

Grover

Shor's

① Polynomial Speed

Exponential Speed

② optional not optional

→ No classical alg can supercede Grover algorithm

③ wide spread applications

→ ML alg. we use Grover's algorithm or its variants.

$$f: \{0,1\}^n \rightarrow \{0,1\}$$

$$\text{find } x \in \{0,1\}^n \text{ s.t. } f(x) = 1$$

$$\text{Let } n=3$$

f-domain

f-range

000

0

$$f(000) = 0$$

001

0

010

0

~~011~~

1

$$f(011) = 1$$

011

-

100

0

101

0

$$f(111) = 0$$

110

0

111

0

classical algorithm  $O(2^n)$  complexity

Grover's alg.  $O(\sqrt{2^n})$

Still exponential but polynomial improvement.

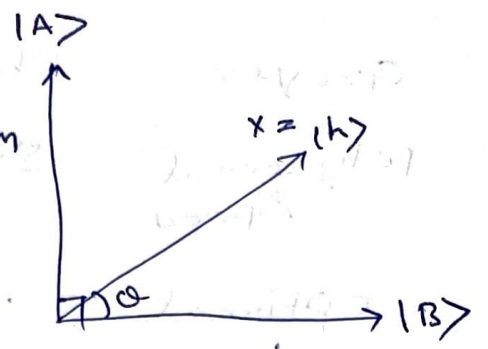
# Key Idea

$|A\rangle$  represents equal superposition of all the inputs.

$\forall x \text{ s.t. } f(x) = 1$

Ex:

f-domain	f-range
000	1
001	0
010	0
011	1
100	0
101	0
110	1
111	0



Let  $|A\rangle = \frac{1}{\sqrt{3}} \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$

Let  $|B\rangle = \frac{1}{\sqrt{5}} \begin{bmatrix} 0 \\ 1 \\ 0 \\ 1 \\ 0 \\ 1 \\ 0 \\ 1 \end{bmatrix}$

Above: three ops yield true  
remaining ops yield false

Vectors  $|A\rangle$  and  $|B\rangle$  are orthogonal to each other  $\langle A|B\rangle = 0$

$|h\rangle$  Created using Hadamard gates.

Let the angle between  $|B\rangle$  and  $|h\rangle$  is  $\theta$

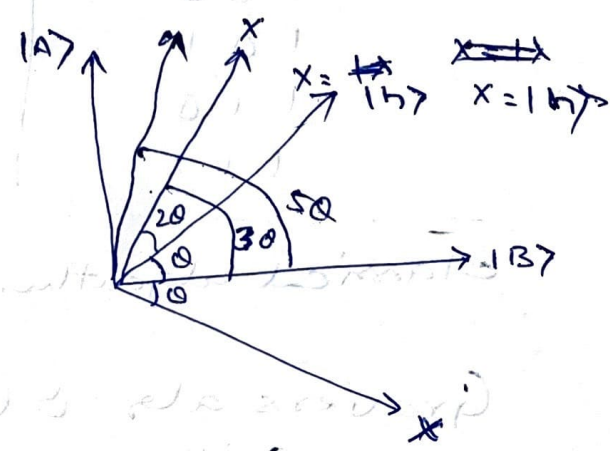
Our goal is to move register  $x$  to move closer to  $|A\rangle$

Action 1: Reflect  $x$  over  $|B\rangle$

2. Reflect  $x$  over  $|h\rangle$

Now  $x$  moved towards  $|A\rangle$

Repeat Action 1 and Action 2



Steps in Grover's algorithm:

① Create equal superposition of all possible inputs in  $|h\rangle$

② for ( $i = 1$  to  $\lfloor \frac{\pi}{4} \sqrt{N} \rfloor$ ) {

$X = \text{Reflect } X \text{ over } |B\rangle$

$X = \text{Reflect } X \text{ over } |h\rangle$

}

③ Measure  $X$

// high probs to get the answer  
> 99% accuracy

Yet to do

① Why we have  $\lfloor \frac{\pi}{4} \sqrt{N} \rfloor$  Grover rotations.

② Math of Grover's rotation

③ Quantum Circuit

④ Calculate prob. of success

⑤ Working Example