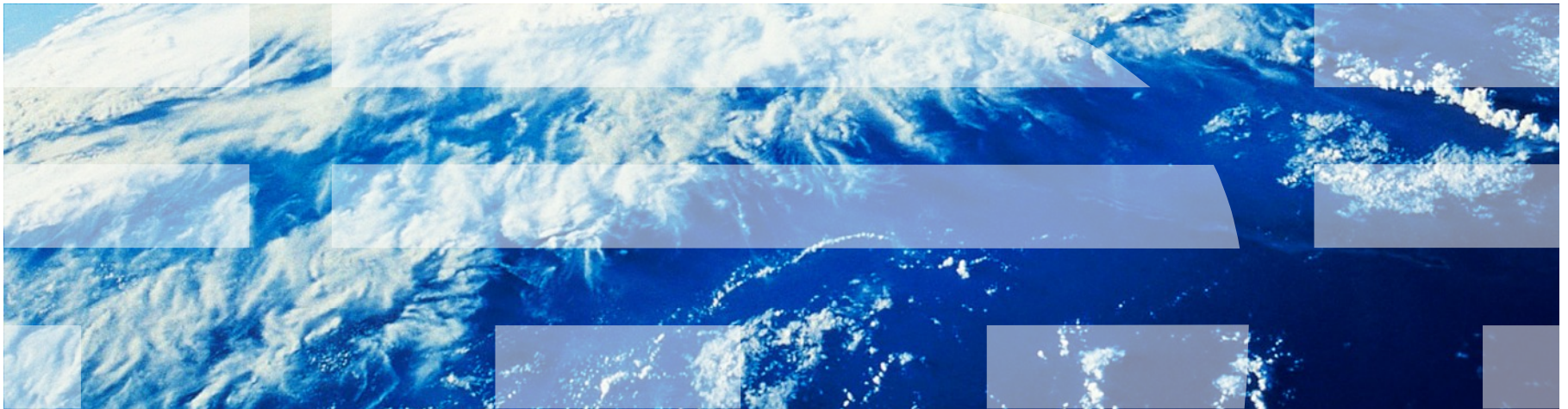


# Computer Systems for Data Science

## Topic 2

### Relational Model and SQL



---

## What we'll cover in this topic

- **Intuition**

- Basic relational model
- Map-filter-reduce concept

- **Intro to SQL**

- Schemas, query structure of SELECT-FROM-WHERE, JOINS

---

# Relational Model: Intuition



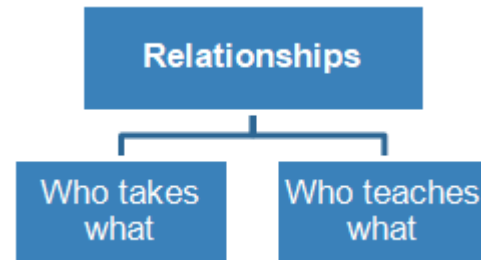
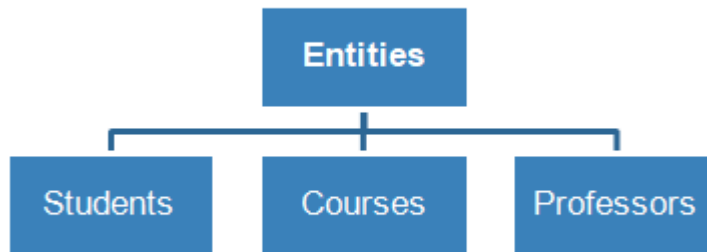
---

## A Motivating Example

### A basic Course Management System (CMS):

*Entities* (e.g., Students, Courses)

*Relationships* (e.g., Alice is enrolled in CSEE 4121)



## Intuition: Spreadsheet Tables

### Logical Schema

Student(cuid: string, name: string,  
gpa: float)

Courses(cid: string, c-name: string,  
room: string)

Enrolled(cuid: string, cid: string, grade:  
string)

Students			Enrolled		
CUID	Name	GPA	CUID	CID	Grade
as2121	Alice Smith	4.3	as2121	4121	A+
jg9999	Jay Goodwin	1.2	jg9999	4121	C
mc2312	Min Chang	2.2	mc2312	3292	A+
zb1111	Zorn Bjorn	3.8	zb1111	2999	D
Courses					
CID	C-Name	Room			
4121	Computer Sy	CEPSR			
3292	Databases	MUDD 1			
2999	Algorithms	MUDD 2			

### Queries [“compute” over tables]

Alice’s GPA?

Jay’s classes?

AVG student GPA?

AVG student GPA in CSEE 4121?

---

# Relational Model and Schemas

- Definition: Data model
  - Organizing principle of data + operations
- Relational model (aka tables)
  - Simple and most popular
  - Elegant algebra (E.F. Codd et al)
- Definition: Schema
  - Describes blueprint of table (s)
- Every relation has a schema
  - Logical Schema: describes types, names
  - Physical Schema: describes data layout
  - Virtual Schema (Views): derived tables

---

## Data Independence

1. Can we add a new column or attribute without rewriting the application?
  - **Logical Data Independence**
    - Protection from changes in the logical structure of the data
2. Do you need to care which disks/machines are the data stored on?
  - **Physical Data Independence**
    - Protection from Physical Layout Changes

---

# Python Operating on Lists

## Basic types

- Int
- Long
- String

## Map + Filter

- `map(function, list)`
- `filter(function, list)`

**Map** applies function to input list

**Filter** returns sub list that satisfies filter condition

## Reduce/Aggregate

- `reduce(...)`

**Reduce** runs a computation on a list and returns a result  
E.g., SUM, AVG, MAX



## SQL Queries on Tables (Lists of Rows)

### Basic types

- Int32, Int64
- Char[n]
- Float32, Float64

### Map + Filter

Single Table Query

```
SELECT c1, c2
FROM T
WHERE condition;
```

Multi Table JOIN

```
SELECT c1, c2
FROM T1, T2
WHERE condition;
```

### Reduce/Aggregate

```
SELECT SUM(c1*c2)
FROM T
WHERE condition
GROUP BY c3;
```

Map-Filter-Reduce pattern: Same simple/powerful idea in MapReduce, Hadoop, Spark, etc.

# SQL Cheat Sheet ([www.sqltutorial.org/sql-cheat-sheet](http://www.sqltutorial.org/sql-cheat-sheet))

## SQL CHEAT SHEET <http://www.sqltutorial.org>



### QUERYING DATA FROM A TABLE

**SELECT c1, c2 FROM t;**  
Query data in columns c1, c2 from a table

**SELECT \* FROM t;**  
Query all rows and columns from a table

**SELECT c1, c2 FROM t**  
**WHERE condition;**  
Query data and filter rows with a condition

**SELECT DISTINCT c1 FROM t**  
**WHERE condition;**  
Query distinct rows from a table

**SELECT c1, c2 FROM t**  
**ORDER BY c1 ASC [DESC];**  
Sort the result set in ascending or descending order

**SELECT c1, c2 FROM t**  
**ORDER BY c1**  
**LIMIT n OFFSET offset;**  
Skip *offset* of rows and return the next *n* rows

**SELECT c1, aggregate(c2)**  
**FROM t**  
**GROUP BY c1;**  
Group rows using an aggregate function

**SELECT c1, aggregate(c2)**  
**FROM t**  
**GROUP BY c1**  
**HAVING condition;**  
Filter groups using HAVING clause

### QUERYING FROM MULTIPLE TABLES

**SELECT c1, c2**  
**FROM t1**  
**INNER JOIN t2 ON condition;**  
Inner join t1 and t2

**SELECT c1, c2**  
**FROM t1**  
**LEFT JOIN t2 ON condition;**  
Left join t1 and t2

**SELECT c1, c2**  
**FROM t1**  
**RIGHT JOIN t2 ON condition;**  
Right join t1 and t2

**SELECT c1, c2**  
**FROM t1**  
**FULL OUTER JOIN t2 ON condition;**  
Perform full outer join

**SELECT c1, c2**  
**FROM t1**  
**CROSS JOIN t2;**  
Produce a Cartesian product of rows in tables

**SELECT c1, c2**  
**FROM t1, t2;**  
Another way to perform cross join

**SELECT c1, c2**  
**FROM t1 A**  
**INNER JOIN t2 B ON condition;**  
Join t1 to itself using INNER JOIN clause

### USING SQL OPERATORS

**SELECT c1, c2 FROM t1**  
**UNION [ALL]**  
**SELECT c1, c2 FROM t2;**  
Combine rows from two queries

**SELECT c1, c2 FROM t1**  
**INTERSECT**  
**SELECT c1, c2 FROM t2;**  
Return the intersection of two queries

**SELECT c1, c2 FROM t1**  
**MINUS**  
**SELECT c1, c2 FROM t2;**  
Subtract a result set from another result set

**SELECT c1, c2 FROM t1**  
**WHERE c1 [NOT] LIKE pattern;**  
Query rows using pattern matching %, \_

**SELECT c1, c2 FROM t**  
**WHERE c1 [NOT] IN value\_list;**  
Query rows in a list

**SELECT c1, c2 FROM t**  
**WHERE c1 BETWEEN low AND high;**  
Query rows between two values

**SELECT c1, c2 FROM t**  
**WHERE c1 IS [NOT] NULL;**  
Check if values in a table is NULL or not

---

# Intro to SQL



---

## SQL Introduction

- ✦ SQL is a standard language for querying and manipulating data
- SQL is a **very high-level** programming language
  - This works because it is optimized well!

**SQL** stands for  
**Structured  
Query  
Language**

---

## SQL is a...

- Data Manipulation Language (DML)

- Query one or more tables

- Insert/delete/modify tuples in tables

- Data Definition Language (DDL)

- Define relational schemata

- Create/alter/delete tables and their attributes

---

## Basic Set Algebra Concepts

- List:  $[1, 1, 2, 3]$       Ordered, duplicates
- Set:  $\{2, 1, 3\}$       Unordered, no duplicates
- Multiset:  $\{2, 1, 3, 1\}$       Unordered, duplicates
  
- Unions:
  - Set:  $\{2, 1, 3\} \cup \{2, 3\} = \{2, 1, 3\}$
  - Multiset:  $\{2, 1, 3\} \cup \{2, 3\} = \{2, 1, 3, 2, 3\}$
  
- Cross-product:
  - $\{1, 1, 2, 3\} * \{y, z\} = \{1, y\}, \{1, y\}, \{2, y\}, \{3, y\}, \{1, z\}, \{1, z\}, \{2, z\}, \{3, z\}$

---

# Tables in SQL

## Product

PName	Price	Manuf
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

A relation or table is a multiset of tuples/rows having the attributes specified by the schema

# Tables in SQL

## Product

PName	Price	Manuf
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

An **attribute** (or **column**) is a typed data entry present in each tuple in the relation

*NB: Attributes must have an **atomic** type in standard SQL, i.e. not a list, set, etc.*



---

# Tables in SQL

## Product

PName	Price	Manuf
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

A **tuple** or **row** or **record** is a single entry in the table having the attributes specified by the schema

# Tables in SQL

## Product

PName	Price	Manuf
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

The number of tuples is the **cardinality** of the relation

The number of attributes is the **arity** of the relation

---

## Data Types in SQL

### Atomic types:

Characters: CHAR(20), VARCHAR(50)

Numbers: INT, BIGINT, SMALLINT, FLOAT

Others: MONEY, DATETIME...

Every attribute must have an atomic type

---

# Table Schemas

The **schema** of a table is the table name, its attributes, and their types:

Product(Pname: *string*, Price: *float*, Category: *string*, Manufacturer: *string*)

A **key** is an attribute whose values are unique; we underline a key

Product(Pname: *string*, Price: *float*, Category: *string*, Manufacturer: *string*)

---

# Key constraints

A **key** is a **minimal subset of attributes** that acts as a unique identifier for tuples in a relation

- A key is an implicit constraint on which tuples can be in the relation
- i.e. if two tuples agree on the values of the key, then they must be the same tuple!

`Students(cuid:string, name:string, gpa: float)`

1. Which would you select as a key?
2. Is a key always guaranteed to exist?
3. Can we have more than one key?

---

# Declaring Schema

Students(cuid: *string*, name: *string*, gpa: *float*)

```
CREATE TABLE Students (  
  cuid CHAR(20),  
  name VARCHAR(50),  
  gpa float,  
  PRIMARY KEY (cuid),  
)
```

---

# NULL and NOT NULL

- To say “don’t know the value” we use **NULL**

Students(cuid:string, name:string, gpa: float)

cuid	name	gpa
123	Alice	3.9
143	Jim	NULL

*Say, Jim just enrolled in his first class.*

In SQL, we may constrain a column to be NOT NULL, e.g., “name” in this table

---

# General Constraints

- We can actually specify arbitrary assertions  
E.g. *“There cannot be 25 people in the DB class”*
- In practice, we don’t specify many such constraints. Why?  
**Performance!**

Whenever we do something ugly (or avoid doing something convenient) it’s for the sake of performance



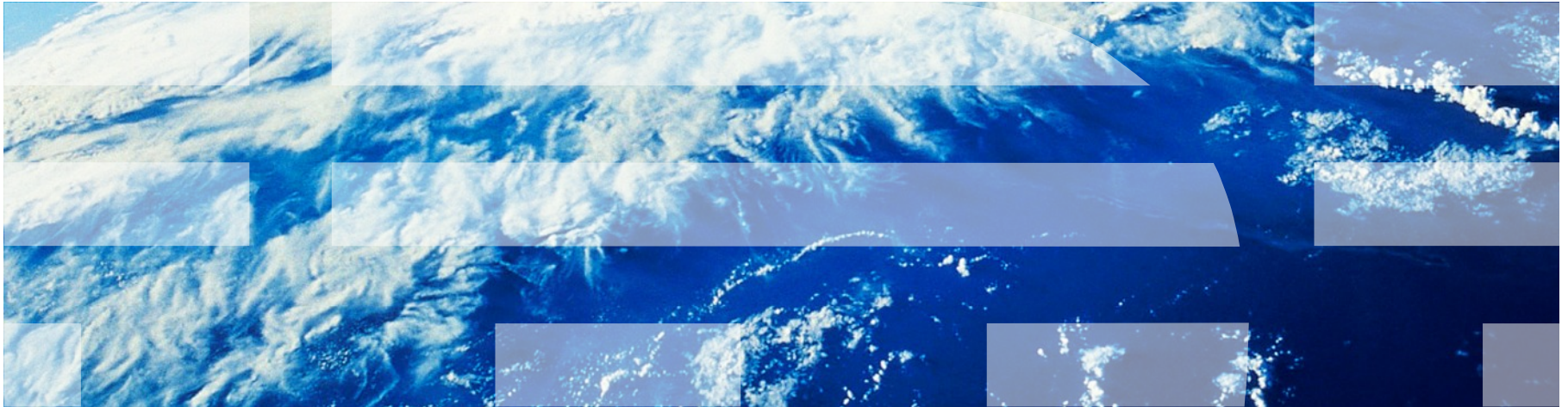
---

# Summary of Schema Information

- Schema and Constraints are how databases understand the semantics (meaning) of data
- SQL supports general constraints:  
**Keys and foreign keys are most important**

---

# Single Table Query



---

# SQL Query

Basic form (there are many many more bells and whistles)

**SELECT** <attributes>

**FROM** <one or more relations>

**WHERE** <conditions>

Call this a **SFW** query.

# Simple SQL Query: Selection

**Selection** is the operation of filtering a relation's tuples on some condition

```
SELECT *  
FROM Product  
WHERE Category = 'Gadgets'
```

PName	Price	Category	Manuf
Gizmo	\$19.99	Gadgets	GWorks
Powergizmo	\$29.99	Gadgets	GWorks

SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi



PName	Price	Category	Manuf
Gizmo	\$19.99	Gadgets	GWorks
Powergizmo	\$29.99	Gadgets	GWorks

# Simple SQL Query: Projection

**Projection** is the operation of producing an output table with tuples that have a subset of their prior attributes

```
SELECT Pname, Price, Manufacturer
FROM Product
WHERE Category = 'Gadgets'
```

PName	Price	Category	Manuf
Gizmo	\$19.99	Gadgets	GWorks
Powergizmo	\$29.99	Gadgets	GWorks

SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi



PName	Price	Manuf
Gizmo	\$19.99	GWorks
Powergizmo	\$29.99	GWorks

---

# Notation

Input Schema

Product(PName, Price, Category, Manufacturer)

```
SELECT Pname, Price, Manufacturer  
FROM   Product  
WHERE  Category = 'Gadgets'
```



Output Schema

Answer(PName, Price, Manufacturer)

---

## A Few Details

- **SQL commands** are case insensitive:

Same: SELECT, Select, select

Same: Product, product

- **Values are not:**

Different: 'Seattle', 'seattle'

- Use single quotes for constants:

'abc' - yes

"abc" - no

---

## LIKE: Simple String Pattern Matching

```
SELECT *  
FROM Products  
WHERE PName LIKE '%gizmo%'
```

- s **LIKE** p: pattern matching on strings
- p may contain two special symbols:
  - % = any sequence of characters
  - \_ = any single character



---

# DISTINCT: Eliminating Duplicates

```
SELECT DISTINCT Category  
FROM Product
```



Category

Gadgets

Photography

Household

**Versus**

```
SELECT Category  
FROM Product
```



Category

Gadgets

Gadgets

Photography

Household

---

## ORDER BY: Sorting the Results

SELECT	PName, Price, Manufacturer
FROM	Product
WHERE	Category='gizmo' AND Price > 50
ORDER BY	Price, PName

Ties are broken by the second attribute on the ORDER BY list, etc.

Ordering is ascending, unless you specify the DESC keyword.

---

# Multi-Table Query



# Foreign Key constraints

- Suppose we have the following schema :

```
Students(cuid: string, name: string, gpa: float)
Enrolled(student_id: string, cid: string, grade: string)
```

- And we want to impose the following constraint:  
Only bona fide students may enroll in courses i.e. a student must appear in the Students table to enroll in a class

Students		
cuid	name	gpa
102	Bob	3.9
123	Mary	3.8

Enrolled		
student_id	cid	grade
102	564	A
123	537	A+

We say that cuid is a foreign key that refers to Students

---

# Declaring Foreign Keys

Students(cuid: *string*, name: *string*, gpa: *float*)  
Enrolled(student\_id: *string*, cid: *string*, grade: *string*)

```
CREATE TABLE Enrolled (  
  student_id CHAR(20),  
  cid CHAR(20),  
  grade CHAR(10),  
  PRIMARY KEY (student_id, cid),  
  FOREIGN KEY (student_id) REFERENCES Students(cuid)  
)
```

---

## Foreign Keys and update operations

Students(cuid: string, name: string, gpa: float)

Enrolled(student\_id: string, cid: string, grade: string)

- What if we insert a tuple into Enrolled, but no corresponding student?

INSERT is rejected (foreign keys are constraints)!

- What if we delete a student?

1. Disallow the delete
2. Remove all of the courses for that student
3. *SQL allows a third via NULL*

*DBA chooses*

---

# Keys and Foreign Keys

## Company

<b><u>CName</u></b>	<b>StockPrice</b>	<b>Country</b>
GizmoWorks	25	USA
Canon	65	Japan
Hitachi	15	Japan

What is a foreign key vs. a key here?

## Product

<b><u>PName</u></b>	<b>Price</b>	<b>Category</b>	<b>Manufacturer</b>
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

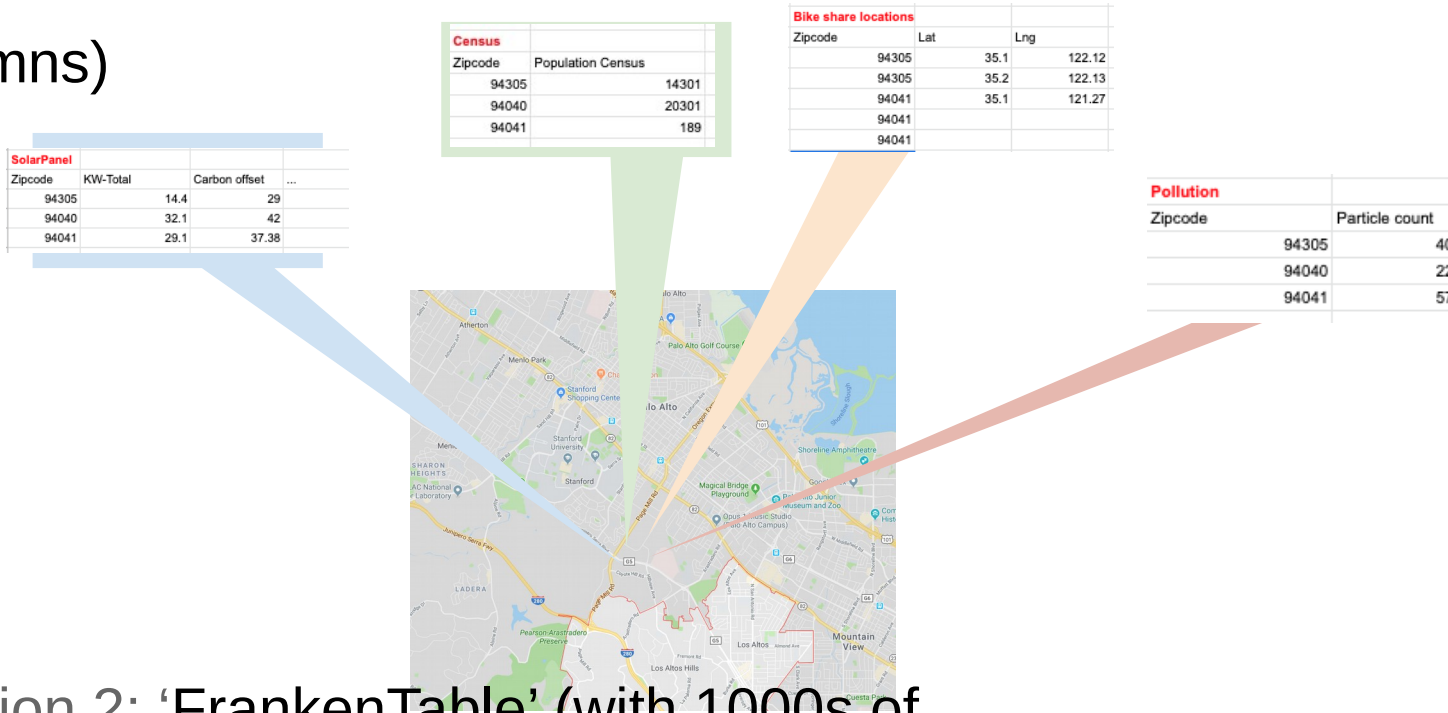
---

# JOINS and Aggregations





Option 1: Organized tables, with 10s-100s of columns)



Option 2: 'FrankenTable' (with 1000s of columns)

Omnidata	SolarPanel			Census	Pollution	Bike share locations	
Zipcode	KW-Total	Carbon offset	...	Population	Particle count	Lat	Lng
94305	14.4	29		14301	40	35.1	122.12
94305						35.2	122.13
94305						35.2	122.1
94305						35.1	122.12
94305						...	...

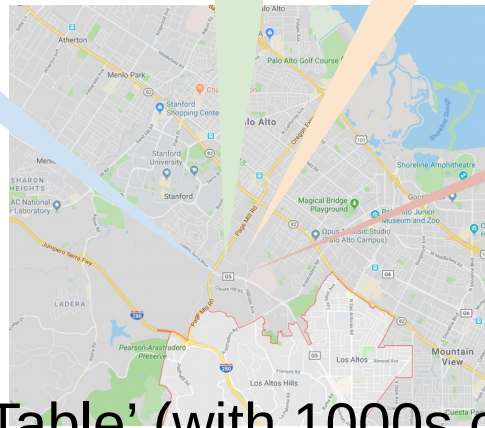
## Option 1: Organized tables, with 10s-100s of columns)

SolarPanel			
Zipcode	KW-Total	Carbon offset	...
94305	14.4	29	
94040	32.1	42	
94041	29.1	37.38	

Census		
Zipcode	Population Census	
94305	14301	
94040	20301	
94041	189	

Bike share locations			
Zipcode	Lat	Lng	
94305	35.1	122.12	
94305	35.2	122.13	
94041	35.1	121.27	
94041			

Pollution		
Zipcode	Particle count	
94305	40	
94040	22	
94041	57	



Trade offs?

- Reads? Writes?
- 100s - thousands of applications reading/writing data

⇒ Hybrids: 1 column → all columns

(Week 7: What's a good schema design?)

## Option 2: 'FrankenTable' (with 1000s of columns)

Omnidata							
SolarPanel				Census		Pollution	
Zipcode	KW-Total	Carbon offset	...	Population	Particle count	Lat	Lng
94305	14.4	29		14301	40	35.1	122.12
94305	14.4	29		14301	40	35.2	122.13
94305	14.4	29		14301	40	35.2	122.1
94305	14.4	29		14301	40	35.1	122.12

# Joins

Product(PName, Price, Category, Manufacturer)

Company(CName, StockPrice, Country)

*Ex:* Find all products under \$200 manufactured in Japan; return their names and prices.

```
SELECT PName, Price
FROM   Product, Company
WHERE  Manufacturer = CName
        AND Country='Japan'
        AND Price <= 200
```

A **join** between tables returns all unique combinations of their tuples **which meet some specified join condition**

# Joins

## Product

<u>PName</u>	Price	Category	Manufacturer
Gizmo	\$19	Gadgets	GizmoWorks
Powergizmo	\$29	Gadgets	GizmoWorks
SingleTouch	\$149	Photography	Canon
MultiTouch	\$203	Household	Hitachi

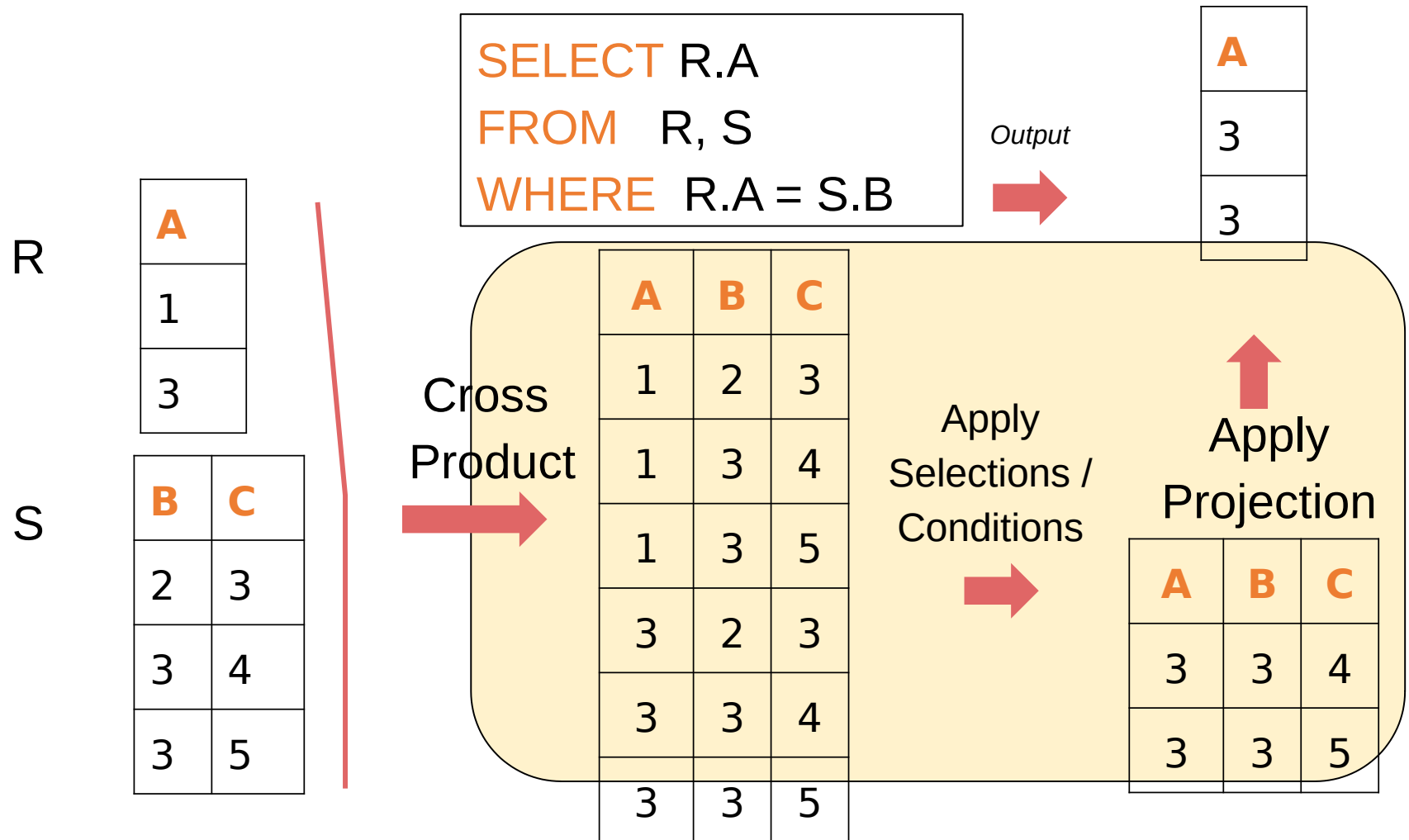
## Company

<u>CName</u>	Stock Price	Country
GizmoWorks	25	USA
Canon	65	Japan
Hitachi	15	Japan

```
SELECT PName, Price
FROM Product, Company
WHERE Manufacturer = CName
      AND Country='Japan'
      AND Price <= 200
```

PName	Price
SingleTouch	\$149

# An example of SQL semantics



---

Note: we say “semantics” not “execution order”

- The preceding slides show *what a join means*
- Not actually how the DBMS executes it under the covers

---

# Aggregation

```
SELECT AVG(price)
FROM Product
WHERE maker = "Toyota"
```

```
SELECT COUNT(*)
FROM Product
WHERE year > 1995
```

- SQL supports several **aggregation** operations:
  - SUM, COUNT, MIN, MAX, AVG

*Except COUNT, all aggregations  
apply to a single attribute*

---

# More Examples

```
Purchase(product, date, price, quantity)
```

```
SELECT SUM(price * quantity)  
FROM Purchase
```

```
SELECT SUM(price * quantity)  
FROM Purchase  
WHERE product = 'bagel'
```

What do these mean?



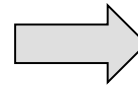
---

# Simple Aggregations

Purchase

Product	Date	Price	Quantity
bagel	10/21	1	20
banana	10/3	0.5	10
banana	10/10	1	10
bagel	10/25	1.50	20

```
SELECT SUM(price * quantity)
FROM Purchase
WHERE product = 'bagel'
```



50 (= 1\*20 + 1.50\*20)

---

# Grouping and Aggregation

Purchase(product, date, price, quantity)

```
SELECT product,  
        SUM(price * quantity) AS TotalSales  
FROM Purchase  
WHERE date > '10/1/2005'  
GROUP BY product
```

Find total sales  
after 10/1/2005  
per product.

Let's see what this means...

---

# Grouping and Aggregation

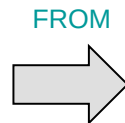
```
SELECT product,  
        SUM(price * quantity) AS TotalSales  
FROM Purchase  
WHERE date > '10/1/2005'  
GROUP BY product
```

## Semantics of the query:

1. Compute the **FROM** and **WHERE** clauses
2. Group by the attributes in the **GROUP BY**
3. Compute the **SELECT** clause: grouped attributes and aggregates

# 1. Compute the **FROM** and **WHERE** clauses

```
SELECT product, SUM(price*quantity) AS TotalSales  
FROM Purchase  
WHERE date > '10/1/2005'  
GROUP BY product
```



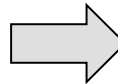
Product	Date	Price	Quantity
Bagel	10/21	1	20
Bagel	10/25	1.50	20
Banana	10/3	0.5	10
Banana	10/10	1	10

## 2. Group by the attributes in the **GROUP BY**

```
SELECT product, SUM(price*quantity) AS TotalSales
FROM Purchase
WHERE date > '10/1/2005'
GROUP BY product
```

Product	Date	Price	Quantity
Bagel	10/21	1	20
Bagel	10/25	1.50	20
Banana	10/3	0.5	10
Banana	10/10	1	10

GROUP BY



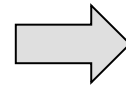
Product	Date	Price	Quantity
Bagel	10/21	1	20
	10/25	1.50	20
Banana	10/3	0.5	10
	10/10	1	10

### 3. Compute the **SELECT** clause: grouped attributes and aggregates

```
SELECT product, SUM(price*quantity) AS TotalSales
FROM Purchase
WHERE date > '10/1/2005'
GROUP BY product
```

Product	Date	Price	Quantity
Bagel	10/21	1	20
	10/25	1.50	20
Banana	10/3	0.5	10
	10/10	1	10

SELECT



Product	TotalSales
Bagel	50
Banana	15

---

# HAVING Clause

```
SELECT product, SUM(price*quantity)
FROM Purchase
WHERE date > '10/1/2005'
GROUP BY product
HAVING SUM(quantity) > 100
```

Same query as before, except that we consider only products that have more than 100 buyers

HAVING clauses contains conditions on **aggregates**

Whereas *WHERE* clauses condition on ***individual tuples...***

---

# RECAP: Joins

By default, joins in SQL are “inner joins”:

```
Product(name, category)
Purchase(prodName, store)
```

```
SELECT Product.name, Purchase.store
FROM   Product, Purchase
WHERE  Product.name = Purchase.prodName
```

```
SELECT Product.name, Purchase.store
FROM   Product
JOIN   Purchase ON Product.name = Purchase.prodName
```

Both equivalent:  
Both INNER JOINS!



---

# Outer Joins

- An **outer join** returns tuples from the joined relations that don't have a corresponding tuple in the other relations
  - I.e. If we join relations A and B on  $a.X = b.X$ , and there is an entry in A with  $X=5$ , but none in B with  $X=5$ ...
  - A LEFT OUTER JOIN will return a tuple (a, NULL)!
- Left outer joins in SQL:

```
SELECT Product.name, Purchase.store  
FROM Product  
LEFT OUTER JOIN Purchase ON  
Product.name = Purchase.prodName
```

Now we'll get products even if they didn't sell

# INNER JOIN

Product

name	category
iphone	media
Tesla	car
Ford Pinto	car

Purchase

prodName	store
iPhone	Apple store
Tesla	car
iPhone	Apple store

```
SELECT Product.name, Purchase.store
FROM Product
  INNER JOIN Purchase
    ON Product.name = Purchase.prodName
```

Note: another equivalent way to write an INNER JOIN!



name	store
iPhone	Apple store
iPhone	Apple store
Tesla	car

# LEFT OUTER JOIN

**Product**

name	category
iphone	media
Tesla	car
Ford Pinto	car

**Purchase**

prodName	store
iPhone	Apple store
Tesla	car
iPhone	Apple store

```
SELECT Product.name, Purchase.store
FROM Product
LEFT OUTER JOIN Purchase
ON Product.name = Purchase.prodName
```



name	store
iPhone	Apple store
iPhone	Apple store
Tesla	car
Ford Pinto	NULL

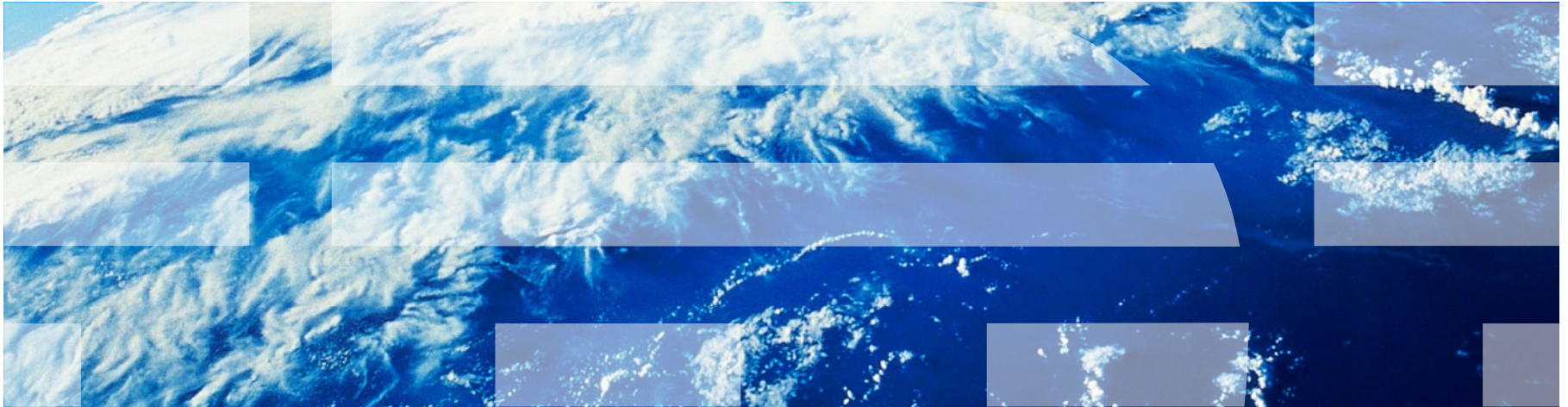
---

# Other Outer Joins

- Left outer join:
  - Include the left tuple even if there's no match
- Right outer join:
  - Include the right tuple even if there's no match
- Full outer join:
  - Include the both left and right tuples even if there's no match

---

# Nested Queries



---

## SQL is Compositional

Can construct powerful query chains (e.g.,  $f(g(\dots(x)))$ )

Inputs / outputs are multisets


⇒ Output of one query can be input to another (nesting)!

⇒ Including on same table

---

## Nested queries: Sub-queries Return Relations

Company(name, city)  
Product(name, maker)  
Purchase(id, product, buyer)



```
SELECT pr.maker
FROM Purchase p, Product pr
WHERE p.product = pr.name
      AND p.buyer = 'Mickey')
```

1. Companies making products bought by Mickey”
2. Location of companies?

---

## Nested queries: Sub-queries Return Relations

Company(name, city)  
Product(name, maker)  
Purchase(id, product, buyer)

```
SELECT c.city
FROM   Company c
WHERE  c.name IN (
    SELECT pr.maker
    FROM   Purchase p, Product pr
    WHERE  p.product = pr.name
          AND p.buyer = 'Mickey')
```

1. Companies making products bought by Mickey”
2. Location of companies?



# Subqueries Return Relations

You can also use operations of the form:

- $s > \text{ALL } R$
- $s < \text{ANY } R$
- $\text{EXISTS } R$

Ex: `Product(name, price, category, maker)`

```
SELECT name
FROM Product
WHERE price > ALL(
    SELECT price
    FROM Product
    WHERE maker = 'Gizmo-Works')
```

Find products that are more expensive than all those produced by “Gizmo-Works”

```
SELECT p1.name
FROM Product p1
WHERE p1.maker = 'Gizmo-Works'
AND EXISTS(
    SELECT p2.name
    FROM Product p2
    WHERE p2.maker <> 'Gizmo-Works'
    AND p1.name = p2.name)
```

`<>` means `!=`

Find ‘copycat’ products, i.e. products made by competitors with the same names as products made by “Gizmo-Works”

Note the scoping of the variables!

---

## Example: Complex Correlated Query

Product(name, price, category, maker, year)

```
SELECT DISTINCT x.name, x.maker
FROM   Product AS x
WHERE  x.price > ALL(
    SELECT y.price
    FROM   Product AS y
    WHERE  x.maker = y.maker
           AND y.year < 1972)
```

Find products (and their manufacturers) that are more expensive than all products made by the same manufacturer before 1972

Can be very powerful (also much harder to optimize)