

# Brute Force: Common Strategies

Keep it simple, keep it neat

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UTEC - Competitive Programming

## Fixing Variables

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# Fixing Variables

- One of the most intuitive techniques.
- Usually this problems involve an equation of  $m$  variables.
- You will *fix* a variable by making it constant and then try to find the solution under this new constrain.
- You will iterate through all possible values of the fixed variable and see how this fixed solution contributes to the general solution of the problem.

## Fixing Variables - Example

Given numbers  $n$ ,  $a$ ,  $b$  and  $c$  find:

$$|\{(x, y, z) | x, y, z \in \mathbb{Z}^{+0} \wedge x \leq a \wedge y \leq b \wedge z \leq c \wedge \frac{x}{2} + y + 2z = n\}|$$

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- Fix  $y$ :  $2z = n_1$
- Check if there exists a value of  $z$  that satisfies the equation.
- This solution has a complexity of  $O(ab)$ . Can we do better?



## **Simulation with Brute Force**

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# Simulation

- Problems that can be solved using this strategy usually have an underlying pattern or rules that we can compute easily.
- To find the answer we just need to let our simulation run until we get the expected results.
- **Problem:** Consider the following function:

$$f(n) = \begin{cases} \frac{n}{2} & n \equiv 0 \pmod{2} \\ 3n + 1 & \text{otherwise} \end{cases}$$

The cycle of a number  $n$  is the number of times we need to apply function  $f$  so that  $n = 1$ . Find the maximum cycle between two given numbers  $i$  and  $j$ .  $i, j < 10^5$

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- Run that function for all  $x$  between  $i$  and  $j$  and save the maximum.
- Complexity? Smallest cycle smaller than  $10^5 \approx 400$ , therefore  $O(400m)$ . Where  $m = \max\{i - j\}$

## Weak Constrains

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- This problems usually seem hard until you realize an underlying constrain that reduces significantly the search space.
- Even the problem might have some constrains, this might be very far from the constrains that can be deduced by the statement.
- After finding this new constrains we can use simple brute-force.



## Weak Constrains: Example

Given  $n$  and  $s$  ( $n, s \leq 10^{18}$ ), find:

$$|\{x | x \leq n \wedge x - f(x) \geq s\}|$$

Where  $f(x)$  is the sum of digits of  $x$ .

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- Observation 2:  $f(x) \leq 9 \cdot 18$
- Observation 3:  $x \geq s + 9 \cdot 18 \rightarrow x - f(x) \geq s$
- We only need to check  $x \in [s, s + 9 \cdot 18]$ !
- Answer will be  $\text{count}(s, s + 9 \cdot 18) + n - (s + 9 \cdot 18)$ .

**Thanks for Listening!**

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