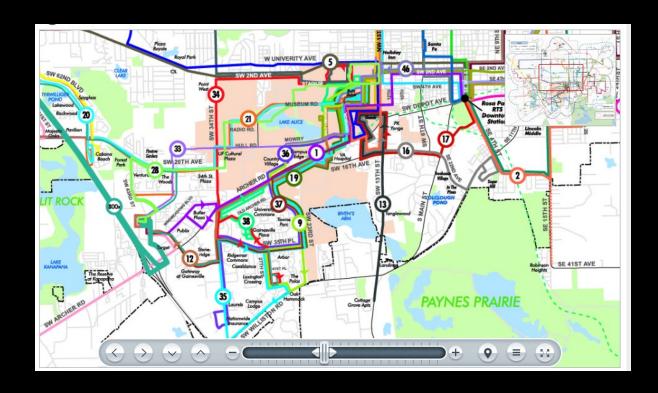
Example



Undirected
Directed
Cyclic
Connected

Graphs

Example



Undirected
Directed
Cyclic
Connected

Common Examples:

- Social Networks
- World Wide Web
- Maps

Weighted? Directed?

Common Examples:

- Social Networks (Unweighted, Undirected)
- World Wide Web (Unweighted, Directed)
- Maps (Weighted, Undirected)

There are lots of interesting questions we can ask about a graph:

- What is the shortest route from S to T? What is the longest without cycles?
- Are there cycles?
- Is there a tour you can take that only uses each node (station) exactly once?
- Is there a tour that uses each edge exactly once?



Some well-known graph problems and their common names:

- s-t Path. Is there a path between vertices s and t?
- Connectivity. Is the graph connected, i.e. is there a path between all vertices?
- Biconnectivity. Is there a vertex whose removal disconnects the graph?
- Shortest s-t Path. What is the shortest path between vertices s and t?
- Cycle Detection. Does the graph contain any cycles?
- Euler Tour. Is there a cycle that uses every edge exactly once?
- Hamilton Tour. Is there a cycle that uses every vertex exactly once?
- Planarity. Can you draw the graph on paper with no crossing edges?
- Isomorphism. Are two graphs isomorphic (the same graph in disguise)?

Often can't tell how difficult a graph problem is without very deep consideration.



Questions

Graph Implementations

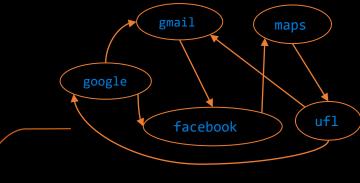


Graph API

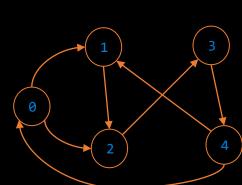
- No common ADT for Graphs
- Graphs were present before Object Oriented Programming
- API must include Graph methods, including their signatures and behaviors
- Defines how Graph client programmers must think.
- An underlying data structure to represent our graphs.
- Our choices can have profound implications on:
 - Runtime
 - Memory usage
 - Difficulty of implementing various graph algorithms

Common Convention

- Map labels to numbers, e.g. If node is called "google.com", assign it a number, say 0.
- Use a map data structure to achieve this: map<string, int>

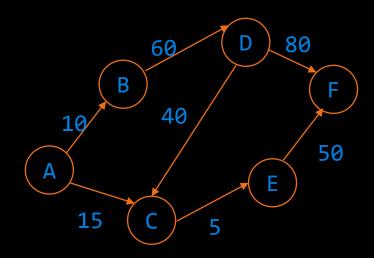


Label	Graph_Index
<pre>google.com</pre>	0
gmail.com	1
facebook.com	2
maps.com	3
ufl.edu	4





Common Operations

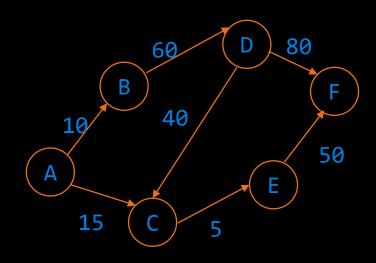


Connectedness

Neighborhood or Adjacency

G

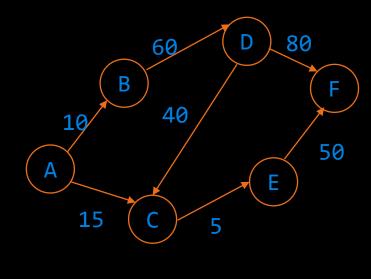
Common Representations



■ Edge List

Adjacency Matrix

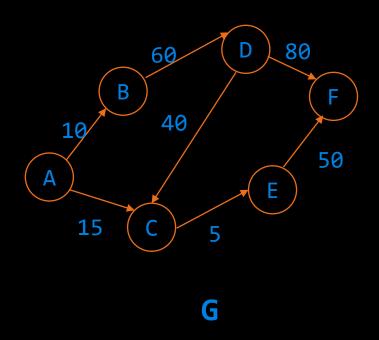
Adjacency List



G

$$G = \{(A,B), (A,C), (B,D), (D,C), (D,F), (E,F), (C,E)\}$$

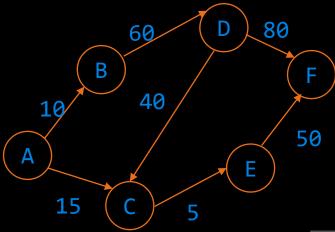




Α	В	10
Α	С	15
В	D	60
D	С	40
D	F	80
Е	F	50
С	Е	5

 $G = \{(A,B), (A,C), (B,D), (D,C), (D,F), (E,F), (C,E)\}$





G

Α	В	10
Α	С	15
В	D	60
D	С	40
D	F	80
Е	F	50
С	Е	5

Common Operations:

1. Connectedness

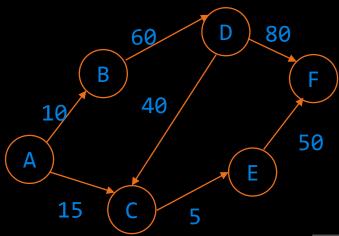
Is A connected to B?

2. Adjacency

What are A's adjacent nodes?

Space: ?





G

Α	В	10
Α	С	15
В	D	60
D	С	40
D	F	80
Е	F	50
С	Е	5

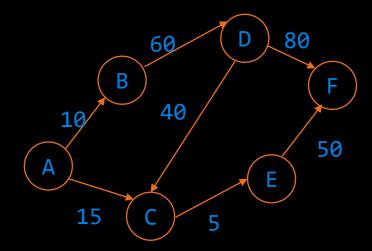
Common Operations:

1. Connectedness

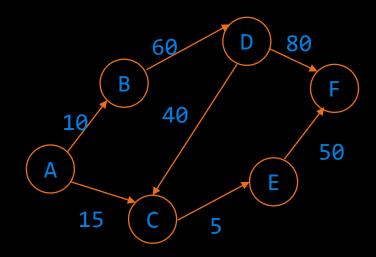
2. Adjacency

Space: O(E)



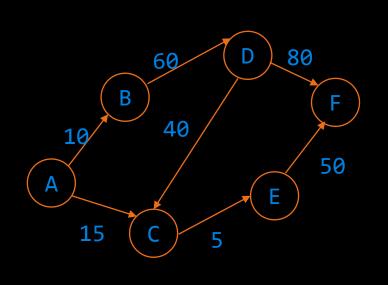


G



G



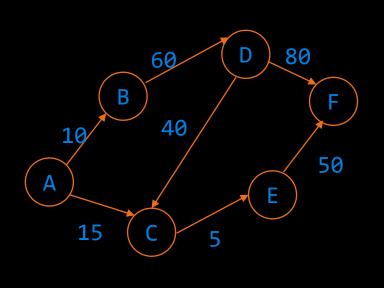


А	В	C	D	Е	F	
0	10	15	0	0	0	
0	0	0	60	0	0	
0	0	0	0	5	0	
0	0	40	0	0	80	
0	0	0	0	0	50	
0	0	0	0	0	0	

G

```
G[from][to] = weight; (if there is an edge, "from" -> "to")
G[from][to] = 0; (otherwise)
```



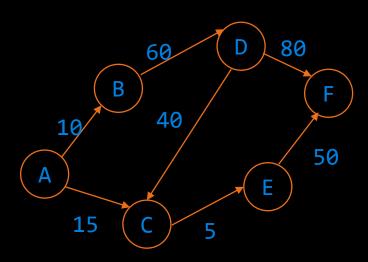


Α	В	C	D	Е	F
0	10	15	0	0	0
0	0	0	60	0	0
0	0	0	0	5	0
0	0	40	0	0	80
0	0	0	0	0	50
0	0	0	0	0	0

```
G[from][to] = weight; (if there is an edge, "from" -> "to")
G[from][to] = 0; (otherwise)
```



Adjacency Matrix Implementation



Map

Insertion:

```
G[from][to] = weight; (if there is an edge, "from" -> "to")
G[from][to] = 0; (otherwise)
```

Т	n	n	П	+
_	ш	Ρ	u	J
		•		

A B 10 A C 15

B D 60

D C 40

C E 5

D F 80

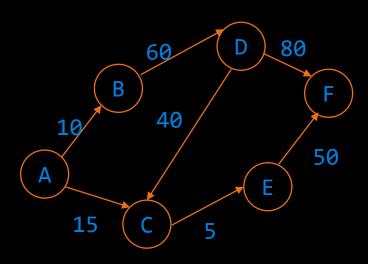
E F 50

0	1	2	3	4	

0	0	10	15	0	0	0
1	0	0	0	60	0	0
2	0	0	0	0	5	0
3	0	0	40	0	0	80
4	0	0	0	0	0	50
5	0	0	0	0	0	0



Adjacency Matrix Implementation



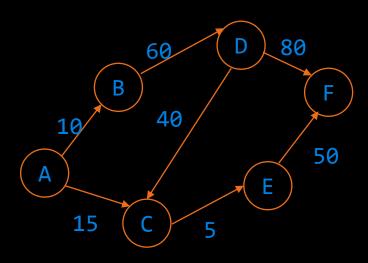
Ιr	าрเ	ıt			
7					G
Д	В	10			
Д	C	15			
В	D	60			
D	C	40			
C	Е	5			
D	F	80			
		ГΩ			

		0	1	2	3	4	5
Мар	0	0	10	15	0	0	0
A 0	1	0	0	0	60	0	0
B 1 C 2	2	0	0	0	0	5	0
D 3	3	0	0	40	0	0	80
E 4	4	0	0	0	0	0	50
F 5	5	0	0	0	0	0	0

```
G[from][to] = weight; (if there is an edge, "from" -> "to")
G[from][to] = 0; (otherwise)
```

```
#include <iostream>
    #include<map>
    #define VERTICES 6
    using namespace std;
    int main()
06
           int no lines, wt, j=0;
           string from, to;
           int graph [VERTICES][VERTICES] = {0};
          map<string, int> mapper;
10
           cin >> no lines;
11
12
13
14
15
16
17
19
20
21
           return 0;
```

Adjacency Matrix Implementation

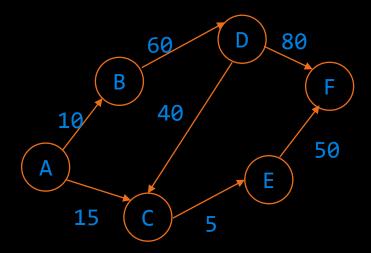


Ιr	าрเ	ıt	
7			
Д	В	10	
Д	C	1 5	
В	D	60	
D	C	40	
C	Е	5	
D	F	80	
	Е	EQ	

		0	1	2	3	4	5
Мар	0	0	10	15	0	0	0
A 0	1	0	0	0	60	0	0
B 1 C 2	2	0	0	0	0	5	0
D 3	3	0	0	40	0	0	80
E 4	4	0	0	0	0	0	50
F 5	5	0	0	0	0	0	0

```
G[from][to] = weight; (if there is an edge, "from" -> "to")
G[from][to] = 0; (otherwise)
```

```
#include <iostream>
    #include<map>
    #define VERTICES 6
    using namespace std;
    int main()
06
           int no lines, wt, j=0;
           string from, to;
           int graph [VERTICES][VERTICES] = {0};
10
           map<string, int> mapper;
           cin >> no lines;
11
12
           for(int i = 0; i < no lines; i++)</pre>
13
                 cin >> from >> to >> wt;
14
                 if (mapper.find(from) == mapper.end())
15
                        mapper[from] = j++;
                 if (mapper.find(to) == mapper.end())
                        mapper[to] = j++;
                 graph[mapper[from]][mapper[to]] = wt;
19
20
21
           return 0;
```



G

Map A 0 B 1 C 2 D 3 E 4 F 5
 0
 1
 2
 3
 4
 5

 0
 0
 10
 15
 0
 0
 0

 1
 0
 0
 0
 60
 0
 0

 2
 0
 0
 0
 0
 5
 0

 3
 0
 0
 40
 0
 0
 80

 4
 0
 0
 0
 0
 0
 0

 5
 0
 0
 0
 0
 0
 0

Common Operations:

Connectedness

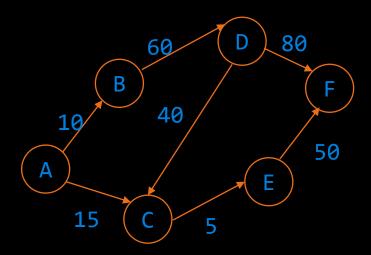
Is A connected to B?

2. Adjacency

What are A's adjacent nodes?

Space: ?





G

Ma	ар
Α	0
В	1
C	2
D	3
Е	4
F	5

 0
 1
 2
 3
 4
 5

 0
 0
 10
 15
 0
 0
 0

 1
 0
 0
 0
 60
 0
 0

 2
 0
 0
 0
 0
 5
 0

 3
 0
 0
 40
 0
 0
 80

 4
 0
 0
 0
 0
 0
 50

 5
 0
 0
 0
 0
 0
 0

Common Operations:

Connectedness

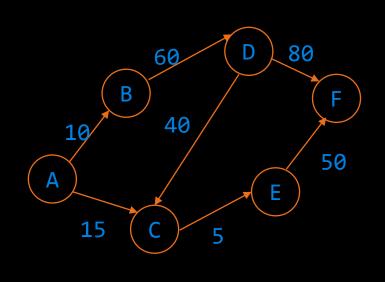
2. Adjacency

What are A's adjacent nodes?

Space: **0(|V| * |V|)**



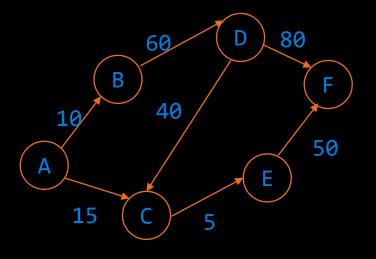
Adjacency Matrix Problem



Α	В	C	D	Е	F
0	10	15	0	0	0
0	0	0	60	0	0
0	0	0	0	5	0
0	0	40	0	0	80
0	0	0	0	0	50
0	0	0	0	0	0

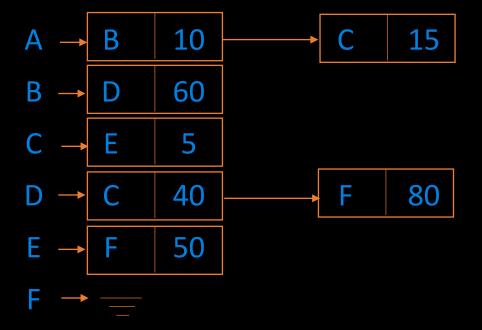
G

Adjacency List



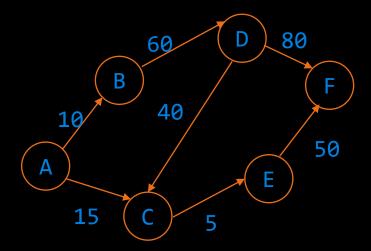
G

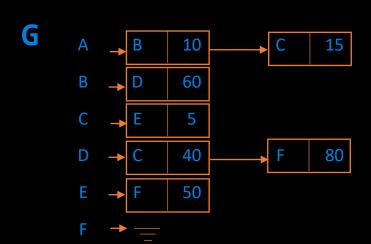
А	В	C	D	E	F	
0	10	15	0	0	0	
0	0	0	60	0	0	
0	0	0	0	5	0	
0	0	40	0	0	80	
0	0	0	0	0	50	
0	0	0	0	0	0	





Adjacency List





Common Operations:

Connectedness

Is A connected to B?

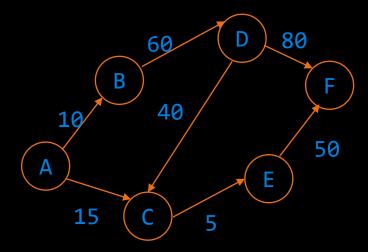
2. Adjacency

What are A's adjacent nodes?

Space: ?

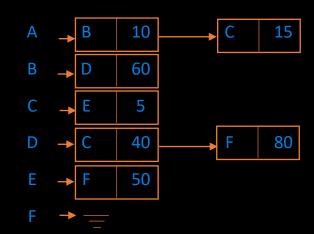


Adjacency List



G

Sparse Graph:
Edges ~ Vertices



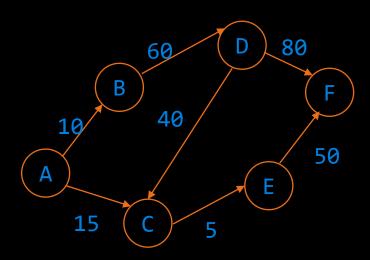
Common Operations:

Connectedness

2. Adjacency

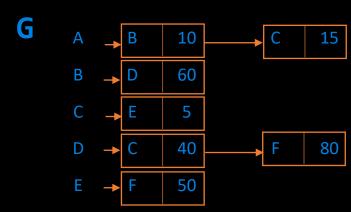


Adjacency List Implementation



Input

7
A B 10
A C 15
B D 60
D C 40
C E 5
D F 80
E F 50



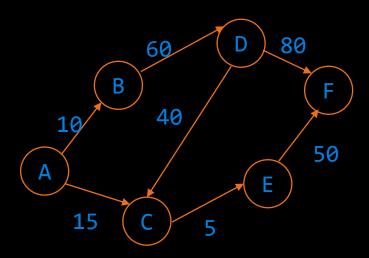
Insertion:

If to or from vertex not present add vertex Otherwise add edge at the end of the list

```
#include <iostream>
    #include<map>
    #include<vector>
    #include<iterator>
    using namespace std;
07
    int main()
           int no lines;
           string from, to, wt;
11
           map<string, vector<pair<string,int>>> graph;
           cin >> no_lines;
12
           for(int i = 0; i < no lines; i++)</pre>
13
14
15
16
17
18
19
20
```



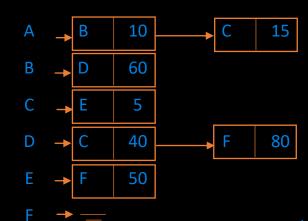
Adjacency List Implementation



Input

7
A B 10
A C 15
B D 60
D C 40
C E 5
D F 80
E F 50

G



Insertion:

If to or from vertex not present add vertex Otherwise add edge at the end of the list

```
#include <iostream>
    #include<map>
    #include<vector>
    #include<iterator>
    using namespace std;
07
    int main()
           int no lines;
           string from, to, wt;
11
           map<string, vector<pair<string,int>>> graph;
           cin >> no_lines;
12
           for(int i = 0; i < no lines; i++)</pre>
13
14
15
                  cin >> from >> to >> wt;
                  graph[from].push back(make pair(to, stoi(wt)));
16
17
                  if (graph.find(to)==graph.end())
18
                         graph[to] = {};
19
20
```



Graph Implementation

	Edge List	Adjacency Matrix	Adjacency List
Time Complexity: Connectedness	O(E)	0(1)	O(outdegree(V))
Time Complexity: Adjacency	O(E)	0(V)	O(outdegree(V))
Space Complexity	O(E)	O(V*V)	O(V+E)

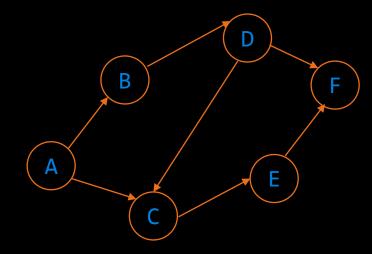


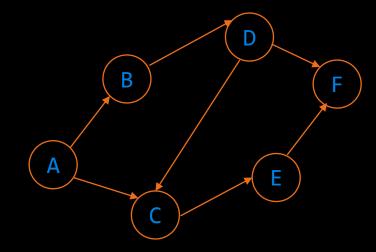
One Graph API

```
class Graph
   private:
     //Graph Data Structure
  public:
     Graph();
     Graph(int V); //Creates graph with v vertices
     int V(); //Returns number of vertices
     int E(); //Returns number of edges
     void insertEdge(int from, int to, int weight);
     bool isEdge(int from, int to);
     int getWeight(int from, int to);
     vector<int> getAdjacent(int vertex);
     void printGraph();
```

Graph Traversal



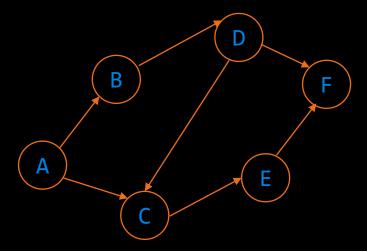




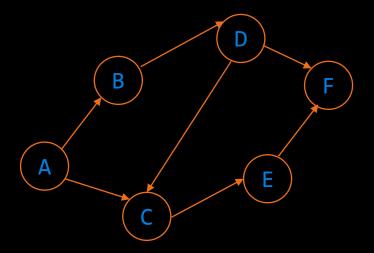
Valid BFS:



```
    Take an arbitrary start vertex, mark it identified, and place it in a queue.
    while the queue is not empty
    Take a vertex, u, out of the queue and visit u.
    for all vertices, v, adjacent to this vertex, u
    if v has not been identified or visited
    Mark it identified
    Insert vertex v into the queue.
    We are now finished visiting u.
```



```
Algorithm for Breadth-First Search
      Take an arbitrary start vertex, mark it identified,
1.
     and place it in a queue.
      while the queue is not empty
3.
           Take a vertex, u, out of the queue and visit u.
           for all vertices, v, adjacent to this vertex, u
4.
                 if \nu has not been identified or visited
6.
                       Mark it identified
7.
                       Insert vertex \nu into the queue.
8.
          We are now finished visiting u.
```



```
string source = "A";
    std::set<string> visited;
    std::queue<string> q;
04
    visited.insert(source);
    q.push(source);
07
    cout<<"BFS: ";</pre>
08
    while(!q.empty())
10
11
          string u = q.front();
12
          cout << u;
13
          q.pop();
14
          vector<string> neighbors = graph[u];
          std::sort(neighbors.begin(), neighbors.begin() + neighbors.size());
          for(string v: neighbors)
16
17
                if(visited.count(v) == 0)
18
19
20
                       visited.insert(v);
21
                      q.push(v);
22
23
24
```

Breadth First Search: Alternate way (7.2.2)

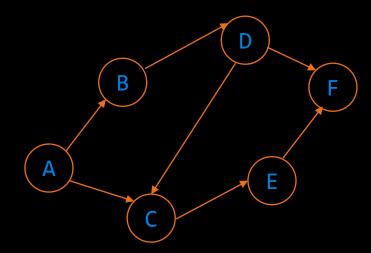
```
Algorithm for Breadth-First Search
      Take an arbitrary start vertex, mark it identified,
     and place it in a queue.
      while the queue is not empty
2.
           Take a vertex, u, out of the queue and visit u.
3.
           for all vertices, v, adjacent to this vertex, u
5.
                 if v has not been identified or visited
6.
                       Mark it identified
7.
                       Insert vertex \nu into the queue.
          We are now finished visiting u.
8.
```

```
// Visited Vertices Alternate
set<string> visited;
visited.insert(source);
if(visited.count(v)==0)
    visited.insert(v);
```

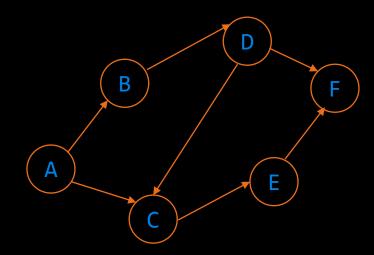
```
void bfs(const Graph& graph, int src)
02
        vector<bool> visited(graph.numVertices);
03
        queue<int> q;
        visited[src] = true;
        q.push(src);
        while (!q.empty())
10
            int u = q.front();
11
            cout << u << " ";
12
            q.pop();
14
            for (int v : graph.adjList[u])
15
16
17
                 if (!visited[v])
18
19
                     visited[v] = true;
20
                     q.push(v);
21
22
23
24
```



Depth First Search



Depth First Search

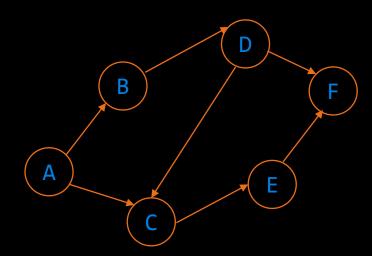


Valid DFS: A, B, D, C, E, F

Depth First Search

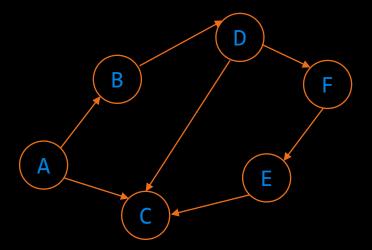
```
Algorithm for Depth-First Search

1. Take an arbitrary start vertex, mark it visited, and place it in a stack.
2. while the stack is not empty
3. the item on top of the stack is u
4. if there is a vertex, v, adjacent to this vertex, u, that has not been visited
5. Mark v visited
6. Push vertex v onto the top of the stack
7. else
8. pop stack
```



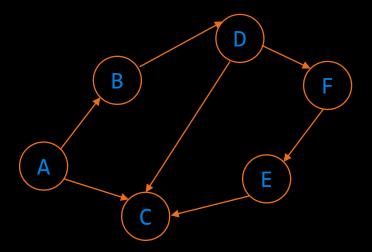
Depth First Search - Modified BFS

```
    Take an arbitrary start vertex, mark it identified, and place it in a stack.
    while the stack is not empty
    Take a vertex, u, out of the stack and visit u.
    for all vertices, v, adjacent to this vertex, u
    if v has not been identified or visited
    Mark it identified
    Insert vertex v into the stack.
    We are now finished visiting u.
```



Depth First Search - Modified BFS

```
Algorithm for Depth-First Search
      Take an arbitrary start vertex, mark it identified,
1.
     and place it in a stack.
      while the stack is not empty
3.
           Take a vertex, u, out of the stack and visit u.
4.
           for all vertices, v, adjacent to this vertex, u
                 if \nu has not been identified or visited
6.
                       Mark it identified
7.
                       Insert vertex \nu into the stack.
          We are now finished visiting u.
8.
```



```
string source = "A";
    std::set<string> visited;
    std::stack<string> s;
04
05
    visited.insert(source);
    s.push(source);
    cout<<"DFS: ";</pre>
08
    while(!s.empty())
10
          string u = s.top();
11
12
          cout<<u;
13
          s.pop();
14
          vector<string> neighbors = graph[u];
15
          for(string v: neighbors)
16
                 if(visited.count(v)==0)
17
18
                       visited.insert(v);
19
20
                       s.push(v);
21
22
23
```



BFS vs DFS

```
string source = "A";
    std::set<string> visited;
    std::queue<string> q;
03
04
05
    visited.insert(source);
    q.push(source);
06
07
    cout<<"BFS: ";</pre>
08
09
    while(!q.empty())
10
          string u = q.front();
11
12
          cout<<u;
13
          q.pop();
          vector<string> neighbors = graph[u];
14
15
          for(string v: neighbors)
16
17
18
                       visited.insert(v);
19
20
                      q.push(v);
21
22
23
```

```
string source = "A";
    std::set<string> visited;
    std::stack<string> s;
04
    visited.insert(source);
    s.push(source);
07
    cout<<"DFS: ";</pre>
08
    while(!s.empty())
10
11
          string u = s.top();
12
          cout<<u;
13
          s.pop();
          vector<string> neighbors = graph[u];
14
15
          for(string v: neighbors)
16
17
                if(visited.count(v)==0)
18
                       visited.insert(v);
19
20
                       s.push(v);
21
22
23
```



Mentimeter

3176 5158

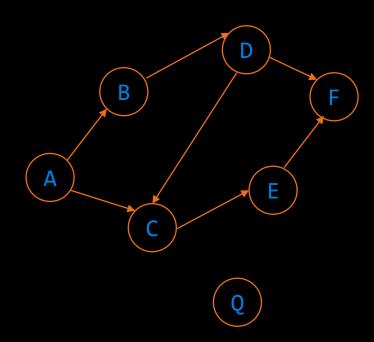




Questions

s-t Path

Is there a path between vertices s and t?

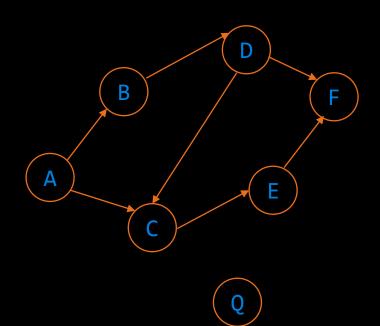


Is there a path between vertices A and C? - Yes
Is there a path between vertices A and Q? - No



s-t Path

Is there a path between vertices s and t?

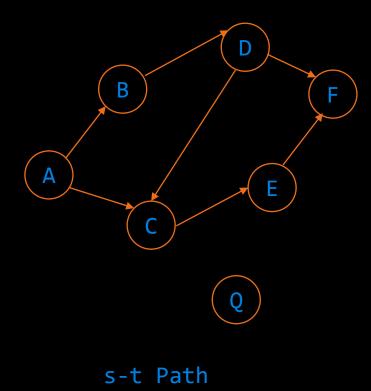


Is there a path between vertices A and C? - Yes
Is there a path between vertices A and Q? - No

Solution

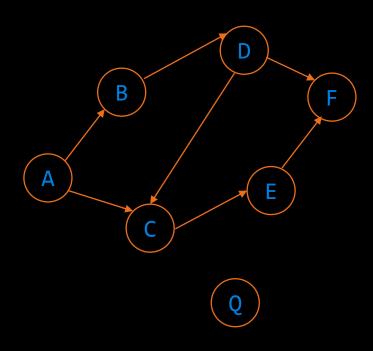
Perform DFS or BFS with source "s" and if we encounter "t" in the path/traversal, then return True otherwise False

7.2.1 DFS to Find Whether a Given Vertex is Reachable (Iterative)



```
bool dfs(const Graph& graph, int src, int dest)
        set<int> visited;
        stack<int> s;
        visited.insert(src);
        s.push(src);
        while(!s.empty())
            int u = s.top();
            s.pop();
            for(auto v: graph.adjList[u])
11.
12.
                if(v == dest)
                     return true;
                if ((visited.find(v) == visited.end()))
                     visited.insert(v);
                     s.push(v);
        return false;
23. }
```

7.2.1 DFS to Find Whether a Given Vertex is Reachable (Recursive)



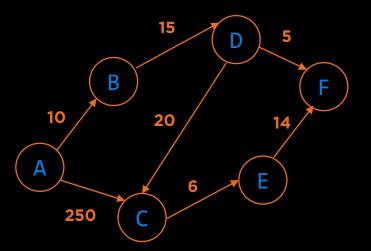
s-t Path: Recursive

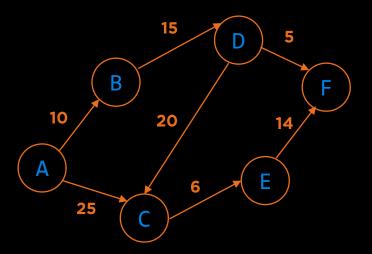
```
bool dfs helper(const Graph& graph, int src, int dest, vector<bool>& visited)
        visited[src] = true;
        if (src == dest)
            return true;
        for (int neighbor : graph.adjList[src]) {
            if (!visited[neighbor]) {
                if (dfs_helper(graph, neighbor, dest, visited))
11.
                    return true;
12.
        return false;
    bool dfs(const Graph& graph, int src, int dest)
        vector<bool> visited(graph.numVertices);
        return dfs helper(graph, src, dest, visited);
21. }
```

Problem with s-t Path

What if the edges are weighted?

The algorithms do not consider the weights.





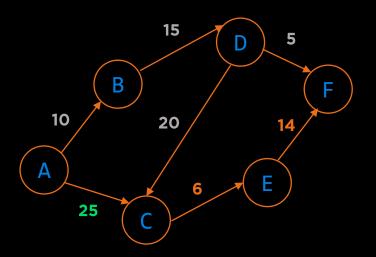
Problem with s-t Path

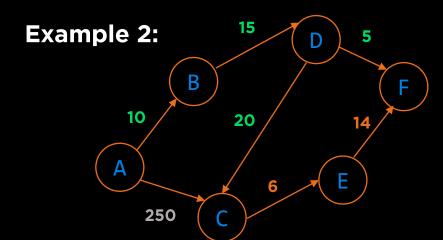
What if the edges are weighted?

The algorithms do not consider the weights.

Example 1: Path for A to C will be A-B-D-C for a DFS traversal which will have a total cost of 45 against 25 for the path directly from A-C.

Example 1:

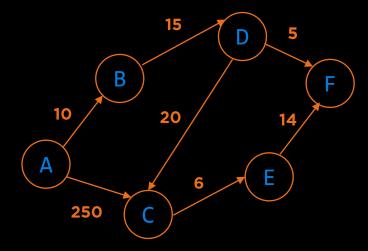




Example 2: Path for A to C will be A-C for a BFS traversal which might have a total cost of 250 against 45 for the path directly from A-B-D-C.

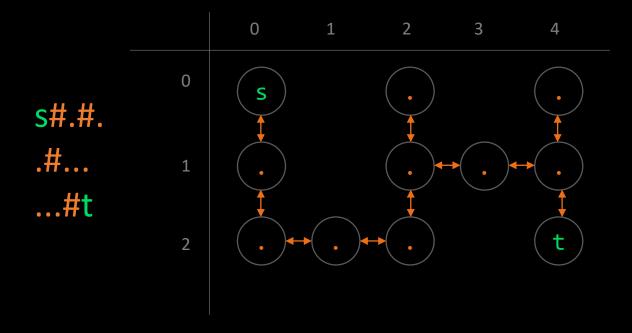


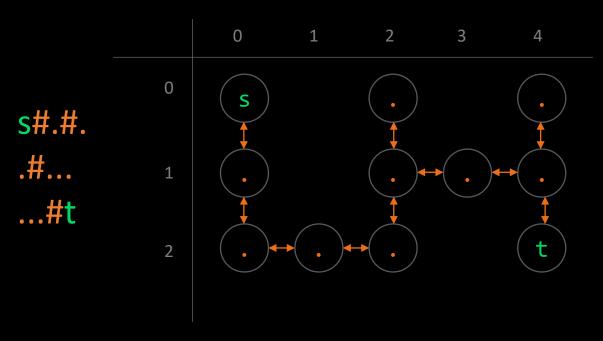
Shortest Weighted s-t Path



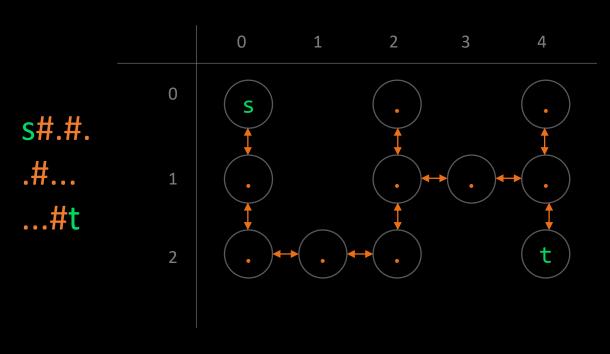
```
s#.#.
.#...
```



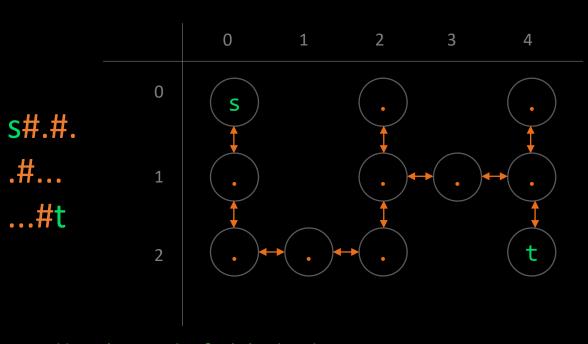




```
General BFS Algorithm
      Start from vertex, 's' (0,0), mark it identified, and
1.
     place it in a queue.
      while the queue is not empty
2.
3.
           Take a vertex, u, out of the queue and visit u.
4.
           for all vertices, v, adjacent to this vertex, u
                 if v has not been identified or visited
6.
                       Mark it identified
7.
                       Insert vertex \nu into the queue.
          We are now finished visiting u.
```



```
    General BFS Algorithm
    Start from vertex, 's' (0,0), mark it identified, and place it in a queue.
    while the queue is not empty
    Take a vertex, u, out of the queue and visit u.
    for all vertices, v, adjacent to this vertex, u
```

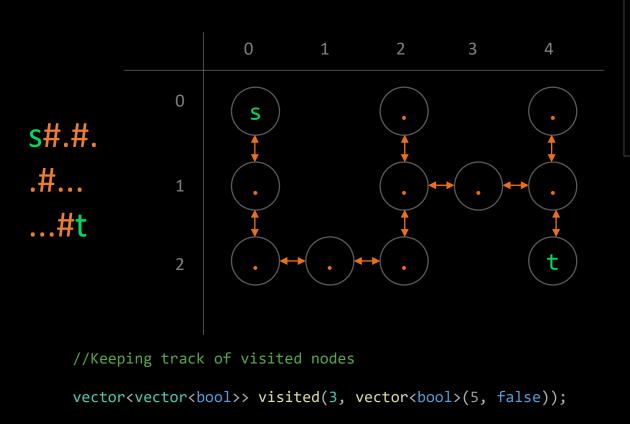


```
    General BFS Algorithm
    Start from vertex, 's' (0,0), mark it identified, and place it in a queue.
    while the queue is not empty
    Take a vertex, u, out of the queue and visit u.
    for all vertices, v, adjacent to this vertex, u
```

```
//Keeping track of visited nodes

vector<vector<bool>> visited(3, vector<bool>(5, false));
Visited[0][0] = true;
```

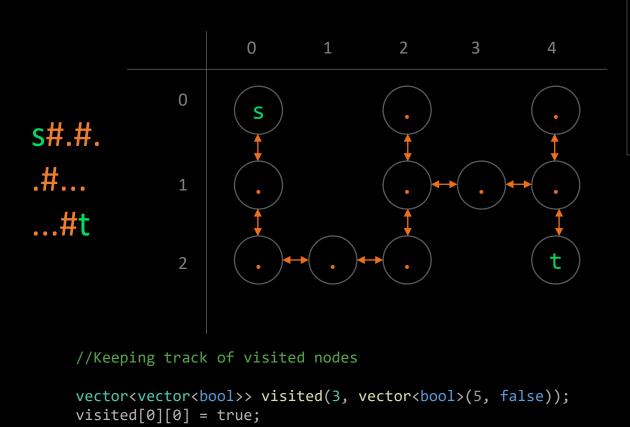




```
    General BFS Algorithm
    Start from vertex, 's' (0,0), mark it identified, and place it in a queue.
    while the queue is not empty
    Take a vertex, u, out of the queue and visit u.
    for all vertices, v, adjacent to this vertex, u
```

	0	1	2	3	4	
0	F	F	F	F	F	
1	F	F	F	F	F	
2	F	F	F	F	F	

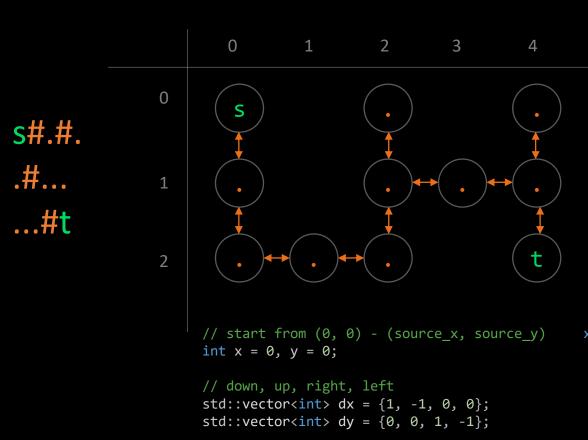




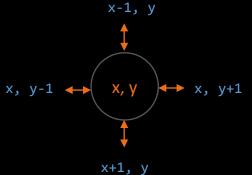
```
    General BFS Algorithm
    Start from vertex, 's' (0,0), mark it identified, and place it in a queue.
    while the queue is not empty
    Take a vertex, u, out of the queue and visit u.
    for all vertices, v, adjacent to this vertex, u
```

	0	1	2	3	4	
0	Т	F	F	F	F	
1	F	F	F	F	F	
2	F	F	F	F	F	

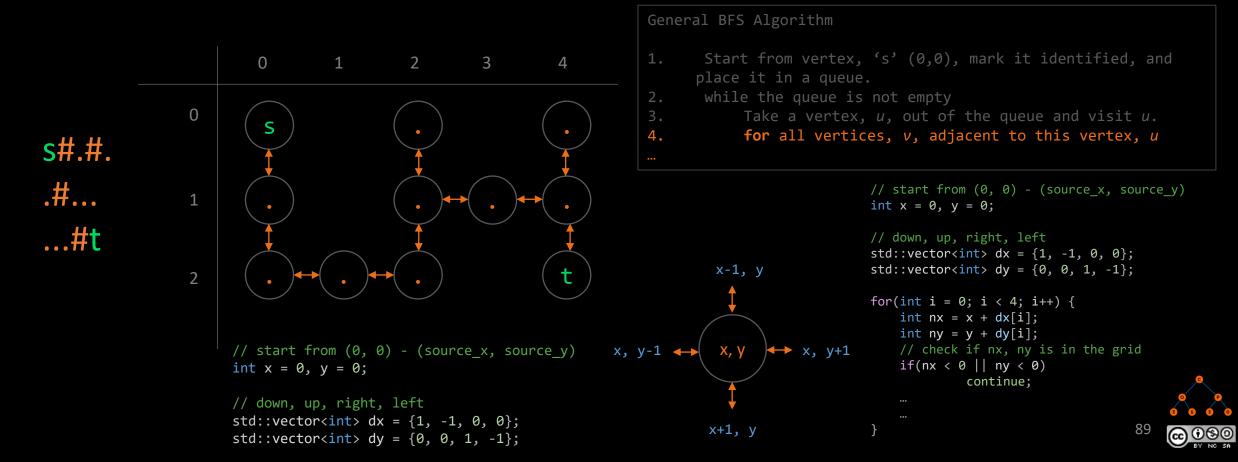




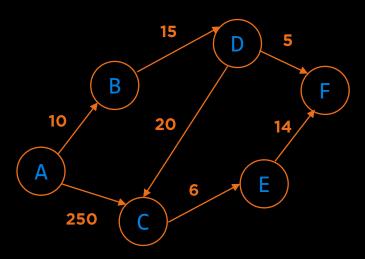
```
    General BFS Algorithm
    Start from vertex, 's' (0,0), mark it identified, and place it in a queue.
    while the queue is not empty
    Take a vertex, u, out of the queue and visit u.
    for all vertices, v, adjacent to this vertex, u
```







Shortest Weighted s-t Path



- Dijkstra's Algorithm
 - Single Source: Path to all vertices
 - Directed Graphs
 - No negative weights allowed
 - No negative weight cycles allowed
- Bellman Ford
 - Single Source: Path to all vertices
 - Negative Weights allowed
 - No negative weight cycles allowed
- Floyd-Warshall
 - All pair shortest paths
- A* Search