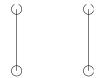
## 1 Build-up error

A graph is *connected* iff there is a path between every pair of its vertices.

**False Claim.** If every vertex in a graph has positive degree, then the graph is connected.

1. Prove that this Claim is indeed false by providing a counterexample. here are many counterexamples; here is one:



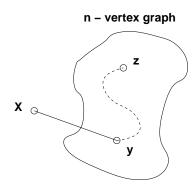
2. Since the Claim is false, there must be a logical mistake in the following bogus proof. Pinpoint the *first* logical mistake (unjustified step) in the proof.

*Proof.* We prove the Claim above by induction. Let P(n) be the proposition that if every vertex in an n-vertex graph has positive degree, then the graph is connected.

**Base cases**:  $(n \leq 2)$ . In a graph with 1 vertex, that vertex cannot have positive degree, so P(1) holds vacuously.

P(2) holds because there is only one graph with two vertices of positive degree, namely, the graph with an edge between the vertices, and this graph is connected.

**Inductive step**: We must show that P(n) implies P(n+1) for all  $n \geq 2$ . Consider an n-vertex graph in which every vertex has positive degree. By the assumption P(n), this graph is connected; that is, there is a path between every pair of vertices. Now we add one more vertex x to obtain an (n+1)-vertex graph:



All that remains is to check that there is a path from x to every other vertex z. Since x has positive degree, there is an edge from x to some other vertex, y. Thus, we can obtain a path from x to z by going from x to y and then following the path from y to z. This proves P(n+1).

By the principle of induction, P(n) is true for all  $n \geq 0$ , which proves the Claim.

his one is tricky: the proof is actually a good proof of something else. The first error in the proof is only in the final statement of the inductive step: "This proves P(n+1)".

The issue is that to prove P(n+1), every (n+1)-vertex positive-degree graph must be shown to be connected. But the proof doesn't show this. Instead, it shows that every (n+1)-vertex positive-degree graph that can be built up by adding a vertex of positive degree to an n-vertex connected graph, is connected.

The problem is that not every (n + 1)-vertex positive-degree graph can be built up in this way. The counterexample above illustrates this: there is no way to build that 4-vertex positive-degree graph from a 3-vertex positive-degree graph.

More generally, this is an example of "buildup error". This error arises from a faulty assumption that every size n + 1 graph with some property can be "built up" in some particular way from a size n graph with the same property. (This assumption is correct for some properties, but incorrect for others— such as the one in the argument above.)

One way to avoid an accidental build-up error is to use a "shrink down, grow back" process in the inductive step: start with a size n+1 graph, remove a vertex (or edge), apply the inductive hypothesis P(n) to the smaller graph, and then add back the vertex (or edge) and argue that P(n+1) holds. Let's see what would have happened if we'd tried to prove the claim above by this method:

Inductive step: We must show that P(n) implies P(n+1) for all  $n \geq 1$ . Consider an (n+1)-vertex graph G in which every vertex has degree at least 1. Remove an arbitrary vertex v, leaving an n-vertex graph G' in which every vertex has degree... uh-oh!

The reduced graph G' might contain a vertex of degree 0, making the inductive hypothesis P(n) inapplicable! We are stuck— and properly so, since the claim is false!