

# Quantum Algorithms

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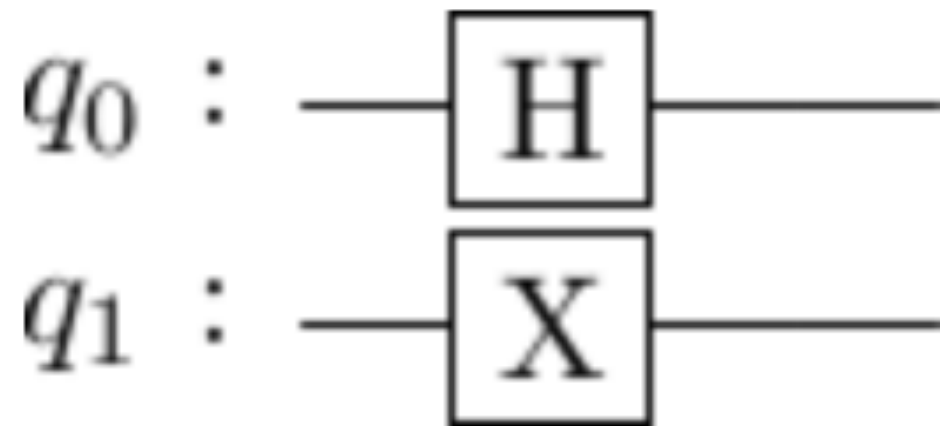
## Recap: we can make the whole quantum circuit as a transition matrix

qc = QuantumCircuit(2)  $X|q_1\rangle \otimes H|q_0\rangle = (X \otimes H)|q_1q_0\rangle$

qc.h(0)

qc.x(1)

qc.draw()



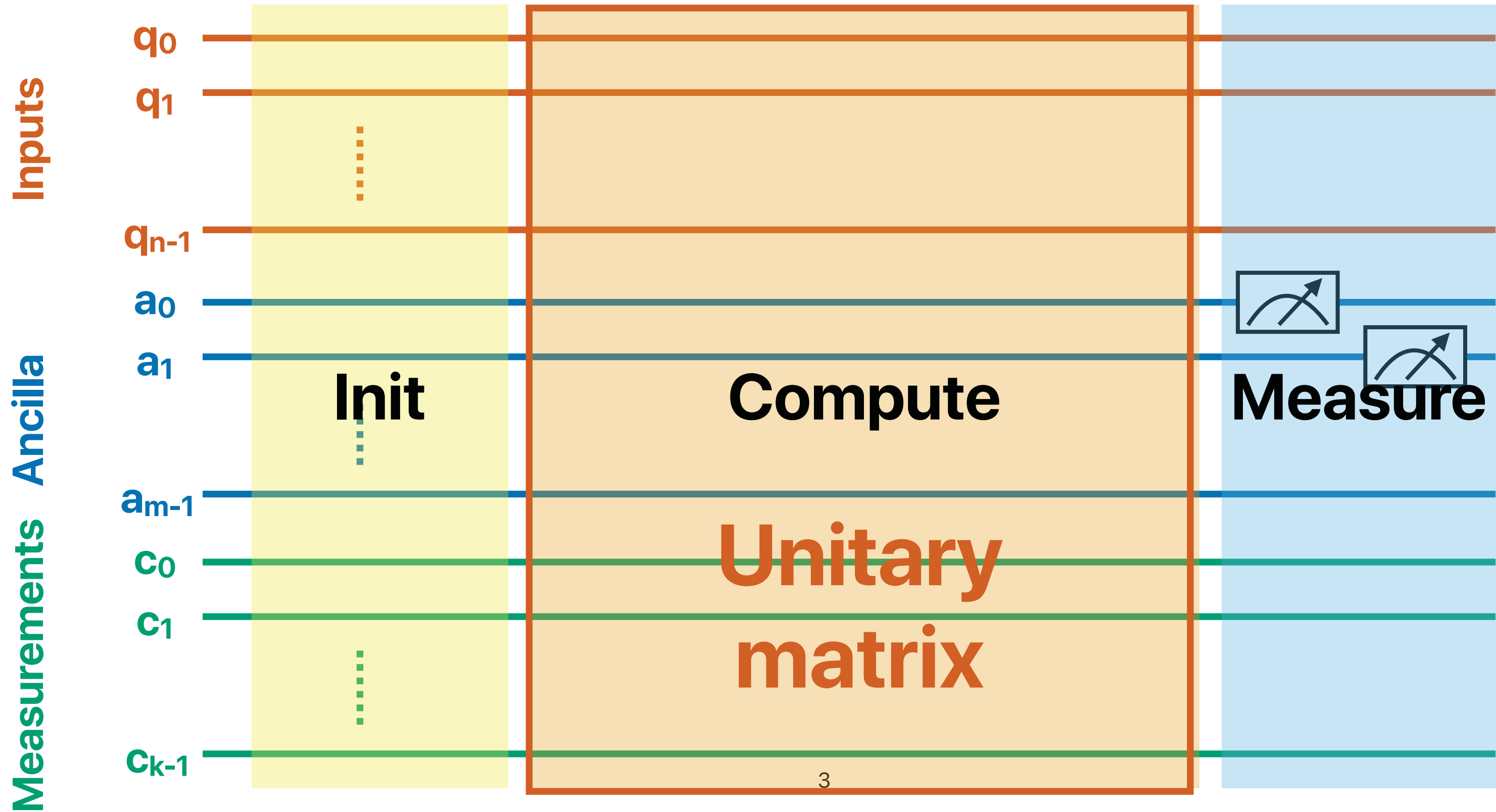
$$X \otimes H = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$= \frac{1}{\sqrt{2}} \begin{bmatrix} 0 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} & 1 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \\ 1 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} & 0 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \end{bmatrix}$$

$$= \frac{1}{\sqrt{2}} \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \\ 1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & H \\ H & 0 \end{bmatrix}$$

Unitary matrix

# Recap: Quantum circuits



# Recap: a state vector can be composed by two qubits

$$|ba\rangle = \begin{bmatrix} b_0a_0 \\ b_0a_1 \\ b_1a_0 \\ b_1a_1 \end{bmatrix} = \begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ c_3 \end{bmatrix}$$

$$\frac{1}{2} \begin{bmatrix} 1 \\ -1 \\ -1 \\ 1 \end{bmatrix}$$

$$b_0a_0 = \frac{1}{2}$$

$$b_0a_1 = -\frac{1}{2}$$

$$b_1a_0 = -\frac{1}{2}$$

$$b_1a_1 = \frac{1}{2}$$

$$\frac{a_0}{a_1} = -1$$

$$|a_0|^2 + |a_1|^2 = 1$$

$$\frac{b_0}{b_1} = -1$$

$$|b_0|^2 + |b_1|^2 = 1$$

$$a_0 = \frac{1}{\sqrt{2}}$$

$$a_1 = \frac{-1}{\sqrt{2}}$$

$$b_0 = \frac{-1}{\sqrt{2}}$$

$$b_1 = \frac{1}{\sqrt{2}}$$

# Recap: not all state vectors can be composed by two qubits

$$|ba\rangle = \begin{bmatrix} b_0a_0 \\ b_0a_1 \\ b_1a_0 \\ b_1a_1 \end{bmatrix} = \begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ c_3 \end{bmatrix}$$

$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

$$b_0a_0 = \frac{1}{\sqrt{2}}$$

$$b_0a_1 = 0$$

$$b_1a_0 = 0$$

$$b_1a_1 = \frac{1}{\sqrt{2}}$$

$$\frac{a_0}{a_1} = \infty$$

$$\frac{b_0}{b_1} = NaN$$

**This is not possible to be a product of two states!**

# Recap: entanglement

- The state is not a product of qubits
- Measuring one will tell us the state of the other and collapse its superposition — because the state cannot stand by itself
- Two or more quantum systems (or quantum particles) have a non-classical correlation, or shared quantum state, even if they are separated by a large distance.

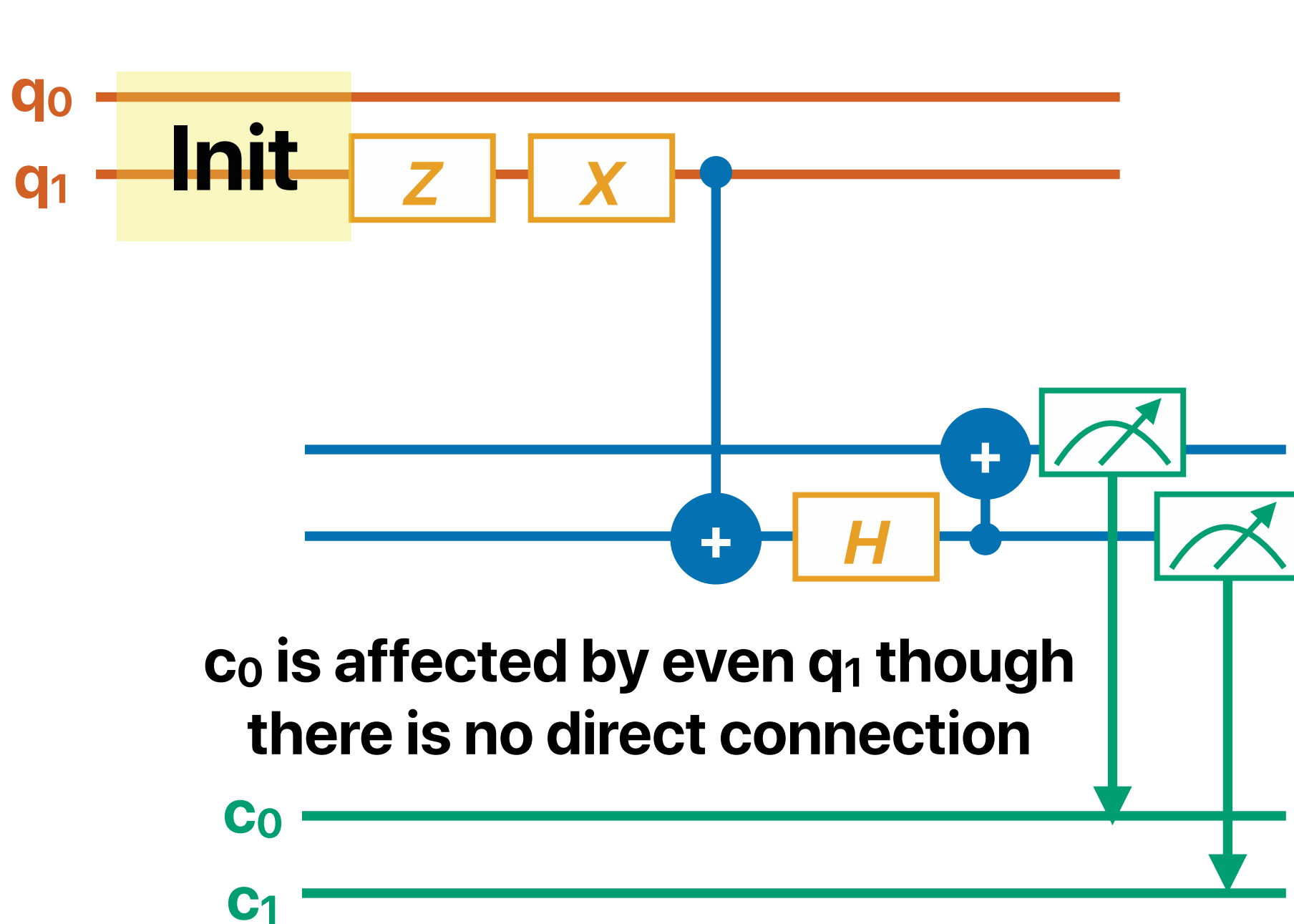
If we initialize the state as  $|\Phi^+\rangle$

$$|\Phi^+\rangle = \sqrt{\frac{1}{2}}(|0\cancel{0}\rangle + |1\cancel{1}\rangle) \quad 01 \rightarrow |\Psi^+\rangle = \sqrt{\frac{1}{2}}(|01\rangle + |10\rangle)$$

$$|\Phi^+\rangle = \sqrt{\frac{1}{2}}(|0\boxed{0}\rangle + |1\boxed{1}\rangle) \quad 10 \rightarrow |\Phi^-\rangle = \sqrt{\frac{1}{2}}(|00\rangle - |11\rangle)$$

$$|\Phi^+\rangle = \sqrt{\frac{1}{2}}(|0\boxed{0}\rangle + |1\boxed{1}\rangle) \quad 11 \rightarrow |\Psi^-\rangle = \sqrt{\frac{1}{2}}(|10\rangle - |01\rangle)$$

# The new implementation



$$00 \rightarrow |\Phi^+\rangle = \sqrt{\frac{1}{2}}(|00\rangle + |11\rangle)$$

$$01 \rightarrow |\Psi^+\rangle = \sqrt{\frac{1}{2}}(|01\rangle + |10\rangle)$$

$$10 \rightarrow |\Phi^-\rangle = \sqrt{\frac{1}{2}}(|00\rangle - |11\rangle)$$

$$11 \rightarrow |\Psi^-\rangle = \sqrt{\frac{1}{2}}(|10\rangle - |01\rangle)$$



# The Deutsch-Jozsa Algorithm

# The Deutsch-Jozsa Problem

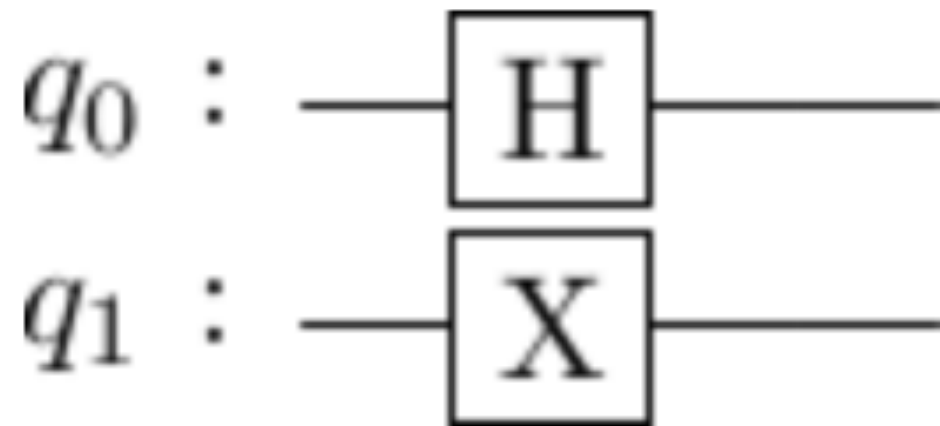
- Given a hidden Boolean function  $f$ , which takes as input a string of bits, and returns either 0 or 1, that is:  
 $f(\{x_0, x_1, x_2, \dots\}) \rightarrow 0 \text{ or } 1$ , where  $x_n$  is 0 or 1
- The given Boolean function is that it is guaranteed to either be balanced or constant
  - A constant function returns all 0s or all 1s for any input
  - A balanced function returns 0s for exactly half of all inputs and 1s for the other half
- Our task is to determine whether the given function is balanced or constant

# The classical solution

- Let's start by choosing two numbers and test their outputs
  - if  $f(0,0,0,\dots) \rightarrow 0$  and  $f(1,0,0,\dots) \rightarrow 1$ , then we know the given one is a balanced one!
  - What if  $f(0,0,0,\dots) \rightarrow 0$  and  $f(1,0,0,\dots) \rightarrow 0$ ? We have to try one more run...
- The worst case will need to go through exactly half of the input space + 1, that is  $2^{n-1} + 1$ , numbers
- The classical solution is therefore

## Recap: we can make the whole quantum circuit as a transition matrix

```
qc = QuantumCircuit(2)
qc.h(0)
qc.x(1)
qc.draw()
```



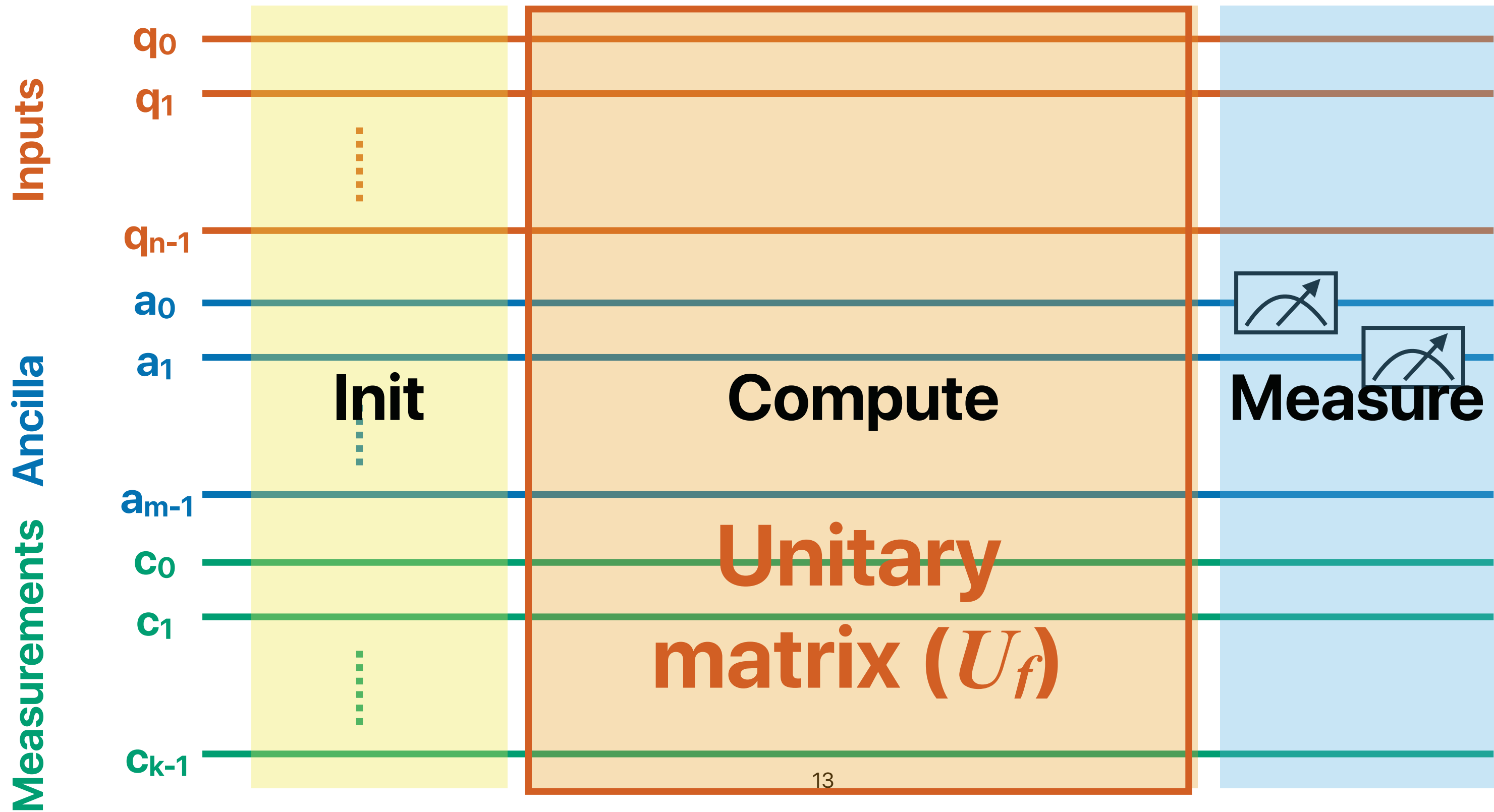
$$X \otimes H = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$= \frac{1}{\sqrt{2}} \begin{bmatrix} 0 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} & 1 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \\ 1 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} & 0 \times \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \end{bmatrix}$$

$$= \frac{1}{\sqrt{2}} \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \\ 1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & H \\ H & 0 \end{bmatrix}$$

Unitary matrix

# Recap: Quantum circuits



**Let's revisit CNOT gate, again**

# What's CNOT

- Flip the state of  $q_0$  if  $q_1$  is set  
 $q_1 q_0$

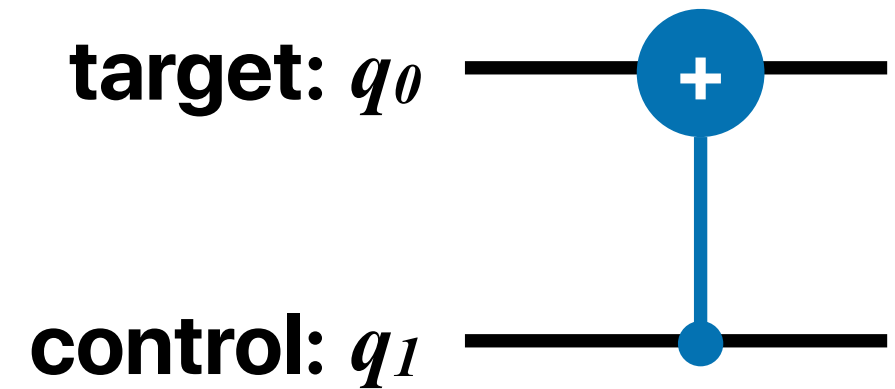
$$|00\rangle \rightarrow |00\rangle$$

$$|01\rangle \rightarrow |01\rangle$$

$$|10\rangle \rightarrow |11\rangle$$

$$|11\rangle \rightarrow |10\rangle$$

- Or can we say it's swapping the amplitudes of  $|10\rangle$  and  $|11\rangle$



$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

# What's CNOT

- Flip the state of  $q_1$  if  $q_0$  is set  
 $q_1 q_0$

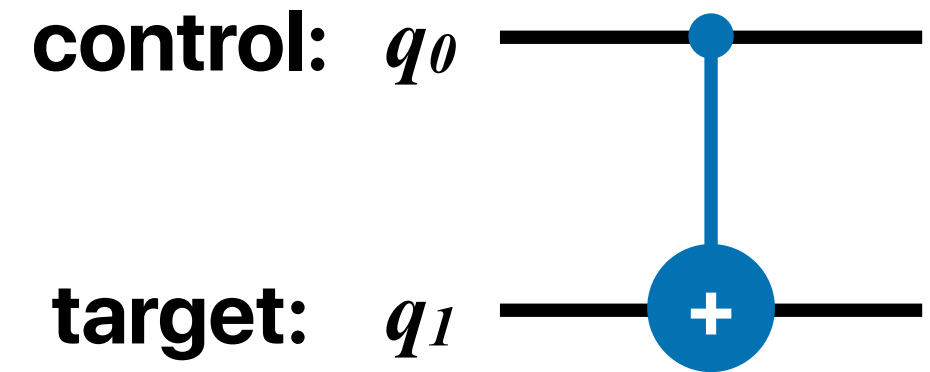
$$|00\rangle \rightarrow |00\rangle$$

$$|01\rangle \rightarrow |11\rangle$$

$$|10\rangle \rightarrow |10\rangle$$

$$|11\rangle \rightarrow |01\rangle$$

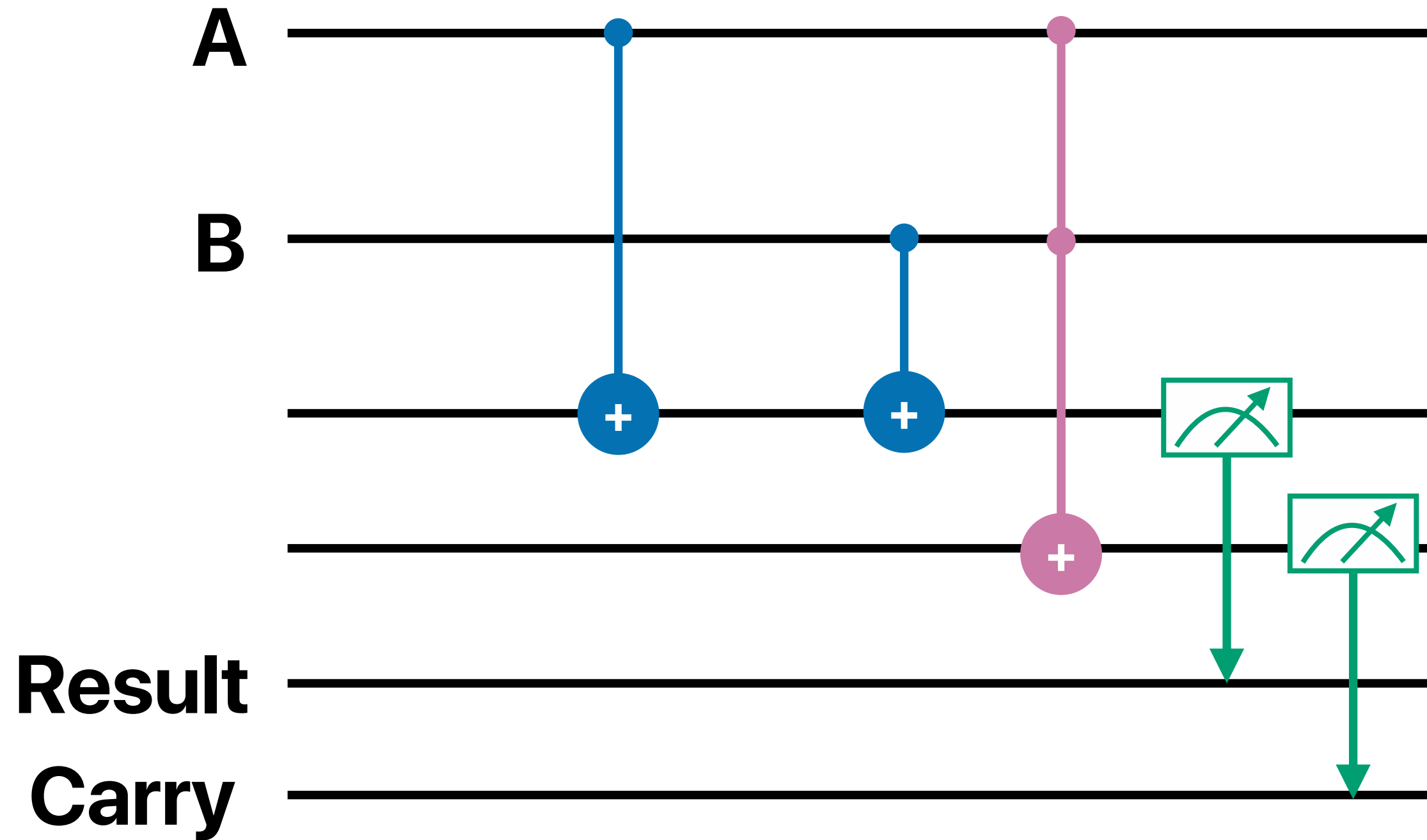
- Or can we say it's swapping the amplitudes of  $|01\rangle$  and  $|11\rangle$



$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$



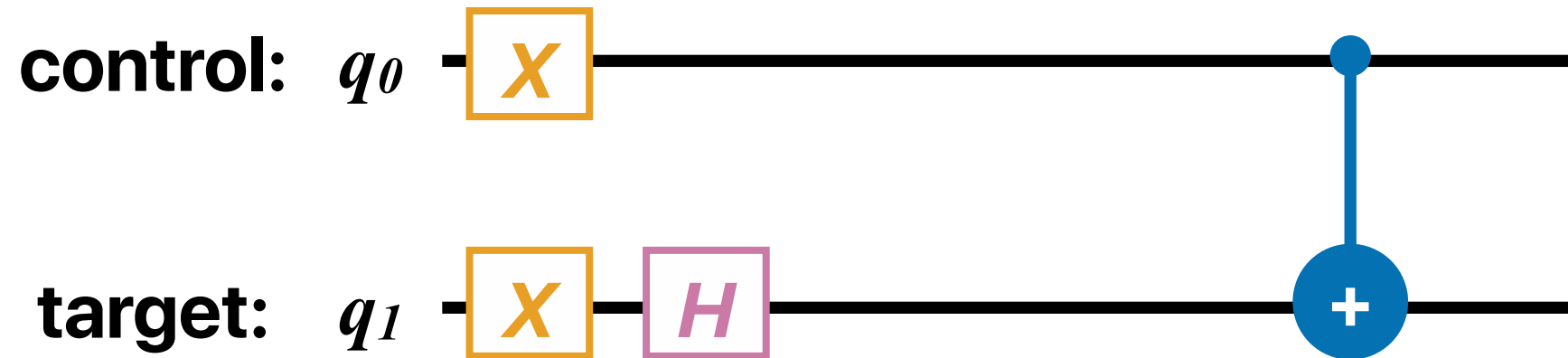
# 1-bit half adder



A	B	carry	result
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

# Let's look into this circuit

$$\begin{aligned} \cdot |+\rangle &= \sqrt{\frac{1}{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \\ \cdot |-\rangle &= \sqrt{\frac{1}{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix} \end{aligned}$$



$$\text{CNOT} | - 1 \rangle = - | - 1 \rangle$$

$$\begin{aligned} &= \begin{bmatrix} 0 \\ 1 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes \begin{bmatrix} 0 \\ 1 \end{bmatrix} \\ &= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix} \otimes \begin{bmatrix} 0 \\ 1 \end{bmatrix} \\ &= \frac{1}{\sqrt{2}} \begin{bmatrix} 0 \\ 1 \\ 0 \\ -1 \end{bmatrix} \end{aligned}$$

$$\begin{aligned} &= \frac{1}{\sqrt{2}} \begin{bmatrix} 0 \\ 1 \\ -1 \\ 0 \end{bmatrix} \end{aligned}$$

$$\begin{aligned} &= \frac{1}{\sqrt{2}} \begin{bmatrix} 0 \\ 1 \\ 0 \\ 1 \end{bmatrix} \end{aligned}$$

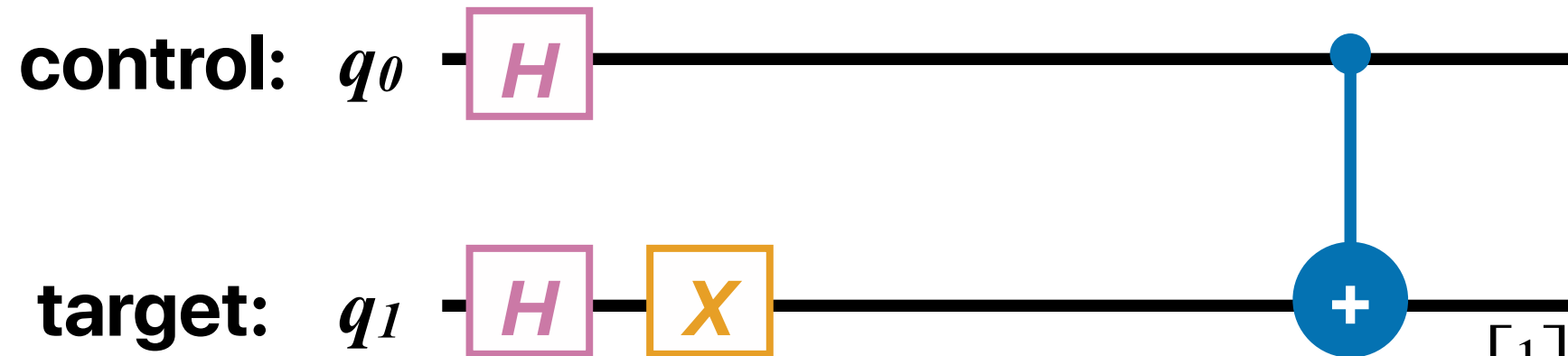
$q_0$  is changed!  $-1$  goes to  $q_0$ !

$q_1$  is not changed!

**Phase kickback**

# Let's look into this circuit

$$\begin{aligned} \cdot |+\rangle &= \sqrt{\frac{1}{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \\ \cdot |-\rangle &= \sqrt{\frac{1}{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix} \end{aligned}$$



$$= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$= \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

$$= \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

$$= \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

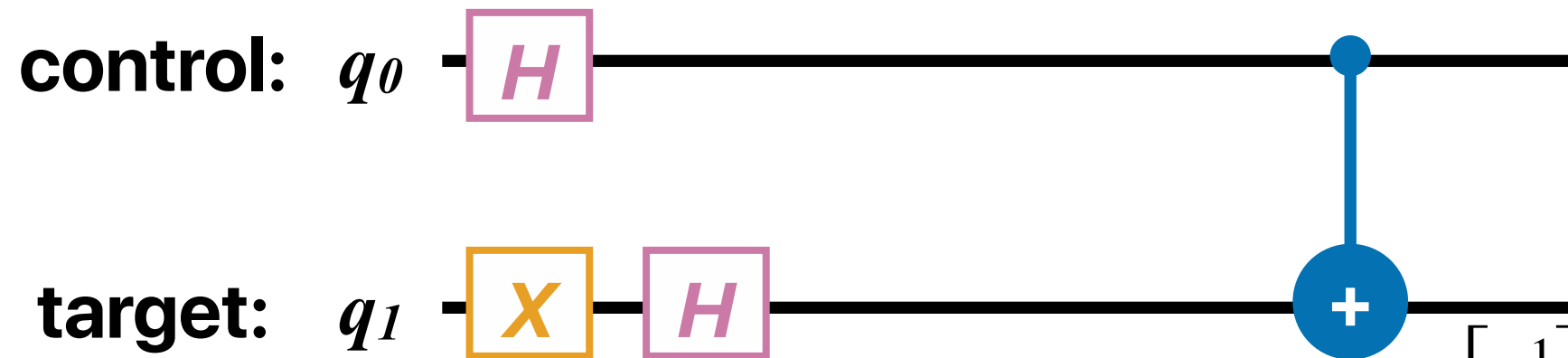
$$\text{CNOT} |++\rangle = \frac{1}{2}(|00\rangle + |01\rangle + |10\rangle + |11\rangle)$$

$$= |++\rangle$$

# Let's look into another circuit.

$$|+\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$|-\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$



$$= \begin{bmatrix} 0 \\ 1 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ -1 \\ -1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 1 \\ -1 \\ -1 \\ 1 \end{bmatrix}$$

$$= \frac{1}{2} \begin{bmatrix} 1 \\ 1 \\ -1 \\ -1 \end{bmatrix}$$

$$\text{CNOT} | - + \rangle = \frac{1}{2} (|00\rangle - |01\rangle - |10\rangle + |11\rangle)$$

$$= | - - \rangle \text{CNOT} | - + \rangle$$

$$= \frac{1}{2} (|00\rangle - |01\rangle - |10\rangle + |11\rangle) = | - - \rangle$$

# Phase kickback

- Kickback is where the eigenvalue added by a gate to a qubit is 'kicked back' into a different qubit via a controlled operation.

- For example, we saw that performing an X-gate on a  $|-\rangle$  qubit gives it the phase  $-1$
- When our control qubit is in either  $|0\rangle$  or  $|1\rangle$ , this phase affects the whole state, however it is a global phase and has no observable effects

$$\begin{aligned}\text{CNOT } | - 0 \rangle &= | - \rangle \otimes | 0 \rangle \\ &= | - 0 \rangle\end{aligned}$$

$$\begin{aligned}\text{CNOT } | - 1 \rangle &= X | - \rangle \otimes | 1 \rangle \\ &= - | - \rangle \otimes | 1 \rangle \\ &= - | - 1 \rangle\end{aligned}$$

- This can then be written as the two separable qubit states:

$$\begin{aligned}\text{CNOT } | - + \rangle &= | - \rangle \otimes \frac{1}{\sqrt{2}}(| 0 \rangle - | 1 \rangle) \\ &= | - - \rangle\end{aligned}$$

# An overview of Deutsch-Jozsa Algorithm

- Initialize  $n + 1$  qubits
- Transform these qubits into Hadamard basis: making each qubit 50%-50% of being 0 or 1
- Encode the given function as an unitary matrix (i.e., oracle)
- Return the qubits for measurements
- Measure the qubits to obtain the solution

# An overview of Deutsch-Jozsa Algorithm

- Initialize  $n + 1$  qubits
- Transform these qubits into Hadamard basis: making each qubit 50%-50% of being 0 or 1

$$|\psi_1\rangle = \sum_{x \in \{0,1\}^n} H|x\rangle = \sum_{x \in \{0,1\}^n} \frac{1}{\sqrt{2^{n+1}}} |x\rangle(|0\rangle - |1\rangle)$$

- Encode the given function as an unitary matrix (i.e., oracle)

$$U_f|\psi_1\rangle = |\psi_2\rangle = \sum_{x \in \{0,1\}^n} \frac{1}{\sqrt{2^{n+1}}} (-1)^{f(x)} |x\rangle(|0\rangle - |1\rangle)$$

- Return the qubits for measurements

- Measure the qubits to obtain the solution

$$\begin{aligned} |\psi_3\rangle &= \frac{1}{2^n} \sum_{x=0}^{2^n-1} (-1)^{f(x)} \left[ \sum_{y=0}^{2^n-1} (-1)^{x \cdot y} |y\rangle \right] \\ &= \frac{1}{2^n} \sum_{y=0}^{2^n-1} \left[ \sum_{x=0}^{2^n-1} (-1)^{f(x)} (-1)^{x \cdot y} \right] |y\rangle \end{aligned}$$

# Why this works

- Constant: When the oracle is constant, it has no effect (up to a global phase) on the input qubits, and the quantum states before and after querying the oracle are the same. Since the H-gate is its own inverse, in  $|\psi_3\rangle$  we reverse  $|\psi_1\rangle$  to obtain the initial quantum state of  $|00\dots 0\rangle$  in the first register.

$$H^{\otimes n} \begin{bmatrix} 1 \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \frac{1}{\sqrt{2^n}} \begin{bmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} \xrightarrow{\text{after } U_f} H^{\otimes n} \frac{1}{\sqrt{2^n}} \begin{bmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

- Balanced: After creating  $|\psi_1\rangle$ , our input register is an equal superposition of all the states in the computational basis. When the oracle is balanced, phase kickback adds a negative phase to exactly half these states

$$U_f \frac{1}{\sqrt{2^n}} \begin{bmatrix} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} = \frac{1}{\sqrt{2^n}} \begin{bmatrix} -1 \\ 1 \\ -1 \\ \vdots \\ 1 \end{bmatrix}$$



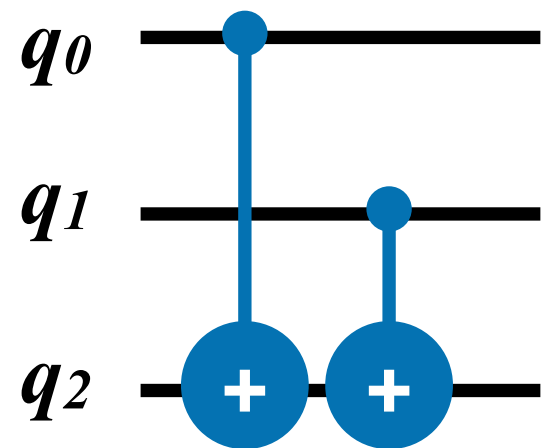
# An example of the algorithm

$$f(0,0) = 0$$

$$f(0,1) = 1$$

$$f(1,0) = 1$$

$$f(1,1) = 0$$

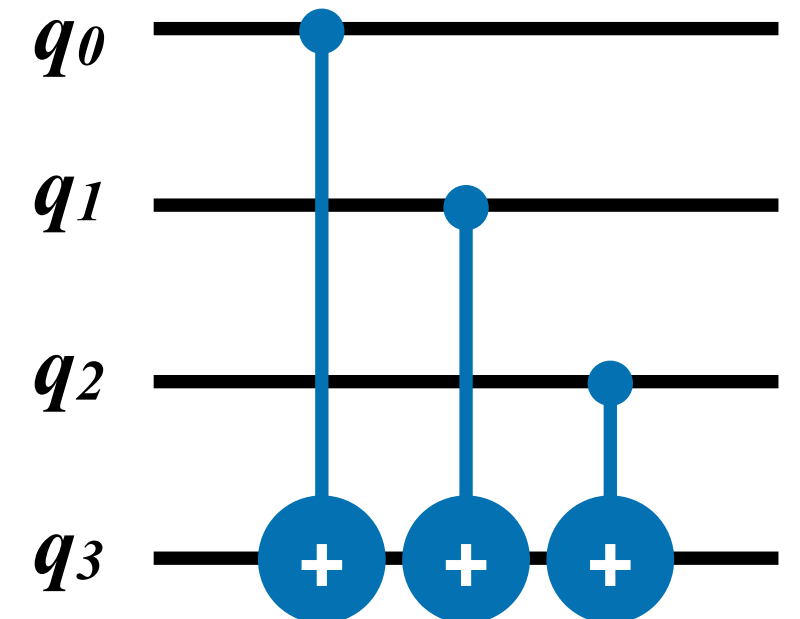


$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

# Create the oracle

# Create the oracle for Deutsch-Jozsa algorithm

- Constant
  - if  $f(x) = 0$ , then apply the  $I$  gate to the qubit in register 2.
  - if  $f(x) = 1$ , then apply the  $X$  gate to the qubit in register 2.
- Balanced
  - Many solutions. One is performing a CNOT for each qubit in register 1, with the qubit in register 2 as the target.



# Announcement

- Assignments will be reduced from 4 to 2
  - You may consider them as take-home midterm and final
  - Assignment 1 — release on 2/11, due 2/18
  - Assignment 2 — release on 3/4, due 3/11
- No lectures on 2/13, 3/4, and 3/6