# Competitive Programming Library

# rwx 2025/2026

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# 1 Templates

# 1.1 Setup

#### 1.1.1 IO Manipulation

## Input/Output

```
#include <bits/stdc++.h>
freopen("input.txt", "r", stdin);
freopen("output.txt", "w", stdout);

#define fastI0  \
   ios_base::sync_with_stdio(false), cin.tie(nullptr), cout.tie(nullptr);
```

#### 1.1.2 GCC Compiler Optimization (Vectorization)

# GCC Opt

```
// Ref: USACO guide
// will make GCC auto-vectorize for loops and optimizes floating points
   better (assumes associativity and turns off denormals).
#pragma GCC optimize ("Ofast")
// can double performance of vectorized code, but causes crashes on old
   machines.
#pragma GCC target ("avx,avx2,fma")

// slows down run time but throws a Runtime Error if an overflow occured
#pragma GCC optimize("trapy")
```

# 1.2 MOD Template

```
constexpr int MOD = 1e9+7; // must be a prime number
int add(int a, int b) {
   int res = a+b;
   if(res >= MOD) return res -= MOD;
}
```

```
int sub(int a, int b) {
    int res = a-b;
    if (res < 0) return res += MOD;
int power(int a, int e) {
    int res = 1;
    while(e) {if(e & 1) res = res * a % MOD; a = a * a % MOD;
    return res;
// note -> from Fermat little theorem, we can deduce that
// a^(BIG) % prime = (a ^ (BIG%(prime - 1)))%prime
// because every a^(prime - 1) is equal to 1
// from the identity says a^(prime - 1) % prime = 1
int inverse(int a) {
    return power(a, MOD-2);
int div(int a, int b) {
    return a * inverse(b) % MOD;
                       MOD Template
```

#### 1.3 Macros

#### Macros

```
#define getBit(n, k) (n >> k) // & 1
#define ON(n, idx) (n | (111 << idx))
#define OFF(n, idx) (n & ~(111 << idx))
#define toggle(n, idx) ((n) ^ (111<<(idx)))
#define gray(n) (n ^ (n >> 1))
#define bitCount(x) (__builtin_popcountl1(x))
```

```
#define clz(x) (__builtin_clzll(x))
#define ctz(x) (__builtin_ctzll(x))
#define uniq(x) x.resize(unique(x.begin(), x.end())-x.begin());
#define angle(a) (atan2((a).imag(), (a).real()))
//#define vec(a, b) ((b)-(a))
#define same(v1, v2) (dp(vec(v1,v2),vec(v1,v2)) < EPS)
#define dotProduct(a, b) ((conj(a)*(b)).real()) // a*b cos(T), if zero ->
#define crossProduct(a, b) ((conj(a)*(b)).imag()) // a*b sin(T), if zero
   -> parallel
//#define length(a) (hypot((a).imag(), (a).real()))
#define normalize(a) ((a)/length(a))
#define rotateO(v, ang) ((v)*exp(point(0,ang)))
#define rotateA(p, ang, about) (rotateO(vec(about,p),ang)+about)
#define reflectO(v, m) (conj((v)/(m))*(m))
#define ceil_i(a, b) (((ll)(a)+(ll)(b-1))/(ll)(b))
#define floor_i(a, b) (a/b)
#define round_i(a, b) ((a+(b/2))/b) // if a>0
#define round_m(a, b) ((a-(b/2))/b) // if a<0
#define round_multiple(n, m) round_i(n,m)*m // round to multiple if
   specified element
const double PI = acos(-1.0);
```

## 1.4 Grid Navigation

#### Grid Nav

```
// knight moves on a chess board
int dx[] = { -2, -1, 1, 2, -2, -1, 1, 2 };
int dy[] = { -1, -2, -2, -1, 1, 2, 2, 1 };

// Grid up, down, right, left (Moves for Chess Rook)
int dx[4] = {1, -1, 0, 0};
int dy[4] = {0, 0, 1, -1};

// Grid cell all neighbours
const int dx[8] = {1, 0, -1, 0, 1, 1, -1, -1}
const int dy[8] = {0, 1, 0, -1, -1, 1, -1, 1};
```

```
// Grid Diagonal (Moves for Chess Bishop)
int dx[] = {1, 1, -1, -1};
int dy[] = {1, -1, 1, -1};
```

#### 1.5 Integer 128

i128

```
typedef __int128 i128;
int128 read() {
   \_int128 x = 0, f = 1;
   char ch = getchar();
   while (ch < '0' || ch > '9') {
      if (ch == '-') f = -1;
       ch = getchar();
   }
   while (ch >= '0' && ch <= '9') {
      x = x * 10 + ch - '0';
       ch = getchar();
   }
   return x * f;
void print(__int128 x) {
   if (x < 0) {
      putchar('-');
      x = -x;
   }
   if (x > 9) print(x / 10);
   putchar(x % 10 + '0');
bool cmp(__int128 x, __int128 y) { return x > y; }
```

### 1.6 Matrix Expo

MExpo

```
constexpr int MAX_N = 201, MOD = 1e9+7;
```

```
struct Matrix { int mat[MAX_N][MAX_N]; }; // we return a 2D array
11 mod(11 a, 11 m) { return ((a%m)+m) % m; } // ensure positive answer
Matrix matMul(Matrix a, Matrix b) { // normally O(n^3)
   Matrix ans{}; // but O(1) as n = 2
   for (auto & i : ans.mat)
       for (int & j : i)
           i = 0;
   for (int i = 0; i < MAX_N; ++i)</pre>
       for (int k = 0; k < MAX_N; ++k) {
           if (a.mat[i][k] == 0) continue; // optimization
           for (int j = 0; j < MAX_N; ++j) {
               ans.mat[i][j] += mod(a.mat[i][k], MOD) * mod(b.mat[k][j],
              ans.mat[i][j] = mod(ans.mat[i][j], MOD);
           }
   return ans;
}
Matrix matPow(Matrix base, int p) { // normally O(n^3 log p)
   Matrix ans{}; // but O(\log p) as n = 2
   for (int i = 0; i < MAX_N; ++i)</pre>
       for (int j = 0; j < MAX_N; ++j)
           ans.mat[i][j] = (i == j); // prepare identity matrix
   while (p) { // iterative D&C version
       if (p&1) // check if p is odd
           ans = matMul(ans, base); // update ans
       base = matMul(base, base); // square the base
       p >>= 1; // divide p by 2
   }
   return ans;
Matrix matMul(Matrix a, Matrix b, int p, int q, int r) { // O(pqr)
   Matrix c{};
   for (int i = 0; i < p; ++i)
       for (int j = 0; j < r; ++j) {
           c.mat[i][j] = 0;
           for (int k = 0; k < q; ++k)
              c.mat[i][j] += a.mat[i][k] + b.mat[k][j];
       }
   return c;
}
```

# 2 Dynamic Programming

## 2.1 Some dp patterns

#### Maximumu/Minimum path cost

```
const int MAX = 21;
int grid[MAX][MAX];
int mem[MAX][MAX];
int n = 20;
bool valid(int r, int c){
 return r >= 0 \&\& r < n \&\& c >= 0 \&\& c < n:
int maxPathSum(int r, int c){
   if(!valid(r,c)){
       return 0:
   }
   if(r == n-1 \&\& c == n-1){
       return mem[r][c] = grid[r][c];
   // available moves
   int path1 = maxPathSum(r+1,c);
   int path2 = maxPathSum(r,c+1);
   return grid[r][c] + max(path1,path2);
```

## add operators between numbers to get max prod/sum

```
// put +, -, between sequence of numbers such that the sum is divisible by
    k, and maximum as possible
const int MAX = 21;
long long mem[MAX][MAX];
const int n = 20;
int k = 4; // example
int v[20];
int fix(int a){
    return (a % k + k) % k;
}
long long tryAll(int pos, int mod){
```

```
long long &ret = mem[pos][mod];
if(ret != -1){
    return ret;
}
if(pos == n){
    return ret = mod == 0;
}
if(tryAll(pos+1,fix(mod + v[pos])) || tryAll(pos+1,fix(mod-v[pos]))){
    return ret = 1;
}
return ret = 0;
}
```

#### pick choices with no two similar consecutive choices

```
// pick minimum of choinces costs with no two similar consecutive choices
const int choices = 4;
const int n = 20;
int MAX = n;
int mem[MAX][choices];
const int 00 = 1e6+1;
int minCost(int pos, int lastChoice){
   if(pos == n){
       return 0; // invalid move
   }
   int &ret = mem[pos][lastChoice];
   if(ret != -1){
       return ret;
   }
   ret = 00; // want to minimze
   // let choices are 0, 1, 2
   if(lastChoice != 0){
       ret = min(ret, minCost(pos+1,0));
   }
   if(lastChoice != 1){
      ret = min(ret, minCost(pos+1,1));
   if(lastChoice != 2){
       ret = min(ret, minCost(pos+1,2));
   }
```

```
return ret;
```

#### sum S and max/min Product

```
int maxK:
ll mem[21][101]; // k, and s
// You are given an integer s and an integer k. Find k positive integers
   a1, a2, ..., ak
// such that their sum is equal to s and their product is the maximal
   possible. Return their product.
11 maxProd(int k, int rem)
if(k == maxK){
 // base case
 if(rem == 0)
 return 1;
 return 0;
if(rem == 0) // invalid case
 return 0;
11 &ret = mem[k][rem];
if(ret != -1)
 return ret;
ret = 0;
for (int i = 1; i <= rem; ++i) {</pre>
 ll sol = maxProd(k+1, rem - i) * i;
 ret = max(ret, sol);
return ret;
```

#### 2.2 DP solutions

#### 2.2.1 Max Subarray sum (Kadane's Algorithm)

#### Max Subarray sum

```
int maxSubarraySum(vector<int>& arr, int len) {
   int ans = INT_MIN, cur = 0;

   for (int i = 0; i < len; i++) {
      cur = cur + arr[i];
      if (ans < cur)
            ans = cur;

      if (cur < 0)
            cur = 0;
   }

   return ans;
}</pre>
```

#### 2.2.2 Maximum Subarray Alternating Sum

#### Maximum Subarray Alternating Sum

```
/* REF: GeeksForGeeks
Input: arr[] = \{-4, -10, 3, 5\}
Output: 9
Explanation: Subarray \{arr[0], arr[2]\} = \{-4, -10, 3\}. Therefore, the sum
     of this subarray is 9.
*/
int maxSubarraySumALT(vector<int>& a, int len) {
   int ans = INT_MIN, cur = 0;
   for (int i = 0; i < len; i++) {</pre>
       if (i % 2 == 0)
           cur = max(cur + a[i], a[i]);
       else
           cur = max(cur - a[i], -a[i]);
       ans = max(ans, cur);
   }
   cur = 0;
```

```
for (int i = 0; i < len; i++) {
    if (i % 2 == 1)
        cur = max(cur + a[i], a[i]);
    else
        cur = max(cur - a[i], -a[i]);
    ans = max(ans, cur);
}
return ans;
}</pre>
```

# 2.2.3 Count number of DISTINCT ordered ways to produce coins sums to x

#### Count distinct

```
For example, if the coins are \{2,3,5\} and the desired sum is 9, there
    are 3 ways:
2+2+5
3+3+3
2+2+2+3
*/
int n, x;
cin >> n >> x;
vector<int> coins(n);
read(coins);
vector dp(x + 1, 0);
dp[0] = 1;
for (int i = 0; i < n; ++i) {</pre>
for (int j = coins[i]; j <= x; ++j) {</pre>
 dp[j] = add(dp[j], dp[j - coins[i]]);
cout << dp[x] << el;</pre>
```

# 2.2.4 Min absolute difference between 2 elements from (L, R) (DP Ranges)

#### Min absolute difference

```
const int N = 1e4 + 1;
int dp[N][N];
int n;
cin >> n;
vector<int> a(n);
read(a);
for (int i = 0; i < n; ++i) dp[i][i] = 1e6; // INF, you can't take the
    element with it self
for (int i = 1; i < n; ++i) dp[i - 1][i] = abs(a[i] - a[i - 1]);
for (int len = 3; len <= n; ++len) {</pre>
for (int l = 0, r = len - 1; r < n; ++1, ++r) {
 dp[l][r] = min(dp[l][r - 1], dp[l + 1][r]);
 dp[l][r] = min(dp[l][r], abs(a[l] - a[r]));
}
}
int q;
cin >> q;
while (q--) {
int 1, r;
 cin >> 1 >> r;
--1, --r;
 cout << dp[l][r] << el;
```

## 2.2.5 Longest common subsequence between 2 Strings

$$dp[i][j] = \begin{cases} max(dp[i-1][j], dp[i][j-1]) & \text{if } A_i \neq B_j \\ dp[i-1][j-1] + 1 & \text{if } A_i = B_j \end{cases}$$

#### LIS 2 Strings

```
// REF: USACO guide
int longestCommonSubsequence(string a, string b) {
int dp[a.size()][b.size()];
for (int i = 0; i < a.size(); i++) { fill(dp[i], dp[i] + b.size(), 0); }</pre>
for (int i = 0; i < a.size(); i++) {</pre>
 if (a[i] == b[0]) dp[i][0] = 1;
 if (i != 0) dp[i][0] = max(dp[i][0], dp[i - 1][0]);
for (int i = 0; i < b.size(); i++) {</pre>
 if (a[0] == b[i]) dp[0][i] = 1;
 if (i != 0) dp[0][i] = max(dp[0][i], dp[0][i - 1]);
for (int i = 1; i < a.size(); i++) {</pre>
 for (int j = 1; j < b.size(); j++) {</pre>
  if (a[i] == b[j]) {
   dp[i][j] = dp[i - 1][j - 1] + 1;
  } else {
   dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
 }
return dp[a.size() - 1][b.size() - 1];
```

#### **2.2.6** Longest common subsequence $\mathcal{O}(n^2)$

#### LIS

```
// REF: cp-algorithms
int lis(vector<int> const& a) {
   int n = a.size();
   vector<int> d(n, 1);
   for (int i = 0; i < n; i++) {
      for (int j = 0; j < i; j++) {
        if (a[j] < a[i])
            d[i] = max(d[i], d[j] + 1);
      }
}
int ans = d[0];
   for (int i = 1; i < n; i++) {</pre>
```

```
ans = max(ans, d[i]);
   }
   return ans;
}
// Restoring
vector<int> lis(vector<int> const& a) {
   int n = a.size();
   vector\langle int \rangle d(n, 1), p(n, -1);
   for (int i = 0; i < n; i++) {</pre>
       for (int j = 0; j < i; j++) {</pre>
           if (a[j] < a[i] && d[i] < d[j] + 1) {</pre>
               d[i] = d[j] + 1;
               p[i] = j;
           }
   }
   int ans = d[0], pos = 0;
   for (int i = 1; i < n; i++) {</pre>
       if (d[i] > ans) {
           ans = d[i];
           pos = i;
       }
   }
   vector<int> subseq;
   while (pos !=-1) {
       subseq.push_back(a[pos]);
       pos = p[pos];
   reverse(subseq.begin(), subseq.end());
   return subseq;
```

#### 2.2.7 Longest common subsequence Binary Search $\mathcal{O}(n + \log N)$

#### LIS

```
int lisBS(vector<int> const& a) {
  int n = a.size();
  const int INF = 1e9;
  vector<int> d(n+1, INF);
```

```
d[0] = -INF;

for (int i = 0; i < n; i++) {
    int l = upper_bound(d.begin(), d.end(), a[i]) - d.begin();
    if (d[1-1] < a[i] && a[i] < d[1])
        d[1] = a[i];
}

int ans = 0;
for (int l = 0; l <= n; l++) {
    if (d[1] < INF)
        ans = 1;
}
return ans;</pre>
```

# 3 Bit Manipulation

# 3.1 Subset Operations

## count subsets with give sum

```
int countDistinctSubsetsWithSum(vector<int>& arr, int n, int k) {
    // Count distinct subsets of array arr that sum up to k
    vector<int> dp(k + 1, 0);
    dp[0] = 1;
    for (int i = 0; i < n; ++i) {
        for (int j = k; j >= arr[i]; --j) {
            dp[j] += dp[j - arr[i]];
        }
    }
    return dp[k]; // Number of distinct subsets with sum k
}
```

#### max xor of any subset of elements in the array

```
int maximalSubsetXOR(vector<int>& arr, int n) {
    // Find the maximum XOR of any subset of elements in array arr
    int maxXor = 0;
    for (int mask = 0; mask < (1 << n); ++mask) {</pre>
```

```
int xorSum = 0;
for (int i = 0; i < n; ++i) {
    if (mask & (1 << i)) {
        xorSum ^= arr[i];
    }
}
maxXor = max(maxXor, xorSum);
}
return maxXor;</pre>
```

#### min xor of any subset

```
int minimumSubsetXOR(vector<int>& arr, int n) {
    // Find the minimum XOR of any pair of elements in array arr
    int minSubsetXor = INT_MAX;
    for (int mask = 0; mask < (1 << n); ++mask) {
        int xorSum = 0;
        for (int i = 0; i < n; ++i) {
            if (mask & (1 << i)) {
                xorSum ^= arr[i];
            }
        }
        minSubsetXor = min(minSubsetXor, xorSum);
    }
    return minSubsetXor;
}</pre>
```

## $subset\ generation$

```
void subsetGeneration(int x, int n) {
    // Generate all non-empty subsets of a set represented by an integer x
    for (int subset = x; subset > 0; subset = (subset - 1) & x) {
        // Process subset
        cout << subset << endl;
    }
}</pre>
```

check if subset of elements in the array sum up to k
void subsetSumCheck(vector<int>& arr, int n, int k) {

```
// Check if a subset of elements in array arr sums up to k
for (int subset = 0; subset < (1 << n); ++subset) {
   int sum = 0;
   for (int i = 0; i < n; ++i) {
      if (subset & (1 << i)) {
        sum += arr[i];
      }
   }
   if (sum == k) {
      // Found subset with sum k
      cout << "Subset with sum " << k << ": " << subset << endl;
   }
}</pre>
```

#### max subset sum mod m

```
int subsetWithMaxSumModuloM(vector<int>& arr, int n, int m) {
   // Find the maximum subset sum modulo m
   vector<int> dp(m, -1);
   dp[0] = 0;
   int currentMod = 0;
   for (int i = 0; i < n; ++i) {</pre>
       currentMod = (currentMod + arr[i]) % m;
       for (int j = 0; j < m; ++j) {
           if (dp[j] != -1) {
              dp[(j + currentMod) % m] = max(dp[(j + currentMod) % m], dp
                  [j] + arr[i]);
          }
       }
       dp[currentMod] = max(dp[currentMod], arr[i]);
   }
   return dp[0]; // Maximum subset sum modulo m
```

#### iterate over all supersets represented by x

```
void iterateOverSupersets(int x, int n) {
    // Iterate over all supersets of a set represented by x
    int subset = x;
    do {
        // Process subset
```

```
cout << subset << endl;
subset = (subset + 1) | x;
} while (subset <= (1 << n) - 1);</pre>
```

# 4 Algorithms

#### 4.1 FFT

#### FFT Algorithm

```
constexpr 11 mod = 998244353, root = 3;
ll modpow(ll b, ll e, ll m) {
   ll ans = 1;
   for (; e; b = b * b % m, e /= 2)
       if (e \& 1) ans = ans * b % m;
   return ans;
}
// Primitive Root of the mod of form 2^a * b + 1
int generator () {
   vector<int> fact;
   int phi = mod-1, n = phi;
   for (int i=2; i*i<=n; ++i)</pre>
       if (n % i == 0) {
           fact.push_back (i);
           while (n \% i == 0)
              n /= i;
       }
   if (n > 1)
       fact.push_back (n);
   for (int res=2; res<=mod; ++res) {</pre>
       bool ok = true;
       for (size_t i=0; i<fact.size() && ok; ++i)</pre>
           ok &= modpow(res, phi / fact[i], mod) != 1;
       if (ok) return res;
   }
   return -1;
}
```

```
void ntt(vector<ll> &a) {
   int n = a.size(), L = 31 - __builtin_clz(n);
   static vector<11> rt(2, 1); // erase the static if you want to use two
   for (static int k = 2, s = 2; k < n; k *= 2, s++) { // erase the
       static if you want to use two moduli;
       rt.resize(n);
       ll z[] = \{1, modpow(root, mod >> s, mod)\};
       for (int i = k; i < 2*k; ++i) rt[i] = rt[i / 2] * z[i & 1] % mod;
   }
   vector<int> rev(n);
   for (int i = 0; i < n; ++i) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
   for (int i = 0; i < n; ++i) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
   for (int k = 1; k < n; k *= 2) {
       for (int i = 0; i < n; i += 2 * k) {
          for (int j = 0; j < k; ++j) {
              ll z = rt[j + k] * a[i + j + k] % mod, &ai = a[i + j];
              a[i + j + k] = ai - z + (z > ai ? mod : 0);
              ai += (ai + z >= mod ? z - mod : z):
          }
       }
   }
vector<ll> multi(const vector<ll> &a, const vector<ll> &b) {
   if (a.empty() || b.empty()) return {};
   int s = static_cast<11>(a.size()) + static_cast<11>(b.size()) - 1, B =
        32 - \_builtin\_clz(s), n = 1 << B;
   int inv = modpow(n, mod - 2, mod);
   vector<ll> L(a), R(b), out(n);
   L.resize(n), R.resize(n);
   ntt(L), ntt(R);
   for (int i = 0; i < n; ++i) out[-i & (n - 1)] = L[i] * R[i] % mod *
       inv % mod;
   ntt(out);
   return {out.begin(), out.begin() + s};
vector<int> poly_pow(vector<int> poly, int p, int limit = 1e9) {
   vector<int> ans{1};
   while (p) {
       if(p&1) ans = multi(ans, poly);
       poly = multi(poly, poly);
```

```
ans.resize(limit + 1);
  poly.resize(limit + 1);
  p >>= 1;
}
return ans;
}
```

#### 4.2 MO

#### MO Algorithm

```
// MO
           -> O(N+Q SQRT(N)) <= 10^5
const int N = 1e5+5, M = 1e5+5;
int n, m;
int nums[N], q_ans[M];
struct query {
   int idx, block_idx, l, r;
   query() = default;
   query(int _1, int _r, int _idx) {
       idx = _idx;
       r = _r - 1;
       1 = _1 - 1;
       block_idx = _l / sqrt(n);
   }
   bool operator <(const query & y) const {</pre>
       if(y.block_idx == block_idx) return r < y.r;</pre>
       return block_idx < y.block_idx;</pre>
   }
};
int freq[N], ans;
void add(int idx) {
   freq[nums[idx]]++;
   if (freq[nums[idx]] == 2) ans++;
}
void remove(int idx) {
```

```
freq[nums[idx]]--;
    if (freq[nums[idx]] == 1) ans--;
cin >> n >> m:
for (int i = 0; i < n; ++i) cin >> nums[i];
vector<query> Query(m);
for (int i = 0; i < m; ++i) {</pre>
   int 1, r; cin >> 1 >> r;
    Query[i] = query(1, r, i);
sort(Query.begin(), Query.end());
int 10 = 1, r0 = 0;
for (int i = 0; i < m; ++i) {</pre>
    while (10 < Query[i].1) remove(10++);</pre>
    while (10 > Query[i].1) add(--10);
    while (r0 < Query[i].r) add(++r0);</pre>
    while (r0 > Query[i].r) remove(r0--);
    q_ans[Query[i].idx] = ans;
for (int i = 0; i < m; ++i) {</pre>
    cout << q_ans[i] << '\n';
```

#### 4.3 Intervals

## 4.3.1 Prefix Sum (L, R) intervals

#### Prefix Sum (L, R) intervals

```
// NOTE: works fine with small n or with large memory
int main() {
   int n, k;
   cin >> n >> k;

   vector<int> a(n + 1);
   vector<vector<int>> rangesPrefix(n + 1, vector<int>(n + 1, 0));
   for (int i = 1; i <= n; ++i)</pre>
```

```
cin >> a[i];
int 1 = 1, r = 1, sum = 0;
// validate your intervals
// here the intervals are the ones that have a sum of k
while (r \le n) {
    sum += a[r];
    while (sum > k) {
       sum -= a[1];
       ++1;
    }
    while (1 \le r \&\& a[1] == 0) {
       if (sum != k)
           break;
       rangesPrefix[r][l]++;
       ++1;
    }
    if (sum == k) {
       rangesPrefix[r][1]++;
    }
    ++r;
}
// prefix sum the columns
for (int i = 1; i <= n; ++i) {</pre>
   for (int j = n - 1; j \ge 0; ---j) {
       rangesPrefix[i][j] += rangesPrefix[i][j + 1];
}
// prefix sum the rows
for (int i = 0; i <= n; ++i) {</pre>
    for (int j = 1; j \le n; ++j) {
       rangesPrefix[j][i] += rangesPrefix[j - 1][i];
   }
}
```

```
int q; cin >> q;

while (q--) {
    cin >> 1 >> r;
    // answer the number of intervals (X, Y) X <= Y that are included
        between L, R
    cout << rangesPrefix[r][l] - rangesPrefix[l - 1][l] << el;
}</pre>
```

4.3.2 Find subarrays intervals that sum to K Using Map

Find subarray intervals that sum to K Using Map

```
int n, k;
cin >> n >> k;
vector < int > a(n + 1);
vector<pair<int, int>> rng;
for (int i = 1; i <= n; ++i)</pre>
   cin >> a[i];
map<int, set<int>> prev;
int currSum = 0;
for (int i = 1; i <= n; ++i) {</pre>
    currSum += a[i];
   if (currSum == k) {
       rng.push_back({1, i});
   if (prev.find(currSum - k) != prev.end()) {
       for (auto &j : prev[currSum - k]) {
           rng.push_back(\{j + 1, i\});
   }
   prev[currSum].insert(i);
}
```

#### 4.4 Ad-hoc

#### 4.4.1 Find duplicate

## Find duplicate using XOR

```
int findDuplicate(int arr[] , int n)
{
   int answer=0;
    //XOR all the elements with 0
   for(int i=0; i<n; i++){
     answer=answer^arr[i];
   }
   //XOR all the elements with no from 1 to n
   // i.e answer^0 = answer
   for(int i=1; i<n; i++){
     answer=answer^i;
   }
   return answer;
}</pre>
```

# 4.5 Sorting Algorithms

#### 4.5.1 Radix Sort

#### Radix Sort

```
// O(n + b), where n is the number of elements and b is the base of the
   number system

// A function to do counting sort of arr[] according to the digit
   represented by exp.

void countingSort(vector<int>& arr, int exp) {
   int n = arr.size();
   vector<int> output(n); // output array
   int count[10] = {0};

   // Store count of occurrences in count[]
   for (int i = 0; i < n; i++)
        count[(arr[i] / exp) % 10]++;

   // Change count[i] so that count[i] now contains the actual
   // position of this digit in output[]
   for (int i = 1; i < 10; i++)</pre>
```

```
count[i] += count[i - 1];
   // Build the output array
   for (int i = n - 1; i \ge 0; i--) {
       output[count[(arr[i] / exp) % 10] - 1] = arr[i];
       count[(arr[i] / exp) % 10]--;
   }
   // Copy the output array to arr[], so that arr now
   // contains sorted numbers according to the current digit
   for (int i = 0; i < n; i++)</pre>
       arr[i] = output[i];
// The main function to that sorts arr[] of size n using Radix Sort
void radixSort(vector<int>& arr) {
   // Find the maximum number to know the number of digits
   int mx = *max_element(arr.begin(), arr.end());
   // Do counting sort for every digit. Note that instead
   // of passing the digit number, exp is passed. exp is 10^i
   // where i is the current digit number
   for (int exp = 1; m / exp > 0; exp *= 10)
       countingSort(arr, exp);
```

#### 4.5.2 Counting Sort

#### Counting Sort

```
// O(N+M), where N and M are the size of inputArray[] and countArray[]
// The main function that sorts arr[] of size n using Counting Sort
void countingSort(vector<int>& arr) {
   int maxElement = *max_element(arr.begin(), arr.end());
   int minElement = *min_element(arr.begin(), arr.end());
   int range = maxElement - minElement + 1;

   vector<int> count(range), output(arr.size());
   for (int i = 0; i < arr.size(); i++)
        count[arr[i] - minElement]++;

   for (int i = 1; i < count.size(); i++)
        count[i] += count[i - 1];</pre>
```

```
for (int i = arr.size() - 1; i >= 0; i--) {
    output[count[arr[i] - minElement] - 1] = arr[i];
    count[arr[i] - minElement]--;
}

for (int i = 0; i < arr.size(); i++)
    arr[i] = output[i];
}</pre>
```

# 4.6 Apply permutation k times

#### $permutation \ k \ times$

```
// Applying a permutation k times
// n log k
vector<int> applyPermutation(vector<int> sequence, vector<int> permutation
   ) {
   vector<int> newSequence(sequence.size());
   for(int i = 0; i < sequence.size(); i++) {</pre>
       newSequence[i] = sequence[permutation[i]];
   }
   return newSequence;
}
vector<int> permute(vector<int> sequence, vector<int> permutation, long
   long k) {
   while (k > 0) {
       if (k & 1) {
           sequence = applyPermutation(sequence, permutation);
       permutation = applyPermutation(permutation, permutation);
       k >>= 1;
   return sequence;
}
```

## 5 Data Structures

#### 5.1 Monotonuc

#### 5.1.1 Min Aggregation Deque

## Min Aggregation Deque

```
struct minstack {
   stack<pair<int, int>> st;
   int getmin() { return st.top().second; }
   bool empty() { return st.empty(); }
   int size() { return st.size(); }
   void push(int x) {
       int mn = x;
       if (!empty()) mn = min(mn, getmin());
       st.push({x, mn});
   void pop() { st.pop(); }
   int top() { return st.top().first; }
   void swap(minstack &x) { st.swap(x.st); }
struct mindeque {
   minstack l, r, t;
   void rebalance() {
       bool f = false;
       if (r.empty()) {
          f = true;
          1.swap(r);
       int sz = r.size() / 2;
       while (sz--) {
          t.push(r.top());
          r.pop();
       }
       while (!r.empty()) {
          1.push(r.top());
          r.pop();
       }
       while (!t.empty()) {
          r.push(t.top());
          t.pop();
```

```
if (f) 1.swap(r);
   }
   int getmin() {
       if (l.empty()) return r.getmin();
       if (r.empty()) return l.getmin();
       return min(l.getmin(), r.getmin());
   }
   bool empty() { return l.empty() && r.empty(); }
   int size() { return l.size() + r.size(); }
   void push_front(int x) { l.push(x); }
   void push_back(int x) { r.push(x); }
   void pop_front() {
       if (1.empty()) rebalance();
       1.pop();
   void pop_back() {
       if (r.empty()) rebalance();
       r.pop();
   }
   int front() {
       if (l.empty()) rebalance();
       return 1.top();
   }
   int back() {
       if (r.empty()) rebalance();
       return r.top();
   }
   void swap(mindeque &x) {
       1.swap(x.1);
      r.swap(x.r);
};
```

## 5.2 Strings

#### 5.2.1 Trie (Prefix Tree)

#### Basic Implementation

```
#define MAX_CHAR 26
struct TrieNode {
```

```
TrieNode *pTrieNode[MAX_CHAR]{};
   bool isWord;
   TrieNode() {
       isWord = false;
       fill(pTrieNode, pTrieNode + 26, (TrieNode *) NULL);
   }
   virtual ~TrieNode() = default;
};
class Trie {
private:
   TrieNode *root;
public:
   Trie() {
       root = new TrieNode();
   }
   virtual ~Trie() = default;
   TrieNode *getTrieNode() {
       return this->root;
   }
   void insert(const string &word) {
       TrieNode *current = root;
       for (char c: word) {
           int i = c - 'a';
           if (current->pTrieNode[i] == nullptr)
              current->pTrieNode[i] = new TrieNode();
           current = current->pTrieNode[i];
       current->isWord = true;
   }
   bool search(const string &word) {
       TrieNode *current = root;
       int ch = 0;
       for (char c: word) {
           ch = c - 'a';
          if (current->pTrieNode[ch] == nullptr)
              return false:
           current = current->pTrieNode[ch];
```

```
return current->isWord;
}

bool startsWith(const string &prefix) {
    TrieNode *current = root;
    int ch = 0;
    for (char c: prefix) {
        ch = c - 'a';
        if (current->pTrieNode[ch] == nullptr)
            return false;
        current = current->pTrieNode[ch];
    }
    return true;
}

};
```

# 5.3 Range Queries

## 5.3.1 Segment Tree

#### Basic Implementation

```
struct Node {
   long long val;
};

struct SegTree {
   private:
        const Node NEUTRAL = {INT_MIN};

        static Node merge(const Node& x1, const Node& x2) {
            return {x1.val + x2.val};
        }

        void set(const int& idx, const int& val, int x, int lx, int rx) {
        if (rx - lx == 1) return void(values[x].val = val);

        int mid = (rx + lx) / 2;

        if (idx < mid)
            set(idx, val, 2 * x + 1, lx, mid);
        return void(values);
        int mid = (rx + lx) / 2;
        if (idx < mid)
            set(idx, val, 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid);
        int mid = (rx + lx) / 2 * x + 1, lx, mid * x + 1, lx,
```

```
else
       set(idx, val, 2 * x + 2, mid, rx);
    values[x] = merge(values[2 * x + 1], values[2 * x + 2]);
}
Node query(const int& 1, const int& r, int x, int lx, int rx) {
   if (lx >= r || l >= rx) return NEUTRAL;
   if (lx >= 1 && rx <= r) return values[x];</pre>
   int mid = (rx + lx) / 2;
   return merge(query(1, r, 2 * x + 1, 1x, mid), query(1, r, 2 * x + 1
        2, mid, rx));
void build(vector<int> &a, int x, int lx, int rx) {
   if (rx - 1x == 1) {
       if (lx < a.size()) {</pre>
           values[x].val = a[lx];
       }
       return;
   }
   int m = (1x + rx) / 2:
   build(a, 2 * x + 1, 1x, m);
   build(a, 2 * x + 2, m, rx);
   values[x] = merge(values[2 * x + 1], values[2 * x + 2]);
}
 void assign_range(int 1, int r, int node, int lx, int rx, int time,
     int val) {
   if (lx > r \mid | l > rx) return;
   if (lx >= 1 && rx <= r) {
       lazy[node] = {time, val};
       return;
   int mid = (lx+rx) / 2;
   assign_range(1, r, 2*node+1, lx, mid, time, val);
   assign_range(l, r, 2*node+2, mid+1, rx, time, val);
}
pair<int, int> point_query(int lx, int rx, int node, int idx) {
   if(rx == lx) return lazy[node];
   int mid = (1x+rx) / 2:
```

```
if(idx <= mid) {</pre>
           auto x = point_query(lx, mid, 2*node+1, idx);
           if(x.first > lazy[node].first) return x;
           return lazy[node];
       }
       auto x = point_query(mid+1, rx, 2*node+2, idx);
       if(x.first > lazy[node].first) return x;
       return lazy[node];
   }
public:
   int size{};
   vector<Node> values;
   void build(vector<int> &a) {
       build(a, 0, 0, size);
   }
   void init(int _size) {
       size = 1;
       while (size < _size) size *= 2;</pre>
       values.assign(2 * size, NEUTRAL);
   }
   void set(int idx, int val) {
       set(idx, val, 0, 0, size);
   }
   Node query(const int& 1, const int& r) {
       return query(1, r, 0, 0, size);
   }
};
```

#### 5.3.2 Lazy Propegation

#### Lazy Propegation

```
struct SegTree {
private:
   void propegate(int lx, int rx, int node) {
      if(!lazy[node]) return;
```

```
if(lx != rx) {
           lazy[2*node+1] = lazy[node];
           lazy[2*node+2] = lazy[node];
       values[node] = lazy[node] * (rx - lx + 1);
       lazy[node] = 0;
   }
   // assign val in range [1, r]
   void update_range(int 1, int r, int node, int lx, int rx, int val,
       bool f) {
       propegate(lx, rx, node);
       if (1x > r \mid | 1 > rx) return;
       if (lx >= 1 && rx <= r) {</pre>
           lazy[node] = val;
           propegate(lx, rx, node);
           return;
       }
       int mid = (lx+rx) / 2;
       update_range(1, r, 2*node+1, lx, mid, val, f);
       update_range(l, r, 2*node+2, mid+1, rx, val, f);
       values[node] = values[2*node+1] + values[2*node+2];
   }
   // get sum in range [1, r]
   int range_query(int 1, int r, int lx, int rx, int node) {
       propegate(lx, rx, node);
       if (lx > r || l > rx) return 0;
       if (lx >= 1 && rx <= r) return values[node];</pre>
       int mid = (1x+rx) / 2;
       return range_query(1, r, lx, mid, 2*node+1) + range_query(1, r, mid
           +1, rx, 2*node+2);
   }
public:
   int size{};
   vector<int> values, lazy;
   void init(int _size) {
       size = 1:
       while (size < _size) size *= 2;</pre>
       values.assign(2 * size, 0);
```

```
lazy.assign(2 * size, 0);
}

void update_range(int 1, int r, int v, bool f) {
    update_range(1, r, 0, 0, size-1, v, f);
}

int range_query(int 1, int r) {
    return range_query(1, r, 0, size-1, 0);
}
};
```

#### 5.3.3 Fenwick Tree

#### Fenwick Tree

```
struct Fenwick {
   // One Based
   vector<int> tree;
   explicit Fenwick(int n) {tree.assign(n + 5, {});}
   // Computes the prefix sum from [1, i], O(log(n))
   int query(int i) {
       int res = 0;
       while (i > 0) {
          res += tree[i];
          i \&= (i \& -i);
       }
       return res;
   }
   int query(int 1, int r) {
       return query(r) - query(1-1);
   }
   // Get the value at index i
   int get(int i) {
       return query(i, i);
   }
   // Add 'v' to index 'i', O(\log(n))
   void update(int i, int v) {
```

```
while (i < tree.size()) {
        tree[i] += v;
        i += (i & -i);
    }
}

// Update range, Point query
// To get(k) do prefix sum [1, k] and in insert update_range(i, i, a[i ])

void update_range(int l, int r, int v) {
        update(l, v);
        update(r+1, -v);
    }
};</pre>
```

#### 5.3.4 Fenwick UpdateRange

#### $BIT\ UpdateRange$

```
struct BITUpdateRange {
private:
   int n;
   vector<int> B1, B2;
   void add(vector<int> &b, int idx, int x) {
       while (idx <= n) {</pre>
           b[idx] += x;
           idx += idx & -idx;
       }
   }
   int sum(vector<int> &b, int idx) {
       int total = 0;
       while (idx > 0) {
           total += b[idx];
           idx &= (idx & -idx);
       }
       return total;
   }
   int prefix(int idx) {
       return sum(B1, idx) * idx - sum(B2, idx);
   }
```

```
public:
   explicit BITUpdateRange(int n) : n(n) {
       B1.assign(n + 1, {});
       B2.assign(n + 1, {});
   }
   void update(int 1, int r, int x) {
       add(B1, 1, x);
       add(B1, r + 1, -x);
       add(B2, 1, x * (1 - 1));
       add(B2, r + 1, -x * r);
   }
   int query(int i) {
       return prefix(i) - prefix(i - 1);
   int query(int 1, int r) {
       return prefix(r) - prefix(l - 1);
   }
};
```

#### 5.3.5 2D BIT

#### 2D BIT

```
struct BIT2D {
   int n, m;
   vector<vector<int>> bit;

BIT2D(int n, int m) : n(n), m(m) {
      bit.assign(n + 2, vector<int>(m + 2));
   }

void update(int x, int y, int val) {
   for (; x <= n; x += x & -x) {
      for (int i = y; i <= m; i += i & -i) {
      bit[x][i] += val;
    }
   }
}</pre>
```

```
int prefix(int x, int y) {
       int res = 0;
       for (; x > 0; x &= (x & -x)) {
          for (int i = y; i > 0; i &= ~(i & -i)) {
              res += bit[x][i];
          }
      }
       return res;
   }
   int query(int sx, int sy, int ex, int ey) {
       int ans = 0;
       ans += prefix(ex, ey);
       ans -= prefix(ex, sy - 1);
       ans -= prefix(sx - 1, ey);
       ans += prefix(sx - 1, sy - 1);
       return ans;
   }
};
```

## 5.3.6 Sparse Table

#### Impl with the index

```
// storing the index also
struct SNode {
    int val;
    int index;
};

class SparseTable {
    private:
        vector<vector<SNode>> table;

    function<SNode(const SNode&, const SNode&)> merge;

    static SNode StaticMerge(const SNode& a, const SNode& b) {
        return a.val < b.val ? a : b;
    }

public:
    explicit SparseTable(const vector<int>& arr, const function<SNode(
        const SNode&, const SNode&)>& mergeFunc = StaticMerge) {
```

```
int n = static_cast<int>(arr.size());
       int log_n = static_cast<int>(log2(n)) + 1;
       this->merge = mergeFunc;
       table.resize(n, vector<SNode>(log_n));
       for (int i = 0; i < n; i++) {
           table[i][0] = {arr[i], i};
       for (int j = 1; (1 << j) <= n; j++) {
           for (int i = 0; i + (1 << j) <= n; i++) {
              table[i][j] = mergeFunc(table[i][j - 1], table[i + (1 << (j</pre>
                   - 1))][j - 1]);
          }
   }
   SNode query(int left, int right) {
       int j = static_cast<int>(log2(right - left + 1));
       return merge(table[left][j], table[right - (1 << j) + 1][j]);</pre>
   }
   // query in O(log(n)) if its could't apply to Sparse Table directly
   T query_log(int 1, int r){
     int len = r - 1 + 1;
     T ans:
     for(int i = 0; 1 <= r; i++){</pre>
         if (len & (1 << i)){</pre>
            ans = merge(ans, table[i][1]);
            1+=(1 << i);
        }
     }
};
int main(void) {
   int n:
   cin >> n;
   vector<int> arr(n);
   for (auto& element : arr) cin >> element;
   SparseTable minSt(arr, [](const SNode& a, const SNode& b) -> SNode {
```

```
return a.val < b.val ? a : b;
});

SparseTable maxSt(arr, [](const SNode& a, const SNode& b) -> SNode {
    return a.val > b.val ? a : b;
});
```

#### 5.3.7 Sparse Table (Shorter Version)

#### Impl

```
template < class T>
struct Rmq {
vector<vector<T>> jmp;
   // vector is 0-based indexed
 Rmg(const vector<T>& V) : jmp(1, V) {
 for (int pw = 1, k = 1; pw * 2 <= V.size(); pw *= 2, ++k) {</pre>
   jmp.emplace_back(V.size() - pw * 2 + 1);
  for(int j = 0; j < jmp[k].size(); j++)</pre>
   jmp[k][j] = (jmp[k-1][j] | jmp[k-1][j+pw]);
}
   // query (a, b] inclusive execlusive
T query(int a, int b) {
 // assert(a < b);</pre>
 int dep = 31 - __builtin_clz(b - a);
 return (jmp[dep][a] | jmp[dep][b - (1 << dep)]);</pre>
};
```

#### 5.3.8 Sqrt Decomp

#### Impl

```
#define 00 1e18

const int sz = 2e5 + 5;
const int sq = 500;
```

```
int arr[sz] = {};
int ans[sq] = {};
void pre() {
   fill(ans, ans + sq, 00);
   for(int i = 0; i < sz; i++) {</pre>
       ans[i / sq] = min(ans[i / sq], arr[i]);
}
// can be O(1) with inversable operations, else it's sqrt
void update(int k, int x) {
  arr[k] = x;
  int start = k / sq;
  ans[start] = *min_element(arr + (start * sq), arr + (start + 1) * sq);
}
int query(int 1, int r){
   int ret = 00;
   while(1 <= r) {</pre>
       if(1 \% sq == 0 \&\& 1 + sq <= r) {
          ret = min(ret, ans[1 / sq]);
          1 += sq;
       }else {
           ret = min(ret, arr[1]);
           1++;
       }
   }
   return ret;
```

#### 5.4 Ordered Set

#### Ordered Set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename T>
```

```
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
   tree_order_statistics_node_update>;
template <class T>
using ordered_multiset = tree<T, null_type, CUSTUM_COMPARE, rb_tree_tag,</pre>
   tree_order_statistics_node_update>;
void erase_set(ordered_set &os, int v) {
   // Number of elements less than v
   int rank = os.order_of_key(v);
   auto it = os.find_by_order(rank);
   os.erase(it);
// Returns iterator to 0-th
// largest element in the set
cout << *S.find_by_order(0) << " ";</pre>
// Returns iterator to 2-nd
// largest element in the set
cout << *S.find_by_order(2);</pre>
```

# 5.5 Custom Compare Functions

#### Custom Compare functions

```
template < class T>
struct custom_compare {
   bool operator()(const T& a, const T& b) const {
      if (a == b) return true; // Keep duplicates
      return a > b;
   }
};

//REF: GFG
class CustomComparator {
public:
   CustomComparator(int baseValue) : baseValue_(baseValue) {}
   bool operator()(int a, int b) const {
```

```
// Custom comparison logic involving state
    return (a % baseValue_) < (b % baseValue_);
}

private:
    int baseValue_;
};

// OR through capture by reference (capture clauses)
auto compare = [&](char a, char b) { return localStructure[a] >
    localStructure[b]; };
```

# 6 Counting Principles

#### 6.1 nCr

$$C(n,k) = \frac{n!}{(n-k)!k!} = \frac{n*(n-1)*(n-2)*...*(n-k+1)}{k!}$$

#### 6.1.1 Fast nCr

$$C(n,k) = \frac{n*(n-1)*(n-2)*\dots*(n-k+1)}{1*2*3*\dots*k} = \prod_{i=0}^{k-1} \frac{n-i}{i+1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} \prod_{i=0}^{i} (i != j) \frac{n!}{n!} \prod_{i=0}^{k-1} (i != j) \frac{$$

#### Fast nCr

```
int nCr(const int& n, const int& r) {
   double res = 1;
   for (int i = 1; i <= r; ++i)
      res = res * (n - r + i) / i;
   return (int)(res + 0.01);
}</pre>
```

# 6.1.2 Method 1: Pascal's Triangle (Dynamic Programming) - $\mathcal{O}(n^2)$

```
nCk using dp
// REF: USACO guide
/** @return nCk mod p using dynamic programming */
int binomial(int n, int k, int p) {
// dp[i][j] stores iCj
vector<vector<int>> dp(n + 1, vector<math><int>(k + 1, 0));
// base cases described above
for (int i = 0; i <= n; i++) {
  * i choose 0 is always 1 since there is exactly one way
  * to choose O elements from a set of i elements
  * (don't choose anything)
  */
 dp[i][0] = 1;
  * i choose i is always 1 since there is exactly one way
  * to choose i elements from a set of i elements
  * (choose every element in the set)
 if (i <= k) { dp[i][i] = 1; }</pre>
for (int i = 0; i <= n; i++) {</pre>
 for (int j = 1; j <= min(i, k); j++) {</pre>
return dp[n][k]; // returns nCk modulo p
```

**6.1.3** Method 2: Factorial Definition (Modular Inverses) -  $O(n + \log MOD)$ 

nCk using Modular Inverses

// REF: USACO guide

```
const int MAXN = 1e6;
long long fac[MAXN + 1];
long long inv[MAXN + 1];
/** @return x^n modulo m in O(log p) time. */
long long exp(long long x, long long n, long long m) {
x \% = m; // note: m * m must be less than 2^63 to avoid 11 overflow
long long res = 1;
while (n > 0) {
 if (n % 2 == 1) { res = res * x % m; }
 x = x * x % m;
 n /= 2;
return res;
/** Precomputes n! from 0 to MAXN. */
void factorial(long long p) {
fac[0] = 1;
for (int i = 1; i <= MAXN; i++) { fac[i] = fac[i - 1] * i % p; }
* Precomputes all modular inverse factorials
* from 0 to MAXN in O(n + log p) time
 */
void inverses(long long p) {
inv[MAXN] = exp(fac[MAXN], p - 2, p);
for (int i = MAXN; i >= 1; i--) { inv[i - 1] = inv[i] * i % p; }
}
/** @return nCr mod p */
long long choose(long long n, long long r, long long p) {
return fac[n] * inv[r] % p * inv[n - r] % p;
int main() {
factorial();
inverses();
int n;
 cin >> n:
for (int i = 0; i < n; i++) {</pre>
```

```
int a, b;
cin >> a >> b;
cout << choose(a, b) << '\n';
}</pre>
```

# 7 Graph Theory

# 7.1 Shortest Path algorithms

#### 7.1.1 Dijkstra Algorithm

#### Dijkstra Implementation

```
#define INF (1e18) // for int defined as 11
int n, m;
vector<vector<pair<int, int>>> adj;
vector<int> cost;
vector<int> parent;
void dijkstra(int startNode = 1) {
   priority_queue<pair<11, int>, vector<pair<11, int>>, greater<>> pq;
   cost[startNode] = 0;
   pq.emplace(0, startNode);
   while (!pq.empty()) {
       int u = pq.top().second;
       11 d = pq.top().first;
       pq.pop();
       if (d > cost[u]) continue;
       for (auto &p: adj[u]) {
          int v = p.first;
          int w = p.second;
          if (cost[v] > cost[u] + w) {
              cost[v] = cost[u] + w;
              parent[v] = u;
              pq.emplace(cost[v], v);
```

```
}
}
void run_test_case(int testNum) {
   cin >> n >> m:
   adj.assign(n + 1, {});
   cost.assign(n + 1, INF);
   parent.assign(n + 1, -1);
   while (m--) {
       // Read Edges
   dijkstra();
   if (cost[n] == INF) {
       cout << -1 << el; // not connected {Depends on you use case}</pre>
       return;
   }
   stack<int> ans;
   for (int v = n; v != -1; v = parent[v]) ans.push(v);
   while (!ans.empty()) { // printing the path
       cout << ans.top() << ' ';</pre>
       ans.pop();
   }
   cout << el;</pre>
```

#### 7.1.2 Floyd Warshal Algorithm

## $FloydWarshal\ Implementation$

```
int main() {
   int n, m; cin >> n >> m;
   vector <vector <int>> adj(n + 1, vector <int> (n + 1, 2e9));
   for (int i = 0; i < n; i++) adj[i][i] = 0;
   while(m--) {</pre>
```

```
int u, v, w;
cin >> u >> v >> w;
    adj[u][v] = min(adj[u][v], w);
    adj[v][u] = min(adj[v][u], w);
}

for (int mid = 1; mid <= n; mid++) {
    for (int start = 1; start <= n; start++) {
        for (int end = 1; end <= n; end++) {
            adj[start][end] = min(adj[start][end], adj[start][mid] +
            adj[mid][end]);
        }
    }
}

return 0;
</pre>
```

#### 7.1.3 Bellman Ford Algorithm

#### BellmanFord Implementation

```
vector <vector <pair<int, int>>> &adj
vector <long long> BellmanFord(int src) {
   int n = (int)adj.size();
   vector <long long> dist(n, 2e18);
   dist[src] = 0;
   for (int it = 0; it < n-1; it++) {</pre>
       bool in = false;
       for (int i = 0; i < n; i++) { // iterate on the edges
           for (auto &[j, w] : adj[i]) {
              if (dist[j] > dist[i] + w) {
                  in = true;
                  dist[j] = dist[i] + w;
              }
       }
       if (!in) return dist;
   }
   for (int i = 0; i < n; i++) {
```

```
for (auto &[j, w] : adj[i]) {
    if (dist[j] > dist[i] + w) { //negative cycle
        return vector <long long> (n, -1); // or any flag
    }
}
return dist;
```

# 7.2 Cycle Detection

}

#### 7.2.1 DFS Implementation

#### DFS Implementation

```
// return true with number of nodes in the cycle, either odd cycle or even
bool cycle_detection(unordered_map<int, vector<int>> &graph, int source,
    int par, unordered_map<int,bool> vis, int c){
    if(vis[source]) return true;

    vis[source] = true;

    for(int v: graph[source]){
        if(v != par){
            c++;
            if(dfs(graph,v, source, vis, c)) return true;
        }
    }
    return false;
}
```

#### 7.2.2 Another way for undirected graphs

#### Another way for undirected graphs

```
// this is true only for undirected graphs
bool dfs1(int cur, int par) {
   bool ret = false;
   vis[cur] = true;
```

```
for (auto &i : adj[cur]) {
    if (!vis[i]) ret|=dfs1(i, cur);
    else if (par != i) ret = true;
}
return ret;
```

#### 7.2.3 General Way

#### General Way

```
// general algorithm
vector <bool> cyc;
bool dfs(int cur, int par) {
   bool ret = false;
   vis[cur] = cyc[cur] = true;
   for (auto &i : adj[cur]) {
      if (par == i) continue;
      if (!vis[i]) ret|=dfs(i, cur);
      else if (cyc[i]) ret = true;
   }
   cyc[cur] = false;
   return ret;
}
```

#### 7.2.4 DSU Implementation

#### $DSU\ Implementation$

```
#include <iostream>
#include <vector>

class UnionFind {
  public:
    UnionFind(int n) {
      parent.resize(n);
      rank.resize(n, 0);
      for (int i = 0; i < n; ++i) {
         parent[i] = i;
      }
}</pre>
```

```
int find(int u) {
       if (parent[u] != u) {
           parent[u] = find(parent[u]);
       return parent[u];
   }
   void unionSets(int u, int v) {
       int rootU = find(u);
       int rootV = find(v);
       if (rootU != rootV) {
           if (rank[rootU] > rank[rootV]) {
               parent[rootV] = rootU;
           } else if (rank[rootU] < rank[rootV]) {</pre>
               parent[rootU] = rootV;
           } else {
               parent[rootV] = rootU;
               ++rank[rootU];
           }
       }
   }
private:
   std::vector<int> parent;
    std::vector<int> rank;
};
bool detectCycle(const std::vector<std::pair<int, int>>& edges, int n) {
   UnionFind uf(n);
   for (const auto& edge : edges) {
       int u = edge.first;
       int v = edge.second;
       if (uf.find(u) == uf.find(v)) {
           return true;
       uf.unionSets(u, v);
   }
   return false:
}
```

```
int main() {
    std::vector<std::pair<int, int>> edges = { {0, 1}, {1, 2}, {2, 3}, {3, 0} };
    int n = 4; // Number of vertices

if (detectCycle(edges, n)) {
    std::cout << "Cycle detected" << std::endl;
} else {
    std::cout << "No cycle detected" << std::endl;
}

return 0;
}</pre>
```

# 7.3 Algorithms

#### 7.3.1 Heavy Light Decomposition

#### Basic HLD Impl

```
struct Node {
   int val;
};

const Node nullNode = {0};

const int N = 2e5 + 5, S = 1 << 19;
int n, q;
int val[N];
int sz[N], par[N], dep[N], id[N], top[N];
vector<int> adj[N];

Node st[S];

Node merge(const Node& a, const Node& b) {
   return {a.val + b.val};
}

void update(int idx, Node val) {
   st[idx += n] = val;
```

```
for (idx /= 2; idx; idx /= 2) st[idx] = merge(st[idx * 2], st[idx * 2
       + 1]);
}
Node query(int lo, int hi) {
   Node ra = nullNode, rb = nullNode;
   for (lo += n, hi += n + 1; lo < hi; lo /= 2, hi /= 2) {
       if (lo & 1) ra = merge(ra, st[lo++]);
       if (hi & 1) rb = merge(st[--hi], rb);
   }
   return merge(ra, rb);
}
int dfs_size(const int& node, const int& parent) {
   sz[node] = 1;
   par[node] = parent;
   for (const int& ch : adj[node]) {
       if (ch == parent) continue;
       dep[ch] = dep[node] + 1;
       par[ch] = node;
       sz[node] += dfs_size(ch, node);
   }
   return sz[node];
}
int curId = 0;
void dfs_hld(const int& cur, const int& parent, const int& curTop) {
   id[cur] = curId++;
   top[cur] = curTop;
   update(id[cur], {val[cur]});
   int heavyChild = -1, heavyMax = -1;
   for (const int& ch : adj[cur]) {
       if (ch == parent) continue;
       if (sz[ch] > heavyMax) {
          heavyMax = sz[ch];
          heavyChild = ch;
       }
   }
   if (heavyChild == -1) return;
   dfs_hld(heavyChild, cur, curTop);
```

```
for (int ch : adj[cur]) {
       if (ch == parent || ch == heavyChild) continue;
       dfs_hld(ch, cur, ch);
   }
Node path(int u, int v) {
   Node ans = nullNode;
   while (top[u] != top[v]) {
       if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
       ans = merge(ans, query(id[top[u]], id[u]));
       u = par[top[u]];
   }
   if (dep[u] > dep[v]) swap(u, v);
   ans = merge(ans, query(id[u], id[v]));
   return ans;
void init() {
   for (int i = 0; i < S; i++) st[i] = nullNode;</pre>
   dfs_size(1, 1);
   dfs_hld(1, 1, 1);
int main() {
   cin >> n >> q;
   for (int i = 1; i <= n; i++) cin >> val[i];
   int a, b;
   for (int i = 2; i <= n; i++) {</pre>
       cin >> a >> b;
       adj[a].pb(b);
       adj[b].pb(a);
   }
   init(); // <----- DON'T FORGET TO CALL THIS FUNCTION
   int type;
   while (q--) {
       cin >> type;
       if (type == 1) {
```

```
cin >> a >> b;
    val[a] = b;
    update(id[a], {val[a]});
}
else {
    cin >> a;
    cout << path(1, a).val << el;
}
}</pre>
```

#### 7.3.2 Heavy Light Decomposition with lazy SegTree

#### Basic HLD Impl

```
struct Node {
   int val;
};
const Node nullNode = {0};
const int N = 2e5 + 5, S = 1 << 19;
int n, q;
int val[N];
int sz[N], par[N], dep[N], id[N], top[N];
vector<int> adj[N];
Node st[S];
int lazy[S];
Node merge(const Node& a, const Node& b) {
   return {a.val + b.val};
}
void push(int idx, int 1, int r) {
   if (lazy[idx] != 0) {
       st[idx].val += lazy[idx] * (r - l + 1);
       if (1 != r) {
          lazy[idx * 2] += lazy[idx];
          lazy[idx * 2 + 1] += lazy[idx];
       }
       lazy[idx] = 0;
   }
```

```
void update_range(int lo, int hi, int l, int r, int idx, int value) {
   push(idx, l, r);
   if (lo > r || hi < 1) return;</pre>
   if (lo <= 1 && r <= hi) {</pre>
       lazy[idx] += value;
       push(idx, l, r);
       return;
   }
   int mid = (1 + r) / 2;
   update_range(lo, hi, l, mid, idx * 2, value);
   update_range(lo, hi, mid + 1, r, idx * 2 + 1, value);
   st[idx] = merge(st[idx * 2], st[idx * 2 + 1]);
void update(int idx, Node val) {
   update_range(idx, idx, 0, n - 1, 1, val.val);
void update_range(int lo, int hi, int value) {
   update_range(lo, hi, 0, n - 1, 1, value);
Node query(int lo, int hi, int l, int r, int idx) {
   push(idx, 1, r);
   if (lo > r || hi < 1) return nullNode;</pre>
   if (lo <= 1 && r <= hi) return st[idx];</pre>
   int mid = (1 + r) / 2;
   return merge(query(lo, hi, l, mid, idx * 2), query(lo, hi, mid + 1, r,
        idx * 2 + 1));
Node query(int lo, int hi) {
   return query(lo, hi, 0, n - 1, 1);
int dfs_size(const int& node, const int& parent) {
   sz[node] = 1;
   par[node] = parent;
   for (const int& ch : adj[node]) {
       if (ch == parent) continue;
       dep[ch] = dep[node] + 1;
       par[ch] = node;
```

```
sz[node] += dfs_size(ch, node);
   }
   return sz[node];
}
int curId = 0;
void dfs_hld(const int& cur, const int& parent, const int& curTop) {
   id[cur] = curId++;
   top[cur] = curTop;
   update(id[cur], {val[cur]});
   int heavyChild = -1, heavyMax = -1;
   for (const int& ch : adj[cur]) {
       if (ch == parent) continue;
       if (sz[ch] > heavyMax) {
           heavyMax = sz[ch];
           heavyChild = ch;
       }
   }
   if (heavyChild == -1) return;
   dfs_hld(heavyChild, cur, curTop);
   for (int ch : adj[cur]) {
       if (ch == parent || ch == heavyChild) continue;
       dfs_hld(ch, cur, ch);
}
int get(int u) {
   return query(id[u], id[u]).val;
}
void path(int u, int v, int val) {
   // Node ans = nullNode;
   while (top[u] != top[v]) {
       if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
       // ans = merge(ans, query(id[top[u]], id[u]));
       update_range(id[top[u]], id[u], val);
       u = par[top[u]];
   }
   if (dep[u] > dep[v]) swap(u, v);
```

```
// ans = merge(ans, query(id[u], id[v]));
   update_range(id[u], id[v], val);
   // return ans;
void init() {
   for (int i = 0; i < S; i++) st[i] = nullNode;</pre>
   memset(lazy, 0, sizeof(lazy));
   dfs_size(1, 1);
   dfs_hld(1, 1, 1);
int main(void) {
   cin >> n >> q;
   for (int i = 1; i <= n; i++) val[i] = 0;</pre>
   int a, b;
   for (int i = 2; i <= n; i++) {
       cin >> a >> b;
       adj[a].push_back(b);
       adj[b].push_back(a);
   }
   init(); // <----- DON'T FORGET TO CALL THIS FUNCTION
   int v;
   while (q--) {
       cin >> a >> b >> v;
       path(a, b, v);
   }
   for (int i = 1; i <= n; i++) {
       cout << get(i) << " ";
   }
   cout << el;</pre>
```

#### 7.3.3 LCA functions using Binary Lifting

#### LCA functions using Binary Lifting

```
const int N = 2e5 + 15, M = 23;
int ancestors[N][M], depth[N], parent[N], val[N];
vector<vector<int>> adj;
//int tin[N], tout[N], timer;
```

```
void dfs_LCA(const int &node, const int &par) {
// tin[node] = timer++;
   parent[node] = par;
   ancestors[node][0] = par;
   depth[node] = depth[par] + 1;
   for (int i = 1; i < M; i++) {
       int p = ancestors[node][i - 1];
       ancestors[node][i] = ancestors[p][i - 1];
   }
   for (const int &v: adj[node]) {
       if (v == par) continue;
       dfs_LCA(v, node);
   }
     tout[node] = timer++;
//bool is_ancestor(int u, int v) {
     return tin[u] <= tin[v] && tout[u] >= tout[v];
//}
int findKth(int u, int k) {
   if (depth[u] <= k) return -1;</pre>
   for (int i = M - 1; i >= 0; i--) {
       if (k & (1 << i)) {
           u = ancestors[u][i];
       }
   }
   return u;
}
int getLCA(int u, int v) {
   if (depth[u] < depth[v])</pre>
       swap(u, v);
   u = findKth(u, depth[u] - depth[v]);
   if (u == v) return u;
   for (int i = M - 1; i >= 0; i--) {
       if (ancestors[u][i] == ancestors[v][i]) continue;
       u = ancestors[u][i]:
       v = ancestors[v][i]:
```

```
return ancestors[u][0];
int getDistance(int u, int v) {
   int lca = getLCA(u, v);
   return (depth[u] + depth[v]) - (2 * depth[lca]);
int dfs_accumulate(const int &node, const int &par) {
   for (const int& ch: adj[node]) {
       if (ch == par) continue;
       val[node] += dfs_accumulate(ch, node);
   }
   return val[node];
void applyOpOnPath(const int a, const int b, const int w) {
   // adding w to each node on the path a to b
   val[a] += w;
   val[b] += w;
   int lca = getLCA(a, b);
   val[lca] -= w;
   val[parent[lca]] -= w;
int main(void) {
   int n, q;
   cin >> n >> q;
   adj.resize(n + 1);
   int u, v;
   for (int i = 2; i <= n; ++i) {
       cin >> u >> v;
       adj[u].push_back(v);
       adj[v].push_back(u);
   }
   dfs_LCA(1, 1);
   parent[1] = -1;
   int w:
   for (int i = 0; i < q; ++i) {</pre>
```

```
cin >> u >> v;
    cout << getDistance(u, v) << el;
}

// dfs_accumulate(1, 0);

// for (int i = 1; i <= n; i++) {
    cout << val[i] << " ";
}
// cout << el;
}</pre>
```

#### 7.3.4 Topological Sort

#### Topological Sort Using DFS

```
//REF: USACO Guide
vector<int> top_sort;
vector<vector<int>> graph;
vector<bool> visited;
void dfs(int node) {
for (int next : graph[node]) {
 if (!visited[next]) {
  visited[next] = true;
  dfs(next);
 }
top_sort.push_back(node);
int main() {
int n, m; // The number of nodes and edges respectively
std::cin >> n >> m;
 graph = vector<vector<int>>(n);
for (int i = 0; i < m; i++) {</pre>
 int a, b;
 std::cin >> a >> b;
 graph[a - 1].push_back(b - 1);
}
visited = vector<bool>(n);
```

```
for (int i = 0; i < n; i++) {</pre>
 if (!visited[i]) {
  visited[i] = true;
  dfs(i);
 }
std::reverse(top_sort.begin(), top_sort.end());
vector<int> ind(n);
for (int i = 0; i < n; i++) { ind[top_sort[i]] = i; }</pre>
// Check if the topological sort is valid
bool valid = true;
for (int i = 0; i < n; i++) {</pre>
 for (int j : graph[i]) {
  if (ind[j] <= ind[i]) {</pre>
   valid = false;
   goto answer;
 }
answer:;
if (valid) {
 for (int i = 0; i < n - 1; i++) { cout << top_sort[i] + 1 << ', '; }</pre>
 cout << top_sort.back() + 1 << endl;</pre>
} else {
 cout << "IMPOSSIBLE" << endl;</pre>
```

# 8 Techniques

## 8.1 Coordinate Compression

```
void coordinate_compress(vector<int> &x, int start=0, int
    step=1) {
    set unique(x.begin(), x.end());
    map<int, int> valPos;
    int idx=0;
```

```
for (auto i: unique) {
  valPos[i] = start + idx * step;
  ++idx;
}
for(auto &i: x) i = valPos[i];
}
```

Coordinate Compression

# 8.2 Binary to decimal

#### Binary to decimal

```
// Function to convert binary to decimal
// 0(32)
int binaryToDecimal(string str)
{
   int dec_num = 0;
   int power = 0;
   int n = str.length();

   for(int i = n-1; i>=0; i--){
    if(str[i] == '1'){
      dec_num += (1<<power);
   }
   power++;
   }

   return dec_num;
}</pre>
```

# 8.3 Decimal to binary

#### Decimal to bianry

```
// Function that convert Decimal to binary
// 0(32)
void decToBinary(int n)
{
```

```
// Size of an integer is assumed to be 32 bits
   for (int i = 31; i >= 0; i--) {
       int k = n >> i;
       if (k & 1)
           cout << "1";
       else
           cout << "0";
   }
// O(logn)
string DecimalToBinary(int num)
   string str;
     while(num){
     if(num & 1) // 1
       str+='1';
     else // 0
       str+='0';
     num>>=1; // Right Shift by 1
     return str;
```

# 9 Number Theory

#### 9.1 Divisors

#### 9.1.1 formulas

### number of divisors

```
int d(int n){
   unordered_map<int, int> factors = pf(n);
   int c = 1;
   for(const auto& factor: factors){
      c *= (factor.second+1);
   }
   return c;
}
```

```
// range Count Divisors backward thinking MAXN = 2e6
 for(int i=1; i <= n; ++i) {</pre>
   for(int j = i; j <= n; j += i) {</pre>
     numFactors[j]++;
   }
 }
int countDivisors(int n) {
 int count = 0;
 for (int i = 1; i * i <= n; ++i) {
   if (n % i == 0) {
       if (i == n / i) {
             count++; // Perfect square
       } else {
             count += 2; // Pair of divisors
   }
   return count;
```

#### sum of divisors

```
int s(int n){
   unordered_map<int,int> factors = pf(n);
   int sum = 1;
   for(const auto& factor: factors){
      int p = factor.first;
      int exp = factor.second;
      sum *= (pow(p,exp+1)-1)/p-1;
   }
   return sum;
}
```

### 9.2 Primes

## prime factorization

```
void factorize(int x, unordered_map<int, int>& factors) {
   while (x % 2 == 0) {
```

```
factors[2]++;
    x /= 2;
}
for (int i = 3; i * i <= x; i += 2) {
    while (x % i == 0) {
        factors[i]++;
        x /= i;
    }
}
if (x > 2) factors[x]++;
}
```

#### number of co-primes with n

```
int eulerTotient(int n){
   int result = n;
   for(int i = 2; SQ(i) \le n; i++){
       if(n\%i == 0){
           while (n\%i == 0) {
              n/=i;
           result -= result/i;
       }
   }
   if(n > 1) result -= result/n;
   return result;
//Phi(n) = n * (1 - 1/P1) * (1 - 1/P2) * ...
//NOTE: summation of Euler function over divisors of n is equal to n
// using seive
void phi_generator() {
   const int MAX = 1000000;
   char primes[MAX];
   int phi[MAX];
   memset(primes, 1, sizeof(primes));
```

```
for (int k = 0; k < MAX; ++k)
       phi[k] = 1;
   for (int i = 2; i <= MAX; ++i) {</pre>
       if (primes[i]) {
           phi[i] = i - 1; // phi(prime) = p-1
           for (int j = i * 2; j <= MAX; j += i) {</pre>
               primes[j] = 0;
               int n = j, pow = 1;
               while (n % i == 0) {
                   pow *= i;
                   n /= i;
               phi[j] *= (pow / i) * (i - 1);
           }
       }
}
// \text{ phi}(N!) = (N \text{ is prime } ? N-1 : N) * \text{phi}((N-1)!)
ll phi_factn(int n) {
   ll ret = 1;
   for (int i = 2; i <= n; ++i)
       ret = ret * (isprime(i) ? i - 1 : i);
   return ret;
}
```

#### Prime Check

```
vector<bool> isPrime(MAXN, true);

void sieve() {
  isPrime[0] = isPrime[1] = false;

for (int i=2; i * i <= isPrime.size(); ++i) {
   if(isPrime[i]) {
     for (int j = 2 * i; i <= isPrime.size(); j += i)
        prime[j] = false;
   }
}</pre>
```

```
bool Prime(int n) {
 if(n == 2) return true;
 if (n < 2 \mid | n \% 2 == 0) return false;
 for(int i=3; i * i <= n; i += 2) {
   if(n % i == 0) return false;
 return true;
// Generate Primes
const int sz = sqrt(MAXN);
vector<int> prime;
vector<bool> vis(sz);
void pre() {
   prime.push_back(2);
   for (int j = 4; j < sz; j += 2) vis[j] = true;</pre>
   for (int i = 3; i < sz; i += 2) {
       if (vis[i]) continue;
       prime.push_back(i);
       for (int j = i * i; j < sz; j += i) vis[j] = true;</pre>
   }
// Preprocessing Prime Factorization of range numbers
constexpr int N = 5e6+1;
int a[N];
for(int i=2; i < N; ++i) {</pre>
   if(!a[i]) {
       for(int j=1; i*j < N; ++j) {</pre>
           for(int k=i*j; k%i==0; k/=i) a[i*j]++;
   }
   a[i] += a[i-1];
```

#### Math 9.3

## 9.3.1 Vieta's Formula for a Polynomial of Degree n

**Problem**: Given a polynomial of degree n:

$$P(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

with roots  $r_1, r_2, \ldots, r_n$ , express the sums and products of its roots using Vieta's formulas.

Solution: Using Vieta's formulas, we can relate the coefficients of the polynomial to sums and products of its roots:

• Sum of the roots taken one at a time:

$$r_1 + r_2 + \dots + r_n = -\frac{a_{n-1}}{a_n}$$

• Sum of the products of the roots taken two at a time:

$$r_1r_2 + r_1r_3 + \dots + r_{n-1}r_n = \frac{a_{n-2}}{a_n}$$

• Sum of the products of the roots taken three at a time:

$$r_1 r_2 r_3 + r_1 r_2 r_4 + \dots + r_{n-2} r_{n-1} r_n = -\frac{a_{n-3}}{a_n}$$

- Continue this pattern until:
- Product of the roots (for even n):

$$r_1 r_2 \cdots r_n = (-1)^n \frac{a_0}{a_n}$$

#### Example Problem Using Vieta's Formula

**Problem:** Given a quadratic equation  $x^2 + bx + c = 0$  with roots  $r_1$  and  $r_2$ , find  $r_1 + r_2$  and  $r_1 r_2$ .

**Solution**: Using Vieta's formulas for a quadratic equation  $ax^2+bx+c=0$ :

• Sum of the roots:

$$r_1 + r_2 = -\frac{b}{a}$$

• Product of the roots:

$$r_1 r_2 = \frac{\alpha}{\alpha}$$

For the given quadratic  $x^2 + bx + c = 0$  (where a = 1):

Also: To find the roots of the quadratic equation, we can use the discriminant formula,  $D = b^2 - 4ac$ . The roots will then be  $x_1 = \frac{b - \sqrt{D}}{2}$  and  $x_2 = \frac{b+\sqrt{D}}{2}$ .

#### Example Vieta's Formula for Cubic Equation

When considering a cubic equation in the form of  $f(x) = ax^3 + bx^2 + cx + d$ , Vieta's formula states that if the equation f(x) = 0 has roots  $r_1, r_2$ , and  $r_3$ , then:

- $r_1 + r_2 + r_3 = -\frac{b}{a}$   $r_1r_2 + r_2r_3 + r_3r_1 = \frac{c}{a}$   $r_1r_2r_3 = -\frac{d}{a}$

#### Phi Function

- Count integers i < n such that gcd(i, n) = 1
- $gcd(a, b) = 1 \Rightarrow$  then coprimes: gcd(5, 7), gcd(4, 9)
- gcd(prime, i) = 1 for i < prime
- $\varphi(10) = 4 \Rightarrow 1, 3, 7, 9$
- $\varphi(5) = 4 \Rightarrow 1, 2, 3, 4 \dots \varphi(\text{prime}) = \text{prime} 1$

• If a, b, c are pairwise coprimes, then

$$\varphi(a \cdot b \cdot c) = \varphi(a) \cdot \varphi(b) \cdot \varphi(c)$$

• If k > 1

$$\varphi(p^k) = p^k - p^{k-1} = p^{k-1}(p-1) = p^k \left(1 - \frac{1}{p}\right)$$

#### 9.5 Euler's Totient Numbers

## Online Sequence

- $\varphi(1) = \varphi(2) = 1, \ \varphi(5) = 4$
- $\varphi(n)$  is even for n > 2
- $\sqrt{n} \le \varphi(n) \le n \sqrt{n}$ : Except 2, 6
- $\varphi(n^k) = n^{k-1} \cdot \varphi(n)$
- $n = \sum \varphi(d_i)$  where d are the divisors of n

#### 9.6 Möbius Function

#### Möbius Function

```
int mobius(int n) {
  int p = 0;
  // Handling 2 separately
  if (n%2 == 0){
      n = n/2;
      p++;

      // If 2^2 also divides N
      if (n % 2 == 0)
          return 0;
}

// Check for all other prime factors
```

```
for (int i = 3; i <= sqrt(n); i = i+2) {</pre>
       if (n\%i == 0){
           n = n/i;
           p++;
           if (n % i == 0) return 0;
   }
   return (p % 2 == 0)? -1 : 1;
void mobius_generator() const {
  const int MAX = 1000000;
  char prime[MAX + 1];
 for (11 i = 2; i <= MAX; i++)</pre>
   moebius[i] = -1, prime[i] = 1;
 for (ll i = 2; i <= MAX; ++i)
   if (prime[i]) {
     moebius[i] = 1;
     for (11 j = 2 * i; j \le MAX; j += i)
       prime[j] = 0, moebius[j] = j % (i * i) == 0 ? -moebius[j] : 1;
   }
// Mobius Inclusion Exclusion
// Count triples gcd(a, b, c) = 1
int n = 4;
11 sum = n * n * n;
for (ll i = 2; i <= n; ++i)</pre>
   sum -= moebius[i] * (n / i) * (n / i) * (n / i);
```

### 9.7 Möbius and Inclusion Exclusion

Count the triples (a, b, c) such that  $a, b, c \le n$ , and gcd(a, b, c) = 1

• Reverse thinking, total - (# triples gcd > 1)

- How many triples with gcd multiple of 2:  $(n/2)^3$
- How many triples with gcd multiple of 3:  $(n/3)^3$
- and 4? Ignore any numbers of internal duplicate primes
- and 6? already computed in 2, 3. Remove it:  $-(n/6)^3$

#### 9.8 Totient and Möbius Connection

Sum over divisors d of n

$$\sum_{d} d\mu \left(\frac{n}{d}\right) = \varphi(n)$$

# 9.9 Lagrange's four-square theorem

Lagrange's four-square theorem states that, every positive integer can be expressed as the sum of the squares of four integers.

#### Lagrange's four-square theorem

```
#include <cmath>
bool isPerfectSquare(int n) {
   int rt = sqrt(n);
   return rt * rt == n;
}
// Function to compute the minimum number of perfect squares
int numSquares(int n) {
   // Case 1:
   if (isPerfectSquare(n)) return 1;
   // Case 2: Check if n can be expressed as the sum of two perfect
       squares
   for (int i = 1; i * i <= n; i++) {
       int d = n - i * i;
       if (isPerfectSquare(d)) return 2;
   }
   // Case 3:
```

# 10 Geometry

# 10.1 Linearity

#### 10.1.1 co-linear points

#### check if two points are co-linear

```
bool co_linear(int x1, int y1, int x2, int y2, int x3, int y3){
  int area = x1*(y2-y3) + x2*(y3-y1) + x3*(y1-y2);
  return area == 0;
}
```

## 10.2 Polygons

#### 10.2.1 Polygon formation

### check if can form polygon with given angle

```
bool possible(double angle){
  if(angle <= 0 || angle >= 180) return false;

double sides = 360.0/(180.0-angle);

return (sides == static_cast<int>(sides) && sides >= 3);
}
```

#### 10.2.2 Polygon Area

#### area of any polygon with x vertices

```
double shoelace(vector<pair<double, double>>> points) {
  double leftSum = 0.0;
  double rightSum = 0.0;

for (int i = 0; i < points.size(); ++i) {
  int j = (i + 1) % points.size();
  leftSum += points[i].first * points[j].second;
  rightSum += points[j].first * points[i].second;
}

return 0.5 * abs(leftSum - rightSum);
}</pre>
```

#### 10.3 Intersections

#### 10.3.1 Rectangle

#### intersection area between 2 rectangles

```
struct Rectangle {
    int x1, y1; // Bottom-left corner
    int x2, y2; // Top-right corner
};

int intersectionArea(const Rectangle& rect1, const Rectangle& rect2){
    int x_left = max(rect1.x1, rect2.x1);
    int y_bottom = max(rect1.y1, rect2.y1);
    int x_right = min(rect1.x2, rect2.x2);
    int y_top = min(rect1.y2, rect2.y2);

    int intersection_width = x_right - x_left;
    int intersection_height = y_top - y_bottom;

if (intersection_width > 0 && intersection_height > 0) {
        return intersection_width * intersection_height;
    }

    return 0;
```

```
10.3.2 Circle
```

#### intersection area between 2 circles

```
double area(int x0, int y0, int r0, int x1, int y1, int r1){
   const double PI = 3.14159265358979323846;
   double rr0 = r0 * r0;
   double rr1 = r1 * r1;
   double d = sqrt((x1 - x0) * (x1 - x0) + (y1 - y0) * (y1 - y0));
   if(d >= r0 + r1){
      return 0;
   }
   double phiAngle = (rr0 + (d * d) - rr1) / (2 * r0 * d);
   double phi = acos(phiAngle) * 2;
   double thetaAngle = (rr1 + (d * d) - rr0) / (2 * r1 * d);
   double theta = acos(thetaAngle) * 2;
   double area1 = 0.5 * theta * rr1 - 0.5 * rr1 * sin(theta);
   double area2 = 0.5 * phi * rr0 - 0.5 * rr0 * sin(phi);
   return area1+area2;
}
```

#### 10.3.3 Triangle

## intersection area between 2 triangles

```
struct Point {
    double x, y;
};

typedef vector<Point> Polygon;

bool inside(const Point &p, const pair<Point, Point> &edge) {
    Point a = edge.first, b = edge.second;
    return (b.x - a.x) * (p.y - a.y) - (b.y - a.y) * (p.x - a.x) >= 0;
}

Point compute_intersection(const Point &p1, const Point &p2, const pair
    Point, Point> &edge) {
    Point a = edge.first, b = edge.second;
    double A1 = b.y - a.y, B1 = a.x - b.x, C1 = b.x * a.y - a.x * b.y;
```

```
double A2 = p2.y - p1.y, B2 = p1.x - p2.x, C2 = p2.x * p1.y - p1.x *
       p2.y;
   double det = A1 * B2 - A2 * B1:
   if (det == 0) return {numeric_limits<double>::quiet_NaN(),
       numeric_limits<double>::quiet_NaN()); // parallel
   double x = (B1 * C2 - B2 * C1) / det:
   double y = (A2 * C1 - A1 * C2) / det;
   return {x, y};
}
Polygon HodgmanClip(const Polygon &subjectPolygon, const Polygon &
    clipPolygon) {
   Polygon outputList = subjectPolygon;
   for (size_t i = 0; i < clipPolygon.size(); ++i) {</pre>
       Point a = clipPolygon[i];
       Point b = clipPolygon[(i + 1) % clipPolygon.size()];
       Polygon inputList = outputList;
       outputList.clear();
       if (inputList.empty()) break;
       Point s = inputList.back();
       for (const auto &e : inputList) {
           if (inside(e, {a, b})) {
              if (!inside(s, {a, b})) {
                  Point intersection = compute_intersection(s, e, {a, b});
                  if (!std::isnan(intersection.x)) {
                      outputList.push_back(intersection);
                  }
              }
              outputList.push_back(e);
           } else if (inside(s, {a, b})) {
              Point intersection = compute_intersection(s, e, {a, b});
              if (!std::isnan(intersection.x)) {
                  outputList.push_back(intersection);
              }
           }
           s = e;
   }
   return outputList;
}
// area of polygon
double shoelace(const Polygon &vertices) {
   double area = 0:
```

```
for (size_t i = 0; i < vertices.size(); ++i) {
    Point p1 = vertices[i];
    Point p2 = vertices[(i + 1) % vertices.size()];
    area += p1.x * p2.y - p2.x * p1.y;
}
return fabs(area) / 2.0;</pre>
```

#### 10.3.4 Rectangle & Circle

intersection area between Rectangle and Circle

```
struct Point {
   double x, y;
};
struct Circle {
   Point center;
   double radius;
};
struct Rectangle {
   Point bottomLeft, topRight;
};
const double PI = acos(-1.0);
bool point_inside_circle(const Point &p, const Circle &c) {
   double dx = p.x - c.center.x;
   double dy = p.y - c.center.y;
   return (dx * dx + dy * dy) <= (c.radius * c.radius);</pre>
double rect_area(const Rectangle &rect) {
   return (rect.topRight.x - rect.bottomLeft.x) * (rect.topRight.y - rect
       .bottomLeft.y);
bool rect_in_circle(const Circle &circle, const Rectangle &rect) {
   vector<Point> corners = {
       rect.bottomLeft,
```

```
{rect.topRight.x, rect.bottomLeft.y},
       rect.topRight,
       {rect.bottomLeft.x, rect.topRight.y}
   };
   for (const auto &corner : corners) {
       if (!isPointInsideCircle(corner, circle)) {
          return false:
       }
   }
   return true;
}
bool circle_in_rect(const Circle &circle, const Rectangle &rect) {
return circle.center.x - circle.radius >= rect.bottomLeft.x &&
          circle.center.x + circle.radius <= rect.topRight.x &&</pre>
          circle.center.y - circle.radius >= rect.bottomLeft.y &&
          circle.center.y + circle.radius <= rect.topRight.y;</pre>
}
double intersection(const Circle &circle, const Rectangle &rect) {
   if (rect_in_circle(circle, rect)) {
       return rect_area(rect);
   }
   if (circle_in_rect(circle, rect)) {
       return PI * circle.radius * circle.radius;
   }
   double intersectionArea = 0.0;
   double dx1 = max(rect.bottomLeft.x, circle.center.x - circle.radius);
   double dx2 = min(rect.topRight.x, circle.center.x + circle.radius);
   double dy1 = max(rect.bottomLeft.y, circle.center.y - circle.radius);
   double dy2 = min(rect.topRight.y, circle.center.y + circle.radius);
   for (double x = dx1; x < dx2; x += 0.001) {
       for (double y = dy1; y < dy2; y += 0.001) {
          Point p = \{x, y\};
          if (point_inside_circle(p, circle)) {
              intersectionArea += 0.001 * 0.001;
          }
       }
   }
```

```
return intersectionArea;
}
```

#### 10.3.5 Line & Circle

#### intersection points between Line and Circle

```
struct Point {
   double x, y;
};
struct Circle {
   Point center;
   double radius;
};
struct Line {
   double slope;
   double intercept;
};
vector<Point> intersect_points(const Circle &circle, const Line &line) {
   vector<Point> intersections;
   // solve y = mx + b with (x - h)^2 + (y - k)^2 = r^2
   double h = circle.center.x;
   double k = circle.center.y;
   double r = circle.radius;
   double m = line.slope;
   double b = line.intercept;
   // Quadratic coefficients
   double A = 1 + m * m;
   double B = 2 * (m * b - m * k - h);
   double C = k * k - r * r + h * h - 2 * b * k + b * b;
   // Discriminant
   double discriminant = B * B - 4 * A * C;
   if (discriminant < 0) {</pre>
       // No intersection
       return intersections:
   } else if (discriminant == 0) {
```

```
// One intersection (tangent line)
       double x = -B / (2 * A);
       double y = m * x + b;
       intersections.push_back({x, y});
   } else {
       // Two intersections
       double sqrtDiscriminant = sqrt(discriminant);
       double x1 = (-B + sqrtDiscriminant) / (2 * A);
       double y1 = m * x1 + b;
       double x2 = (-B - sqrtDiscriminant) / (2 * A);
       double y2 = m * x2 + b;
       intersections.push_back({x1, y1});
       intersections.push_back({x2, y2});
   }
   return intersections;
}
```

## 10.4 Linear Operations

#### 10.5 2D Points

#### Basic Operations

```
#include<bits/stdc++.h>
using namespace std;

#define ftype double

struct Point2D {
   ftype x, y;
   Point2D() {}
   Point2D(ftype x, ftype y): x(x), y(y) {}
   Point2D& operator+=(const Point2D &p2) {
        x += p2.x;
        y += p2.y;
        return *this;
   }
   Point2D& operator-=(const Point2D &p2) {
        x -= p2.x;
        y -= p2.y;
        return *this;
   }
```

```
return *this;
   }
   Point2D& operator*=(const Point2D p2) {
       x *= p2.x;
       y *= p2.y;
       return *this;
   Point2D& operator/=(const Point2D p2) {
       x \neq p2.x;
       y /= p2.y;
       return *this;
   }
   Point2D& operator+(const Point2D p2) {
       return Point2D(*this) += p2;
   Point2D& operator-(const Point2D p2) {
       return Point2D(*this) -= p2;
   Point2D& operator/(const Point2D p2) {
       return Point2D(*this) /= p2;
   Point2D& operator*(const Point2D p2) {
       return Point2D(*this) *= p2;
   }
};
Point2D operator*(ftype a, Point2D b) {
       return a * b;
int main() {
   Point2D p1(1, 2);
   Point2D p2(3, 4);
   Point2D p3 = p1 * p2;
```

#### 10.6 3D Points

#### Basic Operations

```
#include<bits/stdc++.h>
#define ftype double
struct Point3D {
   ftype x, y, z;
   Point3D() {}
   Point3D(ftype x, ftype y, ftype z): x(x), y(y), z(z) {}
   Point3D& operator+=(const Point3D &p2) {
       x += p2.x;
       y += p2.y;
       z += p2.z;
       return *this;
   Point3D& operator-=(const Point3D &p2) {
       x = p2.x;
       y = p2.y;
       z = p2.z;
       return *this;
   }
   Point3D& operator*=(const Point3D p2) {
       x *= p2.x;
       y *= p2.y;
       z *= p2.z;
       return *this;
   Point3D& operator/=(const Point3D p2) {
       x \neq p2.x;
       y /= p2.y;
       z /= p2.z;
       return *this;
   }
   Point3D& operator+(const Point3D p2) {
       return Point3D(*this) += p2;
   }
   Point3D& operator-(const Point3D p2) {
       return Point3D(*this) -= p2;
   Point3D& operator/(const Point3D p2) {
```

```
return Point3D(*this) /= p2;
}
Point3D& operator*(const Point3D p2) {
    return Point3D(*this) *= p2;
}
};

Point3D operator*(ftype a, Point3D b) {
    return a * b;
}

int main() {
    Point3D p1(1, 2);
    Point3D p2(3, 4);

    Point3D p3 = p1 * p2;
}
```

## 11 Miscellaneous

## 11.1 Faster implementations

#### 11.1.1 hashes

#### custom hash

```
#define safe hash unordered_map<type, type, custom_hash> // same for
    gp_hash_table
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        // http://xorshift.di.unimi.it/splitmix64.c
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
    size_t operator()(uint64_t x) const {
```

```
rwx (Alexandria University)
```

## gb hash table

```
//policy based ds (faster hash table)
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
gp_hash_table<int, int> table;
```

## 11.1.2 Binary Search the value

#### $nearest\ sqrt$

```
long long my_sqrt(long long a)
{
    long long l=0,r=5000000001;
    while(r-1>1)
    {
        long long mid=(1+r)/2;
        if(111*mid*mid<=a)1=mid;
        else r=mid;
    }
    return 1;
}</pre>
```