

CSCI 239—Discrete Structures of Computer Science

Lab 10—Counting and Probabilities

This lab will explore ways to use Haskell to solve counting and probability problems.

Objectives:

- to gain more experience writing recursive Haskell functions
- to solve counting problems using a computer
- to solve an interesting probability problem

Part 1: Haskell methods for counting

In this part you will implement Haskell functions for arrangements and selections, i.e. the $p(n, r)$ and $c(n, r)$ functions.

A generalized factorial function

The p and c functions we want to implement are both defined using the factorial function. To make things more efficient, however, it would be helpful to have a function that computes the product of an arbitrary range of integers from m to n . We can use this function to compute the factorial of n by setting m to 1. We need to be careful, however, because we want zero to be factored in as one as well. The following recurrence relation will compute what we want:

$$fact(m, n) = \begin{cases} fact(m, -1) & n = 0 \\ 1 & m > n \\ n \cdot fact(m, n - 1) & \text{otherwise} \end{cases}$$

Implement this recurrence relation as a curried Haskell function. (That is, it takes two parameters, not a single parameter that is a 2-tuple.) Be careful with the cases; you can use pattern matching for the first case, but you have to use the vertical bar selection for the next two; you may use the vertical bar selection for all three if you prefer.

Use the `:type` command to determine the type of the function. Test your function on enough small cases (ranges from negative to positive) to be confident it is working correctly.

A function to compute arrangements (permutations)

The function $p(n, r)$ is defined as:

$$p(n, r) = \frac{n!}{(n - r)!}$$

Using a little algebra, we can simplify this to:

$$\begin{aligned} p(n, r) &= (n - r + 1)(n - r + 2) \dots (n - 1)n \\ &= fact(n - r + 1, n) \end{aligned}$$

Implement this last formula as a Haskell function, get its type, and test it. (Note that there is no more recursion here if you use the `fact` function.) We use the simplified formula to reduce the number of multiplications required by a factor of $2(n - r)$.

Two functions to compute selections

The function $c(n, r)$ is defined as:

$$c(n, r) = \binom{n}{r} = \frac{n!}{r!(n-r)!}$$

Using simple algebra, we can simplify this to two equivalent formulas:

$$\begin{aligned} c(n, r) &= \frac{(r+1)(r+2) \dots (n-1)n}{(n-r)!} = \frac{\text{fact}(r+1, n)}{\text{fact}(1, n-r)} \\ &= \frac{(n-r+1)(n-r+2) \dots (n-1)n}{r!} = \frac{\text{fact}(n-r+1, n)}{\text{fact}(1, r)} \end{aligned}$$

We again want to use a simplified formula to reduce the number of multiplications, but which one is better? Whichever has fewer multiplications, of course! And which is that? It depends. If $r > n/2$ then the first formula will have fewer multiplications; otherwise, the second formula has fewer (or the same if $r = n/2$). Implement a curried Haskell function $c1$ using these two formulas, choosing the formula that is more efficient depending on the value of r relative to n .

There is another, recursive definition for $c(n, r)$, based on our definition of Pascal's triangle:

$$c(n, r) = \begin{cases} 1 & r = 0, r = n \\ c(n-1, r) + c(n-1, r-1) & 0 < r < n \\ 0 & \text{otherwise} \end{cases}$$

Implement a curried Haskell function $c2$ using this definition. Notice that there are only additions, and no multiplications or divisions in this definition, so it should be faster in principle.

Test both functions for all values of n from 0 up to 8 and r from 0 to n . Make a table of your results. If you aren't getting the same values for both functions, at least one of them is not implemented correctly. Keep working on it until you get the same value for both functions.

Now try computing $c1\ 20\ 10$ and $c2\ 20\ 10$. Which is faster?

Use the p function and the faster of the c functions you developed to compute the answers to Exercises 10.4.2.b, 10.4.3.c, 10.5.2.d, 10.5.3.f, and 10.5.5 from the zyBooks textbook.

Use the p function and the faster of the c functions you developed to compute the answers to Exercises 12.1.5.c, 12.2.4.c, and 12.3.5 (compute the two sides of the final inequality).

Part 2: The Monty Hall problem

(Note: Monty Hall died on September 30, 2017, age 96. We do this problem to honor his life.)

In the 1960s and 70s (and maybe into the 80s), Monty Hall was the host of the TV game show *Let's Make a Deal*. (You can see reruns on the *Buzzr* and *Game Show* networks.) One of the games on the show captured the interest of statisticians everywhere. The game works like this: there are three giant doors on the stage labeled 1, 2, and 3; behind two of the doors are goats (UGH!!!); behind the other door is a ...

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[illegible]

(Back in the 60s, a new car was pretty much the most fabulous prize anyone could imagine.)

Anyway, back to the description. Monty would offer the contestant a choice of one of the three doors. Then, since he knew everything, Monty would have one of the other two doors opened, always revealing a goat. (UGH!!!) At this point, one of the doors still hides a goat. (UGH!!!) And behind the other door hides the

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Now comes the fun part: Monty tells the contestant she may either keep the door she chose in the first place, or she may switch to the other unopened door. Should she stick to her first choice, switch, or doesn't it matter? Your task is to compute the probability of winning when she doesn't switch and when she does to answer this question.

Before you go on, what does your intuition tell you? After Monty reveals a door hiding a goat (UGH!!!), two doors remain. Does switching matter, given that two doors remain, one with the fabulous prize and the other with the not-so-fabulous prize?

First, consider the parameters of the sample space: which door hides the ...

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(thereby determining which doors hide the goats (UGH!!!)), which door the contestant chooses, which door Monty reveals (does this one affect the outcome?), and whether the contestant switches doors or not. How big is this sample space of four or three parameters? What determines the outcome? Note that if the contestant chooses the right door to start with, she wins only if she doesn't switch. If she chooses the wrong door at first, when does she win? In both cases, when does she lose? (Answering these questions correctly should confirm whether the door Monty reveals is relevant to the sample space. Assume that the goats (UGH!!!) are indistinguishable.)

Consider the sample space parameters in turn, and give your answers to each question below:

- a) What are the ways to configure the “prizes” behind the three doors? There will be one

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and two goats (UGH!!!), one behind each door. How many ways are there? What is the mathematical formula that explains why you got those answers?

- b) What are the contestant's choices? How many choices are there? Again, what is the mathematical formula that explains why you got those answers?
c) What are Monty's choices when he reveals what's behind one of the doors? Remember, he never reveals the door the contestant chose, and he never reveals the door with the

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Does he always have a choice? If not, does he ever have a choice? If so, when does he have a choice? In that case, does his choice affect the outcome? If not, you don't have to include Monty's choice in the sample space.

Do not go beyond this point until you've verified with the TA or lab instructor that you have the correct answer here.

Part 3: A Monte Carlo/Monty Hall Simulation

You can use the supplied Haskell file to help you write a simulation of the problem. Write a function *makeTuple* that construct an outcome as a tuple, using parameters that pertain to each outcome component. Now write a function *makeTupleList* that combines equal length lists of outcome components into a list of outcome tuples.

Next, write a function *isWin* that identifies whether an outcome tuple is a win or not, another function *isSwitch* that identifies whether the contestant switched doors in an outcome, and a third function *isWinWithSwitch* that identifies whether an outcome is a win with a door switch. Finally, write three functions that respectively count the number of winning outcomes (*countWins*), the number of door switches (*countSwitches*), and the number of winning outcomes with door switches (*countWinsWithSwitches*) in a list of outcomes.

Now use the supplied functions (*selectDoors* twice and *selectSwitches*) and the *makeTupleList* function you wrote to generate a list of a few thousand random outcomes and bind that list to an identifier. Apply the last three functions to that list. The ratio of total wins to total outcomes should be close to the value you calculated for $p(\text{WIN})$, and the ratio of wins with switches to total outcomes with switches should be close to the value you calculated for $p(\text{WIN} \mid \text{SW})$. If not, consult with the instructor or TA to figure out what's going on.

NOTE:

The seed parameter for the two supplied random-list-building functions can be any integer. Every time you use the same seed, you get the same sequence of elements, even if the lists are different lengths. (You just get fewer with a smaller length.) Change the seed to get a different sample. Try your Banner ID (or the last four digits). Just use a different seed from others you're working with.