Competitive Programming Library

rwx 2025/2026

Contents					Bit 3.1	Manipulation Subset Operations	10 10
1 1 1 1 1 1 1 2 E 2	1.1.1 1.1.2 2 MOD 3 Macro 4 Grid I 5 Intege 6 Matri: 0ynamic .1 Some	IO Manipulation GCC Compiler Optimization (Vectorization) Template Solution Navigation T 128 X Expo Programming dp patterns lutions Max Subarray sum (Kadane's Algorithm) Maximum Subarray Alternating Sum Count number of DISTINCT ordered ways to produce coins sums to x Min absolute difference between 2 elements from (L, R) (DP Ranges) Longest common subsequence between 2 Strings	3 3 3 4 4 5 5 6 6 7 7 7 8 8	4	Alg 4.1 4.2 4.3 4.4 4.5 4.6 4.7	FFT MO Intervals 4.3.1 Prefix Sum (L, R) intervals 4.3.2 Find subarrays intervals that sum to K Using Map Ad-hoc 4.4.1 Find duplicate Sorting Algorithms 4.5.1 Radix Sort 4.5.2 Counting Sort Apply permutation k times Strings 4.7.1 Strings Hashing Utils 4.7.2 KMP 4.7.3 Manachers Algorithm Ea Structures Monotonuc 5.1.1 Min Aggregation Deque	11 11 12 13 13 14 14 14 15 15 16 16 16 17 18 18
	2.2.6 2.2.7	Longest common subsequence $\mathcal{O}(n^2)$ Longest common subsequence Binary Search $\mathcal{O}(n+\log N)$	9 10		5.25.3	Strings	19 20 20

		5.3.3 Fenwick Tree	22	8	Techniques	39
		5.3.4 Fenwick UpdateRange	22		8.1 Coordinate Compression	39
		5.3.5 2D BIT	23		8.2 Binary to decimal	39
		5.3.6 Sparse Table	24		8.3 Decimal to binary	39
		5.3.7 Sparse Table (Shorter Version)	24			
		5.3.8 Sqrt Decomp	25	9	Number Theory	4 0
	5.4	Ordered Set	25		9.1 Divisors	
	5.5	Custom Compare Functions	26		9.1.1 formulas	
	5.6	Median Finder	26		9.2 Primes	
	5.7	Binary Trie	27		9.3 Math	
					9.3.1 Vieta's Formula for a Polynomial of Degree $n \ldots \ldots$	
6		inting Principles	28		9.4 Phi Function	
	6.1	nCr	28		9.5 Euler's Totient Numbers	
		6.1.1 Fast nCr	28		9.6 Möbius Function	
		6.1.2 Method 1: Pascal's Triangle (Dynamic Programming)			9.7 Möbius and Inclusion Exclusion	
		- $\mathcal{O}(n^2)$	28		9.8 Totient and Möbius Connection	
		6.1.3 Method 2: Factorial Definition (Modular Inverses) -			9.9 Lagrange's four-square theorem	44
		$\mathcal{O}(n + \log MOD)$	28			
		- (- 1	10	0.00	4 -
7	C	,		10	Geometry 10.1 Linearity	45
7		ph Theory	29	10	10.1 Linearity	45
7	Gra 7.1	aph Theory Shortest Path algorithms	29 29	10	10.1 Linearity	$45 \\ 45$
7		hph Theory Shortest Path algorithms	29 29 29	10	10.1 Linearity	45 45 45
7		Shortest Path algorithms	29 29 29 30	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation	45 45 45 45
7	7.1	Shortest Path algorithms	29 29 29 30 30	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area	45 45 45 45 45
7		Shortest Path algorithms	29 29 29 30 30 31	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections	45 45 45 45 45 45
7	7.1	Shortest Path algorithms	29 29 29 30 30 31 31	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3.1 Rectangle	45 45 45 45 45 45
7	7.1	Shortest Path algorithms	29 29 29 30 30 31 31 31	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle	45 45 45 45 45 45 45
7	7.1	Shortest Path algorithms	29 29 29 30 30 31 31 31	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle 10.3.3 Triangle 10.3.3 Triangle	45 45 45 45 45 45 45 45 45
7	7.1	Shortest Path algorithms 7.1.1 Dijkstra Algorithm 7.1.2 Floyd Warshal Algorithm 7.1.3 Bellman Ford Algorithm Cycle Detection 7.2.1 DFS Implementation 7.2.2 Another way for undirected graphs 7.2.3 General Way 7.2.4 DSU Implementation	29 29 29 30 30 31 31 31 31	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle 10.3.3 Triangle 10.3.4 Rectangle & Circle	45 45 45 45 45 45 45 46 47
7	7.1	Shortest Path algorithms	29 29 29 30 30 31 31 31 31 31 32	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle 10.3.3 Triangle 10.3.4 Rectangle & Circle 10.3.5 Line & Circle 10.3.5 Line & Circle	45 45 45 45 45 45 45 46 47 48
7	7.1	Shortest Path algorithms 7.1.1 Dijkstra Algorithm 7.1.2 Floyd Warshal Algorithm 7.1.3 Bellman Ford Algorithm Cycle Detection 7.2.1 DFS Implementation 7.2.2 Another way for undirected graphs 7.2.3 General Way 7.2.4 DSU Implementation Algorithms 7.3.1 Heavy Light Decomposition	29 29 29 30 30 31 31 31 31 32 32	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle 10.3.2 Circle 10.3.4 Rectangle & Circle 10.3.4 Rectangle & Circle 10.3.5 Line & Circle 10.3.5 Line & Circle 10.4 Linear Operations 10.3 Linear Operations	45 45 45 45 45 45 45 46 47 48
7	7.1	Shortest Path algorithms	29 29 30 30 31 31 31 31 32 32 34	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle 10.3.3 Triangle 10.3.4 Rectangle & Circle 10.3.5 Line & Circle 10.4 Linear Operations 10.5 2D Points	45 45 45 45 45 45 45 46 47 48 48 48
7	7.1	Shortest Path algorithms	29 29 29 30 30 31 31 31 31 32 32 34 36	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle 10.3.2 Circle 10.3.4 Rectangle & Circle 10.3.4 Rectangle & Circle 10.3.5 Line & Circle 10.3.5 Line & Circle 10.4 Linear Operations 10.3 Linear Operations	45 45 45 45 45 45 45 46 47 48 48 48
7	7.1	Shortest Path algorithms	29 29 30 30 31 31 31 31 32 32 34 36 37	10	10.1 Linearity 10.1.1 co-linear points 10.2 Polygons 10.2.1 Polygon formation 10.2.2 Polygon Area 10.3 Intersections 10.3 Intersections 10.3.1 Rectangle 10.3.2 Circle 10.3.3 Triangle 10.3.4 Rectangle & Circle 10.3.5 Line & Circle 10.4 Linear Operations 10.5 2D Points	45 45 45 45 45 45 45 46 47 48 48 48

LI	1 Miscenaneous														อ			
	11.1	Faster	impleme	entations														5
		11.1.1	hashes															5
		11.1.2	Binary S	Search th	e v	alue	,											Ξ.

1 Templates

1.1 Setup

1.1.1 IO Manipulation

```
Input/Output
```

```
#include <bits/stdc++.h>
freopen("input.txt", "r", stdin);
freopen("output.txt", "w", stdout);

#define fastI0  \
   ios_base::sync_with_stdio(false), cin.tie(nullptr), cout.tie(nullptr);
```

1.1.2 GCC Compiler Optimization (Vectorization)

GCC Opt

```
// Ref: USACO guide
// will make GCC auto-vectorize for loops and optimizes floating points
  better (assumes associativity and turns off denormals).
#pragma GCC optimize ("Ofast")
// can double performance of vectorized code, but causes crashes on old
  machines.
#pragma GCC target ("avx,avx2,fma")

// slows down run time but throws a Runtime Error if an overflow occured
#pragma GCC optimize("trapy")
```

1.2 MOD Template

```
constexpr int MOD = 1e9+7; // must be a prime number
int add(int a, int b) {
   int res = a+b;
   if(res >= MOD) return res -= MOD;
```

```
}
int sub(int a, int b) {
    int res = a-b;
    if (res < 0) return res += MOD;
}
int power(int a, int e) {
    int res = 1;
    while(e) {if(e & 1) res = res * a % MOD; a = a * a % MOD;
    e >>= 1;}
    return res;
}
// note -> from Fermat little theorem, we can deduce that
// a^(BIG) % prime = (a ^ (BIG%(prime - 1)))%prime
// because every a^(prime - 1) is equal to 1
// from the identity says a^(prime - 1) % prime = 1
int inverse(int a) {
    return power(a, MOD-2);
}
int div(int a, int b) {
    return a * inverse(b) % MOD;
}
```

MOD Template

1.3 Macros

Macros

```
#define getBit(n, k) (n >> k) // & 1
#define ON(n, idx) (n | (111 << idx))
#define OFF(n, idx) (n & ~(111 << idx))</pre>
```

```
#define toggle(n, idx) ((n) ^ (111<<(idx)))</pre>
#define gray(n) (n ^{(n)} (n >> 1))
#define bitCount(x) (__builtin_popcountll(x))
#define clz(x) (__builtin_clzll(x))
#define ctz(x) (__builtin_ctzll(x))
#define uniq(x) x.resize(unique(x.begin(), x.end())-x.begin());
#define angle(a) (atan2((a).imag(), (a).real()))
//\#define vec(a, b) ((b)-(a))
#define same(v1, v2) (dp(vec(v1,v2),vec(v1,v2)) < EPS)
#define dotProduct(a, b) ((conj(a)*(b)).real()) // a*b cos(T), if zero ->
#define crossProduct(a, b) ((conj(a)*(b)).imag()) // a*b sin(T), if zero
    -> parallel
//#define length(a) (hypot((a).imag(), (a).real()))
#define normalize(a) ((a)/length(a))
#define rotateO(v, ang) ((v)*exp(point(0,ang)))
#define rotateA(p, ang, about) (rotateO(vec(about,p),ang)+about)
#define reflectO(v, m) (conj((v)/(m))*(m))
#define ceil_i(a, b) (((ll)(a)+(ll)(b-1))/(ll)(b))
#define floor_i(a, b) (a/b)
#define round_i(a, b) ((a+(b/2))/b) // if a>0
#define round_m(a, b) ((a-(b/2))/b) // if a<0
#define round_multiple(n, m) round_i(n,m)*m // round to multiple if
    specified element
const double PI = acos(-1.0);
```

1.4 Grid Navigation

Grid Nav

```
// knight moves on a chess board
int dx[] = { -2, -1, 1, 2, -2, -1, 1, 2 };
int dy[] = { -1, -2, -2, -1, 1, 2, 2, 1 };

// Grid up, down, right, left (Moves for Chess Rook)
int dx[4] = {1, -1, 0, 0};
int dy[4] = {0, 0, 1, -1};

// Grid cell all neighbours
```

```
const int dx[8] = {1, 0, -1, 0, 1, 1, -1, -1}
const int dy[8] = {0, 1, 0, -1, -1, 1, -1, 1};

// Grid Diagonal (Moves for Chess Bishop)
int dx[] = {1, 1, -1, -1};
int dy[] = {1, -1, 1, -1};
```

1.5 Integer 128

i128

```
typedef __int128 i128;
__int128 read() {
   _{-}int128 x = 0, f = 1;
   char ch = getchar();
   while (ch < '0' || ch > '9') {
       if (ch == '-') f = -1;
       ch = getchar();
   }
   while (ch >= '0' && ch <= '9') {
       x = x * 10 + ch - '0';
       ch = getchar();
   }
   return x * f;
}
void print(__int128 x) {
   if (x < 0) {
       putchar('-');
       x = -x;
   }
   if (x > 9) print(x / 10);
   putchar(x \% 10 + ^{\prime}0^{\prime});
}
bool cmp(__int128 x, __int128 y) { return x > y; }
```

1.6 Matrix Expo

MExpo

```
constexpr int MAX_N = 201, MOD = 1e9+7;
struct Matrix { int mat[MAX_N][MAX_N]; }; // we return a 2D array
ll mod(ll a, ll m) { return ((a%m)+m) % m; } // ensure positive answer
Matrix matMul(Matrix a, Matrix b) { // normally O(n^3)
    Matrix ans{}; // but O(1) as n = 2
   for (auto & i : ans.mat)
       for (int & j : i)
           j = 0;
   for (int i = 0; i < MAX_N; ++i)</pre>
       for (int k = 0; k < MAX_N; ++k) {
           if (a.mat[i][k] == 0) continue; // optimization
           for (int j = 0; j < MAX_N; ++j) {
               ans.mat[i][j] += mod(a.mat[i][k], MOD) * mod(b.mat[k][j],
                  MOD);
               ans.mat[i][j] = mod(ans.mat[i][j], MOD);
           }
       }
   return ans;
Matrix matPow(Matrix base, int p) { // normally O(n^3 log p)
   Matrix ans{}; // but O(\log p) as n = 2
   for (int i = 0; i < MAX_N; ++i)</pre>
       for (int j = 0; j < MAX_N; ++j)</pre>
           ans.mat[i][j] = (i == j); // prepare identity matrix
    while (p) { // iterative D&C version
       if (p&1) // check if p is odd
           ans = matMul(ans, base); // update ans
       base = matMul(base, base); // square the base
       p >>= 1; // divide p by 2
   }
   return ans;
Matrix matMul(Matrix a, Matrix b, int p, int q, int r) { // O(pqr)
   Matrix c{};
   for (int i = 0; i < p; ++i)</pre>
       for (int j = 0; j < r; ++j) {
           c.mat[i][j] = 0;
           for (int k = 0; k < q; ++k)
```

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```

```
c.mat[i][j] += a.mat[i][k] + b.mat[k][j];
}
return c;
```

2 Dynamic Programming

2.1 Some dp patterns

Maximumu/Minimum path cost

```
const int MAX = 21;
int grid[MAX][MAX];
int mem[MAX][MAX];
int n = 20;
bool valid(int r, int c){
 return r >= 0 \&\& r < n \&\& c >= 0 \&\& c < n;
}
int maxPathSum(int r, int c){
   if(!valid(r,c)){
       return 0;
   }
   if(r == n-1 \&\& c == n-1){
       return mem[r][c] = grid[r][c];
   // available moves
   int path1 = maxPathSum(r+1,c);
   int path2 = maxPathSum(r,c+1);
   return grid[r][c] + max(path1,path2);
}
```

add operators between numbers to get max prod/sum

```
// put +, -, between sequence of numbers such that the sum is divisible by k, and maximum as possible
```

```
const int MAX = 21;
long long mem[MAX][MAX];
const int n = 20;
int k = 4; // example
int v[20]:
int fix(int a){
 return (a % k + k) % k;
long long tryAll(int pos, int mod){
   long long &ret = mem[pos][mod];
   if(ret != -1){
       return ret;
   }
   if(pos == n){
       return ret = mod == 0;
   if(tryAll(pos+1,fix(mod + v[pos])) || tryAll(pos+1,fix(mod-v[pos]))){
       return ret = 1;
   return ret = 0;
```

pick choices with no two similar consecutive choices

```
// pick minimum of choinces costs with no two similar consecutive choices
const int choices = 4;
const int n = 20;
int MAX = n;
int mem[MAX][choices];
const int 00 = 1e6+1;
int minCost(int pos, int lastChoice){
   if(pos == n){
        return 0; // invalid move
   }
   int &ret = mem[pos][lastChoice];

   if(ret != -1){
        return ret;
   }
   ret = 00; // want to minimze
```

```
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```

```
7
```

```
// let choices are 0, 1, 2
if(lastChoice != 0){
    ret = min(ret, minCost(pos+1,0));
}
if(lastChoice != 1){
    ret = min(ret, minCost(pos+1,1));
}
if(lastChoice != 2){
    ret = min(ret, minCost(pos+1,2));
}
return ret;
```

sum S and max/min Product

```
int maxK;
11 mem[21][101]; // k, and s
// You are given an integer s and an integer k. Find k positive integers
// such that their sum is equal to s and their product is the maximal
    possible. Return their product.
11 maxProd(int k, int rem)
{
if(k == maxK){
 // base case
 if(rem == 0)
  return 1;
 return 0;
}
 if(rem == 0) // invalid case
 return 0;
11 &ret = mem[k][rem];
 if(ret != -1)
 return ret;
ret = 0:
```

```
for (int i = 1; i <= rem; ++i) {
    ll sol = maxProd(k+1, rem - i) * i;
    ret = max(ret, sol);
}
return ret;
}</pre>
```

2.2 DP solutions

2.2.1 Max Subarray sum (Kadane's Algorithm)

Max Subarray sum

```
int maxSubarraySum(vector<int>& arr, int len) {
   int ans = INT_MIN, cur = 0;

   for (int i = 0; i < len; i++) {
      cur = cur + arr[i];
      if (ans < cur)
           ans = cur;

      if (cur < 0)
           cur = 0;
   }

   return ans;
}</pre>
```

2.2.2 Maximum Subarray Alternating Sum

Maximum Subarray Alternating Sum

```
/* REF: GeeksForGeeks
Input: arr[] = {-4, -10, 3, 5}
Output: 9
Explanation: Subarray {arr[0], arr[2]} = {-4, -10, 3}. Therefore, the sum of this subarray is 9.
```

```
*/
int maxSubarraySumALT(vector<int>& a, int len) {
   int ans = INT_MIN, cur = 0;
   for (int i = 0; i < len; i++) {
       if (i % 2 == 0)
           cur = max(cur + a[i], a[i]);
       else
           cur = max(cur - a[i], -a[i]);
       ans = max(ans, cur);
   }
   cur = 0;
   for (int i = 0; i < len; i++) {</pre>
       if (i % 2 == 1)
           cur = max(cur + a[i], a[i]);
       else
           cur = max(cur - a[i], -a[i]);
       ans = max(ans, cur);
   }
   return ans;
}
```

2.2.3 Count number of DISTINCT ordered ways to produce coins sums to x

Count distinct

```
/*
For example, if the coins are \{2,3,5\} and the desired sum is 9, there
    are 3 ways:

2+2+5
3+3+3
2+2+2+3
*/
int n, x;
cin >> n >> x;
```

```
vector<int> coins(n);
read(coins);

vector dp(x + 1, 0);

dp[0] = 1;
for (int i = 0; i < n; ++i) {
  for (int j = coins[i]; j <= x; ++j) {
    dp[j] = add(dp[j], dp[j - coins[i]]);
  }
}

cout << dp[x] << el;</pre>
```

2.2.4 Min absolute difference between 2 elements from (L, R) (DP Ranges)

Min absolute difference

```
const int N = 1e4 + 1;
int dp[N][N];
int n;
cin >> n;
vector<int> a(n);
read(a);

for (int i = 0; i < n; ++i) dp[i][i] = 1e6; // INF, you can't take the element with it self
for (int i = 1; i < n; ++i) dp[i - 1][i] = abs(a[i] - a[i - 1]);

for (int len = 3; len <= n; ++len) {
  for (int l = 0, r = len - 1; r < n; ++l, ++r) {
    dp[l][r] = min(dp[l][r - 1], dp[l + 1][r]);
    dp[l][r] = min(dp[l][r], abs(a[l] - a[r]));
}
int q;
cin >> q;
```

```
while (q--) {
  int 1, r;
  cin >> 1 >> r;
  --1, --r;

cout << dp[1][r] << el;
}</pre>
```

2.2.5 Longest common subsequence between 2 Strings

$$dp[i][j] = \begin{cases} max(dp[i-1][j], dp[i][j-1]) & \text{if } A_i \neq B_j \\ dp[i-1][j-1] + 1 & \text{if } A_i = B_j \end{cases}$$

LIS 2 Strings

```
// REF: USACO guide
int longestCommonSubsequence(string a, string b) {
int dp[a.size()][b.size()];
for (int i = 0; i < a.size(); i++) { fill(dp[i], dp[i] + b.size(), 0); }</pre>
for (int i = 0; i < a.size(); i++) {</pre>
 if (a[i] == b[0]) dp[i][0] = 1;
 if (i != 0) dp[i][0] = max(dp[i][0], dp[i - 1][0]);
}
for (int i = 0; i < b.size(); i++) {</pre>
 if (a[0] == b[i]) dp[0][i] = 1;
 if (i != 0) dp[0][i] = max(dp[0][i], dp[0][i - 1]);
}
 for (int i = 1; i < a.size(); i++) {</pre>
 for (int j = 1; j < b.size(); j++) {</pre>
  if (a[i] == b[j]) {
   dp[i][j] = dp[i - 1][j - 1] + 1;
  } else {
   dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
 }
return dp[a.size() - 1][b.size() - 1];
```

2.2.6 Longest common subsequence $O(n^2)$

LIS

```
// REF: cp-algorithms
int lis(vector<int> const& a) {
   int n = a.size();
   vector<int> d(n, 1);
   for (int i = 0; i < n; i++) {</pre>
       for (int j = 0; j < i; j++) {
           if (a[i] < a[i])</pre>
               d[i] = max(d[i], d[j] + 1);
       }
   }
   int ans = d[0];
   for (int i = 1; i < n; i++) {</pre>
       ans = max(ans, d[i]);
   return ans;
// Restoring
vector<int> lis(vector<int> const& a) {
   int n = a.size();
   vector\langle int \rangle d(n, 1), p(n, -1);
   for (int i = 0; i < n; i++) {</pre>
       for (int j = 0; j < i; j++) {</pre>
           if (a[j] < a[i] && d[i] < d[j] + 1) {</pre>
               d[i] = d[j] + 1;
               p[i] = j;
       }
   }
   int ans = d[0], pos = 0;
   for (int i = 1; i < n; i++) {
       if (d[i] > ans) {
           ans = d[i];
           pos = i;
       }
   }
```

```
vector<int> subseq;
while (pos != -1) {
    subseq.push_back(a[pos]);
    pos = p[pos];
}
reverse(subseq.begin(), subseq.end());
return subseq;
```

2.2.7 Longest common subsequence Binary Search $O(n + \log N)$

LIS

```
int lisBS(vector<int> const& a) {
   int n = a.size();
   const int INF = 1e9;
   vector<int> d(n+1, INF);
   d[0] = -INF;
   for (int i = 0; i < n; i++) {
       int 1 = upper_bound(d.begin(), d.end(), a[i]) - d.begin();
       if (d[l-1] < a[i] && a[i] < d[l])</pre>
           d[1] = a[i];
   }
   int ans = 0;
   for (int 1 = 0; 1 <= n; 1++) {
       if (d[1] < INF)</pre>
           ans = 1:
   }
   return ans;
}
```

3 Bit Manipulation

3.1 Subset Operations

count subsets with give sum

```
int countDistinctSubsetsWithSum(vector<int>& arr, int n, int k) {
    // Count distinct subsets of array arr that sum up to k
    vector<int> dp(k + 1, 0);
    dp[0] = 1;
    for (int i = 0; i < n; ++i) {
        for (int j = k; j >= arr[i]; --j) {
            dp[j] += dp[j - arr[i]];
        }
    }
    return dp[k]; // Number of distinct subsets with sum k
}
```

max xor of any subset of elements in the array

```
int maximalSubsetXOR(vector<int>& arr, int n) {
    // Find the maximum XOR of any subset of elements in array arr
    int maxXor = 0;
    for (int mask = 0; mask < (1 << n); ++mask) {
        int xorSum = 0;
        for (int i = 0; i < n; ++i) {
            if (mask & (1 << i)) {
                 xorSum ^= arr[i];
            }
        }
        maxXor = max(maxXor, xorSum);
    }
    return maxXor;
}</pre>
```

min xor of any subset

subset generation

```
void subsetGeneration(int x, int n) {
    // Generate all non-empty subsets of a set represented by an integer x
    for (int subset = x; subset > 0; subset = (subset - 1) & x) {
        // Process subset
        cout << subset << endl;
    }
}</pre>
```

check if subset of elements in the array sum up to k

```
void subsetSumCheck(vector<int>& arr, int n, int k) {
    // Check if a subset of elements in array arr sums up to k
    for (int subset = 0; subset < (1 << n); ++subset) {
        int sum = 0;
        for (int i = 0; i < n; ++i) {
            if (subset & (1 << i)) {
                sum += arr[i];
            }
        }
        if (sum == k) {
                // Found subset with sum k
                cout << "Subset with sum " << k << ": " << subset << endl;
        }
    }
}</pre>
```

$max \ subset \ sum \ mod \ m$

```
int subsetWithMaxSumModuloM(vector<int>& arr, int n, int m) {
    // Find the maximum subset sum modulo m
```

iterate over all supersets represented by x

```
void iterateOverSupersets(int x, int n) {
    // Iterate over all supersets of a set represented by x
    int subset = x;
    do {
        // Process subset
        cout << subset << endl;
        subset = (subset + 1) | x;
    } while (subset <= (1 << n) - 1);
}</pre>
```

4 Algorithms

4.1 FFT

FFT Algorithm

```
constexpr 11 mod = 998244353, root = 3;

11 modpow(11 b, 11 e, 11 m) {
    l1 ans = 1;
    for (; e; b = b * b % m, e /= 2)
```

```
if (e & 1) ans = ans * b % m:
   return ans:
}
// Primitive Root of the mod of form 2^a * b + 1
int generator () {
   vector<int> fact;
   int phi = mod-1, n = phi;
   for (int i=2; i*i<=n; ++i)</pre>
       if (n % i == 0) {
           fact.push_back (i);
           while (n \% i == 0)
              n /= i:
       }
   if (n > 1)
       fact.push_back (n);
   for (int res=2; res<=mod; ++res) {</pre>
       bool ok = true;
       for (size_t i=0; i<fact.size() && ok; ++i)</pre>
           ok &= modpow(res, phi / fact[i], mod) != 1;
       if (ok) return res;
   }
   return -1;
}
void ntt(vector<ll> &a) {
   int n = a.size(), L = 31 - __builtin_clz(n);
   static vector<11> rt(2, 1); // erase the static if you want to use two
        moduli;
   for (static int k = 2, s = 2; k < n; k *= 2, s++) { // erase the
       static if you want to use two moduli;
       rt.resize(n);
       ll z[] = \{1, modpow(root, mod >> s, mod)\};
       for (int i = k; i < 2*k; ++i) rt[i] = rt[i / 2] * z[i & 1] % mod;
   }
   vector<int> rev(n);
   for (int i = 0; i < n; ++i) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
   for (int i = 0; i < n; ++i) if (i < rev[i]) swap(a[i], a[rev[i]]);
   for (int k = 1; k < n; k *= 2) {
       for (int i = 0; i < n; i += 2 * k) {
           for (int j = 0; j < k; ++j) {
              ll z = rt[j + k] * a[i + j + k] % mod, &ai = a[i + j];
```

```
a[i + j + k] = ai - z + (z > ai ? mod : 0);
              ai += (ai + z >= mod ? z - mod : z):
          }
       }
   }
vector<ll> multi(const vector<ll> &a, const vector<ll> &b) {
   if (a.empty() || b.empty()) return {};
   int s = static_cast<ll>(a.size()) + static_cast<ll>(b.size()) - 1, B =
        32 - \_builtin\_clz(s), n = 1 << B;
   int inv = modpow(n, mod - 2, mod);
   vector<ll> L(a), R(b), out(n);
   L.resize(n), R.resize(n);
   ntt(L), ntt(R);
   for (int i = 0; i < n; ++i) out[-i & (n - 1)] = L[i] * R[i] % mod *
       inv % mod;
   ntt(out);
   return {out.begin(), out.begin() + s};
vector<int> poly_pow(vector<int> poly, int p, int limit = 1e9) {
   vector<int> ans{1};
   while (p) {
       if(p&1) ans = multi(ans, poly);
       poly = multi(poly, poly);
       ans.resize(limit + 1);
       poly.resize(limit + 1);
       p >>= 1;
   }
   return ans;
```

4.2 MO

MO Algorithm

```
int nums[N], q_ans[M];
struct query {
   int idx, block_idx, l, r;
   query() = default;
   query(int _1, int _r, int _idx) {
       idx = _idx;
       r = _r - 1;
       1 = _1 - 1;
       block_idx = _l / sqrt(n);
   }
   bool operator <(const query & y) const {</pre>
       if(y.block_idx == block_idx) return r < y.r;</pre>
       return block_idx < y.block_idx;</pre>
};
int freq[N], ans;
void add(int idx) {
   freq[nums[idx]]++;
   if (freq[nums[idx]] == 2) ans++;
}
void remove(int idx) {
   freq[nums[idx]]--;
   if (freq[nums[idx]] == 1) ans--;
}
cin >> n >> m;
for (int i = 0; i < n; ++i) cin >> nums[i];
vector<query> Query(m);
for (int i = 0; i < m; ++i) {</pre>
   int 1, r; cin >> 1 >> r;
   Query[i] = query(1, r, i);
}
sort(Query.begin(), Query.end());
int 10 = 1, r0 = 0;
for (int i = 0: i < m: ++i) {
```

```
while (10 < Query[i].1) remove(10++);
while (10 > Query[i].1) add(--10);
while (r0 < Query[i].r) add(++r0);
while (r0 > Query[i].r) remove(r0--);
q_ans[Query[i].idx] = ans;
}
for (int i = 0; i < m; ++i) {
   cout << q_ans[i] << '\n';
}</pre>
```

4.3 Intervals

4.3.1 Prefix Sum (L, R) intervals

Prefix Sum (L, R) intervals

```
// NOTE: works fine with small n or with large memory
int main() {
   int n, k;
   cin >> n >> k;
   vector < int > a(n + 1);
   vector<vector<int>> rangesPrefix(n + 1, vector<int>(n + 1, 0));
   for (int i = 1; i <= n; ++i)
       cin >> a[i]:
   int 1 = 1, r = 1, sum = 0;
   // validate your intervals
   // here the intervals are the ones that have a sum of k
   while (r \le n) {
       sum += a[r];
       while (sum > k) {
           sum -= a[1];
           ++1;
       while (1 <= r && a[1] == 0) {</pre>
           if (sum != k)
```

}

```
break;
       rangesPrefix[r][1]++;
       ++1;
    }
    if (sum == k) {
       rangesPrefix[r][1]++;
    ++r;
// prefix sum the columns
for (int i = 1; i <= n; ++i) {</pre>
   for (int j = n - 1; j \ge 0; --j) {
       rangesPrefix[i][j] += rangesPrefix[i][j + 1];
    }
}
// prefix sum the rows
for (int i = 0; i <= n; ++i) {
    for (int j = 1; j \le n; ++j) {
       rangesPrefix[j][i] += rangesPrefix[j - 1][i];
    }
}
int q; cin >> q;
while (q--) {
    cin >> 1 >> r;
    // answer the number of intervals (X, Y) X <= Y that are included
       between L, R
    cout << rangesPrefix[r][l] - rangesPrefix[l - 1][l] << el;</pre>
```

4.3.2 Find subarrays intervals that sum to K Using Map

Find subarray intervals that sum to K Using Map

```
int n, k;
cin >> n >> k;
vector<int> a(n + 1);
vector<pair<int, int>> rng;
for (int i = 1; i <= n; ++i)</pre>
   cin >> a[i];
map<int, set<int>> prev;
int currSum = 0;
for (int i = 1; i <= n; ++i) {</pre>
    currSum += a[i];
   if (currSum == k) {
       rng.push_back({1, i});
   if (prev.find(currSum - k) != prev.end()) {
       for (auto &j : prev[currSum - k]) {
           rng.push_back(\{j + 1, i\});
   }
   prev[currSum].insert(i);
}
```

4.4 Ad-hoc

4.4.1 Find duplicate

Find duplicate using XOR

```
int findDuplicate(int arr[] , int n)
{
   int answer=0;
    //XOR all the elements with 0
   for(int i=0; i<n; i++){
      answer=answer^arr[i];
   }
   //XOR all the elements with no from 1 to n
   // i.e answer^0 = answer
   for(int i=1; i<n; i++){</pre>
```

```
answer=answer^i;
}
return answer;
}
```

4.5 Sorting Algorithms

4.5.1 Radix Sort

Radix Sort

```
// O(n + b), where n is the number of elements and b is the base of the
    number system
// A function to do counting sort of arr[] according to the digit
    represented by exp.
void countingSort(vector<int>& arr, int exp) {
   int n = arr.size();
   vector<int> output(n); // output array
   int count[10] = {0};
   // Store count of occurrences in count[]
   for (int i = 0; i < n; i++)</pre>
       count[(arr[i] / exp) % 10]++;
   // Change count[i] so that count[i] now contains the actual
   // position of this digit in output[]
   for (int i = 1; i < 10; i++)
       count[i] += count[i - 1];
   // Build the output array
   for (int i = n - 1; i \ge 0; i--) {
       output[count[(arr[i] / exp) % 10] - 1] = arr[i];
       count[(arr[i] / exp) % 10]--;
   }
   // Copy the output array to arr[], so that arr now
   // contains sorted numbers according to the current digit
   for (int i = 0; i < n; i++)</pre>
       arr[i] = output[i];
}
```

4.5.2 Counting Sort

Counting Sort

```
// O(N+M), where N and M are the size of inputArray[] and countArray[]
// The main function that sorts arr[] of size n using Counting Sort
void countingSort(vector<int>& arr) {
   int maxElement = *max_element(arr.begin(), arr.end());
   int minElement = *min_element(arr.begin(), arr.end());
   int range = maxElement - minElement + 1;
   vector<int> count(range), output(arr.size());
   for (int i = 0; i < arr.size(); i++)</pre>
       count[arr[i] - minElement]++;
   for (int i = 1; i < count.size(); i++)</pre>
       count[i] += count[i - 1];
   for (int i = arr.size() - 1; i >= 0; i--) {
       output[count[arr[i] - minElement] - 1] = arr[i];
       count[arr[i] - minElement]--;
   for (int i = 0; i < arr.size(); i++)</pre>
       arr[i] = output[i];
```

4.6 Apply permutation k times

$permutation \ k \ times$

```
// Applying a permutation k times
// n log k
vector<int> applyPermutation(vector<int> sequence, vector<int> permutation
   ) {
   vector<int> newSequence(sequence.size());
   for(int i = 0; i < sequence.size(); i++) {</pre>
       newSequence[i] = sequence[permutation[i]];
   }
   return newSequence;
}
vector<int> permute(vector<int> sequence, vector<int> permutation, long
   long k) {
   while (k > 0) {
       if (k & 1) {
           sequence = applyPermutation(sequence, permutation);
       permutation = applyPermutation(permutation, permutation);
   }
   return sequence;
}
```

4.7 Strings

4.7.1 Strings Hashing Utils

Hashing

```
// #define int long long
struct PreHash {
   vector<int> pow, invPow;
   const int MOD, BASE;

  int modExp(int base, int exp) const {
    int res = 1;
```

```
while (exp > 0) {
           if (exp % 2 == 1)
              res = 1LL * res * base % MOD;
           base = 1LL * base * base % MOD;
           exp /= 2;
       }
       return res;
   }
   explicit PreHash(const int n, const int MOD, const int BASE) : MOD(MOD
       ), BASE(BASE) {
       pow.resize(n);
       invPow.resize(n);
       pow[0] = 1;
       invPow[0] = 1;
       const int invBase = modExp(BASE, MOD - 2);
       for (int i = 1; i < n; ++i) {</pre>
           pow[i] = 1LL * pow[i - 1] * BASE % MOD;
           invPow[i] = 1LL * invPow[i - 1] * invBase % MOD;
   }
};
struct Hash {
   vector<int> pref;
   const PreHash* pre;
   int n;
   explicit Hash(const string& s, const PreHash& p) : pre(&p) {
       n = s.size();
       if (s.empty()) return;
       pref.resize(n);
       pref[0] = s[0] % pre->MOD;
       for (int i = 1; i < n; ++i) {
          pref[i] = (pref[i - 1] + 1LL * s[i] * pre->pow[i]) % pre->MOD;
       }
   }
   int getHashRange(int 1, int r) const {
```

```
int ret = pref[r];
       if (1 > 0) ret = (ret - pref[1 - 1] + pre->MOD) % pre->MOD;
       // if (ret < 0) ret += pre->MOD;
       return 1LL * ret * pre->invPow[1] % pre->MOD;
};
struct HashWhole {
   int pref, suff;
   const PreHash* pre;
   int n;
   explicit HashWhole(const string& s, const PreHash& p) : pre(&p) {
       n = s.size();
       if (s.empty()) return;
       pref = s[0] % pre->MOD;
       for (int i = 1; i < n; ++i) {</pre>
           pref = (pref + 1LL * s[i] * pre->pow[i]) % pre->MOD;
       suff = s[n - 1] \% pre->MOD;
       for (int i = n - 2; i \ge 0; --i) {
           suff = (suff + 1LL * s[i] * pre->pow[n - 1 - i]) % pre->MOD;
       }
   }
   void mergeHash(const HashWhole& other) {
       pref = (1LL * pref) + (other.pref * 1LL * (pre->pow[n] % pre->MOD)
           % pre->MOD);
       pref %= pre->MOD;
       suff = ((other.suff * 1LL) + 1LL * suff * (pre->pow[other.n] % pre
           ->MOD)) % pre->MOD;
       suff %= pre->MOD;
       n += other.n;
   }
   bool isPalindrome() const {
       return suff == pref;
   }
```

```
};
struct Hasher {
    const PreHash* pre;
    explicit Hasher(const PreHash& p) : pre(&p) {}
    int hash(const string& s, int st, const int end) const {
       if (s.empty() || st > end) return 0;
       int pref = s[st] % pre->MOD, i = 1;
       for (st = st + 1; st < end; ++st) {</pre>
           pref = (pref + 1LL * s[st] * pre->pow[i++]) % pre->MOD;
       return pref;
   }
};
constexpr int N = 3e5 + 10;
const PreHash pre2(N, 1e9 + 9, 107);
const PreHash pre1(N, 1e9 + 7, 37);
// #undef int long long
```

4.7.2 KMP

KMP

```
vector<int> kmp(const string& pat) {
   int n = pat.size();
   vector<int> pi(n);
   pi[0] = -1;

for (int i = 1; i < n; i++){
      int j = pi[i - 1];
      while (j >= 0 && pat[j] != pat[i - 1]) j = pi[j];
           pi[i] = j + 1;
   }

return pi;
```

```
vector<int> find_matches(const string& text, const string& pat) {
   int n = pat.size(), m = text.size();
   string s = pat + "$" + text + "$";

   vector<int> pi = kmp(s), ans;
   for (int i = 0; i <= m; i++)
        if (pi[i + n + 1] == n)
            ans.push_back(i - n);

   return ans;
}
</pre>
```

4.7.3 Manachers Algorithm

Manachers Algorithm

```
struct Manacher {
 vector<int> p[2];
 // p[1][i] = (max odd length palindrome centered at i) / 2 [floor
     division]
 // p[0][i] = same for even, it considers the right center
 // e.g. for s = "abbabba", p[1][3] = 3, p[0][2] = 2
 Manacher(string s) {
   int n = s.size();
   p[0].resize(n + 1);
   p[1].resize(n);
   for (int z = 0; z < 2; z++)
     for (int i = 0, l = 0, r = 0; i < n; i++) {
       int t = r - i + !z;
       if (i < r) p[z][i] = min(t, p[z][1 + t]);
       int L = i - p[z][i], R = i + p[z][i] - !z;
       while (L >= 1 \&\& R + 1 < n \&\& s[L - 1] == s[R + 1]) p[z][i]++, L--,
           R++;
       if (R > r) 1 = L, r = R;
     }
 bool is_palindrome(int 1, int r) {
   int mid = (1 + r + 1) / 2, len = r - 1 + 1;
   return 2 * p[len % 2][mid] + len % 2 >= len;
```

} };

5 Data Structures

5.1 Monotonuc

5.1.1 Min Aggregation Deque

Min Aggregation Deque

```
struct minstack {
    stack<pair<int, int>> st;
   int getmin() { return st.top().second; }
   bool empty() { return st.empty(); }
    int size() { return st.size(); }
    void push(int x) {
       int mn = x;
       if (!empty()) mn = min(mn, getmin());
       st.push({x, mn});
   void pop() { st.pop(); }
    int top() { return st.top().first; }
   void swap(minstack &x) { st.swap(x.st); }
};
struct mindeque {
   minstack l, r, t;
    void rebalance() {
       bool f = false;
       if (r.empty()) {
           f = true;
           1.swap(r);
       int sz = r.size() / 2;
       while (sz--) {
           t.push(r.top());
           r.pop();
       while (!r.empty()) {
```

```
1.push(r.top());
          r.pop();
       while (!t.empty()) {
          r.push(t.top());
          t.pop();
       if (f) 1.swap(r);
   }
   int getmin() {
       if (l.empty()) return r.getmin();
       if (r.empty()) return l.getmin();
       return min(l.getmin(), r.getmin());
   }
   bool empty() { return l.empty() && r.empty(); }
   int size() { return l.size() + r.size(); }
   void push_front(int x) { l.push(x); }
   void push_back(int x) { r.push(x); }
   void pop_front() {
       if (l.empty()) rebalance();
       1.pop();
   }
   void pop_back() {
       if (r.empty()) rebalance();
       r.pop();
   }
   int front() {
       if (1.empty()) rebalance();
       return 1.top();
   }
   int back() {
       if (r.empty()) rebalance();
       return r.top();
   }
   void swap(mindeque &x) {
       1.swap(x.1);
       r.swap(x.r);
   }
};
```

5.2 Strings

5.2.1 Trie (Prefix Tree)

$Basic\ Implementation$

```
#define MAX_CHAR 26
struct TrieNode {
   TrieNode *pTrieNode[MAX_CHAR]{};
   bool isWord;
   TrieNode() {
       isWord = false;
       fill(pTrieNode, pTrieNode + 26, (TrieNode *) NULL);
   virtual ~TrieNode() = default;
};
class Trie {
private:
   TrieNode *root;
public:
   Trie() {
       root = new TrieNode();
   }
   virtual ~Trie() = default;
   TrieNode *getTrieNode() {
       return this->root;
   }
   void insert(const string &word) {
       TrieNode *current = root:
       for (char c: word) {
           int i = c - 'a';
           if (current->pTrieNode[i] == nullptr)
              current->pTrieNode[i] = new TrieNode();
           current = current->pTrieNode[i];
       current->isWord = true;
```

```
}
   bool search(const string &word) {
       TrieNode *current = root;
       int ch = 0:
       for (char c: word) {
          ch = c - 'a';
          if (current->pTrieNode[ch] == nullptr)
              return false;
           current = current->pTrieNode[ch];
       return current->isWord;
   }
   bool startsWith(const string &prefix) {
       TrieNode *current = root;
       int ch = 0;
       for (char c: prefix) {
          ch = c - 'a';
          if (current->pTrieNode[ch] == nullptr)
              return false;
          current = current->pTrieNode[ch];
       }
       return true;
   }
};
```

5.3 Range Queries

5.3.1 Segment Tree

$Basic\ Implementation$

```
struct Node {
    long long val;
};

struct SegTree {
private:
    const Node NEUTRAL = {INT_MIN};
```

```
static Node merge(const Node& x1, const Node& x2) {
   return {x1.val + x2.val};
}
void set(const int& idx, const int& val, int x, int lx, int rx) {
   if (rx - lx == 1) return void(values[x].val = val);
   int mid = (rx + lx) / 2;
   if (idx < mid)</pre>
       set(idx, val, 2 * x + 1, lx, mid);
       set(idx, val, 2 * x + 2, mid, rx);
   values[x] = merge(values[2 * x + 1], values[2 * x + 2]);
}
Node query(const int& 1, const int& r, int x, int lx, int rx) {
   if (lx >= r || l >= rx) return NEUTRAL;
   if (lx >= 1 && rx <= r) return values[x];</pre>
   int mid = (rx + lx) / 2;
   return merge(query(1, r, 2 * x + 1, lx, mid), query(1, r, 2 * x +
        2, mid, rx));
}
void build(vector<int> &a, int x, int lx, int rx) {
   if (rx - lx == 1) {
       if (lx < a.size()) {</pre>
           values[x].val = a[lx];
       return;
   int m = (1x + rx) / 2;
   build(a, 2 * x + 1, lx, m);
   build(a, 2 * x + 2, m, rx);
   values[x] = merge(values[2 * x + 1], values[2 * x + 2]);
}
 void assign_range(int 1, int r, int node, int lx, int rx, int time,
     int val) {
   if (lx > r || l > rx) return;
   if (lx >= 1 && rx <= r) {
```

```
lazy[node] = {time, val};
           return;
       int mid = (lx+rx) / 2;
       assign_range(l, r, 2*node+1, lx, mid, time, val);
       assign_range(l, r, 2*node+2, mid+1, rx, time, val);
   }
   pair<int, int> point_query(int lx, int rx, int node, int idx) {
       if(rx == lx) return lazy[node];
       int mid = (1x+rx) / 2;
       if(idx <= mid) {</pre>
           auto x = point_query(lx, mid, 2*node+1, idx);
           if(x.first > lazy[node].first) return x;
           return lazy[node];
       }
       auto x = point_query(mid+1, rx, 2*node+2, idx);
       if(x.first > lazy[node].first) return x;
       return lazy[node];
   }
public:
   int size{};
   vector<Node> values;
   void build(vector<int> &a) {
       build(a, 0, 0, size);
   }
   void init(int _size) {
       size = 1;
       while (size < _size) size *= 2;</pre>
       values.assign(2 * size, NEUTRAL);
   }
   void set(int idx, int val) {
       set(idx, val, 0, 0, size);
   }
   Node query(const int& 1, const int& r) {
       return query(1, r, 0, 0, size);
```

```
};
```

5.3.2 Lazy Propegation

Lazy Propegation

```
struct SegTree {
private:
   void propegate(int lx, int rx, int node) {
       if(!lazy[node]) return;
       if(lx != rx) {
          lazy[2*node+1] = lazy[node];
          lazy[2*node+2] = lazy[node];
       values[node] = lazy[node] * (rx - lx + 1);
       lazy[node] = 0;
   }
   // assign val in range [l, r]
   void update_range(int 1, int r, int node, int lx, int rx, int val,
       bool f) {
       propegate(lx, rx, node);
       if (lx > r \mid | l > rx) return;
       if (lx >= 1 && rx <= r) {
          lazy[node] = val;
          propegate(lx, rx, node);
          return;
       int mid = (1x+rx) / 2;
       update_range(1, r, 2*node+1, lx, mid, val, f);
       update_range(l, r, 2*node+2, mid+1, rx, val, f);
       values[node] = values[2*node+1] + values[2*node+2];
   }
   // get sum in range [1, r]
   int range_query(int 1, int r, int lx, int rx, int node) {
       propegate(lx, rx, node);
       if (lx > r || l > rx) return 0;
```

```
if (lx >= 1 && rx <= r) return values[node];</pre>
       int mid = (1x+rx) / 2;
       return range_query(1, r, lx, mid, 2*node+1) + range_query(1, r, mid
           +1, rx, 2*node+2);
   }
public:
   int size{};
   vector<int> values, lazy;
   void init(int _size) {
       size = 1;
       while (size < _size) size *= 2;</pre>
       values.assign(2 * size, 0);
       lazy.assign(2 * size, 0);
   void update_range(int 1, int r, int v, bool f) {
       update_range(1, r, 0, 0, size-1, v, f);
   }
   int range_query(int 1, int r) {
       return range_query(1, r, 0, size-1, 0);
   }
};
```

5.3.3 Fenwick Tree

Fenwick Tree

```
struct Fenwick {
    // One Based
    vector<int> tree;

explicit Fenwick(int n) {tree.assign(n + 5, {});}

// Computes the prefix sum from [1, i], O(log(n))
int query(int i) {
    int res = 0;
    while (i > 0) {
```

```
res += tree[i];
       i &= ~(i & -i);
   return res;
}
int query(int 1, int r) {
   return query(r) - query(1-1);
}
// Get the value at index i
int get(int i) {
    return query(i, i);
}
// Add 'v' to index 'i', O(log(n))
void update(int i, int v) {
   while (i < tree.size()) {</pre>
       tree[i] += v;
       i += (i & -i);
}
// Update range, Point query
// To get(k) do prefix sum [1, k] and in insert update_range(i, i, a[i
void update_range(int 1, int r, int v) {
   update(1, v);
   update(r+1, -v);
}
```

5.3.4 Fenwick UpdateRange

$BIT\ UpdateRange$

```
struct BITUpdateRange {
private:
   int n;
   vector<int> B1, B2;
```

```
void add(vector<int> &b, int idx, int x) {
       while (idx <= n) {</pre>
           b[idx] += x;
           idx += idx & -idx;
   }
   int sum(vector<int> &b, int idx) {
       int total = 0;
       while (idx > 0) {
          total += b[idx];
           idx &= (idx & -idx);
       return total;
   int prefix(int idx) {
       return sum(B1, idx) * idx - sum(B2, idx);
   }
public:
   explicit BITUpdateRange(int n) : n(n) {
       B1.assign(n + 1, {});
       B2.assign(n + 1, {});
   }
   void update(int 1, int r, int x) {
       add(B1, 1, x);
       add(B1, r + 1, -x);
       add(B2, 1, x * (1 - 1));
       add(B2, r + 1, -x * r);
   }
   int query(int i) {
       return prefix(i) - prefix(i - 1);
   }
   int query(int 1, int r) {
       return prefix(r) - prefix(l - 1);
   }
};
```

5.3.5 2D BIT

2D BIT

```
struct BIT2D {
   int n, m;
   vector<vector<int>> bit;
   BIT2D(int n, int m) : n(n), m(m) {
       bit.assign(n + 2, vector<int>(m + 2));
   }
   void update(int x, int y, int val) {
       for (; x \le n; x += x & -x) {
          for (int i = y; i <= m; i += i & -i) {</pre>
              bit[x][i] += val;
       }
   }
   int prefix(int x, int y) {
       int res = 0;
       for (; x > 0; x &= (x & -x)) {
          for (int i = y; i > 0; i &= (i & -i)) {
              res += bit[x][i];
       }
       return res;
   }
   int query(int sx, int sy, int ex, int ey) {
       int ans = 0;
       ans += prefix(ex, ey);
       ans -= prefix(ex, sy - 1);
       ans -= prefix(sx - 1, ey);
       ans += prefix(sx - 1, sy - 1);
       return ans;
```

5.3.6 Sparse Table

Impl with the index

```
// storing the index also
struct SNode {
   int val;
   int index;
};
class SparseTable {
private:
   vector<vector<SNode>> table;
   function<SNode(const SNode&, const SNode&)> merge;
   static SNode StaticMerge(const SNode& a, const SNode& b) {
       return a.val < b.val ? a : b;</pre>
   }
public:
   explicit SparseTable(const vector<int>& arr, const function<SNode(</pre>
       const SNode&, const SNode&)>& mergeFunc = StaticMerge) {
       int n = static_cast<int>(arr.size());
       int log_n = static_cast<int>(log2(n)) + 1;
       this->merge = mergeFunc;
       table.resize(n, vector<SNode>(log_n));
       for (int i = 0; i < n; i++) {</pre>
           table[i][0] = {arr[i], i};
       }
       for (int j = 1; (1 << j) <= n; j++) {
           for (int i = 0; i + (1 << j) <= n; i++) {
               table[i][j] = mergeFunc(table[i][j - 1], table[i + (1 << (j
                   - 1))][i - 1]);
           }
   SNode query(int left, int right) {
       int j = static_cast<int>(log2(right - left + 1));
```

```
return merge(table[left][j], table[right - (1 << j) + 1][j]);</pre>
   }
   // query in O(log(n)) if its could't apply to Sparse Table directly
   T query_log(int 1, int r){
     int len = r - l + 1;
     T ans;
     for(int i = 0; 1 <= r; i++){
         if (len & (1 << i)){</pre>
             ans = merge(ans, table[i][1]);
             1+=(1 << i);
  }
};
int main(void) {
   int n;
   cin >> n;
   vector<int> arr(n);
   for (auto& element : arr) cin >> element;
   SparseTable minSt(arr, [](const SNode& a, const SNode& b) -> SNode {
       return a.val < b.val ? a : b;</pre>
   });
   SparseTable maxSt(arr, [](const SNode& a, const SNode& b) -> SNode {
       return a.val > b.val ? a : b;
   });
```

5.3.7 Sparse Table (Shorter Version)

Impl

```
template < class T>
struct Rmq {
  vector < vector < T>> jmp;

  // vector is 0-based indexed
```

```
Rmq(const vector<T>& V) : jmp(1, V) {
  for (int pw = 1, k = 1; pw * 2 <= V.size(); pw *= 2, ++k) {
    jmp.emplace_back(V.size() - pw * 2 + 1);
    for(int j = 0; j < jmp[k].size(); j++)
        jmp[k][j] = (jmp[k - 1][j] | jmp[k - 1][j + pw]);
  }
}

// query (a, b] inclusive execlusive
T query(int a, int b) {
  // assert(a < b);
  int dep = 31 - __builtin_clz(b - a);
  return (jmp[dep][a] | jmp[dep][b - (1 << dep)]);
}
};</pre>
```

5.3.8 Sqrt Decomp

Impl

```
#define 00 1e18
const int sz = 2e5 + 5;
const int sq = 500;
int arr[sz] = {};
int ans[sq] = {};
void pre() {
   fill(ans, ans + sq, 00);
   for(int i = 0; i < sz; i++) {</pre>
       ans[i / sq] = min(ans[i / sq], arr[i]);
}
// can be O(1) with inversable operations, else it's sqrt
void update(int k, int x) {
  arr[k] = x;
  int start = k / sq;
  ans[start] = *min_element(arr + (start * sq), arr + (start + 1) * sq);
}
```

```
int ret = 00;
while(1 <= r) {
    if(1 % sq == 0 && 1 + sq <= r) {
        ret = min(ret, ans[1 / sq]);
        1 += sq;
    }else {
        ret = min(ret, arr[1]);
        1++;
    }
}
return ret;
}</pre>
```

5.4 Ordered Set

int query(int 1, int r){

Ordered Set

```
// largest element in the set
cout << *S.find_by_order(0) << " ";

// Returns iterator to 2-nd
// largest element in the set
cout << *S.find_by_order(2);</pre>
```

5.5 Custom Compare Functions

Custom Compare functions

```
template<class T>
struct custom_compare {
   bool operator()(const T& a, const T& b) const {
       if (a == b) return true; // Keep duplicates
       return a > b;
   }
};
//REF: GFG
class CustomComparator {
public:
   CustomComparator(int baseValue) : baseValue_(baseValue) {}
   bool operator()(int a, int b) const {
       // Custom comparison logic involving state
       return (a % baseValue_) < (b % baseValue_);</pre>
   }
private:
   int baseValue_;
};
// OR through capture by reference (capture clauses)
auto compare = [&](char a, char b) { return localStructure[a] >
   localStructure[b]; };
```

5.6 Median Finder

Median Finder

```
struct MedianFinder {
   multiset<ll> low, up;
   11 \text{ sumLow} = 0, \text{ sumUp} = 0;
   void insert(ll val) {
       if (low.empty() || val <= *low.rbegin()) {</pre>
           low.insert(val);
           sumLow += val;
       } else {
           up.insert(val);
           sumUp += val;
       rebalance();
   }
   void erase(ll val) {
       if (low.find(val) != low.end()) {
           low.erase(low.find(val));
           sumLow -= val;
       } else if (up.find(val) != up.end()) {
           up.erase(up.find(val));
           sumUp -= val;
       rebalance();
   }
   void rebalance() {
       // Ensure low has (n+1)/2 elements
       while (low.size() > up.size() + 1) {
           ll val = *low.rbegin();
           low.erase(prev(low.end()));
           sumLow -= val;
           up.insert(val);
           sumUp += val;
       while (low.size() < up.size()) {</pre>
           ll val = *up.begin();
           up.erase(up.begin());
           sumUp -= val;
```

5.7 Binary Trie

Binary Trie

```
template <typename IntegerType = int, int INT_SIZE = 31, int capacity = 1
    e5>
struct BinaryTrie {
    int nxt[capacity][2]{}, freq[capacity]{}, nodesCount;

BinaryTrie() {
        nodesCount = 0;
    }

    void add(const IntegerType& x) {
        int node = 0;
        for (int i = INT_SIZE; i >= 0; --i) {
            bool bit = (x >> i) & 1;

        if (nxt[node][bit] == 0)
            nxt[node][bit] = ++nodesCount;

        node = nxt[node][bit];
        ++freq[node];
```

```
}
}
void erase(const IntegerType& x) {
   int node = 0;
   for (int i = INT_SIZE; i >= 0; --i) {
       bool bit = (x >> i) & 1;
       node = nxt[node][bit];
       --freq[node];
}
IntegerType getMin(const IntegerType& x) {
   IntegerType res = 0;
   int node = 0;
   for (int i = INT_SIZE; i >= 0; --i) {
       bool bit = (x >> i) & 1;
       if (nxt[node][bit] != 0 && freq[nxt[node][bit]]) {
           node = nxt[node][bit];
       }
       else {
           node = nxt[node][bit ^ 1];
          res |= (1LL << i);
       }
   }
   return res;
}
IntegerType getMax(const IntegerType& x) {
   IntegerType res = 0;
   int node = 0;
   for (int i = INT_SIZE; i >= 0; --i) {
       bool bit = (x >> i) & 1;
       if (nxt[node][bit ^ 1] != 0 && freq[nxt[node][bit ^ 1]]) {
           node = nxt[node][bit ^ 1];
           res |= (1LL << i);
       }
       else {
           node = nxt[node][bit];
       }
   }
```

```
return res;
}
```

6 Counting Principles

6.1 nCr

$$C(n,k) = \frac{n!}{(n-k)!k!} = \frac{n*(n-1)*(n-2)*...*(n-k+1)}{k!}$$

6.1.1 Fast nCr

$$C(n,k) = \frac{n*(n-1)*(n-2)*\dots*(n-k+1)}{1*2*3*\dots*k} = \prod_{i=0}^{k-1} \frac{n-i}{i+1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{if (i != j) { // skips over the base cases}}{dp[i][j] = (dp[i-1][j-1] + dp[i-1][j]) \% p;}$$

Fast nCr

```
int nCr(const int& n, const int& r) {
   double res = 1;
   for (int i = 1; i <= r; ++i)
      res = res * (n - r + i) / i;
   return (int)(res + 0.01);
}</pre>
```

6.1.2 Method 1: Pascal's Triangle (Dynamic Programming) - $\mathcal{O}(n^2)$

nCk using dp

```
// REF: USACO guide

/** @return nCk mod p using dynamic programming */
int binomial(int n, int k, int p) {
   // dp[i][j] stores iCj
   vector<vector<int>> dp(n + 1, vector<int>(k + 1, 0));

   // base cases described above
```

```
for (int i = 0; i <= n; i++) {</pre>
  * i choose 0 is always 1 since there is exactly one way
  * to choose O elements from a set of i elements
  * (don't choose anything)
  */
 dp[i][0] = 1;
  * i choose i is always 1 since there is exactly one way
  * to choose i elements from a set of i elements
  * (choose every element in the set)
if (i <= k) { dp[i][i] = 1; }</pre>
for (int i = 0; i <= n; i++) {</pre>
for (int j = 1; j <= min(i, k); j++) {</pre>
return dp[n][k]; // returns nCk modulo p
```

6.1.3 Method 2: Factorial Definition (Modular Inverses) - $O(n + \log MOD)$

nCk using Modular Inverses

```
// REF: USACO guide
const int MAXN = 1e6;
long long fac[MAXN + 1];
long long inv[MAXN + 1];
/** @return x^n modulo m in O(log p) time. */
long long exp(long long x, long long n, long long m) {
```

```
x \% = m; // note: m * m must be less than 2^63 to avoid 11 overflow
long long res = 1;
while (n > 0) {
 if (n % 2 == 1) { res = res * x % m; }
 x = x * x % m:
 n /= 2:
}
return res;
/** Precomputes n! from 0 to MAXN. */
void factorial(long long p) {
fac[0] = 1;
for (int i = 1; i <= MAXN; i++) { fac[i] = fac[i - 1] * i % p; }</pre>
/**
* Precomputes all modular inverse factorials
* from 0 to MAXN in O(n + log p) time
void inverses(long long p) {
inv[MAXN] = exp(fac[MAXN], p - 2, p);
for (int i = MAXN; i >= 1; i--) { inv[i - 1] = inv[i] * i % p; }
}
/** @return nCr mod p */
long long choose(long long n, long long r, long long p) {
return fac[n] * inv[r] % p * inv[n - r] % p;
}
int main() {
factorial();
inverses();
int n;
 cin >> n;
 for (int i = 0; i < n; i++) {</pre>
 int a, b;
 cin >> a >> b:
 cout << choose(a, b) << '\n';</pre>
}
}
```

7 Graph Theory

7.1 Shortest Path algorithms

7.1.1 Dijkstra Algorithm

$Dijkstra\ Implementation$

```
#define INF (1e18) // for int defined as 11
int n, m;
vector<vector<pair<int, int>>> adj;
vector<int> cost;
vector<int> parent;
void dijkstra(int startNode = 1) {
   priority_queue<pair<11, int>, vector<pair<11, int>>, greater<>> pq;
   cost[startNode] = 0;
   pq.emplace(0, startNode);
   while (!pq.empty()) {
       int u = pq.top().second;
       11 d = pq.top().first;
       pq.pop();
       if (d > cost[u]) continue;
       for (auto &p: adj[u]) {
          int v = p.first;
          int w = p.second;
          if (cost[v] > cost[u] + w) {
              cost[v] = cost[u] + w;
              parent[v] = u;
              pq.emplace(cost[v], v);
void run_test_case(int testNum) {
   cin >> n >> m;
```

```
adj.assign(n + 1, {});
cost.assign(n + 1, INF);
parent.assign(n + 1, -1);
while (m--) {
    // Read Edges
}
dijkstra();
if (cost[n] == INF) {
    cout << -1 << el; // not connected {Depends on you use case}</pre>
    return:
}
stack<int> ans;
for (int v = n; v != -1; v = parent[v]) ans.push(v);
while (!ans.empty()) { // printing the path
    cout << ans.top() << ' ';</pre>
    ans.pop();
}
cout << el;</pre>
```

7.1.2 Floyd Warshal Algorithm

$FloydWarshal\ Implementation$

```
int main() {
   int n, m; cin >> n >> m;
   vector <vector <int>> adj(n + 1, vector <int>> (n + 1, 2e9));
   for (int i = 0; i < n; i++) adj[i][i] = 0;

while(m--) {
    int u, v, w;
   cin >> u >> v >> w;
    adj[u][v] = min(adj[u][v], w);
   adj[v][u] = min(adj[v][u], w);
}
```

7.1.3 Bellman Ford Algorithm

BellmanFord Implementation

```
vector <vector <pair<int, int>>> &adj
vector <long long> BellmanFord(int src) {
   int n = (int)adj.size();
   vector <long long> dist(n, 2e18);
   dist[src] = 0;
   for (int it = 0; it < n-1; it++) {</pre>
       bool in = false;
       for (int i = 0; i < n; i++) { // iterate on the edges</pre>
           for (auto &[j, w] : adj[i]) {
              if (dist[j] > dist[i] + w) {
                  in = true;
                  dist[j] = dist[i] + w;
              }
       if (!in) return dist;
   for (int i = 0; i < n; i++) {</pre>
       for (auto &[j, w] : adj[i]) {
           if (dist[j] > dist[i] + w) { //negative cycle
              return vector <long long> (n, -1); // or any flag
```

```
}
}
return dist;
```

7.2 Cycle Detection

7.2.1 DFS Implementation

DFS Implementation

```
// return true with number of nodes in the cycle, either odd cycle or even
bool cycle_detection(unordered_map<int, vector<int>>> &graph, int source,
    int par, unordered_map<int,bool> vis, int c){
    if(vis[source]) return true;

    vis[source] = true;

    for(int v: graph[source]){
        if(v != par){
            c++;
            if(dfs(graph,v, source, vis, c)) return true;
        }
    }
    return false;
}
```

7.2.2 Another way for undirected graphs

Another way for undirected graphs

```
// this is true only for undirected graphs
bool dfs1(int cur, int par) {
   bool ret = false;
   vis[cur] = true;
   for (auto &i : adj[cur]) {
```

```
if (!vis[i]) ret|=dfs1(i, cur);
    else if (par != i) ret = true;
}
return ret;
}
```

7.2.3 General Way

General Way

```
// general algorithm
vector <bool> cyc;
bool dfs(int cur, int par) {
   bool ret = false;
   vis[cur] = cyc[cur] = true;
   for (auto &i : adj[cur]) {
      if (par == i) continue;
      if (!vis[i]) ret|=dfs(i, cur);
      else if (cyc[i]) ret = true;
   }
   cyc[cur] = false;
   return ret;
}
```

7.2.4 DSU Implementation

DSU Implementation

```
#include <iostream>
#include <vector>

class UnionFind {
  public:
    UnionFind(int n) {
      parent.resize(n);
      rank.resize(n, 0);
      for (int i = 0; i < n; ++i) {
         parent[i] = i;
      }
}</pre>
```

```
int find(int u) {
       if (parent[u] != u) {
           parent[u] = find(parent[u]);
       return parent[u];
   }
   void unionSets(int u, int v) {
       int rootU = find(u);
       int rootV = find(v);
       if (rootU != rootV) {
           if (rank[rootU] > rank[rootV]) {
               parent[rootV] = rootU;
           } else if (rank[rootU] < rank[rootV]) {</pre>
               parent[rootU] = rootV;
           } else {
               parent[rootV] = rootU;
               ++rank[rootU];
           }
       }
   }
private:
   std::vector<int> parent;
    std::vector<int> rank;
};
bool detectCycle(const std::vector<std::pair<int, int>>& edges, int n) {
   UnionFind uf(n);
   for (const auto& edge : edges) {
       int u = edge.first;
       int v = edge.second;
       if (uf.find(u) == uf.find(v)) {
           return true:
       uf.unionSets(u, v);
   }
   return false:
```

```
int main() {
    std::vector<std::pair<int, int>> edges = { {0, 1}, {1, 2}, {2, 3}, {3, 0} };
    int n = 4; // Number of vertices

if (detectCycle(edges, n)) {
    std::cout << "Cycle detected" << std::endl;
} else {
    std::cout << "No cycle detected" << std::endl;
}

return 0;
}</pre>
```

7.3 Algorithms

7.3.1 Heavy Light Decomposition

Basic HLD Impl

```
struct Node {
    int val;
};

const Node nullNode = {0};

const int N = 2e5 + 5, S = 1 << 19;
int n, q;
int val[N];
int sz[N], par[N], dep[N], id[N], top[N];
vector<int> adj[N];

Node st[S];

Node merge(const Node& a, const Node& b) {
    return {a.val + b.val};
}

void update(int idx, Node val) {
```

```
st[idx += n] = val;
   for (idx /= 2; idx; idx /= 2) st[idx] = merge(st[idx * 2], st[idx * 2])
       + 1]);
}
Node query(int lo, int hi) {
   Node ra = nullNode, rb = nullNode;
   for (lo += n, hi += n + 1; lo < hi; lo /= 2, hi /= 2) {
       if (lo & 1) ra = merge(ra, st[lo++]);
       if (hi & 1) rb = merge(st[--hi], rb);
   }
   return merge(ra, rb);
}
int dfs_size(const int& node, const int& parent) {
   sz[node] = 1;
   par[node] = parent;
   for (const int& ch : adj[node]) {
       if (ch == parent) continue;
       dep[ch] = dep[node] + 1;
       par[ch] = node;
       sz[node] += dfs_size(ch, node);
   return sz[node];
}
int curId = 0;
void dfs_hld(const int& cur, const int& parent, const int& curTop) {
   id[cur] = curId++;
   top[cur] = curTop;
   update(id[cur], {val[cur]});
   int heavyChild = -1, heavyMax = -1;
   for (const int& ch : adj[cur]) {
       if (ch == parent) continue;
       if (sz[ch] > heavyMax) {
          heavyMax = sz[ch];
          heavyChild = ch;
       }
   }
```

```
if (heavyChild == -1) return;
   dfs_hld(heavyChild, cur, curTop);
   for (int ch : adj[cur]) {
       if (ch == parent || ch == heavyChild) continue;
       dfs_hld(ch, cur, ch);
   }
Node path(int u, int v) {
   Node ans = nullNode;
   while (top[u] != top[v]) {
       if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
       ans = merge(ans, query(id[top[u]], id[u]));
       u = par[top[u]];
   }
   if (dep[u] > dep[v]) swap(u, v);
   ans = merge(ans, query(id[u], id[v]));
   return ans;
void init() {
   for (int i = 0; i < S; i++) st[i] = nullNode;</pre>
   dfs_size(1, 1);
   dfs_hld(1, 1, 1);
int main() {
   cin >> n >> q;
   for (int i = 1; i <= n; i++) cin >> val[i];
   int a, b;
   for (int i = 2; i <= n; i++) {
       cin >> a >> b;
       adj[a].pb(b);
       adj[b].pb(a);
   }
   init(); // <----- DON'T FORGET TO CALL THIS FUNCTION
   int type;
```

```
while (q--) {
    cin >> type;
    if (type == 1) {
        cin >> a >> b;
        val[a] = b;
        update(id[a], {val[a]});
    }
    else {
        cin >> a;
        cout << path(1, a).val << el;
    }
}</pre>
```

7.3.2 Heavy Light Decomposition with lazy SegTree

Basic HLD Impl

```
struct Node {
   int val;
};
const Node nullNode = {0};
const int N = 2e5 + 5, S = 1 << 19;
int n, q;
int val[N];
int sz[N], par[N], dep[N], id[N], top[N];
vector<int> adj[N];
Node st[S];
int lazy[S];
Node merge(const Node& a, const Node& b) {
   return {a.val + b.val};
}
void push(int idx, int 1, int r) {
   if (lazy[idx] != 0) {
       st[idx].val += lazy[idx] * (r - l + 1);
       if (1 != r) {
```

```
lazy[idx * 2] += lazy[idx];
           lazy[idx * 2 + 1] += lazy[idx];
       lazy[idx] = 0;
   }
void update_range(int lo, int hi, int l, int r, int idx, int value) {
   push(idx, l, r);
   if (lo > r || hi < 1) return;</pre>
   if (lo <= 1 && r <= hi) {</pre>
       lazy[idx] += value;
       push(idx, l, r);
       return;
   int mid = (1 + r) / 2;
   update_range(lo, hi, l, mid, idx * 2, value);
   update_range(lo, hi, mid + 1, r, idx * 2 + 1, value);
   st[idx] = merge(st[idx * 2], st[idx * 2 + 1]);
void update(int idx, Node val) {
   update_range(idx, idx, 0, n - 1, 1, val.val);
void update_range(int lo, int hi, int value) {
   update_range(lo, hi, 0, n - 1, 1, value);
Node query(int lo, int hi, int l, int r, int idx) {
   push(idx, l, r);
   if (lo > r || hi < 1) return nullNode;</pre>
   if (lo <= 1 && r <= hi) return st[idx];</pre>
   int mid = (1 + r) / 2;
   return merge(query(lo, hi, l, mid, idx * 2), query(lo, hi, mid + 1, r,
        idx * 2 + 1));
Node query(int lo, int hi) {
   return query(lo, hi, 0, n - 1, 1);
int dfs_size(const int& node, const int& parent) {
```

```
sz[node] = 1;
   par[node] = parent;
   for (const int& ch : adj[node]) {
       if (ch == parent) continue;
       dep[ch] = dep[node] + 1;
       par[ch] = node;
       sz[node] += dfs_size(ch, node);
   }
   return sz[node];
}
int curId = 0;
void dfs_hld(const int& cur, const int& parent, const int& curTop) {
   id[cur] = curId++;
   top[cur] = curTop;
   update(id[cur], {val[cur]});
   int heavyChild = -1, heavyMax = -1;
   for (const int& ch : adj[cur]) {
       if (ch == parent) continue;
       if (sz[ch] > heavyMax) {
           heavyMax = sz[ch];
          heavyChild = ch;
       }
   }
   if (heavyChild == -1) return;
   dfs_hld(heavyChild, cur, curTop);
   for (int ch : adj[cur]) {
       if (ch == parent || ch == heavyChild) continue;
       dfs_hld(ch, cur, ch);
}
int get(int u) {
   return query(id[u], id[u]).val;
}
void path(int u, int v, int val) {
   // Node ans = nullNode;
   while (top[u] != top[v]) {
```

```
if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
       // ans = merge(ans, query(id[top[u]], id[u]));
       update_range(id[top[u]], id[u], val);
       u = par[top[u]];
   }
   if (dep[u] > dep[v]) swap(u, v);
   // ans = merge(ans, query(id[u], id[v]));
   update_range(id[u], id[v], val);
   // return ans;
void init() {
   for (int i = 0; i < S; i++) st[i] = nullNode;</pre>
   memset(lazy, 0, sizeof(lazy));
   dfs_size(1, 1);
   dfs_hld(1, 1, 1);
int main(void) {
   cin >> n >> q;
   for (int i = 1; i <= n; i++) val[i] = 0;</pre>
   int a, b;
   for (int i = 2; i <= n; i++) {
       cin >> a >> b;
       adj[a].push_back(b);
       adj[b].push_back(a);
   }
   init(); // <----- DON'T FORGET TO CALL THIS FUNCTION
   int v;
   while (q--) {
       cin >> a >> b >> v;
       path(a, b, v);
   }
   for (int i = 1; i <= n; i++) {</pre>
       cout << get(i) << " ";
   }
   cout << el;</pre>
```

7.3.3 LCA functions using Binary Lifting

LCA functions using Binary Lifting

```
const int N = 2e5 + 15, M = 23;
int ancestors[N][M], depth[N], parent[N], val[N];
vector<vector<int>> adj;
//int tin[N], tout[N], timer;
void dfs_LCA(const int &node, const int &par) {
// tin[node] = timer++;
   parent[node] = par;
    ancestors[node][0] = par;
   depth[node] = depth[par] + 1;
   for (int i = 1; i < M; i++) {</pre>
       int p = ancestors[node][i - 1];
       ancestors[node][i] = ancestors[p][i - 1];
   }
   for (const int &v: adj[node]) {
       if (v == par) continue;
       dfs_LCA(v, node);
     tout[node] = timer++;
//bool is_ancestor(int u, int v) {
     return tin[u] <= tin[v] && tout[u] >= tout[v];
//}
int findKth(int u, int k) {
   if (depth[u] <= k) return -1;</pre>
   for (int i = M - 1; i >= 0; i--) {
       if (k & (1 << i)) {
           u = ancestors[u][i];
       }
   }
   return u;
}
int getLCA(int u, int v) {
   if (depth[u] < depth[v])</pre>
```

```
swap(u, v);
   u = findKth(u, depth[u] - depth[v]);
   if (u == v) return u;
   for (int i = M - 1; i >= 0; i--) {
       if (ancestors[u][i] == ancestors[v][i]) continue;
       u = ancestors[u][i];
       v = ancestors[v][i];
   }
   return ancestors[u][0];
int getDistance(int u, int v) {
   int lca = getLCA(u, v);
   return (depth[u] + depth[v]) - (2 * depth[lca]);
int dfs_accumulate(const int &node, const int &par) {
   for (const int& ch: adj[node]) {
       if (ch == par) continue;
       val[node] += dfs_accumulate(ch, node);
   return val[node];
void applyOpOnPath(const int a, const int b, const int w) {
   // adding w to each node on the path a to b
   val[a] += w;
   val[b] += w;
   int lca = getLCA(a, b);
   val[lca] -= w;
   val[parent[lca]] -= w;
int main(void) {
   int n, q;
   cin >> n >> q;
   adj.resize(n + 1);
   int u, v;
   for (int i = 2; i <= n; ++i) {
```

```
cin >> u >> v;
       adj[u].push_back(v);
       adj[v].push_back(u);
   }
   dfs_LCA(1, 1);
   parent[1] = -1;
   int w;
   for (int i = 0; i < q; ++i) {</pre>
       cin >> u >> v;
       cout << getDistance(u, v) << el;</pre>
   }
     dfs_accumulate(1, 0);
     for (int i = 1; i <= n; i++) {
         cout << val[i] << " ";
//
     cout << el;</pre>
```

7.3.4 Topological Sort

Topological Sort Using DFS

```
//REF: USACO Guide
vector<int> top_sort;
vector<vector<int>> graph;
vector<bool> visited;

void dfs(int node) {
  for (int next : graph[node]) {
    if (!visited[next]) {
     visited[next] = true;
     dfs(next);
    }
}
top_sort.push_back(node);
}
```

```
int main() {
int n, m; // The number of nodes and edges respectively
std::cin >> n >> m;
graph = vector<vector<int>>(n);
for (int i = 0; i < m; i++) {</pre>
 int a, b;
 std::cin >> a >> b;
 graph[a - 1].push_back(b - 1);
visited = vector<bool>(n);
for (int i = 0; i < n; i++) {</pre>
 if (!visited[i]) {
 visited[i] = true;
  dfs(i);
std::reverse(top_sort.begin(), top_sort.end());
vector<int> ind(n);
for (int i = 0; i < n; i++) { ind[top_sort[i]] = i; }</pre>
// Check if the topological sort is valid
bool valid = true;
for (int i = 0; i < n; i++) {</pre>
 for (int j : graph[i]) {
  if (ind[i] <= ind[i]) {</pre>
   valid = false;
   goto answer;
answer:;
if (valid) {
for (int i = 0; i < n - 1; i++) { cout << top_sort[i] + 1 << ' '; }</pre>
 cout << top_sort.back() + 1 << endl;</pre>
} else {
 cout << "IMPOSSIBLE" << endl;</pre>
```

7.3.5 Tarjan Strongly Connected Components

Tarjan SCC

```
struct TarjanSCC {
   vector<vector<int>> adj;
   vector<int> idx, lowLink;
   vector<bool> onStack;
   stack<int> st;
   int _size, components, timer;
   vector<vector<int>> sccs;
   explicit TarjanSCC(const int& n, const vector<vector<int>>& graph) :
       adj(graph), _size(n) {
       components = 0;
       timer = 0;
       onStack.assign(n + 1, false);
       idx.assign(n + 1, -1);
       lowLink.assign(n + 1, -1);
       for (int node = 1; node <= n; ++node) if (idx[node] == -1) dfs(node
          );
       ranges::sort(all(sccs), [](const vector<int>& a, const vector<int>&
           b) {
          return a.size() > b.size();
       });
   }
   void dfs(const int& node) {
       idx[node] = timer;
       lowLink[node] = timer++;
       onStack[node] = true;
       st.push(node);
       for (const int& ch : adj[node]) {
          if (idx[ch] == -1) {
              dfs(ch);
              lowLink[node] = min(lowLink[node], lowLink[ch]);
          } else if (onStack[ch]) {
```

```
lowLink[node] = min(lowLink[node], lowLink[ch]);
       }
       if (idx[node] == lowLink[node]) {
           vector<int> scc:
           int cur = -1;
           while (cur != node) {
              cur = st.top();
              st.pop();
              onStack[cur] = false;
              scc.push_back(cur);
           sccs.push_back(scc);
           components++;
   }
   int getLargestCompSize() const {
       return sccs.front().size();
   }
   vector<int> getLargestComp() const {
       return sccs.front();
   }
   vector<vector<int>> getAllComp() const {
       return this->sccs;
   }
};
void run_test_case(int testNum) {
   int n, m;
   cin >> n >> m;
   vector<vector<int>> adj(n + 1);
   int a, b;
   while (m--) {
       cin >> a >> b;
       adj[a].push_back(b);
   }
   const TarjanSCC tar(n, adj);
```

```
cout << tar.components << el;
for (const int &i : tar.getLargestComp()) {
    cout << i << " ";
}</pre>
```

8 Techniques

8.1 Coordinate Compression

```
void coordinate_compress(vector<int> &x, int start=0, int
    step=1) {
    set unique(x.begin(), x.end());
    map<int, int> valPos;

int idx=0;
    for (auto i: unique) {
       valPos[i] = start + idx * step;
       ++idx;
    }
    for(auto &i: x) i = valPos[i];
}
```

Coordinate Compression

8.2 Binary to decimal

Binary to decimal

```
// Function to convert binary to decimal
// 0(32)
int binaryToDecimal(string str)
{
   int dec_num = 0;
   int power = 0;
```

```
int n = str.length();

for(int i = n-1; i>=0; i--){
   if(str[i] == '1'){
      dec_num += (1<<power);
   }
   power++;
}

return dec_num;</pre>
```

8.3 Decimal to binary

Decimal to bianry

```
// Function that convert Decimal to binary
// 0(32)
void decToBinary(int n)
   // Size of an integer is assumed to be 32 bits
   for (int i = 31; i >= 0; i--) {
       int k = n >> i;
       if (k & 1)
           cout << "1";
       else
           cout << "0";
   }
// O(logn)
string DecimalToBinary(int num)
   string str;
     while(num){
     if(num & 1) // 1
       str+='1';
     else // 0
       str+='0':
     num>>=1; // Right Shift by 1
```

```
return str;
}
```

9 Number Theory

9.1 Divisors

9.1.1 formulas

number of divisors

```
int d(int n){
   unordered_map<int, int> factors = pf(n);
   int c = 1;
   for(const auto& factor: factors){
       c *= (factor.second+1);
   }
   return c;
}
// range Count Divisors backward thinking MAXN = 2e6
 for(int i=1; i <= n; ++i) {</pre>
   for(int j = i; j <= n; j += i) {</pre>
     numFactors[j]++;
   }
 }
int countDivisors(int n) {
 int count = 0;
 for (int i = 1; i * i <= n; ++i) {
   if (n % i == 0) {
       if (i == n / i) {
            count++; // Perfect square
       } else {
            count += 2; // Pair of divisors
   }
   return count;
```

sum of divisors

```
int s(int n){
   unordered_map<int,int> factors = pf(n);
   int sum = 1;
   for(const auto& factor: factors){
      int p = factor.first;
      int exp = factor.second;
      sum *= (pow(p,exp+1)-1)/p-1;
   }
   return sum;
}
```

9.2 Primes

prime factorization

number of co-primes with n

```
int eulerTotient(int n){
  int result = n;
```

```
for(int i = 2; SQ(i) \le n; i++){
       if(n\%i == 0){
           while (n\%i == 0) {
              n/=i;
          }
           result -= result/i;
       }
   }
   if(n > 1) result -= result/n;
   return result;
}
//Phi(n) = n * (1 - 1/P1) * (1 - 1/P2) * ...
//NOTE: summation of Euler function over divisors of n is equal to n
// using seive
void phi_generator() {
    const int MAX = 1000000;
   char primes[MAX];
   int phi[MAX];
   memset(primes, 1, sizeof(primes));
   for (int k = 0; k < MAX; ++k)
       phi[k] = 1;
   for (int i = 2; i <= MAX; ++i) {</pre>
       if (primes[i]) {
           phi[i] = i - 1; // phi(prime) = p-1
           for (int j = i * 2; j <= MAX; j += i) {</pre>
              primes[j] = 0;
              int n = j, pow = 1;
              while (n % i == 0) {
                  pow *= i;
                  n /= i;
              phi[j] *= (pow / i) * (i - 1);
           }
```

```
}
}

// phi(N!) = (N is prime ? N-1 : N) * phi((N-1)!)

ll phi_factn(int n) {
    ll ret = 1;
    for (int i = 2; i <= n; ++i)
        ret = ret * (isprime(i) ? i - 1 : i);
    return ret;
}</pre>
```

Prime Check

```
vector<bool> isPrime(MAXN, true);
void sieve() {
 isPrime[0] = isPrime[1] = false;
 for (int i=2; i * i <= isPrime.size(); ++i) {</pre>
   if(isPrime[i]) {
     for (int j = 2 * i; i <= isPrime.size(); j += i)</pre>
       prime[j] = false;
   }
 }
bool Prime(int n) {
 if(n == 2) return true;
 if(n < 2 || n % 2 == 0) return false;</pre>
 for(int i=3; i * i <= n; i += 2) {</pre>
   if(n % i == 0) return false;
 return true;
// Generate Primes
const int sz = sqrt(MAXN);
vector<int> prime;
vector<bool> vis(sz);
```

```
void pre() {
   prime.push_back(2);
   for (int j = 4; j < sz; j += 2) vis[j] = true;
   for (int i = 3; i < sz; i += 2) {</pre>
       if (vis[i]) continue;
       prime.push_back(i);
       for (int j = i * i; j < sz; j += i) vis[j] = true;</pre>
   }
}
// Preprocessing Prime Factorization of range numbers
constexpr int N = 5e6+1;
int a[N];
for(int i=2; i < N; ++i) {</pre>
   if(!a[i]) {
       for(int j=1; i*j < N; ++j) {</pre>
           for(int k=i*j; k%i==0; k/=i) a[i*j]++;
   a[i] += a[i-1];
}
```

9.3 Math

9.3.1 Vieta's Formula for a Polynomial of Degree n

Problem: Given a polynomial of degree n:

$$P(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

with roots r_1, r_2, \ldots, r_n , express the sums and products of its roots using Vieta's formulas.

Solution: Using Vieta's formulas, we can relate the coefficients of the polynomial to sums and products of its roots:

• Sum of the roots taken one at a time:

$$r_1 + r_2 + \dots + r_n = -\frac{a_{n-1}}{a_n}$$

• Sum of the products of the roots taken two at a time:

$$r_1r_2 + r_1r_3 + \dots + r_{n-1}r_n = \frac{a_{n-2}}{a_n}$$

• Sum of the products of the roots taken three at a time:

$$r_1 r_2 r_3 + r_1 r_2 r_4 + \dots + r_{n-2} r_{n-1} r_n = -\frac{a_{n-3}}{a_n}$$

- Continue this pattern until:
- Product of the roots (for even n):

$$r_1 r_2 \cdots r_n = (-1)^n \frac{a_0}{a_n}$$

Example Problem Using Vieta's Formula

Problem: Given a quadratic equation $x^2 + bx + c = 0$ with roots r_1 and r_2 , find $r_1 + r_2$ and r_1r_2 .

Solution: Using Vieta's formulas for a quadratic equation $ax^2 + bx + c = 0$:

• Sum of the roots:

$$r_1 + r_2 = -\frac{b}{a}$$

• Product of the roots:

$$r_1r_2 = \frac{c}{a}$$

For the given quadratic $x^2 + bx + c = 0$ (where a = 1):

Also: To find the roots of the quadratic equation, we can use the discriminant formula, $D = b^2 - 4ac$. The roots will then be $x_1 = \frac{b - \sqrt{D}}{2}$ and $x_2 = \frac{b + \sqrt{D}}{2}$.

- $r_1 + r_2 = -b$
- $r_1r_2=c$

Example Vieta's Formula for Cubic Equation

When considering a cubic equation in the form of $f(x) = ax^3 + bx^2 + cx + d$, Vieta's formula states that if the equation f(x) = 0 has roots r_1, r_2 , and r_3 , then:

- $r_1 + r_2 + r_3 = -\frac{b}{a}$
- $r_1 r_2 + r_2 r_3 + r_3 r_1 = \frac{c}{a}$
- $r_1r_2r_3 = -\frac{d}{a}$

9.4 Phi Function

- Count integers i < n such that gcd(i, n) = 1
- $gcd(a, b) = 1 \Rightarrow$ then coprimes: gcd(5, 7), gcd(4, 9)
- gcd(prime, i) = 1 for i < prime
- $\varphi(10) = 4 \Rightarrow 1, 3, 7, 9$
- $\varphi(5) = 4 \Rightarrow 1, 2, 3, 4 \dots \varphi(\text{prime}) = \text{prime} 1$
- If a, b, c are pairwise coprimes, then

$$\varphi(a \cdot b \cdot c) = \varphi(a) \cdot \varphi(b) \cdot \varphi(c)$$

• If k > 1

$$\varphi(p^k) = p^k - p^{k-1} = p^{k-1}(p-1) = p^k \left(1 - \frac{1}{p}\right)$$

9.5 Euler's Totient Numbers

Online Sequence

• $\varphi(1) = \varphi(2) = 1, \ \varphi(5) = 4$

- $\varphi(n)$ is even for n > 2
- $\sqrt{n} \le \varphi(n) \le n \sqrt{n}$: Except 2, 6
- $\varphi(n^k) = n^{k-1} \cdot \varphi(n)$
- $n = \sum \varphi(d_i)$ where d are the divisors of n

9.6 Möbius Function

Möbius Function

```
int mobius(int n) {
   int p = 0;
   // Handling 2 separately
   if (n\%2 == 0){
       n = n/2;
       p++;
       // If 2^2 also divides N
       if (n % 2 == 0)
          return 0;
   }
   // Check for all other prime factors
   for (int i = 3; i <= sqrt(n); i = i+2) {</pre>
       if (n\%i == 0){
           n = n/i;
           p++;
           if (n % i == 0) return 0;
   }
void mobius_generator() const {
 const int MAX = 1000000;
  char prime[MAX + 1];
```

```
for (ll i = 2; i <= MAX; i++)
    moebius[i] = -1, prime[i] = 1;

for (ll i = 2; i <= MAX; ++i)
    if (prime[i]) {
        moebius[i] = 1;

        for (ll j = 2 * i; j <= MAX; j += i)
            prime[j] = 0, moebius[j] = j % (i * i) == 0 ? -moebius[j] : 1;
        }
}

// Mobius Inclusion Exclusion
// Count triples gcd(a, b, c) = 1
int n = 4;
ll sum = n * n * n;
for (ll i = 2; i <= n; ++i)
        sum -= moebius[i] * (n / i) * (n / i) * (n / i);</pre>
```

9.7 Möbius and Inclusion Exclusion

Count the triples (a, b, c) such that $a, b, c \le n$, and gcd(a, b, c) = 1

- Reverse thinking, total (# triples gcd > 1)
- How many triples with gcd multiple of 2: $(n/2)^3$
- How many triples with gcd multiple of 3: $(n/3)^3$
- and 4? Ignore any numbers of internal duplicate primes
- and 6? already computed in 2, 3. Remove it: $-(n/6)^3$

9.8 Totient and Möbius Connection

Sum over divisors d of n

$$\sum_{d} d\mu \left(\frac{n}{d}\right) = \varphi(n)$$

9.9 Lagrange's four-square theorem

Lagrange's four-square theorem states that, every positive integer can be expressed as the sum of the squares of four integers.

Lagrange's four-square theorem

```
#include <cmath>
bool isPerfectSquare(int n) {
   int rt = sqrt(n);
   return rt * rt == n;
// Function to compute the minimum number of perfect squares
int numSquares(int n) {
   // Case 1:
   if (isPerfectSquare(n)) return 1;
   // Case 2: Check if n can be expressed as the sum of two perfect
   for (int i = 1; i * i <= n; i++) {
       int d = n - i * i;
       if (isPerfectSquare(d)) return 2;
   // Case 3:
   // If n can be reduced to the form 4^a * (8b + 7), then it requires 4
   int m = n:
   while (m \% 4 == 0) m /= 4;
   if (m % 8 == 7) return 4;
   // Case 4:
   return 3;
```

10 Geometry

10.1 Linearity

10.1.1 co-linear points

check if two points are co-linear

```
bool co_linear(int x1, int y1, int x2, int y2, int x3, int y3){
  int area = x1*(y2-y3) + x2*(y3-y1) + x3*(y1-y2);
  return area == 0;
}
```

10.2 Polygons

10.2.1 Polygon formation

check if can form polygon with given angle

```
bool possible(double angle){
  if(angle <= 0 || angle >= 180) return false;

double sides = 360.0/(180.0-angle);

return (sides == static_cast<int>(sides) && sides >= 3);
}
```

10.2.2 Polygon Area

area of any polygon with x vertices

```
double shoelace(vector<pair<double, double>>> points) {
  double leftSum = 0.0;
  double rightSum = 0.0;

for (int i = 0; i < points.size(); ++i) {
  int j = (i + 1) % points.size();
  leftSum += points[i].first * points[j].second;
  rightSum += points[j].first * points[i].second;</pre>
```

```
return 0.5 * abs(leftSum - rightSum);
}
```

10.3 Intersections

10.3.1 Rectangle

intersection area between 2 rectangles

```
struct Rectangle {
    int x1, y1; // Bottom-left corner
    int x2, y2; // Top-right corner
};

int intersectionArea(const Rectangle& rect1, const Rectangle& rect2){
    int x_left = max(rect1.x1, rect2.x1);
    int y_bottom = max(rect1.y1, rect2.y1);
    int x_right = min(rect1.x2, rect2.x2);
    int y_top = min(rect1.y2, rect2.y2);

    int intersection_width = x_right - x_left;
    int intersection_height = y_top - y_bottom;

    if (intersection_width > 0 && intersection_height > 0) {
        return intersection_width * intersection_height;
    }

    return 0;
}
```

10.3.2 Circle

intersection area between 2 circles

```
double area(int x0, int y0, int r0, int x1, int y1, int r1){
   const double PI = 3.14159265358979323846;
```

```
double rr0 = r0 * r0;
double rr1 = r1 * r1;
double d = sqrt((x1 - x0) * (x1 - x0) + (y1 - y0) * (y1 - y0));
if(d >= r0 + r1){
    return 0;
}
double phiAngle = (rr0 + (d * d) - rr1) / (2 * r0 * d);
double phi = acos(phiAngle) * 2;
double thetaAngle = (rr1 + (d * d) - rr0) / (2 * r1 * d);
double theta = acos(thetaAngle) * 2;
double area1 = 0.5 * theta * rr1 - 0.5 * rr1 * sin(theta);
double area2 = 0.5 * phi * rr0 - 0.5 * rr0 * sin(phi);
return area1+area2;
}
```

10.3.3 Triangle

intersection area between 2 triangles

```
struct Point {
   double x, y;
};
typedef vector<Point> Polygon;
bool inside(const Point &p, const pair<Point, Point> &edge) {
   Point a = edge.first, b = edge.second;
   return (b.x - a.x) * (p.y - a.y) - (b.y - a.y) * (p.x - a.x) >= 0;
}
Point compute_intersection(const Point &p1, const Point &p2, const pair<
   Point, Point> &edge) {
   Point a = edge.first, b = edge.second;
   double A1 = b.y - a.y, B1 = a.x - b.x, C1 = b.x * a.y - a.x * b.y;
   double A2 = p2.y - p1.y, B2 = p1.x - p2.x, C2 = p2.x * p1.y - p1.x *
       p2.v;
   double det = A1 * B2 - A2 * B1;
   if (det == 0) return {numeric_limits<double>::quiet_NaN(),
       numeric_limits<double>::quiet_NaN()}; // parallel
   double x = (B1 * C2 - B2 * C1) / det;
   double y = (A2 * C1 - A1 * C2) / det;
   return {x, y};
```

```
Polygon HodgmanClip(const Polygon &subjectPolygon, const Polygon &
   clipPolygon) {
   Polygon outputList = subjectPolygon;
   for (size_t i = 0; i < clipPolygon.size(); ++i) {</pre>
       Point a = clipPolygon[i];
       Point b = clipPolygon[(i + 1) % clipPolygon.size()];
       Polygon inputList = outputList;
       outputList.clear();
       if (inputList.empty()) break;
       Point s = inputList.back();
       for (const auto &e : inputList) {
           if (inside(e, {a, b})) {
              if (!inside(s, {a, b})) {
                  Point intersection = compute_intersection(s, e, {a, b});
                  if (!std::isnan(intersection.x)) {
                      outputList.push_back(intersection);
                  }
              }
              outputList.push_back(e);
           } else if (inside(s, {a, b})) {
              Point intersection = compute_intersection(s, e, {a, b});
              if (!std::isnan(intersection.x)) {
                  outputList.push_back(intersection);
              }
           s = e;
   }
   return outputList;
// area of polygon
double shoelace(const Polygon &vertices) {
   double area = 0;
   for (size_t i = 0; i < vertices.size(); ++i) {</pre>
       Point p1 = vertices[i];
       Point p2 = vertices[(i + 1) % vertices.size()];
       area += p1.x * p2.y - p2.x * p1.y;
   }
   return fabs(area) / 2.0;
```

10.3.4 Rectangle & Circle

struct Point {

intersection area between Rectangle and Circle

```
double x, y;
};
struct Circle {
   Point center;
   double radius;
};
struct Rectangle {
   Point bottomLeft, topRight;
};
const double PI = acos(-1.0);
bool point_inside_circle(const Point &p, const Circle &c) {
   double dx = p.x - c.center.x;
   double dy = p.y - c.center.y;
   return (dx * dx + dy * dy) <= (c.radius * c.radius);</pre>
}
double rect_area(const Rectangle &rect) {
   return (rect.topRight.x - rect.bottomLeft.x) * (rect.topRight.y - rect
        .bottomLeft.v);
}
bool rect_in_circle(const Circle &circle, const Rectangle &rect) {
   vector<Point> corners = {
       rect.bottomLeft,
       {rect.topRight.x, rect.bottomLeft.y},
       rect.topRight,
       {rect.bottomLeft.x, rect.topRight.y}
   };
   for (const auto &corner : corners) {
       if (!isPointInsideCircle(corner. circle)) {
```

```
return false;
   }
   return true;
bool circle_in_rect(const Circle &circle, const Rectangle &rect) {
return circle.center.x - circle.radius >= rect.bottomLeft.x &&
          circle.center.x + circle.radius <= rect.topRight.x &&
         circle.center.y - circle.radius >= rect.bottomLeft.y &&
          circle.center.y + circle.radius <= rect.topRight.y;</pre>
double intersection(const Circle &circle, const Rectangle &rect) {
   if (rect_in_circle(circle, rect)) {
       return rect_area(rect);
   }
   if (circle_in_rect(circle, rect)) {
       return PI * circle.radius * circle.radius;
   }
   double intersectionArea = 0.0:
   double dx1 = max(rect.bottomLeft.x, circle.center.x - circle.radius);
   double dx2 = min(rect.topRight.x, circle.center.x + circle.radius);
   double dy1 = max(rect.bottomLeft.y, circle.center.y - circle.radius);
   double dy2 = min(rect.topRight.y, circle.center.y + circle.radius);
   for (double x = dx1; x < dx2; x += 0.001) {
       for (double y = dy1; y < dy2; y += 0.001) {
          Point p = \{x, y\};
          if (point_inside_circle(p, circle)) {
              intersectionArea += 0.001 * 0.001;
       }
   }
   return intersectionArea:
```

10.3.5 Line & Circle

intersection points between Line and Circle

```
struct Point {
   double x, y;
};
struct Circle {
   Point center;
   double radius;
};
struct Line {
   double slope;
   double intercept;
};
vector<Point> intersect_points(const Circle &circle, const Line &line) {
   vector<Point> intersections;
   // solve y = mx + b with (x - h)^2 + (y - k)^2 = r^2
   double h = circle.center.x;
   double k = circle.center.v;
   double r = circle.radius;
   double m = line.slope;
   double b = line.intercept;
   // Quadratic coefficients
   double A = 1 + m * m;
   double B = 2 * (m * b - m * k - h);
   double C = k * k - r * r + h * h - 2 * b * k + b * b;
   // Discriminant
   double discriminant = B * B - 4 * A * C;
   if (discriminant < 0) {</pre>
       // No intersection
       return intersections;
   } else if (discriminant == 0) {
       // One intersection (tangent line)
       double x = -B / (2 * A);
       double y = m * x + b;
```

```
intersections.push_back({x, y});
} else {
    // Two intersections
    double sqrtDiscriminant = sqrt(discriminant);
    double x1 = (-B + sqrtDiscriminant) / (2 * A);
    double y1 = m * x1 + b;
    double x2 = (-B - sqrtDiscriminant) / (2 * A);
    double y2 = m * x2 + b;
    intersections.push_back({x1, y1});
    intersections.push_back({x2, y2});
}

return intersections;
```

10.4 Linear Operations

10.5 2D Points

Basic Operations

```
#include<bits/stdc++.h>
using namespace std;

#define ftype double

struct Point2D {
   ftype x, y;
   Point2D() {}
   Point2D(ftype x, ftype y): x(x), y(y) {}
   Point2D& operator+=(const Point2D &p2) {
        x += p2.x;
        y += p2.y;
        return *this;
   }
   Point2D& operator-=(const Point2D &p2) {
        x -= p2.x;
        y -= p2.y;
        return *this;
   }
}
```

```
Point2D& operator*=(const Point2D p2) {
       x *= p2.x;
       y *= p2.y;
       return *this;
   }
   Point2D& operator/=(const Point2D p2) {
       x \neq p2.x;
       y /= p2.y;
       return *this;
   Point2D& operator+(const Point2D p2) {
       return Point2D(*this) += p2;
   Point2D& operator-(const Point2D p2) {
       return Point2D(*this) -= p2;
   }
   Point2D& operator/(const Point2D p2) {
       return Point2D(*this) /= p2;
   Point2D& operator*(const Point2D p2) {
       return Point2D(*this) *= p2;
   }
};
Point2D operator*(ftype a, Point2D b) {
       return a * b;
}
int main() {
   Point2D p1(1, 2);
   Point2D p2(3, 4);
   Point2D p3 = p1 * p2;
}
```

10.6 3D Points

Basic Operations

```
#include<bits/stdc++.h>
#define ftype double
struct Point3D {
   ftype x, y, z;
   Point3D() {}
   Point3D(ftype x, ftype y, ftype z): x(x), y(y), z(z) {}
   Point3D& operator+=(const Point3D &p2) {
       x += p2.x;
       y += p2.y;
       z += p2.z;
       return *this;
   Point3D& operator-=(const Point3D &p2) {
       x -= p2.x;
       y = p2.y;
       z = p2.z;
       return *this;
   Point3D& operator*=(const Point3D p2) {
       x *= p2.x;
       y *= p2.y;
       z *= p2.z;
       return *this;
   Point3D& operator/=(const Point3D p2) {
       x \neq p2.x;
       y /= p2.y;
       z \neq p2.z;
       return *this;
   }
   Point3D& operator+(const Point3D p2) {
       return Point3D(*this) += p2;
   Point3D& operator-(const Point3D p2) {
       return Point3D(*this) -= p2;
```

```
}
Point3D& operator/(const Point3D p2) {
    return Point3D(*this) /= p2;
}
Point3D& operator*(const Point3D p2) {
    return Point3D(*this) *= p2;
}

};

Point3D operator*(ftype a, Point3D b) {
    return a * b;
}

int main() {
    Point3D p1(1, 2);
    Point3D p2(3, 4);

    Point3D p3 = p1 * p2;
}
```

11 Miscellaneous

11.1 Faster implementations

11.1.1 hashes

custom hash

```
#define safe hash unordered_map<type, type, custom_hash> // same for
    gp_hash_table
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        // http://xorshift.di.unimi.it/splitmix64.c
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }
```

```
size_t operator()(uint64_t x) const {
    static const uint64_t FIXED_RANDOM = chrono::steady_clock::now().
        time_since_epoch().count();
    return splitmix64(x + FIXED_RANDOM);
};
};
```

gb hash table

```
//policy based ds (faster hash table)
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
gp_hash_table<int, int> table;
```

11.1.2 Binary Search the value

$nearest\ sqrt$

```
long long my_sqrt(long long a)
{
    long long l=0,r=5000000001;
    while(r-l>1)
    {
        long long mid=(l+r)/2;
        if(1ll*mid*mid<=a)l=mid;
        else r=mid;
    }
    return l;
}</pre>
```