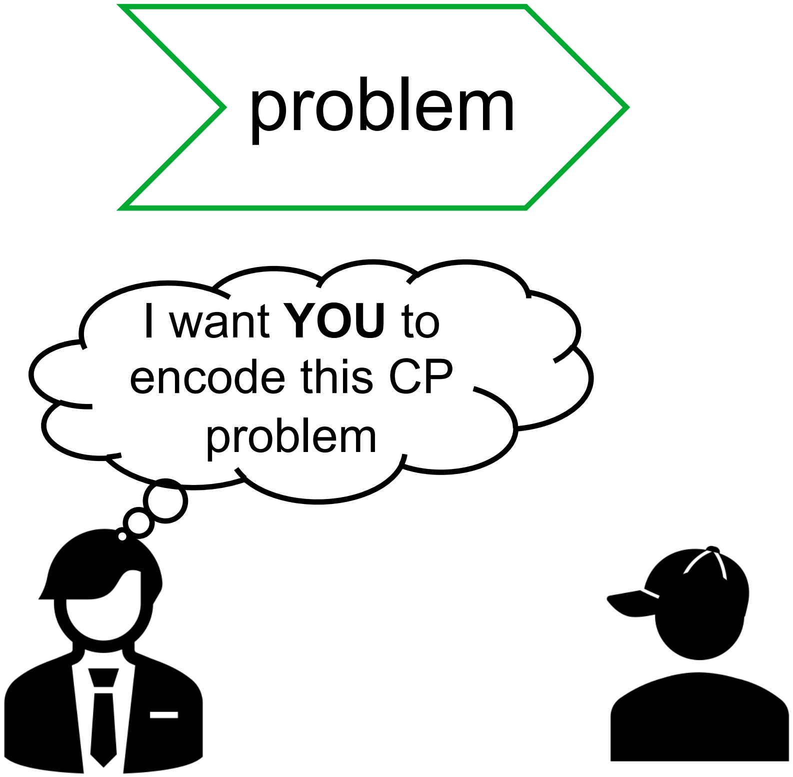
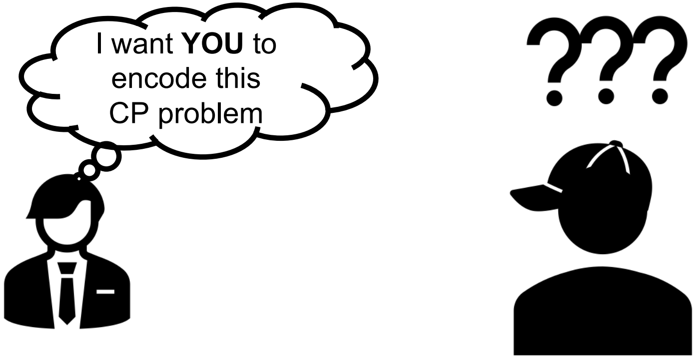
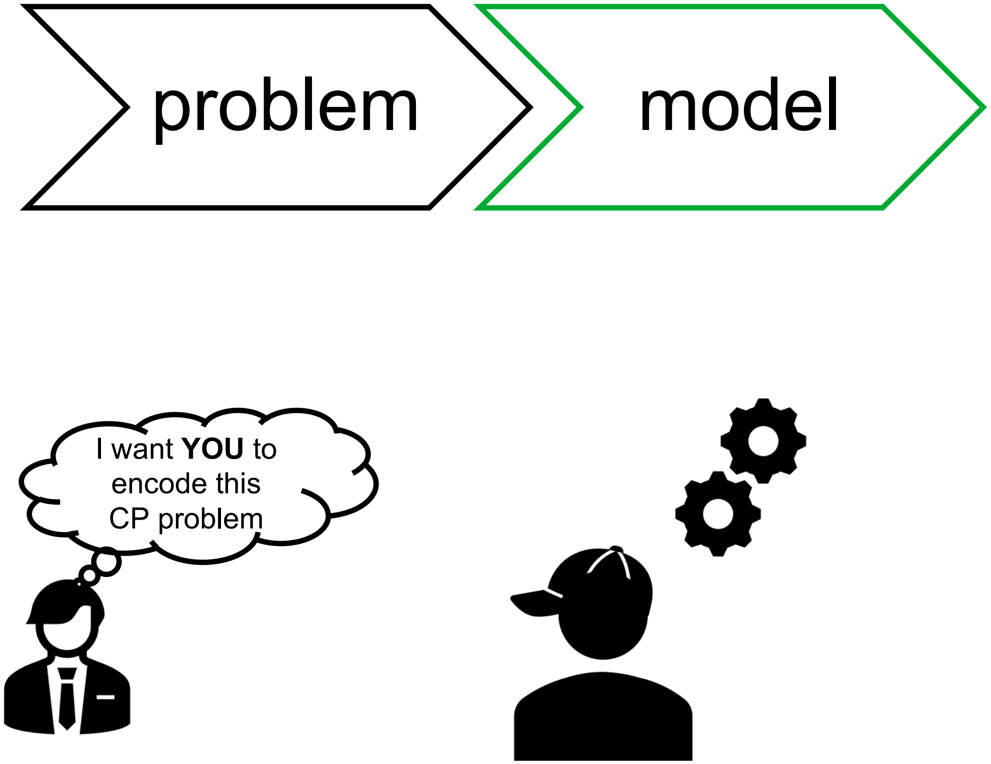
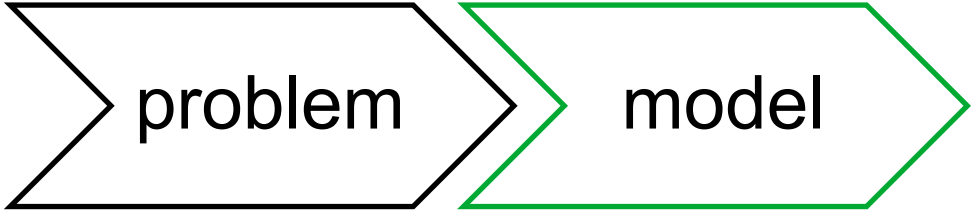

problem
I want **YOU** to encode this CP problem

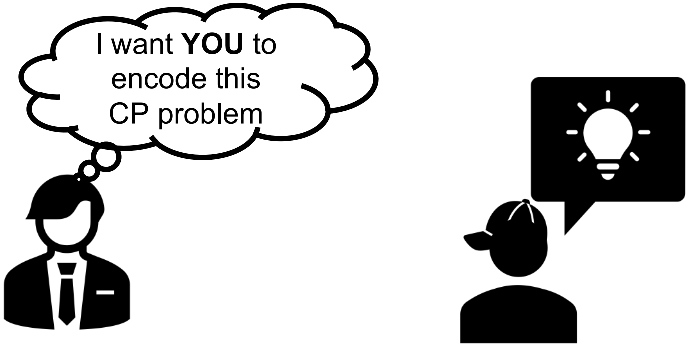
problem

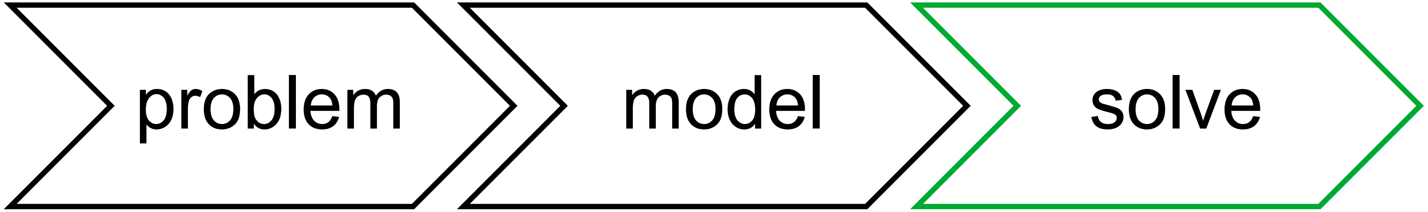


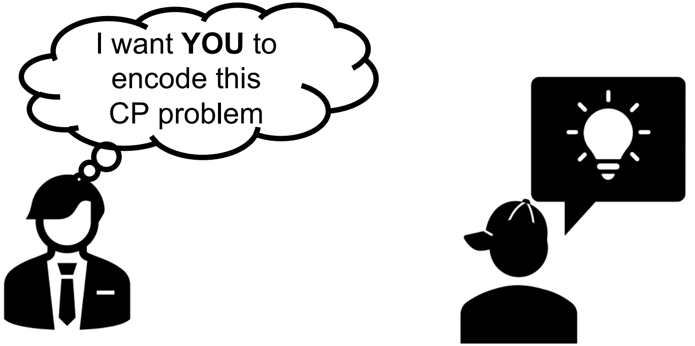

problem
model
I want **YOU** to encode this CP problem




```
x = intvar(-9, 9, name="x")
y = intvar(-9, 9, name="y")
m = Model(
    x < 0,
    x < 1,
    x > 2,
    (x + y > 0) | (y < 0),
    (y >= 0) | (x >= 0),
    (y < 0) | (x < 0),
    (y > 0) | (x < 0),
    AllDifferent(x,y)
)
```








```
x = intvar(-9, 9, name="x")
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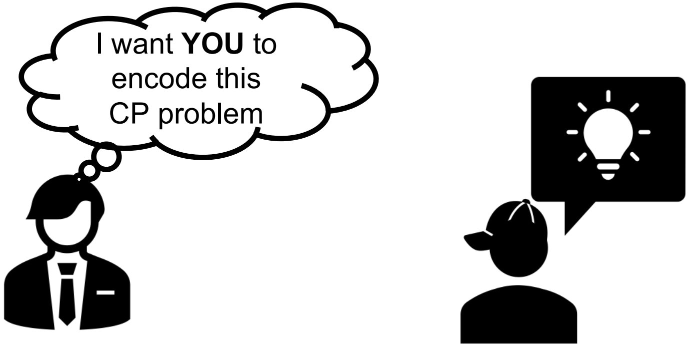
m = Model(
    x < 0,
    x < 1,
    x > 2,
    (x + y > 0) | (y < 0),
    (y >= 0) | (x >= 0),
    (y < 0) | (x < 0),
    (y > 0) | (x < 0),
    AllDifferent(x,y)
)
```

problem → model → solve → UNSAT

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x = intvar(-9, 9, name="x")
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        x < 0,
        x < 1,
        x > 2,
        (x + y > 0) | (y < 0),
        (y >= 0) | (x >= 0),
        (y < 0) | (x < 0),
        (y > 0) | (x < 0),
        AllDifferent(x,y)
)
```

Solver says UNSAT, what now?

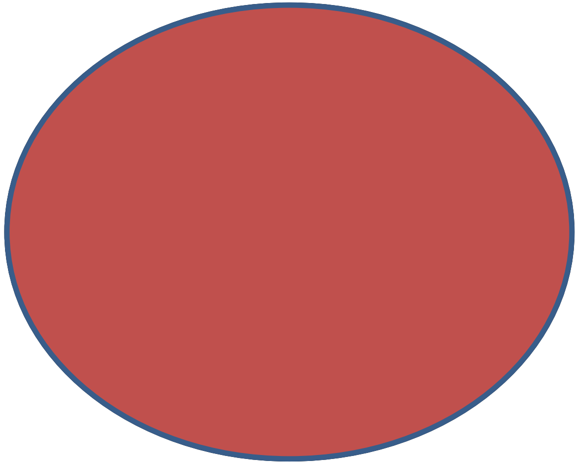
- (Human) modeling error ?
- Problem is over-constrained or unsatisfiable ?

Part 2: Explaining Unsatisfiability

with examples of Master ←→ Sub-problem solving

Debugging a model

= Need for an explanation of UNSAT



Model

Debugging a model

= Need for an explanation of UNSAT

1. Identify conflicting constraints as explanation for UNSAT

   → Extract Minimum Unsatisfiable Subset (MUS)

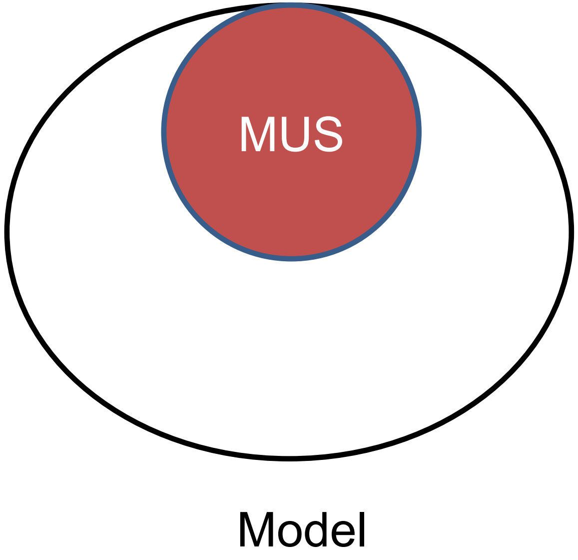
   a.k.a UNSAT Core, Irreducible Inconsistent Subsystem (IIS)

2. Identify Maximal Satisfiable Subset (MSS)

3. “Correct” the infeasibility in the model

   → Extract Minimum Correction Subsets (MCS)

   Complement of some MSS, removal/correction leads to a satisfiable subset



Debugging a model

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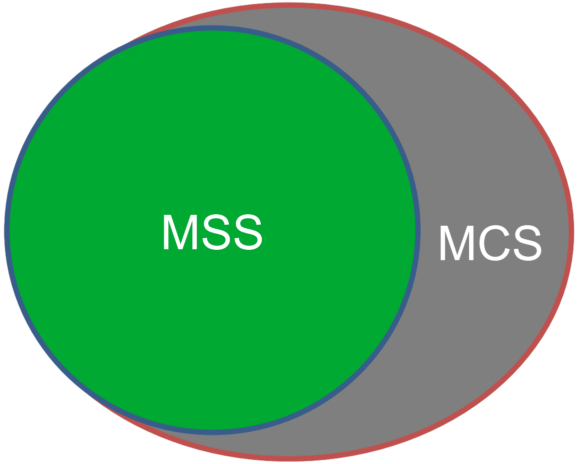

Debugging a model

= Need for an explanation of UNSAT

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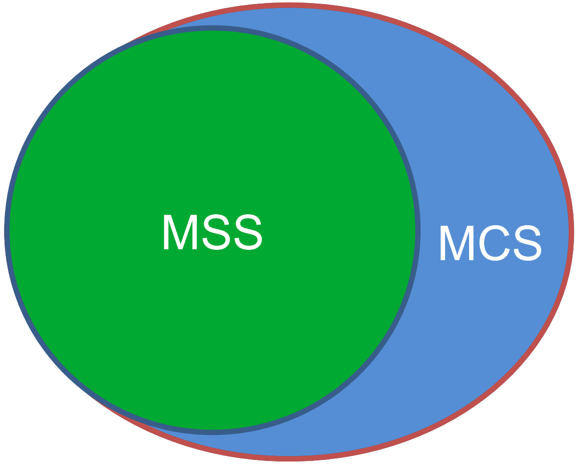

Debugging a model

= Need for an explanation of UNSAT

1. Identify <u>conflicting constraints</u> as explanation for UNSAT

   → Extract <u>Minimum</u> Unsatisfiable Subset (MUS)

   a.k.a UNSAT Core, Irreducible Inconsistent Subsystem (IIS)

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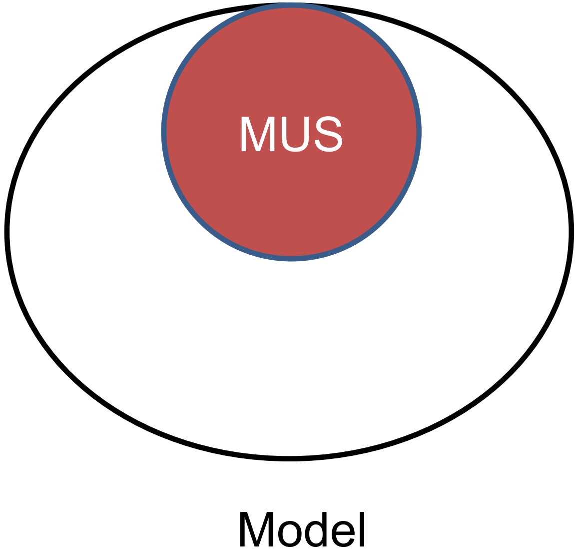
   Complement of some MSS, removal/correction leads to a satisfiable subset



Explaining UNSAT with MUSes

Methods

1. Some solvers provide an implementation for extracting unsatisfiable cores as explanations of UNSAT.
2. **Deletion-based** Minimal unsatisfiable subsets
   - Iterate over constraints
   - Delete constraints if removing them leaves the model UNSAT

```
def mus(constraints):
    m = Model(constraints)
    assert ~m.solve(), "MUS: model must be UNSAT"

    core = m.get_core()  # or all constraints
    i = 0
    while i < len(core):
        subcore = core[:i] + core[i+1:]  # check if all but i makes core SAT

        if Model(subcore).solve():
            i += 1  # removing it makes it SAT, must keep
        else:
            core = subcore  # overwrite core, so core[i] is next one

    return core
```

Joao Marques-Silva. *Minimal Unsatisfiability: Models, Algorithms and Applications*. ISMVL 2010. pp. 9-14

Example of MUS extraction

[*examples/tutorial_ijcai22/3_musx.ipynb*](examples/tutorial_ijcai22/3_musx.ipynb)

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger than', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
]
```

```python
core = m.get_core()  # or all constraints
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger than', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
]
```

```python
core = m.get_core()  # or all constraints
i = 0 # we will dynamically shrink mus vars
while i < len(core):
    # add all other remaining constraints
    subcore = core[:i] + core[i+1:]

    if Model(subcore).solve():
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS", core[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS", core[i])
        # overwrite current 'i' and continue
        core = subcore
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger than', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
]
```

```python
core = m.get_core()  # or all constraints
i = 0 # we will dynamically shrink mus vars
while i < len(core):
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    subcore = core[:i] + core[i+1:]

    if Model(subcore).solve():
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        print("\tSAT so in MUS", core[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS", core[i])
        # overwrite current 'i' and continue
        core = subcore
```

```
SAT so in MUS: (x[0]) > (x[1])
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger than', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
]
```

```python
core = m.get_core()  # or all constraints
i = 0 # we will dynamically shrink mus vars
while i < len(core):
    # add all other remaining constraints
    subcore = core[:i] + core[i+1:]

    if Model(subcore).solve():
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS", core[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS", core[i])
        # overwrite current 'i' and continue
        core = subcore
```

```
SAT so in MUS: (x[0]) > (x[1])
SAT so in MUS: (x[1]) > (x[2])
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger than', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
]
```

```python
core = m.get_core()  # or all constraints
i = 0 # we will dynamically shrink mus vars
while i < len(core):
    # add all other remaining constraints
    subcore = core[:i] + core[i+1:]

    if Model(subcore).solve():
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS", core[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS", core[i])
        # overwrite current 'i' and continue
        core = subcore
```

```
SAT so in MUS: (x[0]) > (x[1])
SAT so in MUS: (x[1]) > (x[2])
SAT so in MUS: (x[2]) > (x[0])
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger than', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
]
```

```python
core = m.get_core()  # or all constraints
i = 0 # we will dynamically shrink mus vars
while i < len(core):
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    if Model(subcore).solve():
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS", core[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS", core[i])
        # overwrite current 'i' and continue
        core = subcore
```

```
SAT so in MUS: (x[0]) > (x[1])
SAT so in MUS: (x[1]) > (x[2])
SAT so in MUS: (x[2]) > (x[0])
UNSAT so not in MUS: (x[3]) > (x[0])
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger than', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
]
```

```python
core = m.get_core()  # or all constraints
i = 0 # we will dynamically shrink mus vars
while i < len(core):
    # add all other remaining constraints
    subcore = core[:i] + core[i+1:]

    if Model(subcore).solve():
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS", core[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS", core[i])
        # overwrite current 'i' and continue
        core = subcore
```

```
SAT so in MUS: (x[0]) > (x[1])
SAT so in MUS: (x[1]) > (x[2])
SAT so in MUS: (x[2]) > (x[0])
UNSAT so not in MUS: (x[3]) > (x[0])
UNSAT so not in MUS: (((x[3]) > (x[1])) -> ((x[3]) > (x[2]))) and ((x[3] == 3) or ((x[1]) == (x[2])))
```

Explaining UNSAT with MUSes

Methods & Insights

- Some solvers provide unsatisfiable cores as a starting point for debugging
- Deletion-based Minimal unsatisfiable subsets

KEY Insights

x Depends on the ordering of the clauses

✓ Faster if the solver supports unsat core extraction and assumptions especially for larger problems

- ❑ Most solvers provide an assumption interface
- ❑ Does not require many modifications

```
i = 0 # we will dynamically shrink mus vars
while i < len(core):
    # add all other remaining constraints
    subcore = core[:i] + core[i+1:]

    if Model(subcore).solve():
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS", core[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS", core[i])
        # overwrite current 'i' and continue
        core = subcore
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger then', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x
]
```

```python
assum_model = Model()
# make assumption indicators, add reified constraints
ind = BoolVar(shape=len(mus_cons), name="ind")
for i,bv in enumerate(ind):
    assum_model += [bv.implies(mus_cons[i])]
# to map indicator variable back to soft_constraints
indmap = dict((v,i) for (i,v) in enumerate(ind))

assum_solver = CPM_ortools(assum_model)
assert (not assum_solver.solve(assumptions=ind)), "Model must be UNSAT"
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger then', UNSAT
mus_cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],

    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x
]
```

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assum_model = Model()
# make assumption indicators, add reified constraints
ind = BoolVar(shape=len(mus_cons), name="ind")
for i,bv in enumerate(ind):
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indmap = dict((v,i) for (i,v) in enumerate(ind))

assum_solver = CPM_ortools(assum_model)
assert (not assum_solver.solve(assumptions=ind)), "Model must be UNSAT"

# unsat core is an unsatisfiable subset
mus_vars = assum_solver.get_core()
print("UNSAT core of size", len(mus_vars))
```

```python
x = intvar(0,3, shape=4, name="x")
# circular 'bigger then', UNSAT
mus_cons = [
    x[0] > x[1],
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    x[2] > x[0],

    x[3] > x[0],
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assum_solver = CPM_ortools(assum_model)
assert (not assum_solver.solve(assumptions=ind)), "Model must be UNSAT"

# unsat core is an unsatisfiable subset
mus_vars = assum_solver.get_core()
print("UNSAT core of size", len(mus_vars))
```

```python
# now we shrink the unsatisfiable subset further
i = 0 # we wil dynamically shrink mus_vars
while i < len(mus_vars):
    # add all other remaining constraints
    assum_vars = mus_vars[:i] + mus_vars[i+1:]

    if assum_solver.solve(assumptions=assum_vars):
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS:", mus_cons[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS:", mus_cons[i])
        # overwrite current 'i' and continue
        mus_cons = testcons
```

```
UNSAT core of size 3
	SAT so in MUS: (x[0]) > (x[1])
	SAT so in MUS: (x[1]) > (x[2])
	SAT so in MUS: (x[2]) > (x[0])
```

Also use new core here

Example of MUS extraction

examples/tutorial_ijcai22/3_musx.ipynb

Explaining UNSAT with MUSes

Methods & Insights

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- Deletion-based Minimal unsatisfiable subsets

KEY Insights

✓ Depends on the ordering of the clauses

✓ Faster if the solver supports unsat core extraction and assumptions especially for larger problems

- ❑ Most solvers provide an assumption interface
- ❑ Does not require many modifications

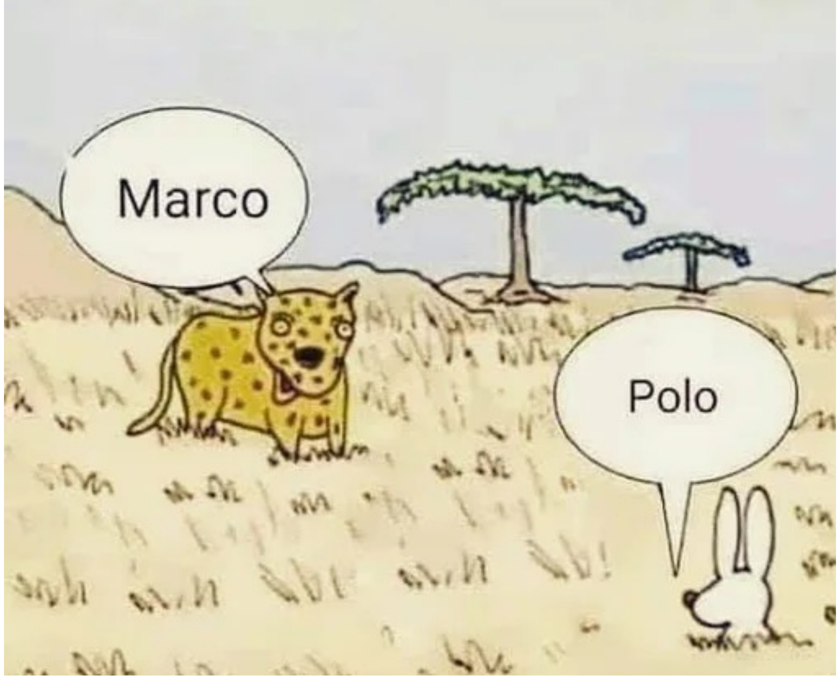
Enumerate all Minimal Unsatisfiable Subsets and Minimum Correction Subets

Master ←→ sub-problem

Marco Polo
(1254 –1324)

Marco-Polo

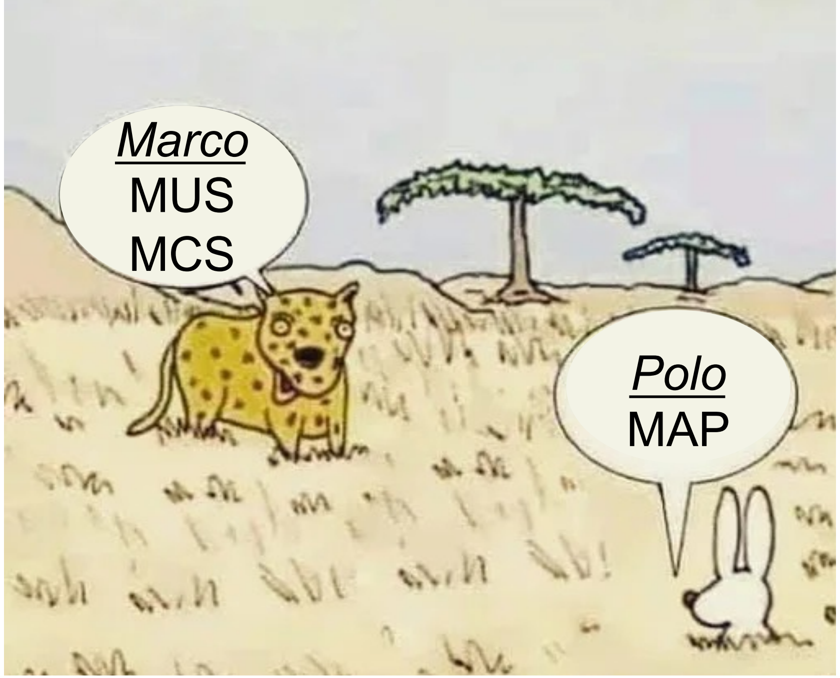
An example of Master ←→ Sub-problem approach



Source: https://9gag.com/gag/a07MMXZ

Marco-Polo

An example of Master ←→ Sub-problem approach



MARCO: **Ma**pping **R**egions of **Co**nstraints sets

➢ MUS enumeration algorithm

POLO: **Po**wer **Lo**gic

➢ "Map" of the powerset as a propositional logic Formula

[Liffiton, Mark H., and Ammar Malik. (2013) "Enumerating infeasibility: Finding multiple MUSes quickly." *CPAIOR 2013*.]

Source: https://9gag.com/gag/a07MMXZ

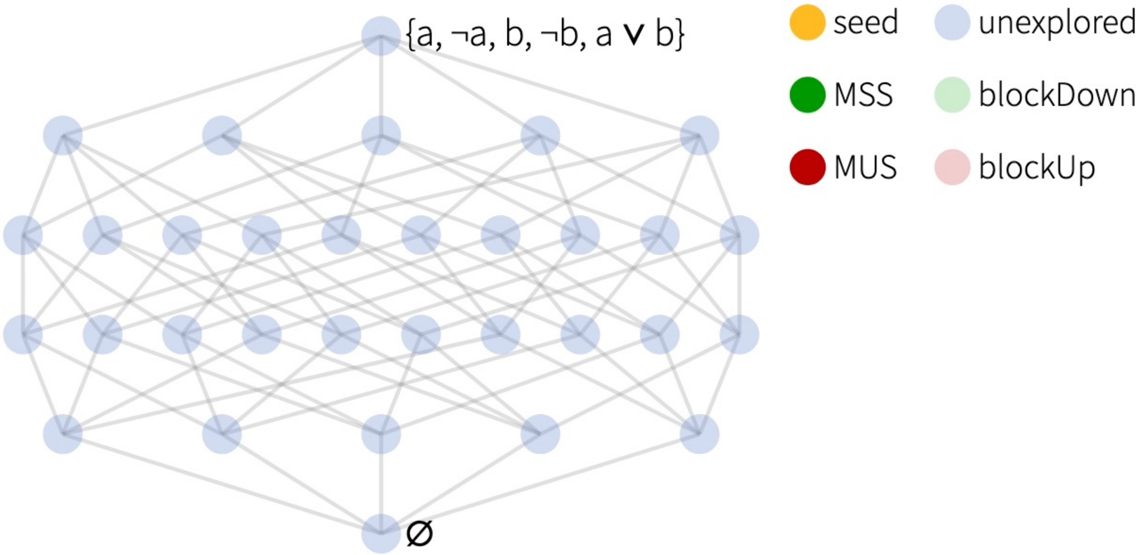
MARCO Algorithm

Input: Constraint system C

Map ← T
while Map is satisfiable:
    seed ← getUnexplored(Map)
    if seed is satisfiable:
        MSS ← grow(seed, C)
        output MSS
        Map ← Map ∧ blockDown(MSS)
    else:
        MUS ← shrink(seed, C)
        output MUS
        Map ← Map ∧ blockUp(MUS)

Example input: C = {a, ¬a, b, ¬b, a ∨ b}

seed :	
Map :	
MUSes :	
MSSes :	



Source: https://sun.iwu.edu/~mliffito/marco-viz/

Demo
MarcoPolo MUS/MCS enumeration

examples/tutorial_ijcai22/4_marco-mus-mcs-enumeration.ipynb

Explaining UNSAT with MUSes

Methods & Insights

- Deletion-based MUS extracts only 1MUS
- Efficiently enumerate all Minimal Unsatisfiable Subsets and Minimum Correction Subets
- No guarantee of subset-minimality (only heuristically)

  → Smallest Minimal Unsatisfiable Subset (SMUS)

  or optimality with weighted constraints

  → Optimal (Constrained) Unsatisfiable Subset (OCUS)

Master ←→ Sub-problem

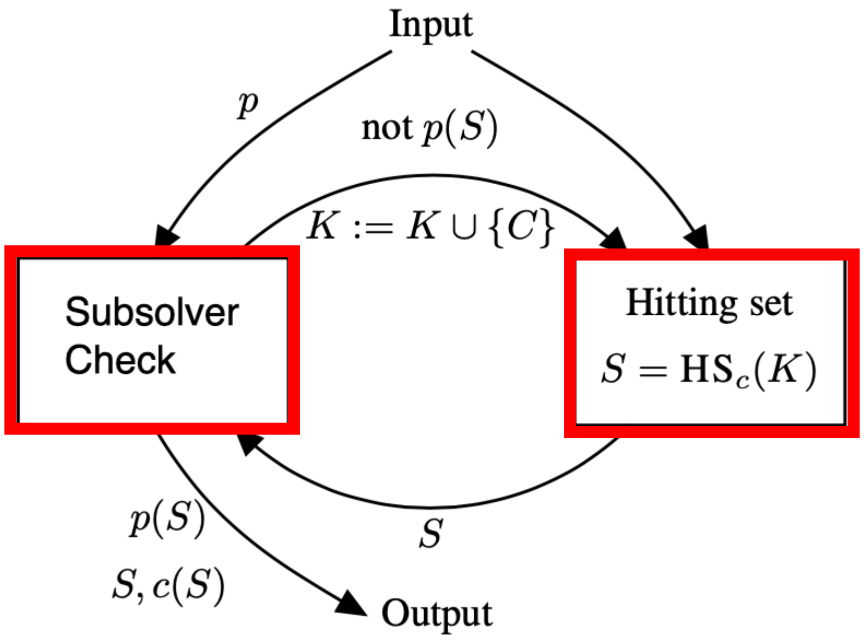
Implicit hitting set algorithms

Master ←→ Sub-problem approach

Implicit hitting set algorithms

General Structure:

1. Find minimum hitting set
2. Call subsolver for checking (SAT, UNSAT ….)



Saikko, Paul, Johannes P. Wallner, and Matti Järvisalo. "Implicit hitting set algorithms for reasoning beyond NP." *Fifteenth International Conference on the Principles of Knowledge Representation and Reasoning*. 2016.

Master ←→ Sub-problem approach

Implicit hitting set algorithms

***Smallest MUS* [1] and *Optimal (C)US* [2]:**

- Deciding whether a MUS of size ≤ k is $\sum_p^2 -complete$
- Extracting a smallest MUS (OCUS/SMUS) is in $FP^{\Sigma_p^2}$

Based on the <u>implicit hitting set</u> duality between **MCSs** and **MUSs**:

- Also used for MaxSAT, the dual of the OCUS-problem, i.e. MaxHS [3]

A set S ⊆ F is a MCS of F if and only if it is a *minimum hitting set* of MUSs(F).

A set S ⊆ F is a MUS of F if and only if it is a *minimum hitting set* of MUSs(F).

[1] E. Gamba, B. Bogaerts, T. Guns, Efficiently Explaining CSPs with Unsatisfiable Subset Optimization, IJCAI 2021
[2] A. Ignatiev, et al. "Smallest MUS extraction with minimal hitting set dualization." CP 2015.
[3] J. Davies, and B.Fahiem. "Exploiting the power of MIP solvers in MAXSAT." SAT 2013.

Optimal Constrained Unsatisfiable Subsets

Let

$\mathcal{F}$ be an unsatisfiable formula,

$f: 2^{\mathcal{F}} \rightarrow \mathbb{N}$ a *cost function*

$p: 2^{\mathcal{F}} \rightarrow \{t, f\}$ a *predicate (constraint)*

We call $\mathcal{S} \subseteq \mathcal{F}$ an OCUS of $\mathcal{F}$ (with respect to f and p) if

- $\mathcal{S}$ is unsatisfiable,
- $p(\mathcal{S})$ is true
- all other unsatisfiable $\mathcal{S}' \subseteq \mathcal{F}$ with $p(\mathcal{S}) = t$ satisfy $f(\mathcal{S}') \geq f(\mathcal{S})$.

[E. Gamba, B. Bogaerts, T. Guns, Efficiently Explaining CSPs with Unsatisfiable Subset Optimization, IJCAI 2021]

Optimal Constrained Unsatisfiable Subsets

Implicit hitting set-based algorithm

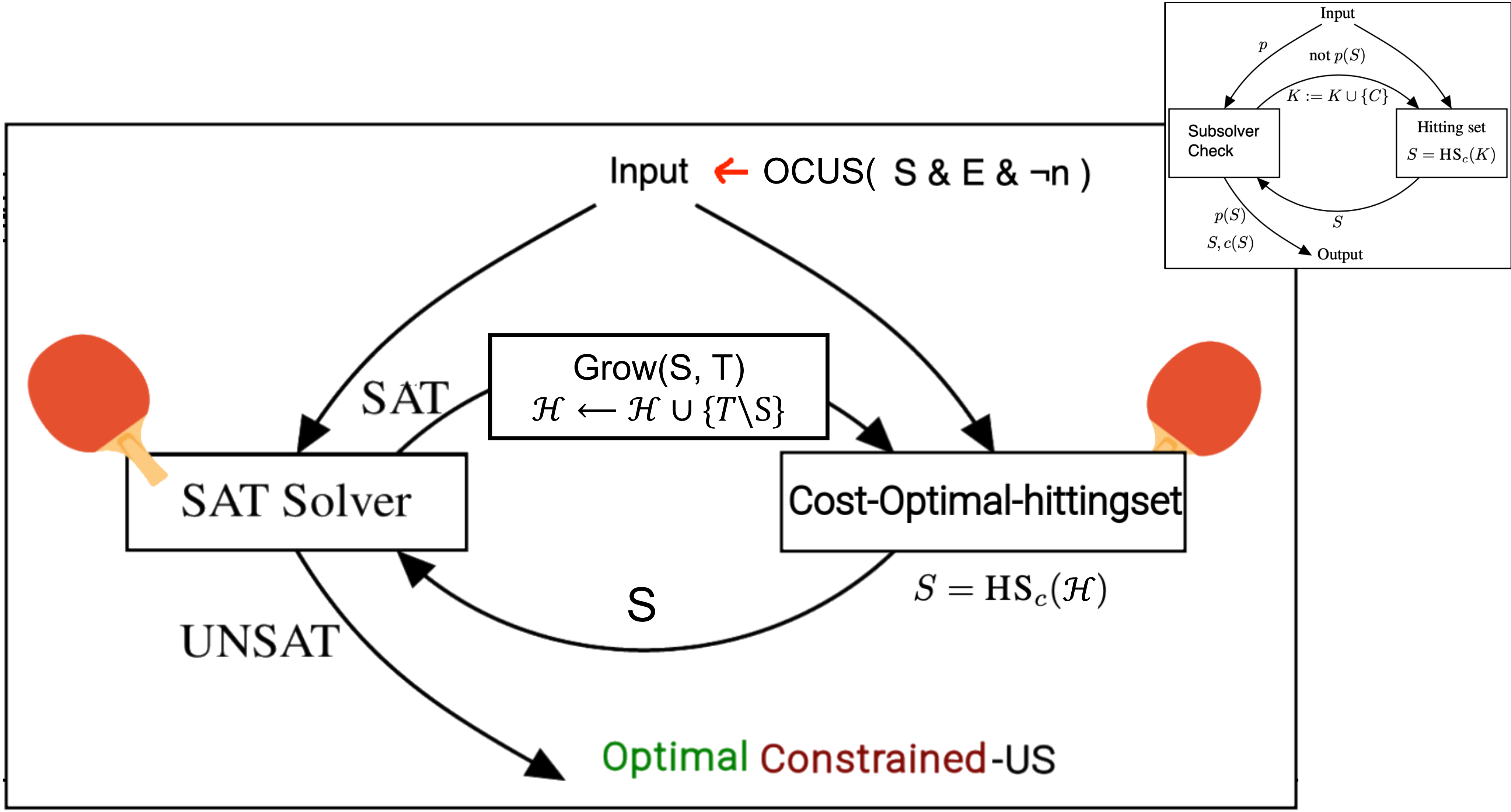
→ uses an implicit hitting set algorithm (like SMUS and MaxHS)

Algorithm 4: $\textsc{ocus}(T, f, p)$

$\mathcal{H} \leftarrow \emptyset$
while true **do**
 $\mathcal{S} \leftarrow \textsc{cost-optimal-hittingset}(\mathcal{H}, f, p)$ → **MIP solver**
 if $\neg\mathrm{SAT}(\mathcal{S})$ **then**
  return $\mathcal{S}$
 end
 $\mathcal{S} \leftarrow \textsc{grow}(\mathcal{S}, T)$
 $\mathcal{H} \leftarrow \mathcal{H} \cup \{T \setminus \mathcal{S}\}$
end

SAT/CP solver
as an oracle

[E. Gamba, B. Bogaerts, T. Guns, Efficiently Explaining CSPs with Unsatisfiable Subset Optimization, IJCAI 2021]



Input ← OCUS(S & E & ¬n)
SAT
Grow(S, T)
$\mathcal{H} \longleftarrow \mathcal{H} \cup \{T \setminus S\}$
SAT Solver
Cost-Optimal-hittingset
S
$S = \mathrm{HS}_c(\mathcal{H})$
UNSAT
Optimal Constrained-US
Input
p
not $p(S)$
$K := K \cup \{C\}$
Subsolver Check
Hitting set
$S = \mathrm{HS}_c(K)$
$p(S)$
$S, c(S)$
S
Output

OCUS *with* assumptions

```python
def OCUS_assum(soft, soft_weights, hard=[], solver='ortools', verbose=1):
    # init with hard constraints
    assum_model = Model(hard)

    # make assumption indicators, add reified constraints
    ind = BoolVar(shape=len(soft), name="ind")
    for i,bv in enumerate(ind):
        assum_model += [bv.implies(soft[i])]
    # to map indicator variable back to soft_constraints
    indmap = dict((v,i) for (i,v) in enumerate(ind))

    assum_solver = SolverLookup.lookup(solver)(assum_model)

    if assum_solver.solve(assumptions=ind):
        return []

    ##
    hs_model = Model(
        # Objective: min sum(x_l * w_l)
        minimize=sum(x_l * w_l for x_l, w_l in zip(ind, soft_weights))
    )

    # instantiate hitting set solver
    hittingset_solver = SolverLookup.lookup(solver)(hs_model)

    while(True):
        hittingset_solver.solve()

        # Get hitting set
        hs = ind[ind.value() == 1]

        if not assum_solver.solve(assumptions=hs):
            return soft[ind.value() == 1]

        # compute complement of model in formula F
        C = ind[ind.value() != 1]

        # Add complement as a new set to hit: sum x[j] * hij >= 1
        hittingset_solver += (sum(C) >= 1)
```

Hitting set Solver

repeatedly compute hitting sets

CP/SAT as an oracle

Extract Correction Subset

MUS extraction

[*examples/tutorial_ijcai22/5_ocus_explanations.ipynb*]

Optimal Constrained Unsatisfiable Subsets

Implicit hitting set-based algorithm

- Example of multi-solver incremental solving
- *Made easier and efficient with assumptions*

1. **Need to repeatedly compute hitting sets.**
   - Problem becomes hard as collection of sets-to-hits expands.
   - Big efficiency gains if incremental (and not restart)!

2. **CP/SAT is used as an oracle**
   - CP/SAT checking satisfiability of a subset
   - Grow solves a MaxSAT problem, s.t. complement is a (small) MCS

[E. Gamba, B. Bogaerts, T. Guns, Efficiently Explaining CSPs with Unsatisfiable Subset Optimization, IJCAI 2021]

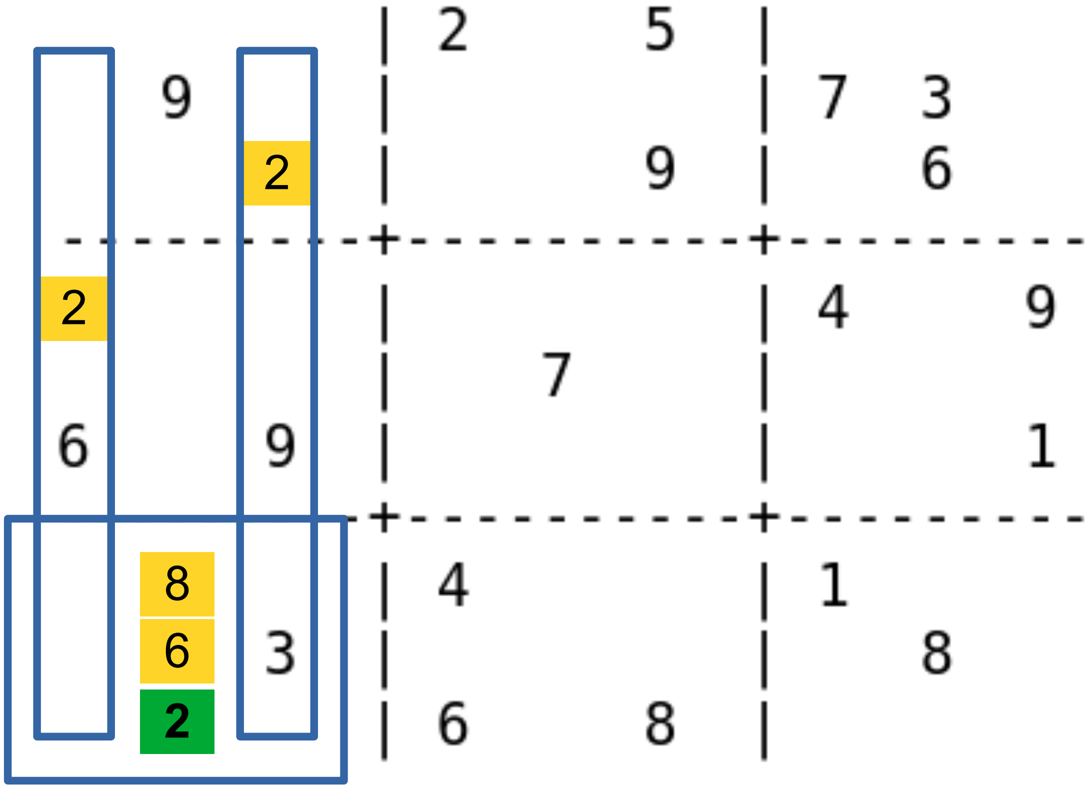
Part 3: Advanced examples of Master ←→ Sub-problem solving

Implicit Hitting set and Cutting plan algorithms

What if a model is SAT?

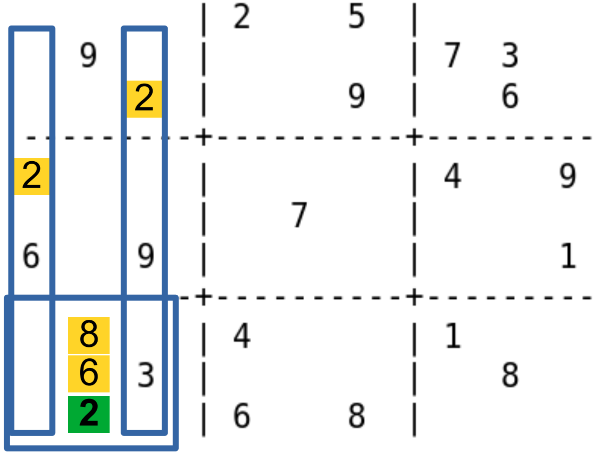
```
         | 2     5  |
   9     |          | 7  3
      2  |       9  |    6
---------+----------+---------
2        |          | 4     9
         |    7     |
6     9  |          |       1
---------+----------+---------
   8     | 4        | 1
   6  3  |          |    8
         | 6     8  |
```

What if a model is SAT?



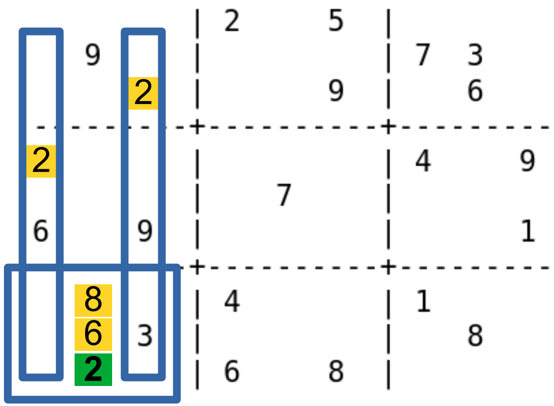
What if a model is SAT?

- User may not understand all derivations
- Or wants to learn about it



***“Explain** in a human-understandable way how to solve constraint satisfaction problems”*

Explanation step



Let *E'* & *S'* => ***n*** be one explanation step.

E' = a subset of previously derived facts E

(Sudoku) Given and derived digits in the grid

S' = a minimal subset of constraints S such that E' & S' => n

(Sudoku) Alldifferent column, row, box constraints

n = a newly derived fact

How? MUS(¬ n & E & S) is a valid explanation step

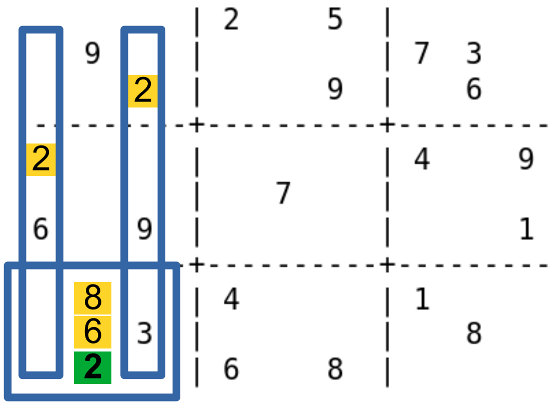
The best/easiest explanation step...

Let *f(S)* be a *cost function* that quantifies how good (e.g. easy to understand) an explanation step is.

Simple MUS-based algo:

```
sol-to-explain = propagate( E & S ) \ E

X_best = None
for n in sol-to-explain:
    X = MUS(~n & E & S )
    if f(X) < f(X_best):
        X_best = X
return X_best
```



MUS gives no guarantees on quality, only subset minimal (SMUS)

The best/easiest explanation step...

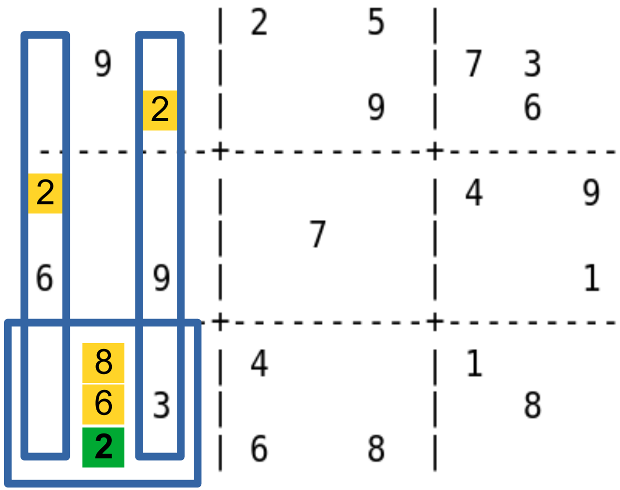
Let *f(S)* be a *cost function* that quantifies how good (e.g. easy to understand) an explanation step is.

Explain 1 step with OCUS

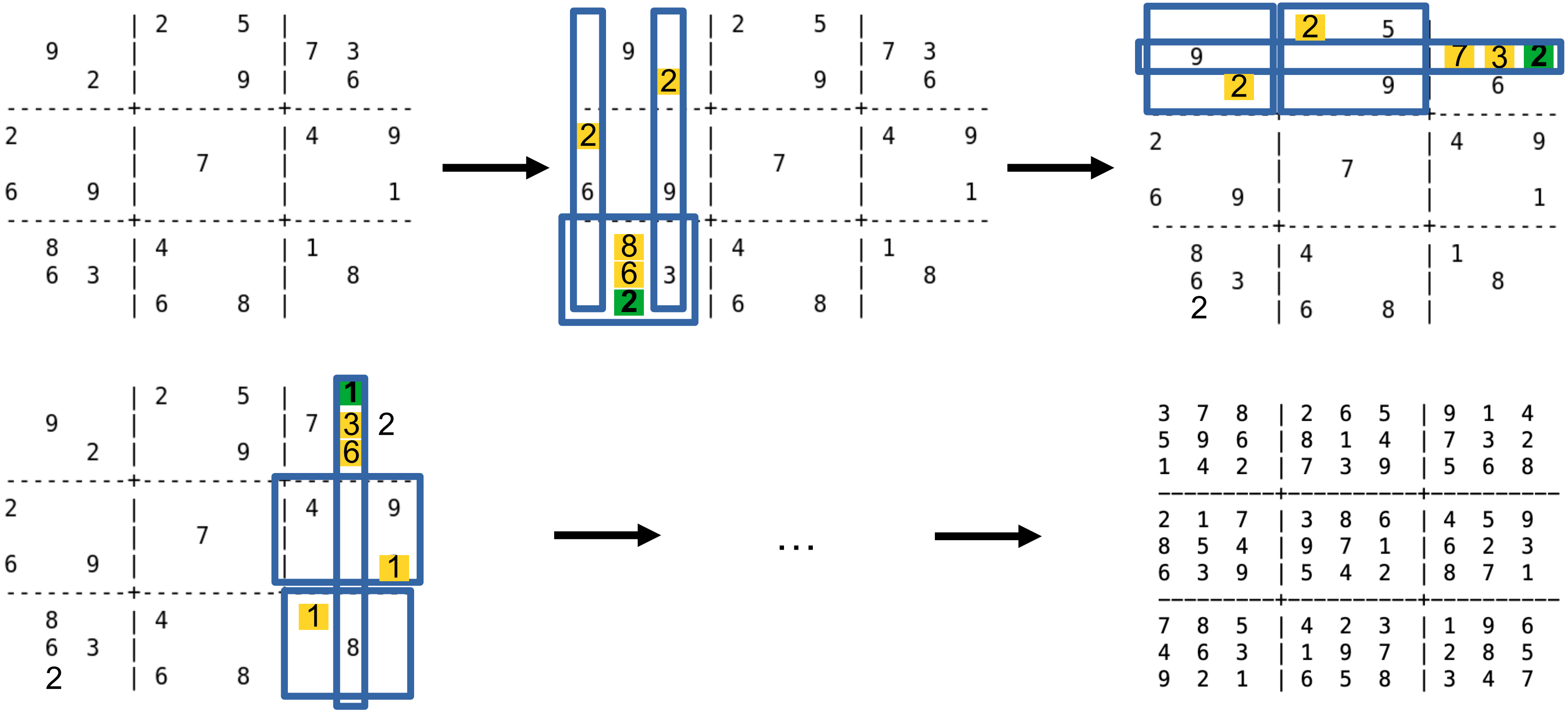
sol-to-explain = propagate(E & S) \ E

p = exactly-one({~n | n ∈ *sol-to-explain*}),

return OCUS(n | n ∈ sol-to-explain} & S & E &{~, **f**, **p**)



A sequence of explanations to explain SAT



Demo

OCUS for explaining SUDOKU

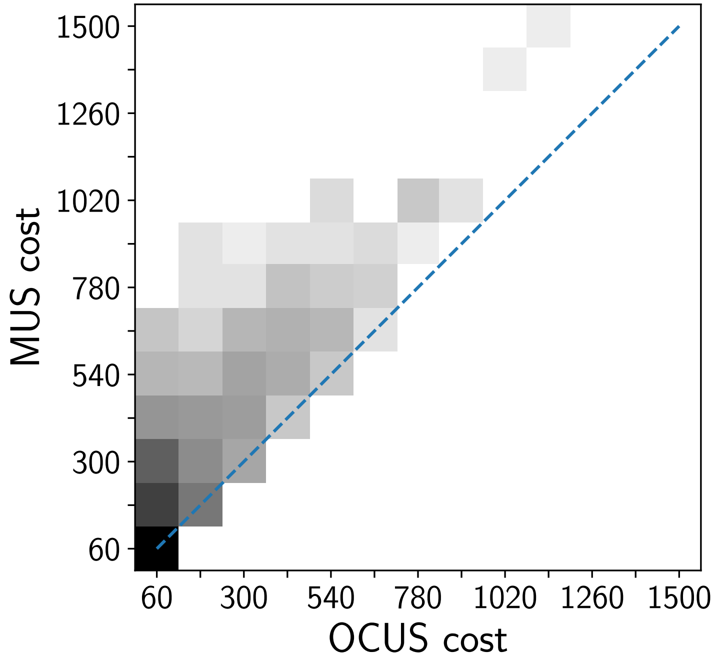
examples/tutorial_ijcai22/6_explain_sudoku.ipynb

The best/easiest explanation step...

sol-to-explain = propagate(*E* & $\boxed{\text{S}}$) \ E

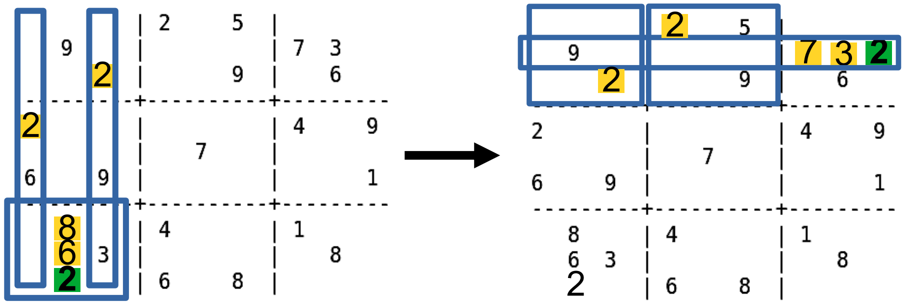
```
X_best = None
for n in sol-to-explain:
    X = MUS(~n & E & S )
    if f(X) < f(X_best):
        X_best = X
return X_best
```

p = exactly-one({~n|n ∈ sol-to-explain}),
OCUS({~n|n ∈ sol-to-explain}& E & S, f, p)



Incrementality at the Sequence-level

$$\textbf{OCUS}(\boxed{S} \;\&\; \boxed{E} \;\&\neg\; \boxed{n}\;, f, p)$$



$\boxed{S}$ Model constraints

- ➢ Do not change from an explanation step to another

$\boxed{E}$ Derived facts of E

- ➢ Precision-increasing!

Non-incremental OCUS - time (s

Incrementality at the Sequence-level

In practice

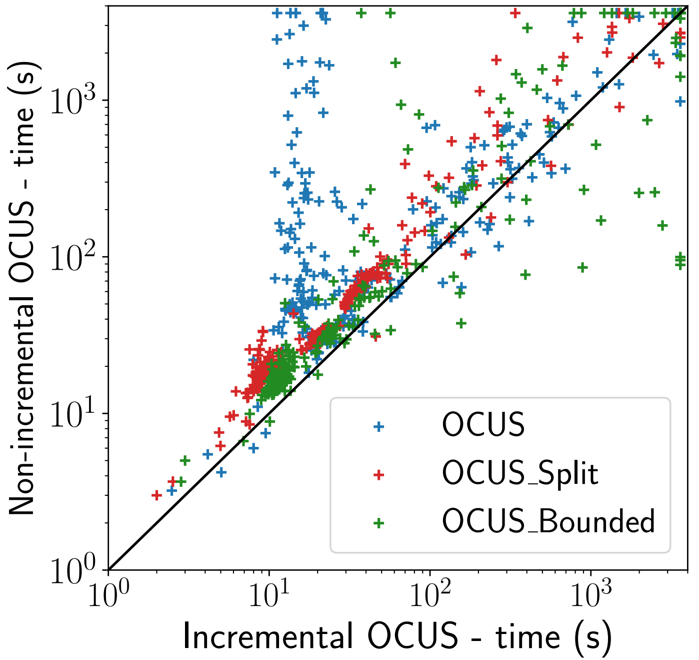
$$\mathrm{OCUS}(\boxed{S} \;\&\; E \;\&\neg\; \boldsymbol{n} \;, f, p)$$

Incremental OCUS works with the full unsatisfiable formula of step 0

$$\boxed{S} \;\&\; E_{end} \;\&\; \{\neg\, \boldsymbol{n} \mid \boldsymbol{n} \in \text{sol-to-explain}\}$$

Initialize hitting set solver **once** and modify objective at every explanation step *i* such that

- Underived facts cannot be taken
- Negated facts ($\neg$ **n**) already explained should <u>not</u> be selected
- Assumptions are used to deactivate unused clauses



Summary observations

(OCUS) Multi-solver Incremental solving

Multi-solver

- MIP – highly effective solving hittingset problem
- CP/SAT is used as an oracle for checking satisfiability of a subset

Incremental

- Need to *repeatedly compute hitting sets*.
- Problem becomes *hard* as collection of sets-to-hits expands.
- Big efficiency gains if incremental (and not restart)!
- *(Sequence)* Sets-to-hit re-used between explanation steps