Competitive Programming Library

Too bad to be Accepted 2023/2024

1 Templates	Contents					4.2.1 Prefix Sum (L, R) intervals	10
1 Templates 3 4.3.1 Find duplicate 11 1.1 Setup 3 4.4.4 Sorting Algorithms 11 1.1.1 IO Manipulation 3 4.4.1 Radix Sort 11 1.2 MOD Template 3 4.4.2 Counting Sort 12 1.3 Macros 3 4.4.2 Counting Sort 12 1.4 Grid Navigation 4 5.1 Strings 12 1.5 Integer 128 4 5.1 Strings 12 2 Dynamic Programming 4 5.1 Strings 12 2.1 Some dp patterns 4 5.2 Range Queries 13 2.2 DP solutions 6 5.2.1 Segment Tree 13 2.2.2 Max Subarray sum (Kadane's Algorithm) 6 5.2.2 Lazy Propegation 14 2.2.2 Maximum Subarray Alternating Sum 6 5.2.2 Enyiek Tree 15 2.2.3 Count number of DISTINCT ordered ways to produce coins sums to x 7 5.2.6 Sparse Table 16 3 Bit Manipulation 8 6 6 5.2.2 D BIT 16 5.2.2 Lorgest common subsequence between 2 Strings 7 6 Counting Sort 17 5.2.6 Sparse Table 16 <						4.2.2 Find subarrays intervals that sum to K Using Map	11
1.1 Setup .					4.3	Ad-hoc	11
1.1.1 IO Manipulation 3 1.1.2 GCC Compiler Optimization (Vectorization) 3 3 1.1.2 GCC Compiler Optimization (Vectorization) 3 3 4.4.2 Counting Sort 12 3 3 3 3 3 3 3 3 3	1	-	- 1			4.3.1 Find duplicate	11
1.1.2 GCC Compiler Optimization (Vectorization) 3 1.2 MOD Template 3 1.3 Macros 3 1.4 Grid Navigation 4 1.5 Integer 128 5 Integer 128 Integer			3		4.4	Sorting Algorithms	11
1.2 MOD Template 3 1.3 Macros 3 1.4 Grid Navigation 3 1.5 Integer 128 5 Data Structures 12 5 1.1 Trie (Prefix Tree) 13 5 1.1 Trie (Prefix Tree) 13 12 12 12 12 13 13 13		1	3			4.4.1 Radix Sort	11
1.3 Macros Macros Integer 128 5 Data Structures 12 Strings 13 Strings 12 Strings 13 Strings 12 Strings 13 Strings 12 Strings 13 Strings 13 Strings 12 Strings 13 Strings 13 Strings 13 Strings 12 Strings 13 Strings 13 Strings 14 Strings 15 Strings <		,	ა ე			4.4.2 Counting Sort	12
1.4 Grid Navigation 4 1.5 Integer 128 12 5.1 Strings 12 5.1 Strings 12 5.1 Trie (Prefix Tree) 12 5.1 Trie (Prefix Tree) 12 5.1 Trie (Prefix Tree) 12 5.2 Range Queries 13 7.1 Trie (Prefix Tree) 12 5.2 Range Queries 13 7.2 Zary Propegation 14 5.2 Zary Propegation 14 5.2 Zary Propegation 14 5.2 Zary Propegation 15 5.2 Zary Propegation 16 Zary Propegation 15 5.2 Zary Propegation 16 Zary Propegation 17 Zary Propegation 17 Zary Propegation 16 Zary Propegation 16 Zary Propegation 17 Zary Propegation 16 Zary Propegation 17 Zary Propegation 17 Zary Propegation 17 Zary Propegation 17 Zary		•	3				
1.5 Integer 128			$\frac{3}{4}$	5			12
2 Dynamic Programming 4 5.2 Range Queries 12			4		5.1		
2.1 Some dp patterns 4 5.2.1 Segment Tree 13 2.2 DP solutions 6 5.2.2 Lazy Propegation 14 2.2.1 Max Subarray sum (Kadane's Algorithm) 6 5.2.3 Femwick Tree 15 2.2.2 Maximum Subarray Alternating Sum 6 5.2.4 Femwick UpdateRange 15 2.2.3 Count number of DISTINCT ordered ways to produce coins sums to x 7 5.2.5 2D BIT 16 2.2.4 Min absolute difference between 2 elements from (L, R) (DP Ranges) 7 5.2.6 Sparse Table 16 3 Bit Manipulation 8 6.1 nCr 18 3.1 Subset Operations 8 6.1.1 Fast nCr 18 4 Algorithms 9 6.1.2 Method 1: Pascal's Triangle (Dynamic Programming) 4 Algorithms 9 6.1.3 Method 2: Factorial Definition (Modular Inverses) -							
2.2 DP solutions	2	Dynamic Programming	4		5.2		
2.2.1 Max Subarray sum (Kadane's Algorithm) 6 5.2.3 Fenwick Tree 15 2.2.2 Maximum Subarray Alternating Sum 6 5.2.4 Fenwick UpdateRange 15 2.2.3 Count number of DISTINCT ordered ways to produce coins sums to x 7 7 2.2.4 Min absolute difference between 2 elements from (L, R) (DP Ranges) 7 7 5.2.6 Sparse Table 16 2.2.5 Longest common subsequence between 2 Strings 7 6 Counting Principles 18 3.1 Subset Operations 8 6.1 nCr 18 6.1.1 Fast nCr 18 6.1.2 Method 1: Pascal's Triangle (Dynamic Programming) - O(n²) 18 6.1.3 Method 2: Factorial Definition (Modular Inverses) -			4			~	
2.2.2 Maximum Subarray Alternating Sum			6				
2.2.3 Count number of DISTINCT ordered ways to produce coins sums to x		· · · · · · · · · · · · · · · · · · ·	- 1				
Coins sums to x Coins sums		v e	6				
2.2.4 Min absolute difference between 2 elements from (L, R) (DP Ranges)		· -	7				
R) (DP Ranges)			'			-	
2.2.5 Longest common subsequence between 2 Strings 7 6 Counting Principles 18 3 Bit Manipulation 3.1 8 6.1 nCr			7		5.3	Ordered Set	17
3 Bit Manipulation 8 6.1 nCr 18 3.1 Subset Operations 8 6.1.1 Fast nCr 18 4 Algorithms 9 - O(n²) 18 4.1 MO 9 6.1.3 Method 2: Factorial Definition (Modular Inverses) 18		- /	7	6	Con	unting Principles	18
3.1 Subset Operations	0						
$6.1.2$ Method 1: Pascal's Triangle (Dynamic Programming) 4 Algorithms 9 $-\mathcal{O}(n^2)$		•			0		
4 Algorithms 9 $-\mathcal{O}(n^2)$		5.1 Subset Operations	$^{\circ}$				
4.1 MO	4	Algorithms				_	18
			9				
4.2 Intervals		4.2 Intervals	10			$\mathcal{O}(n + \log MOD)$	18

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7	Gra	19	11	Mis	cell	
	7.1	Shortest Path algorithms	19		11.1	Fas
		7.1.1 Dijkstra Algorithm	19			11.
		7.1.2 Floyd Warshal Algorithm	20			11.
		7.1.3 Bellman Ford Algorithm	20			
	7.2	Cycle Detection	21			
		7.2.1 DFS Implementation	21			
		7.2.2 Another way for undirected graphs	21			
		7.2.3 General Way	21			
		7.2.4 DSU Implementation	21			
	7.3	Algorithms	22			
		7.3.1 Heavy Light Decomposition	22			
		7.3.2 Heavy Light Decomposition with lazy SegTree	24			
		7.3.3 LCA functions using Binary Lifting	25			
		7.3.4 Topological Sort	27			
8	Tecl	27				
	8.1	Coordinate Compression	27			
	8.2	Binary to decimal	28			
	8.3	Decimal to binary	28			
9	Nur	nber Theory	28			
	9.1	Divisors	28			
		9.1.1 formulas	28			
	9.2	Primes	29			
	9.3	Math	30			
		9.3.1 Vieta's Formula for a Polynomial of Degree n	30			
10	Geo	metry	31			
		Linearity	31			
		10.1.1 co-linear points	31			
	10.2	polygons	31			
		10.2.1 Polygon formation	31			
		10.2.2 Rectangle Intersection	31			

11 Miscellaneous										
	11.1	Faster	implementations	32						
		11.1.1	hashes	32						
		11.1.2	Binary Search the value	32						

1 Templates

1.1 Setup

1.1.1 IO Manipulation

Input/Output

```
#include <bits/stdc++.h>
freopen("input.txt", "r", stdin);
freopen("output.txt", "w", stdout);

#define fastI0  \
   ios_base::sync_with_stdio(false), cin.tie(nullptr), cout.tie(nullptr);
```

1.1.2 GCC Compiler Optimization (Vectorization)

GCC Opt

```
// Ref: USACO guide
// will make GCC auto-vectorize for loops and optimizes floating points
   better (assumes associativity and turns off denormals).
#pragma GCC optimize ("Ofast")
// can double performance of vectorized code, but causes crashes on old
   machines.
#pragma GCC target ("avx,avx2,fma")

// slows down run time but throws a Runtime Error if an overflow occured
#pragma GCC optimize("trapv")
```

1.2 MOD Template

```
constexpr int MOD = 1e9+7; // must be a prime number
int add(int a, int b) {
   int res = a+b;
   if(res >= MOD) return res -= MOD;
}
```

```
int sub(int a, int b) {
   int res = a-b;
   if(res < 0) return res += MOD;
}
int power(int a, int e) {
   int res = 1;
   while(e) {if(e & 1) res = res * a % MOD; a = a * a % MOD;
   e >>= 1;}
   return res;
}
int inverse(int a) {
   return power(a, MOD-2);
}
int div(int a, int b) {
   return a * inverse(b) % MOD;
}
```

MOD Template

1.3 Macros

Macros

```
#define getBit(n, k) (n >> k)
#define ON(n, idx) (n | (111 << idx))
#define OFF(n, idx) (n & ~(111 << idx))
#define toggle(n, idx) ((n) ^ (111<<(idx)))
#define gray(n) (n ^ (n >> 1))
#define bitCount(x) (__builtin_popcountl1(x))
#define clz(x) (__builtin_clzl1(x))
#define ctz(x) (__builtin_ctzl1(x))
#define uniq(x) x.resize(unique(x.begin(), x.end())-x.begin());

#define angle(a) (atan2((a).imag(), (a).real()))
//#define vec(a, b) ((b)-(a))
#define same(v1, v2) (dp(vec(v1,v2),vec(v1,v2)) < EPS)
#define dotProduct(a, b) ((conj(a)*(b)).real()) // a*b cos(T), if zero -> prep
```

```
#define crossProduct(a, b) ((conj(a)*(b)).imag()) // a*b sin(T), if zero
    -> parallel
//#define length(a) (hypot((a).imag(), (a).real()))
#define normalize(a) ((a)/length(a))
#define rotateO(v, ang) ((v)*exp(point(O,ang)))
#define rotateA(p, ang, about) (rotateO(vec(about,p),ang)+about)
#define reflectO(v, m) (conj((v)/(m))*(m))
#define ceil_i(a, b) (((ll)(a)+(ll)(b-1))/(ll)(b))
#define floor_i(a, b) (a/b)
#define round_i(a, b) ((a+(b/2))/b) // if a>0
#define round_m(a, b) ((a-(b/2))/b) // if a<0
#define round_multiple(n, m) round_i(n,m)*m // round to multiple if
    specified element

const double PI = acos(-1.0);</pre>
```

1.4 Grid Navigation

Grid Nav

```
// knight moves on a chess board
int dx[] = { -2, -1, 1, 2, -2, -1, 1, 2 };
int dy[] = { -1, -2, -2, -1, 1, 2, 2, 1 };

// Grid up, down, right, left (Moves for Chess Rook)
int dx[4] = {1, -1, 0, 0};
int dy[4] = {0, 0, 1, -1};

// Grid cell all neighbours
const int dx[8] = {1, 0, -1, 0, 1, 1, -1, -1}
const int dy[8] = {0, 1, 0, -1, -1, 1, -1, 1};

// Grid Diagonal (Moves for Chess Bishop)
int dx[] = {1, 1, -1, -1};
int dy[] = {1, -1, 1, -1};
```

1.5 Integer 128

```
i128
typedef __int128 i128;
__int128 read() {
   \_int128 x = 0, f = 1;
   char ch = getchar();
   while (ch < '0' || ch > '9') {
       if (ch == '-') f = -1;
       ch = getchar();
   }
   while (ch >= '0' && ch <= '9') {
       x = x * 10 + ch - '0';
       ch = getchar();
   }
   return x * f;
void print(__int128 x) {
   if (x < 0) {
       putchar('-');
       x = -x;
   if (x > 9) print(x / 10);
   putchar(x % 10 + '0');
bool cmp(\_int128 x, \_int128 y) { return x > y; }
```

2 Dynamic Programming

2.1 Some dp patterns

Maximumu/Minimum path cost

```
const int MAX = 21;
int grid[MAX][MAX];
int mem[MAX][MAX];
int n = 20;
bool valid(int r, int c){
  return r >= 0 && r < n && c >= 0 && c < n;</pre>
```

```
int maxPathSum(int r, int c){
   if(!valid(r,c)){
      return 0;
   }

   if(r == n-1 && c == n-1){
      return mem[r][c] = grid[r][c];
   }

   // available moves
   int path1 = maxPathSum(r+1,c);
   int path2 = maxPathSum(r,c+1);

   return grid[r][c] + max(path1,path2);
}
```

add operators between numbers to get max prod/sum

```
// put +, -, between sequence of numbers such that the sum is divisible by
    k, and maximum as possible
const int MAX = 21;
long long mem[MAX][MAX];
const int n = 20;
int k = 4; // example
int v[20];
int fix(int a){
 return (a % k + k) % k;
long long tryAll(int pos, int mod){
   long long &ret = mem[pos][mod];
   if(ret != -1){
       return ret;
   if(pos == n){
       return ret = mod == 0;
   if(tryAll(pos+1,fix(mod + v[pos])) || tryAll(pos+1,fix(mod-v[pos]))){
       return ret = 1;
   return ret = 0;
```

pick choices with no two similar consecutive choices

```
// pick minimum of choinces costs with no two similar consecutive choices
const int choices = 4:
const int n = 20;
int MAX = n;
int mem[MAX][choices];
const int 00 = 1e6+1;
int minCost(int pos, int lastChoice){
   if(pos == n){
       return 0; // invalid move
   int &ret = mem[pos][lastChoice];
   if(ret != -1){
       return ret;
   }
   ret = 00: // want to minimze
   // let choices are 0, 1, 2
   if(lastChoice != 0){
       ret = min(ret, minCost(pos+1,0));
   }
   if(lastChoice != 1){
       ret = min(ret, minCost(pos+1,1));
   if(lastChoice != 2){
       ret = min(ret, minCost(pos+1,2));
   return ret;
```

sum S and max/min Product

```
int maxK;

ll mem[21][101]; // k, and s

// You are given an integer s and an integer k. Find k positive integers a1, a2, ..., ak

// such that their sum is equal to s and their product is the maximal possible. Return their product.
```

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```
11 maxProd(int k, int rem)
if(k == maxK){
 // base case
 if(rem == 0)
  return 1;
 return 0;
}
if(rem == 0) // invalid case
 return 0;
11 &ret = mem[k][rem];
if(ret != -1)
 return ret;
ret = 0;
for (int i = 1; i <= rem; ++i) {</pre>
 ll sol = maxProd(k+1, rem - i) * i;
 ret = max(ret, sol);
}
return ret;
}
```

2.2 DP solutions

2.2.1 Max Subarray sum (Kadane's Algorithm)

$Max\ Subarray\ sum$

```
int maxSubarraySum(vector<int>& arr, int len) {
   int ans = INT_MIN, cur = 0;

for (int i = 0; i < len; i++) {
    cur = cur + arr[i];
   if (ans < cur)
       ans = cur;

   if (cur < 0)</pre>
```

```
cur = 0;
}
return ans;
```

2.2.2 Maximum Subarray Alternating Sum

Maximum Subarray Alternating Sum

```
/* REF: GeeksForGeeks
Input: arr[] = \{-4, -10, 3, 5\}
Output: 9
Explanation: Subarray {arr[0], arr[2]} = {-4, -10, 3}. Therefore, the sum
     of this subarray is 9.
int maxSubarraySumALT(vector<int>& a, int len) {
   int ans = INT_MIN, cur = 0;
   for (int i = 0; i < len; i++) {</pre>
       if (i % 2 == 0)
           cur = max(cur + a[i], a[i]);
           cur = max(cur - a[i], -a[i]);
       ans = max(ans, cur);
   }
   cur = 0;
   for (int i = 0; i < len; i++) {</pre>
       if (i % 2 == 1)
           cur = max(cur + a[i], a[i]);
       else
           cur = max(cur - a[i], -a[i]);
       ans = max(ans, cur);
   }
   return ans;
```

sums to x

Count distinct

```
/*
For example, if the coins are \{2,3,5\} and the desired sum is 9, there
    are 3 ways:
2+2+5
3+3+3
2+2+2+3
*/
int n, x;
cin >> n >> x;
vector<int> coins(n);
read(coins);
vector dp(x + 1, 0);
dp[0] = 1;
for (int i = 0; i < n; ++i) {</pre>
for (int j = coins[i]; j <= x; ++j) {</pre>
 dp[j] = add(dp[j], dp[j - coins[i]]);
}
}
cout << dp[x] << el;
```

Min absolute difference between 2 elements from (L, R) (DP 2.2.4Ranges)

Min absolute difference

```
const int N = 1e4 + 1;
int dp[N][N];
int n;
cin >> n;
vector<int> a(n);
read(a);
```

```
Count number of DISTINCT ordered ways to produce coins for (int i = 0; i < n; ++i) dp[i][i] = 1e6; // INF, you can't take the
                                                                         element with it self
                                                                     for (int i = 1; i < n; ++i) dp[i - 1][i] = abs(a[i] - a[i - 1]);
                                                                     for (int len = 3; len <= n; ++len) {</pre>
                                                                     for (int l = 0, r = len - 1; r < n; ++1, ++r) {
                                                                      dp[1][r] = min(dp[1][r - 1], dp[1 + 1][r]);
                                                                      dp[1][r] = min(dp[1][r], abs(a[1] - a[r]));
                                                                     int q;
                                                                     cin >> q;
                                                                     while (q--) {
                                                                     int 1, r;
                                                                     cin >> 1 >> r;
                                                                     --1, --r;
                                                                     cout << dp[l][r] << el;
```

2.2.5 Longest common subsequence between 2 Strings

$$dp[i][j] = \begin{cases} max(dp[i-1][j], dp[i][j-1]) & \text{if } A_i \neq B_j \\ dp[i-1][j-1] + 1 & \text{if } A_i = B_j \end{cases}$$

LIS 2 Strings

```
// REF: USACO guide
int longestCommonSubsequence(string a, string b) {
int dp[a.size()][b.size()];
for (int i = 0; i < a.size(); i++) { fill(dp[i], dp[i] + b.size(), 0); }</pre>
for (int i = 0; i < a.size(); i++) {</pre>
 if (a[i] == b[0]) dp[i][0] = 1;
 if (i != 0) dp[i][0] = max(dp[i][0], dp[i - 1][0]);
for (int i = 0; i < b.size(); i++) {</pre>
 if (a[0] == b[i]) dp[0][i] = 1;
 if (i != 0) dp[0][i] = max(dp[0][i], dp[0][i - 1]);
for (int i = 1; i < a.size(); i++) {</pre>
for (int j = 1; j < b.size(); j++) {</pre>
```

```
if (a[i] == b[j]) {
   dp[i][j] = dp[i - 1][j - 1] + 1;
   } else {
   dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
   }
}
return dp[a.size() - 1][b.size() - 1];
}
```

3 Bit Manipulation

3.1 Subset Operations

count subsets with give sum

```
int countDistinctSubsetsWithSum(vector<int>& arr, int n, int k) {
    // Count distinct subsets of array arr that sum up to k
    vector<int> dp(k + 1, 0);
    dp[0] = 1;
    for (int i = 0; i < n; ++i) {
        for (int j = k; j >= arr[i]; --j) {
            dp[j] += dp[j - arr[i]];
        }
    }
    return dp[k]; // Number of distinct subsets with sum k
}
```

max xor of any subset of elements in the array

```
int maximalSubsetXOR(vector<int>& arr, int n) {
    // Find the maximum XOR of any subset of elements in array arr
    int maxXor = 0;
    for (int mask = 0; mask < (1 << n); ++mask) {
        int xorSum = 0;
        for (int i = 0; i < n; ++i) {
            if (mask & (1 << i)) {
                  xorSum ^= arr[i];
            }
        }
}</pre>
```

```
maxXor = max(maxXor, xorSum);
}
return maxXor;
```

min xor of any subset

```
int minimumSubsetXOR(vector<int>& arr, int n) {
    // Find the minimum XOR of any pair of elements in array arr
    int minSubsetXor = INT_MAX;
    for (int mask = 0; mask < (1 << n); ++mask) {
        int xorSum = 0;
        for (int i = 0; i < n; ++i) {
            if (mask & (1 << i)) {
                xorSum ^= arr[i];
            }
        }
        minSubsetXor = min(minSubsetXor, xorSum);
    }
    return minSubsetXor;
}</pre>
```

$subset\ generation$

```
void subsetGeneration(int x, int n) {
    // Generate all non-empty subsets of a set represented by an integer x
    for (int subset = x; subset > 0; subset = (subset - 1) & x) {
        // Process subset
        cout << subset << endl;
    }
}</pre>
```

check if subset of elements in the array sum up to k

```
void subsetSumCheck(vector<int>& arr, int n, int k) {
    // Check if a subset of elements in array arr sums up to k
    for (int subset = 0; subset < (1 << n); ++subset) {
        int sum = 0;
        for (int i = 0; i < n; ++i) {
            if (subset & (1 << i)) {</pre>
```

```
sum += arr[i];
}

if (sum == k) {
    // Found subset with sum k
    cout << "Subset with sum " << k << ": " << subset << endl;
}
}</pre>
```

max subset sum mod m

```
int subsetWithMaxSumModuloM(vector<int>& arr, int n, int m) {
   // Find the maximum subset sum modulo m
   vector<int> dp(m, -1);
   dp[0] = 0;
   int currentMod = 0;
   for (int i = 0; i < n; ++i) {
       currentMod = (currentMod + arr[i]) % m;
       for (int j = 0; j < m; ++j) {
          if (dp[j] != -1) {
              dp[(j + currentMod) % m] = max(dp[(j + currentMod) % m], dp
                  [j] + arr[i]);
          }
       }
       dp[currentMod] = max(dp[currentMod], arr[i]);
   }
   return dp[0]; // Maximum subset sum modulo m
```

iterate over all supersets represented by x

4 Algorithms

4.1 MO

MO Algorithm

```
// MO
           -> O(N+Q SQRT(N)) <= 10^5
const int N = 1e5+5, M = 1e5+5;
int n, m;
int nums[N], q_ans[M];
struct query {
   int idx, block_idx, l, r;
   query() = default;
   query(int _1, int _r, int _idx) {
       idx = _idx;
       r = _r - 1;
       1 = _1 - 1;
       block_idx = _1 / sqrt(n);
   bool operator <(const query & y) const {</pre>
       if(y.block_idx == block_idx) return r < y.r;</pre>
       return block_idx < y.block_idx;</pre>
   }
};
int freq[N], ans;
void add(int idx) {
   freq[nums[idx]]++;
   if (freq[nums[idx]] == 2) ans++;
void remove(int idx) {
   freq[nums[idx]]--;
   if (freq[nums[idx]] == 1) ans--;
cin >> n >> m;
for (int i = 0; i < n; ++i) cin >> nums[i];
```

```
vector<query> Query(m);
for (int i = 0; i < m; ++i) {</pre>
   int 1, r; cin >> 1 >> r;
   Query[i] = query(1, r, i);
}
sort(Query.begin(), Query.end());
int 10 = 1, r0 = 0;
for (int i = 0; i < m; ++i) {
   while (10 < Query[i].1) remove(10++);</pre>
   while (10 > Query[i].1) add(--10);
   while (r0 < Query[i].r) add(++r0);
   while (r0 > Query[i].r) remove(r0--);
   q_ans[Query[i].idx] = ans;
}
for (int i = 0; i < m; ++i) {
    cout << q_ans[i] << '\n';
}
```

4.2 Intervals

4.2.1 Prefix Sum (L, R) intervals

Prefix Sum (L, R) intervals

```
// NOTE: works fine with small n or with large memory
int main() {
   int n, k;
   cin >> n >> k;

   vector<int>> a(n + 1);
   vector<vector<int>> rangesPrefix(n + 1, vector<int>(n + 1, 0));
   for (int i = 1; i <= n; ++i)
        cin >> a[i];

   int l = 1, r = 1, sum = 0;
   // validate your intervals
   // here the intervals are the ones that have a sum of k
   while (r <= n) {
        sum += a[r];
   }
}</pre>
```

```
while (sum > k) {
       sum -= a[1];
       ++1;
   }
   while (1 \le r \&\& a[1] == 0) {
       if (sum != k)
           break;
       rangesPrefix[r][1]++;
       ++1;
   }
   if (sum == k) {
       rangesPrefix[r][1]++;
    ++r;
}
// prefix sum the columns
for (int i = 1; i <= n; ++i) {
   for (int j = n - 1; j \ge 0; --j) {
       rangesPrefix[i][j] += rangesPrefix[i][j + 1];
   }
}
// prefix sum the rows
for (int i = 0; i <= n; ++i) {
   for (int j = 1; j <= n; ++j) {
       rangesPrefix[j][i] += rangesPrefix[j - 1][i];
   }
}
int q; cin >> q;
while (q--) {
   // answer the number of intervals (X, Y) X <= Y that are included
        between L. R
    cout << rangesPrefix[r][l] - rangesPrefix[l - 1][l] << el;</pre>
```

```
}
```

4.2.2 Find subarrays intervals that sum to K Using Map

Find subarray intervals that sum to K Using Map

```
int n, k;
cin >> n >> k;
vector < int > a(n + 1);
vector<pair<int, int>> rng;
for (int i = 1; i <= n; ++i)
    cin >> a[i];
map<int, set<int>> prev;
int currSum = 0;
for (int i = 1; i <= n; ++i) {
    currSum += a[i];
   if (currSum == k) {
       rng.push_back({1, i});
   if (prev.find(currSum - k) != prev.end()) {
       for (auto &j : prev[currSum - k]) {
           rng.push_back(\{j + 1, i\});
       }
    prev[currSum].insert(i);
```

4.3 Ad-hoc

4.3.1 Find duplicate

Find duplicate using XOR

```
int findDuplicate(int arr[] , int n)
{
   int answer=0;
    //XOR all the elements with 0
```

```
for(int i=0; i<n; i++){
    answer=answer^arr[i];
}
    //XOR all the elements with no from 1 to n
    // i.e answer^0 = answer
for(int i=1; i<n; i++){
    answer=answer^i;
}
return answer;</pre>
```

4.4 Sorting Algorithms

4.4.1 Radix Sort

Radix Sort

```
// O(n + b), where n is the number of elements and b is the base of the
    number system
// A function to do counting sort of arr[] according to the digit
    represented by exp.
void countingSort(vector<int>& arr, int exp) {
   int n = arr.size();
   vector<int> output(n); // output array
   int count[10] = {0};
   // Store count of occurrences in count[]
   for (int i = 0; i < n; i++)</pre>
       count[(arr[i] / exp) % 10]++;
   // Change count[i] so that count[i] now contains the actual
   // position of this digit in output[]
   for (int i = 1; i < 10; i++)</pre>
       count[i] += count[i - 1];
   // Build the output array
   for (int i = n - 1; i >= 0; i--) {
       output[count[(arr[i] / exp) % 10] - 1] = arr[i];
       count[(arr[i] / exp) % 10]--;
   }
   // Copy the output array to arr[], so that arr now
```

4.4.2 Counting Sort

Counting Sort

```
// O(N+M), where N and M are the size of inputArray[] and countArray[]
// The main function that sorts arr[] of size n using Counting Sort
void countingSort(vector<int>& arr) {
   int maxElement = *max_element(arr.begin(), arr.end());
   int minElement = *min_element(arr.begin(), arr.end());
   int range = maxElement - minElement + 1;
   vector<int> count(range), output(arr.size());
   for (int i = 0; i < arr.size(); i++)</pre>
       count[arr[i] - minElement]++;
   for (int i = 1; i < count.size(); i++)</pre>
       count[i] += count[i - 1];
   for (int i = arr.size() - 1; i >= 0; i--) {
       output[count[arr[i] - minElement] - 1] = arr[i];
       count[arr[i] - minElement]--;
   }
   for (int i = 0; i < arr.size(); i++)</pre>
       arr[i] = output[i];
}
```

5 Data Structures

5.1 Strings

5.1.1 Trie (Prefix Tree)

Basic Implementation

```
#define MAX_CHAR 26
struct TrieNode {
   TrieNode *pTrieNode[MAX_CHAR]{};
   bool isWord;
   TrieNode() {
       isWord = false;
       fill(pTrieNode, pTrieNode + 26, (TrieNode *) NULL);
   }
   virtual ~TrieNode() = default;
};
class Trie {
private:
   TrieNode *root;
public:
   Trie() {
       root = new TrieNode();
   }
   virtual ~Trie() = default;
   TrieNode *getTrieNode() {
       return this->root;
   }
   void insert(const string &word) {
       TrieNode *current = root;
       for (char c: word) {
           int i = c - 'a';
           if (current->pTrieNode[i] == nullptr)
```

```
current->pTrieNode[i] = new TrieNode();
           current = current->pTrieNode[i];
       current->isWord = true;
   }
   bool search(const string &word) {
       TrieNode *current = root;
       int ch = 0;
       for (char c: word) {
           ch = c - 'a';
          if (current->pTrieNode[ch] == nullptr)
              return false;
           current = current->pTrieNode[ch];
       return current->isWord;
   }
   bool startsWith(const string &prefix) {
       TrieNode *current = root;
       int ch = 0;
       for (char c: prefix) {
           ch = c - 'a';
           if (current->pTrieNode[ch] == nullptr)
              return false;
           current = current->pTrieNode[ch];
       }
       return true;
};
```

5.2 Range Queries

5.2.1 Segment Tree

Basic Implementation

```
struct Node {
    long long val;
};
struct SegTree {
```

```
private:
   const Node NEUTRAL = {INT_MIN};
   static Node merge(const Node& x1, const Node& x2) {
       return {x1.val + x2.val}:
   }
   void set(const int& idx, const int& val, int x, int lx, int rx) {
       if (rx - lx == 1) return void(values[x].val = val);
       int mid = (rx + lx) / 2;
       if (idx < mid)</pre>
           set(idx, val, 2 * x + 1, lx, mid);
       else
           set(idx, val, 2 * x + 2, mid, rx);
       values[x] = merge(values[2 * x + 1], values[2 * x + 2]);
   }
   Node query(const int& 1, const int& r, int x, int lx, int rx) {
       if (lx >= r || l >= rx) return NEUTRAL;
       if (lx >= 1 && rx <= r) return values[x];</pre>
       int mid = (rx + lx) / 2:
       return merge(query(1, r, 2 * x + 1, 1x, mid), query(1, r, 2 * x + 1
           2, mid, rx));
   void build(vector<int> &a, int x, int lx, int rx) {
       if (rx - lx == 1) {
           if (lx < a.size()) {</pre>
              values[x].val = a[lx];
           }
           return;
       int m = (1x + rx) / 2;
       build(a, 2 * x + 1, 1x, m);
       build(a, 2 * x + 2, m, rx);
       values[x] = merge(values[2 * x + 1], values[2 * x + 2]);
   }
    void assign_range(int 1, int r, int node, int lx, int rx, int time,
        int val) {
```

```
if (lx > r || l > rx) return;
       if (lx >= 1 && rx <= r) {
          lazy[node] = {time, val};
           return;
       }
       int mid = (1x+rx) / 2;
       assign_range(1, r, 2*node+1, lx, mid, time, val);
       assign_range(l, r, 2*node+2, mid+1, rx, time, val);
   }
   pair<int, int> point_query(int lx, int rx, int node, int idx) {
       if(rx == lx) return lazy[node];
       int mid = (1x+rx) / 2;
       if(idx <= mid) {</pre>
           auto x = point_query(lx, mid, 2*node+1, idx);
           if(x.first > lazy[node].first) return x;
           return lazy[node];
       }
       auto x = point_query(mid+1, rx, 2*node+2, idx);
       if(x.first > lazy[node].first) return x;
       return lazy[node];
   }
public:
   int size{};
   vector<Node> values;
   void build(vector<int> &a) {
       build(a, 0, 0, size);
   }
   void init(int _size) {
       size = 1;
       while (size < _size) size *= 2;</pre>
       values.assign(2 * size, NEUTRAL);
   }
   void set(int idx, int val) {
       set(idx, val, 0, 0, size);
   }
   Node query(const int& 1, const int& r) {
```

```
return query(1, r, 0, 0, size);
};
```

5.2.2 Lazy Propegation

Lazy Propegation

```
struct SegTree {
private:
   void propegate(int lx, int rx, int node) {
       if(!lazy[node]) return;
       if(lx != rx) {
           lazy[2*node+1] = lazy[node];
           lazy[2*node+2] = lazy[node];
       values[node] = lazy[node] * (rx - lx + 1);
       lazy[node] = 0;
   }
   // assign val in range [1, r]
   void update_range(int 1, int r, int node, int lx, int rx, int val,
       bool f) {
       propegate(lx, rx, node);
       if (1x > r \mid | 1 > rx) return:
       if (lx >= 1 && rx <= r) {
           lazy[node] = val;
          propegate(lx, rx, node);
           return;
       }
       int mid = (lx+rx) / 2;
       update_range(l, r, 2*node+1, lx, mid, val, f);
       update_range(l, r, 2*node+2, mid+1, rx, val, f);
       values[node] = values[2*node+1] + values[2*node+2];
   }
   // get sum in range [1, r]
   int range_query(int 1, int r, int lx, int rx, int node) {
       propegate(lx, rx, node);
       if (lx > r \mid | l > rx) return 0;
       if (lx >= 1 && rx <= r) return values[node];</pre>
```

```
int mid = (1x+rx) / 2;
       return range_query(1, r, lx, mid, 2*node+1) + range_query(1, r, mid
           +1, rx, 2*node+2);
   }
public:
   int size{};
   vector<int> values, lazy;
   void init(int _size) {
       size = 1;
       while (size < _size) size *= 2;</pre>
       values.assign(2 * size, 0);
       lazy.assign(2 * size, 0);
   }
   void update_range(int 1, int r, int v, bool f) {
       update_range(1, r, 0, 0, size-1, v, f);
   }
   int range_query(int 1, int r) {
       return range_query(1, r, 0, size-1, 0);
   }
};
```

5.2.3 Fenwick Tree

Fenwick Tree

```
struct Fenwick {
    // One Based
    vector<int> tree;

explicit Fenwick(int n) {tree.assign(n + 5, {});}

// Computes the prefix sum from [1, i], O(log(n))
int query(int i) {
    int res = 0;
    while (i > 0) {
        res += tree[i];
        i &= ~(i & -i);
    }
}
```

```
return res;
   }
   int query(int 1, int r) {
       return query(r) - query(1-1);
   }
   // Get the value at index i
   int get(int i) {
       return query(i, i);
   }
   // Add 'v' to index 'i', O(log(n))
   void update(int i, int v) {
       while (i < tree.size()) {</pre>
           tree[i] += v;
           i += (i & -i);
       }
   }
   // Update range, Point query
   // To get(k) do prefix sum [1, k] and in insert update_range(i, i, a[i
       1)
   void update_range(int 1, int r, int v) {
       update(1, v);
       update(r+1, -v);
   }
};
```

5.2.4 Fenwick UpdateRange

$BIT\ UpdateRange$

```
struct BITUpdateRange {
private:
   int n;
   vector<int> B1, B2;

  void add(vector<int> &b, int idx, int x) {
     while (idx <= n) {
        b[idx] += x;
        idx += idx & -idx;
   }</pre>
```

```
}
   int sum(vector<int> &b, int idx) {
       int total = 0;
       while (idx > 0) {
          total += b[idx];
           idx &= (idx & -idx);
       }
       return total;
   }
   int prefix(int idx) {
       return sum(B1, idx) * idx - sum(B2, idx);
   }
public:
   explicit BITUpdateRange(int n) : n(n) {
       B1.assign(n + 1, {});
       B2.assign(n + 1, {});
   }
   void update(int 1, int r, int x) {
       add(B1, 1, x);
       add(B1, r + 1, -x);
       add(B2, 1, x * (1 - 1));
       add(B2, r + 1, -x * r);
   }
   int query(int i) {
       return prefix(i) - prefix(i - 1);
   }
   int query(int 1, int r) {
       return prefix(r) - prefix(l - 1);
   }
};
```

5.2.5 2D BIT

2D BIT

```
struct BIT2D {
   int n, m;
```

```
vector<vector<int>> bit;
   BIT2D(int n, int m) : n(n), m(m) {
       bit.assign(n + 2, vector<int>(m + 2));
   }
   void update(int x, int y, int val) {
       for (; x <= n; x += x & -x) {
           for (int i = y; i <= m; i += i & -i) {
              bit[x][i] += val;
       }
   }
   int prefix(int x, int y) {
       int res = 0;
       for (; x > 0; x &= (x & -x)) {
          for (int i = y; i > 0; i &= (i & -i)) {
              res += bit[x][i];
       }
       return res;
   }
   int query(int sx, int sy, int ex, int ey) {
       int ans = 0;
       ans += prefix(ex, ey);
       ans -= prefix(ex, sy - 1);
       ans -= prefix(sx - 1, ey);
       ans += prefix(sx - 1, sy - 1);
       return ans;
   }
};
```

5.2.6 Sparse Table

Impl with the index

```
// storing the index also
struct SNode {
   int val;
   int index;
};
```

```
class SparseTable {
private:
   vector<vector<SNode>> table;
   function<SNode(const SNode&, const SNode&)> merge;
   static SNode StaticMerge(const SNode& a, const SNode& b) {
       return a.val < b.val ? a : b;</pre>
   }
public:
   explicit SparseTable(const vector<int>& arr, const function<SNode(</pre>
       const SNode&, const SNode&)>& mergeFunc = StaticMerge) {
       int n = static_cast<int>(arr.size());
       int log_n = static_cast<int>(log2(n)) + 1;
       this->merge = mergeFunc;
       table.resize(n, vector<SNode>(log_n));
       for (int i = 0; i < n; i++) {</pre>
           table[i][0] = {arr[i], i};
       }
       for (int j = 1; (1 << j) <= n; j++) {
           for (int i = 0; i + (1 << j) <= n; i++) {
               table[i][j] = mergeFunc(table[i][j - 1], table[i + (1 << (j
                   - 1))][j - 1]);
           }
   SNode query(int left, int right) {
       int j = static_cast<int>(log2(right - left + 1));
       return merge(table[left][j], table[right - (1 << j) + 1][j]);</pre>
   }
   // query in O(log(n)) if its could't apply to Sparse Table directly
   T query_log(int 1, int r){
     int len = r - 1 + 1;
     T ans:
     for(int i = 0; 1 <= r; i++){</pre>
         if (len & (1 << i)){
             ans = merge(ans, table[i][1]);
```

```
l+= (1 << i);
}
}

int main(void) {
  int n;
  cin >> n;
  vector<int> arr(n);
  for (auto& element : arr) cin >> element;

SparseTable minSt(arr, [](const SNode& a, const SNode& b) -> SNode {
    return a.val < b.val ? a : b;
});

SparseTable maxSt(arr, [](const SNode& a, const SNode& b) -> SNode {
    return a.val > b.val ? a : b;
});
}
```

5.3 Ordered Set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace __gnu_pbds;

template < typename T >
    using ordered_set = tree < T, null_type, less < T >, rb_tree_tag,
        tree_order_statistics_node_update >;

void erase_set(ordered_set & os, int v) {
        // Number of elements less than v
        int rank = os.order_of_key(v);

        auto it = os.find_by_order(rank);
        os.erase(it);
}
```

Ordered Set

Counting Principles

6.1nCr

$$C(n,k) = \frac{n!}{(n-k)!k!} = \frac{n*(n-1)*(n-2)*...*(n-k+1)}{k!}$$

6.1.1 Fast nCr

Fast nCr

```
int nCr(const int& n, const int& r) {
   double res = 1;
   for (int i = 1; i <= r; ++i)
       res = res * (n - r + i) / i;
   return (int)(res + 0.01);
}
```

Method 1: Pascal's Triangle (Dynamic Programming) - $\mathcal{O}(n^2)$

nCk using dp

```
// REF: USACO guide
/** Creturn nCk mod p using dynamic programming */
int binomial(int n, int k, int p) {
// dp[i][j] stores iCj
vector<vector<int>> dp(n + 1, vector<int>(k + 1, 0));
// base cases described above
for (int i = 0; i <= n; i++) {</pre>
  * i choose 0 is always 1 since there is exactly one way
```

```
* to choose O elements from a set of i elements
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  * (don't choose anything)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             dp[i][0] = 1;
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  * i choose i is always 1 since there is exactly one way
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  * to choose i elements from a set of i elements
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                 * (choose every element in the set)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        if (i <= k) { dp[i][i] = 1; }</pre>
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     for (int i = 0; i <= n; i++) {</pre>
C(n,k) = \frac{n*(n-1)*(n-2)*\dots*(n-k+1)}{1*2*3*\dots*k} = \prod_{i=0}^{k-1} \frac{n-i}{i+1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ uses the recurrence relation above } \frac{1}{1} = \prod_{i=0}^{k-1} (n-i)(i+1)^{-1} / \text{ use } \frac{1}{1} = \prod_{i=0}^{k-1} 
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     return dp[n][k]; // returns nCk modulo p
```

6.1.3 Method 2: Factorial Definition (Modular Inverses) - O(n + $\log MOD$

nCk using Modular Inverses

```
// REF: USACO guide
const int MAXN = 1e6;
long long fac[MAXN + 1];
long long inv[MAXN + 1];
/** @return x^n modulo m in O(log p) time. */
long long exp(long long x, long long n, long long m) {
x \%= m; // note: m * m must be less than 2^63 to avoid ll overflow
long long res = 1;
while (n > 0) {
 if (n % 2 == 1) { res = res * x % m; }
```

```
x = x * x % m;
 n /= 2:
return res;
/** Precomputes n! from 0 to MAXN. */
void factorial(long long p) {
fac[0] = 1;
for (int i = 1; i <= MAXN; i++) { fac[i] = fac[i - 1] * i % p; }</pre>
/**
* Precomputes all modular inverse factorials
* from 0 to MAXN in O(n + log p) time
void inverses(long long p) {
inv[MAXN] = exp(fac[MAXN], p - 2, p);
for (int i = MAXN; i >= 1; i--) { inv[i - 1] = inv[i] * i % p; }
}
/** @return nCr mod p */
long long choose(long long n, long long r, long long p) {
return fac[n] * inv[r] % p * inv[n - r] % p;
int main() {
factorial();
inverses();
int n;
 cin >> n;
 for (int i = 0; i < n; i++) {</pre>
 int a, b;
 cin >> a >> b;
 cout << choose(a, b) << '\n';</pre>
}
}
```

7 Graph Theory

7.1 Shortest Path algorithms

7.1.1 Dijkstra Algorithm

$Dijkstra\ Implementation$

```
#define INF (1e18) // for int defined as 11
int n, m;
vector<vector<pair<int, int>>> adj;
vector<int> cost;
vector<int> parent;
void dijkstra(int startNode = 1) {
   priority_queue<pair<11, int>, vector<pair<11, int>>, greater<>> pq;
   cost[startNode] = 0;
   pq.emplace(0, startNode);
   while (!pq.empty()) {
       int u = pq.top().second;
       11 d = pq.top().first;
       pq.pop();
       if (d > cost[u]) continue;
       for (auto &p: adj[u]) {
          int v = p.first;
          int w = p.second;
           if (cost[v] > cost[u] + w) {
              cost[v] = cost[u] + w;
              parent[v] = u;
              pq.emplace(cost[v], v);
void run_test_case(int testNum) {
   cin >> n >> m;
   adj.assign(n + 1, {});
```

```
cost.assign(n + 1, INF);
parent.assign(n + 1, -1);
while (m--) {
    // Read Edges
}
dijkstra();
if (cost[n] == INF) {
    cout << -1 << el; // not connected {Depends on you use case}</pre>
    return;
}
stack<int> ans;
for (int v = n; v != -1; v = parent[v]) ans.push(v);
while (!ans.empty()) { // printing the path
    cout << ans.top() << ' ';</pre>
    ans.pop();
}
cout << el;</pre>
```

7.1.2 Floyd Warshal Algorithm

$FloydWarshal\ Implementation$

```
int main() {
    int n, m; cin >> n >> m;
    vector <vector <int>> adj(n + 1, vector <int>> (n + 1, 2e9));
    for (int i = 0; i < n; i++) adj[i][i] = 0;

while(m--) {
        int u, v, w;
        cin >> u >> v >> w;
        adj[u][v] = min(adj[u][v], w);
        adj[v][u] = min(adj[v][u], w);
    }

for (int mid = 1; mid <= n; mid++) {
        for (int start = 1; start <= n; start++) {
            for (int end = 1; end <= n; end++) {</pre>
```

7.1.3 Bellman Ford Algorithm

BellmanFord Implementation

```
vector <vector <pair<int, int>>> &adj
vector <long long> BellmanFord(int src) {
   int n = (int)adj.size();
   vector <long long> dist(n, 2e18);
   dist[src] = 0;
   for (int it = 0; it < n-1; it++) {</pre>
       bool in = false;
       for (int i = 0; i < n; i++) { // iterate on the edges</pre>
           for (auto &[j, w] : adj[i]) {
              if (dist[j] > dist[i] + w) {
                  in = true;
                  dist[j] = dist[i] + w;
              }
           }
       }
       if (!in) return dist;
   }
   for (int i = 0; i < n; i++) {</pre>
       for (auto &[j, w] : adj[i]) {
           if (dist[j] > dist[i] + w) { //negative cycle
              return vector <long long> (n, -1); // or any flag
       }
   }
   return dist;
```

7.2 Cycle Detection

7.2.1 DFS Implementation

DFS Implementation

```
// return true with number of nodes in the cycle, either odd cycle or even
bool cycle_detection(unordered_map<int, vector<int>> &graph, int source,
    int par, unordered_map<int,bool> vis, int c){
    if(vis[source]) return true;

    vis[source] = true;

    for(int v: graph[source]){
        if(v != par){
            c++;
            if(dfs(graph,v, source, vis, c)) return true;
        }
    }
    return false;
}
```

7.2.2 Another way for undirected graphs

Another way for undirected graphs

```
// this is true only for undirected graphs
bool dfs1(int cur, int par) {
   bool ret = false;
   vis[cur] = true;
   for (auto &i : adj[cur]) {
      if (!vis[i]) ret|=dfs1(i, cur);
      else if (par != i) ret = true;
   }
   return ret;
}
```

7.2.3 General Way

General Way

```
// general algorithm
vector <bool> cyc;
bool dfs(int cur, int par) {
   bool ret = false;
   vis[cur] = cyc[cur] = true;
   for (auto &i : adj[cur]) {
      if (par == i) continue;
      if (!vis[i]) ret|=dfs(i, cur);
      else if (cyc[i]) ret = true;
   }
   cyc[cur] = false;
   return ret;
}
```

7.2.4 DSU Implementation

DSU Implementation

```
#include <iostream>
#include <vector>
class UnionFind {
public:
   UnionFind(int n) {
       parent.resize(n);
       rank.resize(n, 0);
       for (int i = 0; i < n; ++i) {</pre>
           parent[i] = i;
       }
   }
   int find(int u) {
       if (parent[u] != u) {
           parent[u] = find(parent[u]);
       }
       return parent[u];
   }
   void unionSets(int u, int v) {
```

```
int rootU = find(u);
       int rootV = find(v);
       if (rootU != rootV) {
           if (rank[rootU] > rank[rootV]) {
              parent[rootV] = rootU;
           } else if (rank[rootU] < rank[rootV]) {</pre>
              parent[rootU] = rootV;
           } else {
              parent[rootV] = rootU;
              ++rank[rootU];
          }
   }
private:
   std::vector<int> parent;
   std::vector<int> rank;
};
bool detectCycle(const std::vector<std::pair<int, int>>& edges, int n) {
   UnionFind uf(n);
   for (const auto& edge : edges) {
       int u = edge.first;
       int v = edge.second;
       if (uf.find(u) == uf.find(v)) {
           return true;
       }
       uf.unionSets(u, v);
   }
   return false;
int main() {
   std::vector<std::pair<int, int>> edges = { {0, 1}, {1, 2}, {2, 3}, {3,
        0} }:
   int n = 4; // Number of vertices
   if (detectCycle(edges, n)) {
       std::cout << "Cycle detected" << std::endl;</pre>
   } else {
```

```
std::cout << "No cycle detected" << std::endl;
}
return 0;</pre>
```

7.3 Algorithms

7.3.1 Heavy Light Decomposition

Basic HLD Impl

```
struct Node {
   int val;
};
const Node nullNode = {0};
const int N = 2e5 + 5, S = 1 << 19;
int n, q;
int val[N];
int sz[N], par[N], dep[N], id[N], top[N];
vector<int> adj[N];
Node st[S];
Node merge(const Node& a, const Node& b) {
   return {a.val + b.val};
void update(int idx, Node val) {
   st[idx += n] = val;
   for (idx /= 2; idx; idx /= 2) st[idx] = merge(st[idx * 2], st[idx * 2]
       + 1]);
Node query(int lo, int hi) {
   Node ra = nullNode, rb = nullNode;
   for (lo += n, hi += n + 1; lo < hi; lo /= 2, hi /= 2) {
       if (lo & 1) ra = merge(ra, st[lo++]);
       if (hi & 1) rb = merge(st[--hi], rb);
```

```
}
   return merge(ra, rb);
}
int dfs_size(const int& node, const int& parent) {
   sz[node] = 1;
   par[node] = parent;
   for (const int& ch : adj[node]) {
       if (ch == parent) continue;
       dep[ch] = dep[node] + 1;
       par[ch] = node;
       sz[node] += dfs_size(ch, node);
   }
   return sz[node];
}
int curId = 0;
void dfs_hld(const int& cur, const int& parent, const int& curTop) {
   id[cur] = curId++;
   top[cur] = curTop;
   update(id[cur], {val[cur]});
   int heavyChild = -1, heavyMax = -1;
   for (const int& ch : adj[cur]) {
       if (ch == parent) continue;
       if (sz[ch] > heavyMax) {
          heavyMax = sz[ch];
           heavyChild = ch;
       }
   }
   if (heavyChild == -1) return;
   dfs_hld(heavyChild, cur, curTop);
   for (int ch : adj[cur]) {
       if (ch == parent || ch == heavyChild) continue;
       dfs_hld(ch, cur, ch);
   }
}
Node path(int u, int v) {
   Node ans = nullNode;
```

```
while (top[u] != top[v]) {
       if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
       ans = merge(ans, query(id[top[u]], id[u]));
       u = par[top[u]];
   }
   if (dep[u] > dep[v]) swap(u, v);
   ans = merge(ans, query(id[u], id[v]));
   return ans;
void init() {
   for (int i = 0; i < S; i++) st[i] = nullNode;</pre>
   dfs_size(1, 1);
   dfs_hld(1, 1, 1);
int main() {
   cin >> n >> q;
   for (int i = 1; i <= n; i++) cin >> val[i];
   int a, b;
   for (int i = 2; i <= n; i++) {
       cin >> a >> b;
       adj[a].pb(b);
       adj[b].pb(a);
   }
   init(); // <----- DON'T FORGET TO CALL THIS FUNCTION</pre>
   int type;
   while (q--) {
       cin >> type;
       if (type == 1) {
           cin >> a >> b;
           val[a] = b;
           update(id[a], {val[a]});
       }
       else {
           cin >> a;
           cout << path(1, a).val << el;</pre>
   }
```

7.3.2 Heavy Light Decomposition with lazy SegTree

Basic HLD Impl

```
struct Node {
   int val;
};
const Node nullNode = {0};
const int N = 2e5 + 5, S = 1 << 19;
int n, q;
int val[N];
int sz[N], par[N], dep[N], id[N], top[N];
vector<int> adj[N];
Node st[S];
int lazy[S];
Node merge(const Node& a, const Node& b) {
   return {a.val + b.val};
}
void push(int idx, int 1, int r) {
   if (lazy[idx] != 0) {
       st[idx].val += lazy[idx] * (r - l + 1);
       if (1 != r) {
           lazy[idx * 2] += lazy[idx];
           lazy[idx * 2 + 1] += lazy[idx];
       lazy[idx] = 0;
   }
}
void update_range(int lo, int hi, int l, int r, int idx, int value) {
   push(idx, l, r);
   if (lo > r || hi < 1) return;
   if (lo <= 1 && r <= hi) {
       lazy[idx] += value;
       push(idx, l, r);
       return;
   }
```

```
int mid = (1 + r) / 2:
   update_range(lo, hi, l, mid, idx * 2, value);
   update_range(lo, hi, mid + 1, r, idx * 2 + 1, value);
   st[idx] = merge(st[idx * 2], st[idx * 2 + 1]);
void update(int idx, Node val) {
   update_range(idx, idx, 0, n - 1, 1, val.val);
void update_range(int lo, int hi, int value) {
   update_range(lo, hi, 0, n - 1, 1, value);
Node query(int lo, int hi, int l, int r, int idx) {
   push(idx, l, r);
   if (lo > r || hi < 1) return nullNode;</pre>
   if (lo <= 1 && r <= hi) return st[idx];</pre>
   int mid = (1 + r) / 2;
   return merge(query(lo, hi, l, mid, idx * 2), query(lo, hi, mid + 1, r,
        idx * 2 + 1));
Node query(int lo, int hi) {
   return query(lo, hi, 0, n - 1, 1);
int dfs_size(const int& node, const int& parent) {
   sz[node] = 1;
   par[node] = parent;
   for (const int& ch : adj[node]) {
       if (ch == parent) continue;
       dep[ch] = dep[node] + 1;
       par[ch] = node;
       sz[node] += dfs_size(ch, node);
   }
   return sz[node];
int curId = 0;
void dfs_hld(const int& cur, const int& parent, const int& curTop) {
   id[cur] = curId++:
   top[cur] = curTop;
```

```
update(id[cur], {val[cur]});
   int heavyChild = -1, heavyMax = -1;
   for (const int& ch : adj[cur]) {
       if (ch == parent) continue;
       if (sz[ch] > heavyMax) {
           heavyMax = sz[ch];
           heavyChild = ch;
       }
   }
   if (heavyChild == -1) return;
   dfs_hld(heavyChild, cur, curTop);
   for (int ch : adj[cur]) {
       if (ch == parent || ch == heavyChild) continue;
       dfs_hld(ch, cur, ch);
}
int get(int u) {
   return query(id[u], id[u]).val;
}
void path(int u, int v, int val) {
   // Node ans = nullNode;
   while (top[u] != top[v]) {
       if (dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
       // ans = merge(ans, query(id[top[u]], id[u]));
       update_range(id[top[u]], id[u], val);
       u = par[top[u]];
   }
   if (dep[u] > dep[v]) swap(u, v);
   // ans = merge(ans, query(id[u], id[v]));
   update_range(id[u], id[v], val);
    // return ans;
}
void init() {
   for (int i = 0; i < S; i++) st[i] = nullNode;</pre>
   memset(lazy, 0, sizeof(lazy));
   dfs_size(1, 1);
   dfs_hld(1, 1, 1);
```

```
int main(void) {
   cin >> n >> q;
   for (int i = 1; i <= n; i++) val[i] = 0;</pre>
   int a. b:
   for (int i = 2; i <= n; i++) {</pre>
       cin >> a >> b;
       adj[a].push_back(b);
       adj[b].push_back(a);
   }
   init(); // <----- DON'T FORGET TO CALL THIS FUNCTION
   int v;
   while (q--) {
       cin >> a >> b >> v;
       path(a, b, v);
   }
   for (int i = 1; i <= n; i++) {
       cout << get(i) << " ";
   }
   cout << el;</pre>
```

7.3.3 LCA functions using Binary Lifting

LCA functions using Binary Lifting

```
const int N = 2e5 + 15, M = 23;
int ancestors[N][M], depth[N], parent[N], val[N];
vector<vector<int>> adj;
//int tin[N], tout[N], timer;

void dfs_LCA(const int &node, const int &par) {
    // tin[node] = timer++;
    parent[node] = par;
    ancestors[node][0] = par;
    depth[node] = depth[par] + 1;

    for (int i = 1; i < M; i++) {
        int p = ancestors[node][i - 1];
        ancestors[node][i] = ancestors[p][i - 1];
    }
}</pre>
```

```
}
   for (const int &v: adj[node]) {
       if (v == par) continue;
       dfs_LCA(v, node);
   }
     tout[node] = timer++;
//bool is_ancestor(int u, int v) {
     return tin[u] <= tin[v] && tout[u] >= tout[v];
//}
int findKth(int u, int k) {
   if (depth[u] <= k) return -1;</pre>
   for (int i = M - 1; i >= 0; i--) {
       if (k & (1 << i)) {
           u = ancestors[u][i];
   }
   return u;
}
int getLCA(int u, int v) {
   if (depth[u] < depth[v])</pre>
       swap(u, v);
   u = findKth(u, depth[u] - depth[v]);
   if (u == v) return u;
   for (int i = M - 1; i >= 0; i--) {
       if (ancestors[u][i] == ancestors[v][i]) continue;
       u = ancestors[u][i];
       v = ancestors[v][i];
   return ancestors[u][0];
}
int getDistance(int u, int v) {
   int lca = getLCA(u, v);
   return (depth[u] + depth[v]) - (2 * depth[lca]);
}
int dfs_accumulate(const int &node, const int &par) {
```

```
for (const int& ch: adj[node]) {
       if (ch == par) continue;
       val[node] += dfs_accumulate(ch, node);
   }
   return val[node];
void applyOpOnPath(const int a, const int b, const int w) {
   // adding w to each node on the path a to b
   val[a] += w;
   val[b] += w;
   int lca = getLCA(a, b);
   val[lca] -= w;
   val[parent[lca]] -= w;
int main(void) {
   int n, q;
   cin >> n >> q;
   adj.resize(n + 1);
   int u, v;
   for (int i = 2; i <= n; ++i) {
       cin >> u >> v;
       adj[u].push_back(v);
       adj[v].push_back(u);
   }
   dfs_LCA(1, 1);
   parent[1] = -1;
   int w;
   for (int i = 0; i < q; ++i) {</pre>
       cin >> u >> v;
       cout << getDistance(u, v) << el;</pre>
   }
     dfs_accumulate(1, 0);
     for (int i = 1; i <= n; i++) {
         cout << val[i] << " ";
     }
     cout << el:</pre>
```

}

7.3.4 Topological Sort

Topological Sort Using DFS

```
//REF: USACO Guide
vector<int> top_sort;
vector<vector<int>> graph;
vector<bool> visited;
void dfs(int node) {
for (int next : graph[node]) {
 if (!visited[next]) {
  visited[next] = true;
  dfs(next);
top_sort.push_back(node);
int main() {
int n, m; // The number of nodes and edges respectively
std::cin >> n >> m;
 graph = vector<vector<int>>(n);
for (int i = 0; i < m; i++) {</pre>
 int a, b;
 std::cin >> a >> b;
 graph[a - 1].push_back(b - 1);
}
visited = vector<bool>(n);
for (int i = 0; i < n; i++) {</pre>
 if (!visited[i]) {
  visited[i] = true;
  dfs(i);
 std::reverse(top_sort.begin(), top_sort.end());
vector<int> ind(n);
for (int i = 0; i < n; i++) { ind[top_sort[i]] = i; }</pre>
```

```
// Check if the topological sort is valid
bool valid = true;
for (int i = 0; i < n; i++) {
  for (int j : graph[i]) {
    if (ind[j] <= ind[i]) {
     valid = false;
     goto answer;
    }
  }
}
answer:;

if (valid) {
  for (int i = 0; i < n - 1; i++) { cout << top_sort[i] + 1 << ' '; }
  cout << top_sort.back() + 1 << endl;
} else {
  cout << "IMPOSSIBLE" << endl;
}
</pre>
```

8 Techniques

8.1 Coordinate Compression

```
void coordinate_compress(vector<int> &x, int start=0, int
    step=1) {
    set unique(x.begin(), x.end());
    map<int, int> valPos;

    int idx=0;
    for (auto i: unique) {
       valPos[i] = start + idx * step;
       ++idx;
    }
    for(auto &i: x) i = valPos[i];
}
```

$Coordinate\ Compression$

8.2 Binary to decimal

Binary to decimal

```
// Function to convert binary to decimal
// 0(32)
int binaryToDecimal(string str)
{
   int dec_num = 0;
   int power = 0;
   int n = str.length();

   for(int i = n-1; i>=0; i--){
    if(str[i] == '1'){
      dec_num += (1<<power);
   }
   power++;
   }

   return dec_num;
}</pre>
```

8.3 Decimal to binary

Decimal to bianry

```
// Function that convert Decimal to binary
// 0(32)
void decToBinary(int n)
{
    // Size of an integer is assumed to be 32 bits
    for (int i = 31; i >= 0; i--) {
        int k = n >> i;
        if (k & 1)
            cout << "1";
        else
            cout << "0";
    }
}</pre>
```

```
// O(logn)
string DecimalToBinary(int num)
{
    string str;
    while(num){
    if(num & 1) // 1
        str+='1';
    else // 0
        str+='0';
    num>>=1; // Right Shift by 1
    }
    return str;
}
```

9 Number Theory

9.1 Divisors

9.1.1 formulas

number of divisors

```
int d(int n){
    unordered_map<int, int> factors = pf(n);
    int c = 1;
    for(const auto& factor: factors){
        c *= (factor.second+1);
    }
    return c;
}

// range Count Divisors backward thinking MAXN = 2e6
    for(int i=1; i <= n; ++i) {
        for(int j = i; j <= n; j += i) {
            numFactors[j]++;
        }
    int countDivisors(int n) {
        int count = 0;
    }
}</pre>
```

```
for (int i = 1; i * i <= n; ++i) {
   if (n % i == 0) {
      if (i == n / i) {
           count++; // Perfect square
      } else {
           count += 2; // Pair of divisors
      }
   }
}
return count;</pre>
```

sum of divisors

```
int s(int n){
   unordered_map<int,int> factors = pf(n);
   int sum = 1;
   for(const auto& factor: factors){
      int p = factor.first;
      int exp = factor.second;
      sum *= (pow(p,exp+1)-1)/p-1;
   }
   return sum;
}
```

9.2 Primes

prime factorization

number of co-primes with n

Prime Check

```
vector<bool> isPrime(MAXN, true);

void sieve() {
  isPrime[0] = isPrime[1] = false;

for (int i=2; i * i <= isPrime.size(); ++i) {
    if(isPrime[i]) {
      for (int j = 2 * i; i <= isPrime.size(); j += i)
          prime[j] = false;
    }
  }
}
bool Prime(int n) {
  if(n == 2) return true;
  if(n < 2 || n % 2 == 0) return false;</pre>
```

```
for(int i=3; i * i <= n; i += 2) {
   if(n % i == 0) return false;
 return true;
// Generate Primes
const int sz = sqrt(MAXN);
vector<int> prime;
vector<bool> vis(sz);
void pre() {
   prime.push_back(2);
   for (int j = 4; j < sz; j += 2) vis[j] = true;
   for (int i = 3; i < sz; i += 2) {</pre>
       if (vis[i]) continue;
       prime.push_back(i);
       for (int j = i * i; j < sz; j += i) vis[j] = true;</pre>
}
// Preprocessing Prime Factorization of range numbers
constexpr int N = 5e6+1;
int a[N];
for(int i=2; i < N; ++i) {</pre>
   if(!a[i]) {
       for(int j=1; i*j < N; ++j) {</pre>
           for(int k=i*j; k%i==0; k/=i) a[i*j]++;
   a[i] += a[i-1];
```

9.3 Math

9.3.1 Vieta's Formula for a Polynomial of Degree n

Problem: Given a polynomial of degree n:

$$P(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

with roots r_1, r_2, \ldots, r_n , express the sums and products of its roots using Vieta's formulas.

Solution: Using Vieta's formulas, we can relate the coefficients of the polynomial to sums and products of its roots:

• Sum of the roots taken one at a time:

$$r_1 + r_2 + \dots + r_n = -\frac{a_{n-1}}{a_n}$$

• Sum of the products of the roots taken two at a time:

$$r_1r_2 + r_1r_3 + \dots + r_{n-1}r_n = \frac{a_{n-2}}{a_n}$$

• Sum of the products of the roots taken three at a time:

$$r_1r_2r_3 + r_1r_2r_4 + \dots + r_{n-2}r_{n-1}r_n = -\frac{a_{n-3}}{a_n}$$

- Continue this pattern until:
- Product of the roots (for even n):

$$r_1 r_2 \cdots r_n = (-1)^n \frac{a_0}{a_n}$$

Example Problem Using Vieta's Formula

Problem: Given a quadratic equation $x^2 + bx + c = 0$ with roots r_1 and r_2 , find $r_1 + r_2$ and r_1r_2 .

Solution: Using Vieta's formulas for a quadratic equation $ax^2 + bx + c = 0$:

• Sum of the roots:

$$r_1 + r_2 = -\frac{b}{a}$$

• Product of the roots:

$$r_1r_2 = \frac{\alpha}{\alpha}$$

For the given quadratic $x^2 + bx + c = 0$ (where a = 1):

Also: To find the roots of the quadratic equation, we can use the discriminant formula, $D = b^2 - 4ac$. The roots will then be $x_1 = \frac{b - \sqrt{D}}{2}$ and $x_2 = \frac{b + \sqrt{D}}{2}$.

- $r_1 + r_2 = -b$
- $r_1r_2 = c$

Example Vieta's Formula for Cubic Equation

When considering a cubic equation in the form of $f(x) = ax^3 + bx^2 + cx + d$, Vieta's formula states that if the equation f(x) = 0 has roots r_1, r_2 , and r_3 , then:

- $r_1 + r_2 + r_3 = -\frac{b}{a}$
- $r_1r_2 + r_2r_3 + r_3r_1 = \frac{c}{a}$
- $\bullet \ r_1 r_2 r_3 = -\frac{d}{a}$

10 Geometry

10.1 Linearity

10.1.1 co-linear points

check if two points are co-linear

```
bool co_linear(int x1, int y1, int x2, int y2, int x3, int y3){
  int area = x1*(y2-y3) + x2*(y3-y1) + x3*(y1-y2);
  return area == 0;
}
```

10.2 polygons

10.2.1 Polygon formation

check if can form polygon with given angle

```
bool possible(double angle){
  if(angle <= 0 || angle >= 180) return false;

double sides = 360.0/(180.0-angle);

return (sides == static_cast<int>(sides) && sides >= 3);
}
```

10.2.2 Rectangle Intersection

intersection area between 2 rectangles

```
struct Rectangle {
    int x1, y1; // Bottom-left corner
    int x2, y2; // Top-right corner
};

int intersectionArea(const Rectangle& rect1, const Rectangle& rect2){
    int x_left = max(rect1.x1, rect2.x1);
    int y_bottom = max(rect1.y1, rect2.y1);
    int x_right = min(rect1.x2, rect2.x2);
    int y_top = min(rect1.y2, rect2.y2);

    int intersection_width = x_right - x_left;
    int intersection_height = y_top - y_bottom;

    if (intersection_width > 0 && intersection_height > 0) {
        return intersection_width * intersection_height;
    }

    return 0;
}
```

11 Miscellaneous

11.1 Faster implementations

11.1.1 hashes

custom hash

```
#define safe hash unordered_map<type, type, custom_hash> // same for
    gp_hash_table
struct custom_hash {
   static uint64_t splitmix64(uint64_t x) {
       // http://xorshift.di.unimi.it/splitmix64.c
       x += 0x9e3779b97f4a7c15;
       x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
       x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
       return x ^ (x >> 31);
   }
   size_t operator()(uint64_t x) const {
       static const uint64_t FIXED_RANDOM = chrono::steady_clock::now().
           time_since_epoch().count();
       return splitmix64(x + FIXED_RANDOM);
   }
};
```

gb hash table

```
//policy based ds (faster hash table)
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
gp_hash_table<int, int> table;
```

11.1.2 Binary Search the value

$nearest \ sqrt$

```
long long my_sqrt(long long a)
{
    long long l=0,r=5000000001;
    while(r-l>1)
```

```
{
    long long mid=(1+r)/2;
    if(111*mid*mid<=a)1=mid;
    else r=mid;
}
return 1;
}</pre>
```