# Certified Denfense to Image Transformation via Randomized Smoothing

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#### Contribution

- First generalization of randomized smoothing to image transformations
- Presented several certified defenses allowing for both distributional and individual guarantees (relying on statistical error bounds or on efficient inverse computation)

#### Attack of Parametric Transformations

















# Randomized Smoothing (RS)

A smoothed classifier  $g : \mathbb{R}^m \mapsto \mathcal{Y}$  can be constructed out of an ordinary classifier  $f : \mathbb{R}^m \mapsto \mathcal{Y}$ , by calculating the most probable result of  $f(\boldsymbol{x} + \epsilon)$  where  $\epsilon \sim \mathcal{N}(0, \sigma^2 \mathbb{1})$ :

$$g(\boldsymbol{x}) := \underset{c}{\arg\max} \mathbb{P}_{\epsilon \sim \mathcal{N}(0, \sigma^2 \mathbb{1})}(f(\boldsymbol{x} + \epsilon) = c).$$

One then obtains the following robustness guarantee:

**Theorem 3.1** (From  $\boxed{7}$ ). Suppose  $c_A \in \mathcal{Y}$ ,  $\underline{p_A}$ ,  $\overline{p_B} \in [0, 1]$ . If

$$\mathbb{P}_{\epsilon}(f(\boldsymbol{x}+\epsilon)=c_A) \geq \underline{p_A} \geq \overline{p_B} \geq \max_{c \neq c_A} \mathbb{P}_{\epsilon}(f(\boldsymbol{x}+\epsilon)=c),$$

then 
$$g(x + \delta) = c_A$$
 for all  $\delta$  satisfying  $\|\delta\|_2 \leq \frac{\sigma}{2}(\Phi^{-1}(\underline{p_A}) - \Phi^{-1}(\overline{p_B})) =: r_\delta$ .

### randomized smoothing for parametric transformations (SPT)

to randomized smoothing for parametric transformations (SPT):

$$g(x) = \operatorname{argmax}_{c \in \mathcal{Y}} \mathbb{P}(f(\psi_{\beta}(x)) = c)$$
  
for classifier  $f$ , noise  $\beta \sim \mathcal{N}(0, \sigma^2 I)$   
Then  $g(\psi_{\gamma}(x)) = g(x)$  for  $\|\gamma\|_2 \leq r_{\gamma}$ .

requires 
$$\psi_{\alpha} \circ \psi_{\beta} = \psi_{\alpha+\beta}$$

then  $g \circ \psi_{\gamma}(x) = c_A$  for all  $\gamma$  satisfying  $\|\gamma\|_2 \leq \frac{\sigma}{2}(\Phi^{-1}(\underline{p_A}) - \Phi^{-1}(\overline{p_B})) =: r_{\gamma}$ . Further, if g is evaluated on a proxy classifier f' that behaves like f with probability  $1 - \rho$  and else returns an arbitrary answer, then  $r_{\gamma} := \frac{\sigma}{2}(\Phi^{-1}(\underline{p_A} - \rho) - \Phi^{-1}(\overline{p_B} + \rho))$ .



## Pratical algorithm

```
# evaluate g at x

function PREDICT(f, \sigma, x, n, \alpha)

counts \leftarrow SAMPLEUNDERNOISE(f, x, n, \sigma)
\hat{c}_A, \hat{c}_B \leftarrow top two indices in counts

n_A, n_B \leftarrow counts[\hat{c}_A], counts[\hat{c}_B]

if BINOMPVALUE(n_A, n_A + n_B, 0.5) \leq \alpha return \hat{c}_A
else return ABSTAIN
```

- SAMPLEUNDERNOISE(f, x, num, σ) works as follows:
  - 1. Draw num samples of noise,  $\varepsilon_1 \dots \varepsilon_{\text{num}} \sim \mathcal{N}(0, \sigma^2 I)$ .
  - 2. Run the noisy images through the base classifier f to obtain the predictions  $f(x+\varepsilon_1),\ldots,f(x+\varepsilon_{\text{num}})$ .
  - 3. Return the counts for each class, where the count for class c is defined as  $\sum_{i=1}^{\text{num}} \mathbf{1}[f(x+\varepsilon_i)=c]$ .
- BINOMPVALUE $(n_A, n_A + n_B, p)$  returns the p-value of the two-sided hypothesis test that  $n_A \sim \text{Binomial}(n_A + n_B, p)$ . Using <u>scipt.stats.binom.test</u>, this can be implemented as: binom.test(nA, nA + nB, p).

## Pratical algorithm

```
# certify the robustness of g around x function Certify(f, \sigma, x, n_0, n, \alpha) counts 0 \leftarrow \text{SAMPLEUNDERNOISE}(f, x, n_0, \sigma) \hat{c}_A \leftarrow \text{top index in counts} 0 counts \leftarrow \text{SAMPLEUNDERNOISE}(f, x, n, \sigma) \underline{p}_A \leftarrow \text{LOWERCONFBOUND}(\text{counts}[\hat{c}_A], n, 1 - \alpha) if \underline{p}_A > \frac{1}{2} return prediction \hat{c}_A and radius \sigma \Phi^{-1}(\underline{p}_A) else return ABSTAIN
```

```
LOWERCONFBOUND(k, n, 1-\alpha) returns a one-sided (1-\alpha) lower confidence interval for the Binomial parameter p given that k \sim \text{Binomial}(n,p). In other words, it returns some number \underline{p} for which \underline{p} \leq p with probability at least 1-\alpha over the sampling of k \sim \text{Binomial}(n,p). Following Lecuyer et al. (2019), we chose to use the Clopper-Pearson confidence interval, which inverts the Binomial CDF (Clopper & Pearson, 1934). Using \underline{\text{statsmodels}}, \underline{\text{stats.proportion.proportion.confint}}, this can be implemented as
```

```
proportion_confint(k, n, alpha=2*alpha, method="beta")[0]
```

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# BaseSPT(Heuristic best effort defense)

Table 2: Evaluation of BASESPT. We obtain Acc for b on the test set and evaluate adv. Acc. on 3000 images obtained by the worst-of-100 attack. t denotes the average run time of g.

			Acc.	adv.		
Dataset	$T^{I}$	$\Gamma_{\pm}$	b	b	g	t [s]
MNIST	$R^{I}$	30°	0.99	0.73	0.99	0.97
CIFAR-10	$R^{I}$	$30^{\circ}$	0.91	0.26	0.85	0.95
ImageNet	$R^{I}$	$30^{\circ}$	0.76	0.56	0.76	5.43
MNIST	$\Delta^{I}$	4	0.99	0.03	0.53	0.86
CIFAR-10	$\Delta^{I}$	4	0.91	0.44	0.79	0.95
ImageNet	$\Delta^{I}$	20	0.76	0.65	0.75	6.70

- By applying SPT to image rotation we can obtain a heuristic defense as rotations don't compose as required (discussed next).
- Here we show results for adversarial rotations of up to 30° and translations up 4 or 20.
- worst-of-k attack:randomly sample k parameters;choose the three worst(max loss).

• 
$$n_{\gamma} = 1000, \gamma = \Gamma_{\pm}, \alpha_{\gamma} = 0.01$$

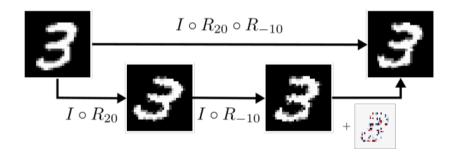


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### Interpolation:Image Rotations don't compose



### Certification in the Presence of Interpolation

- $\text{Aim:} g_E \circ T_{\gamma}^I(x) = g_E(x), \|\gamma\| \le r_{\gamma}$
- $h_E \circ T_{\beta}^I \circ T_{\gamma}^I(\mathbf{x}) = h_E \circ T_{\beta + \gamma}^I(\mathbf{x}), \ \forall \beta, \gamma \in \mathbb{R}^d$

$$g_E \circ T_{\gamma}(\boldsymbol{x}) = \arg\max_{c} \mathbb{P}_{\beta \sim \mathcal{N}(0, \sigma^2 \mathbb{1})} \left( h_E \circ I \circ T_{\beta} \circ T_{\gamma}(\boldsymbol{x}) = c \right)$$

$$f := h_E \circ I, \psi_\beta := T_\beta$$

$$= \arg\max_{c} \mathbb{P}_{\beta} \sim \mathcal{N}(0, \sigma^{2} \mathbb{1}) \left( h_{E} \circ T_{\beta}^{I} \circ T_{\gamma}^{I}(\boldsymbol{x}) = c \right)$$

$$\mathbf{0} g_E(\mathbf{x}) := \arg \max_c \mathbb{P}_{\beta \sim \mathcal{N}(0, \sigma^2 1)} \left( h_E \circ I \circ T_\beta(\mathbf{x}) = c \right)$$

$$=g_{E}\circ T_{\gamma}^{I}(\boldsymbol{x}),$$

Such that 
$$g_E \circ T_{\gamma}(x) = c_A = g_E(x), ||\gamma|| \le r_{\gamma}$$

# Second Step: Construction $h_E$

$$\begin{split} \epsilon(\beta,\gamma,\boldsymbol{x}) &:= T^I_{\beta} \circ T^I_{\gamma}(\boldsymbol{x}) - T^I_{\beta+\gamma}(\boldsymbol{x}), \\ \text{bounded by } E \in \mathbb{R}^{\geq 0} \text{ s.t. } \forall \beta,\gamma \in \mathbb{R}^d. \ \|\epsilon(\beta,\gamma,\boldsymbol{x})\|_2 \leq E \end{split}$$

Thus if  $h_E$  is  $l_2$ -robust with radius E around  $T^I_{\beta+\gamma}(x)$ , interpolation invariance holds.



## Obtaining probabilistic guarantees from Theorem 3.2

$$\mathbb{P}_{\beta \sim \mathcal{N}(0, \sigma^2 1)} \left( \| \epsilon(\beta, \gamma, \mathbf{x}) \|_2 \le E \right) \ge 1 - \rho_E \quad \forall \gamma \in \mathbb{R}^d$$

- $h'_E$ : like  $h_E$  with probability at least  $1 \rho_E$
- $f = h'_E \circ I$  like f with probability at least  $1 \rho_E$

For a fixed  $E \in \mathbb{R}^{\geq 0}$ ,  $\rho_E \in [0,1]$ , the probability that  $\epsilon$  is bounded by E for  $x \sim \mathcal{D}$  is

$$q_E := \mathbb{P}_{\boldsymbol{x} \sim \mathcal{D}}(\mathbb{P}_{\beta \sim \mathcal{N}(0, \sigma^2 \mathbb{1})}(\max_{\gamma \in \Gamma} \|\epsilon(\beta, \gamma, \boldsymbol{x})\|_2 \le E) \ge 1 - \rho_E). \tag{6}$$

# Evaluate $q_E:DistSPT^{D}$

- **1** Sampling multiple realizations of  $\beta$
- ② Computing their corresponding error  $\epsilon$  and checking how many are successfully bounded by E
- **3** Bounding the inner probability using Clopper-Pearson.
- f 0 Counts x and apply Clopper-Pearson and obtain a lower bound  $q_E$  with the desired confidence
- **5** For  $DistSPT^{x}$ , Only compute inner probability  $(q_{E} = 1)$

#### maximization interpolation error over $\gamma$

- method: interval analysis; propagate lower and upper bounds
- ② obtain a lower and upper bound for the norm Calculation, pick maximum:

$$\max_{\gamma \in \Gamma} \|\epsilon(\beta, \gamma, x)\|_2 \le \max \|T_{\beta}^I \circ T_{\Gamma}^I(x) - T_{\beta + \Gamma}^I(x)\|_2$$

**3** Refine: splitting the hyperrectangle  $\Gamma$  into smaller hyperrectangles  $\Gamma_k$ 

$$\max_{\gamma \in \Gamma} \|\epsilon(\beta, \gamma, \mathbf{x})\|_2 \leq \max_{k \in \{0, \dots, N\}} \max \left\| T_{\beta}^I \circ T_{\Gamma_k}^I(\mathbf{x}) - T_{\beta + \Gamma_k}^I(\mathbf{x}) \right\|_2$$

- E:For each sample  $(x, \beta)$ , we simply keep the values;
- $DistSPT^D$ : pick an E that bounds many of these values, choosing  $\rho_E$  to be small, and Theorem hold with  $q_E$ .  $DistSPT^x$ : suitable E

#### $DistSPT^{D}$

Table 3: Evaluation of DISTSPT<sup>D</sup> for  $T^I := R^I$ . We show the test set accuracy of b, certified accuracy of g at different radii  $r_{\gamma_i}$  along with the average run time t. # denotes values obtained by sampling. Each certificate hold with overall confidence 0.99.

				$g$ cert. acc at $r_{\gamma}$					
Dataset	E	$q_E$	b acc.	0°	$10^{\circ}$	$20^{\circ}$	30°	t [s]	$n_{\gamma}$
MNIST	0.45	0.99	0.98	0.89	0.88	0.87	0.85	21.56	200
CIFAR-10	0.55	0.99	0.56	0.31	0.28	0.25	0.19	89.75	50
CIFAR-10	0.55	0.99	0.56	0.32	0.30	0.28	0.25	351.47	200
RImageNet	$1.20^{\#}$	0.97	0.78	0.74	0.72	0.68	0.61	100.73	50
RImageNet	$1.35^{\#}$	0.99	0.78	0.64	0.62	0.56	0.50	100.13	50
ImageNet	$0.95^{\#}$	0.75	0.38	0.30	0.24	0.18	0.12	100.21	50
ImageNet	$1.20^{\#}$	0.97	0.38	0.23	0.19	0.13	0.09	100.73	50
ImageNet	1.35#	0.99	0.38	0.16	0.12	0.08	0.06	100.44	50

- obtain E:Sample the interpolation error using 1000 images
- test for  $\rho_E = 10^{-3}$ , obtain  $q_E$  with  $1 \alpha_E = 0.999$ ,  $(x = 1000, \beta = 8000)$
- optimization over  $\gamma$  for many images is computationally expensive  $\to$  10 sample for  $\gamma$

#### $DistSPT^{x}$

Table 4: Evaluation of DISTSPT\* for  $T^I := R^I$ . We show the test set certified accuracy of g at different radii  $r_{r_0}$ , along with the average E estimated, the average time  $t_E$  to estimate E and average time  $t_{RE}$  to apply randomized smoothing. # denotes values obtained by sampling. \* we use a server with 128 threads on an AMD EPYC 7601 processor, on the same system as the other results these take 766 s. Each certificate hold with overall confidence 0.99.

				$g$ cert. acc at $r_{\gamma}$							
Dataset	$\Gamma_{\pm}$	$\sigma_{\gamma}$	0°	$10^{\circ}$	$20^{\circ}$	30°	50°	avg. $E$	$t_E$ [s]	$t_{RS}$ [s]	$n_{\gamma}$
MNIST	50°	30	0.93	0.92	0.91	0.90	0.82	0.34	53.33	20.56	200
CIFAR-10	$30^{\circ}$	40	0.35	0.30	0.27	0.22	-	0.34	81.83	91.72	50
CIFAR-10	$10^{\circ}$	10	0.43	0.37	-	-	-	0.34	51.12	92.83	50
ImageNet	$30^{\circ}$	30	0.31	0.25	0.17	0.11	-	$0.86^{\#}$	73.58*	100.47	50
ImageNet	$30^{\circ}$	30	0.32	0.29	0.22	0.16	-	$0.86^{\#}$	73.58*	396.50	200

- individual bounds are much lower than distribution bounds for  $E \to \text{better}$  accuracy
- 100 sample of  $\beta$  to guess E; 400 sample of  $\beta$  to test  $\rho_E$
- In theory,  $DistSPT^D$  can perform better if  $\rho_E$  is lower (e.g. when more samples are used)
- However, in practice this is offset by the tighter error bound attained by DistSPT<sup>x</sup>



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#### Flow

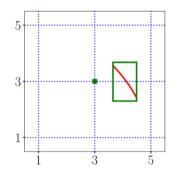
- we are given  $x' := T_{\gamma}^{I}(x)$  but neither the original x nor the parameter  $\gamma \in \Gamma$ , and we would like to certify that  $g_{E}(x') = g_{E}(x)$
- $\mathbb{P}_{\beta \sim \mathcal{N}(0, \sigma^2 1)} (\|\epsilon(\beta, \gamma, \mathbf{x})\|_2 \leq E) \geq 1 \rho_E \quad \forall \gamma \in \mathbb{R}^d$
- compute an upper bound on the max term without having access to x
- obtain  $E: \max_{\gamma \in \Gamma} \|\epsilon(\beta, \gamma, \mathbf{x})\|_{2} \leq \max \|T_{\beta}^{I}(\mathbf{x}') T_{\beta+\gamma}^{I} \circ (T_{\gamma}^{I})^{-1}(\mathbf{x}')\|_{2}$
- KEY:  $\left(T_{\gamma}^{I}\right)^{-1} \rightarrow \left(T_{\Gamma}^{I}\right)^{-1}(\mathbf{x}') := \left\{\mathbf{x} \in \mathbb{R}^{m} \mid T_{\gamma}^{I}(\mathbf{x}) = \mathbf{x}', \gamma \in \Gamma\right\}$
- $\max_{\gamma \in \Gamma} \|\epsilon(\beta, \gamma, \mathbf{x})\|_2 \le \max \|T_{\beta}^I(\mathbf{x}') T_{\beta+\Gamma}^I \circ (T_{\Gamma}^I)^{-1}(\mathbf{x}')\|_2$



#### base

- even image height and width; center it at 0 on  $G := (2\mathbb{Z} + 1) \times (2\mathbb{Z} + 1)$
- pixel at (i,j):  $p_{(i,j)} \in [0,1]$
- (v,w)-interpolation region:  $[v, v+2] \times [w, w+2]$
- Bilinear interpolation:  $I(x, y) = p_{v,w} \frac{2+v-x}{2} \frac{2+w-y}{2} + p_{v,w+2} \frac{2+v-x}{2} \frac{y-w}{2} + p_{v+2,w} \frac{x-v}{2} \frac{2+w-y}{2} + p_{v+2,w+2} \frac{x-v}{2} \frac{y-w}{2}$
- Example: assume attack  $\gamma \in [23^{\circ}, 26^{\circ}]$ , calculate the constraint for pixel (3,3) of the original image.

For every  $(i',j') \in G$ , we over-approximate the region the pixel value  $p'_{i',j'}$  could have been interpolated from, which is  $c_{i',j'} := T_{\Gamma}^{-1}(i',j')$ ,  $C := \{c_{i',j'} \mid (i',j') \in G\}$ . We illustrate the calculation of the set C for  $c_{5,1} := R_{[23^{\circ},26^{\circ}]}^{-1} \binom{5}{1} = \binom{[4.06,4.21]}{[2.85,3.11]}$ .



the pixel value  $p'_{i',i'}$  yields constraints for value  $p_{i,j}$ 

$$[x_l, x_u] \times [y_l, y_u] := c_{i',j'} \cap [i, i+2] \times [j, j+2]$$

plug this into the interpolation I

$$p'_{i',j'} \in I([x_l, x_u], [y_l, y_u]) = p_{i,j} \frac{2 + i - [x_l, x_u]}{2} \frac{2 + j - [y_l, y_u]}{2} + p_{i,j+2} \frac{2 + i - [x_l, x_u]}{2} \frac{[y_l, y_u] - j}{2} + p_{i+2,j} \frac{[x_l, x_u] - i}{2} \frac{2 + j - [y_l, y_u]}{2} + p_{i+2,j+2} \frac{[x_l, x_u] - i}{2} \frac{[y_l, y_u] - j}{2}$$

- Solve for the pixel value of interest  $p_{i,j}$ :replace other three with the (trivial) [0,1] constraint, covering all possible pixel values.
- instantiating  $[x_l, x_u]$  and  $[y_l, y_u]$  with its corner (x, y) furthest from (i, j), yields still a sound but more precise constraint  $q_{i,j}$  for  $p_{i,j}$ .
- choose  $x_u, y_u$ , then  $p_{i,j}$  is:

$$q_{i,j} = \left[p'_{i',j'} - \left(\frac{2+i-x_u}{2} \frac{y_u - j}{2} + \frac{x_u - i}{2} \frac{2+j-y_u}{2} + \frac{x_u - i}{2} \frac{y_u - j}{2}\right), p'_{i',j'}\right] \left(\frac{2+i-x_u}{2} \frac{2+j-y_u}{2}\right)^{-1}$$

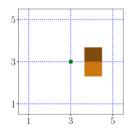


**Step 2** The only non-empty intersections of  $c_{5,1}$  with interpolation regions (blue squares in Fig. [3]), cornering (3, 3) are the (3, 1) and the (3, 3)-interpolation regions, hence we omit  $q_{1,1}$  and  $q_{1,3}$ . The intersection with the (3, 3)-interpolation region yields  $[x_l, x_u] = [4.06, 4.21]$  and  $[y_l, y_u] = [3, 3.11]$  (dark brown rectangle in Fig. [3b), hence at the furthest corner (x, y) = (4.21, 3.11), we get

$$q_{3,3} = [0.73, 2.48] = \left[p'_{5,1} - \left(\frac{5-x}{2}, \frac{y-3}{2} + \frac{x-3}{2}, \frac{5-y}{2} + \frac{x-3}{2}, \frac{y-3}{2}\right), p'_{5,1}\right] \left(\frac{5-x}{2}, \frac{5-y}{2}\right)^{-1},$$

and the intersection with the (3,1)-interpolation region yields  $[x_l,x_u]=[4.06,4.21]$  and  $[y_l,y_u]=[2.85,3]$  (light brown rectangle in Fig. 3b), hence at the furthest corner (x,y)=(4.21,2.85), we get

$$q_{3,1} = \left[0.72, 2.48\right] = \left[p_{5,1}' - \left(\frac{5-x}{2}\frac{3-y}{2} + \frac{x-3}{2}\frac{3-y}{2} + \frac{x-3}{2}\frac{y-1}{2}\right), p_{5,1}'\right] \left(\frac{5-x}{2}\frac{y-1}{2}\right)^{-1}.$$



- In order to be sound, we need to take the union  $q_{i,j}$
- $\bullet$  To gain precision, we can intersect all of those unions and [0,1]
- $[a, b] \sqcup [c, d] := [\min(a, c), \max(b, d)]$

•

$$p_{i,j} \in [0,1] \cap \left( \bigcap_{c_{i',j'} \in C} q_{i-2,j-2} \left( c_{i',j'} \right) \sqcup q_{i,j-2} \left( c_{i',j'} \right) \sqcup q_{i-2,j} \left( c_{i',j'} \right) \sqcup q_{i,j} \left( c_{i',j'} \right) \right)$$

• In example:  $q_{3,1} \sqcup q_{3,1} = [0.72, 2.48] \to \cap [0,1] \to \text{constraints from the other } c'_{i',j'} \to p_{3,3} \in [0.73,1]$ 

## Experiment

- Refine: 10 times
- E:large due to loss of precision
- $\bullet$  g was correct on 82% of attacked images;81% we could certify on MNIST
- ullet analysis of E took on average 0.26s and the randomized smoothing 25.03s
- Challenges on larger datasets

# further Experiment

Table 11: Maximum observed errors and without gaussian blur (G) and without vignetting (V).

Dataset	Both	-V	-G	-V-G
MNIST	0.36	0.36	2.47	2.51
CIFAR-10	0.51	6.08	2.66	18.17
ImageNet	0.91	70.66	9.25	75.69

Table 12: Correct classifications and by the model and verifications by DeepG [11], with and without vignetting (V), out of 100 images.

Model	Correct	[11]	[11]+V
MNIST	98	86	87
CIFAR-10	74	65	32
CIFAR-10+V	78	63	23

we apply a circular or rectangular vignette for rotation and translation respectively, to reduce error estimates in areas of the image where information is lost.

Distributional bounds for IndivSPT Inverse Computation further Experiment

Thank you