Attack Lab Recitation

İlker Işık

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Introduction

- 4 attacks against two targets:
  - 3 code-injection attacks
  - 1 return-oriented programming attack
- Exploiting the buffer overflow security vulnerability.

Introduction

Objectives

Learning Outcomes

- Write safer programs.
- Understand stack and parameter passing mechanisms.
- Understand how x64 instructions are encoded.
- Gain more experience with OBJDUMP and GDB.

Logistics

- You have **2 weeks** starting from tomorrow (November 15th).
- Your target file will be delivered to you as a **feedback file** on the ODTUClass assignment page.

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- Your target file will be delivered to you as a **feedback file** on the ODTUClass assignment page.
- After that, **lab quiz** will be held, similarly to the bomb lab.
- You must get **at least 50 / 100** from the homework to qualify for the lab quiz.

Target Programs

- Two executables named CTARGET and RTARGET.
- Both target read from stardard input with `getbuf` function defined below:

GETBUF Function

```
unsigned getbuf()
{
    char buf[BUFFER_SIZE];
    Gets(buf);
    return 1;
}
```

Target Programs

GETBUF

- `Gets` works similarly to `gets`. Simply reads from `stdin` until it encounters `EOF`.
- Destination is an array `buf`, declared as having `BUFFER_SIZE` bytes.
  - They do not have a way to determine if the array is large enough to store the input.
  - This means that it is possible to overwrite the bounds allocated at destination.

Target Programs

GETBUF

- If the string is short it will return normally:

```
unix> ./ctarget
Cookie: 0x1a7dd803
Type string: Keep it short!
[enter CTRL+D after newline to terminate]
No exploit.  Getbuf returned 0x1
Normal return
```

- Typically an error occurs if you type a long string:

```
unix> ./ctarget
Cookie: 0x1a7dd803
[enter CTRL+D after newline to terminate]
Type string: This is not a very interesting string, but it
has the property ...
Ouch!: You caused a segmentation fault!
Better luck next time
```

Target Programs

GETBUF

- Both targets work in the same way.
- You need to feed special strings to CTARGET and RTARGET to achieve certain results. They are called *exploit* strings.

Target Programs

Arguments

Command line arguments for CTARGET and RTARGET:

- `-h`: Print list of possible command line arguments
- `-q`: Don't send results to the grading server. Offline working option.
- `-i FILE`: Supply input from a file, rather than from standard input

You can also use gdb to make sure your program work as intended:

Example

```
> gdb ./ctarget
(gdb) r -q
(gdb) r -i ctarget.l1.raw
(gdb) r -q -i ctarget.l1.raw
```

Target Programs

Important Points

- Your exploit strings will typically contain byte values that do not correspond to the ASCII values for printing characters. The program HEX2RAW will enable you to generate these *raw* strings.
- HEX2RAW expects two-digit hex values separated by one or more white spaces. So if you want to create a byte with a hex value of 0, you need to write it as 00.

Target Programs

Important Points

- When you have correctly solved one of the levels, your target program will automatically send a notification to the grading server.

Example

```
unix> ./hex2raw < ctarget.l2.txt | ./ctarget
Cookie: 0x1a7dd803
Type string:Touch2!: You called touch2(0xXXXXXXXX, 0xXXXXXXXX)
Valid solution for level 2 with target ctarget
PASSED: Sent exploit string to server to be validated.
NICE JOB!
```

- Unlike the Bomb Lab, there is no penalty for making mistakes in this lab.

Target Programs

Point Distribution

Phase	Program	Level	Method	Function	Points
1	CTARGET	1	CI	`touch1`	10
2	CTARGET	2	CI	`touch2`	25
3	CTARGET	3	CI	`touch3`	30
4	RTARGET	2	ROP	`touch2`	35

CI: Code injection

ROP: Return-oriented programming

Figure: Summary of attack lab phases

Scoreboard

- The scoreboard server will be up and running, using the same address:
  `http://144.122.71.31:15213/scoreboard`
- Your ranking is determined by these criteria in the given order:
  1. Total points (largest to smallest)
  2. Time of the last submission (earliest to latest)
- If you re-submit a phase that you've already solved, you may fall down in ranking.

Main Points I

- Your exploit strings will attack CTARGET in this part.
- Stack positions will be consistent from one run to the next and so that data on the stack can be treated as executable code.
- These features make the program vulnerable to attacks where the exploit strings contain the byte encodings of executable code.
- Function `getbuf` is called within CTARGET by a function `test` having the following C code:

```
void test() {
    int val;
    val = getbuf();
    printf("No exploit.  Getbuf returned 0x%x\n", val);
}
```

- When `getbuf` executes its return statement, the program normally resumes execution within function `test`. You need to change this behaviour.

Level 1 I

- For Phase 1, you will not inject new code. Your exploit string will redirect the program to execute an existing procedure. Its C representation is given below:

```
void touch1() {
    vlevel = 1;  // This is a part of the validation protocol
    printf("Touch1!: You called touch1() but you must not execute this part\n");
    fail(1);
    srand(331);  // Seed the RNG
    if (rand() % 42 != 0) exit(0);

    vlevel = 1;  // This is a part of the validation protocol
    printf("Touch1!: You called touch1() correctly\n");
    validate(1);
    srand(331);
    if (rand() % 42 != 0) exit(0);

    // ...
}
```

Level 1

Some Advice

- Exploit string for this level can be determined by examining a disassembled version of CTARGET. Use `objdump -d` to get this dissembled version.
- Be careful about byte ordering.
- You can use GDB to step the program through the last few instructions of `getbuf`.
- The address of the stack is consistent across runs but it's different for each student. You need to examine dissembled version to determine its position.

Level 2

- Your task is to get CTARGET to execute the code for `touch2` rather than returning to `test`. Its C representation given below:

```
void touch2(unsigned int val1, unsigned int val2, unsigned int val3) {
    vlevel = 2;  // Part of the validation protocol
    if (val1 == cookie && val2 == COMPUTE_VAL2(cookie) && val3 == COMPUTE_VAL3(cookie)) {
        printf("Touch2!: You called touch2(0x%.8x, 0x%.8x, 0x%.8x)\n", val1, val2, val3);
        validate(2);
    } else {
        printf("Misfire: You called touch2(0x%.8x, 0x%.8x, 0x%.8x)\n", val1, val2, val3);
        fail(2);
    }
    exit(0);
}
```

- You will need to return to `touch2` with the appropriate arguments.

Level 2

Some Advice

Some Advice:

- The first three arguments are passed in these registers in the given order:
  1. `%rdi`
  2. `%rsi`
  3. `%rdx`
- Do NOT use `jmp` or `call` instructions in your exploit code.
- You need generate the byte-level representations of instruction sequences for injection.

Level 3 I

- Your task is to get CTARGET to execute the code for `touch3` rather than returning to `test`. Its C representation given below:

```
void touch3(char *sval, unsigned short *nums) {
    vlevel = 3;  // Part of the validation protocol
    if (hexmatch(cookie, sval) && checknums(cookie, nums)) {
        printf("Touch3!: You called touch3(\"%s\")\n", sval);
        validate(3);
    } else {
        printf("Misfire: You called touch3(\"%s\")\n", sval);
        fail(3);
    }
    exit(0);
}
```

- First Argument: Null-terminated string containing the lowercase hexadecimal encoding of your cookie
- Second Argument: The same characters, but the array is processed by a macro and the values are `shorts`

Level 3

Some Advice

Some Advice:

- The cookie string should consist of eight hexadecimal digits (ordered from most to least significant) without a leading `0x`.
- Do not forget to put a 0 at the end of your string.
- Second argument should have 8 unsigned short characters consecutively. Each unsigned short is 2 bytes long.
- The functions `hexmatch`, `checknums` and `strncmp` push data onto the stack, overwriting portions of memory that held the buffer used by `getbuf`. You need be careful where to place your arrays.

Generating Byte Codes I

- You can use GCC as an assembler and OBJDUMP as a disassembler to generate byte codes. For example, suppose you write a file `example.s` containing the following assembly code:

```
   # Example of hand-generated assembly code
pushq   $0xabcdef           # Push value onto stack
addq    $17,%rax            # Add 17 to %rax
movl    %eax,%edx           # Copy lower 32 bits to %edx
```

- You can now assemble and disassemble this file:

```
unix> gcc -c example.s
unix> objdump -d example.o > example.d
```

Generating Byte Codes II

```
example.o:     file format elf64-x86-64


Disassembly of section .text:

0000000000000000 <.text>:
   0: 68 ef cd ab 00          pushq  $0xabcdef
   5: 48 83 c0 11             add    $0x11,%rax
   9: 89 c2                   mov    %eax,%edx
```

- From this file, you can get the byte sequence for the code:

```
68 ef cd ab 00 48 83 c0 11 89 c2
```

Generating Byte Codes III

- You can also add C-style comments to your string before feeding them to HEX2RAW.

```
68 ef cd ab 00   /* pushq  $0xabcdef  */
48 83 c0 11      /* add    $0x11,%rax */
89 c2            /* mov    %eax,%edx  */
```

Using HEX2RAW I

- HEX2RAW takes as input a *hex-formatted* string. In this format, each byte value is represented by two hex digits.
- Hex characters should be separated by whitespace.

Example

"012345" ⇒ 30 31 32 33 34 35 00

Using HEX2RAW I

Examples

There are several ways you can use HEX2RAW:

1. You can set up a series of pipes to pass the string through HEX2RAW.

```
unix> cat exploit.txt | ./hex2raw | ./ctarget
```

2. You can store the raw string in a file and use I/O redirection:

```
unix> ./hex2raw < exploit.txt > exploit-raw.txt
unix> ./ctarget < exploit-raw.txt
```

3. This approach can also be used when running from within GDB:

```
unix> gdb ctarget
(gdb) run < exploit-raw.txt
```

Using HEX2RAW II

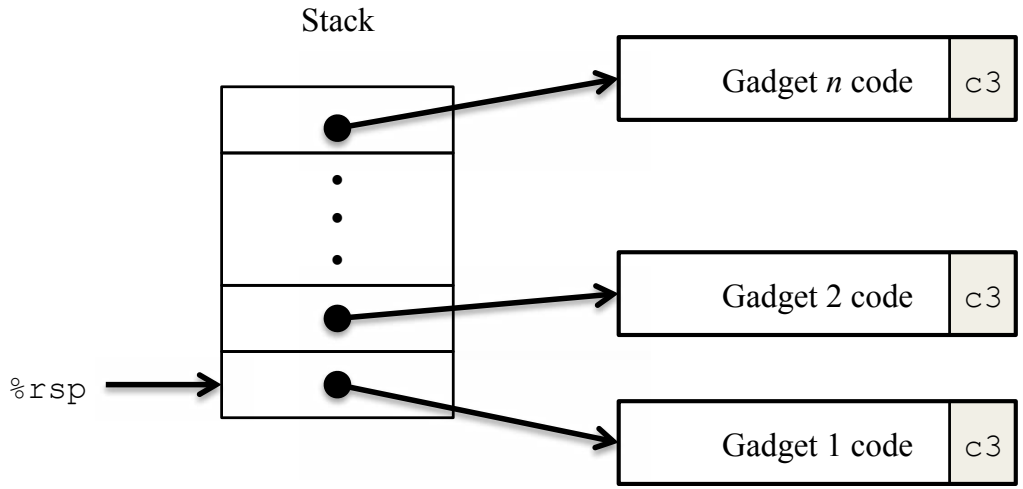
Examples

4. You can store the raw string in a file and provide the file name as a command-line argument:

```
unix> ./hex2raw < exploit.txt > exploit-raw.txt
unix> ./ctarget -i exploit-raw.txt
```

Return Oriented Programming I

- RTARGET uses two techniques to prevent code-injection.
  - Randomizes stack so that its position cannot be determined.
  - Makes the stack non-executable.
- Solution is to use existing code other than injecting new code.
- The strategy of ROP is to identify byte sequences followed by a return instruction. These are called gadgets and they can be chained using return instructions.



Return Oriented Programming II

Examples

- One version of RTARGET contains following code:

```
void setval_210(unsigned *p)
{
    *p = 3347663060U;
}
```

- When we look at the dissambled machine code we encounter:

```
0000000000400f15 <setval_210>:
  400f15:       c7 07 d4 48 89 c7   #movl   $0xc78948d4,(%rdi)
  400f1b:       c3                  #retq
```

where `48 89 c7` encodes the instruction `movq %rax, %rdi` followed by a `ret` instruction.