

Problem Set 6 Solutions

Due: Thursday, October 20

Reading Assignment: Sections 5.4.3, 6.1, 6.2, 7.1-7.9

Problem 1. [20 points]

This problem involves trees, which are a special class of graphs. Suppose we consider a simple connected graph G on n nodes. Then, G is a tree if and only if G is minimally connected (that is removing an edge disconnects G).

(a) [5 pts] Show that our definition of a tree implies that a simple connected graph G is a tree if and only if G contains no cycles.

Rubric

[5pts] Any correct proof.

or

[2.5pts] Show that if G contains a cycle, it is not minimally connected.

[2.5pts] Contrapositive – show that if G is not minimally connected, then it contains a cycle.

Solution. Suppose that a simple connected graph G is a tree, then we know by definition that it is minimally connected. For the sake of contradiction, suppose that G has a cycle and remove an edge (u, v) in the cycle. Now, we claim any two vertices x, y in G are still connected. If x, y were connected in G by some path p that did not use edge (u, v) , then they are still connected after removing edge (u, v) . Now if path p did contain edge (u, v) , then we replace the edge (u, v) by the other longer path from u to v given by the cycle. We now have a walk from x to y , but now if there is a walk from x to y , there must be a path and the graph is still connected after removing edge (u, v) , which is a contradiction.

Now we show the other direction by proving the contrapositive. That is, we want to show that if G is not minimally connected, then it contains a cycle. As G is not minimally connected, there is some edge $e = (u, v)$ such that removing e still leaves our graph connected. Hence, there must be a path p from u to v after removing edge e . However, this means that the path p along with edge e forms a cycle, and so G will contain a cycle. ■

- (b) [5 pts] Prove that a simple connected graph with n nodes and $n - 1$ edges is a tree.

Rubric

[5pts] Any correct proof.

or

[2pts] Proof by contradiction. Assume G has a cycle.

[1pt] State that we want to show it is acyclic.

[2pts] Attempt to remove an edge and show it is not connected.

Solution. To show that a connected graph G with n nodes and $n - 1$ edges is a tree, it is sufficient to show that it is acyclic. Assume to the contrary that there is a cycle in G ; then we can remove an edge from this cycle while preserving connectivity. Then we have produced a connected subgraph G' with n vertices but only $n - 2$ edges; this contradicts connectivity, and we conclude that the original graph must have been a tree. ■

- (c) [10 pts] Prove that in any tree, any two longest paths cross each other (that is any two longest paths share a vertex).

Rubric

[10pts] Any correct proof.

or

[2pts] Proof by contradiction.

[2pts] Assume two longest paths do not cross each other.

[2pts] Since trees are connected, there is some path between the two longest path.

[3pts] Show that the new longest path would be a combination of the two plus the path between the two.

Solution. We prove this by contradiction. Suppose that there are two longest paths in our tree that do not cross each other. Let one path be from vertex u to vertex v and the other be from vertex x to vertex y . Then, as trees are connected, we must have that there is some vertex a along the path from u to v that is connected to some vertex b along the path from x to y by a path p from a to b of length at least 1. Now, let p_1 be the longer of the paths from u to a and from a to v . Similarly, let p_2 be the longer of the paths from x to b and from b to y . Then each of the lengths of p_1 , p_2 is at least half the length of the longest path. Now consider the path $p_1 \cup p \cup p_2$. As the length of p is at least 1, we have that the length of $p_1 \cup p \cup p_2$ is at least the length of the longest path plus 1, which is a contradiction as this is now the longest path in the tree. ■

Problem 2. [10 points] The adjacency matrix A of a graph G with n vertices as defined in lecture is an $n \times n$ matrix in which $A_{i,j}$ is 1 if there is an edge from i to j and 0 if there is not. In lecture we saw how the smallest k where $A_{i,j}^k \neq 0$ describes the length of the shortest path from i to j . Give a combinatorial interpretation of the following statements about the adjacency matrix in terms of connectivity properties of G . For example, the smallest k such that $A_{i,j}^k$ is non-zero means that the distance from i to j is at most k .

Rubric

For each:

[3pts] Correct solution (all-or-nothing).

(a) [3 pts] The trace of A^k (the trace of a matrix is the sum of its diagonal entries).

Solution. This is the number of closed walks of length k in the graph G . ■

(b) [2 pts] $\forall k. A_{i,j}^k = 0$

Solution. There is no path in G from i to j . ■

(c) [2 pts] $\forall i \forall k. A_{i,i}^k = 0$

Solution. The graph is acyclic. ■

(d) [3 pts] $\forall k$, we can write A^k as $\begin{bmatrix} X & 0 \\ 0 & Y \end{bmatrix}$

Solution. There is no path from X to Y and so there are at least two connected components. ■

Problem 3. [15 points] A set of PageRank values is stationary if the amount of PageRank going into each vertex is the same as the amount leaving the vertex on every update. In this problem we will prove that a strongly connected graph has at most one set of PageRank values that are stationary. There always is a set of PageRank values that are stationary, but we are not asking you prove this.

(a) [7 pts] For two sets of values of PageRank, d_1 and d_2 , let γ be defined as $\gamma : \max_{x \in V} \frac{d_1(x)}{d_2(x)}$, the maximum ratio of a value in d_1 over the corresponding value in d_2 . Show that there exists a directed edge from y to z such that $d_1(y)/d_2(y) < \gamma$ and $d_1(z)/d_2(z) = \gamma$.

Rubric

[3pts] Argue that ratio of γ has to exist otherwise values would be the same.

[4pts] Show why being strongly connected is necessary. Give 2 or 3 points if they mention something about vertices being connected.

Solution. γ must be attained at some vertex and not all vertices have a ratio of γ because if they did then the values would be the same. Because the graph is strongly connected, we can follow edges backwards until we arrive at any other vertex and there must be some vertex along the way that does not attain γ . So such an edge must exist. ■

(b) [8pts] Prove that a strongly connected graph has at most one set of PageRank values that are stationary by deriving a contradiction using the edge found in part a.

Rubric

[8pts] Any correct proof.

or

[2pts] d_1 and d_2 are stationary at z

[3pts] Apply a step of page rank

[3pts] Correct math for deriving a contradiction.

[-1pt] Math errors.

Solution. Assume for the sake of contradiction that there are two stationary values d_1 and d_2 . Let γ be defined as $\gamma =: \max_{x \in V} \frac{d_1(x)}{d_2(x)}$, the maximum ratio of a value in d_1 over the corresponding value in d_2 .

Pick an edge from y to z such that $d_1(y)/d_2(y) < \gamma$ and $d_1(z)/d_2(z) = \gamma$. γ must be attained at some vertex and not all vertices have a ratio of γ because if they did then the values would be the same. Because the graph is strongly connected, we can follow edges until we arrive at any other vertex and there must be some vertex along the way that does not attain γ . So such an edge must exist.

Define $p(x, z)$ to be $\frac{\text{Number of edges from } x \text{ to } z}{\text{Number of edges out of } x}$. We apply one step of PageRank to d_1 and d_2 which are stationary values \hat{d}_1 and \hat{d}_2 .

$$d_1(z) = \hat{d}_1(z) = \sum_{x: x \rightarrow z} d_1(x)p(x, z), \quad (1)$$

and

$$d_2(z) = \hat{d}_2(z) = \sum_{x: x \rightarrow z} d_2(x)p(x, z). \quad (2)$$

Substituting $d_1(z) = \gamma d_2(z)$ into (1), and recalling that $d_1(x) \leq \gamma d_2(x)$ for all x , we have

$$\begin{aligned}
 \gamma d_2(z) &= \sum_{x:x \rightarrow z} d_1(x)p(x, z) \\
 &= \left(\sum_{x \neq y: x \rightarrow z} d_1(x)p(x, z) \right) + d_1(y)p(y, z) \\
 &\leq \left(\sum_{x \neq y: x \rightarrow z} \gamma d_2(x)p(x, z) \right) + d_1(y)p(y, z) \\
 &\leq \left(\left(\gamma \sum_{x:x \rightarrow z} d_2(x)p(x, z) \right) - \gamma d_2(y)p(y, z) \right) + d_1(y)p(y, z) \\
 &\leq \gamma d_2(z) - \gamma d_2(y)p(y, z) + d_1(y)p(y, z) \\
 &< \gamma d_2(z) - \gamma d_2(y)p(y, z) + \gamma d_2(y)p(y, z) \\
 &< \gamma d_2(z).
 \end{aligned}$$

The strict inequality at the next to last step of the derivation follows from the assumption that $d_1(y) < \gamma d_2(y)$, and $p(y, z) > 0$. Thus, we have shown $d_2(z) < d_2(z)$, which is a contradiction. ■

Problem 4. [15 points]

(a) [5 pts] For the graph in Figure 1, compute the first two iterations of PageRank, starting from uniform PageRank values across all vertices.

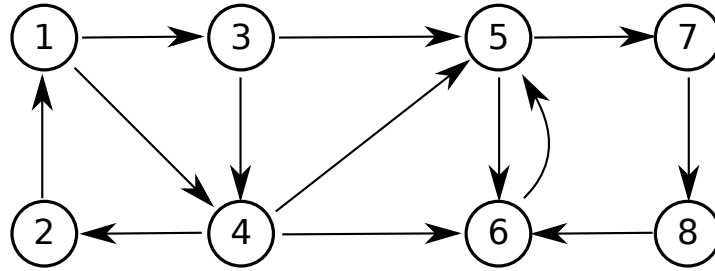


Figure 1: A graph of web pages and links.

Rubric

[-0.5pts] Each incorrect vertex in iter 2.

[-5pts] Obviously incorrect methodology.

Solution. In the table below, we write the PageRank value of each vertex at iterations 0, 1, and 2. We also include, for each vertex, the individual contributions made to that vertex in computing the next stage, which are summed in computing the following PageRank value.

	1	2	3	4	5	6	7	8
iter 0	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$
contrib	$\frac{1}{8}$	$\frac{1}{24}$	$\frac{1}{16}$	$\frac{1}{16}, \frac{1}{8}$	$\frac{1}{8}, \frac{1}{16}, \frac{1}{24}$	$\frac{1}{8}, \frac{1}{16}, \frac{1}{24}$	$\frac{1}{16}$	$\frac{1}{8}$
iter 1	$\frac{1}{8}$	$\frac{1}{24}$	$\frac{1}{16}$	$\frac{1}{8}$	$\frac{11}{48}$	$\frac{11}{48}$	$\frac{1}{16}$	$\frac{1}{8}$
contrib	$\frac{1}{24}$	$\frac{1}{24}$	$\frac{1}{16}$	$\frac{1}{16}, \frac{1}{32}$	$\frac{1}{24}, \frac{1}{32}, \frac{11}{48}$	$\frac{1}{8}, \frac{1}{24}, \frac{11}{96}$	$\frac{11}{96}$	$\frac{1}{16}$
iter 2	$\frac{1}{24}$	$\frac{1}{24}$	$\frac{1}{16}$	$\frac{3}{32}$	$\frac{29}{96}$	$\frac{9}{32}$	$\frac{11}{96}$	$\frac{1}{16}$

■

(b) [10 pts] A strongly connected component of a directed graph is a subgraph which has the property that for every pair of vertices u and v , there exists a path from u to v and one from v to u . Also, every vertex which can be reached by a path starting in the strongly connected component is also in the strongly connected component. Suppose that a graph G consists of exactly two strongly connected components C_1 and C_2 , and that there exist edges from C_1 to C_2 (but not from C_2 to C_1). There is always a stationary set of values which is non-negative. Prove that the stationary PageRank values of this graph are entirely concentrated in C_2 , i.e. that the PageRank values are all zero on C_1 .

Rubric

[10pts] Any correct proof.

or

[2pts] Proof by contradiction.

[2pts] Contradiction is in whether PageRank values are stationary.

[2pts] Use a vertex in C_1 that has an edge to C_2 .

[2pts] Show vertex has contributions to it.

[2pts] No contributions from C_2 to C_1 so not stationary.

Solution. For any vertex v , let $p(v)$ denote the stationary PageRank value at v , so that if we apply another step of the PageRank update rule to these values, they remain unchanged. Suppose for a contradiction that for some vertex $v \in C_1$, $p(v) \neq 0$.

Let w be any other vertex of C_1 ; we first claim that the PageRank value at w is positive. From the hypothesis that C_1 is strongly connected, there exists a path from v to w , of some length k . If we were to apply k steps of the PageRank algorithm, then the non-zero

PageRank value $p(v)$ would yield some positive contribution to the resulting value at w , via the path mentioned above. The resulting PageRank value at w after k steps would then be positive. As the PageRank values were already stationary we can conclude that $p(w) > 0$.

Now, we know that some vertex $w \in C_1$ has an edge leading into C_2 . As the PageRank value $p(w)$ is positive, then if we were to apply another step of the PageRank update rule, vertex w would make a contribution of at least $\frac{p(w)}{\text{outdeg}(w)}$ to the total PageRank of component C_2 . There are however no edges from C_2 to C_1 . As the total PageRank in a graph always sums to 1, we must conclude that the total PageRank in component C_1 decreases by at least $\frac{p(w)}{\text{outdeg}(w)} > 0$ by applying the update rule.

But these PageRank values were assumed to be stationary, and we have argued above that they must change when we apply the update step. This is a contradiction, and so the stationary PageRank values must be zero throughout component C_1 , as desired. ■

Problem 5. [10 points] For each of the following, either prove that it is an equivalence relation and state its equivalence classes, or give an example of why it is not an equivalence relation.

Rubric

For each:

[1pt] Correct true/false for equivalence relation.

[1pt] Correct counterexample or proof.

(a) [2 pts] $R_n := \{(x, y) \in \mathbb{Z} \times \mathbb{Z} \text{ s.t. } x \equiv y \pmod{n}\}$

Solution. It is an equivalence relation. To prove this, we will show that R_n is symmetric, transitive, and reflexive.

- **Reflexive:** $x \equiv x \pmod{n}$. This is because $x = x + 0 \cdot n$.
- **Symmetric:** We want to show that $R_n(x, y) \Rightarrow R_n(y, x)$. If $R_n(x, y)$, then there is some $c \in \mathbb{Z}$ such that $x = y + c \cdot n$. But then, subtracting $c \cdot n$ from both sides, we have that $y = x + (-c) \cdot n$, so $y \equiv x \pmod{n}$. So $R_n(y, x)$, and the symmetric property holds.
- **Transitivity.** Suppose $R_n(x, y)$ and $R_n(y, z)$. From the first statement, we know that there is some $c \in \mathbb{Z}$ such that $x = y + c \cdot n$. From the second, we know that there is some $d \in \mathbb{Z}$ such that $y = z + d \cdot n$. Substituting in this value of y , we see that $x = (z + d \cdot n) + c \cdot n = z + (d + c) \cdot n$. The sum $c + d$ is an integer, so $R_n(x, z)$ holds.

The equivalence classes are then the sets of numbers congruent to the numbers $\{0, 1, \dots, n-1\}$ modulo n . ■

(b) [2 pts] $R := \{(x, y) \in P \times P \text{ s.t. } x \text{ is taller than } y\}$ where P is the set of all people in the world today.

Solution. This is not an equivalence relation, because the concept of symmetry is broken. If y is taller than x , then x is not taller than y . ■

(c) [3 pts] $R := \{(x, y) \in \mathbb{Z} \times \mathbb{Z} \text{ s.t. } \gcd(x, y) = 1\}$

Solution. This is not an equivalence relation, because transitivity is broken. Consider the case when $x = 3$, $y = 7$, and $z = 15$. Then, $\gcd(x, y) = 1$ and $\gcd(y, z) = 1$, but $\gcd(x, z) = 3 \neq 1$. ■

(d) [3 pts] $R_G :=$ the set of $(x, y) \in V \times V$ such that V is the set of vertices of a graph G , and there is a path x, v_1, \dots, v_k, y from x to y along the edges of G .

Solution. This is an equivalence relation. We will show this by proving that it obeys reflexivity, symmetry, and transitivity.

- **Reflexivity:** Any vertex is connected to itself.
- **Symmetry:** If $R_G(x, y)$, then there is a path x, v_1, \dots, v_k, y from x to y . The reverse of this path is y, v_k, \dots, v_1, x , and is a path from y to x . So $R_G(y, x)$.
- **Transitivity:** Suppose $R_G(x, y)$ and $R_G(y, z)$. Then, there is a path from x to y : x, v_1, \dots, v_k, y . Furthermore, there is a path from y to z : y, w_1, \dots, w_l, z . But then, the concatenation of those two is a path $x, v_1, \dots, v_k, y, w_1, \dots, w_l, z$ from x to z . So $R_G(x, z)$.

Thus we have shown that R_G is an equivalence relation on a graph G , and the equivalence classes are the connected components of G . ■

Problem 6. [10 points] Let R_1 and R_2 be two equivalence relations on a set, A . Prove or give a counterexample to the claims that the following are also equivalence relations:

(a) [5 pts] $R_1 \cap R_2$.

Rubric

[1pt] Is an equivalence relation.

[4pts] Correct proof.

Solution. Let $R =: R_1 \cap R_2$. We give two proofs that R is an equivalence relation using different characterizations of equivalence relations.

Proof. We first prove that R is an equivalence relations by showing that R is reflexive, symmetric, and transitive.

Reflexive: R_i is reflexive because it is an equivalence relation, for $i = 1, 2$. So $(a, a) \in R_i$ for $i = 1, 2$ and all $a \in A$. So, $(a, a) \in (R_1 \cap R_2) = R$ for all $a \in A$, that is, R is reflexive.

Transitive: Suppose $(a, b), (b, c) \in R$. Since $R = R_1 \cap R_2$, we have $(a, b), (b, c) \in R_i$ for $i = 1, 2$. But R_i is an equivalence relation, and so is transitive. Therefore, $(a, c) \in R_i$, and so $(a, c) \in R_1 \cap R_2 = R$. This shows that R is transitive.

Symmetric: The proof that R is symmetric follows the same format.

□

Proof. Since R_i is an equivalence relation for $i = 1, 2$, there is by definition a total function, f_i , with domain, A , such that

$$aR_ib \quad \text{iff} \quad f_i(a) = f_i(b).$$

Define the function, f , with domain, A , by

$$f(a) =: (f_1(a), f_2(a)).$$

Clearly f is total, since f_i is total for $i = 1, 2$. Now we have,

$$\begin{array}{llll} aRb & \text{iff} & a(R_1 \cap R_2)b & \text{def. of } R \\ & \text{iff} & aR_ib \text{ for } i = 1, 2 & \text{def. of } \cap \\ & \text{iff} & f_i(a) = f_i(b) \text{ for } i = 1, 2 & \text{def. of } f_i \\ & \text{iff} & (f_1(a), f_2(a)) = (f_1(b), f_2(b)) & \text{def. of } = \text{ pairs} \\ & \text{iff} & f(a) = f(b) & \text{def. of } f. \end{array}$$

That is

$$aRb \quad \text{iff} \quad f(a) = f(b),$$

which proves that R is the equivalence relation \equiv_f .

□

■

(b) [5 pts] $R_1 \cup R_2$.

Rubric

[1pt] Is not an equivalence relation.

[4pts] Correct counterexample.

Solution. We give a counterexample showing that $R_1 \cup R_2$ may not be an equivalence relation. Let R_1 and R_2 be the relations on $\{1, 2, 3\}$ where

$$\begin{aligned} R_1 &=: \{(1, 1)(2, 2)(3, 3)(1, 2)(2, 1)\}, \\ R_2 &=: \{(1, 1)(2, 2)(3, 3)(2, 3)(3, 2)\}. \end{aligned}$$

It's easy to check that R_1 and R_2 are both equivalence relations. But $R_1 \cup R_2$ is not transitive, because $(1, 2), (2, 3) \in R_1 \cup R_2$ and $(1, 3) \notin R_1 \cup R_2$. Therefore $R_1 \cup R_2$ is not an equivalence relation.

■

Problem 7. [20 points] In this problem we study partial orders (posets). Recall that a weak partial order \preceq on a set X is reflexive ($x \preceq x$), anti-symmetric ($x \preceq y \wedge y \preceq x \rightarrow x = y$), and transitive ($x \preceq y \wedge y \preceq z \rightarrow x \preceq z$). Note that it may be the case that neither $x \preceq y$ nor $y \preceq x$. A chain is a list of *distinct* elements x_1, \dots, x_i in X for which $x_1 \preceq x_2 \preceq \dots \preceq x_i$. An antichain is a subset S of X such that for all distinct $x, y \in S$, neither $x \preceq y$ nor $y \preceq x$.

(a) [10 pts] A chain in a poset is called maximal if it cannot be extended (that is no other element of the poset can be added to it). Let B_4 (the Boolean algebra) be the set of subsets of $\{1, 2, 3, 4\}$ where the partial order is defined by the contains operation (the number of elements in B_4 is 2^4 and for subsets S_1, S_2 , $S_1 < S_2$ if $S_1 \subset S_2$). How many maximal chains does B_4 have?

Rubric

[10pts] 24 or 4!

or

[-1pt] Math errors.

[5pts] Use definition of maximal by extending a chain.

[5pts] Permutation on the elements that can be added to the chain.

Solution. B_4 should have $4! = 24$ maximal chains. Let us consider the general case of B_n and show that B_n has $n!$ maximal chains. Note that if we have some non-maximal chain, we can always extend it by adding a new element to the largest subset in the chain if the largest element is not $\{1, 2, \dots, n\}$ or by removing an element from the smallest subset in the chain if the smallest element is not \emptyset . Hence, to find all maximal chains, let us begin with \emptyset and add elements to our chain until we reach $\{1, 2, \dots, n\}$. We have n choices for our first extension to our chain, $n - 1$ choices for the second extension and so on until we are left with exactly 1 choice to extend our chain to $\{1, 2, \dots, n\}$. Hence there are $n!$ maximal chains, and so for B_4 there are 24. ■

(b) [10 pts] Let π be a permutation, and let d be the smallest integer so that π can be written as a union of d increasing sequences. Prove that the longest decreasing subsequence of π consists of d elements. (Hint: use Dilworth's theorem)

Rubric

[10pts] Any correct proof.

or

[3pts] Define a partial order for the poset.

[3pts] Chains in poset are increasing subsequences.

[3pts] Antichains correspond to decreasing subsequences.

[1pt] Dilworth's theorem.

Solution. This is actually a special case of Dilworth's theorem. We just need to define the appropriate partial order $<_P$ for our poset P . Let $\pi = \pi_1\pi_2\ldots\pi_n$. Then we say that $\pi_i <_P \pi_j$ if $\pi_i < \pi_j$ and $i < j$. Then the chains of our poset P are the increasing subsequences in π and so antichains of P correspond to decreasing subsequences in π . The proof then follows by Dilworth's theorem. ■