Final Solutions

Name:	
manne.	

Problem 1. [10 points] Let $a_0 = a_1 = 1$, and let $a_{n+2} = a_{n+1} + 5a_n$ for $n \ge 0$. Prove by strong induction that $a_n \leq 3^n$ for all $n \geq 0$. You do not need to solve the recursion to do this.

Solution. Base Case: $a_0 \le 3^0 = 1$

Inductive Hypothesis: Assume for i = 0, 1, 2, ..., n + 1 that $a_i \le 3^i$. Inductive Step: $a_{n+2} = a_{n+1} + 5a_n \le 3^{n+1} + 5 \cdot 3^n = 8 \cdot 3^n \le 9 \cdot 3^n = 3^{n+2}$

Hence as the statement holds for i = n + 2 given i = 0, 1, 2, ..., n + 1, we conclude that it holds for all $n \geq 0$.

Problem 2. [10 points] The MIT Social Statistics Society (MIT S³) is doing a study to see how social engineers are. So they host a party with 15 students. After the party, 8 students report not having met anyone at the party, 4 students report having met one other student each, 2 students report have met two other students each, and one student reports having met three other students. Is this possible? You must prove your answer. Assume that the "met" relation is symmetric and anti-reflexive.

Solution. This is not possible. The total sum of degrees is 1+1+1+1+2+2+3=11, which implies that there are $\frac{11}{2}$ edges, which is not possible.

Name: _____

Problem 3. [10 points] Find a closed form for $\prod_{i=1}^{n} \prod_{j=i}^{n} 2^{(i-j)}$.

Solution. The solution is:

$$\begin{split} \prod_{i=1}^{n} 2^{(n-i+1)i} 2^{-(i+\ldots+n)} &= \prod_{i=1}^{n} 2^{(n-i+1)i} 2^{-\frac{(n+i)(n-i+1)}{2}} \\ &= 2^{\left(\sum_{i=1}^{n} \left(i(n-i+1) - \frac{(n+i)(n-i+1)}{2}\right)\right)} \\ &= 2^{\left(\sum_{i=1}^{n} \left(i(n-i+1) - \frac{n^2 - i^2}{2} - \frac{n+i}{2}\right)\right)} \\ &= 2^{\left(\frac{n(n+1)^2}{2} - \frac{n(n+1)(2n+1)}{6} - \frac{n^3}{2} + \frac{n(n+1)(2n+1)}{12} - \frac{n^2}{2} - \frac{n(n+1)}{4}\right)} \\ &= 2^{\frac{n-n^3}{6}} \end{split}$$

Name:

Problem 4. [10 points]

(a) [7 pts] Find a solution to $f_n = 4f_{n-1} + 5f_{n-2}$, with $f_0 = 1, f_1 = 1$.

Solution.
$$f_i = (2/3)(-1)^i + (1/3)(5)^i$$

Name.	
manne.	

(b) [8 pts] Give an asymptotic expression for the following recurrence, in Θ notation:

$$T(n) = 8T(\frac{n}{4}) + 18T(\frac{n}{6}) + n^3, \qquad T(1) = 0$$

Solution. Using Akra-Bazzi, we see that $a_1 = 8, b_1 = \frac{1}{4}, a_2 = 18, b_2 = \frac{1}{36}$ and so p = 2. Hence, $T(n) = \Theta\left(n^2\left(1 + \int_1^n \frac{n^3}{n^3}dn\right)\right) = \Theta(n^3)$.

Name:	
1	

Problem 5. [20 points]

(a) [6 pts] A cashier wants to work 5 days a week, but he wants to have at least one of Saturday or Sunday off. How many ways can he choose the days he will work? Your answer should be an integer.

Name:	

Solution. The cashier can choose to work on exactly one of Saturday or Sunday in $2 \cdot {5 \choose 4}$ ways, and the cashier can choose to work on neither Saturday or Sunday in ${5 \choose 5}$ ways. Hence the total number of ways is just 11.

(b) [6 pts] How many permutations of $1, 2, 3, \ldots n$ are there if 1 must precede 2 and 3 must precede 4 (for positive integers $n \ge 4$). Your answer should be in terms of n.

Solution. Let us label both 1, 2 as A, A and 3, 4 as B, B. Then each permutation of the n symbols A, A, B, B, 5, 6, . . . n corresponds to a permutation where 1 precedes 2 and 3 precedes 4. Hence the number of ways to order these n symbols is just $\frac{n!}{2\cdot 2} = \frac{n!}{4}$.

TA T		
Name:		

(c) [8 pts] Let $a_1, a_2, \ldots a_k$ be positive integers with sum at most n (with k > 1). Use a combinatorial argument to show that $a_1!a_2!\ldots a_k! < n!$.

Solution. Let us first increase a_k to a'_k such that $a_1 + a_2 + \ldots + a'_k = n$. Then $\frac{n!}{a_1!a_2!\ldots a'_k!}$ is just the number of ways to order n items of which there are a_1 of them being the same, a_2 of them being the same, ..., a'_k of them being the same. As this is just an enumeration of items, we must have that it is a positive quantity. Hence $\frac{n!}{a_1!a_2!\ldots a'_k!} \geq 1$. Now this implies that $n! \geq a_1!a_2!\ldots a'_k! > a_1!a_2!\ldots a_k!$.

Name:	
manne.	

Problem 6. [10 points] Find the generating function of the number of solutions to

$$x_1 + 2x_2 + 3x_3 + 4x_4 = n$$

where x_1, x_2, x_3, x_4 are positive integers. Express your answer as the inverse of a product of polynomials.

Solution. A generating function for x_1 is just $(x + x^2 + x^3 + \ldots) = \frac{x}{1-x}$. A generating function for $2x_2$ is just $(x^2 + x^4 + x^6 + \ldots) = \frac{x^2}{1-x^2}$. A generating function for $3x_3$ is just $(x^3 + x^6 + x^9 + \ldots) = \frac{x^3}{1-x^3}$. A generating function for $4x_4$ is just $(x^4 + x^8 + x^{12} + \ldots) = \frac{x^4}{1-x^4}$. Hence a generating function for the number of solutions to the above equation is just $\frac{x \cdot x^2 \cdot x^3 \cdot x^4}{(1-x)(1-x^2)(1-x^3)(1-x^4)}$, and the number of solutions to the particular solution above is the coefficient of the x^n term in the expansion of the generating function.

Problem 7. [15 points] Vlad's tiger has wandered into one of two forests overnight. It is in forest A with probability .4 and in forest B with probability .6.

If his tiger is in forest A and Vlad spends a day searching for it in forest A, the conditional probability that he will find his tiger that day is .25. Similarly, if his tiger is in forest B and he spends a day searching for it in forest B, then he will find his tiger that day with conditional probability .15.

You don't have to reduce your answer for the following problems.

(a) [5 pts] In which forest should Vlad look on the first day in order to maximize the probability of finding his tiger that day?

Solution. We just compare the probabilities that he finds the tiger in forest A vs. in forest B. As $.25 \cdot .4 \ge .15 \cdot .6$, he should look in forest A first.

Name:	

(b) [5 pts] Vlad looked in forest A on the first day but didn't find his tiger. What is the probability that the tiger is in forest A?

Solution. We just find the conditional probability that the tiger is in forest A given that Vlad didn't find his tiger and that he searched in forest A. This is just $\frac{.75 \cdot .4}{.75 \cdot .4 + .6 \cdot 1} = \frac{1}{3}$.

Name:
(c) [5 pts] Vlad flips a fair coin to determine where to look on the first day and finds his tiger on the first day. What is the probability that he looked in forest A?
Solution. This is just the conditional probability that Vlad looked in forest A given that he found his tiger. This is given by $\frac{.5 \cdot .4 \cdot .25}{.5 \cdot .4 \cdot .25 + .5 \cdot .6 \cdot .15} = \frac{10}{19}$.

n. T	
Name:	

Problem 8. [10 points] In a permutation of n elements, a pair (i, j) is called an inversion if and only if i < j and i comes after j. For example, the permutation 31542 in the case of n=5 has five inversions: (3,1),(3,2),(5,4),(5,2) and (4,2). What is the expected number of inversions in a uniform random permutation of the number $1, 2, \dots n$?

Hint: Use appropriate indicator variables and linearity of expectation.

Solution. We use $\mathbb{I}(i,j)$ to be an indicator variable that the pair (i,j) is an inversion. Then if T is the number of inversions, by linearity of expectation, the number of inversions

$$\mathbb{E}(T) = \sum_{i=1}^{n-1} \sum_{j=i}^{n} \mathbb{E}(\mathbb{I}(i,j))$$
. Now the probability that (i,j) is an inverted pair is just $\frac{1}{2}$ as we draw a uniform random permutation. Hence $\mathbb{E}(\mathbb{I}(i,j)) = \frac{1}{2}$ and so the $\mathbb{E}(T) = \frac{1}{2} \binom{n}{2}$.

draw a uniform random permutation. Hence $\mathbb{E}(\mathbb{I}(i,j)) = \frac{1}{2}$ and so the $\mathbb{E}(T) = \frac{1}{2} \binom{n}{2}$.

Problem 9. [10 points] Consider tossing a non-fair coin C until one throws a heads. Tossing C results in heads with probability $\frac{1}{3}$. Let X be a random variable corresponding to the number of tosses needed until one throws a heads (so $X \ge 1$).

(a) [5 pts] Calculate $\mathbb{E}(X)$.

Solution. We use a recursion to calculate $\mathbb{E}(X)$. Namely, $\mathbb{E}(X) = \frac{1}{3} + \frac{2}{3}(\mathbb{E}(X) + 1)$. This implies that $\mathbb{E}(X) = 3$.

T. T		
Name:		

(b) [5 pts] Calculate the variance of X.

Solution. We use the hint. We already know $\mathbb{E}(X)$. Now we just calculate $\mathbb{E}(X^2)$ again by using a recursion. Namely, $\mathbb{E}(X^2) = \frac{1}{3} + \frac{2}{3}(\mathbb{E}((X+1)^2))$.

Hence we have,
$$\mathbb{E}(X^2) = \frac{1}{3} + \frac{2}{3}(\mathbb{E}(X^2) + 2\mathbb{E}(X) + 1)$$
, and so $\mathbb{E}(X^2) = 15$.

Thus
$$\mathbb{E}(X^2) - \mathbb{E}(X)^2 = 15 - 9 = 6$$
 is the variance.

NT .	
Name:	

Problem 10. [10 points] One raffle ticket is drawn randomly from a bowl containing four tickets numbered 1, 2, 3, and 4. Consider the following three random variables:

- 1. Let A be a binary random variable that is 1 if a 1 or 2 is drawn and 0 otherwise.
- 2. Let B be a binary random variable that is 1 if a 1 or 3 is drawn and 0 otherwise.
- 3. Let C be a binary random variable that is 1 if a 1 or 4 is drawn and 0 otherwise.

Assume that all raffle tickets have equal probability of being drawn

(a) [5 pts] Are A, B, C mutually independent? Briefly justify your answer.

Solution. No A, B, C are not mutually independent as fixing values for A, B, uniquely determine a value for variable C.

TA T		
Name:		

(b) [5 pts] Are A, B, C pairwise independent? Briefly justify your answer.

Solution. A, B, and C are pairwise independent. We first see that A, B are independent. Now we show that A, C are independent and by symmetry this will imply B, C are independent. We have that $P(C=1|A=1)=\frac{1}{2}=P(A=1)$ as now C is uniquely determined by the value of B and so is 1 with probability $\frac{1}{2}$ and 0 with probability $\frac{1}{2}$. Similarly, $P(C=1|A=0)=\frac{1}{2}=P(A=0)$. Hence we have that C and A are independent events.

Name:	
manne.	

Problem 11. [15 points] Let X be a random variable indicating the runtime of an algorithm on an input of size n. You know that $X \ge 0$ and that $\mathbb{E}(X) \le 10n$.

(a) [5 pts] Give as good a bound as you can on the probability that $X \geq 20n$.

Solution. Using Markov's inequality, we find that $\mathbb{P}(X \geq 20n) \leq \frac{10n}{20n} = \frac{1}{2}$.

(b) [5 pts] Now suppose you are told that $Var(X) \leq 10n$. Use Chebyshev's inequality to bound the probability that $X \geq 20n$.

Solution. By Chebyshev's inequality, we have $\mathbb{P}(X \geq 20n) \leq \mathbb{P}(|X - 10n| \geq 10n) \leq \frac{\operatorname{Var}(X)}{100n^2} \leq \frac{1}{10}$.

Name:	

(c) [5 pts] Now suppose there is another random variable T indicating the runtime of a different algorithm on the same input of size n. You are given that $\mathbb{E}(e^T) \leq e^{10n}$. Write down as strong a bound as you can on the probability that $T \geq 20n$.

Solution. We use the fact that $\mathbb{P}(T \ge 20n) \le \mathbb{P}(e^T \ge e^{20n}) \le \frac{e^{10n}}{e^{20n}} = e^{-10n}$.

Name:	
i (dilic.	

(This space left blank internationally)

Name:	
i (dilic.	

(This space left blank internationally)

Name:	
i (dilic.	

(This space left blank internationally)