10/4/23 Having some over this leterase last week. I wanted to tack-e this problem while the content was still fresh in my mind yet not so new I would be able to call back exactly what

tinition from the READ ME: T(n) EO (f(n)) <=> ]C, no: T(n) ≤ c.f(n) 41 Z no. Recylling that two prove equivalency, we must Show that anythms belongers to the asymptotic ampletity at less (1) also belongs to that of log\_(n) and vice-versy, we can set to w-ork. Let's begin by subbing in our two f(n) expressions. T(n) { O (log, (n)) < => ] (, no: T(n) S (. (los, (n)) \ Ynzh T(n) { O (1093 (n)) < => ] (, no: T(n) 5 ( · (105, Ca)) \ \nzh

you'd thrown up on the board. Let's begin with the previded de.

Change of base formula is quite simple for loss:  $log_6(g) = \frac{log_8(g)}{log_8(b)}$ which can be rewritten as follows: \frac{1}{lug\_x(b)} lug\_x(b). Let's take these sesults back to our expressions above to finish out this problem: T(n) E O (109, (n1) < => ] (, no: T(n) 5 (3) · log (n) \ YNZA

This is a good start. Nov. to prove equivalency, it makes sense to write our expressions in terms at each other. Fortunarely the

Furtherman we see Something Similar shake out for O(los, (n)). T(n) E O (109, (n)) < => ] (, no: T(n) 5 (05, (s) VAZA T(n) & O (109, (n1) < => ] (, no: T(n) 5 d. lox, (n) Ynzh Where I is any arbitrary Constant. Essentially we have shown that T(n) E O los, (n) and Ten) & O los, (n). Expressed Symbolically. YTh) & Ollos, (n)) I(n) EO (loss (n)) HT(n) & O(loss(n)) T(n) & O(loss(n)) That words up this exercise: Fee I free to possible Coviections as there's a decent chance I messed something up with. our referencing the Mes for Thursday. Ictul time: ~ 2 hars

T(n) E O (109, (n1) < => ] (, no: T(n) 5 d. logs (n) Ynzh

Where I is any arbitrary constant.