

OS-Course-Lab0 拆弹实验

phase0

```
0000000000400734 <phase_0>:
  400734: a9bf7bfd    stp x29, x30, [sp, #-16]!
  400738: 910003fd    mov x29, sp
  40073c: 94000126    bl 400bd4 <read_int>
  400740: 90000501    adrp x1, 4a0000 <.got.plt+0x18>
  400744: b9405421    ldr w1, [x1, #84]
  400748: 6b00003f    cmp w1, w0
  40074c: 54000061    b.ne 400758 <phase_0+0x24> // b.any
  400750: a8c17bfd    ldp x29, x30, [sp], #16
  400754: d65f03c0    ret
  400758: 940000e7    bl 400af4 <explode>
  40075c: 17ffffff    b 400750 <phase_0+0x1c>
```

首先需要理解 `read_int` 函数的作用

```
0000000000400bd4 <read_int>:
  400bd4: a9be7bfd    stp x29, x30, [sp, #-32]!
  400bd8: 910003fd    mov x29, sp
  400bdc: 910073e2    add x2, sp, #0x1c
  400be0: 90000321    adrp x1, 464000 <free_mem+0x40>
  400be4: 91214021    add x1, x1, #0x850
  400be8: 94001866    bl 406d80 <__isoc99_sscanf>
  400bec: 7100001f    cmp w0, #0x0
  400bf0: 5400008d    b.le 400c00 <read_int+0x2c>
  400bf4: b9401fe0    ldr w0, [sp, #28]
  400bf8: a8c27bfd    ldp x29, x30, [sp], #32
  400bfc: d65f03c0    ret
  400c00: 97ffffff    bl 400af4 <explode>
  400c04: d503201f    nop
  400c08: d503201f    nop
  400c0c: d503201f    nop
```

这个函数调用了 `__isoc99_sscanf`。`__isoc99_sscanf` 函数从输入中提取一个满足 `x1` 表达式的值，存入 `x2` 对应的内存中。在 `read_int` 函数中，`x2` 保存的地址是 `sp+0x1c`。调用 `__isoc99_sscanf` 后，`read_int` 函数再把 `sp+0x1c` 中的值存入 `w0` 并返回。通过gdb调试，发现 `x1` 读取的字符串是 `%d`，说明 `read_int` 函数的作用是从输入中提取一个数并通过 `x0` 返回。

```
pwndbg> x/s 0x464850
0x464850: "%d"
```

回到 `phase0`，在通过 `read_int` 函数读取了输入中的数字后，程序又从内存读取一个数到 `x1` 中。比较 `x1` 和 `x0`，若值相等，则通过。利用gdb，发现读取的数为 2022。则正确答案为 2022。

```

-----[ BACKTRACE ]-----
▶ 0      0x400748 phase_0+20
  1      0x4006b0 main+28
  2      0x400cb8 __libc_start_call_main+88
  3      0x40100c __libc_start_main_impl+792
  4      0x400570 _start+48

pwndbg> info registers x1
x1          0x7e6          2022

```

phase1

```

0000000000400760 <phase_1>:
400760: a9bf7bfd      stp x29, x30, [sp, #-16]!
400764: 910003fd      mov x29, sp
400768: 90000501      adrp    x1, 4a0000 <.got.plt+0x18>
40076c: f9402c21      ldr x1, [x1, #88]
400770: 94008504      bl 421b80 <strcmp>
400774: 35000060      cbnz    w0, 400780 <phase_1+0x20>
400778: a8c17bfd      ldp x29, x30, [sp], #16
40077c: d65f03c0      ret
400780: 940000dd      bl 400af4 <explode>
400784: 17ffffffd      b 400778 <phase_1+0x18>

```

`phase_1` 从内存中的读取一个字符串到 `x1` 中，和输入的字符串比较，若相同则通过

```

pwndbg> x/s *0x4a0058
0x464810: "KISS: Keep It Simple, Stupid."

```

利用gdb即可得到所需字符串

phase2

```

0000000000400788 <phase_2>:
400788: a9bc7bfd      stp x29, x30, [sp, #-64]!
40078c: 910003fd      mov x29, sp
400790: a90153f3      stp x19, x20, [sp, #16]
400794: 910083e1      add x1, sp, #0x20
400798: 940000f9      bl 400b7c <read_8_numbers>
40079c: b94023e0      ldr w0, [sp, #32] // w0 = number[1]
4007a0: 7100041f      cmp w0, #0x1 // w0==1
4007a4: 54000081      b.ne    4007b4 <phase_2+0x2c> // b.any
4007a8: b94027e0      ldr w0, [sp, #36] // w0 = number[0]
4007ac: 7100041f      cmp w0, #0x1
4007b0: 54000040      b.eq    4007b8 <phase_2+0x30> // b.none
4007b4: 940000d0      bl 400af4 <explode>
4007b8: 910083f3      add x19, sp, #0x20 // x19 <- number[] + 7
4007bc: 9100e3f4      add x20, sp, #0x38 // x20 <- number[] + 1
4007c0: 14000004      b 4007d0 <phase_2+0x48>
4007c4: 91001273      add x19, x19, #0x4 // x19 +=1

```

```

4007c8: eb14027f    cmp x19, x20
4007cc: 54000140    b.eq     4007f4 <phase_2+0x6c>  // b.none
4007d0: b9400260    ldr w0, [x19] // number[i]
4007d4: b9400661    ldr w1, [x19, #4] // number[i-1]
4007d8: 0b010000    add w0, w0, w1 // w0 = number[i]+number[i-1]
4007dc: 11000800    add w0, w0, #0x2 // w0 +=2
4007e0: b9400a61    ldr w1, [x19, #8] // w1 = number[i-2]
4007e4: 6b00003f    cmp w1, w0 // if w1 = w2
4007e8: 54ffffe0    b.eq     4007c4 <phase_2+0x3c>  // b.none
4007ec: 940000c2    bl 400af4 <explode>
4007f0: 17ffffff5    b 4007c4 <phase_2+0x3c>
4007f4: a94153f3    ldp x19, x20, [sp, #16]
4007f8: a8c47bfd    ldp x29, x30, [sp], #64
4007fc: d65f03c0    ret

```

首先调用了 `read_8_numbers` 函数，从输入中读取8个数字，并以数组保存在栈上。

phase2有两重检测

第一层，要求 `number[0]` 和 `number[1]` 都等于1才能通过

第二层是一个略微复杂的循环，改写成 C++ 代码大致如下：

```

for (int i = 7; i > 1; i--)
{
    if (number[i] + number[i - 1] + 2 != number[i - 2])
    {
        explode();
    }
}

```

而我们已知 `number[0]` 和 `number[1]`，很容易就能递推出整个数组为 `1 1 4 7 13 22 37 61`

phase3

```

0000000000400800 <phase_3>:
400800: a9be7bfd    stp x29, x30, [sp, #-32]!
400804: 910003fd    mov x29, sp
400808: 910063e3    add x3, sp, #0x18
40080c: 910073e2    add x2, sp, #0x1c
400810: 90000321    adrp x1, 464000 <free_mem+0x40>
400814: 911f6021    add x1, x1, #0x7d8
400818: 9400195a    bl 406d80 <__isoc99_sscanf>
40081c: 7100081f    cmp w0, #0x2
400820: 54000161    b.ne     40084c <phase_3+0x4c>  // b.any
400824: b9401fe0    ldr w0, [sp, #28]
400828: 7100081f    cmp w0, #0x2
40082c: 54000140    b.eq     400854 <phase_3+0x54>  // b.none
400830: 7100141f    cmp w0, #0x5
400834: 54000280    b.eq     400884 <phase_3+0x84>  // b.none
400838: 7100041f    cmp w0, #0x1
40083c: 54000320    b.eq     4008a0 <phase_3+0xa0>  // b.none
400840: 940000ad    bl 400af4 <explode>

```

```

400844: a8c27bfd    ldp x29, x30, [sp], #32
400848: d65f03c0    ret
40084c: 940000aa    bl 400af4 <explode>
400850: 17ffffff5    b 400824 <phase_3+0x24>
400854: b9401be2    ldr w2, [sp, #24]
400858: 528ccce0    mov w0, #0x6667 // #26215
40085c: 72acccc0    movk w0, #0x6666, lsl #16
400860: 9b207c40    smull x0, w2, w0
400864: 9362fc00    asr x0, x0, #34
400868: 4b827c00    sub w0, w0, w2, asr #31
40086c: 0b000801    add w1, w0, w0, lsl #2
400870: 4b010441    sub w1, w2, w1, lsl #1
400874: 0b000020    add w0, w1, w0
400878: 7100081f    cmp w0, #0x2
40087c: 54ffffe40    b.eq 400844 <phase_3+0x44> // b.none
400880: 9400009d    bl 400af4 <explode>
400884: b9401be0    ldr w0, [sp, #24]
400888: 4a800c00    eor w0, w0, w0, asr #3
40088c: 12000800    and w0, w0, #0x7
400890: b9401fe1    ldr w1, [sp, #28]
400894: 6b01001f    cmp w0, w1
400898: 54fffd60    b.eq 400844 <phase_3+0x44> // b.none
40089c: 94000096    bl 400af4 <explode>
4008a0: b9401be0    ldr w0, [sp, #24]
4008a4: b9401fe1    ldr w1, [sp, #28]
4008a8: 12000802    and w2, w0, #0x7
4008ac: 6b01005f    cmp w2, w1
4008b0: 54fffc40    b.eq 400844 <phase_3+0x44> // b.none
4008b4: d3431400    ubfx x0, x0, #3, #3
4008b8: 6b00003f    cmp w1, w0
4008bc: 54fffc40    b.eq 400844 <phase_3+0x44> // b.none
4008c0: 9400008d    bl 400af4 <explode>
4008c4: 17ffffdf    b 400840 <phase_3+0x40>

```

看起来相当复杂，实际上是三个并列的条件判断。只需满足其中之一即可。这里只选取 `w0 == 2` 的情况（即 `0x400854 - 0x400878`）。

`mov w0, #0x6667` 和 `movk w0, #0x6666, lsl #16` 两条指令是载入一个立即数 `0x66666667`。这个立即数是用来优化除以10算术过程的一个 magic number。结合后续过程，我们可以确定经过 `0x0x400854 - 0x400874` 后，`w0` 存储的实际上是我们输入的第二个数字整除10后的商和余数的和，即

```
w0 = number[1] % 10 + number[1] // 10
```

最后要求 `w0` 中保存的值为2,那么我们很快就能得到满足要求的输入为 `2 11`

phase4

```

00000000004009e4 <phase_4>:
4009e4: a9be7bfd    stp x29, x30, [sp, #-32]!
4009e8: 910003fd    mov x29, sp
4009ec: 90153f3     stp x19, x20, [sp, #16]
4009f0: aa0003f3    mov x19, x0

```

```

4009f4: 97fffe43    bl 400300 <.plt+0x60>
4009f8: aa0003f4    mov x20, x0
4009fc: 7100281f    cmp w0, #0xa
400a00: 540001ec    b.gt 400a3c <phase_4+0x58>
400a04: 2a1403e1    mov w1, w20
400a08: aa1303e0    mov x0, x19
400a0c: 97ffffaf    bl 4008c8 <encrypt_method1>
400a10: 2a1403e1    mov w1, w20
400a14: aa1303e0    mov x0, x19
400a18: 97ffffd3    bl 400964 <encrypt_method2>
400a1c: 90000500    adrp x0, 4a0000 <.got.plt+0x18>
400a20: f9403401    ldr x1, [x0, #104]
400a24: aa1303e0    mov x0, x19
400a28: 94008456    bl 421b80 <strcmp>
400a2c: 350000c0    cbnz w0, 400a44 <phase_4+0x60>
400a30: a94153f3    ldp x19, x20, [sp, #16]
400a34: a8c27bfd    ldp x29, x30, [sp], #32
400a38: d65f03c0    ret
400a3c: 9400002e    bl 400af4 <explode>
400a40: 17ffffff1    b 400a04 <phase_4+0x20>
400a44: 9400002c    bl 400af4 <explode>
400a48: 17ffffffa    b 400a30 <phase_4+0x4c>

```

要求输入的字符串经过两次加密后与某一固定字符串匹配。

```

00000000004008c8 <encrypt_method1>:
4008c8: a9be7bfd    stp x29, x30, [sp, #-32]!
4008cc: 910003fd    mov x29, sp
4008d0: 910043e2    add x2, sp, #0x10
4008d4: 3821c85f    strb wzr, [x2, w1, sxtw]
4008d8: 0b417c23    add w3, w1, w1, lsr #31
4008dc: 13017c63    asr w3, w3, #1
4008e0: 7100043f    cmp w1, #0x1
4008e4: 540003cd    b.le 40095c <encrypt_method1+0x94>
4008e8: aa0203e4    mov x4, x2
4008ec: d2800002    mov x2, #0x0 // #0
4008f0: d37ff845    lsl x5, x2, #1
4008f4: 38656805    ldrb w5, [x0, x5]
4008f8: 38001485    strb w5, [x4], #1
4008fc: 91000442    add x2, x2, #0x1
400900: 6b02007f    cmp w3, w2
400904: 54ffff6c    b.gt 4008f0 <encrypt_method1+0x28>
400908: 7100007f    cmp w3, #0x0
40090c: 1a9fc465    csinc w5, w3, wzr, gt
400910: 6b05003f    cmp w1, w5
400914: 540001cd    b.le 40094c <encrypt_method1+0x84>
400918: 4b050024    sub w4, w1, w5
40091c: 4b0300a2    sub w2, w5, w3
400920: 8b22c402    add x2, x0, w2, sxtw #1
400924: 91000442    add x2, x2, #0x1
400928: d2800001    mov x1, #0x0 // #0
40092c: 910043e3    add x3, sp, #0x10
400930: 8b25c065    add x5, x3, w5, sxtw

```

```

400934: d37ff823    lsl x3, x1, #1
400938: 38636843    ldrb w3, [x2, x3]
40093c: 382168a3    strb w3, [x5, x1]
400940: 91000421    add x1, x1, #0x1
400944: eb04003f    cmp x1, x4
400948: 54ffff61    b.ne 400934 <encrypt_method1+0x6c> // b.any
40094c: 910043e1    add x1, sp, #0x10
400950: 940084dc    bl 421cc0 <strcpy>
400954: a8c27bfd    ldp x29, x30, [sp], #32
400958: d65f03c0    ret
40095c: 52800005    mov w5, #0x0 // #0
400960: 17ffffec    b 400910 <encrypt_method1+0x48>

```

0000000000400964 <encrypt_method2>:

```

400964: 7100003f    cmp w1, #0x0
400968: 540003cd    b.le 4009e0 <encrypt_method2+0x7c>
40096c: a9bd7bfd    stp x29, x30, [sp, #-48]!
400970: 910003fd    mov x29, sp
400974: a90153f3    stp x19, x20, [sp, #16]
400978: a9025bf5    stp x21, x22, [sp, #32]
40097c: aa0003f3    mov x19, x0
400980: 8b21c015    add x21, x0, w1, sxtw
400984: 90000516    adrp x22, 4a0000 <.got.plt+0x18>
400988: 910162d6    add x22, x22, #0x58
40098c: 14000009    b 4009b0 <encrypt_method2+0x4c>
400990: 39400281    ldrb w1, [x20]
400994: f94006c0    ldr x0, [x22, #8]
400998: 8b010000    add x0, x0, x1
40099c: 3859f000    ldurb w0, [x0, #-97]
4009a0: 39000280    strb w0, [x20]
4009a4: 91000673    add x19, x19, #0x1
4009a8: eb15027f    cmp x19, x21
4009ac: 54000120    b.eq 4009d0 <encrypt_method2+0x6c> // b.none
4009b0: aa1303f4    mov x20, x19
4009b4: 39400260    ldrb w0, [x19]
4009b8: 51018400    sub w0, w0, #0x61
4009bc: 12001c00    and w0, w0, #0xff
4009c0: 7100641f    cmp w0, #0x19
4009c4: 54ffffe69    b.ls 400990 <encrypt_method2+0x2c> // b.plast
4009c8: 9400004b    bl 400af4 <explode>
4009cc: 17ffffff1    b 400990 <encrypt_method2+0x2c>
4009d0: a94153f3    ldp x19, x20, [sp, #16]
4009d4: a9425bf5    ldp x21, x22, [sp, #32]
4009d8: a8c37bfd    ldp x29, x30, [sp], #48
4009dc: d65f03c0    ret
4009e0: d65f03c0    ret

```

`encrypt_method1` 将字符串中下标为偶数的字符放在前面，下标为奇数的放在后面。

`encrypt_method2` 要求输入的字符串仅含小写字母，且以字母的ascii码为索引，进行替换。

关键需要获取到 `method2` 中用到的替换表。利用gdb即可。

```
pwndbg> x/s *0x4a0060
0x4647f0:      "qwertyuiopasdfghjklzxcvbnm"
```

再利用gdb获取目标字符串

```
pwndbg> x/s *0x4a0068
0x4647e0:      "isggstsvkw"
```

即可逆向推出输入字符串为 `hellor1b`

phase5

```
000000000400ac0 <phase_5>:
 400ac0: a9bf7bfd      stp x29, x30, [sp, #-16]!
 400ac4: 910003fd      mov x29, sp
 400ac8: 94000043      bl 400bd4 <read_int>
 400acc: 90000501      adrp x1, 4a0000 <.got.plt+0x18>
 400ad0: 91016021      add x1, x1, #0x58
 400ad4: 91006021      add x1, x1, #0x18
 400ad8: 97ffffdd      bl 400a4c <func_5>
 400adc: 7100c1f       cmp w0, #0x3
 400ae0: 54000061      b.ne 400aec <phase_5+0x2c> // b.any
 400ae4: a8c17bfd      ldp x29, x30, [sp], #16
 400ae8: d65f03c0      ret
 400aec: 94000002      bl 400af4 <explode>
 400af0: 17ffffff      b 400ae4 <phase_5+0x24>
```

要求 `func5` 的返回结果等于3

```
000000000400a4c <func_5>:
 400a4c: b4000361      cbz x1, 400ab8 <func_5+0x6c>
 400a50: a9be7bfd      stp x29, x30, [sp, #-32]!
 400a54: 910003fd      mov x29, sp
 400a58: a90153f3      stp x19, x20, [sp, #16]
 400a5c: 2a0003f4      mov w20, w0
 400a60: aa0103f3      mov x19, x1
 400a64: b9400020      ldr w0, [x1]
 400a68: 6b14001f      cmp w0, w20
 400a6c: 54000160      b.eq 400a98 <func_5+0x4c> // b.none
 400a70: b9400260      ldr w0, [x19]
 400a74: 6b14001f      cmp w0, w20
 400a78: 5400014d      b.le 400aa0 <func_5+0x54>
 400a7c: f9400661      ldr x1, [x19, #8]
 400a80: 2a1403e0      mov w0, w20
 400a84: 97ffffff      bl 400a4c <func_5>
 400a88: 531f7800      lsl w0, w0, #1
 400a8c: a94153f3      ldp x19, x20, [sp, #16]
 400a90: a8c27bfd      ldp x29, x30, [sp], #32
 400a94: d65f03c0      ret
 400a98: 94000017      bl 400af4 <explode>
 400a9c: 17ffffff      b 400a70 <func_5+0x24>
```

```

400aa0: f9400a61 ldr x1, [x19, #16]
400aa4: 2a1403e0 mov w0, w20
400aa8: 97ffffe9 bl 400a4c <func_5>
400aac: 531f7800 lsl w0, w0, #1
400ab0: 11000400 add w0, w0, #0x1
400ab4: 17ffffff b 400a8c <func_5+0x40>
400ab8: 52800000 mov w0, #0x0 // #0
400abc: d65f03c0 ret

```

func5 明显有三个条件分支且有递归调用。x1 传入的是一个指针，x0 传入的是我们输入的数。当 x0 与 x1 指向的数相等时会使炸弹爆炸；大于时，x1 的值变为 x1 + 8 对应的指针；小于时，x1 的值变为 x1 + 16 对应的指针；指针为空时，返回0。可以看作在一棵二叉树中搜索。写为 C++ 代码如下

```

int func5(int number, node node_ptr)
{
    if (bTree == NULL)
    {
        return 0;
    }
    else if (number > node_ptr->value)
    {
        node_ptr = node_ptr->left;
        return 2 * func5(number, node_ptr) + 1;
    }
    else if (number < node_ptr->value)
    {
        node_ptr = node_ptr->right;
        return 2 * func5(number, node_ptr);
    }
    else
    {
        explode();
    }
}

```

若要 func5 返回3，则要求在树中搜索的前两次必须为向右搜索，后面不管多少次都必须为向左搜索，即我们输入的数要大于根节点和根节点的右子节点，且小于根节点右子节点的右子节点的右子节点及其之后的所有左子节点直到指针为空。我们只需利用gdb找到根节点和其右子节点的值，并确认右子节点的右子节点是否为空即可。

```

pwndbg> x/3dg 0x4a0070
0x4a0070 <search_tree>: 49      4849800
0x4a0080 <search_tree+16>: 4849824
pwndbg> x/3dg 4849824
0x4a00a0 <search_tree+48>: 88      4849896
0x4a00b0 <search_tree+64>: 4849920
pwndbg> x/3dg 4849920
0x4a0100 <search_tree+144>: 91      0
0x4a0110 <search_tree+160>: 0

```

所以我们的输入 x 满足 $88 < x < 91$ 即可。