



Lecture: BITWISE OPERATIONS
Unit: 9 - DYNAMIC PROGRAMMING II
Instructor: Nelson

INTEGERS \rightarrow BINARY

$$\begin{aligned} 23 &= 16 + \underline{7} = 16 + \underline{4 + 3} = 16 + 4 + 2 + 1 \\ &= 2^4 + 2^2 + 2^1 + 2^0 \\ &= \underline{10111}_{(2)} \end{aligned}$$

BINARY \rightarrow INTEGER

$$\begin{array}{ccccccc} 1 & 0 & 1 & 1 & 1 & = & 2^4 + 2^2 + 2^1 + 2^0 = 16 + 4 + 2 + 1 = 23 \\ 2^4 & 2^3 & 2^2 & 2^1 & 2^0 & & \end{array}$$

LOGICAL OPERATIONS

A	B	!A	A && B	A B	A ^ B
0	0	1	0	0	0
0	1	1	0	1	1
1	0	0	0	1	1
1	1	0	1	1	0

* NOTE: $A != B$ IS EQUIVALENT TO $A \wedge B$

A	B	!=	^
0	0	0	0
0	1	1	1
1	0	1	1
1	1	0	0

BITWISE OPERATIONS

→ COMPUTE LOGICAL OPERATIONS BIT BY BIT

$$23 = 10101_{(2)}$$

$$69 = 1000101_{(2)}$$

$$23 \oplus 69$$

$$\begin{array}{r} 0010101 \\ \oplus 1000101 \\ \hline 0000101 \rightarrow 5 \end{array}$$

BITWISE SHIFTS

① LEFT SHIFT: $a \ll b \Rightarrow a \cdot 2^b$

$$\left. \begin{array}{l} a = 5 = 101_{(2)} \\ b = 3 = 011_{(2)} \end{array} \right\} \begin{array}{l} 5 \ll 3 = 101_{(2)} \ll 3 = 101000_{(2)} = \underline{40} \\ 5 \ll 3 = 5 \cdot 2^3 = \underline{40} \end{array}$$

$$128 = 1000000_{(2)} = 1 \ll 7 = \underline{\underline{1 \cdot 2^7}}$$

② RIGHT SHIFT: $a \gg b \Rightarrow \left\lfloor \frac{a}{2^b} \right\rfloor$

↖ FLOOR FUNCTION

$$\left. \begin{array}{l} a = 5 = 101_{(2)} \\ b = 3 = 011_{(2)} \end{array} \right\} \begin{array}{l} 5 \gg 3 = 101_{(2)} \gg 3 = \cancel{0101} \rightarrow \\ 5 \gg 3 = \left\lfloor \frac{5}{2^3} \right\rfloor = 0 \end{array}$$

$$\left. \begin{array}{l} a = 23 = 10111_{(2)} \\ b = 3 = 011_{(2)} \end{array} \right\} \begin{array}{l} 23 \gg 3 = 10111_{(2)} \gg 3 = 10\cancel{11}\cancel{1}_{(2)} = 2 \\ 23 \gg 3 = \left\lfloor \frac{23}{2^3} \right\rfloor = \left\lfloor \frac{23}{8} \right\rfloor = 2 \end{array}$$

BINARY REPRESENTATION OF A NUMBER

$$x = 101001$$

$$y = 100101$$

$$\boxed{m = 8 = 1 \ll 3} \\ = 1000_{(2)}$$

$$\begin{array}{r} x \ll m \\ \swarrow \\ 101001 \\ \oplus 001000 \\ \hline 001000 \end{array}$$

$$\begin{array}{r} y \ll m \\ \swarrow \\ 100101 \\ \oplus 001000 \\ \hline 000000 \end{array}$$

$$\boxed{x \ll (1 \ll i)} \rightarrow i \text{ IS } \underline{\underline{i\text{TH BIT OF } X \text{ ON?}}}$$

IMPORTANT: OVERFLOW

$(1 \ll i) \rightarrow i$ AT MOST 30 SINCE 1 IS TREATED AS AN INTEGER

\Downarrow solution

$(1 \ll i) \rightarrow i$ AT MOST 62

\Downarrow

$1 \ll i \ll i$

BETTER SOLUTION: (AVOID OVERFLOWS)

\rightarrow CHECK IF $\boxed{(x \gg i) \& 1}$

BITSETS

↳ A BITMASK CAN ONLY HANDLE AS MUCH INTEGERS AS
THE COMPUTER WORD SIZE
↳ ~ 64

∴ WE CAN AT MOST REPRESENT $(1 \leq 63)$

SOLUTION?

↳ BITSETS → `bitset < 64 > a;`



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NO PROBLEM

- "LONGER BINARY MASK"
- "SET WITH QUICK INTERSECTIONS"

→ `bool A[20] : 20 BYTES`

→ `bitset<20> A : 20 BITS` ✓✓